

Lexical Analysis

CMPT 379: Compilers

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anoopsarkar.github.io/compilers-class

Building a Lexical Analyzer

- Token \Rightarrow Pattern
- Pattern \Rightarrow Regular Expression
- Regular Expression \Rightarrow NFA
- NFA \Rightarrow DFA
- DFA \Rightarrow Table-driven implementation of DFA

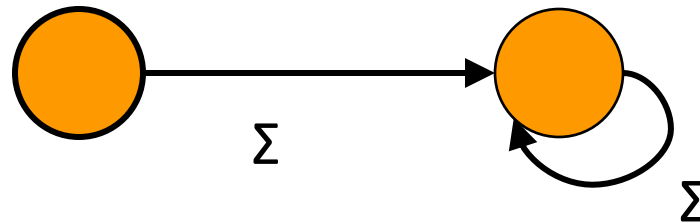
Thompson's construction

- Converts regexps to equivalent NFA
- Six simple rules
 - Empty language
 - Symbols (Σ)
 - Empty String (ϵ)
 - Alternation (r_1 or r_2)
 - Concatenation (r_1 followed by r_2)
 - Repetition (r_1^*)

Used by Ken Thompson for pattern-based search in text editor QED (1968)

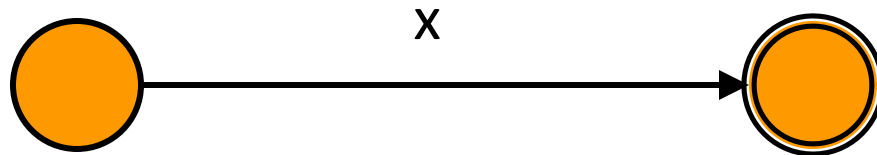
Thompson Rule 0

- For the empty language ϕ (optionally include a *sinkhole* state)



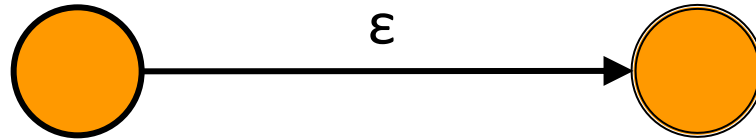
Thompson Rule 1

- For each symbol x of the alphabet, there is a NFA that accepts it



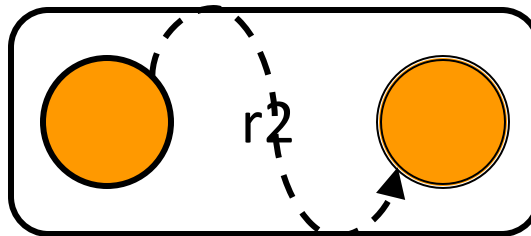
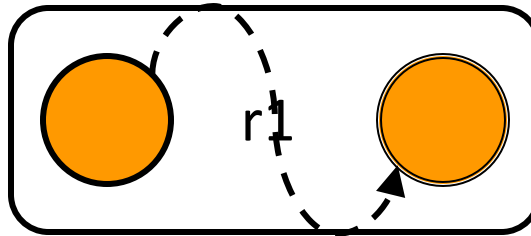
Thompson Rule 2

