# The Draft of MrHeer

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### 1 Throwing eggs from a building

#### Question

Suppose that you have an N-story building and plenty of eggs. Suppose also that an egg is broken if it is thrown off floor F or higher, and unhurt otherwise. First, devise a strategy to determine the value of F such that the number of broken eggs is  $\sim \lg N$  when using  $\sim \lg N$  throws, then find a way to reduce the cost to  $\sim 2 \lg F$ .

#### Answer

 $\sim \lg N$ : start at the middle, always cut search space in half  $\to$  binary search.  $\sim 2 \lg F$ : start at 1, next 2, 4, 8 (i.e.,  $2^i$ ), once the egg breaks after ( $\sim \lg F$  steps) do binary search in the smaller search space (range < F and hence number of searches  $< \sim \lg F$ )  $\to$  exponential search.

$$\begin{split} 2^{\lceil \lg F \rceil - 1} < F \leqslant 2^{\lceil \lg F \rceil} \\ range &= 2^{\lceil \lg F \rceil} - 2^{\lceil \lg F \rceil - 1} = 2^{\lceil \lg F \rceil - 1} < 2^{\lg F} = F \\ range &< F \end{split}$$

### 2 Throwing two eggs from a building

#### Question

Consider the previous question, but now suppose you only have two eggs, and your cost model is the number of throws. Devise a strategy to determine F such that the number of throws is at most  $2\sqrt{N}$ , then find a way to reduce the cost to  $\sim c\sqrt{F}$ . This is analogous to a situation where search hits (egg intact) are much cheaper than misses (egg broken).

#### Answer

Let us make our first attempt on k'th floor. If it breaks, we try remaining (k-1) floors one by one. So in worst case, we make k trials. If it doesn't break, we jump (k-1) floors (Because we have already made one attempt and we don't want to go beyond k attempts. Therefore (k-1) attempts are available), Next floor we try is floor k + (k-1) Similarly, if this drop does not break, next need to jump up to floor k + (k-1) + (k-2), then k + (k-1) + (k-2) + (k-3) and so on. Since the last floor to be tried is k'th floor, sum of series should be k for optimal value of k.

$$k + (k-1) + (k-2) + (k-3) + \dots + 1 \ge F$$

$$\frac{k(k+1)}{2} \ge F$$

$$k \ge \frac{\sqrt{8F+1} - 1}{2}$$

$$k_{min} = \lceil \frac{\sqrt{8F+1} - 1}{2} \rceil \sim \sqrt{2F}$$

#### Official solution

Solution to Part 1: To achieve  $2\sqrt{N}$ , drop eggs at floors  $\sqrt{N}$ ,  $2*\sqrt{N}$ ,  $3*\sqrt{N}$ , ...,  $\sqrt{N}\sqrt{N}$ . (For simplicity, we assume here that  $\sqrt{N}$  is an integer.) Let assume that the egg broke at level  $k\sqrt{N}$ . With the second egg you should then perform a linear search in the interval  $(k-1)\sqrt{N}$  to  $k\sqrt{N}$ . In total you will be able to find the floor F in at most  $2\sqrt{N}$  trials.

Hint for Part 2: 
$$1 + 2 + 3 + \dots + k \sim \frac{1}{2}k^2 \ge F$$
.

### 3 3-collinearity

#### Question

Suppose that you have an algorithm that takes as input N distinct points in the plane and can return the number of triples that fall on the same line. Show that you can use this algorithm to solve the 3-sum problem. Strong hint: Use algebra to show that  $(a, a^3)$ ,  $(b, b^3)$ , and  $(c, c^3)$  are collinear if and only if a + b + c = 0.

#### Answer

*Proof.*  $(a, a^3)$ ,  $(b, b^3)$ , and  $(c, c^3)$  are collinear if and only if a + b + c = 0: We use a formulation of collinearity which equates gradients (assuming our points are distinct)

$$\frac{y_2 - y_1}{x_2 - x_1} = \frac{y_3 - y_1}{x_3 - x_1}$$

This becomes

$$\frac{b^3 - a^3}{b - a} = \frac{c^3 - a^3}{c - a}$$

which leaves us with

$$b^2 + ab + a^2 = c^2 + ac + a^2$$

so that

$$b^2 - c^2 = a(c - b)$$

 $c \neq b$  so we have a = -(b+c)

$$a + b + c = 0$$

# 4 Queue with three stacks

### Question

Implement a queue with three stacks so that each queue operation takes a constant (worst-case) number of stack operations.

### Answer

## 5 Queue with two stacks

### Question

#### Answer

Implement a queue with two stacks so that each queue operation takes a constant amortized number of stack operations. Hint: If you push elements onto a stack and then pop them all, they appear in reverse order. If you repeat this process, they're now back in order.