



南开大学

NANKAI UNIVERSITY, P.R. CHINA 1919

允公允能 日新月异



恶意代码分析与防治技术

## 第13章 数据加密与解密

王志

zwang@nankai.edu.cn

南开大学 网络空间安全学院

2023-2024学年



允公允能 日新月异

# 本章知识点

- The Goal of Analyzing Encoding Algorithms
- Simple Ciphers
  - 重点: XOR、BASE64
- Common Cryptographic Algorithms
  - 难点: 信息熵Entropy
- Custom Encoding
- Decoding
  - 重点: 自解密Self-decoding



南开大学  
Nankai University



南開大學

NANKAI UNIVERSITY, P.R. CHINA 1919

允公允能 日新月異



# The Goal of Analyzing Encoding Algorithms

计算机病毒为什么要使用数据加密算法？

正常使用主观题需2.0以上版本雨课堂

作答



允公允能 日新月异

# Reasons Malware Uses Encoding

- **Hide** configuration information
  - Such as C&C domains
- Save information to a staging file
  - Before **stealing** it
- Store **strings** needed by malware
  - Decode them just before they are needed
- **Disguise** malware as a legitimate tool
  - Hide suspicious strings



南开大学  
Nankai University



允公允能 日新月异

# The Goal

- **Identify** the encoding functions
  - Understand the encoding method
- **Decode** malware secrets
  - Using the encoding knowledge



恶意代码通常会对哪些数据进行加密？

- ☒ A 恶意代码的配置信息
- ☒ B 偷取的数据、文件等
- ☒ C 恶意代码用到重要字符串
- ☒ D 伪装成正常软件所需要隐藏的信息

提交



破解恶意代码的数据加密需要完成哪些工作？

- ☒ A 识别数据加密函数
- ☒ B 分析数据加密方法
- ☐ C 找到恶意代码的文件特征码
- ☒ D 解密被加密的数据

提交







南開大學

NANKAI UNIVERSITY, P.R. CHINA 1919

允公允能 日新月異



# Simple Ciphers

简单数据加密有哪些方法？为什么要使用简单加密方法？

正常使用主观题需2.0以上版本雨课堂

作答



允公允能 日新月异

# Why Use Simple Ciphers?

- They are easily broken, but
  - They are **small**, so they fit into space-constrained environments like exploit shellcode
  - **Less obvious** than more complex ciphers
  - **Low overhead**, little impact on performance
- These are ***obfuscation***, not *encryption*
  - They make it difficult to recognize the data, but can't stop a skilled analyst



南开大学  
Nankai University



允公允能 日新月异

# Caesar Cipher

- **Shift** each letter forward 3 spaces in the alphabet

ABCDEFGHIJKLMNOPQRSTUVWXYZ

DEFGHIJKLMNOPQRSTUVWXYZABC

- Example

ATTACK AT NOON

DWWDFN DW QRRQ



南开大学  
Nankai University



允公允能 日新月异

# XOR

$$0 \text{ xor } 0 = 0$$

$$0 \text{ xor } 1 = 1$$

$$1 \text{ xor } 0 = 1$$

$$1 \text{ xor } 1 = 0$$

- Uses a key to encrypt data
- Uses one bit of data and one bit of the key at a time



对字符串“HI”进行XOR数据加密，密钥是0x3c，“HI”的ASCII码是0x48 0x49，XOR加密后的数据是？

正常使用主观题需2.0以上版本雨课堂

作答



南开大学  
Nankai University



允公允能 日新月异

# XOR

- Example: Encode HI with a key of **0x3c**

HI = 0x48 0x49 (ASCII encoding)


Data:                   0100 1000 0100 1001

Key:                    0011 1100 0011 1100

Result:       **0111 0100 0111 0101**





A	T	T	A	C	K		A	T		N	O	O	N
0x41	0x54	0x54	0x41	0x43	0x4B	0x20	0x41	0x54	0x20	0x4E	0x4F	0x4F	0x4E
													
}	h	h	}	DEL	W	FS	}	H	FS	r	s	s	r
0x7d	0x68	0x68	0x7d	0x7F	0x77	0x1C	0x7d	0x68	0x1C	0x72	0x71	0x71	0x72

*Figure 14-1. The string ATTACK AT NOON encoded with an XOR of 0x3C  
(original string at the top; encoded strings at the bottom)*





允公允能 日新月异

# XOR Reverses Itself

- Example: Encode HI with a key of 0x3c

HI = 0x48 0x49 (ASCII encoding)

Data: 0100 1000 0100 1001

Key: 0011 1100 0011 1100

$$0 \text{ xor } 0 = 0$$

$$0 \text{ xor } 1 = 1$$

$$1 \text{ xor } 0 = 1$$

$$1 \text{ xor } 1 = 0$$

- Encode it again

Result: 0111 0100 0111 0101

Key: 0011 1100 0011 1100

Data: 0100 1000 0100 1001





# Brute-Forcing XOR Encoding

- If the key is a single byte, there are only 256 possible keys
  - Error in book; this should be "a.exe"
  - PE files begin with MZ

*Example 14-1. First bytes of XOR-encoded file a.gif*

5F 48 42 12 10 12 12 12 16 12 1D 12 ED ED 12 12	_HB.....
AA 12 12 12 12 12 12 12 52 12 08 12 12 12 12 12	.....R.....
12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12	.....
12 12 12 12 12 12 12 12 12 12 12 12 12 13 12 12	.....
A8 02 12 1C 0D A6 1B DF 33 AA 13 5E DF 33 82 82	.....3..^.3..
46 7A 7B 61 32 62 60 7D 75 60 73 7F 32 7F 67 61	Fz{a2b`}u`s.2.ga





允公允能 日新月异

MZ = 0x4d 0x5a

Table 14-1. Brute-Force of XOR-Encoded Executable

XOR key value	Initial bytes of file	MZ header found?
Original	5F 48 42 12 10 12 12 12 16 12 10 12 ED ED 12	No
XOR with 0x01	5e 49 43 13 11 13 13 13 17 13 1c 13 ec ec 13	No
XOR with 0x02	5d 4a 40 10 12 10 10 10 14 10 1f 10 ef ef 10	No
XOR with 0x03	5c 4b 41 11 13 11 11 11 15 11 1e 11 ee ee 11	No
XOR with 0x04	5b 4c 46 16 14 16 16 16 12 16 19 16 e9 e9 16	No
XOR with 0x05	5a 4d 47 17 15 17 17 17 13 17 18 17 e8 e8 17	No
...	...	No
XOR with 0x12	4d 5a 50 00 02 00 00 00 04 00 0f 00 ff ff 00	Yes!