- There is an NFA that accepts only ϵ



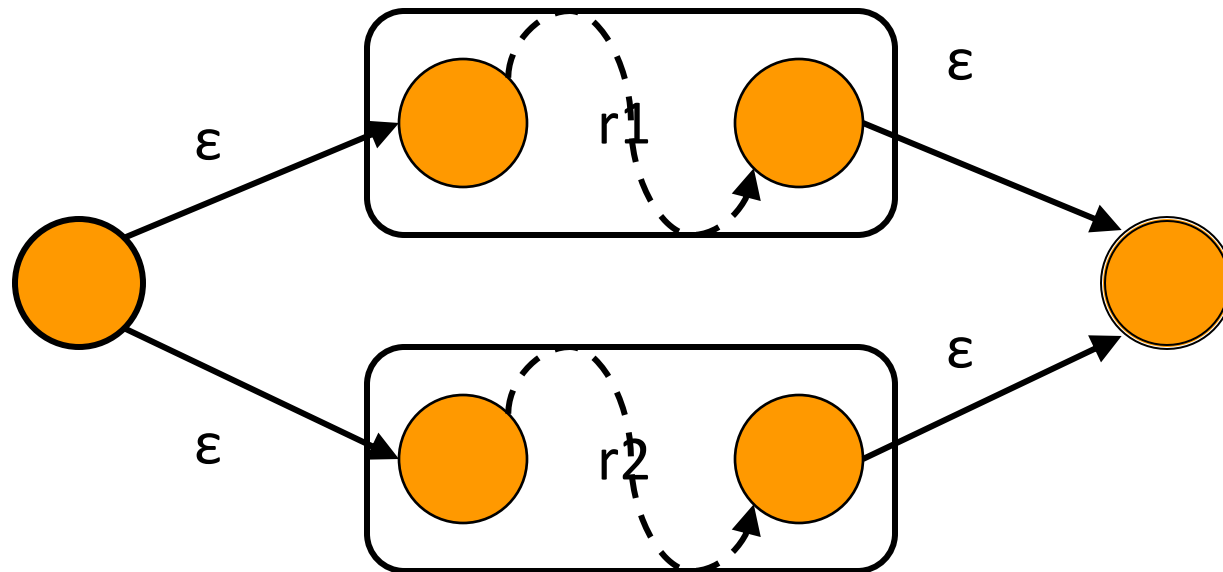
Thompson Rule 3

- Given two NFAs for r_1 , r_2 , there is a NFA that accepts $r_1 | r_2$



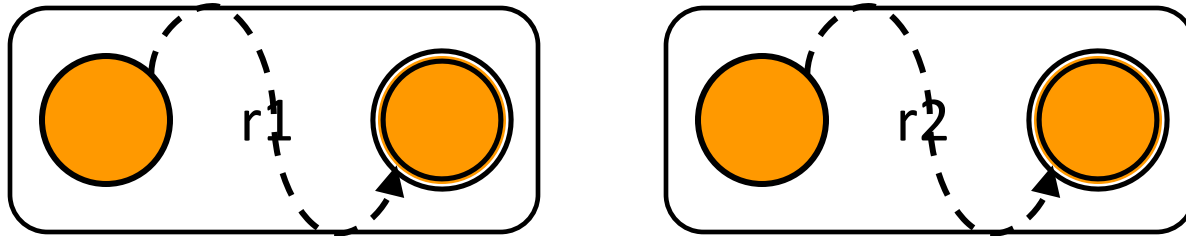
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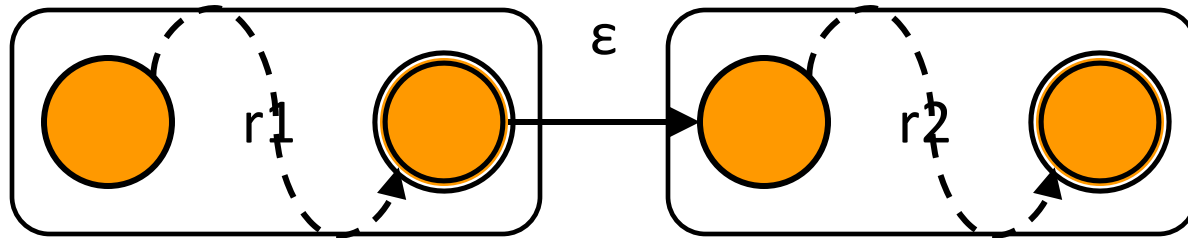
Thompson Rule 4

- Given two NFAs for r_1 , r_2 , there is a NFA that accepts r_1r_2



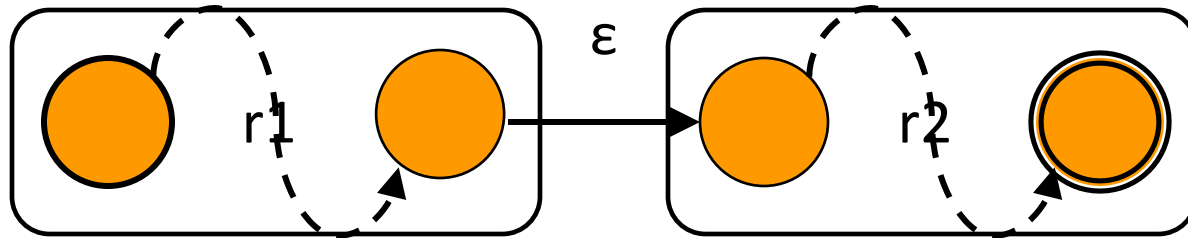
Thompson Rule 4

- Given two NFAs for r_1 , r_2 , there is a NFA that accepts r_1r_2



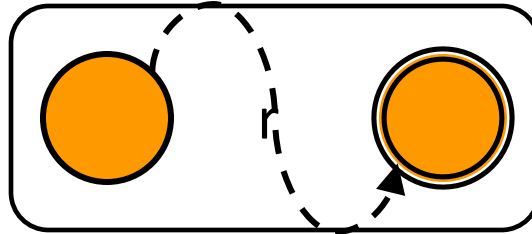
Thompson Rule 4

- Given two NFAs for r_1 , r_2 , there is a NFA that accepts r_1r_2



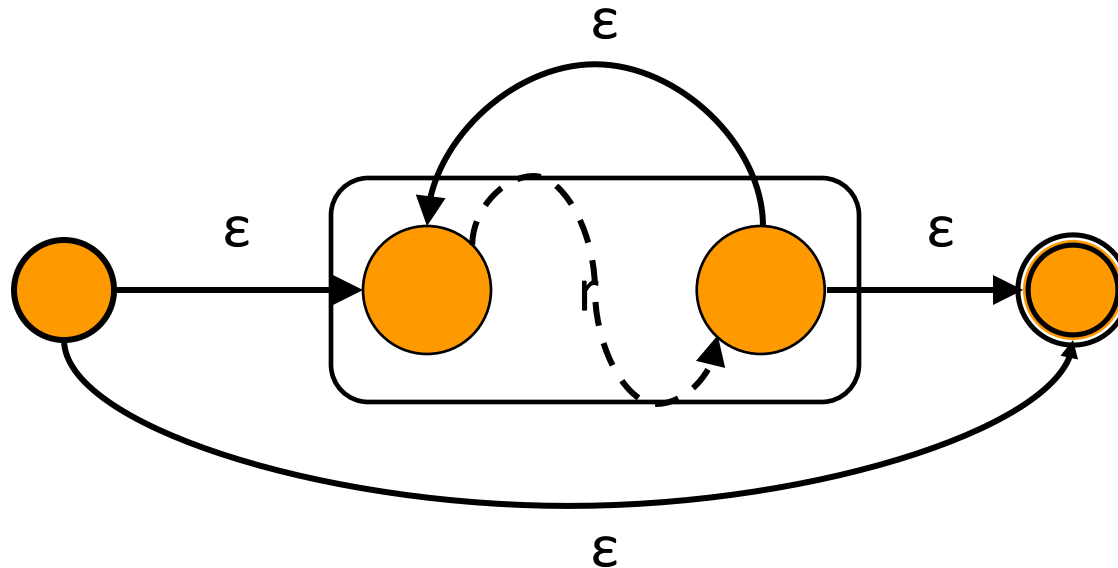
Thompson Rule 5

- Given a NFA for r , there is an NFA that accepts r^*



Thompson Rule 5

- Given a NFA for r , there is an NFA that accepts r^*



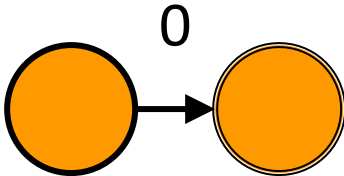
Example

- Set of all binary strings that are divisible by four (include 0 in this set)
- Defined by the regexp: $((0|1)^*00) | 0$
- Apply Thompson's Rules to create an NFA

