Link layer, LANs: outline

- 5. I introduction, services
- 5.2 error detection, correction
- 5.3 multiple access protocols

5.4 LANs

- addressing, ARP
- Ethernet
- switches
- VLANS

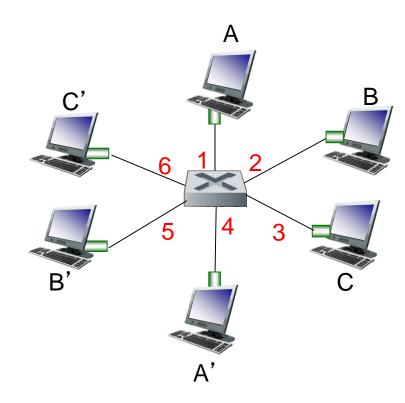
- 5.5 link virtualization: MPLS
- 5.6 data center networking
- 5.7 a day in the life of a web request

Ethernet switch

- link-layer device: takes an active role
 - store, forward Ethernet frames
 - examine incoming frame's MAC address, selectively forward frame to one-or-more outgoing links when frame is to be forwarded on segment, uses CSMA/CD to access segment
- transparent
 - hosts are unaware of presence of switches
- plug-and-play, self-learning
 - switches do not need to be configured

Switch: multiple simultaneous transmissions

- hosts have dedicated, direct connection to switch
- switches buffer packets
- Ethernet protocol used on each incoming link, but no collisions; full duplex (more later)
 - each link is its own collision domain
- * switching: A-to-A' and B-to-B' can transmit simultaneously, without collisions, either in the switch, or in the line.



switch with six interfaces (1,2,3,4,5,6)

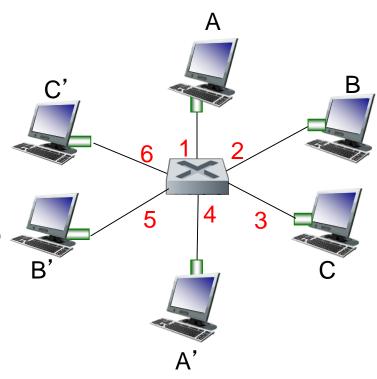
Switch forwarding table

Q: how does switch know A' reachable via interface 4, B' reachable via interface 5?

- A: each switch has a switch table, each entry:
 - (MAC address of host, interface to reach host, time stamp)
 - looks like a routing table!

Q: how are entries created, maintained in switch table?

something like a routing protocol?



switch with six interfaces (1,2,3,4,5,6)

Switch: self-learning

- switch learns which hosts can be reached through which interfaces
 - (only) when frame received, switch "learns" location of sender's incoming LAN segment
 - records sender/location pair in switch table

Dest: A'

Remember: no ACKs when sending Ethernet frames!

| MAC addr | interface | TTL |
|----------|-----------|-----|
| Α | 1 | 60 |
| | | |
| | | |

Switch table (initially empty)

Source: A

Switch: frame filtering/forwarding

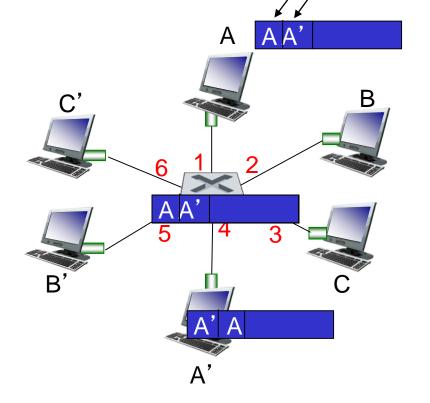
when frame received at switch:

```
I. record incoming link, MAC address of sending host
2. index switch table using MAC destination address
3. if entry found for destination, then
    if destination on segment from which frame arrived
        then drop frame
     else
       forward frame on interface indicated by entry
  else flood // forward on all interfaces except arriving one
```

Self-learning, forwarding: example

Source: A Dest: A'

- frame destination, A', locaton unknown: flood
- destination A location known: selectively send on just one link

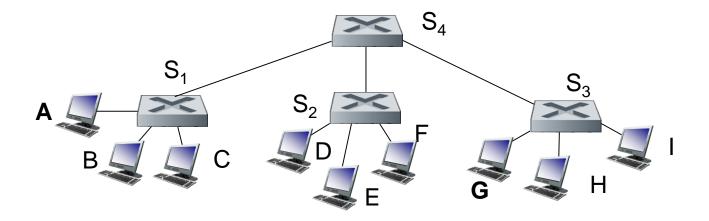


| MAC addr | interface | TTL |
|----------|-----------|-----|
| Α | 1 | 60 |
| Α' | 4 | 60 |
| | | |

switch table (initially empty)

Interconnecting switches

switches can be connected together

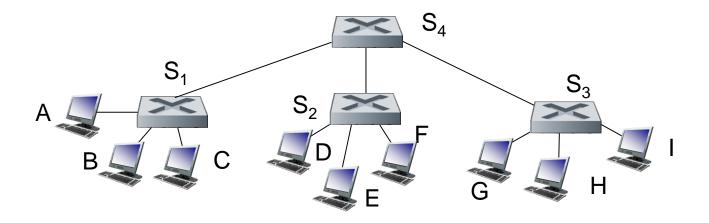


Q: sending from A to G - how does S_1 know to forward frame destined to G via S_4 and S_3 ?

