

CMPT 379

Compilers

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<http://anoopsarkar.github.io/compilers-class/>

Building a compiler

- Programming languages have a lot in common
- Do not write a compiler for each language
- Create a general mathematical model for the **structure** of all languages
- Implement a compiler using this model

Building a compiler

- Each language compiler is built using this general model (so-called compiler compilers)
 - yacc = yet another compiler compiler
- Code optimization ideas can also be shared across languages

Demo: compiler for the expr language

Building a compiler

- The cost of compiling and executing should be managed
- No program that violates the definition of the language should escape
- No program that is valid should be rejected

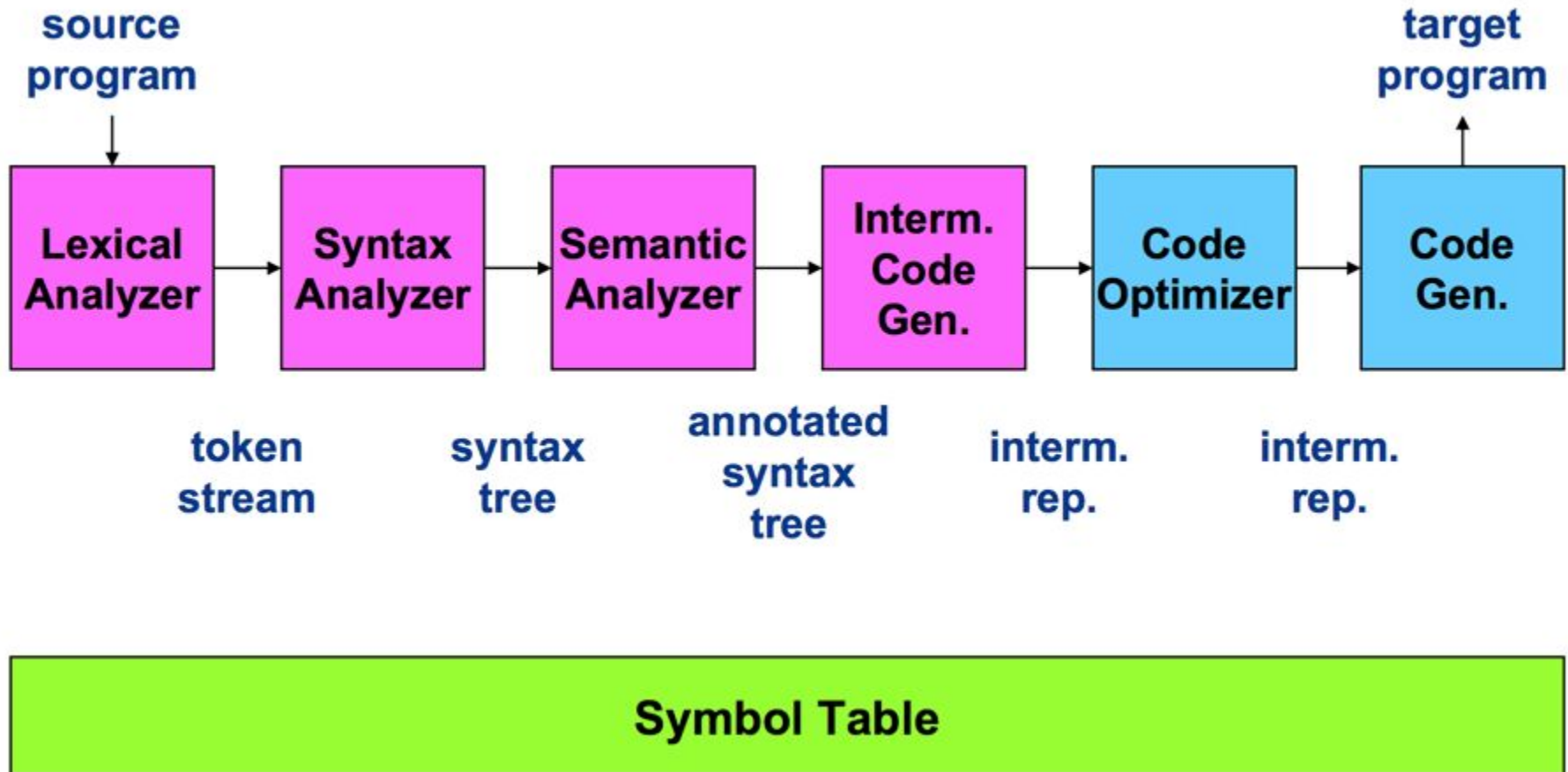
Building a compiler

- Requirements for building a compiler:
 - Symbol-table management
 - Error detection and reporting
- Stages of a compiler:
 - Analysis (front-end)
 - Synthesis (back-end)

Stages of a Compiler

- Analysis (Front-end)
 - Lexical analysis
 - Syntax analysis (parsing)
 - Semantic analysis (type-checking)
- Synthesis (Back-end)
 - Intermediate code generation
 - Code optimization
 - Code generation

Stages of a Compiler



Compiler Front-end

Lexical Analysis

- Also called *scanning*, take input program *string* and convert into tokens

- Example:

