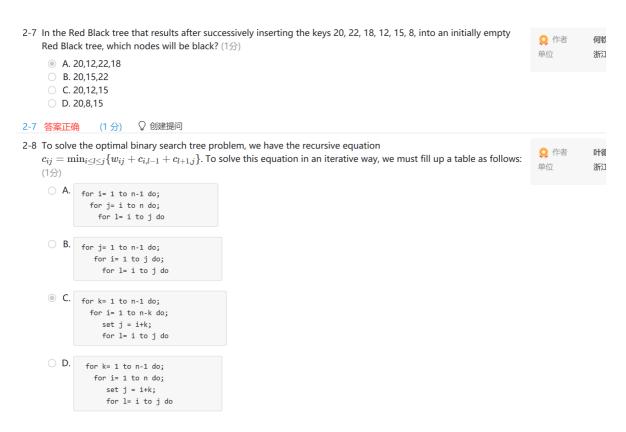
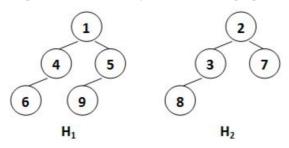
1-1	Let S be the set of activities in Activity Selection Problem. Then the latest start activity ai must be included in some maximum-size subset of mutually compatible activities of S. (1 $\%$)									
	⊚ T	F								
1-1	答案正确	(1分)								
1-2	In a red-black tree, the number of internal nodes in the subtree rooted at x is no more than $2^{bh(x)}-1$ where $bh(x)$ is the black-height of x . (1 $\cancel{\bigcirc}$)									
		F								
1-2	答案错误(〕 (0分)	♀ 创建提问							
1-3	Insert 1, 2, 3, 4, 5, and 6 one by one into an initially empty AVL tree. Then the preorder traversal sequence of the resulting tree must be $\{4, 2, 1, 3, 5, 6\}$. $(1\bigcirc)$									
	○ T	○ T								
1-3	答案错误(〕 (0分)	○ 创建提问							
1-4	To solve a problem by dynamic programming instead of recursions, the key approach is to store the results of computations for the subproblems so that we only have to compute each different subproblem once. Those solutions can be stored in an array or a hash table. (1)									
		F								
1-4	答案正确	(1分)	② 创建提问							
4.4		(4. /\)	O Alizabenza							
	答案正确	(1分)	© 创建提问		et of the					
1-5	When measuring the relevancy of the answer set of a search engine, the precision is low means that most of the relevant documents are not retrieved. (1%)									
	O T	• F								
	答案正确	(1分)	○ 创建提问							
1-6	In a B+ tre	ee, leaves ar	nd nonleaf nodes have some key values in common. (1分)							
	T	F								
1-6	答案正确	(1分)	♀ 创建提问							
2-1		t of accessing tements is FAI	the keys 4 and 8 in order in the splay tree given in the figure, which one of the LSE? (1%)	介書 申位	徐镜春 浙江大学					
		7)							
	Q	`	a							
	0 6	,								
	0									
	O B. 4 and	I 14 are sibling I 11 are sibling ne parent of 7 ne root								
2-1	 安正确	(1 公) 〇合	健福 间							

