LEX5: Regexps to NFA

Lexical Analysis

CMPT 379: Compilers

Instructor: Anoop Sarkar

anoopsarkar.github.io/compilers-class

Building a Lexical Analyzer

- Token ⇒ Pattern
- Pattern ⇒ Regular Expression
- Regular Expression ⇒ NFA
- NFA ⇒ DFA
- DFA ⇒ Table-driven implementation of DFA

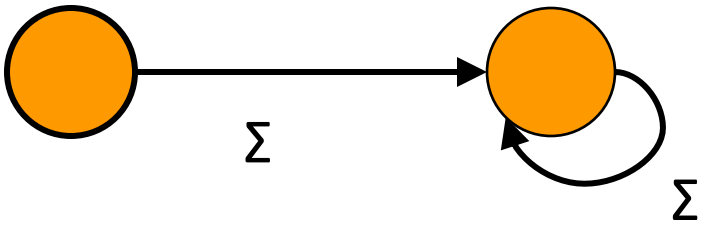
Thompson's construction

- Converts regexps to equivalent NFA
- Six simple rules
  - Empty language
  - Symbols (Σ)
  - Empty String (ε)
  - Alternation (r_1 or r_2)
  - Concatenation (r_1 followed by r_2)
  - Repetition (r_1*)

Used by Ken Thompson for pattern-based search in text editor QED (1968)

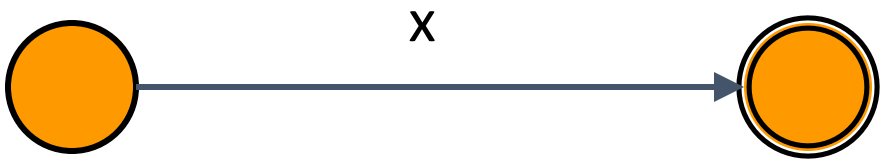
Thompson Rule 0

- For the empty language ϕ
- (optionally include a *sinkhole* state)



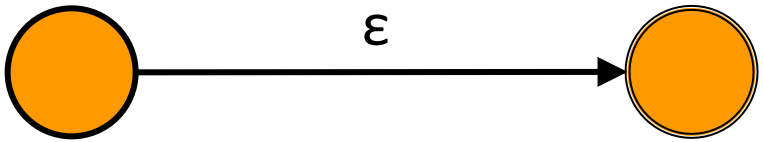
Thompson Rule 1

- For each symbol *x* of the alphabet, there is a NFA that accepts it



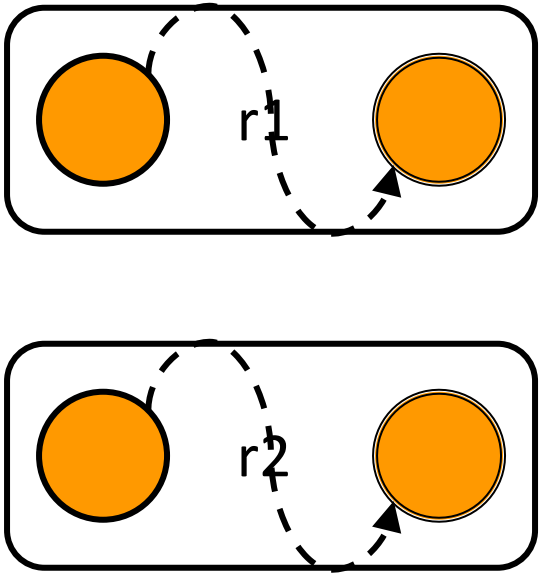
Thompson Rule 2

- There is an NFA that accepts only ε



Thompson Rule 3

- Given 2 NFAs r_1, r_2, there is a NFA that accepts $r_1 | r_2$



Thompson Rule 3

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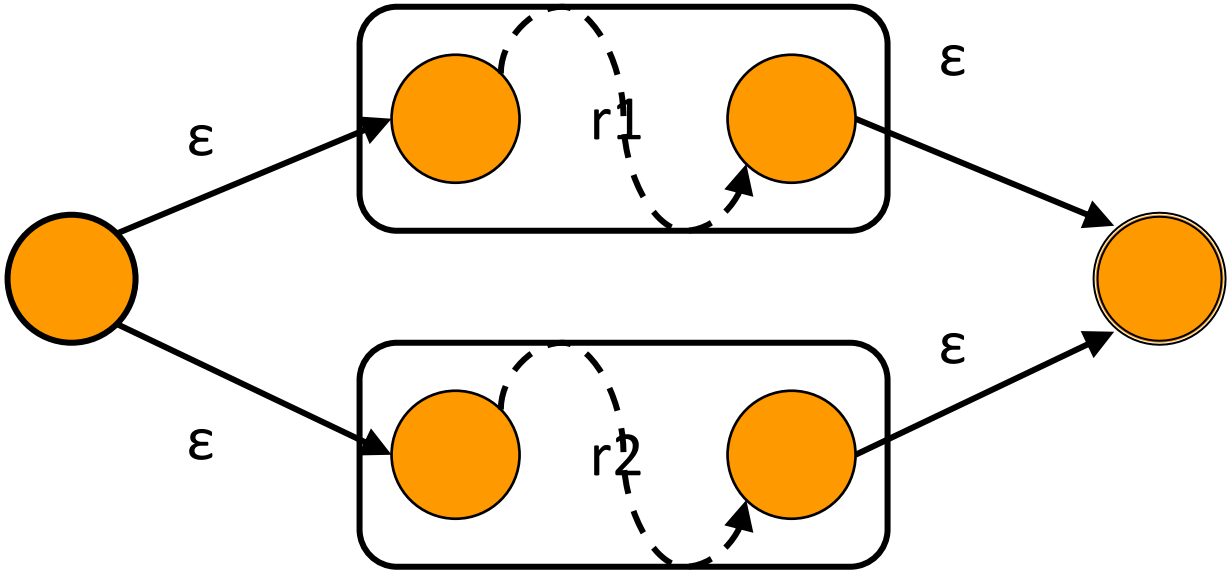
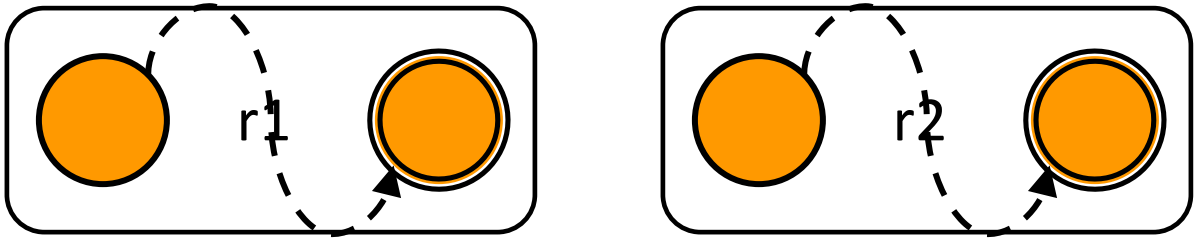


