

# CMPT 379

## Compilers

Anoop Sarkar

<http://anoopsarkar.github.io/compilers-class/>

# Building a compiler

- Programming languages have a lot in common
- Do not write a compiler for each language
- Create a general mathematical model for the **structure** of all languages
- Implement a compiler using this model

# Building a compiler

- Each language compiler is built using this general model (so-called compiler compilers)
  - yacc = yet another compiler compiler
- Code optimization ideas can also be shared across languages

Demo: compiler for the expr language

# Building a compiler

- The cost of compiling and executing should be managed
- No program that violates the definition of the language should escape
- No program that is valid should be rejected

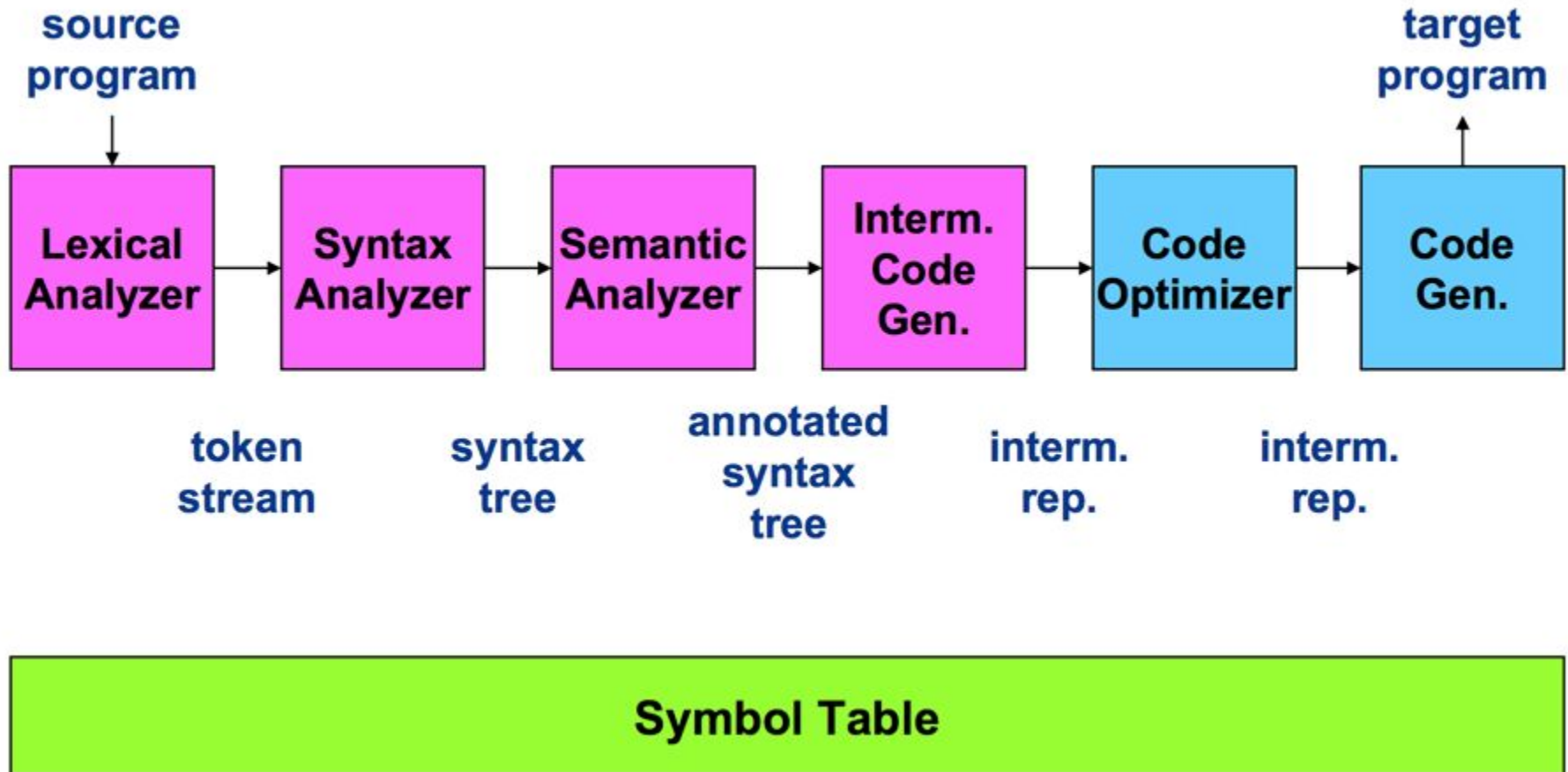
# Building a compiler

- Requirements for building a compiler:
  - Symbol-table management
  - Error detection and reporting
- Stages of a compiler:
  - Analysis (front-end)
  - Synthesis (back-end)