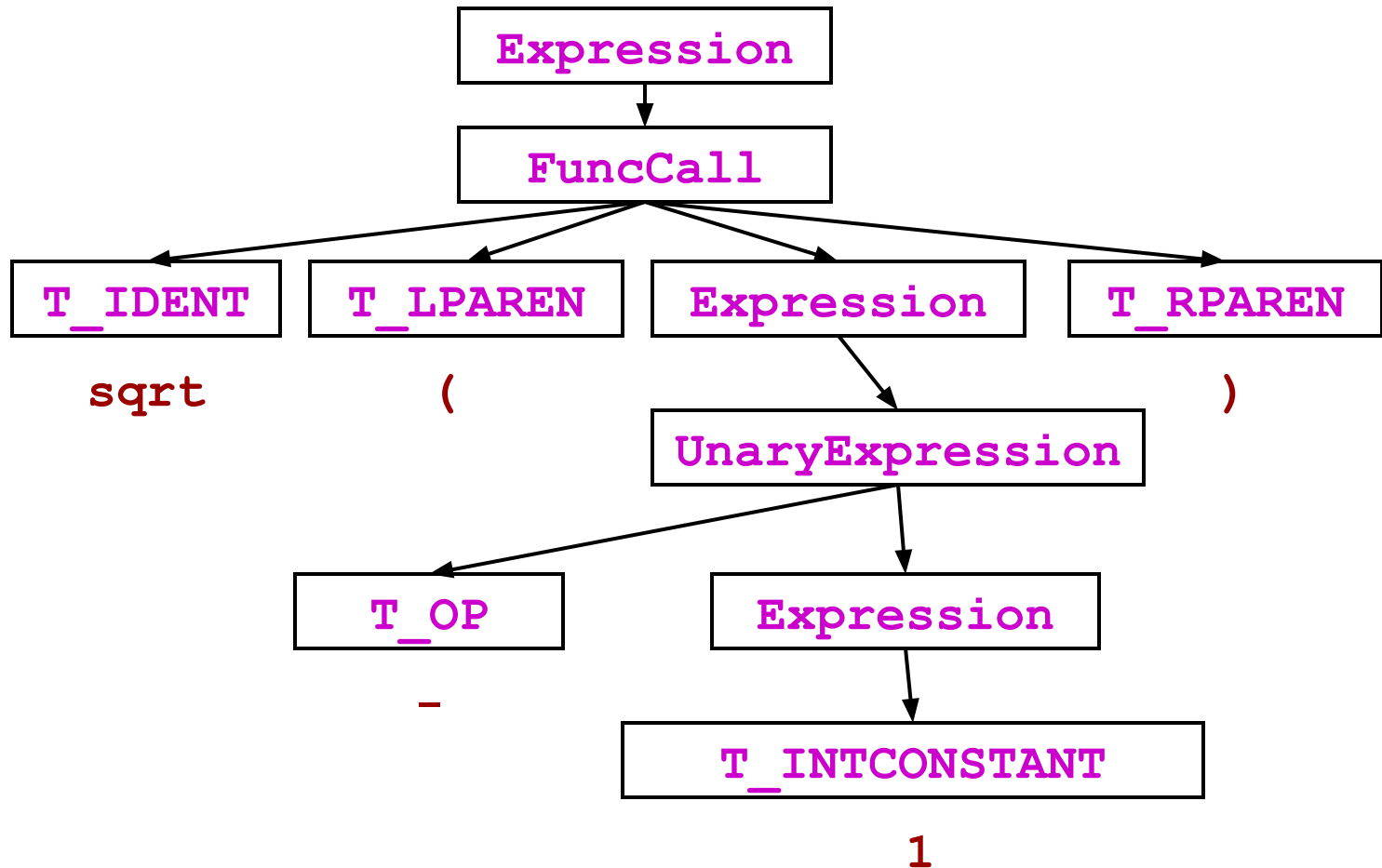
```
double f = sqrt(-1);
```

T_DOUBLE	("double")
T_IDENT	("f")
T_OP	("=")
T_IDENT	("sqrt")
T_LPAREN	("(")
T_OP	("-")
T_INTCONSTANT	("1")
T_RPAREN	(")")
T_SEP	(";")

Syntax Analysis

- Also called *parsing*
- Describe the set of strings that are programs using a grammar
- Structural validation
- Create a parse tree or derivation

Parse tree for `sqrt(-1)`



Abstract Syntax Tree

`sqrt (-1) :=`

`MethodCall (`

`sqrt,`

`UnaryExpr(UnaryMinus,`

`Number(1)`

`)`

`)`

Semantic analysis

- “does it make sense”? Checking semantic rules,
 - Is there a `main` function?
 - Is variable declared?
 - Are operand types compatible? (coercion)
 - Do function arguments match function declarations?
- Type checking
- Static vs. run-time semantic checks
 - Array bounds, return values do not match definition

Compiler Back-end

Source -> abstract syntax tree

```
extern void print_int(int);

class C {
    bool foo() { return(true); }
    int main() {
        if (foo()) {
            print_int(1); }
        }
}
```


Source -> abstract syntax tree

```
Program(  
  ExternFunction(print_int, VoidType, VarDef(IntType)),  
  Class(C,  
    None,  
    Method(foo,  
      BoolType,  
      None,  
      MethodBlock( None,  
                    ReturnStmt(BoolExpr(True))) ),  
  Method(  main,  
          IntType,  
          None,  
          MethodBlock( None,  
                        IfStmt(MethodCall(foo, None),  
                               Block(None,  
                                     MethodCall  
                                     (print_int, Number(1))),  
                        None) ) ) ) )
```