A: self learning! (works exactly the same as in single-switch case!)

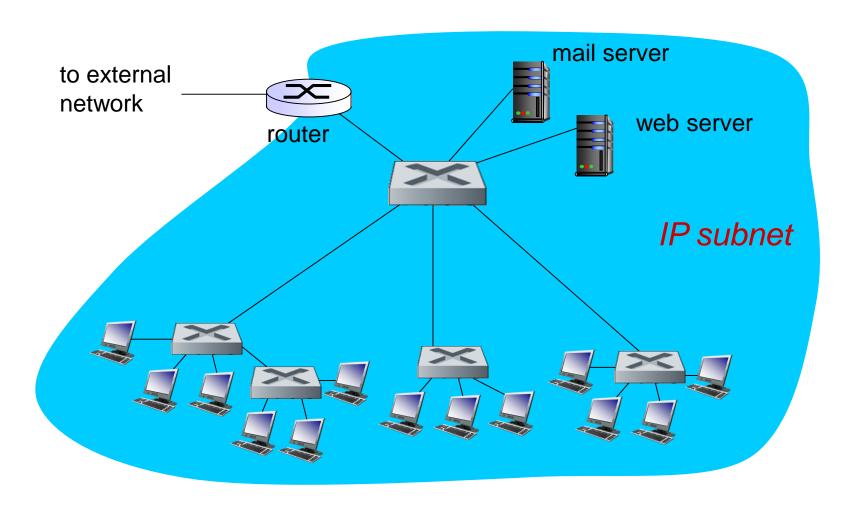
Self-learning multi-switch example

Suppose C sends frame to I, I responds to C



* Q: show switch tables and packet forwarding in S_1 , S_2 , S_3 , S_4

Institutional network



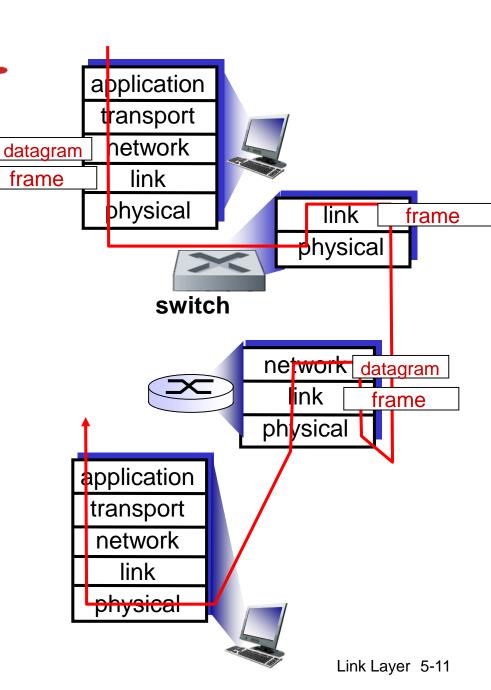
Switches vs. routers

both are store-and-forward:

- routers: network-layer devices (examine networklayer headers)
- switches: link-layer devices (examine link-layer headers)

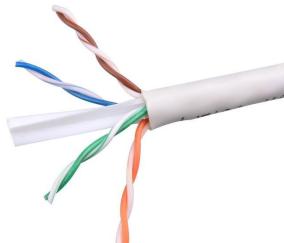
both have forwarding tables:

- routers: compute tables using routing algorithms, IP addresses
- switches: learn forwarding table using flooding, learning, MAC addresses



Collisions and MAC Protocols

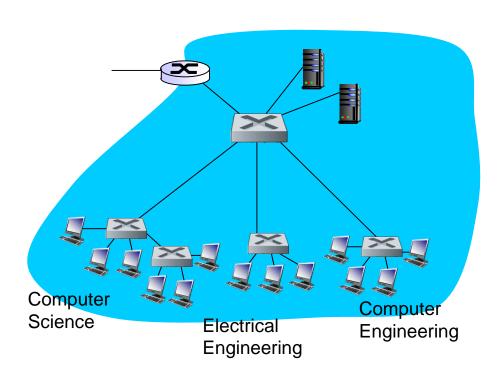
- Remember that MAC protocols involve both controlling multiple access needs and addressing
- Why use MAC protocols with a switch?
 - How many devices are communicating on each "link"?
 - Are collisions possible on a link?
 - Are they full duplex?
 - **2**
 - No, because 2 wires are used for each node out of 8 in the cable (so 4 total)
 - Yes, because collisions aren't possible!