following statements is FALSE? (1分)	♀ 作者	徐镜春
0	单位	浙江大学
0 9		
Ó		
A. 7 and 14 are siblings		
O B. 4 and 11 are siblings		
C. 4 is the parent of 7D. 8 is the root		
2-1 签宏正确 (1 分) ○ 创建棉间		
-2 In a binomial queue with 150 nodes, how many nodes have depth 1 (the root has depth 0)? (1 分)		
● A. 4	♀ ↑ 单位	F者 杨欣舒 浙江之
O B. 8	4-17	מביוומ
C. 14D. Cannot be determined		
-2 答案错误 ① (0 分) ② 创建提问		
1-3 Given 4 characters (u, v, w, x) with some frequencies (f_u, f_v, f_w, f_x) in a text. If the corresponding Huffman codes are u : 00, v : 010, w : 011 and x : 1. Which of the following sets of frequencies is a possible one for	Q 1	F者 徐镜
(f_u, f_v, f_w, f_x) ? (1 $\%$)	单位	浙江
O A. 15, 23, 16, 45		
B. 30, 21, 12, 33		
○ C. 41, 12, 20, 32 ○ D. 55, 22, 18, 46		
0.000		
-3 <mark>答案正确 (1分) ♥ 创建提问 </mark>		
-4 A B+ tree of order 3 with 21 numbers has at most nodes of degree 3. (1分)	Q 1	F者 徐镜
○ A. 1 ○ B. 2	单位	浙江
O C. 3		
D. 4		
-4 <mark>答案正确</mark> (1分) Q 创建提问		
E. In proving the emertized bound of a Morge energtion in clean house, the notantial of a clean house is defined to	ha	
2-5 In proving the amortized bound of a Merge operation in skew heaps, the potential of a skew heap is defined to be	/h-tv	
the total number of right heavy nodes. Then we can prove that, in an N -node skew heap, the amortized cost for a	○ 作者	沈鑫 浙江大学
Merge operation is exactly		
Hint:		
<i>Hint:</i> Define the weight of a node, $w(x)$, to be the number of descendants of x (including x). A non-root node is said to		
Hint:		
Hint: Define the weight of a node, $w(x)$, to be the number of descendants of x (including x). A non-root node is said to be heavy if its weight is greater than half the weight of its parent. • Lemma 1: At most one child is heavy, of all children of any node.		
<i>Hint:</i> Define the weight of a node, $w(x)$, to be the number of descendants of x (including x). A non-root node is said to be <i>heavy</i> if its weight is greater than half the weight of its parent.		
Hint: Define the weight of a node, $w(x)$, to be the number of descendants of x (including x). A non-root node is said to be heavy if its weight is greater than half the weight of its parent. • Lemma 1: At most one child is heavy, of all children of any node. • Lemma 2: On any path from node x down to a descendant y , there are at most $\lfloor log_2 \frac{w(x)}{w(y)} \rfloor$ light nodes, excluding x .		
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Hint: Define the weight of a node, $w(x)$, to be the number of descendants of x (including x). A non-root node is said to be heavy if its weight is greater than half the weight of its parent. • Lemma 1: At most one child is heavy, of all children of any node. • Lemma 2: On any path from node x down to a descendant y , there are at most $\lfloor log_2 \frac{w(x)}{w(y)} \rfloor$ light nodes, excluding x . (1分) A. $\lfloor log_2 N \rfloor + 1$ B. $2 \lfloor log_2 N \rfloor + 1$ C. $3 \lfloor log_2 N \rfloor + 1$ D. $4 \lfloor log_2 N \rfloor + 1$ C. $3 \lfloor log_2 N \rfloor + 1$ D. $4 \lfloor log_2 N \rfloor + 1$	☆ 作者 单位	何钦铭
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Hint: Define the weight of a node, $w(x)$, to be the number of descendants of x (including x). A non-root node is said to be heavy if its weight is greater than half the weight of its parent. • Lemma 1: At most one child is heavy, of all children of any node. • Lemma 2: On any path from node x down to a descendant y , there are at most $\lfloor log_2 \frac{w(x)}{w(y)} \rfloor$ light nodes, excluding x . (1分) A. $\lfloor log_2 N \rfloor + 1$ B. $2 \lfloor log_2 N \rfloor + 1$ C. $3 \lfloor log_2 N \rfloor + 1$ D. $4 \lfloor log_2 N \rfloor + 1$ D. $4 \lfloor log_2 N \rfloor + 1$ C. $3 \lfloor log_2 N \rfloor + 1$ C. $3 \lfloor log_2 N \rfloor + 1$ D. $4 \lfloor log_2 N \rfloor + 1$ A. $2 - 5$ Sesting ① (0 分) ② 创建提问		
 Hint: Define the weight of a node, w(x), to be the number of descendants of x (including x). A non-root node is said to be heavy if its weight is greater than half the weight of its parent. Lemma 1: At most one child is heavy, of all children of any node. Lemma 2: On any path from node x down to a descendant y, there are at most ⌊log2 w(x) / w(y) ⌋ light nodes, excluding x. (1分) A. ⌊log2N⌋ + 1 B. 2 ⌊log2N⌋ + 1 C. 3 ⌊log2N⌋ + 1 D. 4 ⌊log2N⌋ + 1 D. 4 ⌊log2N⌋ + 1 D. 4 ⌊log2N⌋ + 1 D. 4 ⌊log2N⌋ + 1 2-5 答案错误 ③ (0 分) ② 创建提问 2-6 Given the distance set D={1,1,2,2,2,2,3,3,3,4,5,5,6,6,8} in a Turnpike Reconstruction problem, first it can be sure that x1=0 and x6=8. Which of the following possible solutions will be checked next? (1分) A. x2=1, x5=6 B. x2=2, x5=6 		



2-8 答案正确 (1分) ○ 创建提问

2-9 Merge the two leftist heaps in the following figure. Which one of the following statements is FALSE? (1

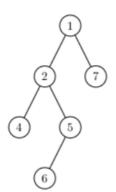


- A. the null path length of 5 is the same as that of 2
- O B. 1 is the root with 4 being its right child
- C. Along the left most path from top down, we have 1, 2, 5, and 7
- O. 8 is the left child of 3

2-9 答案正确 (1分) ♀ 创建提问

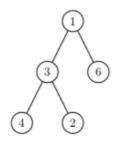
2-10 Which one of the following statements is TRUE? (1分)

A.