Intermediate representation

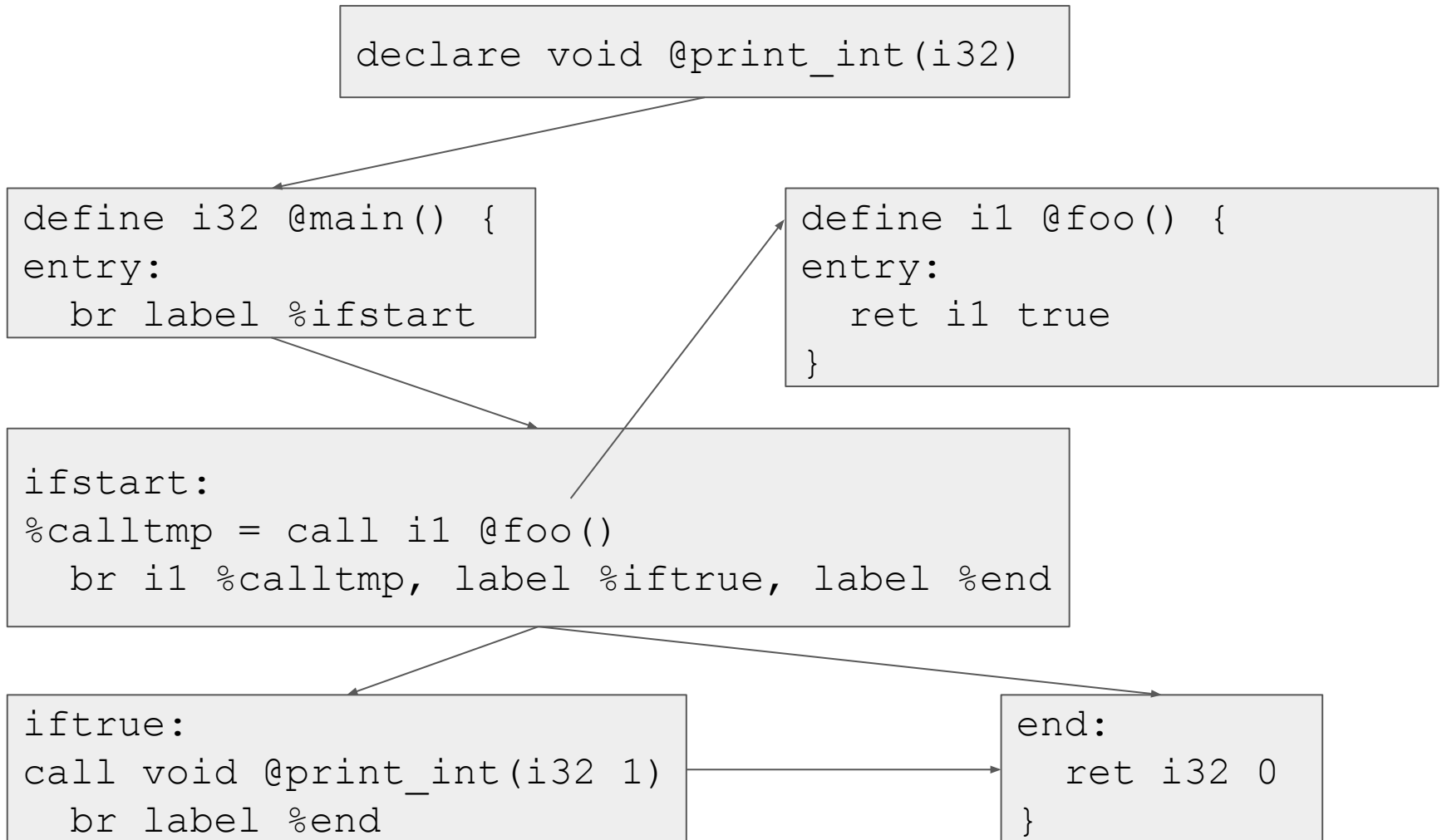
```
; ModuleID = 'C'      define i32 @main() {
                        entry:
                        br label %ifstart
                        ifstart:
                        %calltmp = call i1 @foo()
                        br i1 %calltmp, label %iftrue, label %end
                        iftrue:
                        call void @print_int(i32 1)
                        br label %end
                        end:
                        ret i32 0
                        }

declare void
@print_int(i32)

define i1 @foo() {
entry:
    ret i1 true
}
```