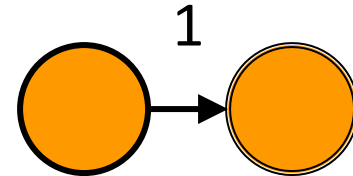
Basic Blocks 0 and 1

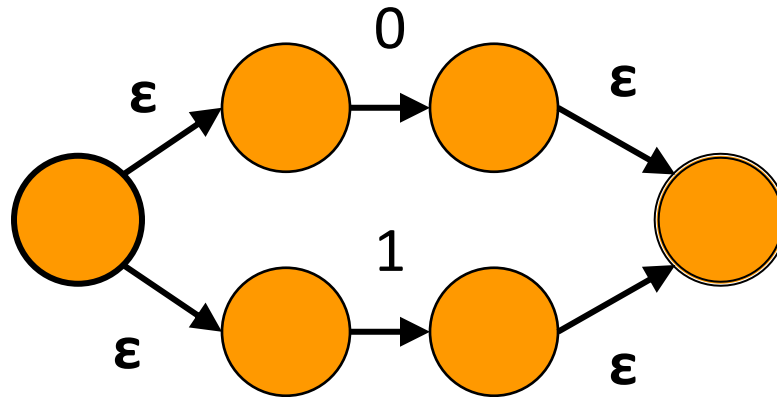
$((0|1)^*00) | 0$

NFA for 0



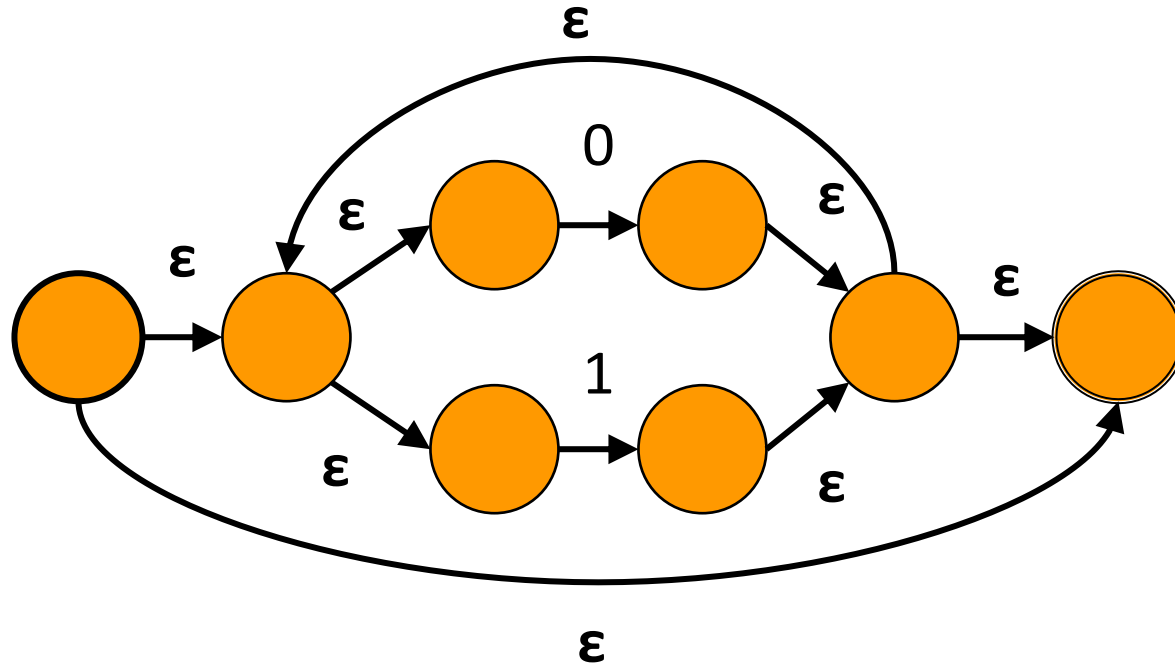
NFA for 1





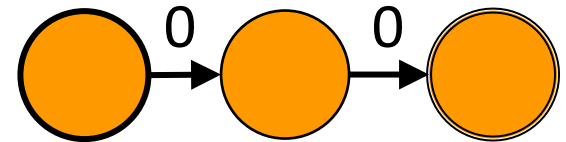
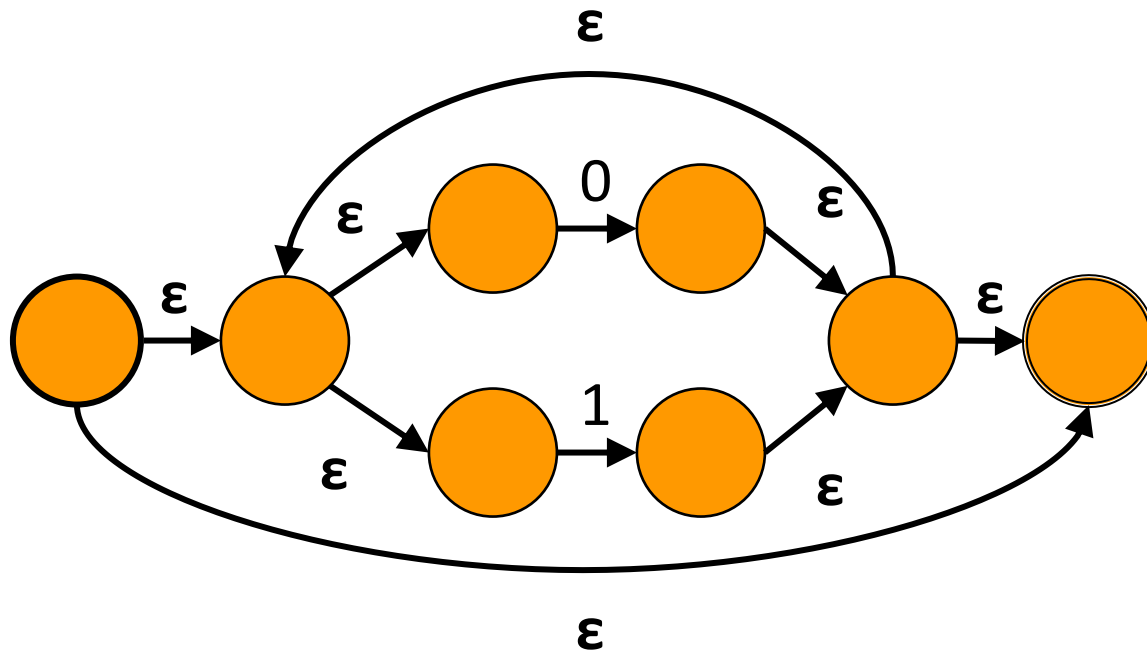
$0|1$

