§1 IMO

Problem 1.1 (IMO 2019 P1)

Find all $f: \mathbb{Z} \to \mathbb{Z}$ such that,

$$f(2x) + 2f(y) = f(f(x+y))$$

Proof. Let P(x, y) the condition generated by x, y. Then,

$$\begin{cases} P(x,0) : f(2x) + 2f(0) = f(f(x)) \\ P(0,x) : f(0) + 2f(x) = f(f(x)) \end{cases}$$

Combining these two we obtain that,

$$f(2x) = 2f(x) - f(0)$$

Now let us consider the following,

$$P(x,1): f(2x) + 2f(1) = f(f(x+1))$$

However from the previous observations we can say that,

$$f(f(x+1)) = 2f(x+1) + f(0)$$

Thus,

$$2f(x) + 2f(1) - f(0) = 2f(x+1) + f(0)$$
$$f(x+1) = f(x) + f(1) - f(0)$$

Thus, we see that $f(x) = f(0) + x \cdot C$ where C is some constant (specifically f(1) - f(0)) for all $x \in \mathbb{Z}$. Now let us consider,

$$P(x,-x): f(2x) + 2f(-x) = f(f(0))$$

We can rewrite this as,

$$f(0) + 2xC + 2f(0) - 2xC = f(0) + f(0) \cdot C$$

 $\implies 2f(0) = f(0) \cdot C$

This implies that either f(0) = 0 or C = 2. Let us consider both cases,

- If f(0) = 0, then, f(x) = cx and it is simple to verify that then f(x) must equal 2x or 0.
- If C = 2, then,

$$f(0) + 4x + 2(f(0) + 2y) = f(f(0) + 2(x+y)) = f(0) + 2f(0) + 4(x+y)$$
$$4x + 2f(0) + 4y = 2f(0) + 4x + 4y$$

Thus, we see the only solutions to this functional equation are f(x) = 2x + c and f(x) = 0.

Problem 1.2 (IMO 2022 P2)

Find all $f: \mathbb{R}^+ \to \mathbb{R}^+$ such that for each $x \in \mathbb{R}^+$ there exists exactly one $y \in \mathbb{R}^+$ such that,

$$xf(y) + yf(x) \le 2$$

Proof. Obviously $f(x) = \frac{1}{x}$ works (from AM-GM). Let us say that $x \sim y$ if and only if the condition from the problem statement is true. Then obviously this is commutative, meaning that if $x \sim y$ then $y \sim x$.

Lemma 1.1 If $f(x) \leq \frac{1}{x}$, then, $x \sim x$.

This is simple to see due to,

$$2xf(x) \le 2x \cdot \frac{1}{x} = 2$$

However, notice that if $x \sim y$, then,

$$xf(y) + yf(x) \le 2$$

and in it must be that $f(x) > \frac{1}{x}$ and $f(y) > \frac{1}{y}$ which would imply that by AM-GM,

$$xf(y) + yf(x) = \frac{x}{y} + \frac{y}{x} \ge 2$$

Thus we obtain that for all $x \in \mathbb{R}^+$ it must be that $x \sim x$. Thus that,

$$f(x) \le \frac{1}{x}$$

for all $x \in \mathbb{R}^+$.

Notice that if $f(a) < \frac{1}{a}$, then, if,

$$f(a) = \frac{1}{a} - \epsilon$$

then,

$$xf(a) + af(x) = x\left(\frac{1}{a} - \epsilon\right) + af(x) \le x\left(\frac{1}{a} - \epsilon\right) + \frac{a}{x} \le 2$$

However the last inequality for $\epsilon > 0$ has multiple solutions for x. Because,

$$x^2 \left(\frac{1}{a} - \epsilon\right) + a - 2x \le 0$$

The only way for this inequality to have a single solution for $x \in \mathbb{R}^+$ is if the equality has only one solution, thus the discriminant is zero,

$$\Delta = 4 - 4 \cdot a \cdot \left(\frac{1}{a} - \epsilon\right)$$

which is obviously never 0 for $\epsilon > 0$. Thus because for each x there is only one such y that $x \sim y$ it must mean that,

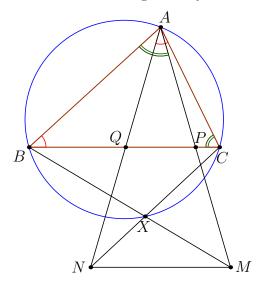
$$f(x) = \frac{1}{x}$$

for all $x \in \mathbb{R}^+$.