# Stages of a Compiler

- Analysis (Front-end)
  - Lexical analysis
  - Syntax analysis (parsing)
  - Semantic analysis (type-checking)
- Synthesis (Back-end)
  - Intermediate code generation
  - Code optimization
  - Code generation

# Stages of a Compiler





# Compiler Front-end

# Lexical Analysis

- Also called *scanning*, take input program *string* and convert into tokens

- Example:

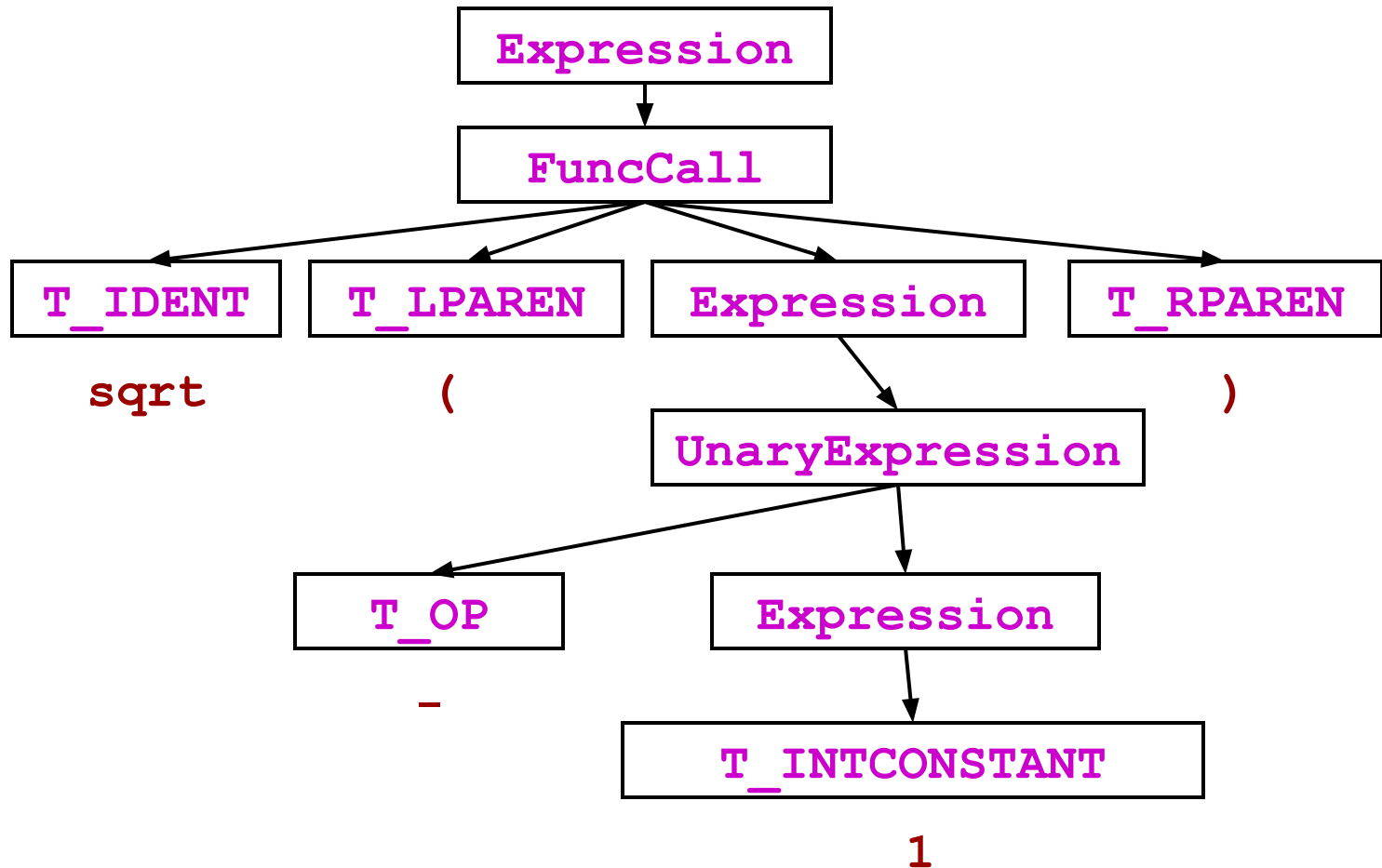
```
double f = sqrt(-1);
```

T_DOUBLE	("double")
T_IDENT	("f")
T_OP	("=")
T_IDENT	("sqrt")
T_LPAREN	("(")
T_OP	("-")
T_INTCONSTANT	("1")
T_RPAREN	(")")
T_SEP	(";")

# Syntax Analysis

- Also called *parsing*
- Describe the set of strings that are programs using a grammar
- Structural validation
- Create a parse tree or derivation

# Parse tree for `sqrt(-1)`



# Abstract Syntax Tree

`sqrt (-1) :=`

`MethodCall (`

`sqrt,`

`UnaryExpr( UnaryMinus,`

`Number(1)`

`)`

`)`

# Semantic analysis

- “does it make sense”? Checking semantic rules,
  - Is there a `main` function?
  - Is variable declared?
  - Are operand types compatible? (coercion)
  - Do function arguments match function declarations?
- Type checking
- Static vs. run-time semantic checks
  - Array bounds, return values do not match definition

# Compiler Back-end

# Source -> abstract syntax tree

```
extern void print_int(int);

class C {
    bool foo() { return(true); }
    int main() {
        if (foo()) {
            print_int(1); }
        }
}
```



# Source -> abstract syntax tree

```
Program(  
  ExternFunction(print_int, VoidType, VarDef(IntType)),  
  Class(C,  
    None,  
    Method(foo,  
      BoolType,  
      None,  
      MethodBlock( None,  
                    ReturnStmt(BoolExpr(True))) ),  
  Method(  main,  
          IntType,  
          None,  
          MethodBlock( None,  
                        IfStmt(MethodCall(foo, None),  
                               Block(None,  
                                     MethodCall  
                                     (print_int, Number(1))),  
                        None) ) ) ) )
```

