Given a binary array nums and an integer k, return the maximum number of consecutive 1's in the array if you can flip at most k 0's.

#### Example 1:

**Input:** nums = [1,1,1,0,0,0,1,1,1,1,0], k = 2

Output: 6

**Explanation:**  $[1,1,1,0,0,\frac{1,1,1,1,1,1}{2}]$ 

Bolded numbers were flipped from 0 to 1. The longest subarray is underlined.

#### Example 2:

**Input**: nums = [0,0,1,1,0,0,1,1,1,0,0,0,0,1,1,1,1], k = 3

Output: 10

**Explanation:**  $[0,0,\underline{1,1,1,1,1,1,1,1,1},0,0,0,1,1,1,1]$ 

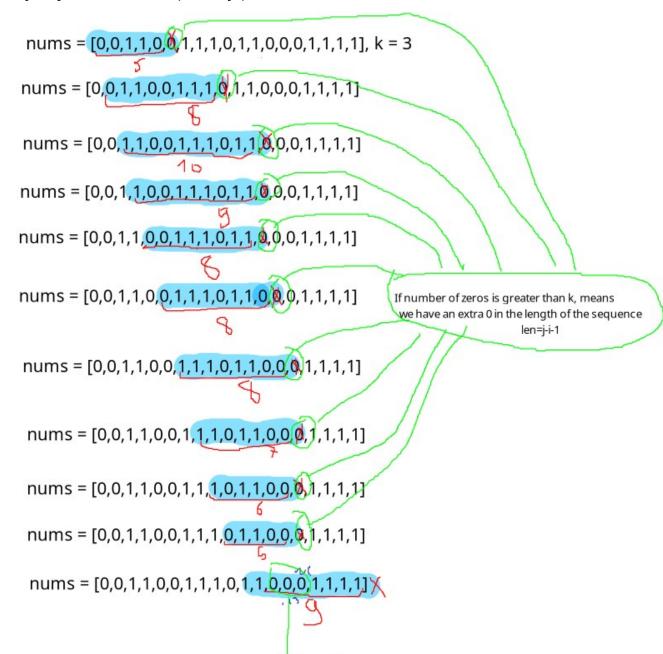
Bolded numbers were flipped from 0 to 1. The longest subarray is underlined.

#### **Constraints:**

- $1 \le nums.length \le 10^5$
- nums[i] is either 0 or 1.
- 0 <= k <= nums.length

## **Brute force approach**

Try all possible windows (subarrays)



If number of zeros is less or equal to k, means no xtra 0 in the length of the sequence is len=j-i

```
Brute Force: try all possibilities
  Time complexity: O(n^2) (TLE)
  Space complexity: O(1)
*/
class Solution {
public:
  int longestOnes(vector<int>& nums, int k) {
     int n=nums.size();
     int ans=INT MIN;
     // For each window
     for(int i=0;i< n;++i){
       int j=i; // starting from i
       // Expand it to the right, while counting the numbers of zeros
       int cnt_zero=0; // Initialize the number of zeros to 0
       // While the end of the array is not reached and the number of zeros is
       // not out of the limit k
       // Otherwise, the actual window can not be expanded any more
       while(j \le k \le cnt_zero \le k){
          cnt_zero+=(nums[i]==0); // Update the number of zeros
         j++; // Expand it to the right
       }
       // If number of zeros is greater than k, means
       // we have an extra 0 in the length of the sequence
       // If number of zeros is less or equal to k, means
       // no xtra 0 in the length of the sequence is len=j-i
       int len=cnt_zero>k?j-i-1:j-i;
       ans=std::max(ans,len); // Maximize the answer
     }
     return ans;
};
```

### **Dynamic sliding window**

Expand the window while number of zeros is less or equal to k, shrink it otherwise.

nums = 
$$[0,0,1,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1,1]$$
, k = 3  
nums =  $[0,0]$ ,1,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1,1], k = 3  
nums =  $[0,0]$ ,1,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1,1], k = 3  
nums =  $[0,0]$ ,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1,1], k = 3  
nums =  $[0,0]$ ,1,1,0,1,1,0,1,1,0,0,0,1,1,1,1], k = 3  
nums =  $[0,0]$ ,1,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1,1], k = 3  
nums =  $[0,0]$ ,1,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1,1], k = 3  
nums =  $[0]$ ,0,1,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1,1], k = 3

```
#0s=4
nums = [0,0,1,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1,1], k = 3
                          #0s=4
nums = [0.0.1.1, 0.0.1.1.1.0, 1.1.0, 0.0.1.1.1.1], k = 3
                          #0s=3,len=8,ans=10
 nums = [0,0,1,1,0,0,1,1,1,0,0,0,1,1,1,1], k = 3
                         #0s=4
nums = [0,0,1,1,0]0,1,1,1,0,1,1,0,00,1,1,1,1,1, k = 3
                          #0s=3,len=8,ans=10
 nums = [0.0,1.1,0.0,1.1,1.0,1.1,0.0,0.0,1.1,1.1], k = 3
                          #0s=4
 nums = [0.0,1,1,0,0,1,1,0,1,1,0,0,0,1,1,1,1], k = 3
  nums = [0.0,1.1,0.0.1,1.1,0.1,1.0,0.0,1.1,1.1], k = 3
                                  #0s=4
  nums = [0,0,1,1,0,0,1,1,1,0,0,0,0,1,1,1,1,1], k = 3
                                   \#0s=4
 nums = [0,0,1,1,0,0,1,1,1,0,1,1,0,0,0] 1,1,1,1], k = 3
                                   #0s=3,len=5,ans=10
 nums = [0,0,1,1,0,0,1,1,1,0,1,1,0,0,0] 1,1,1,1], k = 3
                                  #0s=3,len=6,ans=10
nums = [0,0,1,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1,1], k = 3
                                 #0s=3,len=7,ans=10
nums = [0.0,1,1,0,0,1,1,1,0,\frac{1}{1},\frac{1}{1},0,0,0,1,\frac{1}{1},\frac{1}{1},\frac{1}{1}], k = 3
                                  #0s=3,len=8,ans=10
nums = [0,0,1,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1], k = 3
                                 #0s=3,len=9,ans=10
nums = [0,0,1,1,0,0,1,1,1,0,1,1,0,0,0,1,1,1,1], k = 3
```

```
Dynamic sliding window
  Time complexity: O(n) (AC)
  Space complexity: O(1)
*/
class Solution {
public:
  int longestOnes(vector<int>& nums, int k) {
     int n=nums.size();
     int ans=INT_MIN;
     int cnt_zeros=0;
    // Expand the window to the right
     int l=0;
     for(int r=0;r< n;++r){
       cnt_zeros+=(nums[r]==0); // Add 1, if we encounter a 0
       // While number of zeros is greater than k
       while(cnt_zeros>k){
         cnt_zeros-=(nums[l]==0); // Subtract 1, if we encounter a 0, while ...
         1++; // ... shrinking the window from the left
       }
       ans=std::max(ans,r-l+1); // Maximize the answer
     }
    return ans;
  }
};
```