Problem 1.3 (IMO 2014 P4)

Let P and Q be on segment BC of an acute triangle $\triangle ABC$ such that $\angle PAB = \angle BCA$ and $\angle CAQ = \angle ABC$. Let M and N be the points on AP and AQ, respectively, such that P is the midpoint of AM and Q is the midpoint of AN. Prove that the intersection of BM and CN is on the circumference of triangle $\triangle ABC$.

Proof. Let $\triangle ABC$ be the reference triangle in barycentric coordinates.



Then we can calculate the point P, because $\triangle ABP \sim \triangle CBA$ and $\triangle ACQ \sim \triangle BCA$. Thus,

$$\frac{BA}{BC} = \frac{BP}{BA}$$

$$\implies BP = \frac{BA^2}{BC} = \frac{c^2}{a}$$

consequently, $P=\left(0:a-\frac{c^2}{a}:\frac{c^2}{a}\right)$, analogously we get, $Q=\left(0:\frac{b^2}{a}:a-\frac{b^2}{a}\right)$. Now we can calculate points M and N,

$$M = (-a^2 : 2a^2 - 2c^2 : 2c^2)$$
$$N = (-a^2 : 2b^2 : 2a^2 - 2b^2)$$

Now we simply need to intersect BM and CN which is not hard to do considering that both are cevians in the triangle. It is not hard to see that the following point satisfies these conditions,

$$X = \left(-a^2 : 2b^2 : 2c^2\right)$$

Now we just need to check that this point lies on the circumcircle of $\triangle ABC$ which is given by,

$$a^2yz + b^2xz + c^2xy = 0$$

Thus,

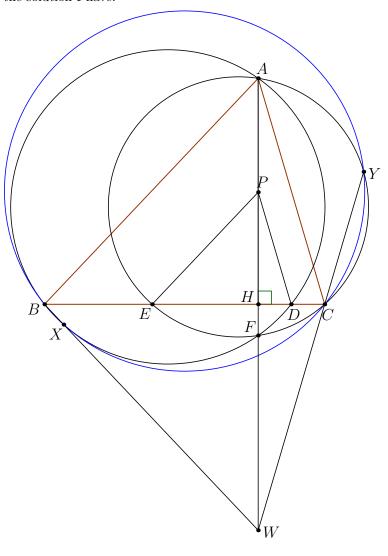
$$a^2yz + b^2xz + c^2xy = 4a^2b^2c^2 - 2a^2b^2c^2 - 2a^2b^2c^2 = 0$$

Consequently, it must be that $X \in (\triangle ABC)$.

Problem 1.4 (IMO 2022 Shortlist, G2)

In the acute-angled triangle $\triangle ABC$, the point F is the foot of the altitude from A, and P is a point on the segment AF. The lines through P parallel to AC and AB meet BC at D and E, respectively. Points $X \neq A$ and $Y \neq A$ lie on the circles ABD and ACE, respectively, such that DA = DX and EA = EY. Prove that B, C, X and Y are concyclic.

Proof. There are a lot of observations that can be done in this problem, however this is the solution I have.



Lemma 1.2 The second intersection of (ABD) and (AEC) lies on AH.

This is true due to basic angle chase, assume that F lies on AH and (ABD), let us prove that then AECF is cyclic. By Power of the Point,

$$HA \cdot HF = HD \cdot HD = HE \cdot HC$$

the last is obviously true because $PE \parallel AB$ and $PD \parallel AC$. Now let us intersect BX with AF at a point W, then the condition that BXCY is cyclic by the radical center theorem is equivelent to showing that W, C and Y are colinear.

Lemma 1.3 BC is the angle bisector of $\angle ABW$

This is because,

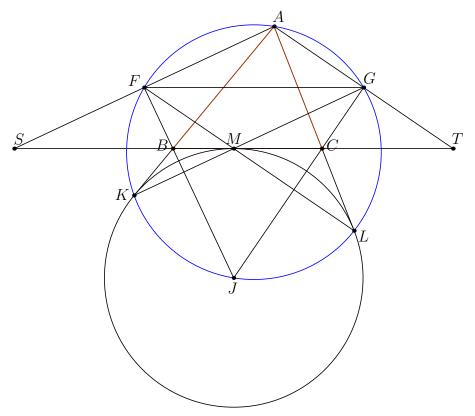
$$\angle CBW = \angle XAD = \angle AXD = \angle ABC$$

because F is the midpoint of the arc between A and X per definition. Thus, W is the reflection of A over H. The exact same argument for C shows that BC is the angle bisector of $\angle YCA$ and thus, CY must intersect AF at W.