*Example 14-2. First bytes of the decrypted PE file*

4D 5A 50 00 02 00 00 00 04 00 0F 00 FF FF 00 00	MZP.....
B8 00 00 00 00 00 00 00 40 00 1A 00 00 00 00 00	.....@.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00	.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 01 00 00	.....
BA 10 00 0E 1F B4 09 CD 21 B8 01 4C CD 21 90 90	.....!...L!...
54 68 69 73 20 70 72 6F 67 72 61 6D 20 6D 75 73	This program mus





# Brute-Forcing Many Files

- Look for a common string, like "This Program"

*Table 14-2. Creating XOR Brute-Force Signatures*

## XOR key value "This program"

Original	54 68 69 73 20 70 72 6f 67 72 61 6d 20
XOR with 0x01	55 69 68 72 21 71 73 6e 66 73 60 6c 21
XOR with 0x02	56 6a 6b 71 22 72 70 6d 65 70 63 6f 22
XOR with 0x03	57 6b 6a 70 23 73 71 6c 64 71 62 6e 23
XOR with 0x04	50 6c 6d 77 24 74 76 6b 63 76 65 69 24
XOR with 0x05	51 6d 6c 76 25 75 77 6a 62 77 64 68 25
...	...
XOR with 0xFF	ab 97 96 8c df 8f 8d 90 98 8d 9e 92 df



# XOR and Nulls

- A null byte reveals the key, because
  - $0x00 \text{ xor KEY} = \text{KEY}$
- Obviously the key here is 0x12

*Example 14-1. First bytes of XOR-encoded file a.gif*

5F 48 42 12 10 12 12 12 16 12 1D 12 ED ED 12 12	_HB.....
AA 12 12 12 12 12 12 12 52 12 08 12 12 12 12	.....R.....
12 12 12 12 12 12 12 12 12 12 12 12 12 12 12	.....
12 12 12 12 12 12 12 12 12 12 12 12 12 13 12 12	.....
A8 02 12 1C 0D A6 1B DF 33 AA 13 5E DF 33 82 82	.....3..^..3..
46 7A 7B 61 32 62 60 7D 75 60 73 7F 32 7F 67 61	Fz{a2b`}u`s.2.ga





# NULL-Preserving Single-Byte XOR Encoding

- Algorithm:
  - Use XOR encoding, EXCEPT
  - If the plaintext is NULL or the key itself, **skip** the byte

*Table 14-3. Original vs. NULL-Preserving XOR Encoding Code*

Original XOR	NULL-preserving XOR
<pre>buf[i] ^= key;</pre>	<pre>if (buf[i] != 0 &amp;&amp; buf[i] != key)     buf[i] ^= key;</pre>





*Example 14-1. First bytes of XOR-encoded file a.gif*

5F 48 42 12 10 12 12 12 16 12 1D 12 ED ED 12 12	_HB.....
AA 12 12 12 12 12 12 12 52 12 08 12 12 12 12 12	.....R.....
12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12	.....
12 12 12 12 12 12 12 12 12 12 12 12 12 13 12 12	.....
A8 02 12 1C 0D A6 1B DF 33 AA 13 5E DF 33 82 82	.....3..^..3..
46 7A 7B 61 32 62 60 7D 75 60 73 7F 32 7F 67 61	Fz{a2b`}u`s.2.ga

*Example 14-3. First bytes of file with NULL-preserving XOR encoding*

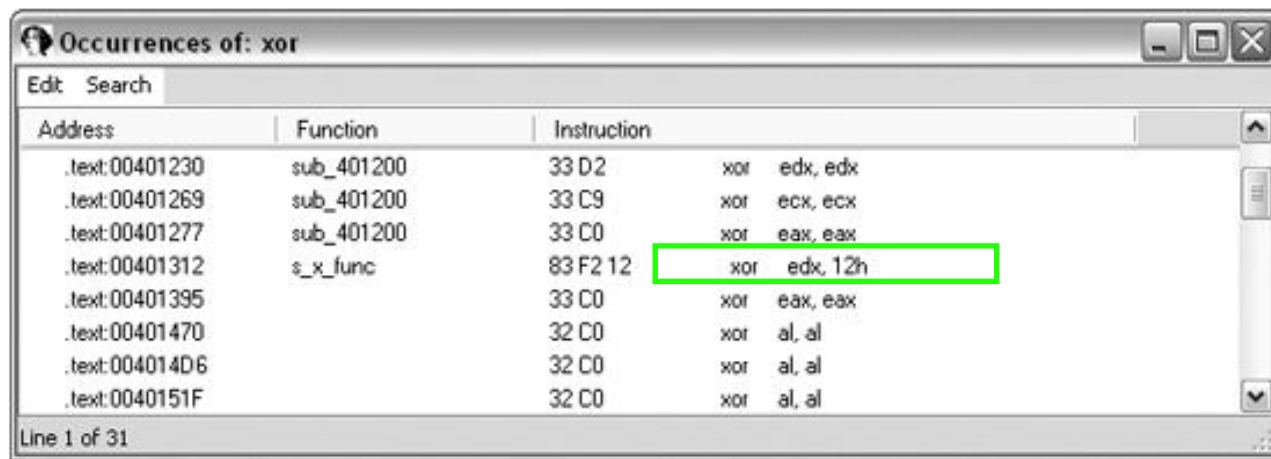
5F 48 42 00 10 00 00 00 16 00 1D 00 ED ED 00 00	_HB.....
AA 00 00 00 00 00 00 00 52 00 08 00 00 00 00 00	.....R.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00	.....
00 00 00 00 00 00 00 00 00 00 00 00 00 13 00 00	.....
A8 02 00 1C 0D A6 1B DF 33 AA 13 5E DF 33 82 82	.....3..^..3..
46 7A 7B 61 32 62 60 7D 75 60 73 7F 32 7F 67 61	Fz{a2b`}u`s.2.ga





# Identifying XOR Loops in IDA

- Small loops with an XOR instruction inside
  1. Start in "IDA View" (seeing code)
  2. Click **Search, Text**
  3. Enter **xor** and **Find all occurrences**





允公允能 日新月异

# Three Forms of XOR

- XOR a **register with itself**, like `xor edx, edx`
  - Innocent, a common way to zero a register
- XOR a **register or memory reference with a constant**
  - May be an encoding loop, and key is the constant
- XOR a **register or memory reference with a different register or memory reference**
  - May be an encoding loop, key less obvious