$((0|1)^*00) | 0$



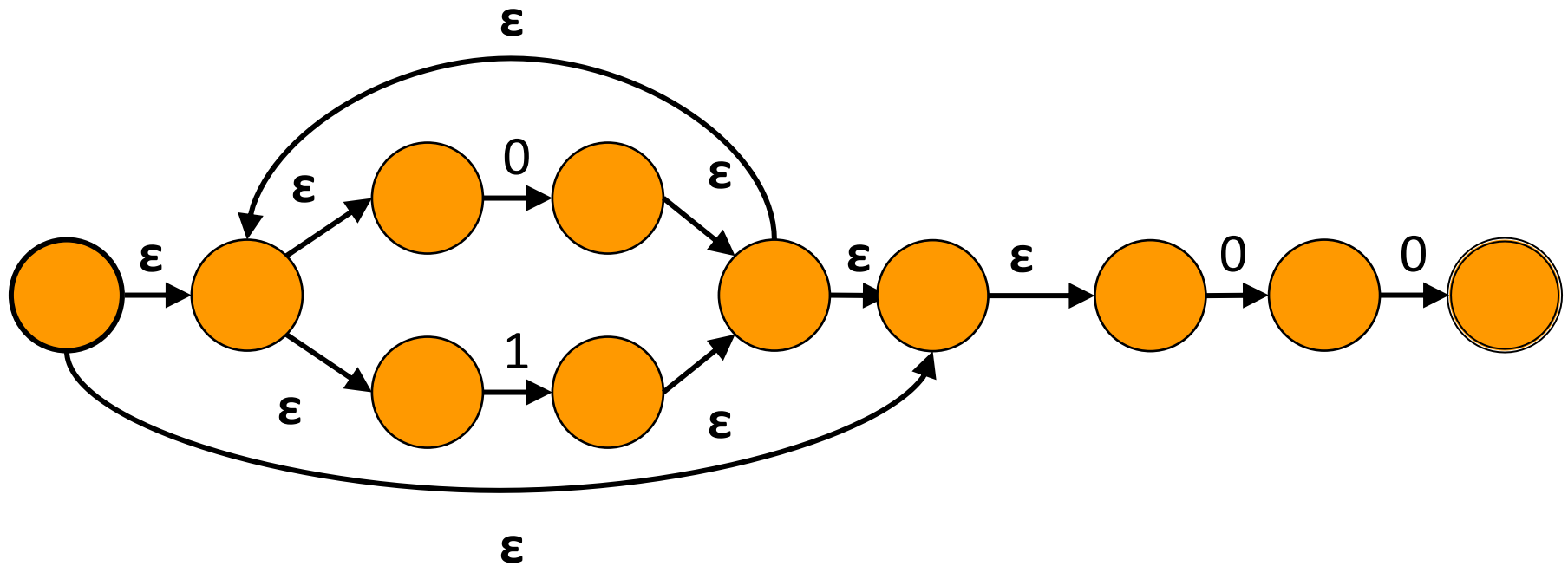
$(0|1)^*$

$((0|1)^*00) | 0$



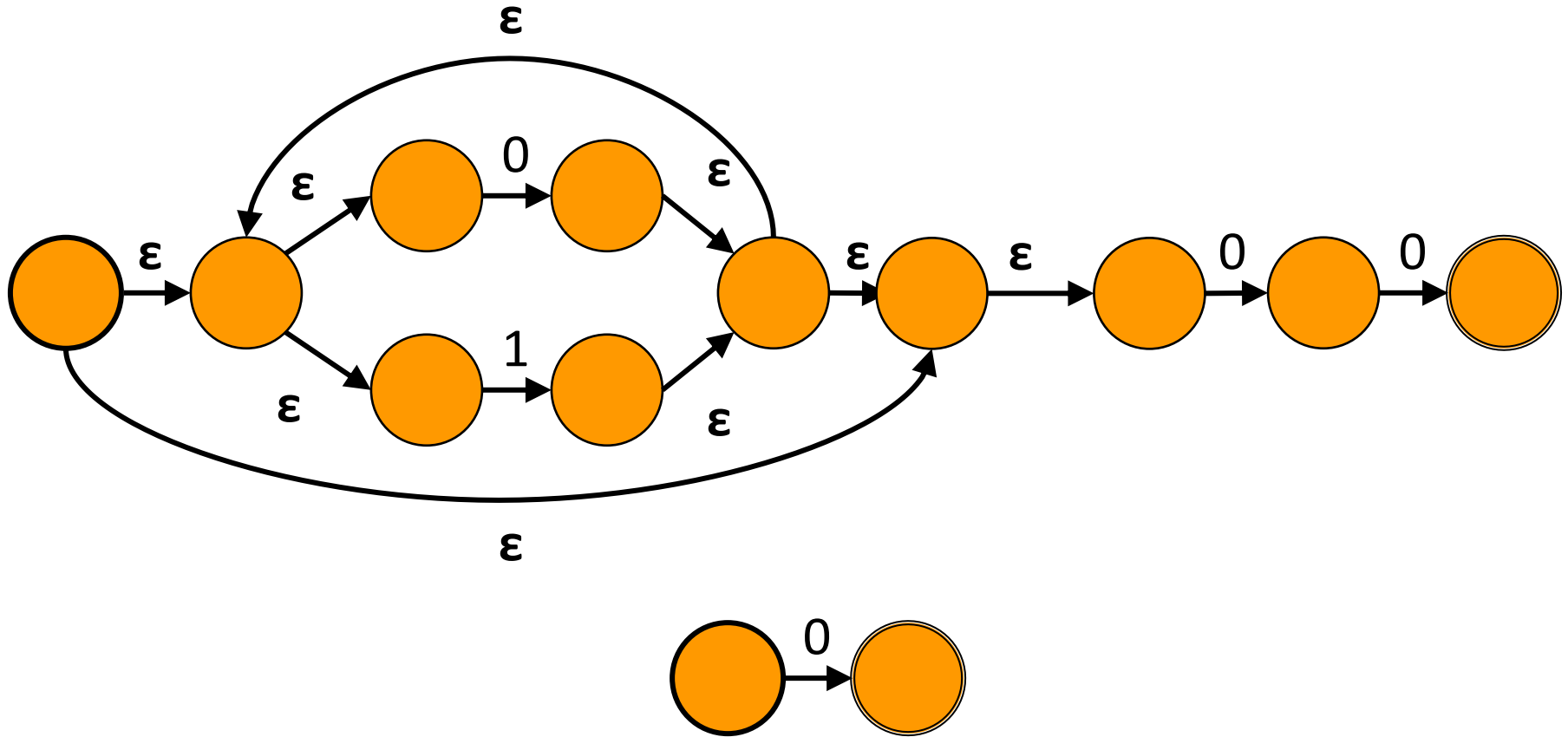
$(0|1)^*00$

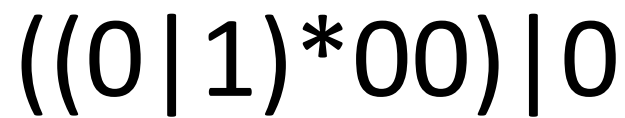
$((0|1)^*00) | 0$



$(0|1)^*00$

$((0|1)^*00) | 0$





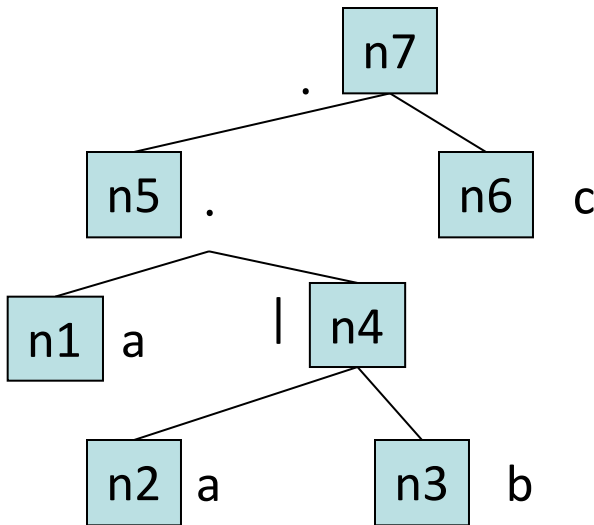
Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree

$(a(a|b))c$
 $aab|.c.$

Post-order traversal of
regexp tree



$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

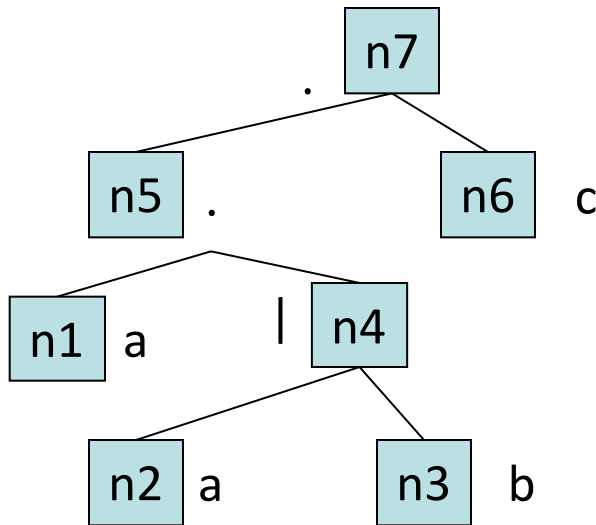
$n6 = \text{nfa}(c)$

$n7 = \text{nfa}(n5, n6, .)$

Thompson's construction

Converts regexps to NFA

