

# Artificial Intelligence II

## Multi-Agent Systems

Working together:  
Distributed Constraint Optimization  
(Complete approaches)

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# Outline

- Introduction
  - DCOP for MAS
  - how to model problems in the DCOP framework
- Solution Techniques for DCOPs
  - Exact algorithms (DCSP, DCOP)
    - DPOP
  - Approximate Algorithms (without/with quality guarantees)
    - DSA, MGM, Max-Sum, k-optimality, bounded max-sum

# Working together

## Coordination problem:

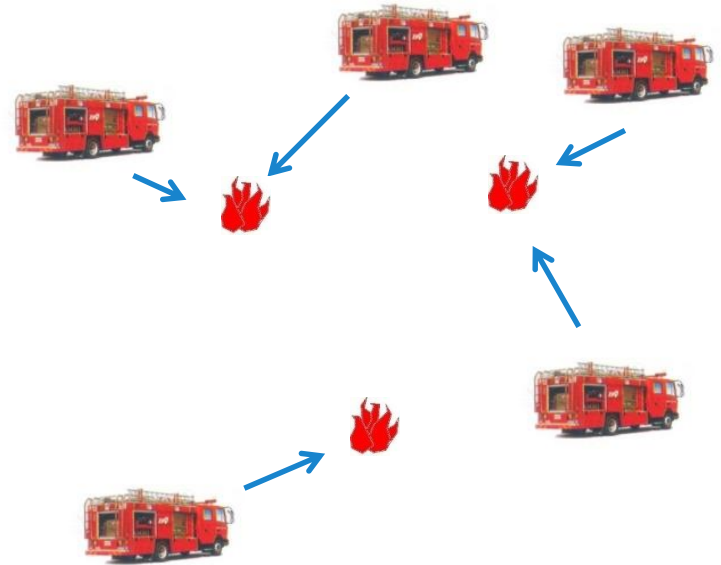
Choose agent's individual actions so to maximise a system-wide objective (benevolent agents)

## Task allocation:

individual actions: which fire to tackle

system-wide objective: minimise total  
extinguish time

solution: a **joint** action



# Decentralised Coordination

- Decentralised coordination: Local decision with local information
- Why Decentralised coordination ?
  - In general no benefit for computation or solution quality
  - Robustness
    - avoid single point of failure
  - Scalability
    - Not enough bandwidth to communicate/process all information
  - Leads to problem decomposition
    - Each agent cares only of **local** neighbours

# Decentralized Coordination: example



## Issues:

No central controller:  
distributed  
knowledge

Information sharing:  
limited  
communication

Complex system:  
difficult to  
design/analyse

# Decentralized Coordination: issues

- No central controller: distributed knowledge
- Information sharing: limited communication
- Complex system: difficult to design/analyse

# DCOPs for Decentralized Coordination

Why DCOPs for decentralized coordination ?

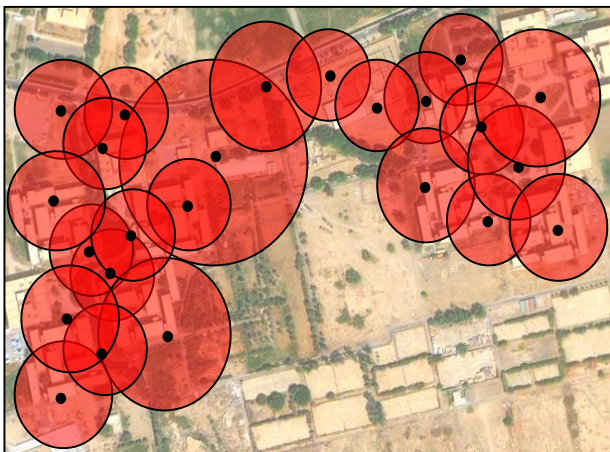
- Well defined problem
  - Clear mathematical formulation that captures most important aspects
  - Many solution techniques
    - Optimal: ABT, ADOPT, DPOP, ...
    - Approximate: DSA, MGM, Max-Sum, ...
- Solution techniques can handle large problems
  - compared for example to CNP and other sequential dec. making (MDP, POMDP)

