

# Complexity

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## Measuring Complexity

1. Just Run a program and see
2. Count the number of instructions it would require
3. Come up with an approximate technique
  - We have to implement an algorithm
  - It might take different amounts for different inputs
  - It will probably be faster on faster computers

Experiment 1:

```
def count2sum(a):
    N = len(a)

    count = 0
    for i in range(N):
        for j in range(N):
            if a[i] + a[j] == 0:
                count += 1
    return count
```

Experiment 2:

```
def counteven(N):
    count = 0
    for i in range(1, N + 1):
        if i % 2 == 0:
            count += 1
    return count
```

Experiment 3:

```
def sum3(a):
    N = len(a)

    count = 0
    for i in range(N):
        for j in range(i + 1, N):
            for k in range(j + 1, N):
                if a[i] + a[j] + a[k] == 0:
                    count += 1
    return count
```

Experiment 4:

```
def counteven(N):  
    # Half the numbers will be even  
    # The other half are odd  
    count = N//2  
    return count
```

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## Classifications

- We can classify algorithms based on their performance
  1. Quadratic
    - Runtime proportional to  $N^2$
  2. Linear
    - Runtime proportional to  $N$
  3. Cubic
    - Runtime proportional to  $N^3$
  4. Linear
    - Runtime is proportional to 1

## Constant Factors

- Generally constant factors are not important, that is one quadratic algorithm could be twice as fast as another
  - They are both still quadratic
  - We are interested in the performance relative to  $N$

## Lower terms

- The runtime could be proportional to  $N^2 + 4N + 20$ 
    - We ignore the  $N$  and the 20 because there is  $N^2$  which is going to dominate as  $N$  gets large
    - Just take the highest power as it is what will dictate the runtime as  $N$  increases
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## Big O Notation

- Big O Notation is used to classify running times
  - $O(N^2)$  - Quadratic
  - $O(N^3)$  - Cubic
  - $O(N)$  - Linear
  - $O(1)$  - Constant

## Worst Case

- In theory, Big O Notation is for the worst case, so saying an algorithm is  $O(N^3)$  means it will not be worse than  $O(N^3)$ 
  - You could say that an  $O(1)$  algorithm is  $O(N^3)$
  - You could not say that an  $O(N^3)$  algorithm is  $O(1)$

## Algorithm Classes

O-Notation	Description	Speed	n = 1000	N = 1000000
$O(1)$	Constant	Instant	$10^{-9}$ secs	$10^{-9}$ secs
$O(n)$	Linear	Fast	$10^{-6}$ secs	$10^{-3}$ secs
$O(N^2)$	Quadratic	Slow for Large N	$10^{-3}$ secs	$10^3$ secs
$O(N^3)$	Cubic	Slow for medium N	1 second	$10^9$ secs

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## Examining Code

- It's easier if to not have to run an experiment everytime you run a program
- Generally, everytime we go over an input, the runtime is multiplied by N.
  - If we examine the input once, the algorithm will be  $O(N)$
  - If for each item in the input, we check every other item in the input, then the algorithm is  $O(N^2)$

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## Logarithmic and Exponential

- There are two more important classes of algorithm
  1. Logarithmic
    - The input is split in half every time (e.g. Binary Search)
  2. Exponential
    - The runtime doubles each time N increases by 1

## Logarithmic

- Very Fast
  - Sample values:

N	Value
1000	10
1000000	20
1000000000	30

The log of a number rises very slowly, Doing a binary search on an array of 1000000 elements is only 20 times longer than a binary search of an array of 1 element.

### **N\*Log(N)**

- $N \log(N)$ ,  $O(N \log N)$ ,  $N \log(N)$ 
  - Linearithmic
  - \* very close to linear

### **Exponential**

- $2^N$ 
  - The runtime doubles each time N is increased by 1

N	Value
10	1000
20	1000000
30	1000000000

- When N is 30, the operation takes the order of a second. When N is 40 it takes 1000 seconds (15 minutes)
- N = 50 requires 1000 \* 15 minutes ( approx 250 hours or 10 days )

O-Notation	Description	Speed	~Time for n=1000	~time for n=1000000
O(1)	Constant	Instant	$10^{-9}$ secs	$10^{-9}$ secs
O(logN)	Logarithmic	Very Fast	$10^{-8}$ secs	$10^{-8}$ secs
O(N)	Linear	Fast	$10^{-6}$ secs	$10^{-3}$ secs
O(NlogN)	Linearithmic	Fast	$10^{-5}$ secs	$10^{-3}$ secs
O(N <sup>2</sup> )	Quadratic	Slow for Large N	$10^{-3}$ secs	$10^3$ secs
O(N <sup>3</sup> )	Cubic	Slow for Medium N	1 second	$10^9$ secs
O(2 <sup>N</sup> )	Exponential	Extremely Slow	$10^{301}$ secs	$\rightarrow \infty$

### **Types of Data**

- The runtime complexity can depend on the actual data
  - Not just the size on the dataset (N)
  - In this case we can talk about best case
  - Or worst case
  - Or average case
- Worst case is usually considered