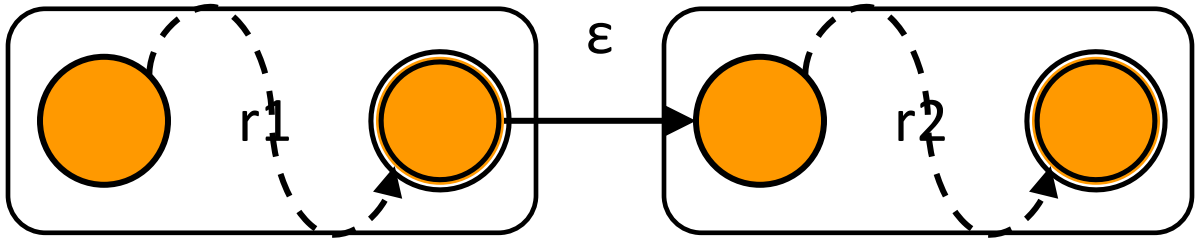
Thompson Rule 4

- Given 2 NFAs r_1, r_2, there is a NFA that accepts $r_1 r_2$



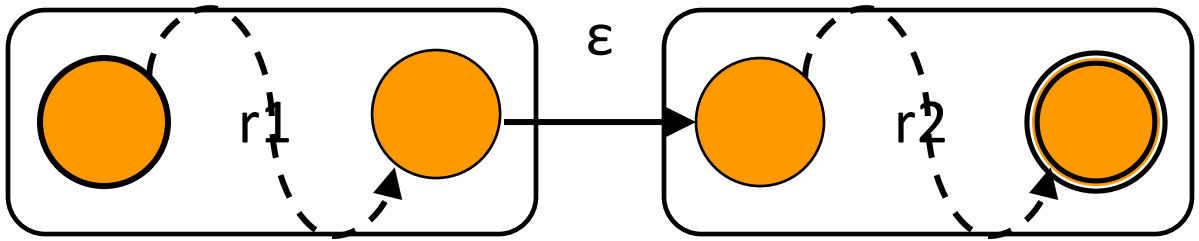
Thompson Rule 4

- Given 2 NFAs r_1, r_2, there is a NFA that accepts $r_1 r_2$



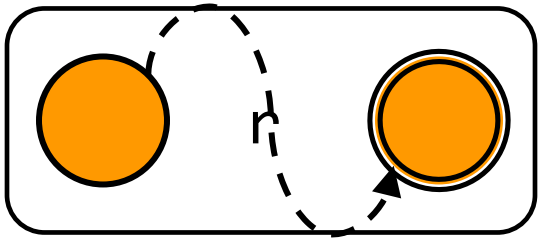
Thompson Rule 4

- Given 2 NFAs r_1, r_2, there is a NFA that accepts $r_1 r_2$



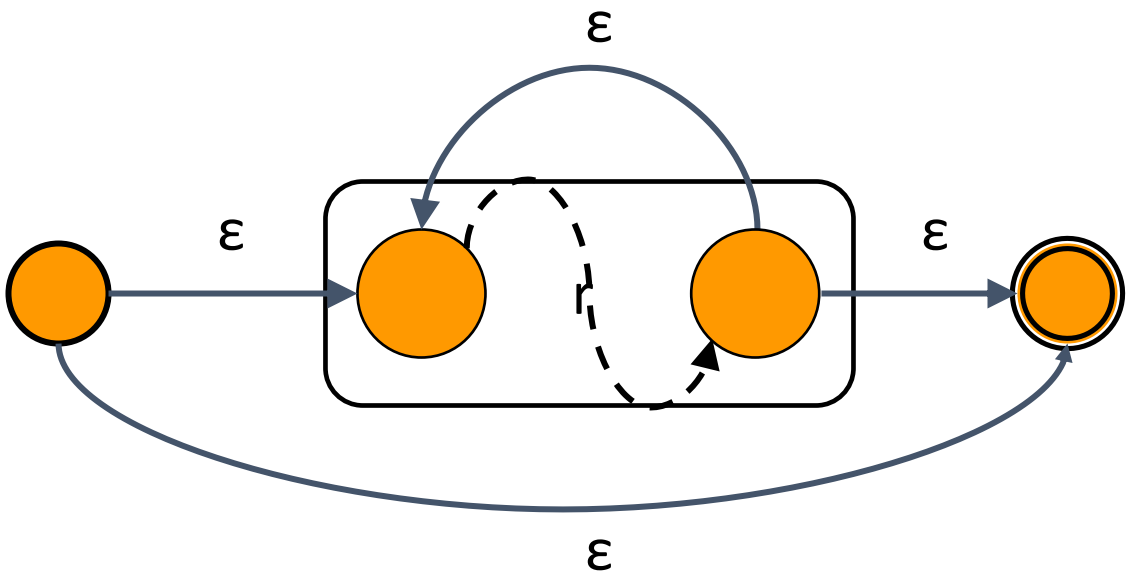
Thompson Rule 5

- Given an NFA for r, there is an NFA that accepts *r**



Thompson Rule 5

- Given an NFA for r, there is an NFA that accepts *r**

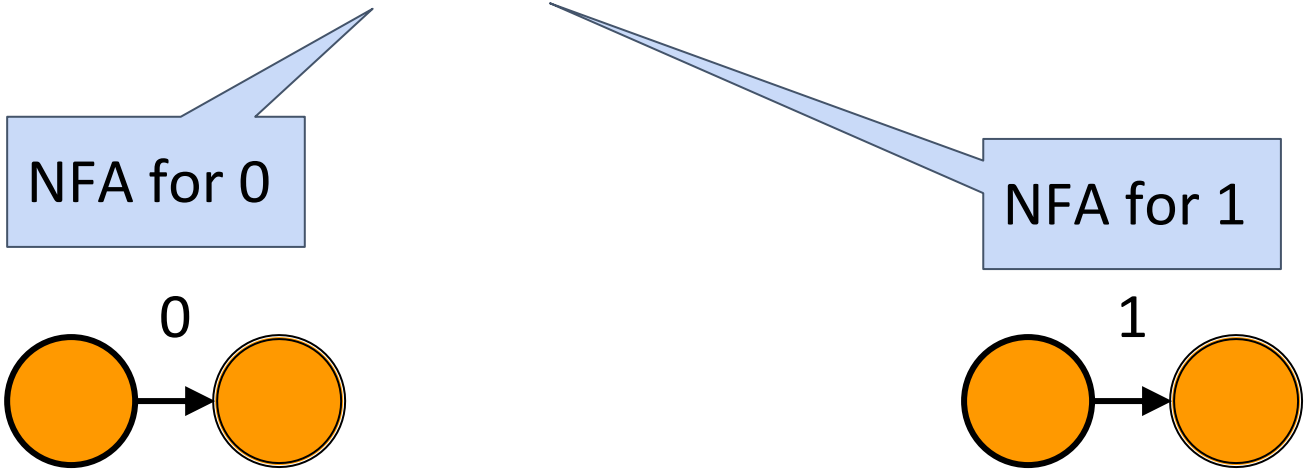


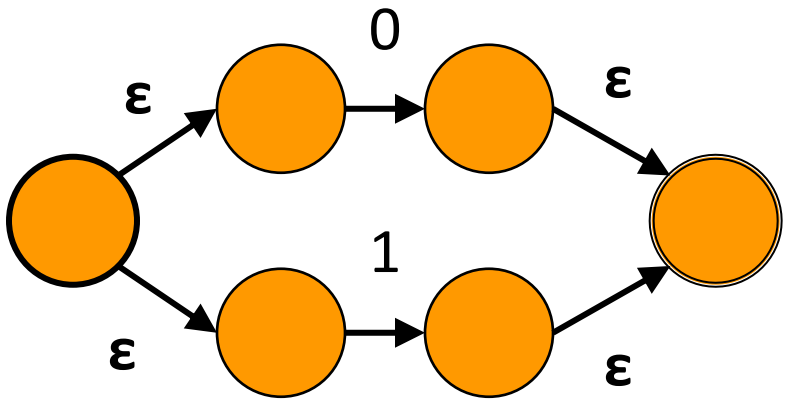
Example

- Set of all binary strings that are divisible by four (include 0 in this set)
- Defined by the regexp: ((0|1)*00) | 0
- Apply Thompson's Rules to create an NFA