Collisions and MAC Protocols

- Why use MAC protocols with a switch?
 - Not necessary!
 - But, the frame format hasn't changed, and the protocol works equally well with wired as well as wireless

VLANs: motivation



consider:

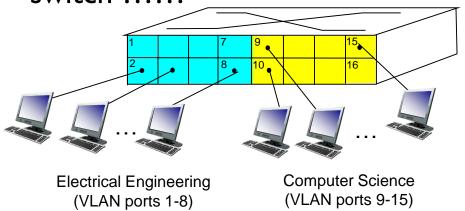
- CS user moves office to EE, but wants connect to CS switch?
- single broadcast domain:
 - all layer-2 broadcast traffic (ARP, DHCP, unknown location of destination MAC address) must cross entire LAN
 - security/privacy, efficiency issues

VLANs

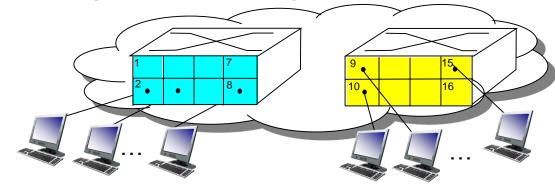
Virtual Local Area Network

switch(es) supporting VLAN capabilities can be configured to define multiple *virtual* LANS over single physical LAN infrastructure.

port-based VLAN: switch ports grouped (by switch management software) so that single physical switch



... operates as multiple virtual switches

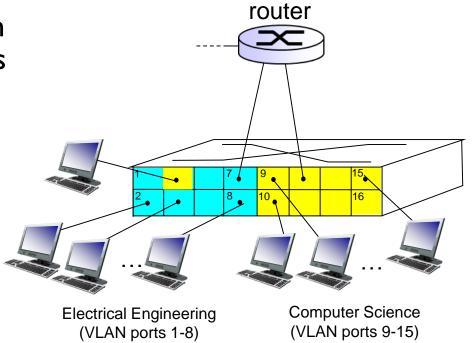


Electrical Engineering (VLAN ports 1-8)

Computer Science (VLAN ports 9-16)

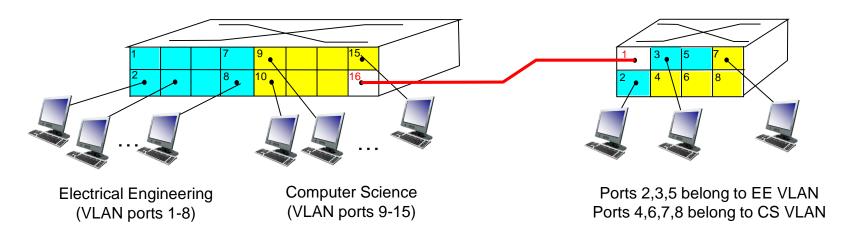
Port-based VLAN

- traffic isolation: frames to/from ports I-8 can only reach ports I-8
 - can also define VLAN based on MAC addresses of endpoints, rather than switch port #
- dynamic membership: ports can be dynamically assigned among VLANs



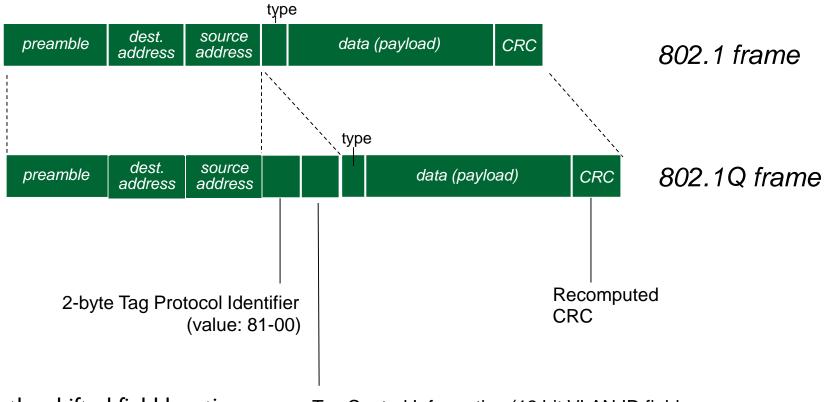
- forwarding between VLANS: done via routing (just as with separate switches)
 - in practice vendors sell combined switches+routers

VLANS spanning multiple switches



- trunk port: carries frames between VLANS defined over multiple physical switches
 - frames forwarded within VLAN between switches can't be vanilla 802.1 frames (must carry VLAN ID info)
 - 802. I q protocol adds/removes additional header fields for frames forwarded between trunk ports
 - Behavior of non-VLAN-aware switches when receiving VLAN tagged frames is undefined(!), and could include dropping, forwarding as-is, forwarding with tag removed (i.e. no longer "secure"), forwarding mangled packets, or crashing!

802. Iq VLAN frame format



Note the shifted field locations: this is what can cause mangled packets, etc.! How well does your switch or router handle a mangled packet?

Tag Control Information (12 bit VLAN ID field, 3 bit priority field like IP TOS)

Link layer, LANs: outline

- 5. I introduction, services
- 5.2 error detection, correction
- 5.3 multiple access protocols
- **5.4** LANs
 - addressing, ARP
 - Ethernet
 - switches
 - VLANS

- 5.5 link virtualization: MPLS
- 5.6 data center networking
- 5.7 a day in the life of a web request