Problem 1.5 (IMO 2012 P1)

Given triangle ABC the point J is the center of the excircle opposite the vertex A. This excircle is tangent to side BC at M, and to lines AB and AC at K and L, respectively. Lines LM and BJ meet at F, and lines KM and CJ meet at G. Let S be the point of intersection of lines AF and BC, and let T be the point of intersection of lines AG and BC. Prove that M is the midpoint of ST.

Proof. Pascal's theorem hints us towards checking whether FAGKJL is cyclic.



Notice, that because $KM \perp FJ$ and $JG \perp ML$, it must be that M is the orthocenter of FGJ. Thus, $FG \parallel BC$, consequently,

$$\angle GFJ = \angle BCJ = \angle MKJ$$

Thus, FKJG is cyclic, analogously we conclude that FKJLG is cyclic and by Pascal's theorem it must be that A lies on this circle as well, thus FAGLJK is cyclic.

Now notice,

$$\angle FGA = \angle ALF = \angle CML = \angle GFL$$

the last step comes from $FG \parallel BC$. Thus, $AT \parallel FM$, the same for $AG \parallel FM$. Consequently, AGMF is a parallelogram. Because of that it must be that FG is the midline of $\triangle AST$ and M is the midpoint of the side ST.

Problem 1.6 (IMO 2022 Shortlist C1)

A ± 1 -sequence is a sequence of 2022 numbers a_1, \ldots, a_{2022} , each equal to either +1 or -1. Determine the largest C such that, for any ± 1 -sequence, there exists an integer k and indices $1 \leq t_1 \leq \ldots \leq t_k \leq 2022$ so that $t_{i+1} - t_i \leq 2$ for all i, and

$$\left| \sum_{i=1}^{k} a_{t_i} \right| \ge C$$

Proof. WLOG, let there be more 1 than -1 in the ± 1 -sequence. Then, let us consider the following algorithm,

Algorithm 1.4 If the next element is -1, then skip to the next one and include it. If the next element is 1, then include it. Continue this process until you have went through the entire sequence.

This algorithm will garantee that we skip at least $\lceil \frac{X}{2} \rceil$, where X is the number of -1 in the sequence. Thus,

$$\left| \sum_{i=1}^{k} a_{t_i} \right| \ge (2022 - X) - \left\lfloor \frac{X}{2} \right\rfloor \ge (2022 - 1011) - \left\lfloor \frac{1011}{2} \right\rfloor = 506$$

Consequently, it must be that $C \geq 506$. Now consider the following sequence,

$$1, 1, -1, -1, 1, 1, \ldots, -1, -1, 1, -1$$

It alternates between 1, 1 and -1, -1 until the very end, where it is 1 and -1. Notice that in each consequtive pair of identical numbers at least one of the numbers is present in our sequence. Thus, no matter our choice of indices, the total sum is bounded by,

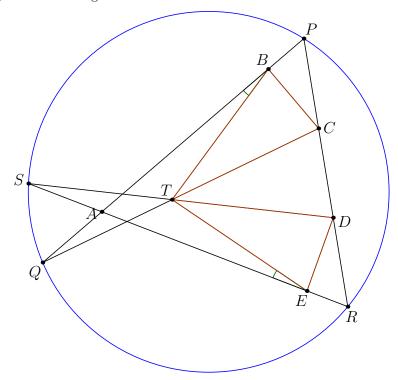
$$\left| \sum_{i=1}^{k} a_{t_i} \right| \le 2 \cdot \frac{1010}{2} - \frac{1010}{2} + 1 = 2 \cdot 505 - 505 + 1 = 506$$

Thus, it must be that C = 506.

Problem 1.7 (IMO 2022 P4)

Let ABCDE be a convex pentagon such that BC = DE. Assume that there is a point T inside ABCDE with TB = TD, TC = TE and $\angle ABT = \angle TEA$. Let line AB intersect lines CD and CT at points P and Q, respectively. Assume that the points P, B, A, Q occur on their line in that order. Let line AE intersect lines CD and DT at points R and S, respectively. Assume that the points R, E, A, S occur on their line in that order. Prove that the points P, P, P, P, P, P0 lie on a circle.

Proof. The first thing which one can notice is that $\triangle TDE \sim \triangle TBC$.



Now, notice that,

Lemma 1.5 $\triangle STE \sim \triangle QTB$ and SQCD is cyclic.

