CFG3: Push-down Automata

Context-Free Grammars

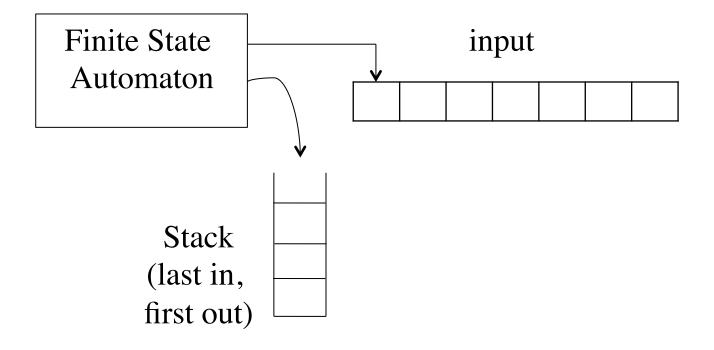
CMPT 379: Compilers

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anoopsarkar.github.io/compilers-class

Context-free languages and Pushdown Automata

- Recall that for each regular language there was an equivalent finite-state automaton
- The FSA was used as a recognizer of the regular language
- For each context-free language there is also an automaton that recognizes it: called a pushdown automaton (pda)



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Context-free languages and Pushdown Automata

- Similar to FSAs there are non-deterministic pda and deterministic pda
- Unlike in the case of FSAs we cannot always convert a npda to a dpda
- Our goal in compiler design will be to choose grammars carefully so that we can always provide a dpda for it
- Similar to the FSA case, a DFA construction provides us with the algorithm for lexical analysis,
- In this case the construction of a dpda will provide us with the algorithm for parsing (take in strings and provide the parse tree)
- We will study later how to convert a given CFG into a ¹⁶⁻⁰⁶ parser by first converting into a PDA

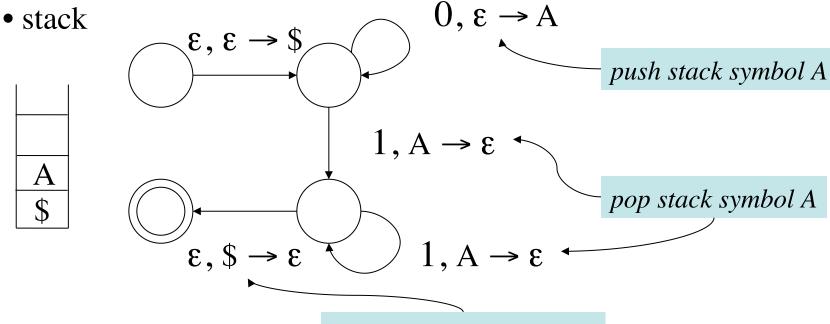
Pushdown Automata

- PDA has
 - an alphabet (terminals),
 - stack symbols (like non-terminals and terminals),
 - a finite-state automaton,

e.g. PDA for language

$$L = \{ 0^n 1^n : n >= 1 \}$$

→ implies a push/pop of stack symbol(s)



Non-CF Languages

$$L_1 = \{wcw \mid w \in (a|b)*\}$$
 $L_2 = \{a^nb^mc^nd^m \mid n \ge 1, m \ge 1\}$
 $L_3 = \{a^nb^nc^n \mid n \ge 0\}$

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CF Languages

$$L_4 = \{wcw^R \mid w \in (a|b)*\}$$
 $S \to aSa \mid bSb \mid c$
 $L_5 = \{a^nb^mc^md^n \mid n \ge 1, m \ge 1\}$
 $S \to aSd \mid aAd$
 $A \to bAc \mid bc$

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Summary

- CFGs can be used describe PL
- Derivations correspond to parse trees
- Parse trees represent structure of programs
- Ambiguous CFGs exist
- Some forms of ambiguity can be fixed by changing the grammar
- CF languages can be recognized using Pushdown Automata

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