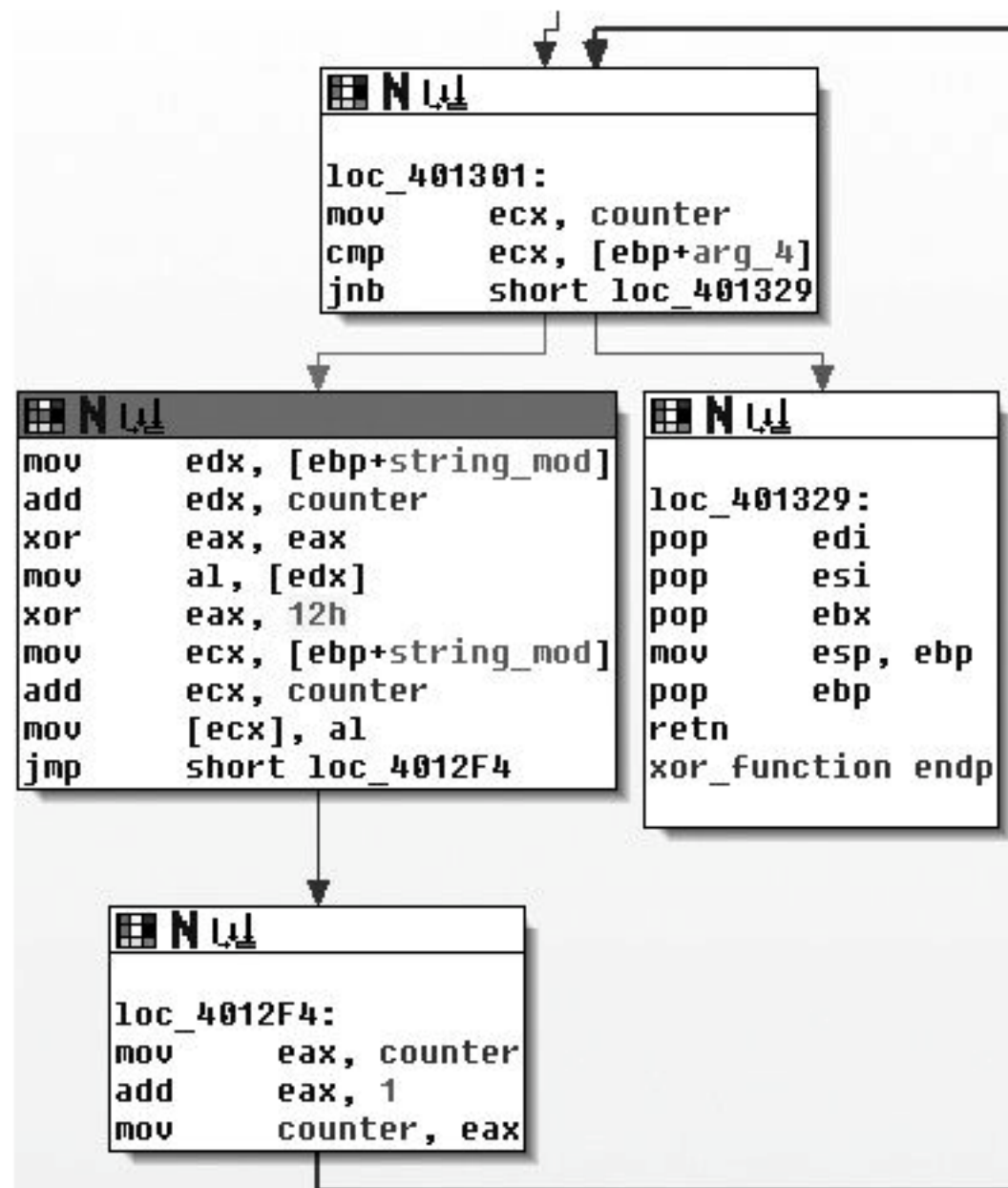


Figure 14-3. Graphical view of single-byte XOR loop



*Table 14-4. Additional Simple Encoding Algorithms*

Encoding Description scheme	
ADD, SUB	Encoding algorithms can use ADD and SUB for individual bytes in a manner that is similar to XOR. ADD and SUB are not reversible, so they need to be used in tandem (one to encode and the other to decode).
ROL, ROR	Instructions rotate the bits within a byte right or left. Like ADD and SUB, these need to be used together since they are not reversible.
ROT	This is the original Caesar cipher. It's commonly used with either alphabetical characters (A-Z and a-z) or the 94 printable characters in standard ASCII.
Multibyte	Instead of a single byte, an algorithm might use a longer key, often 4 or 8 bytes in length. This typically uses XOR for each block for convenience.
Chained or loopback	This algorithm uses the content itself as part of the key, with various implementations. Most commonly, the original key is applied at one side of the plaintext (start or end), and the encoded output character is used as the key for the next character.



讨论：大家经常见到Base64算法，尤其在Web和email的数据文件中，Base64算法是否可以应用到恶意代码的数据加解密？

作答



允公允能 日新月异

# Base64

- Converts 6 bits into one character in a 64-character alphabet
- There are a few versions, but all use these 62 characters:

ABCDEFGHIJKLMNOPQRSTUVWXYZ

abcdefghijklmnopqrstuvwxyz

0123456789

- MIME uses + and /
  - Also = to indicate padding



南开大学  
Nankai University



```
Content-Type: multipart/alternative;  
    boundary="_002_4E36B98B966D7448815A3216ACF82AA201ED633ED1MBX3THNDRBIRD_"  
MIME-Version: 1.0  
--_002_4E36B98B966D7448815A3216ACF82AA201ED633ED1MBX3THNDRBIRD_  
Content-Type: text/html; charset="utf-8"  
Content-Transfer-Encoding: base64
```

SWYgeW91IGFyZSBYZWFKaW5nIHRoaXMsIHlvdSBwcm9lYWJseSBzaG91bGQganVzdCBza2lwIHRoaX  
MgY2hhcHRlcilBhbmQgZ28gdG8gdGhlIG5leHQgb25lLiBEbyB5b3UgcmlvbmVhZGx5IGhhdmUgdGhlIHRp  
bWUgdG8gdHlwZSB0aGlzIHdob2xliHN0cmVudG91IGFyZSBvYnZpb3VzbHkgdGFsZW50ZW  
QuIE1heWJlIHlvdSBzaG91bGQgY29udGFjdCB0aGUgYXV0aG9ycyBhbmQgc2VlIGlmIH





# Transforming Data to Base64

- Use 3-byte chunks (24 bits)
- Break into four 6-bit fields
- Convert each to Base64

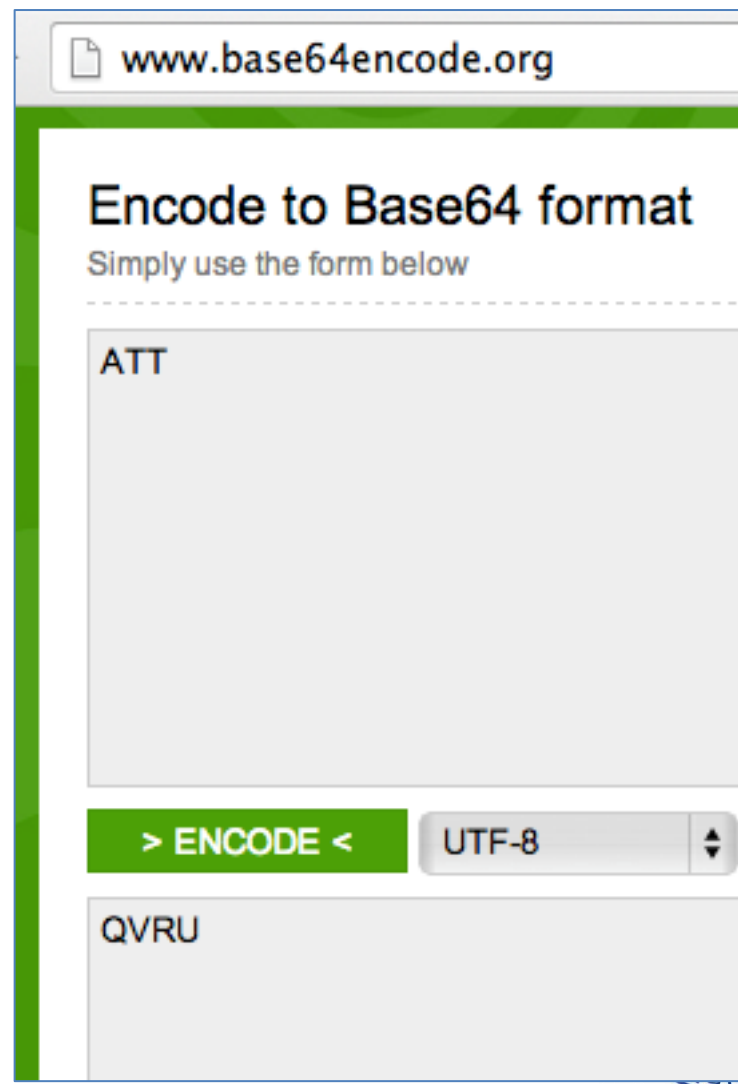
A								T								T							
0x4				0x1				0x5				0x4				0x5				0x4			
0	1	0	0	0	0	0	1	0	1	0	1	0	1	0	0	0	1	0	1	0	1	0	0
16				21				17				20											
Q				V				R				U											

Figure 14-4. Base64 encoding of ATT



# base64encode.org base64decode.org

- 3 bytes encode to 4  
Base64 characters



The screenshot shows a web browser window with the address bar displaying "www.base64encode.org". The page has a green header bar. Below the header, the main content area is titled "Encode to Base64 format" with the subtitle "Simply use the form below". There is a large text input field containing the text "ATT". Below the input field, there is a green button labeled "> ENCODE <" and a dropdown menu currently set to "UTF-8". Below the button and dropdown, there is another text input field containing the text "QVRU".





