Absolute permutation editorial

problem: https://www.hackerrank.com/challenges/absolute-permutation/problem

problem statement

Given N and K, a permutation of the first N natural integer is called an absolute permutation if $|a_i-i|=k$ for every $i\in[1,N]$

We must print the lexicographically smallest absolute permutation for given N and K, or -1 if no such permutation exists.

Bad solution

The brute force solution is to generate all possible permutation and for each permutation check if it's a absolute permutation or not.

The time complexity for this solution is $O(N \times N!)$

```
C++11 code:
  Timeout solution by brute force
  Time complexity: O(t * n * n!)
#include <bits/stdc++.h>
using namespace std;
bool check(int *arr, int n, int k) {
  int i = 0;
  while (i < n \&\& abs(arr[i] - (i+1)) == k)
  return i >= n;
// Time complexity of the function: O(n * n!)
int* absolute permutation(int *arr, int n, int k) {
  int* res = \overline{NULL};
  do {
    if (check(arr, n, k)) {
      res = arr;
      break;
  }while(next_permutation(arr, arr+n));
  return res;
```

```
int main() {
    int t;
    cin >> t;
    cin.ignore(numeric_limits<streamsize>::max(), '\n');nk = q 2k

for (int t_itr = 0; t_itr < t; t_itr++) {
    int n, k;
    cin >> n >> k;

    int *arr = new int[n];
    for (int i = 0; i < n; ++i)
        arr[i] = i + 1;

    int* result = absolute_permutation(arr, n, k);

    if (result == NULL)
        cout << "-1";
    else {
        for (int i = 0; i < n; i++) {
            cout << result[i] <<' ';
        }
    }

    delete [] arr;
    cout << "\n";
}

return 0;
}</pre>
```

Test case 3 ©

Compiler Message

Test case 4 [©]

Terminated due to timeout

Optimized solution

The question is: can we generate directly the absolute permutation, when is possible? Firstable, we gonna check the solvability of the problem.

The solvability of the problem

Greetz to: mancha, int-e, Gilly, salt_ from ##maths on IRC freenode.

Can we generate the lexicographically smallest absolute permutation for given N and K? The formula of an absolute permutation:

$$\sum_{i=1}^{n} |a_i - i| - nk = 0, with \, a_i \in [1, n]$$

• if k = 0

$$\sum_{i=1}^{n} |a_i - i| = 0, <=> a_i = i, \text{ for every } i \in [1, n]$$

• if $k \neq 0$

we know that:

$$a_i - i = k$$
, if $a_i > i$
 $a_i - i = -k$, if $a_i < i$

That implies: $\sum_{i=1}^{n} (a_i - i) = 0$, with $a_i \in [1, n]$

So, we must have the same number of k 's and -k 's.

That leads, that n must be even.

Can we get a relation between n and k to know the solvability of the problem ?

$$n$$
 is even \iff $n=2q$ \iff $nk=q2k$ \iff $n=\frac{q}{k}\times 2k$ \iff $n=m\times 2k$ \iff $2k|n$

To get an absolute permutation for given n **and** k :

- *n* must be even
- 2k|n

A solution

```
So, we have: for every i \in [1,n], |a_i-i|=k <=> a_i=i+k and a_{i+k}=i we must be careful to not write a_i twice, so, we have to check if a_i is empty or not. Algorithm
```

```
for i \in [1, n]
if i is not visited
mark i as visited
mark i + k as visited
a_i = i + k
a_{i+k} = i
```

The complexity of this part of code is O(N)

```
C++ code:
 Optimized solution
 Time complexity: O(t * n)
#include <bits/stdc++.h>
using namespace std;
// Time complexity for the function: O(n)
int* absolute_permutation(int n, int k) {
 bool *visited = new bool[n];
  fill n(visited, n, false);
  // Check solvability
 bool check = (k != 0) ? (n % (2*k) == 0) : true;
  if (!check) return NULL;
  int* res = new int[n];
  for (int i = 1; i \le n
    if (!visited[i-1]) {
      visited[i-1] = true;
      visited[i+k-1] = true;
      res[i-1] = i + k;
      res[i+k-1] = i;
   }
  return res;
int main(){
  int t;
  cin >> t;
  cin.ignore(numeric_limits<streamsize>::max(), '\n');
  for (int t_itr = 0; t_itr < t; t_itr++) {</pre>
   int n, k;
    cin >> n >> k;
    /*int *arr = new int[n];
    for (int i = 0; i < n; ++i)
arr[i] = i + 1;*/
    int* result = absolute_permutation(n, k);
    if (result == NULL)
     cout << "-1";
    else {
     for (int i = 0; i < n; i++) {
    cout << result[i] <<' ';
    }
    cout << "\n";
  return 0;
```