Basic Blocks 0 and 1

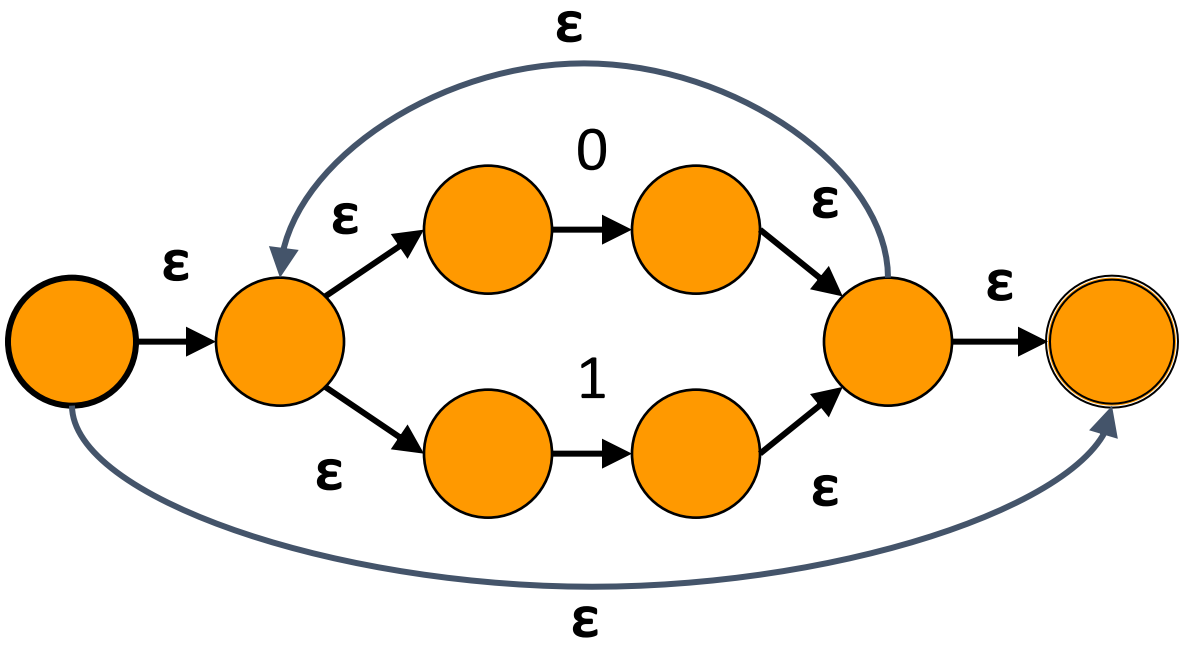
((0|1)*00) | 0





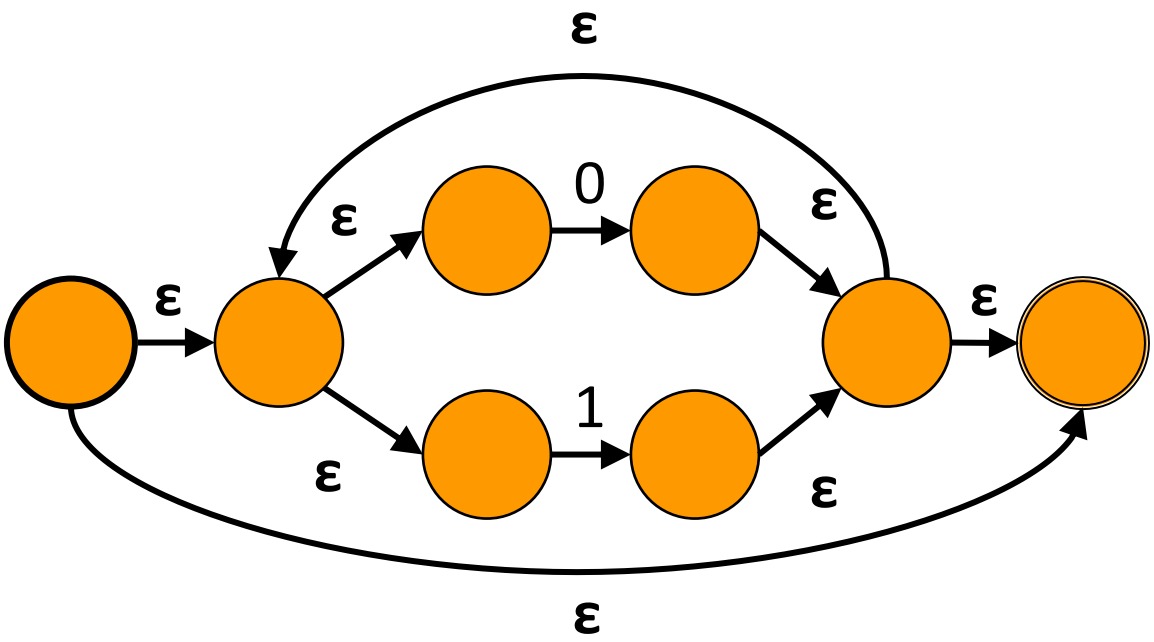
0|1

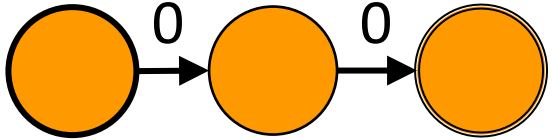
((0|1)*00) | 0



(0|1)*

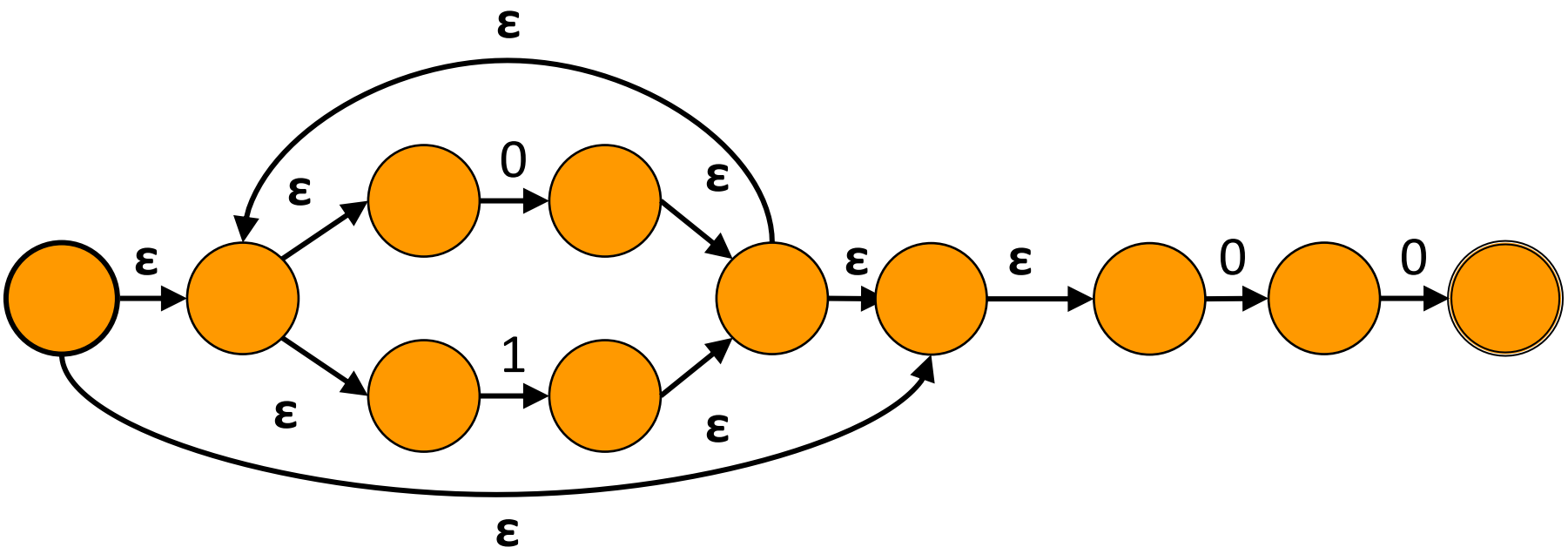
((0|1)*00) | 0





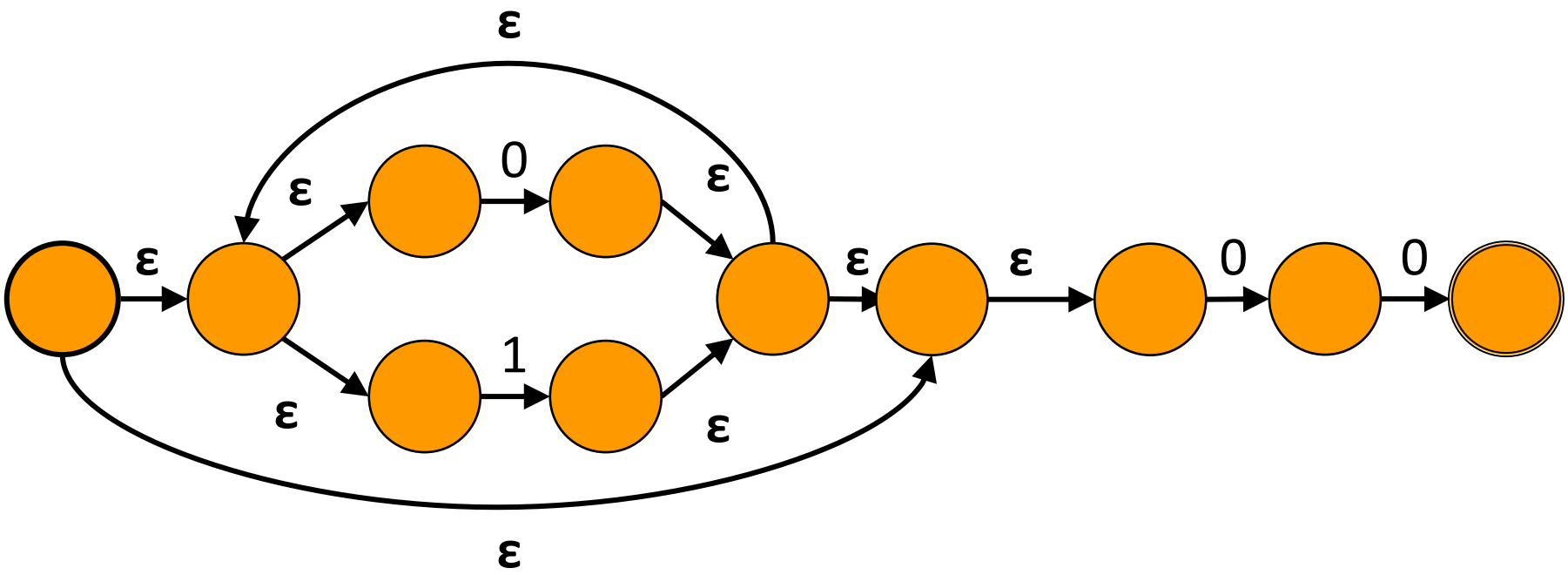
(0|1)*00

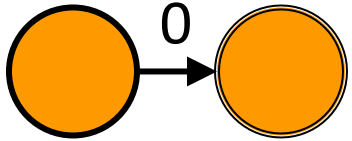
