CFG3: Push-down Automata

Context-Free Grammars

CMPT 379: Compilers

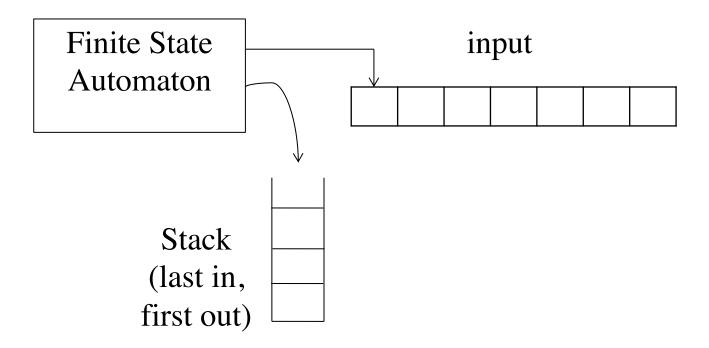
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anoopsarkar.github.io/compilers-class

Context-free languages and Pushdown Automata

- Recall that for each regular language there was an equivalent finite-state automaton
- The FSA was used as a recognizer of the regular language
- For each context-free language there is also an automaton that recognizes it: called a pushdown automaton (pda)

Pushdown Automata



Context-free languages and Pushdown Automata

- Similar to FSAs there are non-deterministic pda and deterministic pda
- Unlike in the case of FSAs we cannot always convert a npda to a dpda
- Our goal in compiler design will be to choose grammars carefully so that we can always provide a dpda for it
- Similar to the FSA case, a DFA construction provides us with the algorithm for lexical analysis,
- In this case the construction of a dpda will provide us with the algorithm for parsing (take in strings and provide the parse tree)
- We will study later how to convert a given CFG into a parser by first converting into a PDA

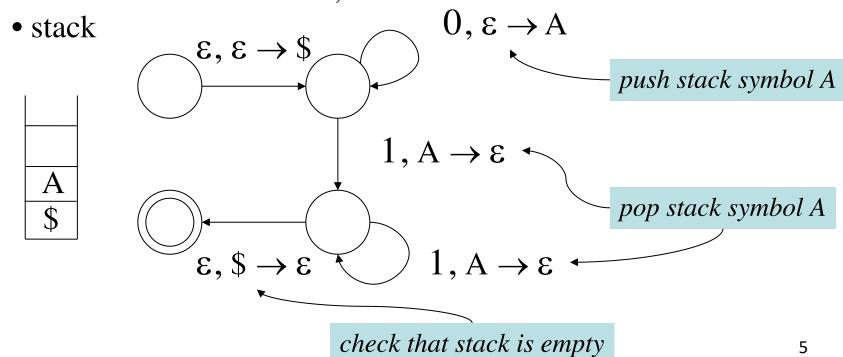
Pushdown Automata

- PDA has
 - an alphabet (terminals),
 - stack symbols (like non-terminals and terminals),
 - a finite-state automaton,

e.g. PDA for language

$$L = \{ 0^n 1^n : n >= 1 \}$$

→ implies a push/pop of stack symbol(s)



Non-CF Languages

$$L_1 = \{wcw \mid w \in (a|b)*\}$$
 $L_2 = \{a^nb^mc^nd^m \mid n \ge 1, m \ge 1\}$
 $L_3 = \{a^nb^nc^n \mid n \ge 0\}$

CF Languages

$$egin{aligned} L_4 &= \{wcw^R \mid w \in (a|b)*\} \ S &\to aSa \mid bSb \mid c \ L_5 &= \{a^nb^mc^md^n \mid n \geq 1, m \geq 1\} \ S &\to aSd \mid aAd \ A &\to bAc \mid bc \end{aligned}$$

Summary

- CFGs can be used describe PL
- Derivations correspond to parse trees
- Parse trees represent structure of programs
- Ambiguous CFGs exist
- Some forms of ambiguity can be fixed by changing the grammar
- CF languages can be recognized using Pushdown Automata