This is because,

$$\angle STE = 180 - \angle DTE = 180 - \angle BTC = \angle QTB$$

Thus,

$$\begin{split} \frac{TS}{TC} &= \frac{TS}{TE} = \frac{TQ}{TB} = \frac{TQ}{TD} \\ \Longrightarrow & TS \cdot TD = TC \cdot TQ \end{split}$$

Consequently SQCD is cyclic. Now,

$$\angle QSR = \angle QSD - \angle RSD = \angle QCD - \angle PQC = \angle QPR$$

which proves that SQPR is cyclic.

Problem 1.8 (IMO 2012 P4)

Find all functions $f: \mathbb{Z} \to \mathbb{Z}$ such that, for all integers a, b, c such that a+b+c=0, the following equality holds,

$$f^{2}(a) + f^{2}(b) + f^{2}(c) = 2f(a)f(b) + 2f(b)f(c) + 2f(a)f(c)$$

Proof. Obviously from a = b = c = 0 we obtain that f(0) = 0.

Notice, if we set into the original functional equation b = -a and c = 0, we will obtain,

$$f^{2}(a) + f^{2}(-a) = 2f(a)f(-a)$$

$$\implies (f(a) - f(-a))^{2} = 0$$

$$\implies f(a) = f(-a)$$

Now, notice that a very common transformation with this type of equation is,

$$x^{2} + y^{2} + z^{2} = 2xy + 2yz + 2xz$$

$$(x+y-z)^{2} = x^{2} + y^{2} + z^{2} + 2xy - 2xz - 2xy$$

$$\implies 4xy = (x+y-z)^{2}$$

$$\implies xy = \left(\frac{x+y-z}{2}\right)^{2}$$

Thus,

$$f(a)f(b) = \left(\frac{f(a) + f(b) - f(c)}{2}\right)^2$$

Now is the perfect time to get rid of the third variable, let c = -a - b and we will obtain, (because f(-a - b) = f(a + b))

$$f(a)f(b) = \left(\frac{f(a) + f(b) - f(a+b)}{2}\right)^2$$

$$f(a+b) = f(a) + f(b) \pm 2\sqrt{f(a)f(b)}$$

The nice thing about this is first of all, because f(a)f(b) is always nonnegative it must be that f is either always ≥ 0 or ≤ 0 . The second thing is that now we obtained an equation equivelent to the original functional equation and we can assume that f is positive (because if f is a solution, then -f is a solution as well). However, we can still do some preliminary steps for example obviously we want to rewrite the equation as,

$$f(a+b) = (\sqrt{f(a)} \pm \sqrt{f(b)})^{2}$$

$$\implies \sqrt{f(a+b)} = \pm \sqrt{f(a)} \pm \sqrt{f(b)}$$

thus if $g(x) = \sqrt{f(x)}$, then we know that,

$$g(x+y) = \pm g(x) \pm g(y)$$

Now, if g(1) = c, then,

$$g(2) = \pm g(1) \pm g(1) = c \pm c = 2c \vee 0$$

Now, let us consider two cases,

Case 1) If g(2) = 2c, then,

$$g(3) = \pm f(1) \pm f(2) = \pm c \pm 2c = 3c \lor c$$

because f is ≥ 0 . Now, we again consider two cases, right now let g(3) = 3c. Then,

$$\begin{cases} g(4) = \pm g(2) \pm g(2) = 4c \lor 0 \\ g(4) = \pm g(1) \pm g(3) = 4c \lor 2c \end{cases}$$
$$\implies g(4) = 4c$$

inductively it is simple to continue this logic and show that g(n) = nc. Thus, one of the solutions for f is $f(n) = cn^2$.

Case 2) If g(2) = 2c and g(3) = c. Then,

$$\begin{cases} g(4) = \pm g(1) \pm g(3) = 0 \lor 2c \\ g(4) = \pm g(2) \pm g(2) = \pm 2c \pm 2c = 0 \lor 4c \end{cases}$$

$$\implies a(4) = 0$$

However, if you repeat this argument for higher value of g(n) you will obtain a periodic function of the following form,

$$g(n) = \begin{cases} c, & \text{if } n \equiv 1 \pmod{2} \\ 2c, & \text{if } n \equiv 2 \pmod{4} \\ 0, & \text{if } n \equiv 0 \pmod{4} \end{cases}$$

which is equivelent to,

$$f(n) = \begin{cases} c, & \text{if } n \equiv 1 \pmod{2} \\ 4c, & \text{if } n \equiv 2 \pmod{4} \\ 0, & \text{if } n \equiv 0 \pmod{4} \end{cases}$$