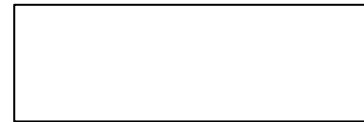
Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$
↑

stack



Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

$n6 = \text{nfa}(c)$

$n7 = \text{nfa}(n5, n6, .)$

Thompson's construction

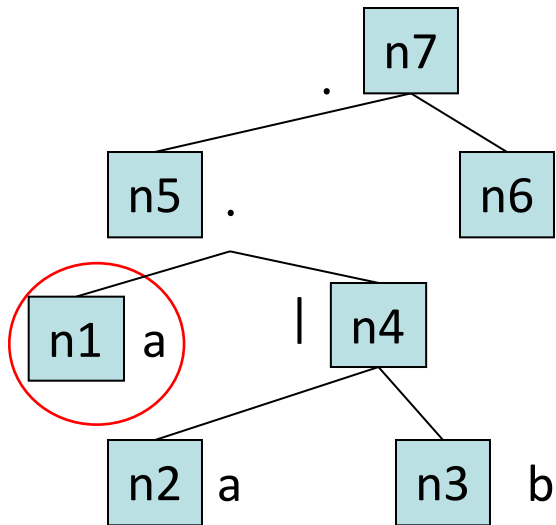
Converts regexps to NFA

Build NFA recursively
from regexp tree

$(a(a|b))c$

$aab|.c.$

Post-order traversal of
regexp tree



c push n1

stack

n1

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

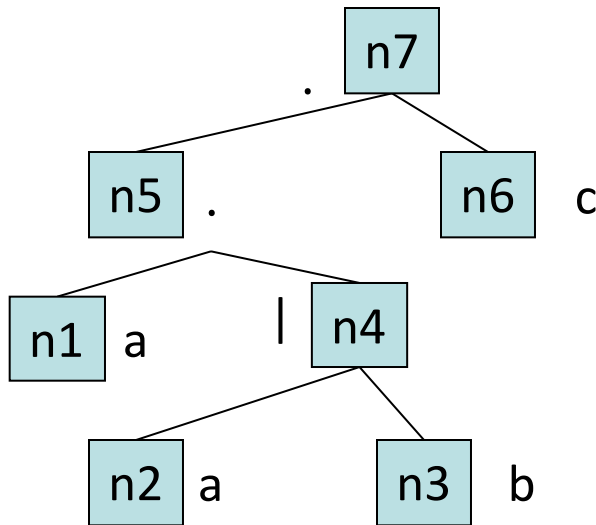
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Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$



