

# SAT Solvers

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# What is a SAT Solver

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- A **SAT solver** is a computer program that aims to solve the **Boolean satisfiability problem**.
- Input: A Boolean expression (We can use a CNF)
- Output: A Boolean to show the selected Boolean expression is satisfiable or not. Also, if the expression is satisfiable we can output an assignment that satisfies the expression.

# Why Not A True Table

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- In large examples, we are unable to try a million random guesses, to return the correct output.
- Algorithms:
  - 1) DPLL
  - 2) CDLL



# Input Integration

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- We can convert any Boolean expression to CNF using the CNF Algorithm.
- So we can have a CNF as input. (If not, we can convert it to CNF)

- A CNF:

$\text{AND}((\text{OR } x1 \ x2) \ (\text{OR NOT}(x3) \ x4) \ (\text{OR NOT}(x2) \ x3))$

$\Rightarrow [[1, 2], [-3, 4], [-2, 3]]$

# DPLL

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1. "Guess" a variable
2. Find all unit clauses created from the last assignment and assign the needed value
3. Iteratively retry step 2 until there is no change (found transitive closure)
4. If the current assignment cannot yield true for all clauses - fold back from recursion and retry a different assignment
5. If it can - "guess" another variable (recursively invoke and return to 1)



```
def DPLL(cnf):
    if not cnf:
        return True
    if any(not clause for clause in cnf):
        return False
    for clause in cnf:
        if len(clause) == 1:
            literal = clause[0]
            cnf = unitPropagate(cnf, literal)
            return DPLL(cnf)
    literal = cnf[0][0]
    cnf = unitPropagate(cnf, literal)
    if DPLL(cnf):
        return True
    cnf = unitPropagate(cnf, -literal)
    return DPLL(cnf)

def unitPropagate(cnf, literal):
    newCnf = [clause for clause in cnf if literal not in clause]
    newCnf = [list(filter(lambda l: l != -literal, clause)) for clause in newCnf]
    return newCnf
```

```
cnf = [
    [-2, 1, 3],
    [-2, 1, 3],
    [2, 1, 3],
]
result = DPLL(cnf)
print(result)
```



# CDCL

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.The Conflict-Driven Clause Learning (CDCL) algorithm is a method used to solve the Boolean satisfiability problem (SAT). Here are the main steps of the CDCL algorithm:

- 1. Select a variable and assign True or False:** This is referred to as the decision state.
- 2. Apply Boolean constraint propagation (unit propagation):** If an unsatisfied clause has all but one of its literals or variables evaluated at False, then the free literal must be True in order for the clause to be True.



# CDCL

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**3. Build the implication graph:** This graph represents the implications of the assignments made.

**4. Check for conflicts:** If there is any conflict, find the cut in the implication graph that led to the conflict.

**5. Derive a new clause and backtrack:** Create a new clause which is the negation of the assignments that led to the conflict. Then, non-chronologically backtrack (also known as “back jump”) to the appropriate decision level, where the first-assigned variable involved in the conflict was assigned.

**6. Repeat the process:** Continue from step 1 until all variable values are assigned.

.This algorithm is widely used in various SAT solvers due to its efficiency in handling complex CNF formulas.



# CDCL

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Sample codes for CDCL are too long thus it was decided to put a reference link:

- CDCL Code:

[sat-solver/implementation/CDCL.py at master · Mr-Ahmadi/sat-solver\(github.com\)](https://github.com/Mr-Ahmadi/sat-solver/blob/master/sat-solver/implementation/CDCL.py)

- Also DPLL Code (Too easy):

[sat-solver/implementation/DPLL.py at master · Mr-Ahmadi/sat-solver\(github.com\)](https://github.com/Mr-Ahmadi/sat-solver/blob/master/sat-solver/implementation/DPLL.py)



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Bye :)