# Reference Applications

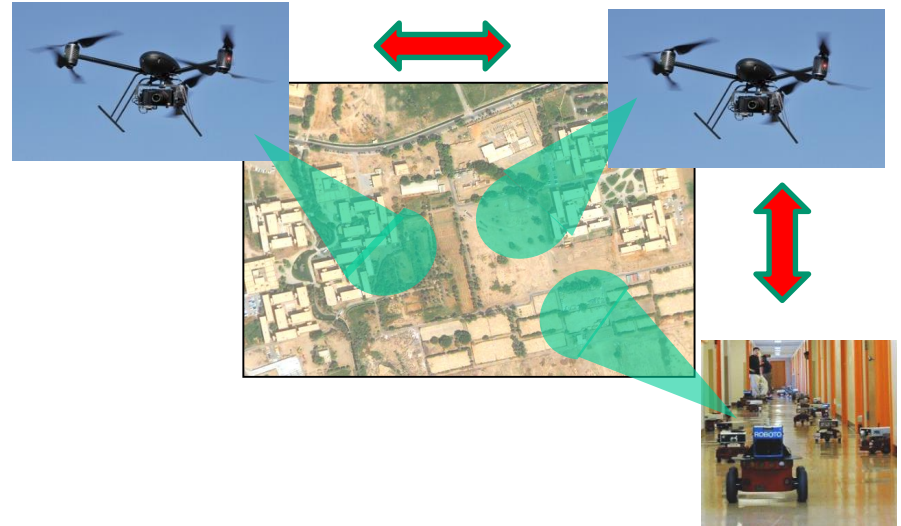
## Incident Management



## Environment monitoring



## Cooperative Exploration



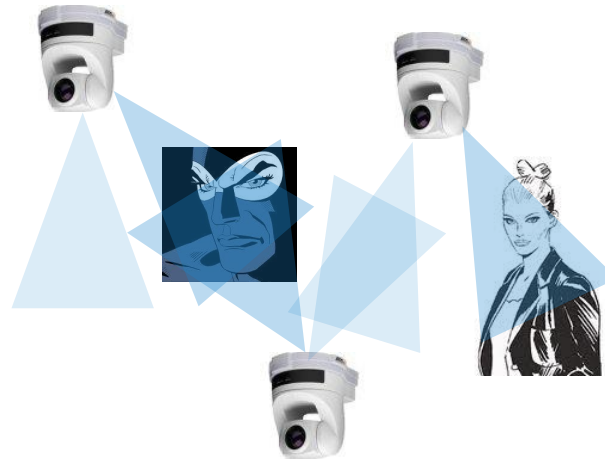
## Energy management



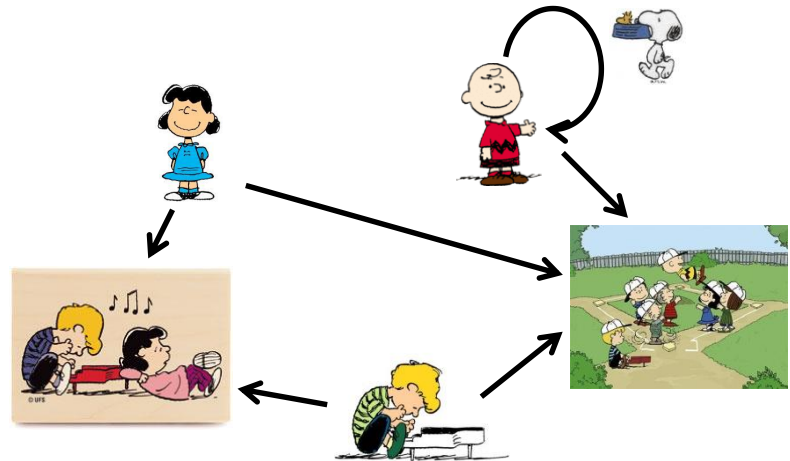


# Modeling Problems as DCOP

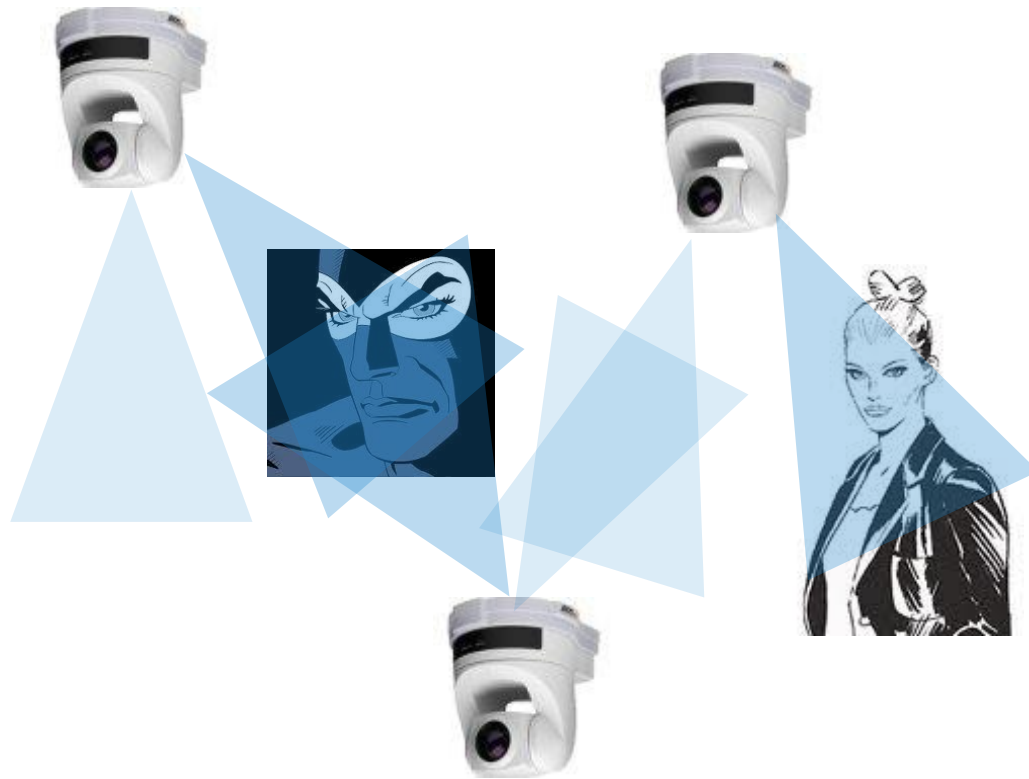
- Surveillance



- Meeting Scheduling



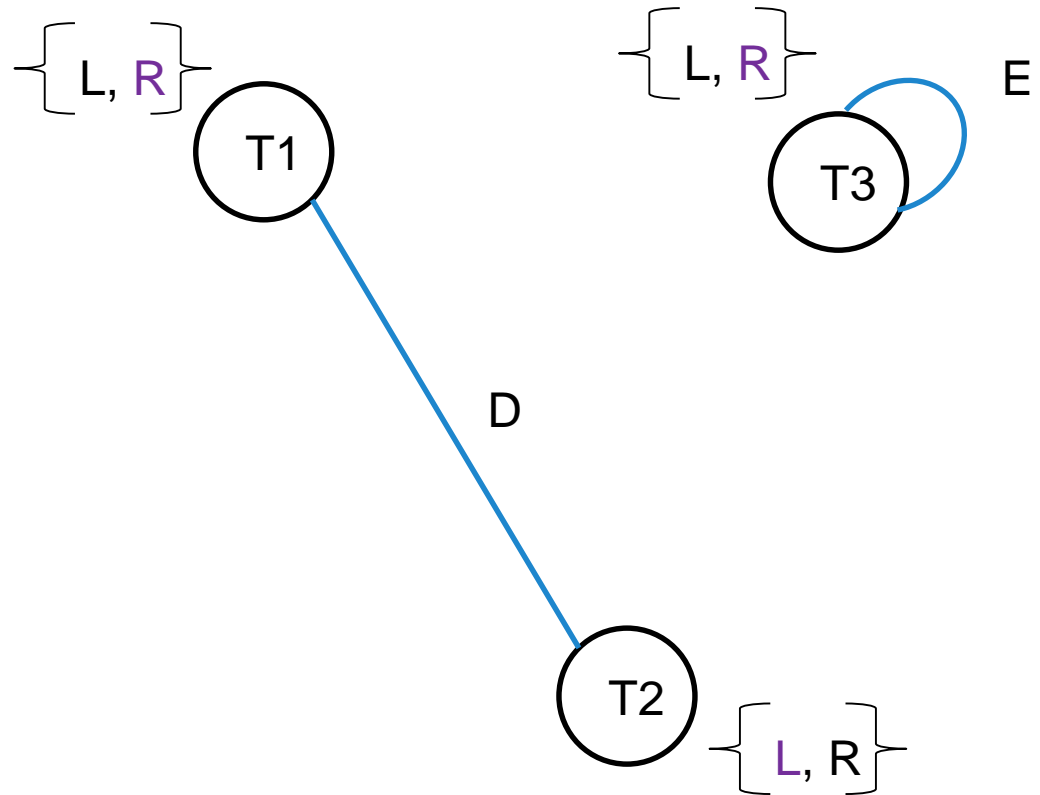
# Target Tracking



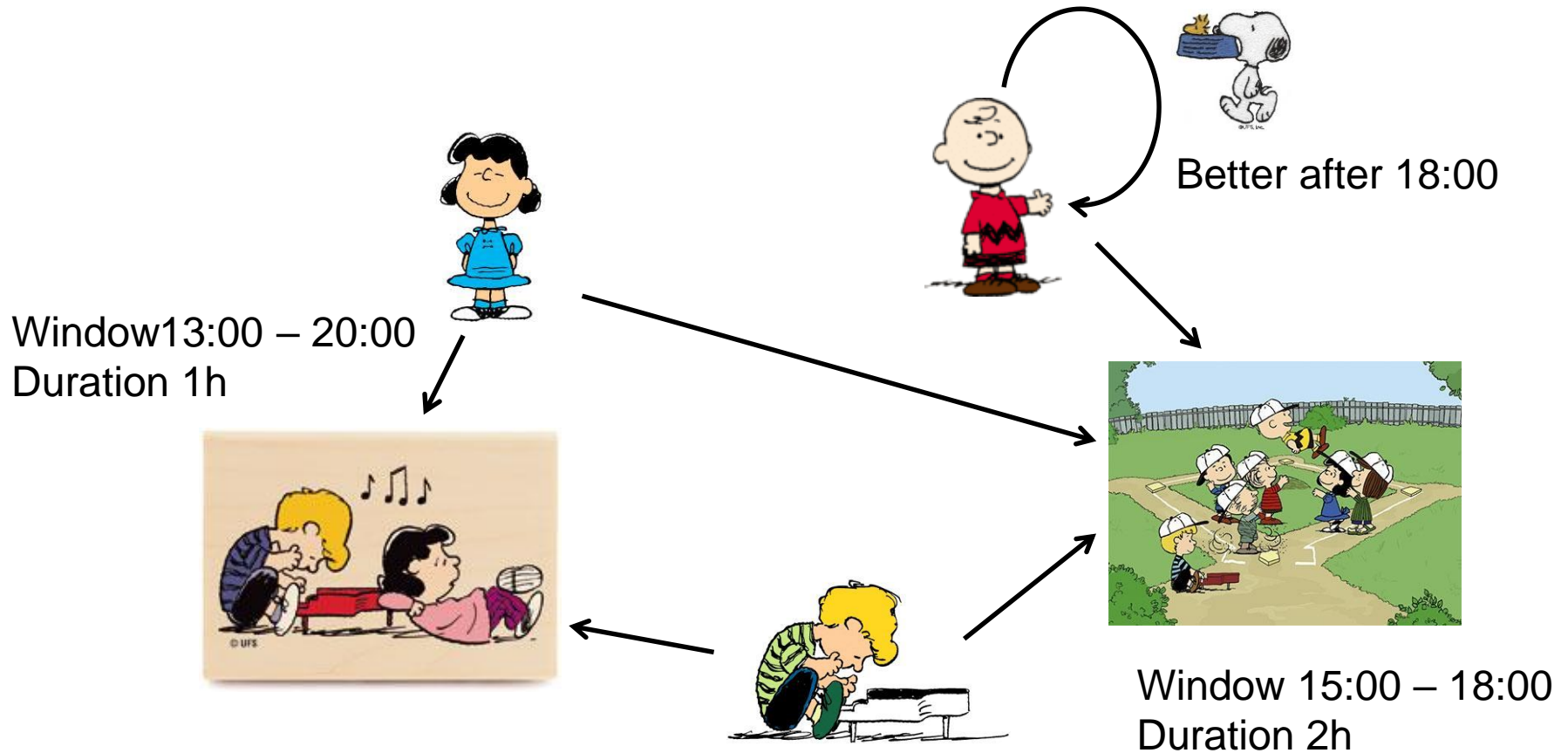
- Why decentralize
  - Robustness to failure and message loss