((0|1)*00) | 0



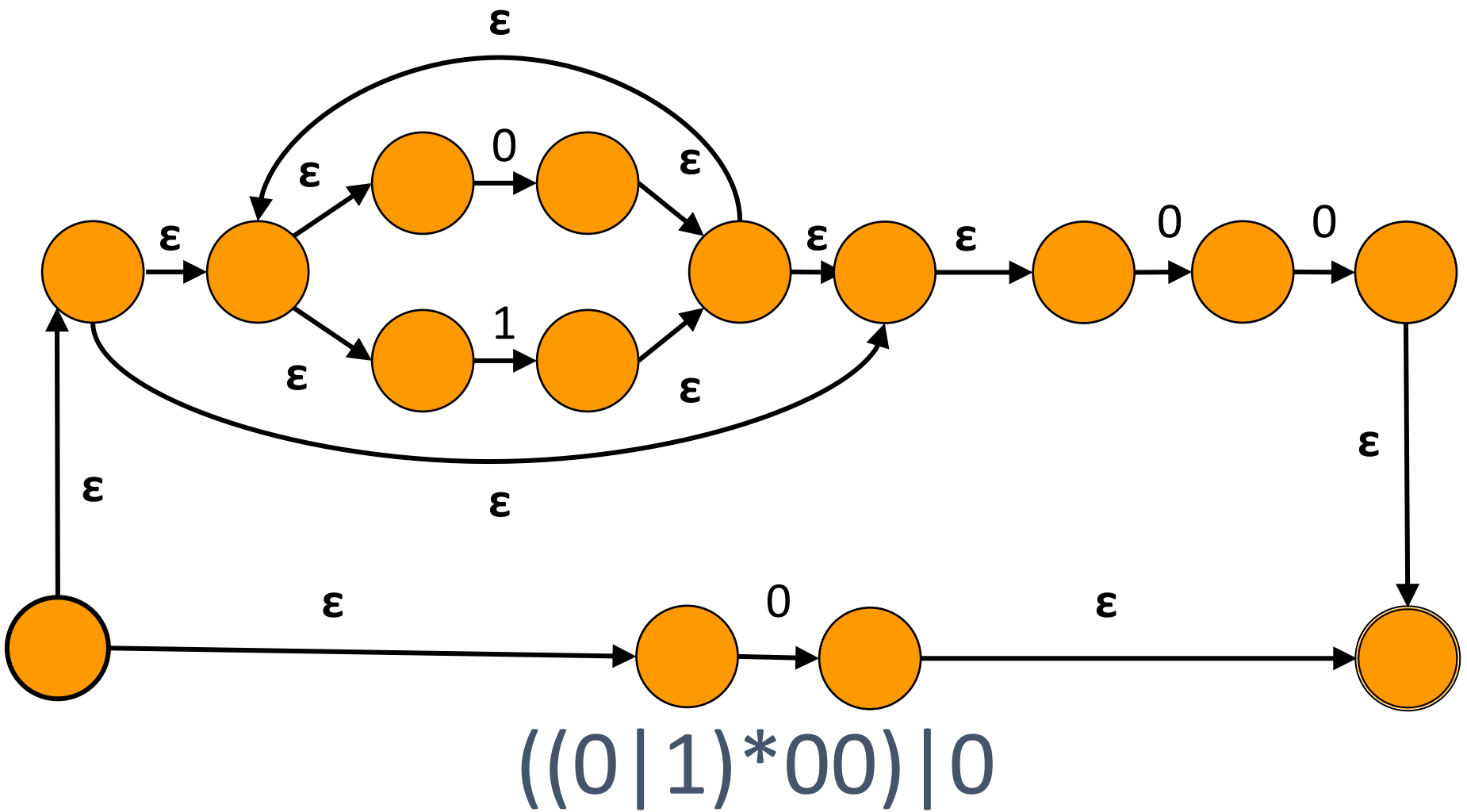
(0|1)*00

((0|1)*00) | 0





((0|1)*00)|0

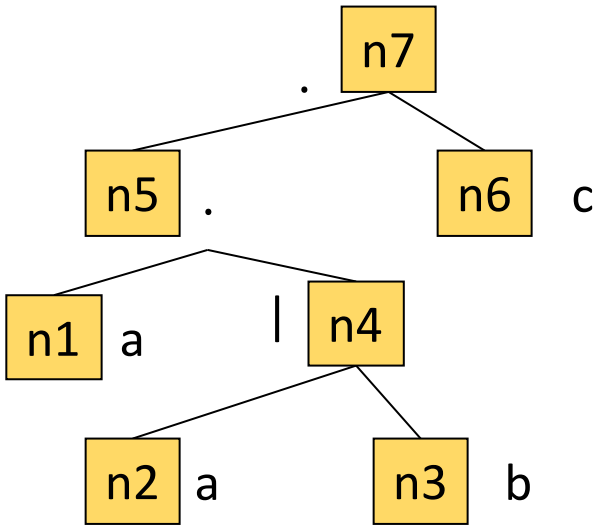


((0|1)*00)|0

Thompson's construction

Converts regexps to NFA

Build NFA recursively from the regexp tree



(a(a|b))c

aab|.c.

Input is the tree in postfix

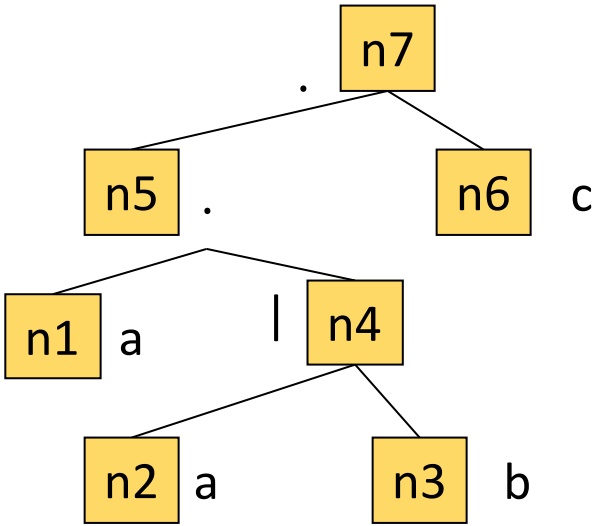