stack

n1

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

$n6 = \text{nfa}(c)$

$n7 = \text{nfa}(n5, n6, .)$

Thompson's construction

Converts regexps to NFA

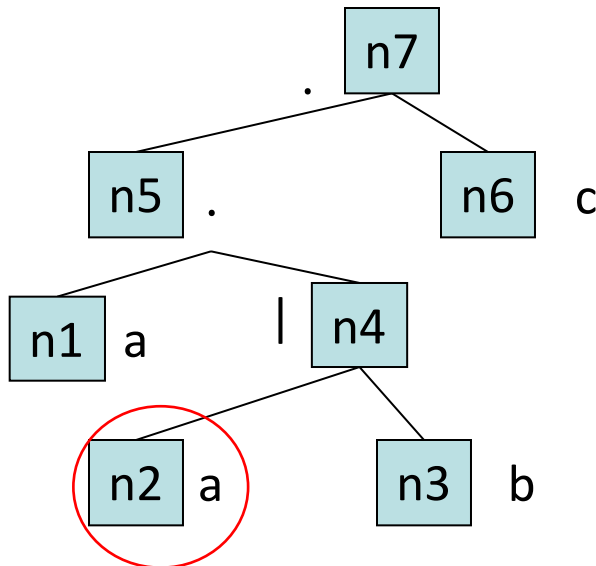
Build NFA recursively
from regexp tree

$(a(a|b))c$

$aab|.c.$



Post-order traversal of
regexp tree



push n2

stack
n2, n1

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

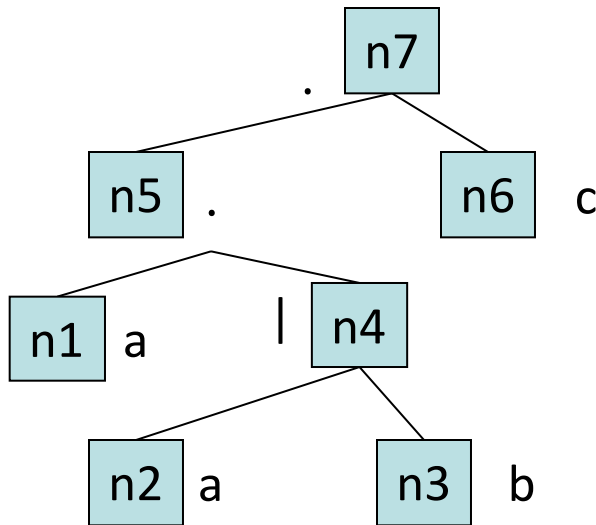
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$n7 = \text{nfa}(n5, n6, .)$

Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$



stack

n2, n1

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

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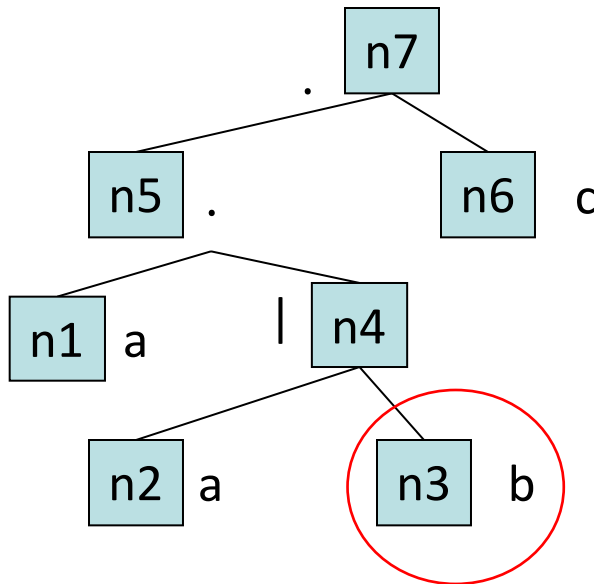
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Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|c.$

stack
push n3
n3, n2, n1

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

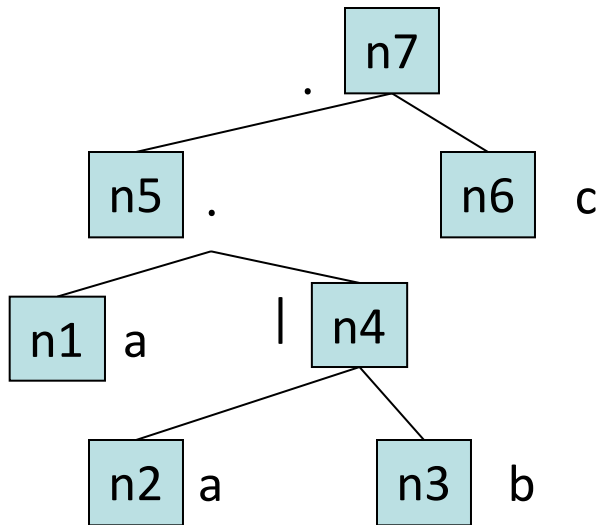
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Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|c.$



stack

n3, n2, n1

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

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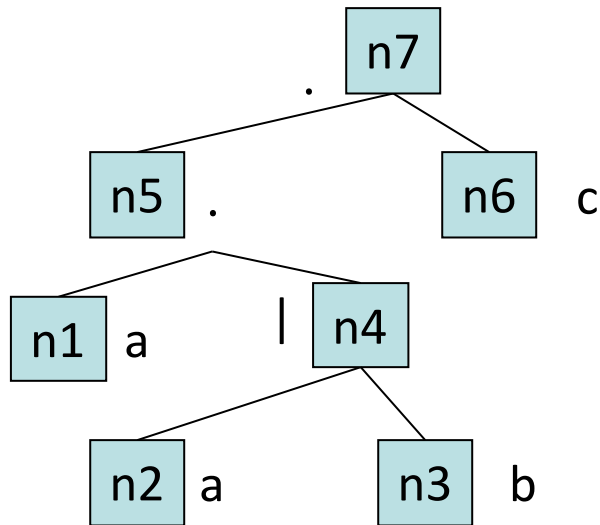
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Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$



