

Training II

2021. Sept 17. Ye Cao.

Review of C++ pointer

&

Used to get the
memory location of
a variable

*

Used to get the
value of the variable
associated the pointer

Question:

```
int x = 10;  
int *p = &x;  
cout << p << endl;
```

1. Is printout equal 10?
2. If not, what's fix to print 10?

Review of pair

```
pair<int, int> p; // represents a point  
p.first; // x-coordinate  
p.second; // y-coordinate
```

Review of sort

<7 5 6 4 10 1 3>
↓
<1 3 4 5 6 7 10>

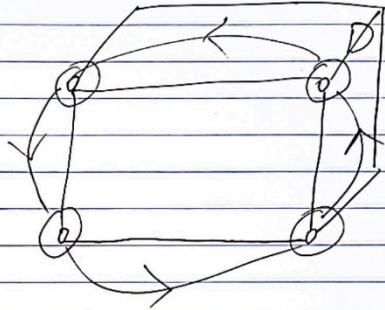
Review of binary search:

1. Define active region
2. Execute search
3. Update active region

Aquarium Tank Problem :

Problem Understanding :

Inputs : 1. Cartesian Coordinate of the convex polygon in counterclockwise order



2. Depth D

3. Water volume L

Computational Procedure :

1. How to process inputs ?

put into pairs contained by an array.

2. How to calculate height of water ?

GUESSING vs Math Formula

Linear guess - $O(N)$

$O(1)$

Binary Search - $O(\log(N))$

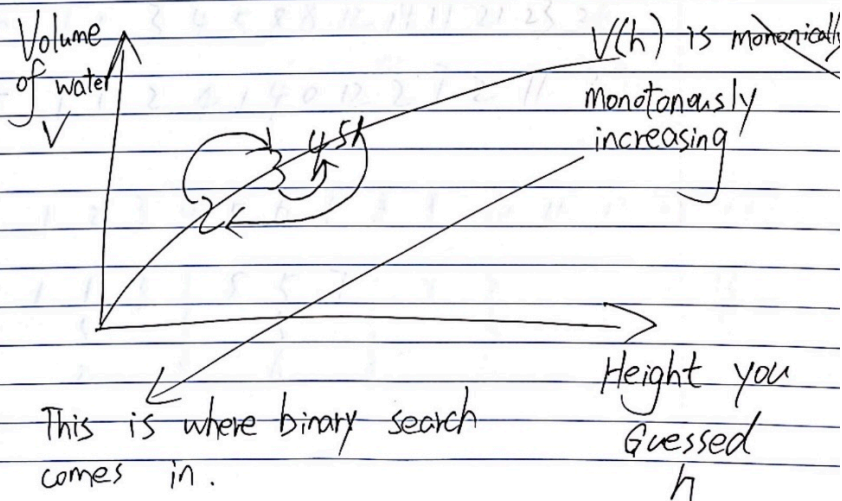
3. Look at input constraints. Guessing is feasible

Tips:

1. Guessing is a very generalized approach in CP.

Normally you don't know the answer, you make a guess and see if you can make a more EDUCATED guess each time.

2. Generally, guessing is very good if each time you make a guess, you are able to eliminate a range of answers from being right ones besides the single guess you made. This is the core idea of binary search.

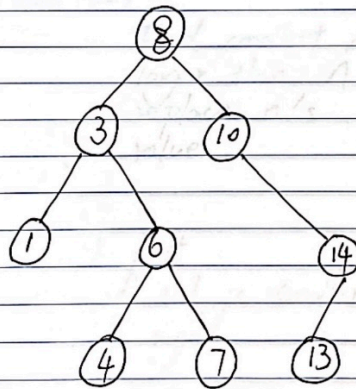


Intro to a tree

In computer science, a tree is an abstract data structure that simulates a hierarchical tree structure, with a root and subtree of children with parent node, represented as a set of linked nodes.

In graph theory, a tree is a connected acyclic graph.

Binary Search Tree



Insertion

Deletion

Retrieval

Insertion

1. Start with root
2. Go left if val is smaller
3. Otherwise go right
4. If reach a leaf node, attach val to leaf node

Deletion .

No children \rightarrow Remove node n

One children \rightarrow Remove node n and connect n 's parent to n 's children

Two children \rightarrow

Find smallest node m that's larger than n , remove m and replace n 's value with m 's value.

Retrieval $O(h)$

1. Start with root
2. Go left if val is smaller
3. Otherwise go right
4. If reach a leaf node, not found

Binary Indexed Tree / Fenwick Tree

Proposed by Peter M. Fenwick

Range Sum Query

Point Update

Consider this problem: There're n boxes and 2 queries

1. add marble to box i
2. sum marbles from box k to box r

How to come up with a data structure that works in time complexity $O(m \cdot \log(n))$?

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
f	1	0	2	1	1	3	0	4	2	5	2	2	3	1	0	2
c	1	1	3	4	5	8	8	12	14	19	21	23	26	27	27	29
T	1	1	2	4	1	4	0	12	2	7	2	11	3	4	0	29

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15

1 1 3 1 5 5 7 1 9 9 11 9 13 13 15

1 1 3 1 5 5 7 1 9 9 11 9 13 13 15

2 4 6 8 10 12 14

Binary Index Tree Implementation Details

1. Isolate last bit

$$\text{Last Bit } (10010) = 10$$

$$(11000) = 1000$$

$$(10001) = 1$$

$$\text{Last bit } (10000) = 10000$$

$a|b$
 ↑
 Last bit
 ↙ ↘
 Before last bit zeros

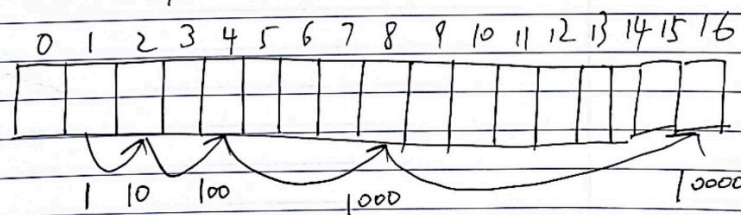
$$\begin{aligned}
 (a|b)^{-1} &= \bar{a} \bar{0} \bar{b}^{-1} + 1 \\
 &= \bar{a}^{-1} 0 (111 \dots 1) + 1 \\
 &= \bar{a}^{-1} b
 \end{aligned}$$

$$\therefore a|b \& (a|b)^{-1} = \text{Last bit}$$

2. Update

Start with current $\text{idx} + 1$ (1-based index)

Continue update to parent index



```
idx++; while (idx <= n)
```

```
idx += idx & (-idx);
```

Getsum

Start with current $\text{idx} + 1$ (1-based index)

Continue update to ~~parent~~ index