Post-order traversal of the regexp tree

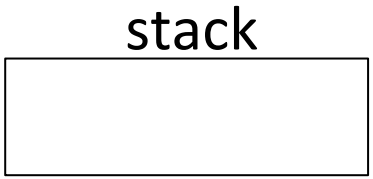
n1= nfa(a)
n2= nfa(a)
n3= nfa(b)
n4= nfa(n2, n3, |)
n5= nfa(n1, n4, .)
n6= nfa(c)
n7= nfa(n5, n6, .)

Thompson's construction

(a(a|b))c

aab|.c.



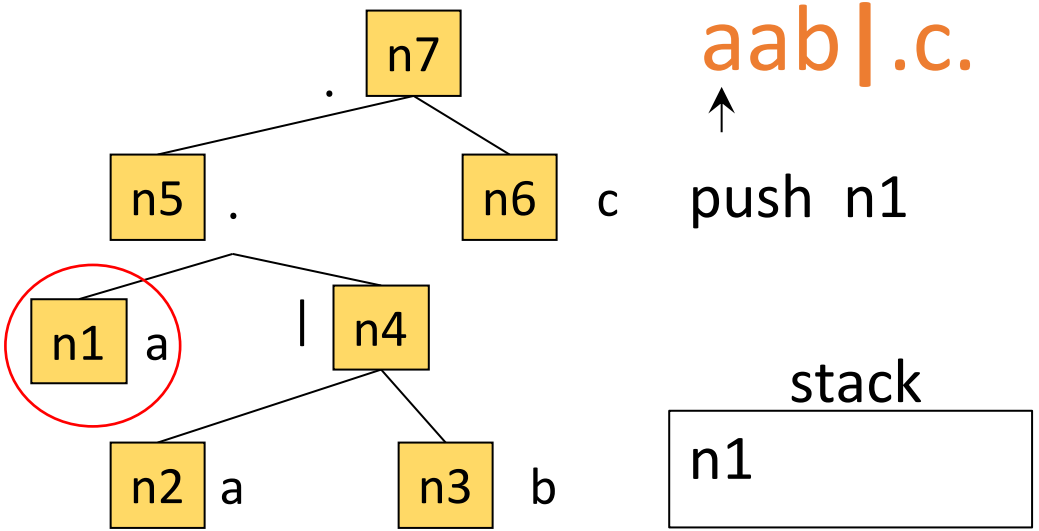


Thompson's construction

(a(a|b))c

aab|.c.

