

The Puzzles of MrHeer

MrHeer

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1 Throwing eggs from a building

Question

Suppose that you have an N -story building and plenty of eggs. Suppose also that an egg is broken if it is thrown off floor F or higher, and unhurt otherwise. First, devise a strategy to determine the value of F such that the number of broken eggs is $\sim \lg N$ when using $\sim \lg N$ throws, then find a way to reduce the cost to $\sim 2 \lg F$.

Answer

$\sim \lg N$: start at the middle, always cut search space in half \rightarrow binary search.
 $\sim 2 \lg F$: start at 1, next 2, 4, 8 (i.e., 2^i), once the egg breaks after ($\sim \lg F$ steps) do binary search in the smaller search space (range $< F$ and hence number of searches $< \sim \lg F$) \rightarrow exponential search.

$$2^{\lceil \lg F \rceil - 1} < F \leq 2^{\lceil \lg F \rceil}$$

$$range = 2^{\lceil \lg F \rceil} - 2^{\lceil \lg F \rceil - 1} = 2^{\lceil \lg F \rceil - 1} < 2^{\lg F} = F$$

$$range < F$$

2 Throwing two eggs from a building

Question

Consider the previous question, but now suppose you only have two eggs, and your cost model is the number of throws. Devise a strategy to determine F such that the number of throws is at most $2\sqrt{N}$, then find a way to reduce the cost to $\sim c\sqrt{F}$. This is analogous to a situation where search hits (egg intact) are much cheaper than misses (egg broken).

Answer

Let us make our first attempt on k 'th floor. If it breaks, we try remaining $(k - 1)$ floors one by one. So in worst case, we make k trials. If it doesn't break, we jump $(k - 1)$ floors (Because we have already made one attempt and we don't want to go beyond k attempts. Therefore $(k - 1)$ attempts are available), Next floor we try is floor $k + (k - 1)$ Similarly, if this drop does not break, next need to jump up to floor $k + (k - 1) + (k - 2)$, then $k + (k - 1) + (k - 2) + (k - 3)$ and so on. Since the last floor to be tried is F 'th floor, sum of series should be F for optimal value of k .

$$k + (k - 1) + (k - 2) + (k - 3) + \cdots + 1 \geq F$$

$$\frac{k(k + 1)}{2} \geq F$$

$$k \geq \frac{\sqrt{8F + 1} - 1}{2}$$

$$k_{min} = \lceil \frac{\sqrt{8F + 1} - 1}{2} \rceil \sim \sqrt{2F}$$

Official solution

Solution to Part 1: To achieve $2\sqrt{N}$, drop eggs at floors \sqrt{N} , $2*\sqrt{N}$, $3*\sqrt{N}$, ..., $\sqrt{N}\sqrt{N}$. (For simplicity, we assume here that \sqrt{N} is an integer.) Let assume that the egg broke at level $k\sqrt{N}$. With the second egg you should then perform a linear search in the interval $(k - 1)\sqrt{N}$ to $k\sqrt{N}$. In total you will be able to find the floor F in at most $2\sqrt{N}$ trials.

Hint for Part 2: $1 + 2 + 3 + \cdots + k \sim \frac{1}{2}k^2 \geq F$.

3 3-collinearity

Question

Suppose that you have an algorithm that takes as input N distinct points in the plane and can return the number of triples that fall on the same line. Show that you can use this algorithm to solve the 3-sum problem. Strong hint: Use algebra to show that (a, a^3) , (b, b^3) , and (c, c^3) are collinear if and only if $a + b + c = 0$.

Answer

Proof. (a, a^3) , (b, b^3) , and (c, c^3) are collinear if and only if $a + b + c = 0$: We use a formulation of collinearity which equates gradients (assuming our points are distinct)

$$\frac{y_2 - y_1}{x_2 - x_1} = \frac{y_3 - y_1}{x_3 - x_1}$$

This becomes

$$\frac{b^3 - a^3}{b - a} = \frac{c^3 - a^3}{c - a}$$

which leaves us with

$$b^2 + ab + a^2 = c^2 + ac + a^2$$

so that

$$b^2 - c^2 = a(c - b)$$

$c \neq b$ so we have $a = -(b + c)$

$$a + b + c = 0$$

□

4 Queue with three stacks [Unsolved]

Question

Implement a queue with three stacks so that each queue operation takes a constant (worst-case) number of stack operations.

Answer

5 Queue with two stacks

Question

Implement a queue with two stacks so that each queue operation takes a constant amortized number of stack operations. Hint: If you push elements onto a stack and then pop them all, they appear in reverse order. If you repeat this process, they're now back in order.

Answer

Transfer

Result: stack1 to stack2 and reverse order

```
while stack1 is not empty do  
    | Item tmp = stack1.pop();  
    | stack2.push(tmp);  
end
```

Enqueue

Input: element item

Result: enqueue

```
stack1.push(item);
```

Dequeue

Result: dequeue

```
if queue is empty then  
    | throw NoSuchElementException;  
end  
if stack2 is empty then  
    | transfer();  
end  
return stack2.pop();
```

Code

```
public class QueueWithTwoStacks<Item> {
    private Stack<Item> stack1;
    private Stack<Item> stack2;

    public QueueWithTwoStacks() {
        stack1 = new Stack<>();
        stack2 = new Stack<>();
    }

    private void transfer() {
        while (!stack1.isEmpty()) {
            Item tmp = stack1.pop();
            stack2.push(tmp);
        }
    }

    public void enqueue(Item item) {
        stack1.push(item);
    }

    public Item dequeue() {
        if (isEmpty()) throw new NoSuchElementException("Queue underflow");
        if (stack2.isEmpty()) transfer();
        return stack2.pop();
    }
}
```

Analysis

For N items enqueue and dequeue.

enqueue: N times push operations.

dequeue:

transfer: N times pop and N times push operations.

dequeue: N times pop operations.

total: $4N$ times stack operations.

