There are some robots and factories on the X-axis. You are given an integer array robot where robot[i] is the position of the ith robot. You are also given a 2D integer array factory where factory[j] = [positionj, limitj] indicates that positionj is the position of the jth factory and that the jth factory can repair at most limitj robots.

The positions of each robot are **unique**. The positions of each factory are also **unique**. Note that a robot can be **in the same position** as a factory initially.

All the robots are initially broken; they keep moving in one direction. The direction could be the negative or the positive direction of the X-axis. When a robot reaches a factory that did not reach its limit, the factory repairs the robot, and it stops moving.

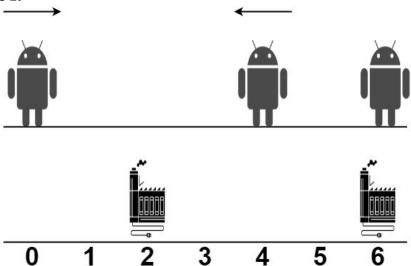
At any moment, you can set the initial direction of moving for **some** robot. Your target is to minimize the total distance traveled by all the robots.

Return *the minimum total distance traveled by all the robots*. The test cases are generated such that all the robots can be repaired.

Note that

- All robots move at the same speed.
- If two robots move in the same direction, they will never collide.
- If two robots move in opposite directions and they meet at some point, they do not collide. They cross each other.
- If a robot passes by a factory that reached its limits, it crosses it as if it does not exist.
- If the robot moved from a position X to a position Y, the distance it moved is | Y X |.

Example 1:



Input: robot = [0,4,6], factory = [[2,2],[6,2]]

Output: 4

Explanation: As shown in the figure:

- The first robot at position ${\tt 0}$ moves in the positive direction. It will be repaired at the first factory.

- The second robot at position 4 moves in the negative direction. It will be repaired at the first factory.

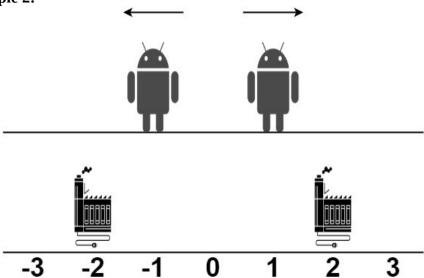
- The third robot at position 6 will be repaired at the second factory. It does not need to move.

The limit of the first factory is 2, and it fixed 2 robots.

The limit of the second factory is 2, and it fixed 1 robot.

The total distance is |2 - 0| + |2 - 4| + |6 - 6| = 4. It can be shown that we cannot achieve a better total distance than 4.

Example 2:



Input: robot = [1,-1], factory = [[-2,1],[2,1]]

Output: 2

Explanation: As shown in the figure:

- The first robot at position 1 moves in the positive direction. It will be repaired at the second factory.

- The second robot at position -1 moves in the negative direction. It will be repaired at the first factory.

The limit of the first factory is 1, and it fixed 1 robot.

The limit of the second factory is 1, and it fixed 1 robot.

The total distance is |2 - 1| + |(-2) - (-1)| = 2. It can be shown that we cannot achieve a better total distance than 2.

Constraints:

- 1 <= robot.length, factory.length <= 100
- factory[j].length == 2
- -109 <= robot[i], positionj <= 109
- 0 <= limitj <= robot.length
- The input will be generated such that it is always possible to repair every robot.

```
Include/Exclude: Recursion+Memoization
  Time complexity: O(nlogn+mlogm+nm)=O(mn)
 Space complexity: O(mn)+O(m+n)
                                         O Runtime
                                                                                 @ Memory
typedef long long ll;
                                                                                 67.92 MB | Beats 23.93%
                                         143 ms | Beats 21.21%
typedef std::vector<int> vi;
typedef std::vector<vi>vvi;
class Solution {
  public:
     ll minimumTotalDistance(vi& robot,vvi& factory){
       // To ensure that the robot goes the the nearest factory
       std::sort(robot.begin(),robot.end());
       std::sort(factory.begin(),factory.end());
          having 1 factory of k limits \leq having k factories of limit 1
          (at the same position)
       vi factories;
       for(auto& f: factory){
          for(int i=1;i \le f[1];++i) factories.push back(f[0]);
       int n=robot.size();
       int m=factories.size();
       std::vector<std::vector<ll>> memo(n,std::vector<ll>(m,-1));
       auto solve=[&](int robot index,int factory index,auto& self)->ll{
          // All robots are repaired
          //(Must be called first)
          if(robot_index==n) return 0;
          // If we still have broken robots and not no more factories
          if(factory index==m) return LONG MAX/2;
          if(memo[robot_index][factory_index]!=-1) return memo[robot_index][factory_index];
          ll include=
            abs(robot[robot_index]-factories[factory_index]) // current robot moves to current factory
            +self(robot index+1,factory index+1,self); // Pass the the next robot and next factory
                                                        // This is why sorting is important
          // Current robot skip the current factory
          // and go the next nearest factoty
          // This is why sorting is important
          ll exclude=self(robot_index,factory_index+1,self);
          return memo[robot_index][factory_index]=std::min(include,exclude);
       };
       return solve(0,0,solve);
                                    }};
```

```
Include/Exclude: bottom up
  Time complexity: O(nlogn+mlogm+nm)=O(mn)
  Space complexity: O(m+mn)
                                        O Runtime
typedef long long ll;
                                        63 ms | Beats 60.61%
                                                                              67.52 MB | Beats 33.74%
typedef std::vector<int> vi;
typedef std::vector<vi>vvi;
class Solution {
  public:
     ll minimumTotalDistance(vi& robot,vvi& factory){
       // To ensure that the robot goes the the nearest factory
       std::sort(robot.begin(),robot.end());
       std::sort(factory.begin(),factory.end());
       vi factories:
       for(auto& f: factory){
         for(int i=1;i <= f[1];++i) factories.push back(f[0]);
       int n=robot.size();
       int m=factories.size();
       std::vector<std::vector<ll>> dp(n+1,std::vector<ll>(m+1,0));
       for(int robot_index=1;robot_index<=n;++robot_index)</pre>
                dp[robot_index][0]=LONG_MAX/2;
       for(int robot_index=1;robot_index<=n;++robot_index){</pre>
         for(int factory index=1;factory index<=m;++factory index){</pre>
            ll include=abs(robot[robot_index-1]-factories[factory_index-1])
                   +dp[robot_index-1][factory_index-1];
            ll exclude=dp[robot_index][factory_index-1];
            dp[robot_index][factory_index]=std::min(include,exclude);
         }
       }
       return dp[n][m];
};
```

```
Include/Exclude: bottom up: space optimization
  Time complexity: O(nlogn+mlogm+nm)=O(mn)
  Space complexity: O(m)
                                          O Runtime
typedef long long ll;
                                                                                66.82 MB | Beats 37.42%
                                          58 ms | Beats 66.67%
typedef std::vector<int> vi;
typedef std::vector<vi>vvi;
class Solution {
  public:
    ll minimumTotalDistance(vi& robot,vvi& factory){
       std::sort(robot.begin(),robot.end());
       std::sort(factory.begin(),factory.end());
       vi factories;
       for(auto& f: factory){
         for(int i=1;i <= f[1];++i) factories.push_back(f[0]);
       }
       int n=robot.size();
       int m=factories.size();
       std::vector<ll> prev_row(m+1,0);
       for(int robot_index=1;robot_index<=n;++robot_index){</pre>
         std::vector<ll> cur row(m+1,LONG MAX/2);
         for(int factory_index=1;factory_index<=m;++factory_index){</pre>
            ll include=abs(robot[robot_index-1]-factories[factory_index-1])
                   +prev_row[factory_index-1];
            ll exclude=cur_row[factory_index-1];
            cur_row[factory_index]=std::min(include,exclude);
         prev_row=cur_row;
       return prev_row[m];
     }
};
```