n1= nfa(a)



push n1

stack

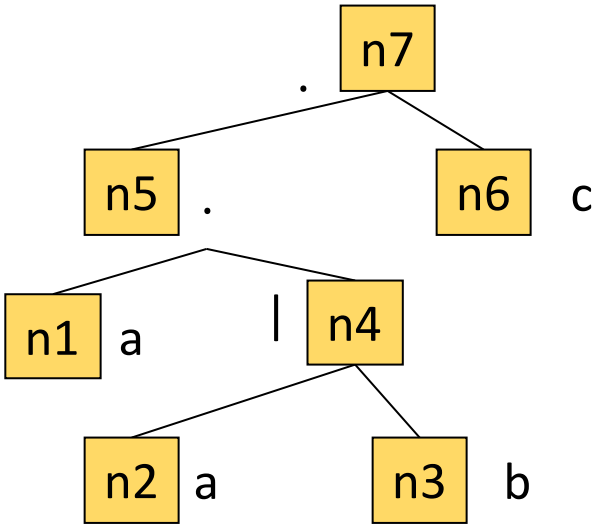
n1

Thompson's construction

(a(a|b))c

aab|.c.

n1= nfa(a)



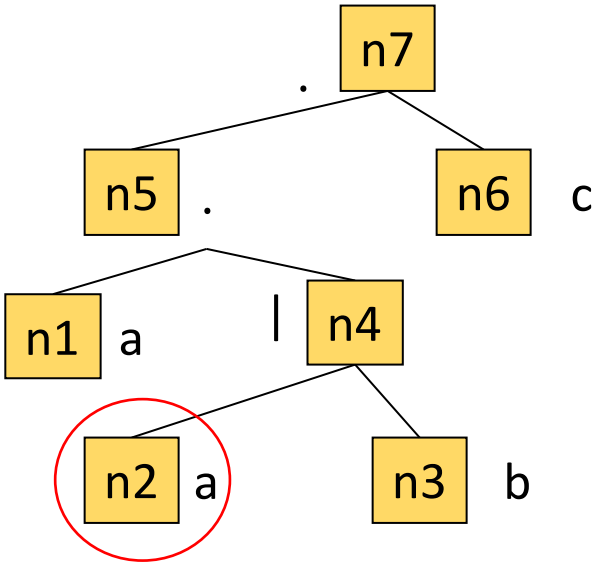
stack

n1

Thompson's construction

(a(a|b))c

aab|.c.



push n2

n1= nfa(a)

n2= nfa(a)

stack

n2, n1

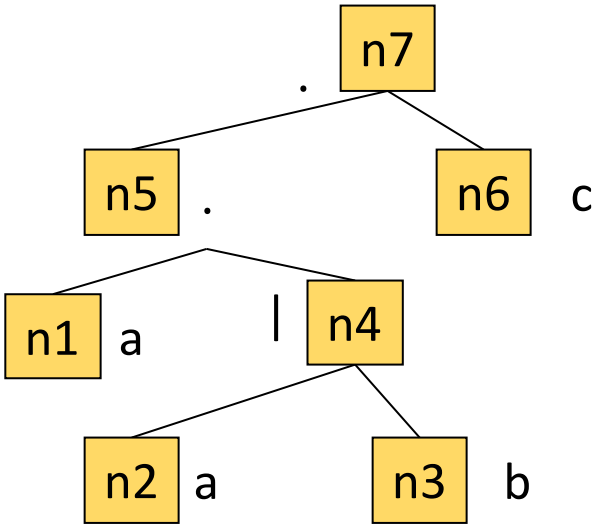
Thompson's construction

(a(a|b))c

aab|.c.

n1= nfa(a)

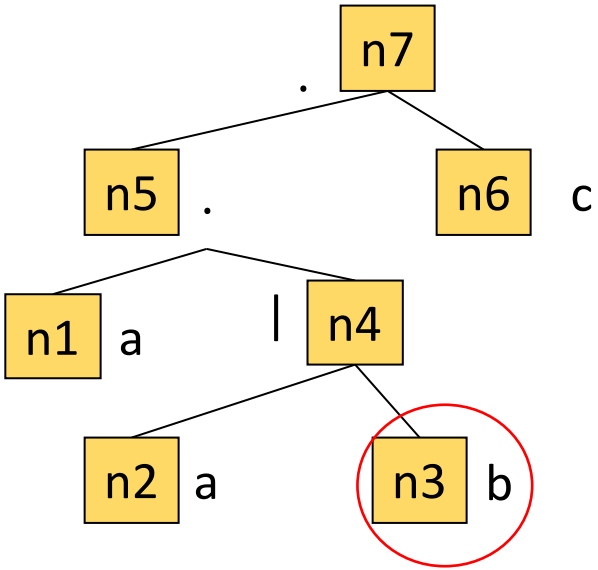
n2= nfa(a)



stack

n2, n1

Thompson's construction



(a(a|b))c

aab|.c.

push n3

n1= nfa(a)

n2= nfa(a)

n3= nfa(b)

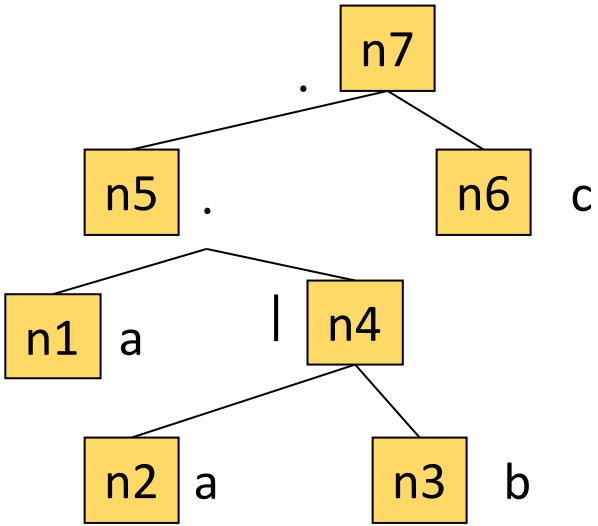
stack

n3, n2, n1

Thompson's construction

(a(a|b))c

aab|.c.



n1= nfa(a)

