### **Beyond Physical Memory: Policies**

- When there is not much memory free the OS has to start paging out pages to make room for actively used pages.
- The replacement policy contains the information to decide which pages to evict.

## Cache Management

- Main memory can be viewed as a cache for virtual memory pages in the sytem.
- The goal of a replacement policy is to minimize the number of cache misses for main memory.
  - Also to maximize cache hits.
- Knowing the number of cache hits and misses allows the calculation of the average memory access time (AMAT) for a program.
  - AMAT = (P\_hit T\_M) + (P\_miss T\_D)
    - T\_M is the cost of accessing memory.
    - T\_D is the cost of accessing disk.
    - P\_hit probability of a hit.
    - P\_miss probability of a miss.

# 22.2 Optimal Replacement Policy

- The optimal replacement policy leads to the fewest number ofm isses overall.
- Replacing a page that will be accessed furthest in the future is the optimal policy.
- · Cold start misses occur when the cache is empty in the beginning.

### 22.3 FIFO

- Some early systems used FIFO replacement.
- FIFO is noticably worse than the optimal replacement policy.

#### 22.4 Random

- Picks a random page to replace in memory.
- Randoms performance depends on how lucky it is.

### 22.5 LRU

- We can use history as a guide for replacement.
- Frequency can be used for historical information.
  - If a page has been accessed many times it should not be replaced.
- Recency can also be used.
  - The more recently a page has been used the more likely it is to be used again.
- Both policies are examples of the principle of locality.
  - Says the programs tend to access certain code sequences and data structures.

# 22.9 Dirty Pages

- When a page has been modified it has to be written to disk to evict it.
- · Referred to as dirty.
  - the modified bit is referred to as the dirty bit.
    - Set at any time a page is written.

#### 22.10 Other VM Policies.

- The OS has to decide when to bring a page into memory.
  - Called page selection.
  - Most OS's use demand paging.
    - bring a page into memory when it is accessed.
  - The OS could use prefetching and guess when the page will be accessed.
- Many OS's tend to collect a number of pending writes in memory and write them together at one time.
  - called clustering.

## 22.11 Thrashing

- Thrashing is when the system is constantly paging.
- Some OS systems would decide not to run a subset of processes if it was constantly paging.
  - Known as admission control
- Some versions of Linux run an out of memory killer.
  - when memory is oversubscribed a daemon chosses a memory intensive process and kills it.