允公允能 日新月异

# Padding

- If input had only 2 characters, an “=” is appended

### Encode to Base64 format

Simply use the form below

AT

> ENCODE <

UTF-8

QVQ=





# Padding

- If input had only 1 character, == is appended

**Encode to Base64 format**  
Simply use the form below

---

A

> ENCODE <

UTF-8

QQ==





允公允能 日新月异

# Example

- URL and cookie are Base64-encoded

*Example 14-5. Sample malware traffic*

```
GET /X29tbVEuYC8=/index.htm
```

```
User-Agent: Mozilla/4.0 (compatible; MSIE 7.0; Windows NT 5.1)
```

```
Host: www.practicalmalwareanalysis.com
```

```
Connection: Keep-Alive
```

```
Cookie: Ym90NTQxNjQ
```

```
GET /c2UsYi1kYWM0cnUjdFlvbiAjb21wbFU0YP==/index.htm
```

```
User-Agent: Mozilla/4.0 (compatible; MSIE 7.0; Windows NT 5.1)
```

```
Host: www.practicalmalwareanalysis.com
```

```
Connection: Keep-Alive
```

```
Cookie: Ym90NTQxNjQ
```



南开大学  
Nankai University



# Cookie: Ym90NTQxNjQ

- This has 11 characters—padding is omitted
- Some Base64 decoders will fail, but this one just automatically adds the missing padding

Decode from Base64 format  
Simply use the form below

Ym90NTQxNjQ

< DECODE > UTF-8 (Y

bot54164





允公允能 日新月异

# Finding the Base64 Function

- Look for this "indexing string"

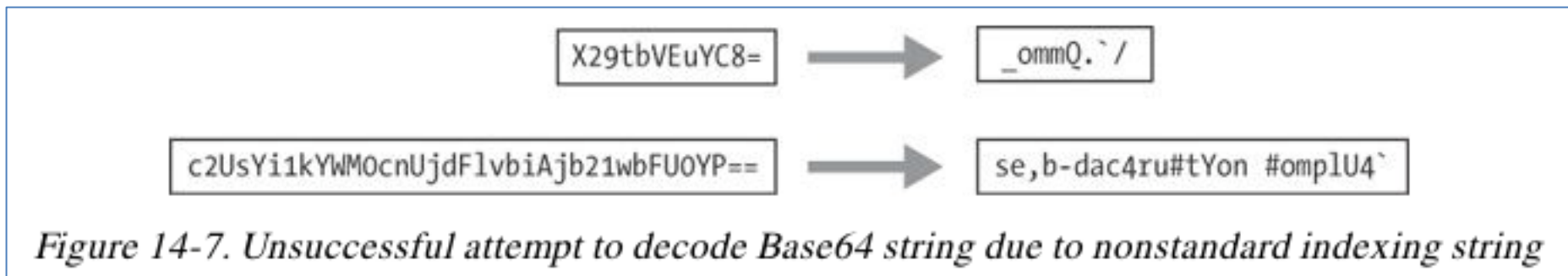
```
ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz  
0123456789+/'
```

- Look for a lone padding character (typically =) hard-coded into the encoding function





# Decoding the URLs



- Custom **indexing string**

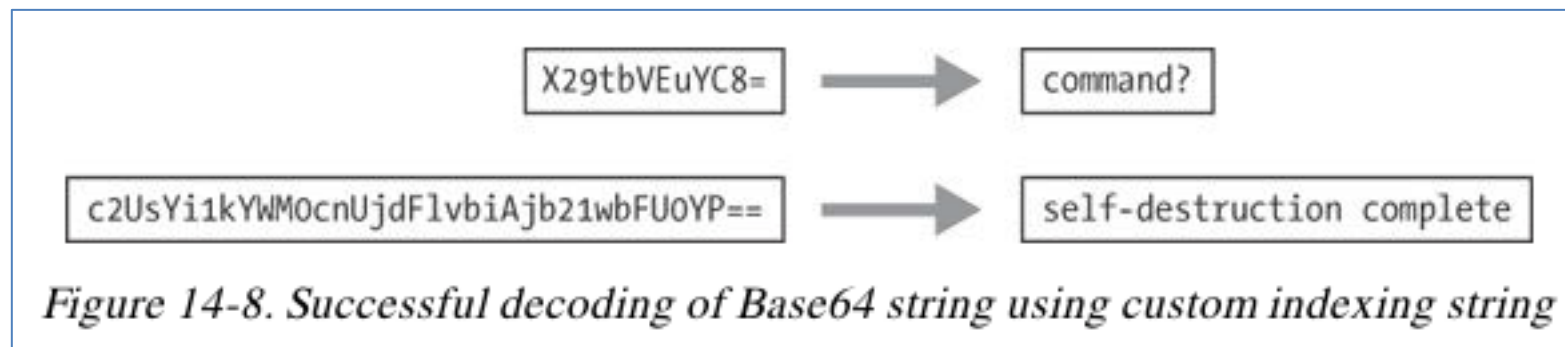
aABCDEFGHIJKLMNOPQRSTUVWXYZbcdefghijklmnopqrs  
tuvwxyz0123456789+/  
=

- Look for a lone **padding character** (typically =) hard-coded into the encoding function





允公允能 日新月异





以下哪些指令可以被恶意代码用于数据加密和解密？

- ☒ A XOR
- ☒ B ADD、SUB
- ☒ C ROL、ROR
- ☐ D MOV

提交



恶意代码常使用哪些简单的数据加解密方案？

☒ A 凯撒密码

☒ B XOR

☒ C Base64

☐ D MD5

提交





南開大學

NANKAI UNIVERSITY, P.R. CHINA 1919

允公允能 日新月異



# Common Cryptographic Algorithms

有哪些常见的高强度加密算法？恶意代码中使用这些算法有哪些优点和缺点？

正常使用主观题需2.0以上版本雨课堂

作答



允公允能 日新月异

# Cryptography

- Cryptography vs. Cipher ?
- Cryptography applications?





允公允能 日新月异

# Strong Cryptography

- Strong enough to **resist brute-force attacks**
  - Ex: SSL, AES, etc.
- Disadvantages of strong encryption
  - **Large** cryptographic libraries required
  - May make code **less portable**
  - Standard cryptographic libraries are **easily detected**
    - Via function imports, function matching, or identification of cryptographic constants
  - Symmetric encryption requires a way to hide the **key**



# Recognizing Strings and Imports

- Strings found in malware encrypted with OpenSSL

```
OpenSSL 1.0.0a
SSLv3 part of OpenSSL 1.0.0a
TLSv1 part of OpenSSL 1.0.0a
SSLv2 part of OpenSSL 1.0.0a
You need to read the OpenSSL FAQ,
http://www.openssl.org/support/faq.html
%s(%d): OpenSSL internal error, assertion failed: %s
AES for x86, CRYPTOAMS by <appro@openssl.org>
```