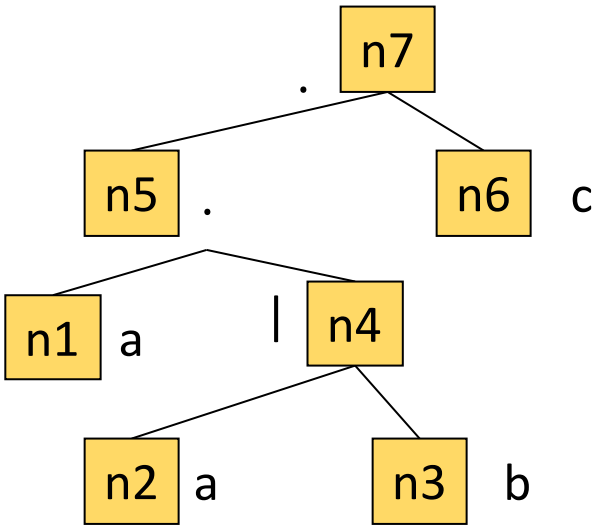
n2= nfa(a)

n3= nfa(b)

stack

n3, n2, n1

Thompson's construction



(a(a|b))c

aab|.c.

pop n3,n2

n1= nfa(a)

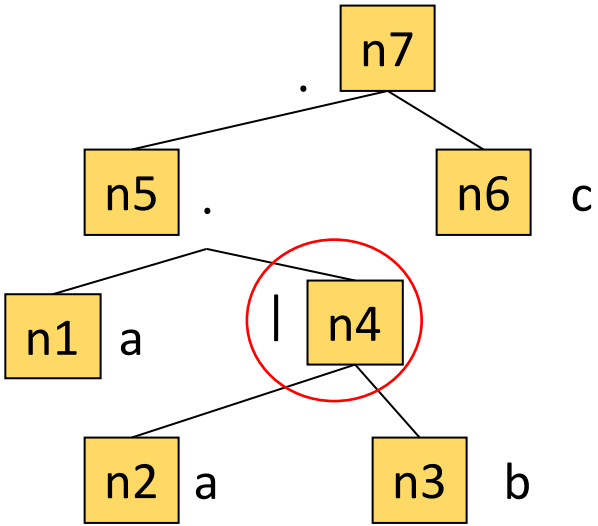
n2= nfa(a)

n3= nfa(b)

stack

n1

Thompson's construction



(a(a|b))c

aab|.c.

push n4

n1= nfa(a)

n2= nfa(a)

n3= nfa(b)

n4= nfa(n2, n3, |)

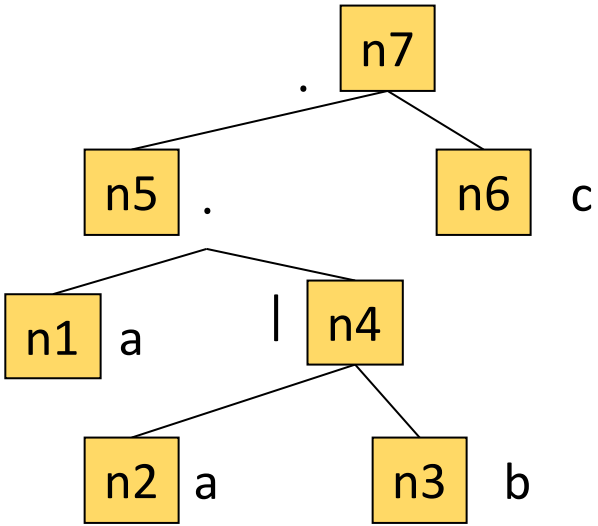
stack

n4, n1

Thompson's construction

(a(a|b))c

aab|.c.



n1= nfa(a)

n4= nfa(n2, n3, |)

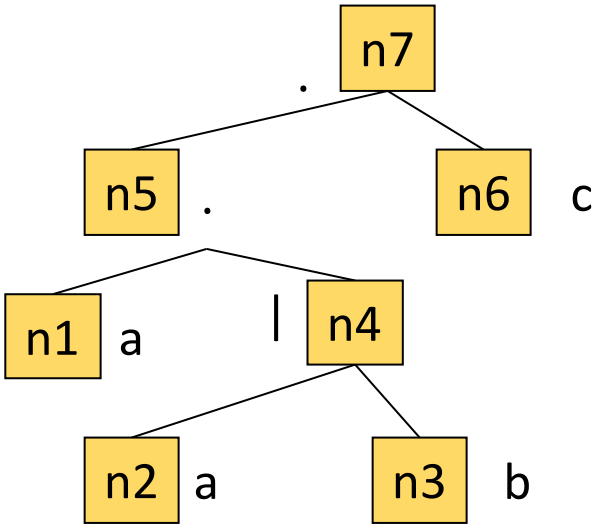
stack

n4, n1

Thompson's construction

(a(a|b))c

aab|.c.



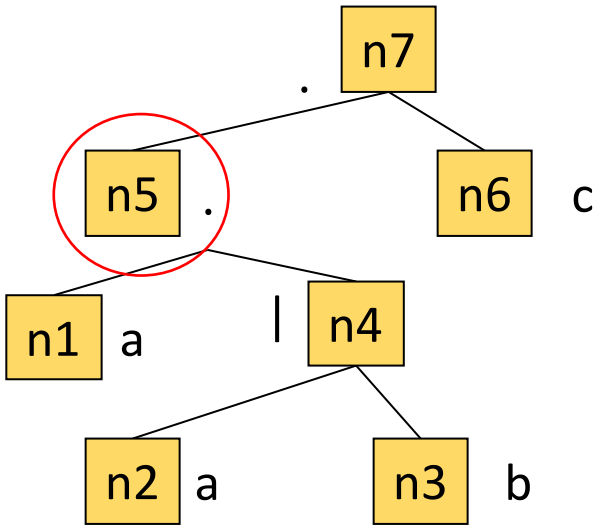
pop n4,n1

n1= nfa(a)

n4= nfa(n2, n3, |)

stack

Thompson's construction



(a(a|b))c

aab|.c.

push n5

stack

n5

n1= nfa(a)

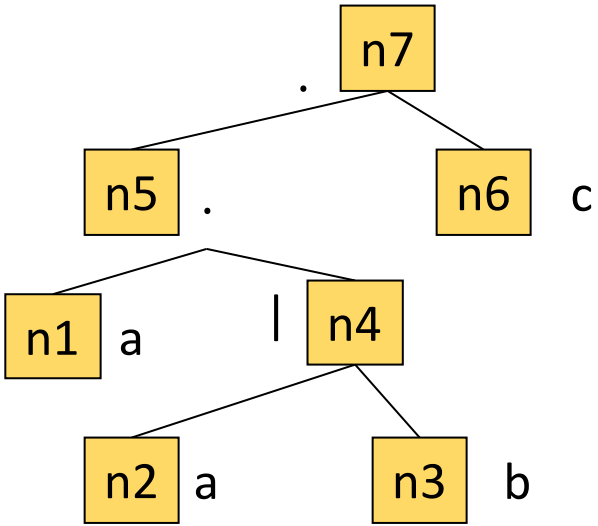
n4= nfa(n2, n3, |)

n5= nfa(n1, n4, .)