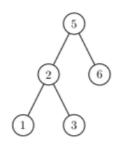
may be a leftist heap

B.



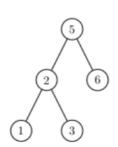
may be a leftist heap

○ C.



may be a leftist heap

O.



may be a skew heap

2-10 答案正确

(1分)

○ 创建提问

2-11 When solving a problem with input size N by divide and conquer, if at each step, the problem is divided into 9 sub-problems and each size of these sub-problems is N/3, and they are conquered in $O(N^2logN)$. Which one of the following is the closest to the overall time complexity? (1分)

○ 作者 单位

- $\ \ \, igotantering A.\ O(N^2log^2N)$
- \bigcirc B. $O(N^2logN)$
- \bigcirc C. $O(N^2)$
- \bigcirc D. $O(N^3logN)$
- 2-11 答案正确 (1分) ♀ 创建提问

5-1 The function BinQueue Merge is to merge two binomial queues H1 and H2, and return H1 as the resulting queue. BinQueue BinQueue Merge(BinQueue H1, BinQueue H2) { BinTree T1, T2, Carry = NULL; int i, j; H1->CurrentSize += H2-> CurrentSize; for (i=0, j=1; j<= H1->CurrentSize; i++, j*=2) { T1 = H1->TheTrees[i]; T2 = H2->TheTrees[i]; switch(4*!!Carry + 2*!!T2 + !!T1) { case 0: case 1: break; (1分); break; case 2: H1->TheTrees[i]=T2;H2->TheTrees[i]=NULL case 4: H1->TheTrees[i] = Carry; Carry = NULL; break; case 3: Carry = CombineTrees(T1, T2); H1->TheTrees[i] = H2->TheTrees[i] = NULL; break; case 5: Carry = CombineTrees(T1, Carry); H1->TheTrees[i] = NULL; break; case 6: Carry = CombineTrees(T2, Carry); H2->TheTrees[i] = NULL; break; case 7: H1->TheTrees[i] = Carry; Carry=CombineTrees(T1,T2) H2->TheTrees[i] = NULL; break; } /* end switch */ } /* end for-loop */ return H1; 5-1 <mark>答案正确</mark> (2分) ♀ 创建提问 5-2 The functions ISRBT is to check if a given binary search tree T is a red-black tree. Return true if T is, or false if not. ○ 作者 陈越 The red-black tree structure is defined as the following: 浙江大学 単位 400 ms typedef enum { red, black } colors; typedef struct RBNode *PtrToRBNode; struct RBNode{ int Data; 时间限制 内存限制 64 MB PtrToRBNode Left, Right, Parent; int BlackHeight; colors Color; typedef PtrToRBNode RBTree; Please fill in the blanks. bool IsRBT(RBTree T) int LeftBH, RightBH; if (!T) return true;
if (T->Color == black) T->BlackHeight = 1; else {
 if (T->Left && (T->Left->Color==red) (1分)) return false; if (T->Right && (T->Right->Color == red)) return false; }
if (!T->Left && !T->Right) return true;
(1分)) { if (IT->Left && IT->Right) return true; if (IsBRT(T->Left)&&IsBRT(T->Right) (1/2) if (IT->Left) LeftBH = T->Left->BlackHeight; else LeftBH = 0; if (IT->Right) RightBH = T->Right->BlackHeight; else RightBH = 0; if (LeftBH == RightBH) { (1分); T->BlackHeight+=LeftBH 截图(Alt + A) return true; else return false; else return false; 5-2 答案正确 (3分) ♀ 创建提问 5-2 The functions ISRBT is to check if a given binary search tree T is a red-black tree. Return true if T is, or false if not. 🔘 作者 陈哉 The red-black tree structure is defined as the following: 单位 浙江大学 typedef enum { red, black } colors; 时间限制 400 ms typedef struct RBNode *PtrToRBNode; 内存限制 64 MB struct RBNode{ int Data;
PtrTORBNode Left, Right, Parent;
int BlackHeight; colors Color; typedef PtrToRBNode RBTree; Please fill in the blanks bool IsRBT(RBTree T) int LeftBH, RightBH; if (!T) return true; if (T->Color == black) T->BlackHeight = 1; else { se { $if (\mbox{ T->Left &\& } (\mbox{T->Left->Color==red})$ if ($\mbox{T->Right && (T->Right->Color == red})) return false; }$ if (IsRBT(T->Left)&&IsRBT(T->Right) | (Tishol(T-Stert)Asishol(T-Skight) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | (17) | 截图(Alt + A) return true: else return false; else return false;

5-2 答案正确 (3分) ♀ 创建提问