Case 3) If g(2) = 0, then,

$$\begin{cases} g(3) = \pm g(1) \pm g(2) = c \\ g(4) = \pm g(2) \pm g(2) = 0 \\ \dots \end{cases}$$

Thus,

$$g(n) = \begin{cases} 0, & \text{if } n \equiv 0 \pmod{2} \\ c, & \text{if } n \equiv 0 \pmod{2} \end{cases}$$

$$\implies f(n) = \begin{cases} 0, & \text{if } n \equiv 0 \pmod{2} \\ c, & \text{if } n \equiv 0 \pmod{2} \end{cases}$$

Thus, we have obtained all solutions (all the solutions above have there negative counterparts as well). $\hfill\Box$

Problem 1.9 (IMO 2013 P2)

A configuration of 4027 points in the plane is called Colombian if it consists of 2013 red points and 2014 blue points, and no three of the points of the configuration are collinear. By drawing some lines, the plane is divided into several regions. An arrangement of lines is good for a Colombian configuration if the following two conditions are satisfied:

- 1. No line passes through any point of the configuration.
- 2. No region contains points of both colors.

Find the least value of k such that for any Colombian configuration of 4027 points, there is a good arrangement of k lines.

Proof. Let us consider the following construction (alternating the colors), I claim one needs at least 2013 lines.



To prove this simply consider all the segments connecting neighboring blue/red points, then each of those segments must be split by some line (else they wouldn't be in different regions), but each line cuts no more than 2 such segments. However, there is a total of $2 \cdot 2013$ such segments, thus a minimum of 2013 lines is required.

Now, we need to prove that 2013 lines is always enough.

Lemma 1.6 Given n points blue and k < n red points, where k is even, then it is possible to satisfy the conditions in the problem statement using k lines.

Consider two red points, then no blue point lies on the line between them. Thus, we can consider two close enough lines parallel to the line between the two red points and divide the plane into three parts with one of the parts containing both of the red points and no other point. Thus, if we repeat this argument for

the $\frac{k}{2}$ pairs of red points we will obtain a construction satisfying the conditions of the problem statement involving only $2 \cdot \frac{k}{2} = k$ lines.

Lemma 1.7 Given n points of one color and n-1 points of another color, a minimum of n-1 lines is required to satisfy the conditions in the problem statement.

If n-1 is even, then we are done, thus let us assume that n-1 is odd.

Let us construct a convex hull given all the points. If it contains two neighboring blue points, then we separate them from the rest of the points using one line and then apply the induction hypothesis on the rest of the points and obtain 1 + n - 2 = n - 1 lines.

If at least one red point is present in the convex hull we can separate it from the rest of the points using one line. Then, we have n-2 red points left, which is an even number, thus we can satisfy the conditions of the problem statement using 1+n-2=n-1 lines.

Obviously the problem follows when n = 2014.

Problem 1.10 (IMO Shortlist 2022 N2)

Find all positive integers n > 2 such that,

$$n! \mid \prod_{p < q \le n} (p + q)$$

Proof. Let r be the biggest prime $\leq n$. Then, obviously $v_r(n!) = 1$. However,

$$v_p\left(\prod_{p< q\le n}(p+q)\right) = \sum_{p+q=r}1$$

because p+q < r+r = 2r. Thus, there must be exactly one solution to p+q = r for prime p, q. However, because r is an odd prime it must be that either p or q is even, i.e. 2. WLOG q = r - 2.

Notice, that it can't be that p, p+2 and p+4 are all prime, except for 3, 5, 7. Let us use the exact same logic for q, assuming that q > 5, which is true for $n \ge 11$. Notice, that,

$$x + y$$

where x, y are prime (due to the above). However, notice that it also can't be that x + y = q where x, y are prime, also due to the argument before. Thus,

$$v_q \left(\prod_{x < y \le n} (x + y) \right) = 0 < v_q(n!)$$

contradiction! Thus, $7 \le n < 11$,

Case 1) If n = 7, (primes are 2, 3, 5, 7)

$$\prod_{p < q \le n} (p+q) = (2+3) \cdot (3+5) \cdot (5+7) \cdot (2+5) \cdot (2+7) \cdot (3+7)$$

$$= 5 \cdot 8 \cdot 12 \cdot 7 \cdot 9 \cdot 10 \quad (1)$$

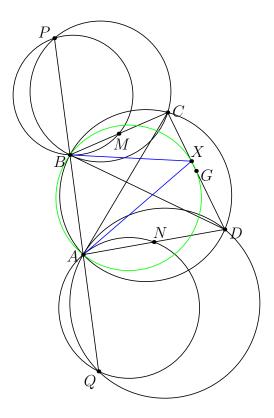
which is obviously divisible by 7!.