# Recognizing Strings and Imports

- Microsoft crypto functions usually start with **Crypt** or **CP** or **Cert**

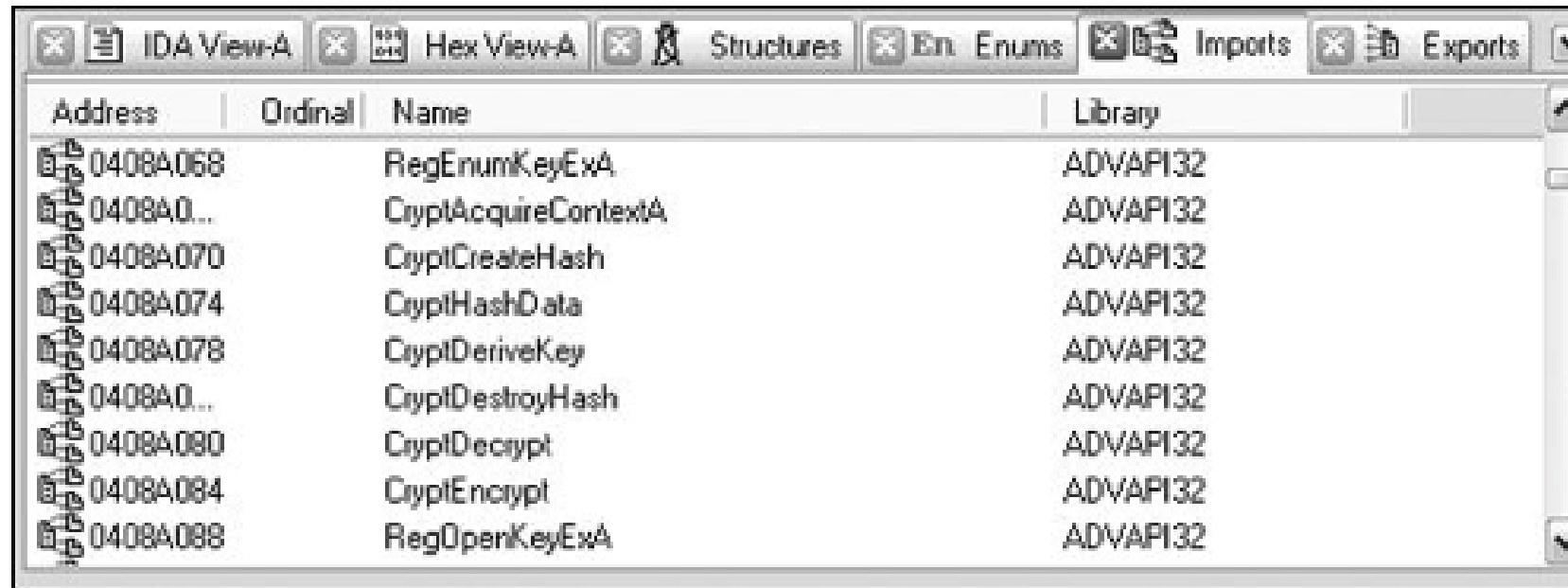


Figure 14-9. IDA Pro imports listing showing cryptographic functions





允公允能 日新月异

# Searching for Cryptographic Constants

- IDA Pro's FindCrypt2 Plug-in
  - Finds *magic constants* (binary signatures of crypto routines)
  - Cannot find RC4 or IDEA routines because they don't use a magic constant
  - **RC4** is commonly used in malware because it's small and easy to implement

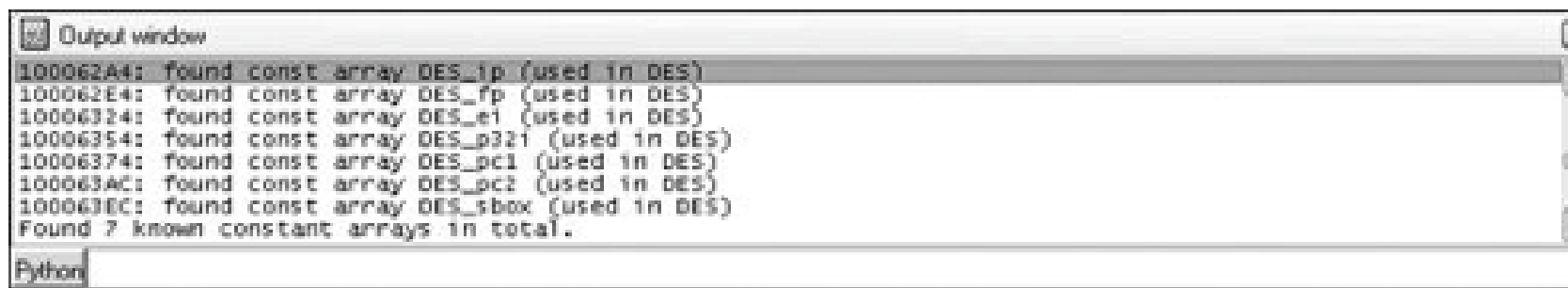


南开大学  
Nankai University



# FindCrypt2

- Runs automatically on any new analysis
- Can be run manually from the Plug-In Menu



```
Output window
100062A4: found const array DES_ip (used in DES)
100062E4: found const array DES_fp (used in DES)
10006324: found const array DES_e1 (used in DES)
10006354: found const array DES_p321 (used in DES)
10006374: found const array DES_pc1 (used in DES)
100063AC: found const array DES_pc2 (used in DES)
100063EC: found const array DES_sbox (used in DES)
Found 7 known constant arrays in total.
Python
```

*Figure 14-10. IDA Pro FindCrypt2 output*

# Krypto ANALyzer (PEiD Plug-in)

- Download from link Ch 13d
- Has wider range of constants than FindCrypt2
  - More false positives
- Also finds **Base64 tables** and crypto function imports

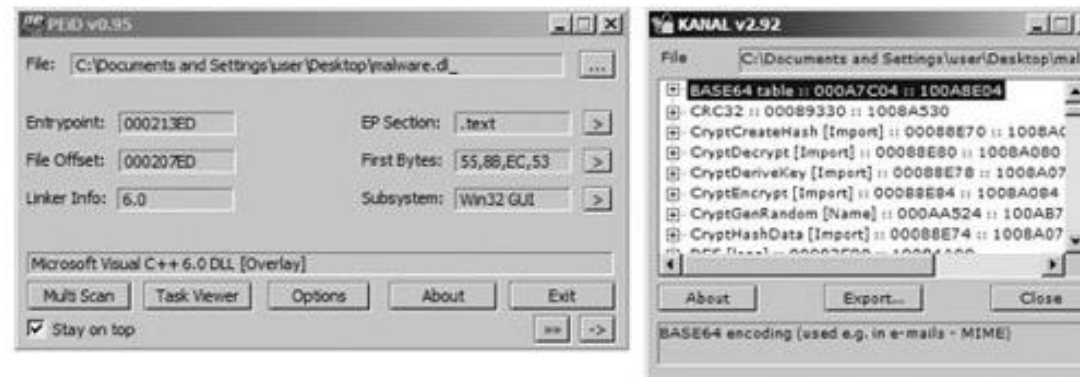


Figure 14-11. PEiD and Krypto ANALyzer (KANAL) output



# Entropy

- Entropy measures disorder

$$P(X = i) = p_i > 0 \quad (i = 1, 2, \dots, n) \quad , \quad \sum_{i=1}^n p_i = 1$$

$$H(X) = - \sum_{i=1}^n p_i \log_2 p_i \leq - \sum_{i=1}^n \frac{1}{n} \log_2 \frac{1}{n} = \log_2 n$$

- 最大熵定理，等概率场的平均不确定性最大
- The number of occurrences of each byte from 0 to 255
  - If all the bytes are equally likely, the entropy is 8 (maximum disorder)
  - If all the bytes are the same, the entropy is 0



# Searching for High-Entropy Content

- IDA Pro Entropy Plugin
- Finds regions of high entropy, indicating encryption (or compression)