# Intermediate representation

```
; ModuleID = 'C'

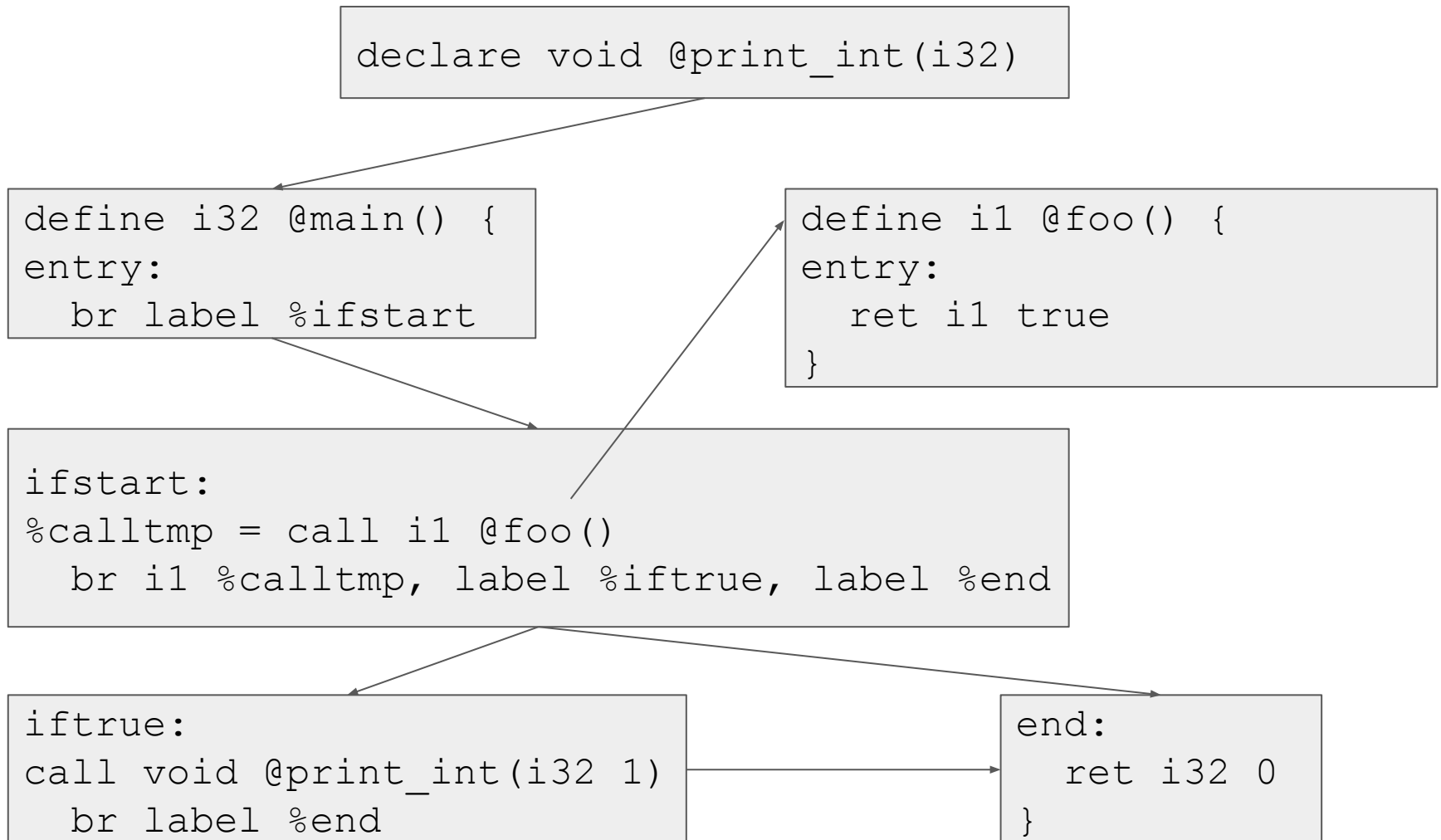
declare void
@print_int(i32)

define i1 @foo() {
entry:
    ret i1 true
}

define i32 @main() {
entry:
    br label %ifstart
ifstart:
    %calltmp = call i1 @foo()
    br i1 %calltmp, label %iftrue, label %end
iftrue:
    call void @print_int(i32 1)
    br label %end
end:
    ret i32 0
}
```

Translation from IR to machine specific assembly

# Intermediate representation



# Assembly language output from IR

x86  
assembly

```
.section __TEXT,__text,  
regular,pure_instructions  
.globl _foo  
.align 4, 0x90  
@foo  
.cfi_startproc  
%entry  
    movl, 1  
    ret  
.cfi_endproc  
  
.globl _main  
.align 4, 0x90
```

```
@main  
.cfi_startproc  
%entry  
    push    rax  
Ltmp0:  
.cfi_def_cfa_offset 16  
    call _foo  
    test al, 1  
    je     LBB1_2  
%iftrue  
    movedi, 1  
    call _print_int  
%end  
    xor    eax, eax  
    pop    rdx  
    ret  
.cfi_endproc
```