# Target Tracking - DCOP

- Variables -> Cameras
- Domains -> Camera actions
  - look left, look right
- Constraints
  - One constraint per target
  - Better assessment with more cameras
- Maximise “observation”

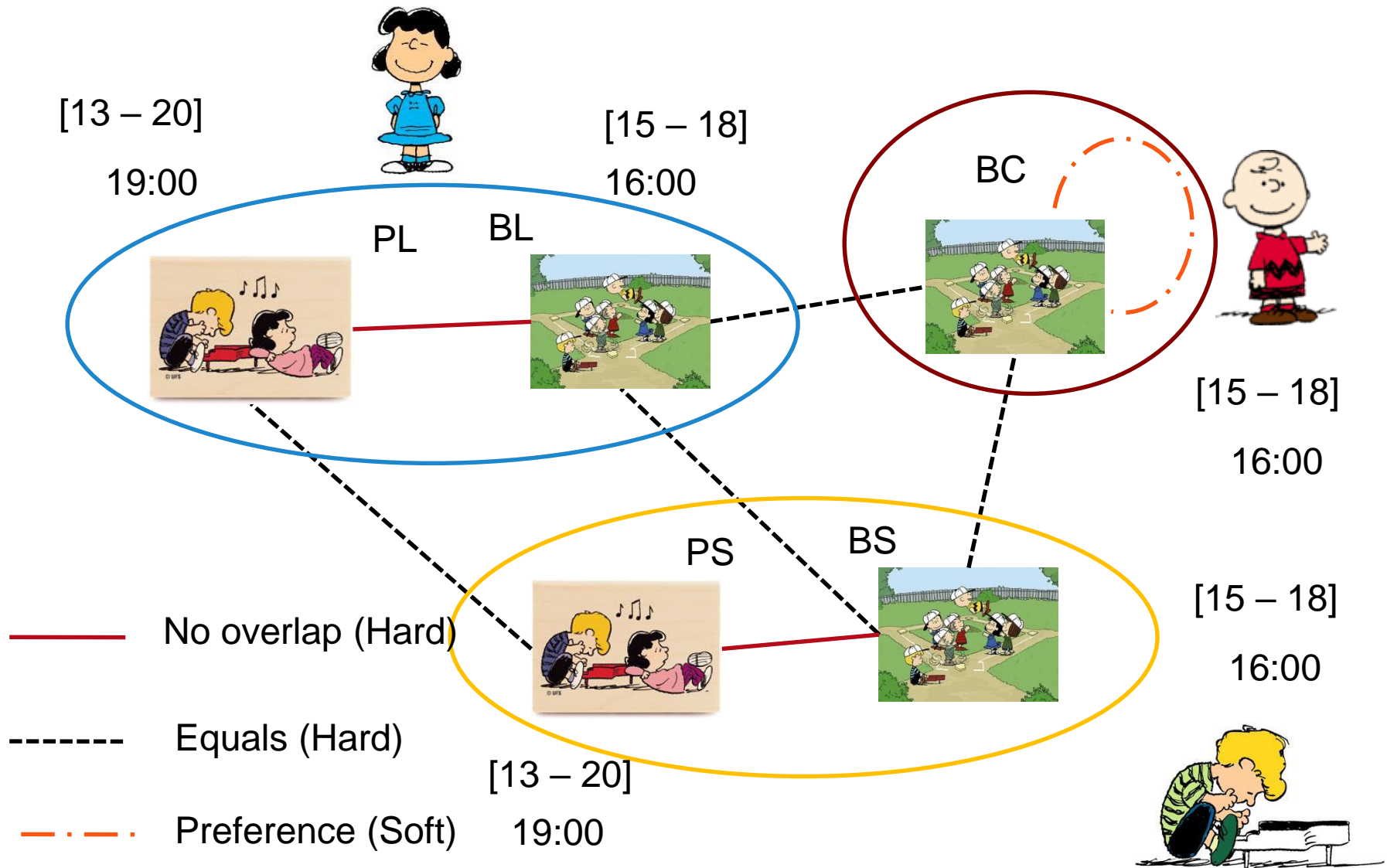


# Meeting Scheduling



- Why decentralize
  - Privacy

# Meeting Scheduling - DCOP



# Constraint Networks

$X = \{X_1, \dots, X_n\}$  a set of variables (e.g. meetings)

$D = \{D_1, \dots, D_n\}$  a set of discrete variable domains (e.g. time slots)

$C = \{C_1, \dots, C_m\}$  a set of constraints (e.g., equality, non overlap, )

$S_i \subseteq X$  Scope of constraint  $C_i$

Hard constraints

$R_i$	$x_j$	$x_k$
	0	1
	1	0

Soft constraint

$F_i$	$x_j$	$x_k$
2	0	0
0	0	1
0	1	0
1	1	1

# Objectives for constraint networks

- Constraint Satisfaction Problem (**CSP**)
  - Objective: find an assignment for all the variables in the network that satisfies all constraints
- Constraint Optimization Problems (**COP**)
  - Objective: find an assignment for all the variables in the network that satisfies all constraints and optimizes a global objective function

$$X^* = \arg \max_X \left( \sum_i F_i(X_i) \right)$$

Global function: an aggregation (i.e., sum) of local functions  $F_i(X_i)$

# Benchmarking problems

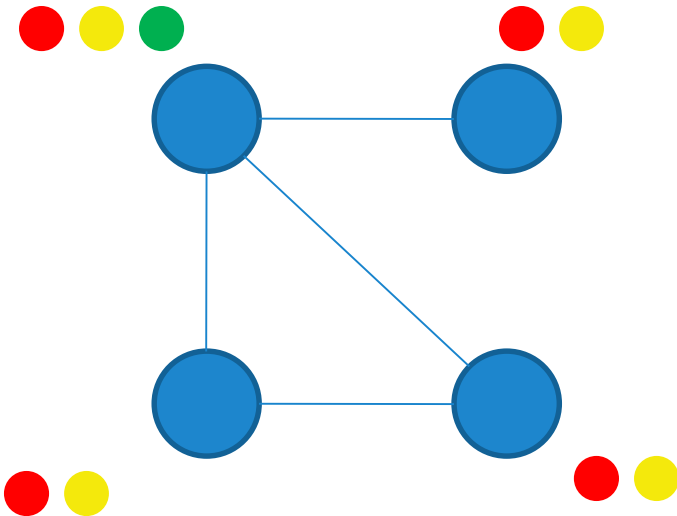
- Motivations
  - Analysis of complexity and optimality is not enough
  - Need to empirically evaluate algorithms on the same problem
- Graph coloring
  - Simple to formalise very hard to solve
  - Well known parameters that influence complexity
    - Number of nodes, number of colors, density (number of link/number of nodes)
  - Many versions of the problem
    - CSP, MaxCSP, COP