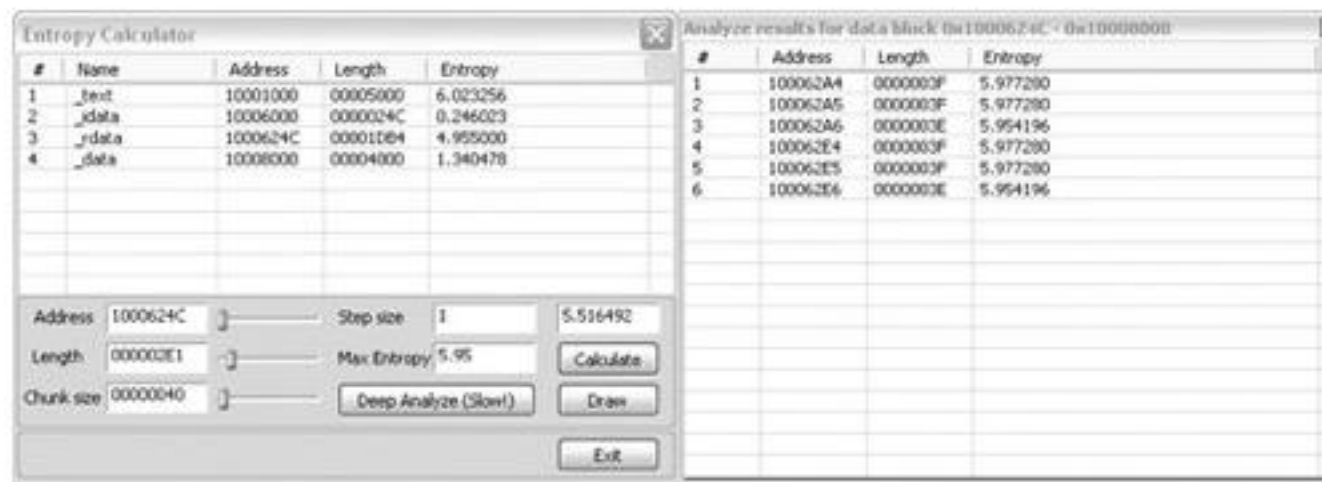


Figure 14-12. IDA Pro Entropy Plugin



允公允能 日新月异

# Recommended Parameters

- Chunk size: 64      Max. Entropy: 5.95
  - Good for finding many constants,
  - Including **Base64-encoding** strings (entropy 6)
- Chunk size: 256      Max. Entropy: 7.9
  - Finds **very random regions**



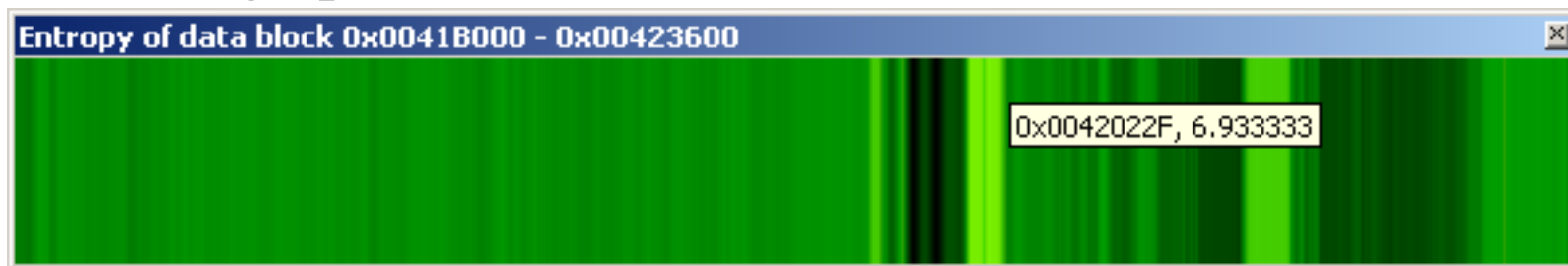
南开大学  
Nankai University



允公允能 日新月异

# Entropy Graph

- IDA Pro Entropy Plugin
  - Download from link Ch 13g
  - Use StandAlone version
  - Double-click region, then Calculate, Draw
  - Lighter regions have high entropy
  - Hover over graph to see numerical value





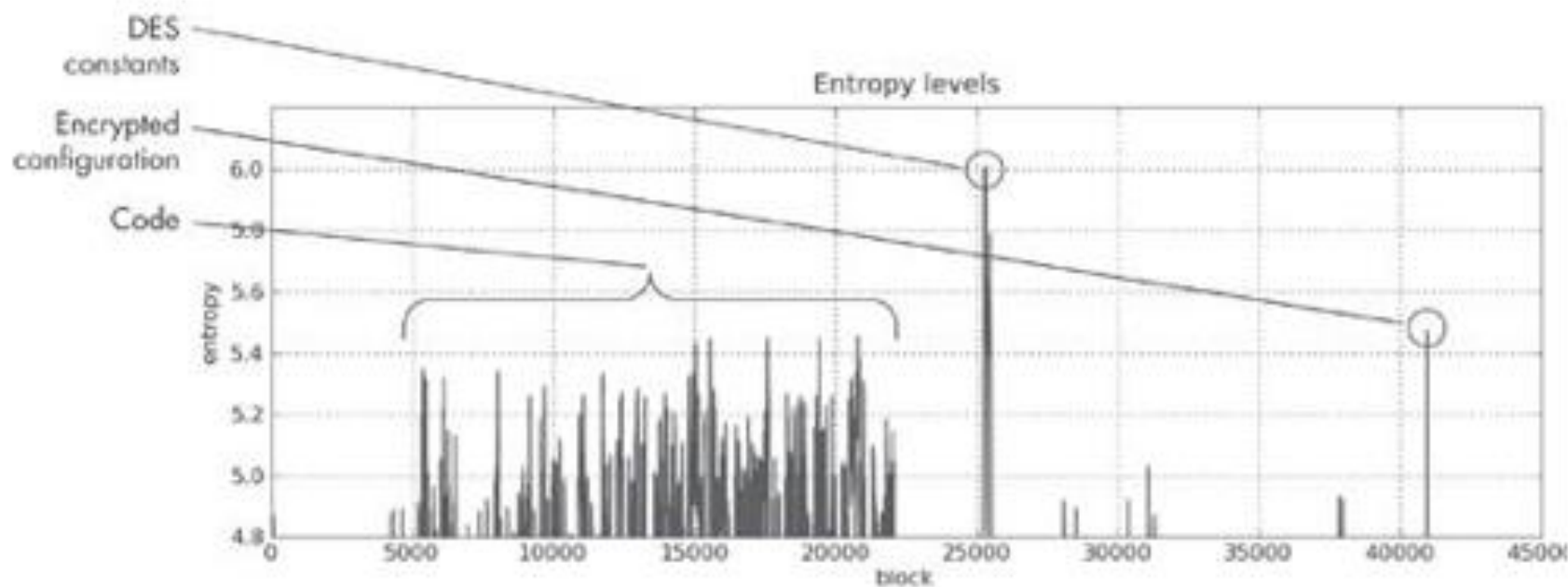


Figure 14-13. Entropy graph for a malicious executable



使用库函数进行数据加密有哪些缺点？

- ☒ A 增加恶意代码体积
- ☒ B 降低移动性
- ☒ C 容易被发现
- ☒ D 对称加密需要隐藏密钥

提交



有哪些检测恶意代码是否使用库函数加密数据的方法?

- ☒ A Entropy
- ☒ B 加密相关的常量
- ☒ C 加密相关函数的名字
- ☒ D 加密相关库的名字

Entropy值高的区域还是低的区域有可能是加密数据？

- ☒ A Entropy值高的区域
- ☐ B Entropy值低的区域



南開大學

NANKAI UNIVERSITY, P.R. CHINA 1919

允公允能 日新月異



# Custom Encoding

讨论：基于标准的加密方法，恶意代码是否可以设计自定义的加密方法？自定义的加密方法有哪些优点？

作答



允公允能 日新月异

# Homegrown Encoding Schemes

- Examples
  - One round of XOR, then Base64
  - Custom algorithm, possibly similar to a published cryptographic algorithm



南開大學  
Nankai University



# Identifying Custom Encoding

*Example 14-6. First bytes of an encrypted file*

88 5B D9 02 EB 07 5D 3A 8A 06 1E 67 D2 16 93 7F	.[....]:...g....
43 72 1B A4 BA B9 85 B7 74 1C 6D 03 1E AF 67 AF	Cr.....t.m...g.
98 F6 47 36 57 AA 8E C5 1D 70 A5 CB 38 ED 22 19	..G6W....p..8..".
86 29 98 2D 69 62 9E C0 4B 4F 8B 05 A0 71 08 50	.)..-ib..KO...q.P
92 A0 C3 58 4A 48 E4 A3 0A 39 7B 8A 3C 2D 00 9E	...XJH...9{.<-..