Case 2) However, already for n > 7, there isn't enough factors in the product (because no new primes appear up until 11).

Thus, the only n satisfying the problem statement is 7.

Problem 1.11 (IMO Shortlist 2022 G3)

Let ABCD be a cyclic quadrilateral. Assume that the points Q, A, B, P are collinear in this order, in such a way that the line AC is tangent to the circle ADQ, and the line BD is tangent to the circle BCP. Let M and N be the midpoints of segments BC and AD, respectively. Prove that the following three lines are concurrent: line CD, the tangent of circle ANQ at point A, and the tangent to circle BMP at point B.



Proof. Notice that,

$$\angle QDA = \angle CAB = \angle BDC$$

and,

$$\angle QAD = \angle BCD$$

Thus, $\triangle ADQ \sim \triangle CDB$, and D is the center of the spiral similarity between the triangles.

Let G be the midpoint of CD. Notice, under the spiral similarity N goes to G. Thus,

$$\angle GBC = \angle NQA = \angle DAX$$

where X is the intersection of the tangent of (ANQ) from A with CD. Analogously we obtain that $\angle DAG = \angle X'BC$.

However, this means that $\angle GBX' = \angle GAX$, thus X and X' are the same point, the intersection of (BAG) with CD.

Problem 1.12 (IMO 2022 P1)

The Bank of Oslo issues two types of coin: aluminum (denoted A) and bronze (denoted B). Marianne has n aluminum coins and n bronze coins arranged in a row in some arbitrary initial order. A chain is any subsequence of consecutive coins of the same type. Given a fixed positive integer $k \leq 2n$, Gilberty repeatedly performs the following operation: he identifies the longest chain containing the k^{th} coin from the left and moves all coins in that chain to the left end of the row. For example, if n=4 and k=4, the process starting from the ordering AABBBABA would be $AABBBABA \rightarrow BBBAAABA \rightarrow AAABBBBAA \rightarrow BBBBAAAAA \rightarrow ...$

Find all pairs (n, k) with $1 \le k \le 2n$ such that for every initial ordering, at some point of the process there will be at most one aluminium coin adjacent to a copper coin.

Proof. Obviously it makes sense to split the row into consecutive blocks of coins of the same type. Then, the fact that at most one aluminium coin is adjacent to a copper coin is equivelent to there being ≤ 2 blocks. Notice,

Lemma 1.8 The number of blocks decreases or remains constant after every operation.

Thus, in order to not obtain a situation with ≤ 2 blocks, it must be that the number of blocks doesn't change. However, the number of blocks doesn't change if and only if, either the pointer is on the first block, or that the pointer is on the last block. Thus, one of these two conditions must be satisfied after a certain point in time for infinitely many operations. Thus, a situation stabalizes (i.e. doesn't change the number of blocks) if and only if the size of each block is less than the distance from the pointer to the right end of the row, or that the pointer points to the first block.

Case 1) Notice, that for $n > \lceil \frac{3n}{2} \rceil$ we can consider four blocks of alternating coin types with the sizes $\lfloor \frac{n}{2} \rfloor, \lfloor \frac{n}{2} \rfloor, \lceil \frac{n}{2} \rceil, \lceil \frac{n}{2} \rceil$, then the blocks will simply rotate and there will never be 2 blocks.

Case 2) If k < n, then the configuration of where the first n-1 coins are of one time, and then the rest is randomly distributed will never change, thus will never achieve ≤ 2 blocks.

Case 3) If $\lceil \frac{3n}{2} \rceil \geq k \geq n$, then if there are ≥ 4 blocks the situation won't be stable (due to the above). However, a configuration with an odd number of blocks is never stable, because at some point blocks of the same type will merge. Thus, the situation will only be stable when the number of blocks is 2, which is what is required.

Thus,
$$n \le k \le \lceil \frac{3n}{2} \rceil$$
.

§2 EGMO

Problem 2.1 (EGMO 2025 P1)

For a positive integer N, let $c_1 < c_2 < \cdots < c_m$ be all positive integers smaller than N that are coprime to N. Find all $N \ge 3$ such that

$$\gcd(N, c_i + c_{i+1}) \neq 1$$

for all $1 \leq i \leq m-1$

Here gcd(a, b) is the largest positive integer that divides both a and b. Integers a and b are coprime if gcd(a, b) = 1.