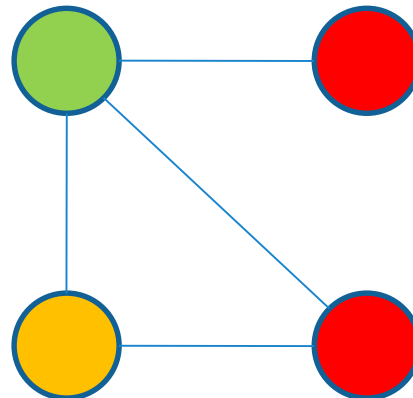
# Graph Coloring

- Network of nodes
- Nodes can take on various colors
- Adjacent nodes should not have the same color
  - If it happens this is a conflict

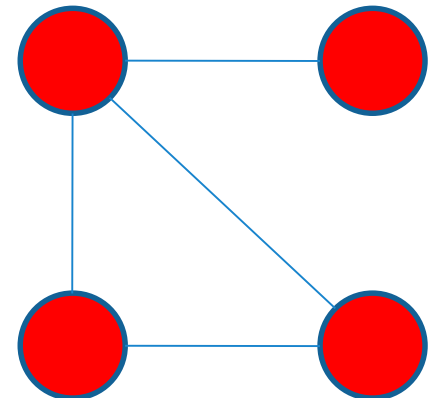
CSP



Yes

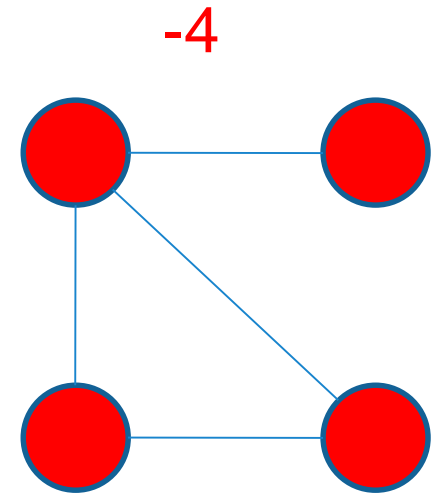
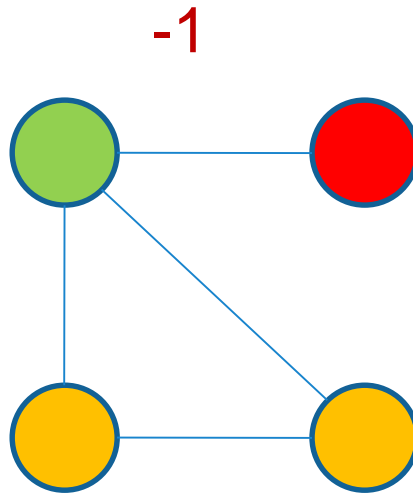
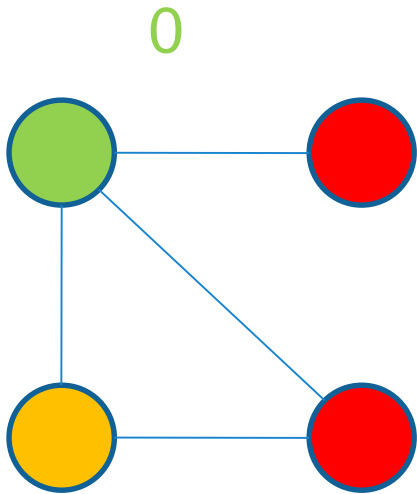


No



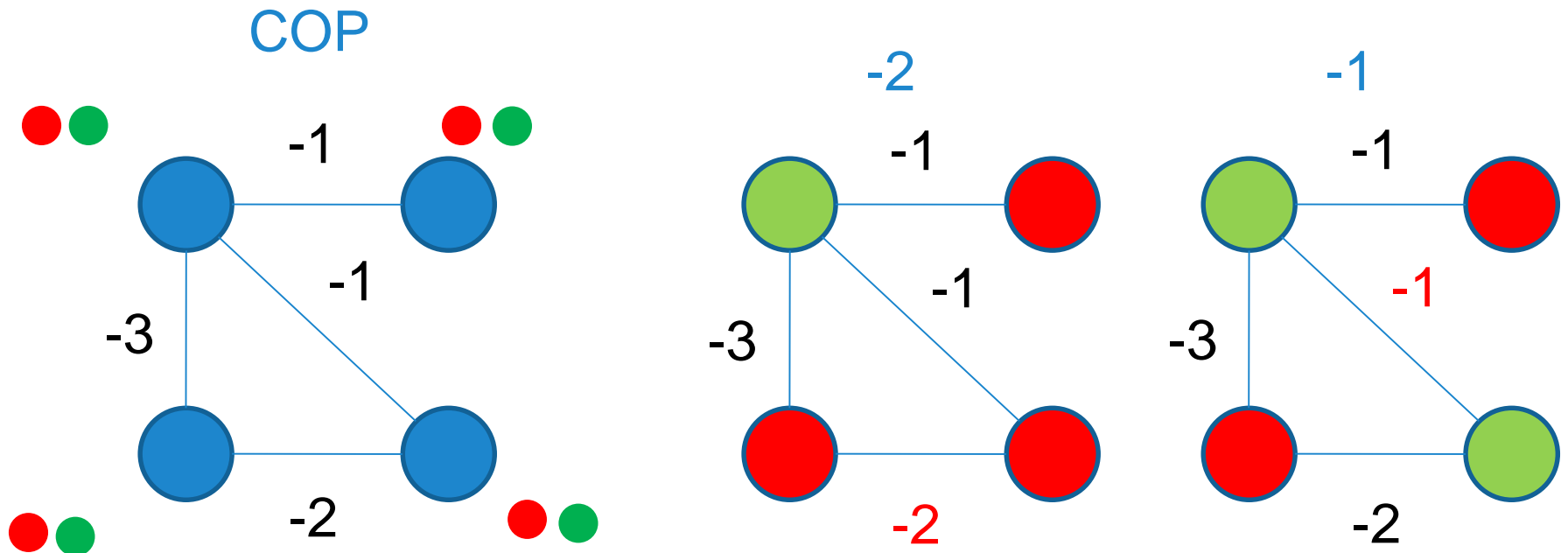
# Graph Coloring - MaxCSP

- Optimization Problem
- Natural extension of CSP
- Minimise number of conflicts



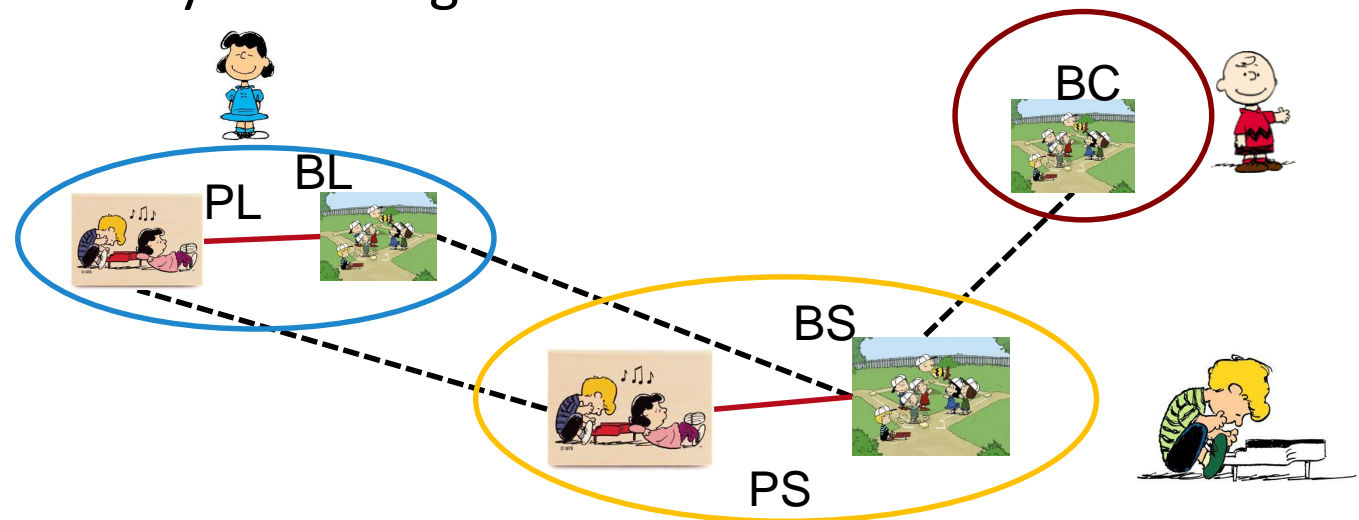
# Weighted Graph Coloring - COP

- Optimization Problem
- Conflicts have a weight
- Maximise the sum of weights of violated constraints