- This sample makes a bunch of 700 KB files
- Figure out the encoding from the code
- Find **CreateFileA** and **WriteFileA**
  - In function **sub\_4011A9**
- Uses XOR with a pseudorandom stream



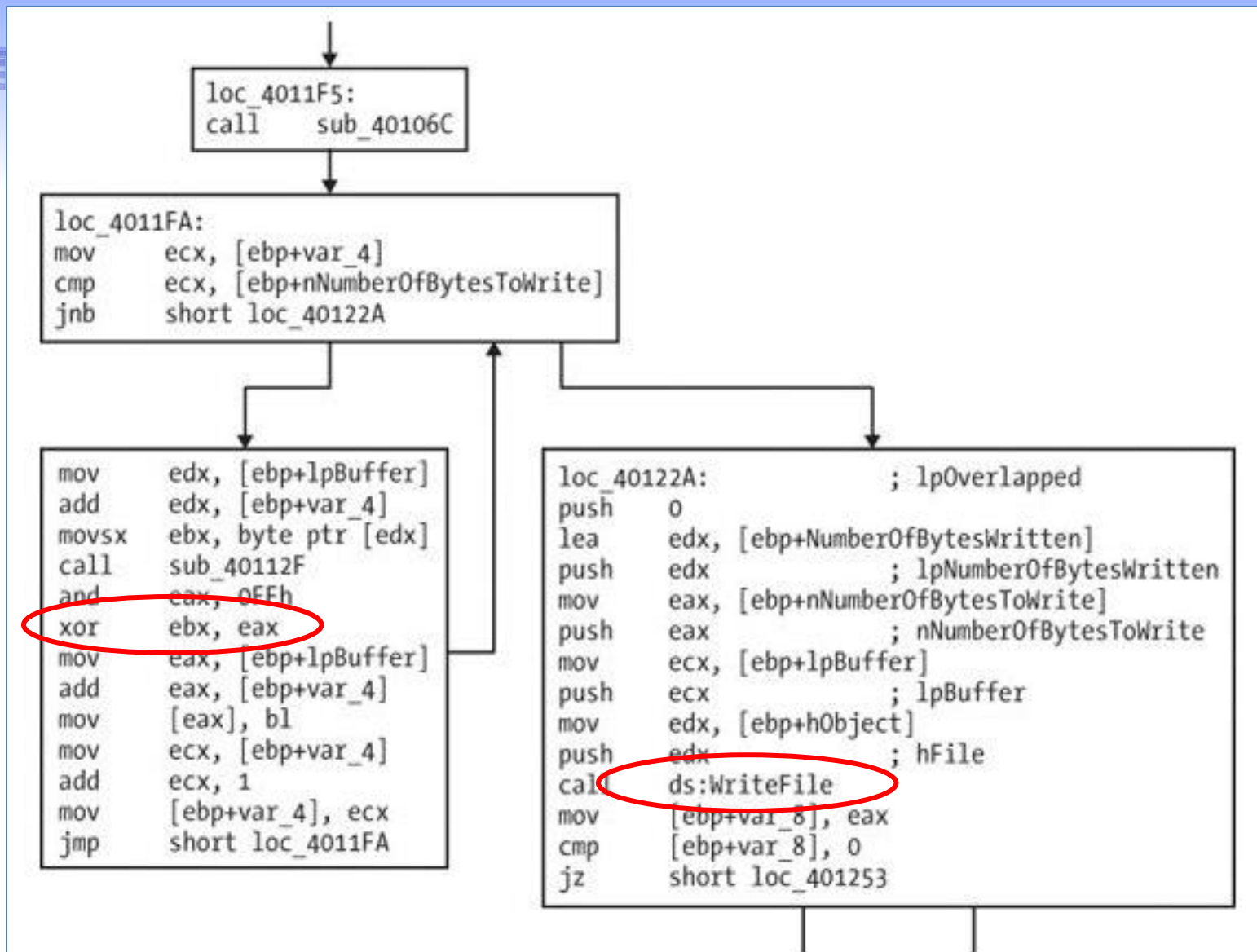


Figure 14-14. Function graph showing an encrypted write







允公允能 日新月异

# Advantages of Custom Encoding to the Attacker

- Can be small and **nonobvious**
- Harder to reverse-engineer
  - key
  - decryption function



南开大学  
Nankai University

自定义的数据加密算法比使用标准库函数加密有哪些优点？

- ☒ A 增加逆向工程的难度
- ☒ B 结合简单加密方案，执行速度更快，体积小
- ☒ C 可以隐藏密钥
- ☒ D 不能直接找到解密函数

提交





南開大學

NANKAI UNIVERSITY, P.R. CHINA 1919

允公允能 日新月異

Decoding

针对恶意代码使用的数据加密方法，逆向分析恶意代码时有哪些有效的解密方法？

作答



允公允能 日新月异

# Two Methods

- Reprogram the functions
- Use the functions in the malware itself





允公允能 日新月异

# Self-Decoding

- Stop the malware in a debugger with data decoded
- Isolate the decryption function and set a **breakpoint** directly after it
- BUT sometimes you can't figure out how to stop it with the data you need decoded





# Manual Programming of Decoding Functions

- Standard functions may be available

*Example 14-7. Sample Python Base64 script*

```
import string
import base64

example_string = 'VGhpcyBpcyBhIHRlc3Qgc3RyaW5n'
print base64.decodestring(example_string)
```

*Example 14-8. Sample Python NULL-preserving XOR script*

```
def null_preserving_xor(input_char, key_char):
    if (input_char == key_char or input_char == chr(0x00)):
        return input_char
    else:
        return chr(ord(input_char) ^ ord(key_char))
```



*Example 14-9. Sample Python custom Base64 script*

```
import string
import base64

s = ""
custom = "9ZABCDEFGHIJKLMNOPQRSTUVWXYabcdefghijklmnopqrstuvwxyz012345678+/"
Base64 = "ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz0123456789+/"

ciphertext = 'TEgobxZobxZgGFPkb20='

for ch in ciphertext:
    if (ch in Base64):
        s = s + Base64[string.find(custom,str(ch))]
    elif (ch == '='):
        s += '='

result = base64.decodestring(s)
```







允公允能 日新月异

# PyCrypto Library

- Good for standard algorithms

*Example 14-10. Sample Python DES script*

```
from Crypto.Cipher import DES
import sys

obj = DES.new('password',DES.MODE_ECB)
cfile = open('encrypted_file','r')
cbuf = f.read()
print obj.decrypt(cbuf)
```



南开大学  
Nankai University



允公允能 日新月异

# How to Decrypt Using Malware

1. Set up the malware in a debugger.
2. Prepare the encrypted file for reading and prepare an output file for writing.
3. Allocate memory inside the debugger so that the malware can reference the memory.
4. Load the encrypted file into the allocated memory region.
5. Set up the malware with appropriate variables and arguments for the encryption function.
6. Run the encryption function to perform the encryption.
7. Write the newly decrypted memory region to the output file.