Thompson's construction

(a(a|b))c

aab|.c.



stack

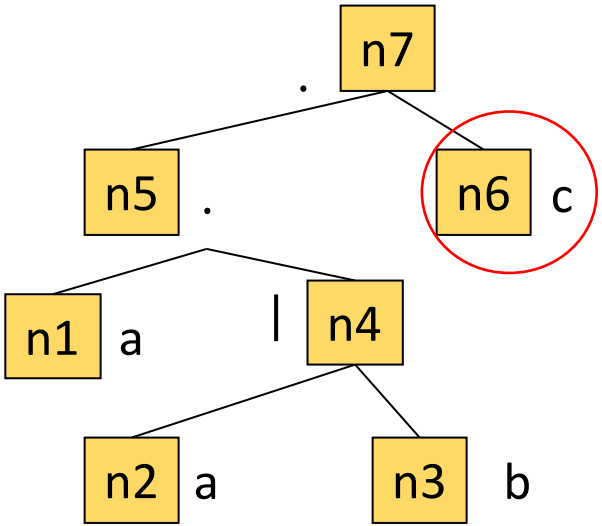
n5

n5= nfa(n1, n4, .)

Thompson's construction

(a(a|b))c

aab|.c.



push n6

stack

n6, n5

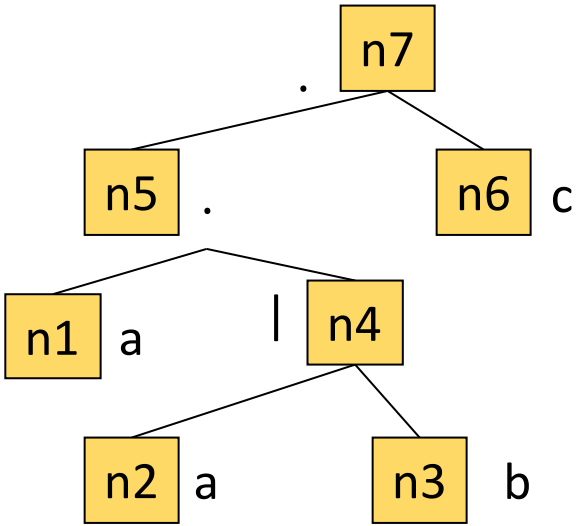
n5= nfa(n1, n4, .)

n6= nfa(c)

Thompson's construction

(a(a|b))c

aab|.c.



stack

n6, n5

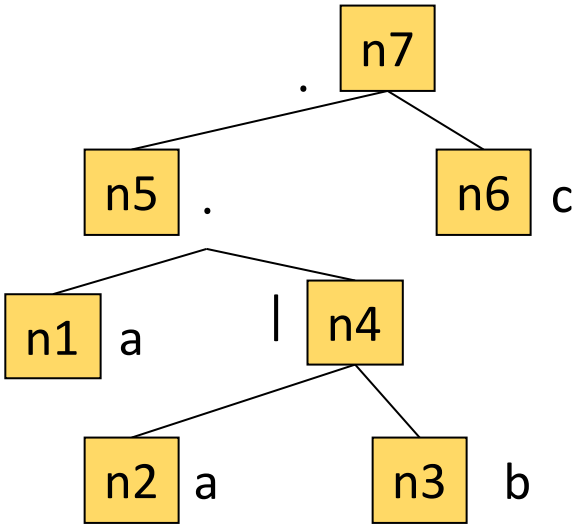
n5= nfa(n1, n4, .)

n6= nfa(c)

Thompson's construction

(a(a|b))c

aab|.c.



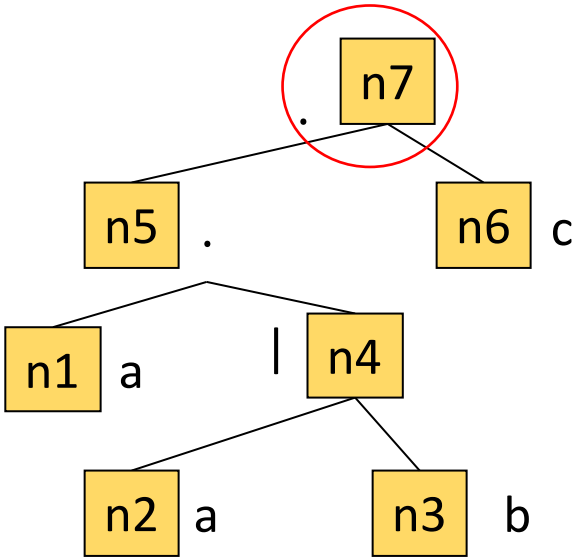
pop n6, n5

stack

n5= nfa(n1, n4, .)

n6= nfa(c)

Thompson's construction



(a(a|b))c

aab|.c.

push n7

stack

n7

n5= nfa(n1, n4, .)

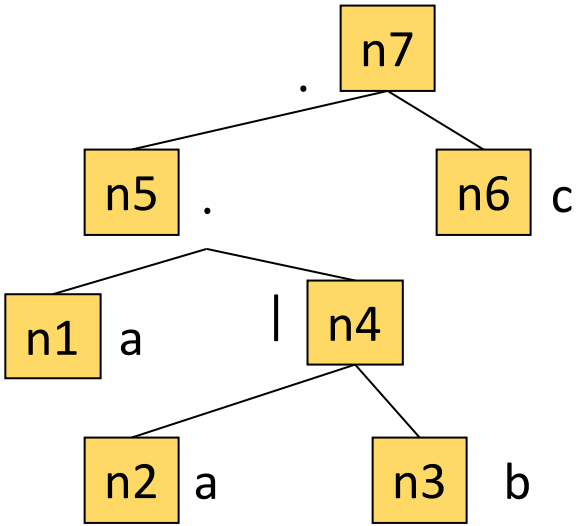
n6= nfa(c)

n7= nfa(n5, n6, .)

Thompson's construction

(a(a|b))c

aab|.c.



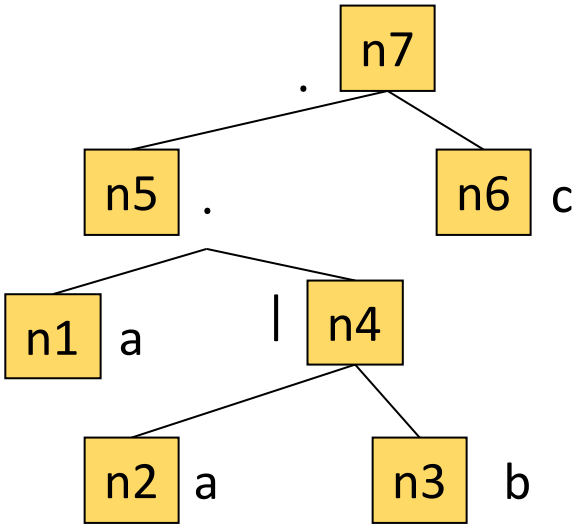
pop n7

stack

n7= nfa(n5, n6, .)

Thompson's construction

Q: Use Thompson's construction to build an NFA for `(0|1)(0|1)*`



(a(a|b))c

aab|.c.

no more input

pop n7

return(n7)

stack

stack is empty