# Distributed COP

- We focus on optimization
- DCOP = Constraint Network + Agents  $A = \{A_1, \dots, A_k\}$
- Where each agent:
  - Controls a subset of the variables (typically just one)
  - Is only aware of constraints that involve the variables it controls
  - Communicates only with neighbors



# Performance measures

- Solution quality (the higher the better)
  - Optimality not always achievable,
  - Optimality Guarantees
- Coordination Overhead (the lower the better)
  - Computation: computation effort (time complexity)
  - Communication: number and **size** of messages (network load)
- Desirable properties (hard to quantify)
  - Robustness to failures, parallelism, flexibility, privacy maintenance, etc.

# DCOP Solution techniques

- Complete approaches
  - Guarantee optimal solution
  - Exponential coordination overhead
  - ADOPT, DPOP, OptAPO
- Heuristics
  - Low coordination overhead
  - No guarantees on optimality
  - DSA, MGM, Max-Sum
- Approximate approaches
  - Low coordination overhead
  - Optimality guarantees
  - Bounded max-sum, k-optimality

# Complete Approaches

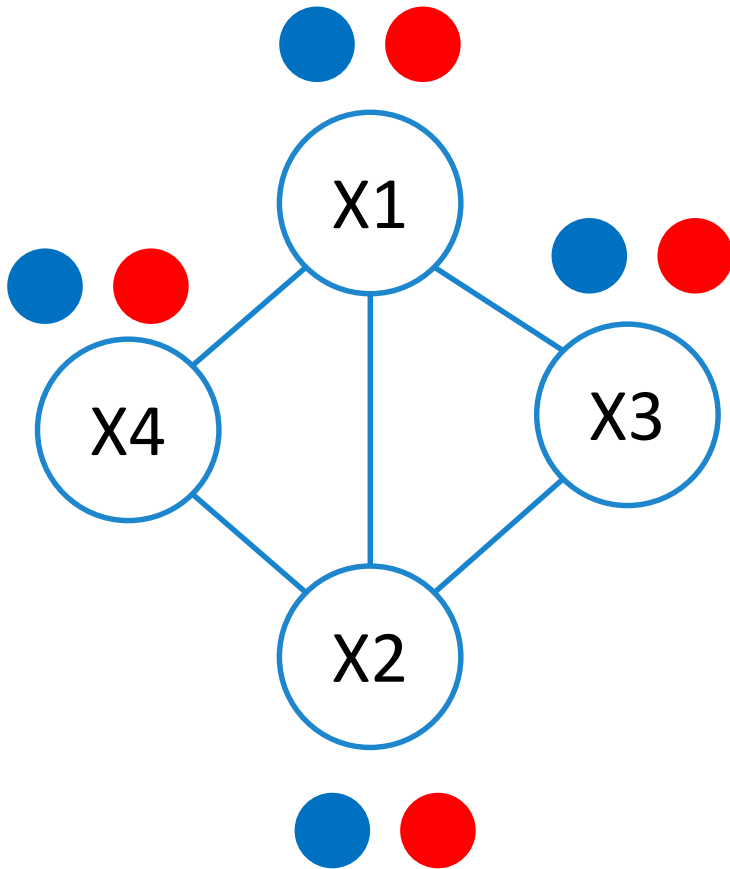
- ADOPT (Search based) [Modi et al 05]
  - Distributed branch and bound
  - Partial order based on a DFS search
  - Asynchronous (high parallelism, flexible)
  - Number of messages exponential in number of agents
- DPOP (Dynamic programming) [Petcu and Faltings 07]
  - Distributed Bucket Elimination
  - Partial order based on a DFS search
  - Linear number of messages
  - Exponential message size (in width of DFS search tree)
  - DFS-tree width typically much less than number of agents

# Dynamic Programming Optimization Protocol

- DFS-tree building (special case of Pseudo tree)
  - Constraint graph  $\rightarrow$  DFS-Tree
  - Token passing
- Utility propagation
  - Compile information to compute optimal value
  - Util messages from leaves to root
- Value Propagation
  - Root chooses optimal value and propagate decision
  - Value messages from root to leaves