# Code optimization

```
; ModuleID = 'C'

declare void @print_int(i32)

define i32 @main() {
entry:
    br label %ifstart

ifstart:
    call void @print_int(i32 1)
    br label %end

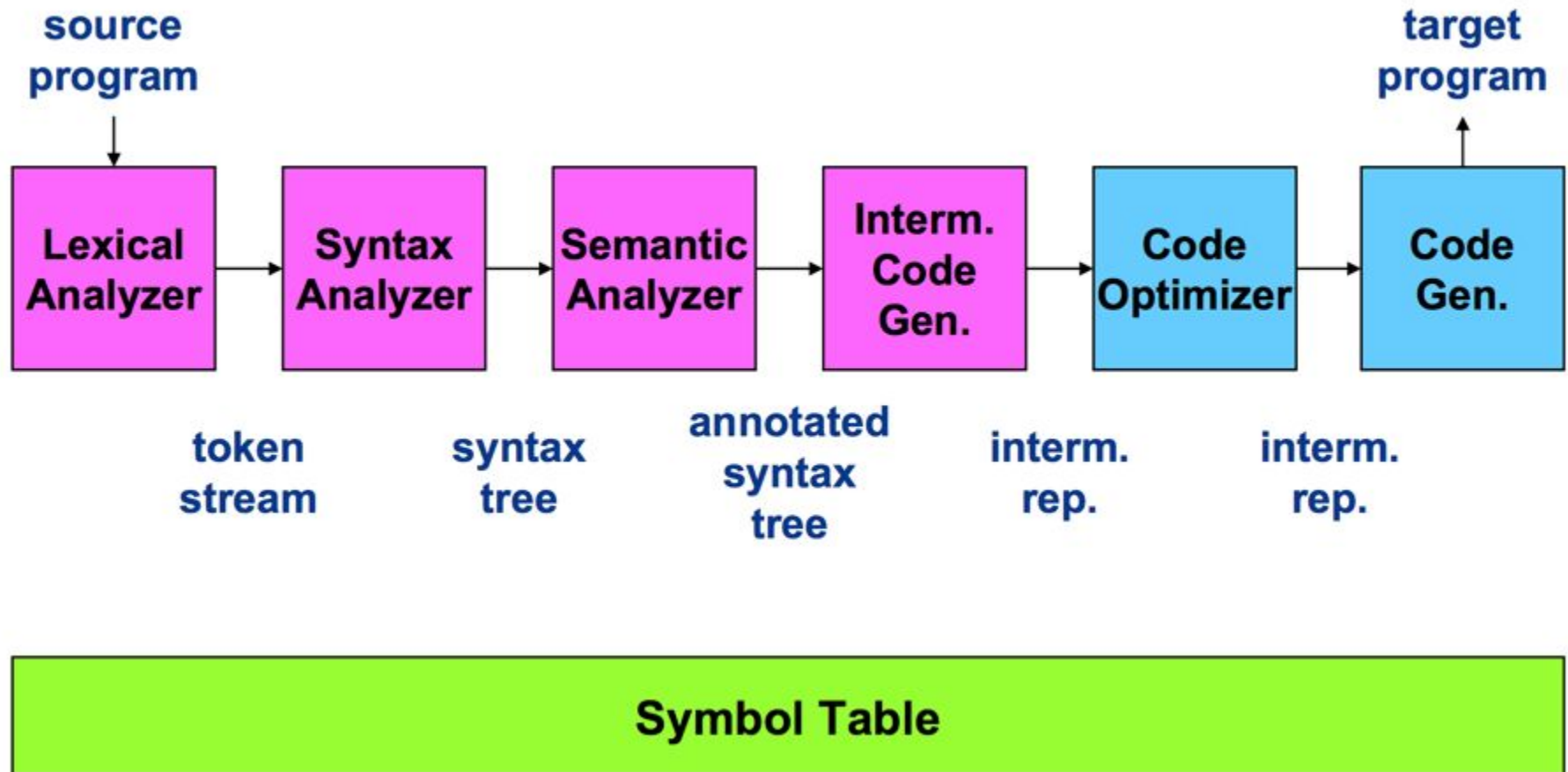
end:
    ret i32 0
}
```

# Code Optimization

x86  
assembly

```
.section __TEXT,__text,regular,  
pure_instructions  
.macosx_version_min 10, 11  
.globl    _main  
.p2align  4, 0x90  
_main:  
.cfi_startproc  
## BB#0:  
    pushq   %rax  
Ltmp0:  
    .cfi_def_cfa_offset 16  
    movl    $1, %edi  
    callq   _print_int  
    xorl    %eax, %eax  
    popq    %rcx  
    retq  
.cfi_endproc
```

# Stages of a Compiler



# Wrap Up

- Analysis/Synthesis
  - Translation from string to executable
- Divide and conquer
  - Build one component at a time
  - Theoretical analysis will ensure we keep things **simple** and **correct**
  - Create a complex piece of software