南开大学  
Nankai University

### Example 14-12. ImmDbg sample decryption script

```
import immlib

def main ():
    imm = immlib.Debugger()
    cfile = open("C:\\encrypted_file","rb") # Open encrypted file for read
    pfile = open("decrypted_file", "w")    # Open file for plaintext
    buffer = cfile.read()                  # Read encrypted file into buffer
    sz = len(buffer)                        # Get length of buffer
    membuf = imm.remoteVirtualAlloc(sz)     # Allocate memory within debugger
    imm.writeMemory(membuf,buffer)          # Copy into debugged process's memory

    imm.setReg("EIP", 0x004011A9)           # Start of function header
    imm.setBreakpoint(0x004011b7)          # After function header
    imm.Run()                              # Execute function header

    regs = imm.getRegs()
    imm.writeLong(regs["EBP"]+16, sz)       # Set NumberOfBytesToWrite stack variable
    imm.writeLong(regs["EBP"]+8, membuf)    # Set lpBuffer stack variable

    imm.setReg("EIP", 0x004011f5)           # Start of crypto
    imm.setBreakpoint(0x0040122a)          # End of crypto loop
    imm.Run()                              # Execute crypto loop

    output = imm.readMemory(membuf, sz)    # Read answer
    pfile.write(output)                    # Write answer
```



针对恶意代码的加密数据，有哪些可行的解密方法？

A

等恶意代码自解密Self-decoding

B

编写解密函数

C

通过插装控制恶意代码来解密

D

暴力破解

提交





南开大学

NANKAI UNIVERSITY, P.R. CHINA 1919

允公允能 日新月异



# 勒索病毒

# Data Encrypted for Impact

## Data Encrypted for Impact

Adversaries may encrypt data on target systems or on large numbers of systems in a network to interrupt availability to system and network resources. They can attempt to render stored data inaccessible by encrypting files or data on local and remote drives and withholding access to a decryption key. This may be done in order to extract monetary compensation from a victim in exchange for decryption or a decryption key (ransomware) or to render data permanently inaccessible in cases where the key is not saved or transmitted.<sup>[1][2][3][4]</sup>

In the case of ransomware, it is typical that common user files like Office documents, PDFs, images, videos, audio, text, and source code files will be encrypted (and often renamed and/or tagged with specific file markers). Adversaries may need to first employ other behaviors, such as [File and Directory Permissions Modification](#) or [System Shutdown/Reboot](#), in order to unlock and/or gain access to manipulate these files.<sup>[5]</sup> In some cases, adversaries may encrypt critical system files, disk partitions, and the MBR.<sup>[3]</sup>

To maximize impact on the target organization, malware designed for encrypting data may have worm-like features to propagate across a network by leveraging other attack techniques like [Valid Accounts](#), [OS Credential Dumping](#), and [SMB/Windows Admin Shares](#).<sup>[2][3]</sup> Encryption malware may also leverage [Internal Defacement](#), such as changing victim wallpapers, or otherwise intimidate victims by sending ransom notes or other messages to connected printers (known as "print bombing").<sup>[6]</sup>

In cloud environments, storage objects within compromised accounts may also be encrypted.<sup>[7]</sup>

ID: T1486

Sub-techniques: No sub-techniques

① **Tactic:** [Impact](#)

① **Platforms:** IaaS, Linux, Windows, macOS

① **Impact Type:** Availability

**Contributors:** ExtraHop; Harshal Tupsamudre, Qualys; Mayuresh Dani, Qualys; Oleg Kolesnikov, Securonix; Travis Smith, Qualys

**Version:** 1.4

**Created:** 15 March 2019

**Last Modified:** 16 June 2022

[Version Permalink](#)

<https://attack.mitre.org/techniques/T1486/>







# Data Encrypted for Impact

## Procedure Examples

ID	Name	Description
G0082	APT38	APT38 has used Hermes ransomware to encrypt files with AES256. <sup>[8]</sup>
G0096	APT41	APT41 used a ransomware called Encryptor RaaS to encrypt files on the targeted systems and provide a ransom note to the user. <sup>[9]</sup>
S0640	Avaddon	Avaddon encrypts the victim system using a combination of AES256 and RSA encryption schemes. <sup>[10]</sup>
S1053	AvosLocker	AvosLocker has encrypted files and network resources using AES-256 and added an .avos, .avos2, or .AvosLinux extension to filenames. <sup>[11][12][13][14]</sup>
S0638	Babuk	Babuk can use ChaCha8 and ECDH to encrypt data. <sup>[15][16][17][18]</sup>
S0606	Bad Rabbit	Bad Rabbit has encrypted files and disks using AES-128-CBC and RSA-2048. <sup>[19]</sup>
S0570	BitPaymer	BitPaymer can import a hard-coded RSA 1024-bit public key, generate a 128-bit RC4 key for each file, and encrypt the file in place, appending .locked to the filename. <sup>[20]</sup>
S1070	Black Basta	Black Basta can encrypt files with the ChaCha20 cypher and using a multithreaded process to increase speed. <sup>[21][22][23][24][25][26][27][28][29]</sup>
S1068	BlackCat	BlackCat has the ability to encrypt Windows devices, Linux devices, and VMWare instances. <sup>[30]</sup>
C0015	C0015	During C0015, the threat actors used Conti ransomware to encrypt a compromised network. <sup>[31]</sup>
C0018	C0018	During C0018, the threat actors used AvosLocker ransomware to encrypt files on the compromised network. <sup>[13][32]</sup>
S0611	Clop	Clop can encrypt files using AES, RSA, and RC4 and will add the ".clon" extension to encrypted files. <sup>[33][34][35]</sup>
S0575	Conti	Conti can use CreateIoCompletionPort(), PostQueuedCompletionStatus(), and GetQueuedCompletionPort() to rapidly encrypt files, excluding those with the extensions of .exe, .dll, and .lnk. It has used a different AES-256 encryption key per file with a bundled RAS-4096 public encryption key that is unique for each victim. Conti can use "Windows Restart Manager" to ensure files are unlocked and open for encryption. <sup>[36][5][37][38][31]</sup>

<https://attack.mitre.org/techniques/T1486/>



讨论：针对勒索病毒，有哪些缓解和检测方法？

作答





# Data Encrypted for Impact

## Mitigations

ID	Mitigation	Description
M1040	Behavior Prevention on Endpoint	On Windows 10, enable cloud-delivered protection and Attack Surface Reduction (ASR) rules to block the execution of files that resemble ransomware. <sup>[97]</sup>
M1053	Data Backup	Consider implementing IT disaster recovery plans that contain procedures for regularly taking and testing data backups that can be used to restore organizational data. <sup>[98]</sup> Ensure backups are stored off system and is protected from common methods adversaries may use to gain access and destroy the backups to prevent recovery. Consider enabling versioning in cloud environments to maintain backup copies of storage objects. <sup>[99]</sup>

## Detection

ID	Data Source	Data Component	Detects
DS0010	Cloud Storage	Cloud Storage Modification	Monitor for changes made in cloud environments for events that indicate storage objects have been anomalously modified.
DS0017	Command	Command Execution	Monitor executed commands and arguments for actions involved in data destruction activity, such as vssadmin, wbadmin, and bcdedit
DS0022	File	File Creation	Monitor for newly constructed files in user directories.
		File Modification	Monitor for changes made to files in user directories.
DS0033	Network Share	Network Share Access	Monitor for unexpected network shares being accessed on target systems or on large numbers of systems.
DS0009	Process	Process Creation	Monitor for newly constructed processes and/or command-lines involved in data destruction activity, such as vssadmin, wbadmin, and bcdedit.





允公允能 日新月异

# 本章知识点

- The Goal of Analyzing Encoding Algorithms
- Simple Ciphers
  - 重点: XOR、BASE64
- Common Cryptographic Algorithms
  - 难点: 信息熵Entropy
- Custom Encoding
- Decoding
  - 重点: 自解密Self-decoding





南开大学

NANKAI UNIVERSITY, P.R. CHINA 1919

允公允能 日新月异



恶意代码分析与防治技术

## 第13章 数据加密与解密

王志

zwang@nankai.edu.cn

南开大学 网络空间安全学院

2023-2024学年