Proof. Notice, that if N is even, then all c_i are odd, however that implies that $c_i + c_{i+1}$ will always be even, thus never relatively prime with N. Consequently, every even N satisfies the conditions of the problem statement. If N is odd, then it is by definition relatively prime to 2, thus,

$$\gcd(N, 1+2) \neq 1 \implies 3 \mid N$$

Thus, $N = 3^{\alpha}x$ (where x is odd), however,

Case 1) If $x \equiv 1 \pmod{3}$, then let us consider two numbers, x + 1 and x + 3, then,

$$\gcd(3^{\alpha}x, x+1) = 1$$

$$\gcd(3^{\alpha}x, x-2) = 1$$

But, then,

$$(3^{\alpha}x, 2x - 1) = (3^{\alpha}, 2x - 1) = 1$$

contradiction!

Case 2) If $x \equiv 2 \pmod{3}$, then consider x + 2 and x - 1, then,

$$\begin{cases} \gcd(3^{\alpha}x, x+2) = 1\\ \gcd(3^{\alpha}x, x-1) = 1 \end{cases}$$

then,

$$gcd(3^{\alpha}x, 2x + 1) = gcd(3^{\alpha}, 2x + 1) = 1$$

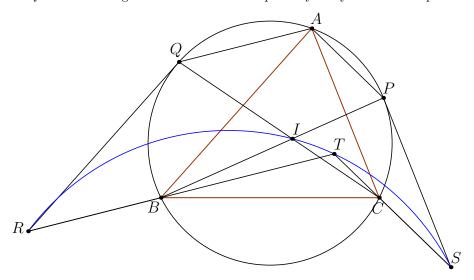
contradiction!

Thus x=1. Consequently, the only solutions are all even numbers and all powers of three.

Problem 2.2 (EGMO 2025 P4)

Let ABC be an acute triangle with incentr I and $AB \neq AC$. Let lines BI and CI intersect the circumcircle of ABC at $P \neq B$ and $Q \neq C$, respectively. Consider points R and S such that AQRB and ACSP are parallelograms (with $AQ \parallel RB, AB \parallel QR, AC \parallel SP$, and $AP \parallel CS$). Let T be the point of intersection of lines RB and SC. Prove that points R, S, T, and I are concyclic.

Proof. At first it might seem unclear how to prove cyclicity of these four points.



However, after playing around with some angles it is not difficult to notice,

Lemma 2.1 BITC is cyclic.

This is because,

$$\angle BTC = 180 - \angle TBC - \angle TCB = 180 - (\angle B - \angle ABT) - (\angle C - \angle TCA)$$
$$= \angle A + \angle QAB + \angle CAP = \angle A + \frac{\angle C}{2} + \frac{\angle B}{2} = 90 + \frac{\angle A}{2} = \angle BIC \quad (2)$$

Now, notice,

Lemma 2.2 $\triangle IBR \sim \triangle ICS$

Because, $(\triangle IBQ \sim \triangle ICP)$

$$\frac{BI}{BR} = \frac{IB}{QA} = \frac{IB}{BQ} = \frac{IC}{CP} = \frac{IC}{PA} = \frac{IC}{CS}$$

and $\angle RBI = \angle ICS$ since $\angle IBT = \angle TCI$ from ITBC being cyclic. But this similarity implies that $\angle BIR = \angle CIS$, thus,

$$\angle RIS = \angle BIC + \angle RIB - \angle ICS = \angle BIC = \angle RTS$$

Consequently, RITS is cyclic.

Problem 2.3 (EGMO 2025 P5)

Let n > 1 be an integer. In a configuration of an $n \times n$ board, each of the n^2 cells contains an arrow, either pointing up, down, left, or right. Given a starting configuration, Turbo the snail starts in one of the cells of the board and travels from cell to cell. In each move, Turbo moves one square unit in the direction indicated by the arrow in her cell (possibly leaving the board). After each move, the arrows in all of the cells rotate 90° counterclockwise. We call a cell good if, starting from that cell, Turbo visits each cell of the board exactly once, without leaving the board, and returns to her initial cell at the end. Determine, in terms of n, the maximum number of good cells over all possible starting configurations.

Proof. Notice, if n is odd, then n^2 is odd and thus it is impossible to make a cycle (due to parity).