Translation from IR to machine specific assembly

Intermediate representation



Assembly language output from IR

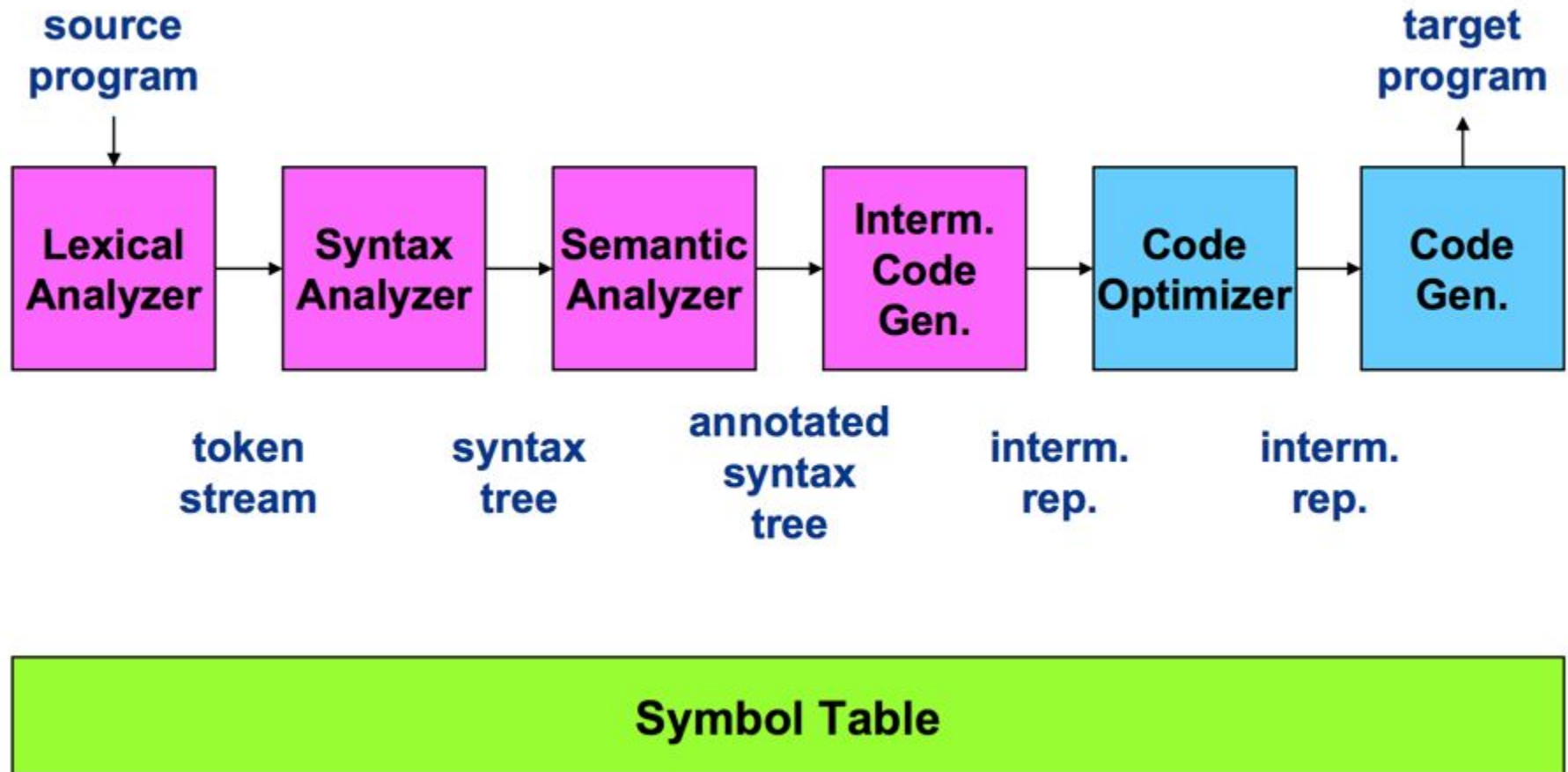
x86
assembly

```
.section __TEXT,__text,
regular,pure_instructions
.globl _foo
.align 4, 0x90
@foo
.cfi_startproc
%entry
    movl, 1
    ret
.cfi_endproc

.globl _main
.align 4, 0x90
```

```
@main
.cfi_startproc
%entry
    push    rax
Ltmp0:
.cfi_def_cfa_offset 16
    call _foo
    test al, 1
    je     LBB1_2
%iftrue
    movedi, 1
    call _print_int
%end
    xor    eax, eax
    pop    rdx
    ret
.cfi_endproc
```

Stages of a Compiler



Wrap Up

- Analysis/Synthesis
 - Translation from string to executable
- Divide and conquer
 - Build one component at a time
 - Theoretical analysis will ensure we keep things **simple** and **correct**
 - Create a complex piece of software