

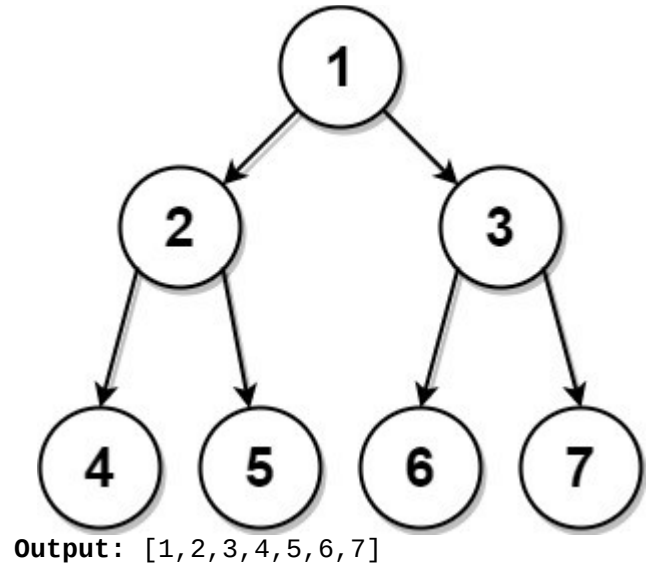
## 889. Construct Binary Tree from Preorder and Postorder Traversal

Given two integer arrays, `preorder` and `postorder` where `preorder` is the preorder traversal of a binary tree of **distinct** values and `postorder` is the postorder traversal of the same tree, reconstruct and return *the binary tree*.

If there exist multiple answers, you can **return any** of them.

### Example 1:

**Input:** `preorder = [1,2,4,5,3,6,7]`,  
`postorder = [4,5,2,6,7,3,1]`



### Example 2:

**Input:** `preorder = [1]`, `postorder = [1]`

**Output:** `[1]`

### Constraints:

- $1 \leq \text{preorder.length} \leq 30$
- $1 \leq \text{preorder}[i] \leq \text{preorder.length}$
- All the values of `preorder` are **unique**.
- $\text{postorder.length} == \text{preorder.length}$
- $1 \leq \text{postorder}[i] \leq \text{postorder.length}$
- All the values of `postorder` are **unique**.
- It is guaranteed that `preorder` and `postorder` are the preorder traversal and postorder traversal of the same binary tree.

## 889. Construct Binary Tree from Preorder and Postorder Traversal

```
/**
 * Definition for a binary tree node.
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode() : val(0), left(nullptr), right(nullptr) {}
 *     TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 *     TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left), right(right) {}
 * };
 */
/*
Recursion: Divide and Conquer
Time complexity: O(n)
Space complexity: O(2n)
*/
class Solution {
public:
    TreeNode* constructFromPrePost(std::vector<int>& preorder, std::vector<int>& postorder){
        int n=preorder.size();

        // nodes's values are distincts, we can map the postorder values a conctant time lookup
        // 1 <= preorder.length <= 30
        // 1 <= preorder[i] <= preorder.length
        std::vector<int> post_index(31);
        for(int i=0;i<n;++i){
            post_index[postorder[i]]=i;
        }
    }
};
```

```

// Recursive function to recover the binary tree from a preorder and a postorder DFS
// traversal
auto recover=[&](int pre_start,int pre_end,int post_start,auto& self)->TreeNode*{
    // Base case: If there are no nodes to process, return NULL
    if(pre_start>pre_end) return nullptr;

    // Base case: If only one node is left, return that node
    if(pre_start==pre_end) return new TreeNode(preorder[pre_start]);

    // Obtain the root of the left subtree from the preorder traversal
    int left_subtree_root=preorder[pre_start+1];

    // Obtain the left subtree size from the postorder traversal
    int left_subtree_size=post_index[left_subtree_root]-post_start+1;

    // Build the root and its left and right subtree
    // Root is obtained from preorder traversal
    // Based on the above data, we can determine the elements (ranges) of both left and right
    // subtrees
    TreeNode* root=new TreeNode(preorder[pre_start],
        self(pre_start+1,pre_start+left_subtree_size,post_start,self),
        self(pre_start+left_subtree_size+1,pre_end,post_start+left_subtree_size,self));

    return root;
};

return recover(0,n-1,0,recover);

}
};

```