Let us consider the very edges of the board, the only way to go through them is using one of the two following algorithms,

	\rightarrow	\leftarrow
		7
. · ·		:



where by \nearrow I mean either \rightarrow or \uparrow . Thus, due to the cycles being uniquely determined after the turn, it must be due to there being only two ways of going through the corners, only two global cycles. However, notice that the two configurations above are not equivelent under the rotation operation, thus there cannot be more than one global cycle. However, given a cycle there are only $\frac{n^2}{4}$ generators (every fourth element in the cycle).

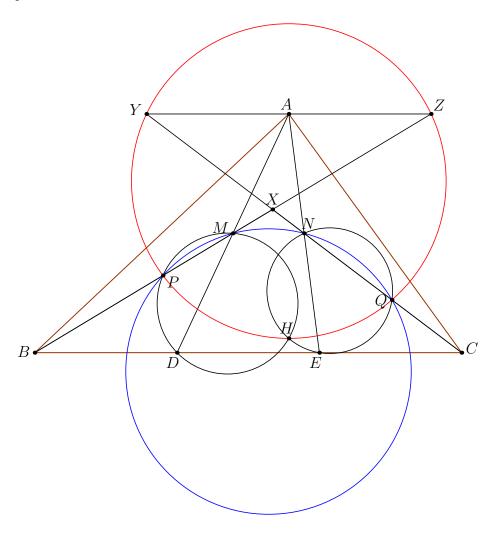
To achieve this bound, simply consider any global cycle through the entire board and appropriatly adjust the arrows so that it is a cycle from a given square. Then, due to the entire board taking the original position after four rotations, it must be that every fourth element in the cycle is a generator as well.

Thus, for all odd n there are 0 good cells, however for even n there are $\frac{n^2}{4}$ good cells.

Problem 2.4 (EGMO 2025 P3)

Let ABC be an acute triangle. Points B, D, E, and C lie on a line in this order and satisfy BD = DE = EC. Let M and N be the midpoints of AD and AE, respectively. Suppose triangle ADE is acute, and let H be its orthocentre. Points P and Q lie on lines BM and CN, respectively, such that D, H, M, and P are concyclic and pairwise different, and E, H, N, and P are concyclic and pairwise different. Prove that P, Q, N, and P are concyclic.

Proof. Let us reflect C and B over N and M, respectively, then we will obtain points Y and Z. Let X be the intersect of BM and CN.



Notice,

Lemma 2.3 AZHE is cyclic.

This is because,

$$\angle AHE = \angle EDA = \angle AZE$$

Now, notice that,

$$\angle AZH = \angle AEH = \angle YQH$$

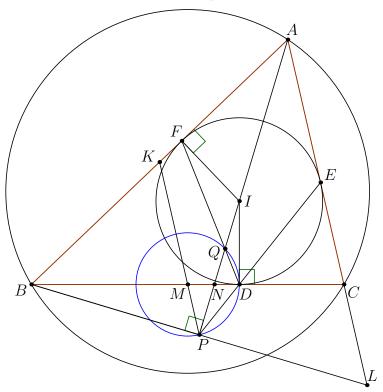
thus YZHQ is cyclic, analogously it must be that YZPH is cyclic, consequently YZPQH is cyclic. However, because $MN\parallel YZ$ by Reim's lemma it must be that MNPQ is cyclic.

§3 Korean National Olympiad

Problem 3.1 (Korea 2025 P4)

Triangle ABC satisfies $\overline{CA} > \overline{AB}$. Let the incenter of triangle ABC be ω , which touches BC, CA, AB at D, E, F, respectively. Let M be the midpoint of BC. Let the circle centered at M passing through D intersect DE, DF at $P(\neq D), Q(\neq D)$, respectively. Let line AP meet BC at N, line BP meet CA at D. Prove that the three lines EQ, FP, NL are concurrent.

Proof. Let us start by proving that A, P, Q are colinear. Let us define P as the intersection of AI and DE, then we must prove that MPD is isoseles.



This can be easily shown by angle chase, let us introduce K, the midpoint of AB it is well known that then K, M, P are colinear and $\angle BPA = 90$. Thus,

$$\angle MPD = \angle APD + \angle KPA = \angle IBC + \angle PAC = \frac{\angle A + \angle B}{2}$$
$$= \frac{180 - \angle C}{2} = \angle EDC = \angle MDP \quad (3)$$

Thus, it must be that P lies on AI, the exact same logic shows that $Q \in AI$, thus, A, I, Q and P are colinear.

Now, notice that under reflection over AI the line FQ goes to EQ and the line LN goes to BN (since $\angle APB = 90$ it must be that B goes to L under reflection) and PF goes to PE. Thus, EQ, LN and PF are concurrent, and the intersection point is the reflection of D over AI.