stack

pop n3,n2

n1

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

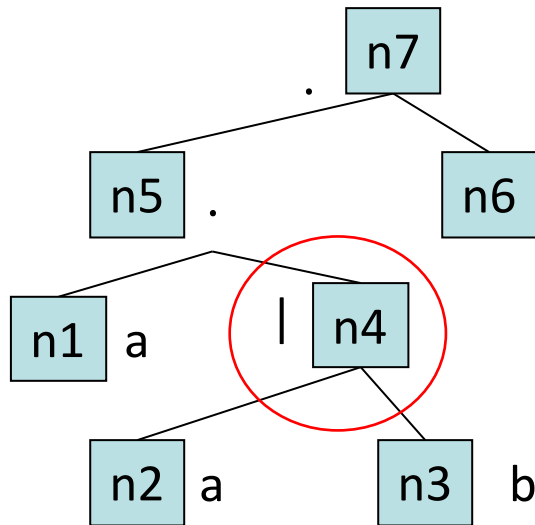
$n6 = \text{nfa}(c)$

$n7 = \text{nfa}(n5, n6, .)$

Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$



c push n4

stack

n4, n1

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

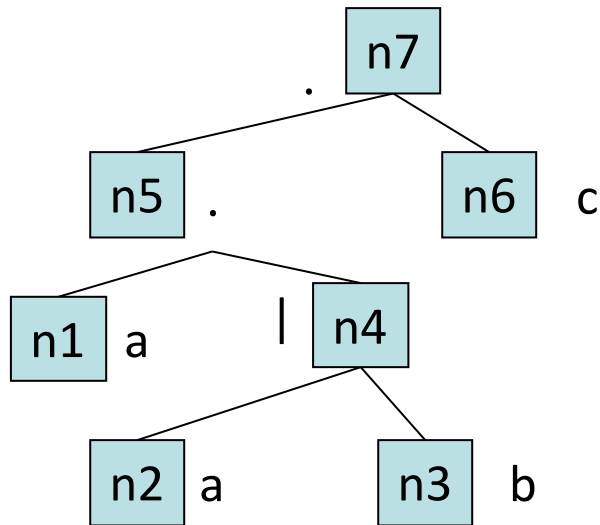
$n6 = \text{nfa}(c)$

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Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$



stack

n4, n1

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

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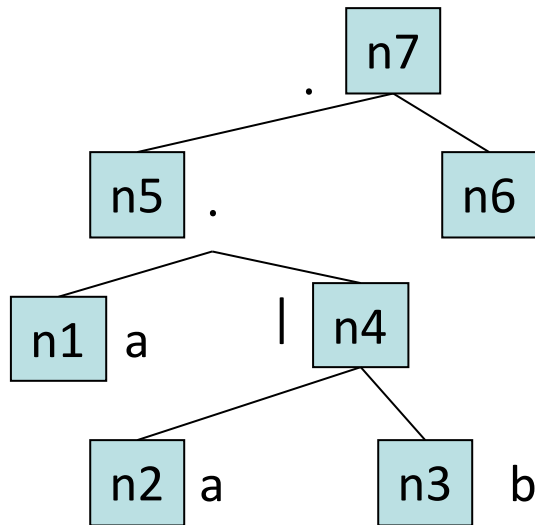
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Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$



stack

c pop n4,n1



Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

$n6 = \text{nfa}(c)$

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Thompson's construction

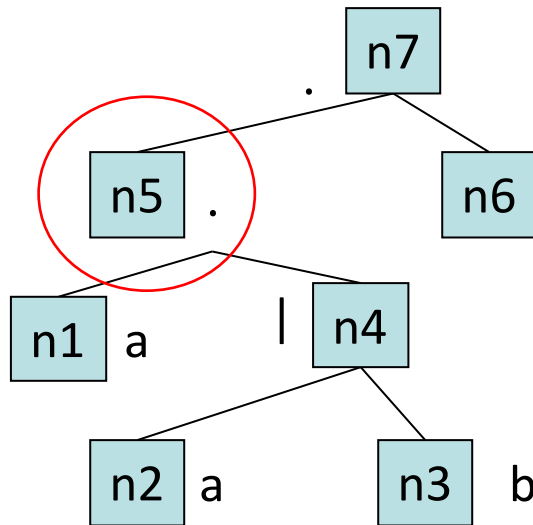
Converts regexps to NFA

Build NFA recursively
from regexp tree

$(a(a|b))c$

$aab|.c.$

Post-order traversal of
regexp tree



c push n5

stack

n5

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

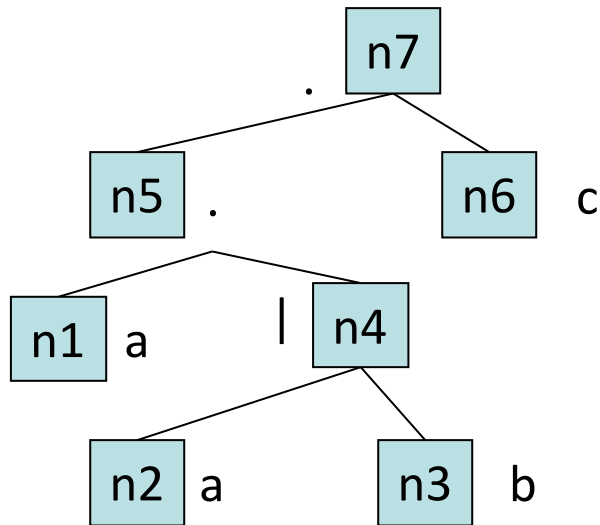
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$n7 = \text{nfa}(n5, n6, .)$

Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$



stack

n5

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

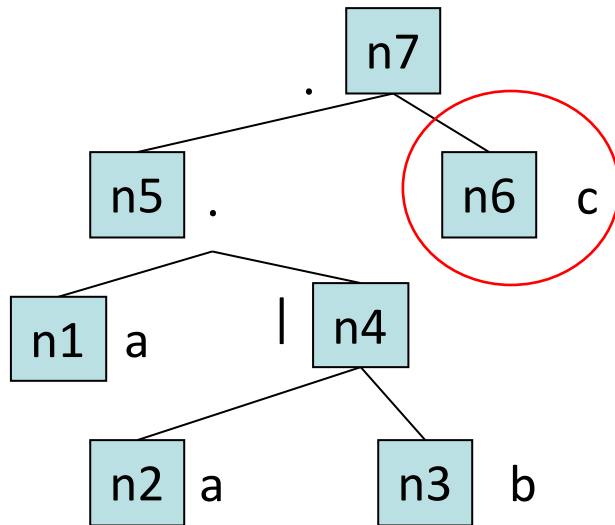
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Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$

push n6



stack

n6, n5

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

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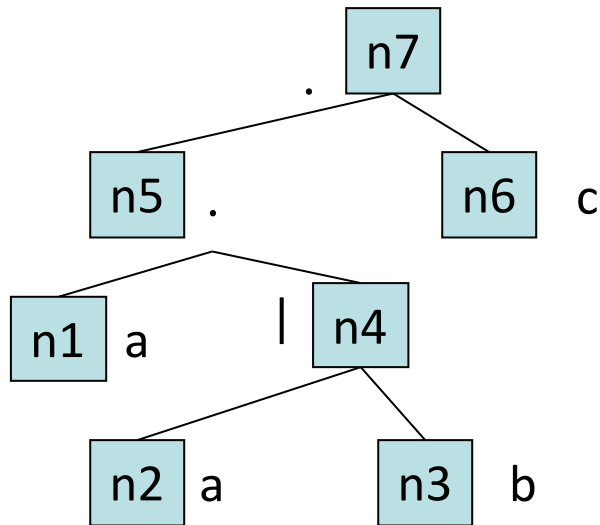
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Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$



stack

n6, n5

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

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Thompson's construction

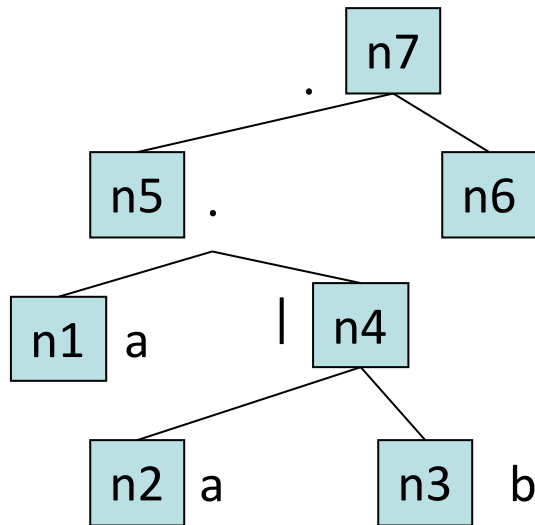
Converts regexps to NFA

Build NFA recursively
from regexp tree

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$aab|.c.$

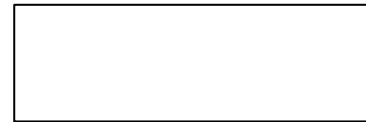
Post-order traversal of
regexp tree



c pop n6, n5



stack



$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

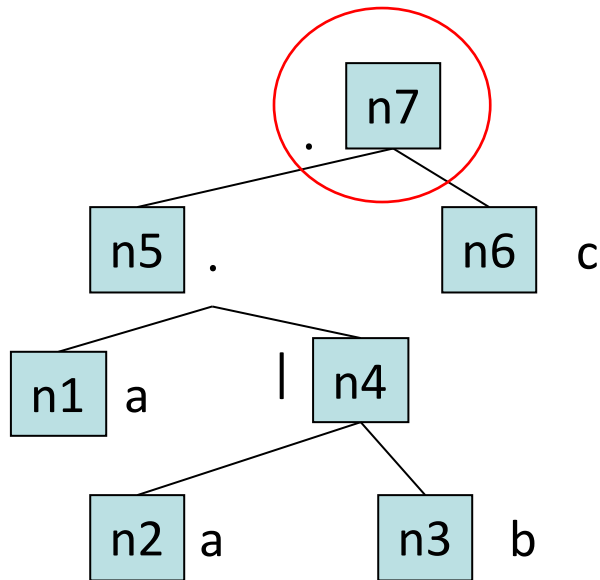
$n6 = \text{nfa}(c)$

$n7 = \text{nfa}(n5, n6, .)$

Thompson's construction

Converts regexps to NFA

Build NFA recursively
from regexp tree



$(a(a|b))c$

$aab|.c.$

push n7

↑
stack

n7

Post-order traversal of
regexp tree

$n1 = \text{nfa}(a)$

$n2 = \text{nfa}(a)$

$n3 = \text{nfa}(b)$

$n4 = \text{nfa}(n2, n3, |)$

$n5 = \text{nfa}(n1, n4, .)$

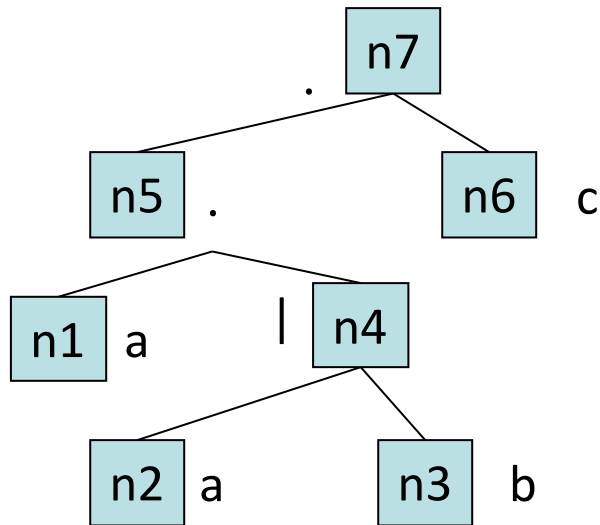
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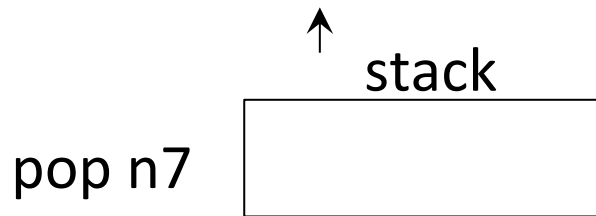
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Post-order traversal of
regexp tree

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 $n5 = \text{nfa}(n1, n4, .)$
 $n6 = \text{nfa}(c)$
 $n7 = \text{nfa}(n5, n6, .)$