Given an array of positive integers nums, return the *maximum possible sum of an ascending subarray* in nums.

A subarray is defined as a contiguous sequence of numbers in an array.

A subarray [numsl, numsl+1, ..., numsr-1, numsr] is **ascending** if for all i where l <= i < r, numsi < numsi+1. Note that a subarray of size 1 is **ascending**.

Example 1:

```
Input: nums = [10, 20, 30, 5, 10, 50]
```

Output: 65

Explanation: [5,10,50] is the ascending subarray with the maximum sum of 65.

Example 2:

```
Input: nums = [10, 20, 30, 40, 50]
```

Output: 150

Explanation: [10,20,30,40,50] is the ascending subarray with the maximum sum of

150.

Example 3:

```
Input: nums = [12, 17, 15, 13, 10, 11, 12]
```

Output: 33

Explanation: [10,11,12] is the ascending subarray with the maximum sum of 33.

Constraints:

- 1 <= nums.length <= 100
- 1 <= nums[i] <= 100

Overview

We need to find the highest possible sum of an ascending subarray in a given array of positive integers. An ascending subarray is a contiguous sequence where each element is strictly smaller than the next (i.e., nums[i] < nums[i+1] for all valid indices). A subarray of size 1 is always considered ascending because there are no adjacent elements to compare.

Note: There is a difference between "ascending" and "non-decreasing." "Ascending" means strictly increasing, where each value is greater than the previous one. On the other hand, "non-decreasing" allows the values to stay the same or increase, so equality is permitted. For example:

- 1 2 3 4 5 is **ascending** (strictly increasing).
- 1 2 2 3 3 4 5 is **non-decreasing** (values can remain the same or increase).

Intuition

Instead of using the brute-force method of checking every possible subarray, which is inefficient, we can solve the problem in a single pass through the array.

The key idea is to keep extending the current subarray as long as it stays ascending. If we encounter an element that isn't greater than the previous one, we stop and compare the current subarray's sum with the largest sum we've found so far and update it if the current element is not greater than the previous one. Then, we reset the current sum to the current element and start a new subarray from there.

This strategy works because, with all numbers being positive, extending a subarray will always increase its sum. Thus, we should never start a new subarray when we can extend the current one. For an added challenge, try solving the similar problem <u>53. Maximum Subarray</u> (Kadane's algorithm), where the numbers are not restricted to be positive.

By following this idea, we only need to go through the array once. At the end of the loop, we perform a final check to ensure we account for the last subarray, just in case it had the largest sum.

The algorithm is visualized below:

```
Single pass, single pointer
  Time complexity: O(n)
  Space complexity: O(1)
typedef std::vector<int> vi;
class Solution {
  public:
    int maxAscendingSum(vi& nums) {
       int n=nums.size();
       int sum=nums[0]; // Prefix sum initial value
       int ans=nums[0]; // Array could be of size 1
       // Iterate over all elements
       for(int i=0;i< n-1;++i){
         // If previous element is greater than the next one
         if(nums[i]<nums[i+1]){</pre>
            // Cumulate the sum
            sum+=nums[i+1];
         // Otherwise, reset the sum the the value that breaks the condition
          else sum=nums[i+1];
         // Maximize the sum
          ans=std::max(ans,sum);
       }
       return ans;
     }
};
```

```
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  Space complexity: O(1)
typedef std::vector<int> vi;
class Solution {
  public:
    int maxAscendingSum(vi& nums) {
       int n=nums.size();
       int sum=nums[0],ans=0;
       for(int i=0;i< n-1;++i){
         if(nums[i] \ge nums[i+1]){
            ans=std::max(ans,sum);
            sum=0;
         sum+=nums[i+1];
       }
       return std::max(ans,sum);
     }
};
```