# DPOP complete example

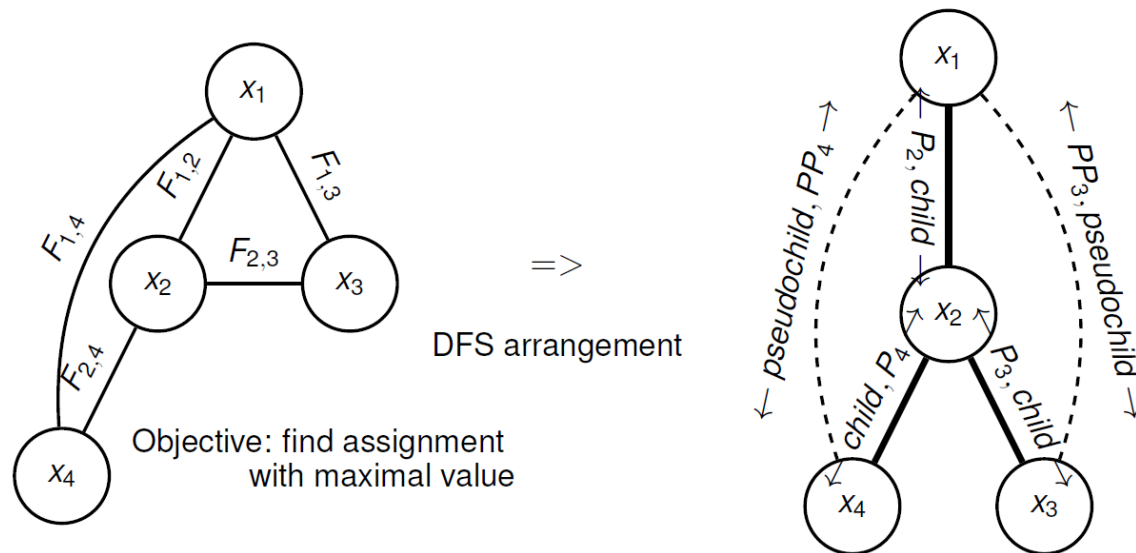


Value of Each constraint:  
same color -1  
different colors 0

Solve this constraint network with DPOP.  
Assume one agent per variable

# Pseudotrees: basic concepts

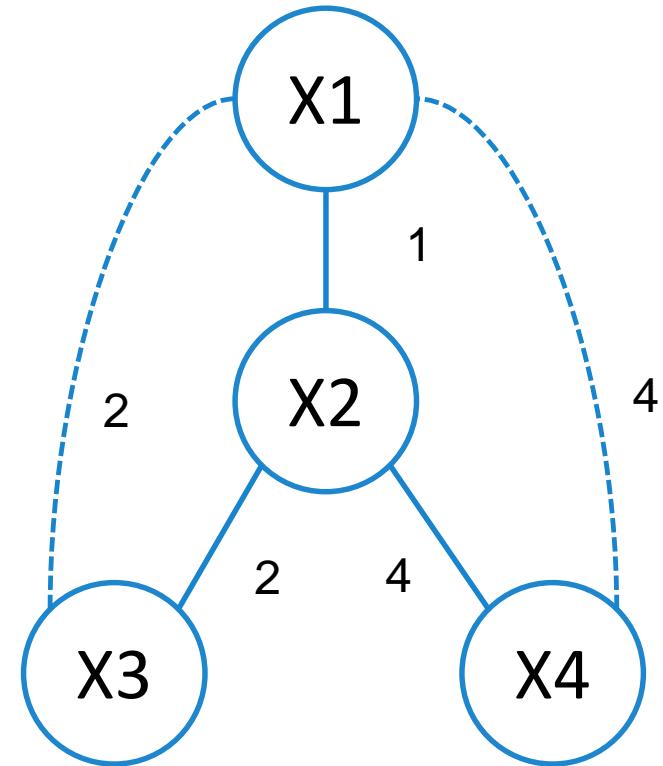
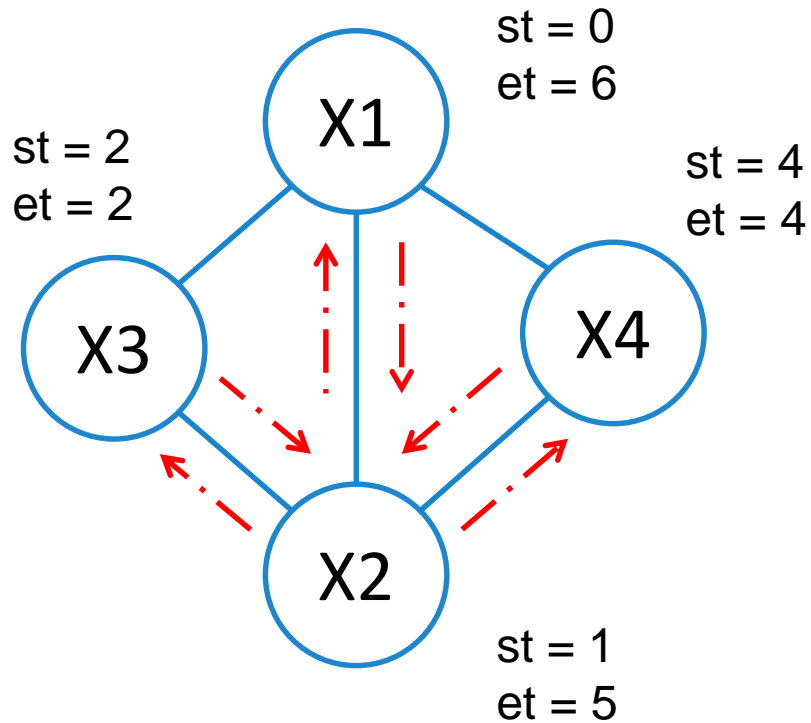
- Pseudotree arrangement of a graph  $G$ 
  - A rooted tree with the same nodes as  $G$
  - Adjacent nodes in  $G$  fall in the same branch of the tree
    - Nodes in different branches do not share direct constraints
  - A DFS visit of a graph induces a Pseudotree
    - Not every pseudotree can be obtained with a DFS visit



# Building a DFS tree

- Traverse the graph using a recursive procedure
- Each time we reach  $X_i$  from  $X_j$  we mark  $X_i$  as visited and state that  $X_j$  is the father of  $X_i$  (and  $X_i$  is a children of  $X_j$ )
- When a node  $X_i$  has a visited neighbor that is not its parent we state that  $X_j$  is a pseudo-parent of  $X_i$  (and  $X_i$  is a pseudochildren of  $X_j$ )
- Can be done with a distributed procedure:
  - Each node need only to communicate with neighbors
  - Token passing to propagate information (e.g., visited nodes)

# Building a DFS-tree: example



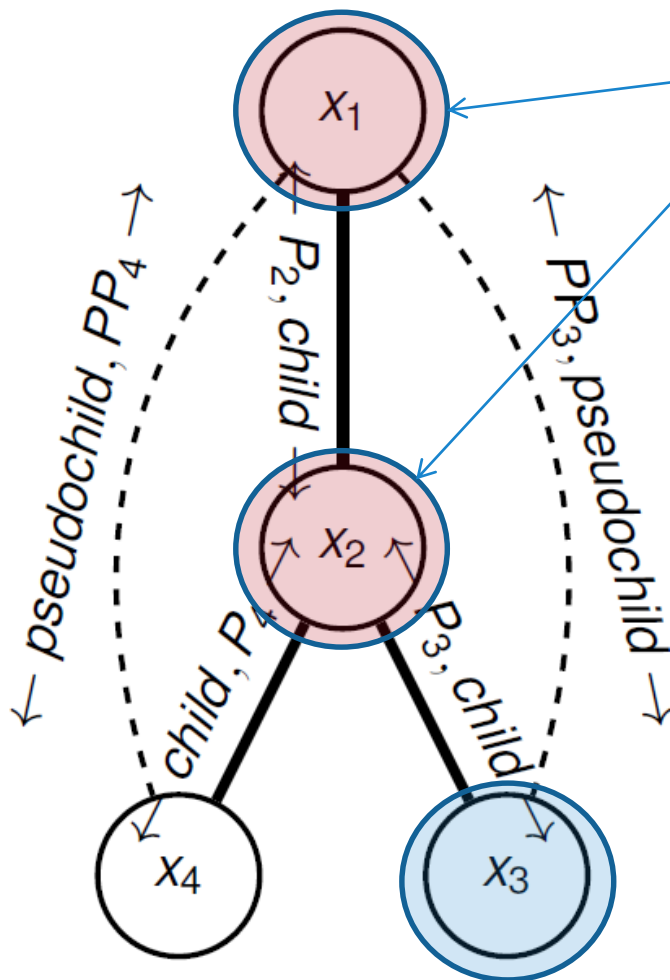
— . → token movement

st: first time node received the token

et: last time node sent the token

time++ = each time token moves

# Pseudotrees and Separator



Separator

## Definition

$Sep_i$  separator of node  $X_i$ : all ancestors (through tree and back edges) which are connected with  $X_i$  and with any descendant of  $X_i$

## Basic Property

$Sep_i$  minimal set of ancestors that, if removed, completely disconnects the subtree rooted at node  $X_i$  from the rest of the problem

## Operative Definition

$$Sep_i = \bigcup_{X_j \in C_i} Sep_j \cup P_i \cup PP_i \setminus X_i$$

$C_i$ : children of node  $i$

# Util Propagation

Aim: build a value function so that root agent can make optimal decision.

Dynamic programming: provide only key information

Each agent computes messages for its parent based on messages received from children and relevant constraints.

Each message projects out  $X_i$  (by maximisation) and aggregates (by summation) functions received from children and all constraints with ancestors (parents and pseudoparents)

# Message Computation

Functions  $\rightarrow$  tables (variable are all discrete)

Aggregation  $\rightarrow$  join operator (relational algebra)