So, the amortized number of stack operations is $\frac{4N}{N} = 4$

6 Deque with a stack and a steque [Unsolved]

Question

Implement a deque with a stack and a steque so that each deque operation takes a constant amortized number of stack and steque operations.

Answer

7 Deque with three stacks [Unsolved]

Question

Implement a deque with three stacks so that each deque operation takes a constant amortized number of stack operations.

Answer

8 Hot or cold

Question

Your goal is to guess a secret integer between 1 and N . You repeatedly guess integers between 1 and N . After each guess you learn if your guess equals the secret integer (and the game stops). Otherwise, you learn if the guess is hotter (closer to) or colder (farther from) the secret number than your previous guess. Design an algorithm that finds the secret number in at most $\sim 2 \lg N$ guesses. Then design an algorithm that finds the secret number in at most $\sim 1 \lg N$ guesses.

Hint: use binary search for the first part. For the second part, first design an algorithm that solves the problem in $\sim 1 \lg N$ guesses assuming you are permitted to guess integers in the range $-N$ to $2N$.

Answer

$\sim 2 \lg N$: Suppose the secret between $[a, b]$, $mid = \frac{a+b}{2}$. First guess a , then guess b , if b hotter a , guess $[mid, b]$, else guess $[a, mid]$...

$\sim 1 \lg N$: Suppose you know that your secret integer is in $[a, b]$, and that your last guess is c . You want to divide your interval by two, and to know whether your secret integer lies in between $[a, m]$ or $[m, b]$, with $m = \frac{a+b}{2}$. The trick is to guess d , such that $\frac{c+d}{2} = \frac{a+b}{2}$.

Proof. guess integers in the range $-N$ to $2N$:

Suppose you know that your secret integer is in $[a_i, b_i]$, $b_i - a_i = d_i$, your last guess is g_{i-1} , this guess is g_i , and $\frac{g_i + g_{i-1}}{2} = \frac{a_i + b_i}{2} = m_i$.

$$g_i + g_{i-1} = 2m_i \tag{1}$$

$$g_{i-1} + g_{i-2} = 2m_{i-1} \tag{2}$$

$$(1) - (2)$$

$$g_i - g_{i-2} = 2(m_i - m_{i-1})$$

We know that $|m_i - m_{i-1}| = \frac{d_i}{2}$

$$|g_i - g_{i-2}| = d_i$$

Because divide interval by two everytime

$$d_i = \frac{1}{2}d_{i-1}$$

We want to know the g_{max} , so i is even and $g_i > g_{i-2}$

$$g_2 - g_0 = \frac{1}{2}d_1 \quad (1)$$

$$g_4 - g_2 = \frac{1}{8}d_1 \quad (2)$$

...

$$g_i - g_{i-2} = \frac{1}{2^{i-1}}d_1 \quad (\frac{i}{2})$$

$$(1) + (2) + \dots + (\frac{i}{2})$$

$$\begin{aligned} g_i - g_0 &= (\frac{1}{2} + \frac{1}{8} + \dots + \frac{1}{2^{i-1}})d_1 \\ &= \frac{2}{3}(1 - \frac{1}{2^i})d_1 \\ &= \frac{2}{3}(1 - \frac{1}{2^i})(N - 1) \end{aligned}$$

$$g_i = \frac{2}{3}(1 - \frac{1}{2^i})(N - 1) + g_0$$

When $g_0 = N$

$$\begin{aligned} g_{max} &= \lim_{i \rightarrow \infty} \frac{2}{3}(1 - \frac{1}{2^i})(N - 1) + N \\ &= \frac{2}{3}(N - 1) + N \end{aligned}$$

$$g_{max} < 2N$$

The same reason

$$g_i = g_0 - \frac{2}{3}(1 - \frac{1}{2^i})(N - 1)$$

When $g_0 = 1$

$$\begin{aligned} g_{max} &= \lim_{i \rightarrow \infty} 1 - \frac{2}{3}(1 - \frac{1}{2^i})(N - 1) \\ &= 1 - \frac{2}{3}(N - 1) \end{aligned}$$

$$g_{min} > -N$$

So, guess integers in the range $-N$ to $2N$.

Of course, we can also use mathematical induction to prove.

Suppose $a_i - b_i < g_{i-1} < 2b_i$, because $b_{i-1} - a_{i-1} > b_i - a_i$, and $b_{i-1} \geq b_i$

$$a_{i-1} - b_{i-1} < a_i - b_i < g_{i-1} < 2b_i \leq 2b_{i-1}$$

$$\implies a_{i-1} - b_{i-1} < g_{i-1} < 2b_{i-1}$$

Because $\frac{g_i + g_{i-1}}{2} = \frac{a_i + b_i}{2}$, so

$$a_i - b_i < a_i + b_i - g_i < 2b_i$$

$$\implies a_i - b_i < g_i < 2b_i$$

So, $a_i - b_i < g_i < 2b_i$, and we know $1 \leq a_i \leq b_i \leq N$

$$1 - N < g_i < 2N$$

So, guess integers in the range $-N$ to $2N$.

□