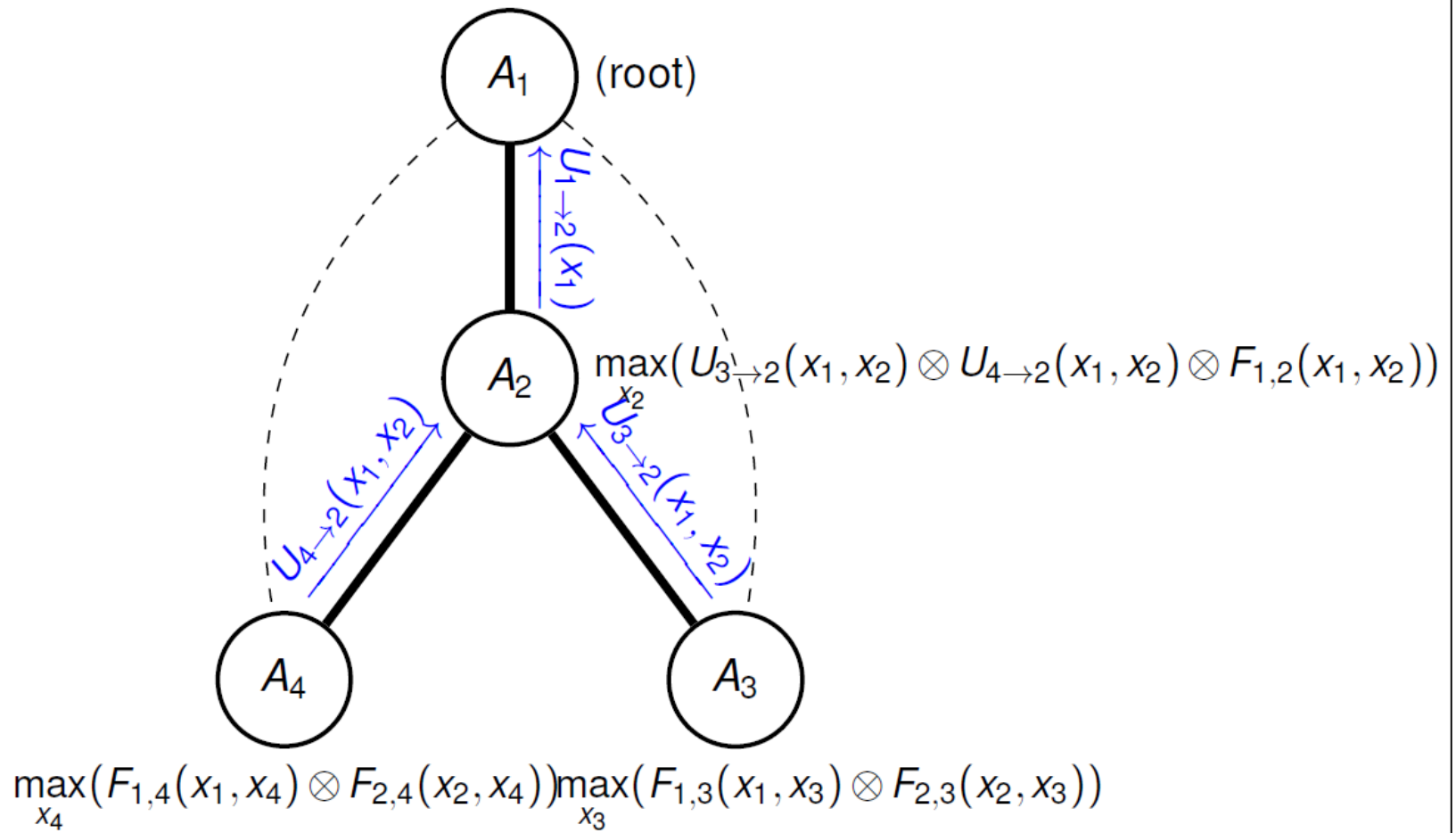
Maximization  $\rightarrow$  projection (keeping most valuable tuples)

The *Util* message  $U_{i \rightarrow j}$  that agent  $A_i$  sends to its parent  $A_j$  can be computed as:

$$U_{i \rightarrow j}(Sep_i) = \max_{x_i} \left( \bigotimes_{A_k \in C_i} U_{k \rightarrow i} \bigotimes \bigotimes_{A_p \in P_i \cup PP_i} F_{i,p} \right)$$

The  $\bigotimes$  operator is a join operator that sums up functions with different but overlapping scores consistently.

# Util Propagation: messages





# Util Propagation: message comp

Consider the computation of the node  $x_3$ :

$$U_{3 \rightarrow 2}(x_1, x_2) = \max_{x_3} (F_{1,3}(x_1, x_3) \oplus F_{2,3}(x_2, x_3))$$

1. Create  $F_{1,3}(x_1, x_3)$  and  $F_{2,3}(x_2, x_3)$  by checking the value of the constraints
2. Build the table for all the 8 combinations of the variables  $x_1, x_2, x_3$  with the values of  $F_{1,3}$  and  $F_{2,3}$
3. Sum the values of  $F_{1,3}$  and  $F_{2,3}$ , (call it  $V$ )
4. Maximize wrt  $x_3$ , i.e. for each pair of  $x_1, x_2$ , select the one with the highest value of  $V$

# DPOP complete example

Value of Each constraint:  
 same color -1  
 different colors 0

X1	X2	X3	F13	F23	F13+F23
0	0	0	-1	-1	-2
0	0	1	0	0	0
0	1	0	-1	0	-1
0	1	1	0	-1	-1
1	0	0	0	-1	-1
1	0	1	-1	0	-1
1	1	0	0	0	0
1	1	1	-1	-1	-2

X1	X3	F13
0	0	-1
0	1	0
1	0	0
1	1	-1

X2	X3	F23
0	0	-1
0	1	0
1	0	0
1	1	-1

X1	X2	U32
0	0	0
0	1	-1
1	0	-1
1	1	0

$$U_{3 \rightarrow 2}(x_1, x_2) = \max_{x_3} (F_{1,3}(x_1, x_3) \oplus F_{2,3}(x_2, x_3))$$

# Value util example

X1	X2	F12	U32	U42	F12+U32+U42
0	0	-1	0	0	-1
0	1	0	-1	-1	-2
1	0	0	-1	-1	-2
1	1	-1	0	0	-1

X1	U12
0	-1
1	-1

$$U_{2 \rightarrow 1}(x_1) = \max_{x_2} (U_{3 \rightarrow 2}(x_1, x_2) \oplus U_{4 \rightarrow 2}(x_1, x_2) \oplus F_{1,2}(x_1, x_2))$$

When propagated to the root a decision can be taken on  $x_1$

# Value Propagation

Aim: inform all agents about decision from above so that they can choose best values for their variables

Root agent  $A_r$  computes  $x_r^*$  which is the argument that maximises the sum of messages received by all children

It sends a message  $V_{r \rightarrow c} = \{X_r = x_r^*\}$  containing this value to all children  $C_r$

The generic agent  $A_i$  sends a message to each child  $A_j$   
 $V_{i \rightarrow j} = \{X_s = x_s^*\} \cup X_i = x_i^*$ , where  $X_s \in Sep_i \cap Sep_j$

# Value Computation

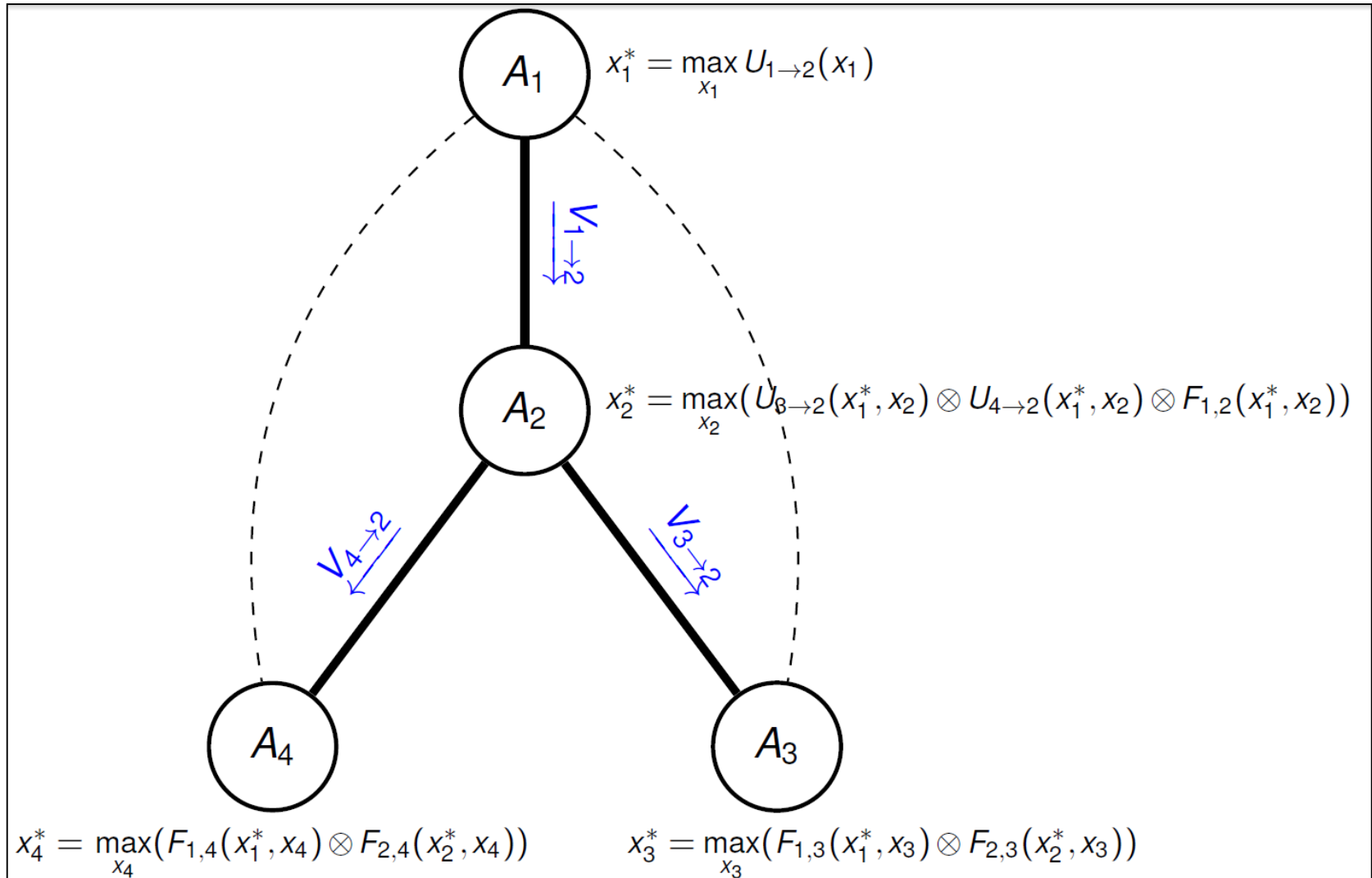
After receiving a value message from its parent, each agent computes the optimal value for its variable as:

$$x_i^* = \underset{x_i}{\operatorname{argmax}} \left( \sum_{A_j \in C_i} U_{j \rightarrow i}(x_p^*) + \sum_{A_j \in P_i \cup PP_i} F_{i,j}(x_i, x_j^*) \right)$$

where  $x_p^* = \bigcup_{A_j \in P_i \cup PP_i} \{x_j^*\}$  is the set of optimal values for  $A_i$ 's parent and pseudoparents received from  $A_i$ 's parent.

Can reuse stored tables for computing util messages

# Value Propagation: messages



# Value propagation

$x1 \leftarrow 0$

X1	X2	F12	U32	U42	F12+U32+F42
0	0	-1	0	0	-1
0	1	0	-1	-1	-2
1	0	0	-1	-1	-2
1	1	-1	0	0	-1

$x2 \leftarrow 0$

# Value propagation

$x1 \leftarrow 0$

$x2 \leftarrow 0$

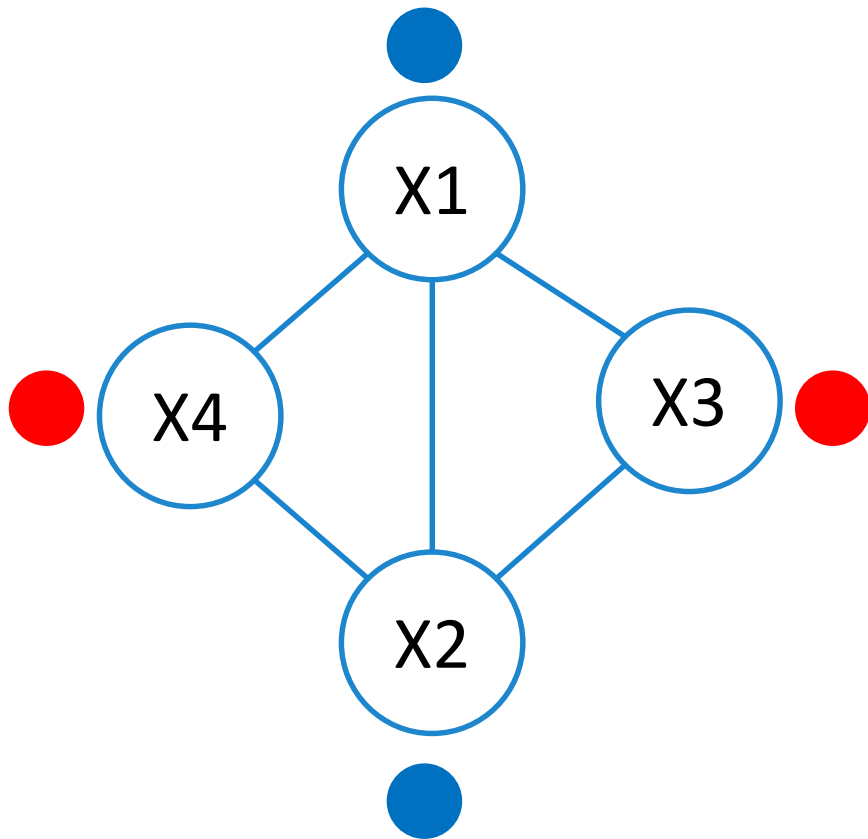
X1	X2	X3	F13	F23	F13+F23
0	0	0	-1	-1	-2
0	0	1	0	0	0
0	1	0	-1	0	-1
0	1	1	0	-1	-1
1	0	0	0	-1	-1
1	0	1	-1	0	-1
1	1	0	0	0	0
1	1	1	-1	-1	-2

$x3 \leftarrow 1$

$x4 \leftarrow 1$



# DPOP complete example



Only 1 constraint is violated

# DPOP analysis

- Synchronous algorithm
- Linear number of messages but exponentially large
- Messages (and computation) is exponential in separators size
- Separator size  $\rightarrow$  graph induced width with DFS ordering

# DFS tree and efficiency

- Depth first order is crucial for DPOP efficiency
- Finding optimal order is hard
  - Optimal  $\rightarrow$  minimize separator size
- Good heuristics:
  - Maximum Connected Node (MCN)
  - Maximum Cardinality Set (MCS) for DFS

# DFS tree Pseudotrees

- DPOP would work on any pseudotree arrangement of primal graph
- But DFS induces only a specific set of orderings:
  - Not all pseudotrees are DFS trees
- We might lose good orderings to keep computation local
  - Trade-off depends on applications

# Summary

- DCOPs, general framework to address Multi-Agent coordination
  - Many solution techniques for (relatively) large scale systems
- Complete approaches
  - Suffer from exponential element (DCOPs are hard problems)
  - ADOPT:
    - search based, asynchronous
    - Small messages but exponentially many
  - DPOP:
    - Dynamic programming based, synchronous
    - Few messages but exponentially large
    - Typically much more efficient than ADOPT

# References

**Multiagent Systems** edited by *G. Weiss* MIT Press, 2013, 2nd edition  
<http://www.the-mas-book.info/index.html>

Chapter 12 **Distributed Constraint Handling and Optimization**

*A. Farinelli, M. Vinyals, A. Rogers, and N.R. Jennings*

## FURTHER READING

### Constraint Network

- **Constraint Processing**, *R. Dechter, Morgan Kaufmann*

### ADOPT

- [Modi et al., 2005] P. J. Modi, W. Shen, M. Tambe, and M. Yokoo. **ADOPT: Asynchronous distributed constraint optimization with quality guarantees**. Artificial Intelligence Journal, (161):149-180, 2005.

### DPOP

- [Petcu 2007] A. Petcu. **A Class of Algorithms for Distributed Constraint Optimization**. PhD. Thesis No. 3942, Swiss Federal Institute of Technology (EPFL), Lausanne (Switzerland), 2007. (Chapters 2, 3 and 4)