

DES EXAMPLE

M = 0123456789ABCDEF

Where **M** is in hexadecimal (base 16) format.

Convert the binary: Rewriting **M** in binary format, we get the 64-bit block of text:

M = 0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 1010 1011 1100 1101 1110 1111

L = 0000 0001 0010 0011 0100 0101 0110 0111

R = 1000 1001 1010 1011 1100 1101 1110 1111

K = 133457799BBCDFF1

Convert the binary: This gives us as the binary key (setting 1 = 0001, 3 = 0011, etc., and grouping together every eight bits, of which the last one in each group will be unused):

K = 00010011 00110100 01010111 01111001 10011011 10111100 11011111 11110001

Step 1: Use permutation Choice-1 (PC-1) to create the permute key (PC-1)

57	49	41	33	25	17	9
1	58	50	42	34	26	18
10	2	59	51	43	35	27
19	11	3	60	52	44	36
63	55	47	39	31	23	15
7	62	54	46	38	30	22
14	6	61	53	45	37	29
21	13	5	28	20	12	4

From the original 64-bit key

K = 0001001**1** 0011010**0** 0101011**1** 0111100**1** 1001101**1** 1011110**0** 1101111**1** 1111000**1**

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63	64
0	0	0	1	0	0	1	1	0	0	1	1	0	1	0	0	0	1	0	1	0	1	1	1	0	1	1	1	1	0	0	1	1	0	0	1	1	0	1	1	0	1	1	1	1	0	1	1	0	1	1	1	1	1	1	1	0	0	0	1				


we get the 56-bit permutation:

K_P = 1111000 0110011 0010101 0101111 0101010 1011001 1001111 0001111


Step 2:

Split this key into left and right halves, **C₀** and **D₀**, where each half has 28 bits.

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48	49	50	51	52	53	54	55	56
1	1	1	1	0	0	0	0	1	1	0	0	1	1	0	0	1	0	1	0	1	0	1	0	1	1	1	1	0	1	0	1	0	1	0	1	1	0	0	1	1	1	1	0	0	0	1	1	1	1	1					



C₀



D₀

$C_0 = 1111000\ 0110011\ 0010101\ 0101111$

$D_0 = 0101010\ 1011001\ 1001111\ 0001111$

Step 3: Rotate shift

- In **rounds $i = 1, 2, 9, 16$** , the two halves are each rotated left by **one bit**.
- In **all other rounds** where the two halves are each rotated left by **two bits**.

From original pair pair C_0 and D_0 we obtain:

$C_0 = 1111000011001100101010101111$

$D_0 = 010101011100110011110001111$

R_i and K_i	Left Part (C)	Right Part (D)
K_0 :	1111000011001100101010101111	010101011100110011110001111
K_1 :LS-1:	1110000110011001010101011111	1010101011001100111100011110
K_2 : LS-1:	1100001100110010101010111111	0101010110011001111000111101
K_3 : LS-2:	0000110011001010101011111111	0101011001100111100011110101
K_4 : LS-2:	0011001100101010101111111100	0101100110011110001111010101
K_5 : LS-2:	1100110010101010111111110000	0110011001111000111101010101
K_6 : LS-2:	0011001010101011111111000011	1001100111100011110101010101
K_7 : LS-2:	1100101010101111111100001100	0110011110001111010101010110
K_8 : LS-2:	0010101010111111110000110011	1001111000111101010101011001
K_9 : LS-1:	01010101011111111100001100110	0011110001111010101010110011
K_{10} : LS-2:	01010101111111110000110011001	1111000111101010101011001100
K_{11} : LS-2:	01010111111111000011001100101	1100011110101010101100110011
K_{12} : LS-2:	01011111111100001100110010101	0001111010101010110011001111
K_{13} : LS-2:	01111111110000110011001010101	0111101010101011001100111100

K₁₄: LS-2:	1111111000011001100101010101	1110101010101100110011110001
K₁₅: LS-2:	1111100001100110010101010111	1010101010110011001111000111
K₁₆: LS-1:	1111000011001100101010101111	0101010101100110011110001111

Step 4: **PC-2**

*In each round (R_i) permuted choice **PC-2***

We now form the keys **K_n**, for 1 ≤ **n** ≤ 16, by applying the following permutation table to each of the concatenated pairs **C_nD_n**. Each pair has 56 bits, but **PC-2** only uses 48 of these.

PC-2

14	17	11	24	1	5	3	28
15	6	21	10	23	19	12	4
26	8	16	7	27	20	13	2
41	52	31	37	47	55	30	40
51	45	33	48	44	49	39	56
34	53	46	42	50	36	29	32

K₁:LS-1:	1110000110011001010101011111	1010101011001100111100011110
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1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48	49	50	51	52	53	54	55	56
1	1	1	0	0	0	0	1	1	0	0	1	1	0	0	1	0	1	0	1	0	1	0	1	1	1	1	1	1	0	1	0	1	0	1	0	1	1	0	0	1	1	1	1	1	0	0	0	1	1	1	1	0			

The following bits are discarded

9 18 22 25 35 38 43 54

K_I = 00011011 00000010 11101111 11111100 1110000 01110010

For the other keys we have

$K_2 = 011110\ 011010\ 111011\ 011001\ 110110\ 111100\ 100111\ 100101$

$K_3 = 010101\ 011111\ 110010\ 001010\ 010000\ 101100\ 111110\ 011001$

$K_4 = 011100\ 101010\ 110111\ 010110\ 110110\ 110011\ 010100\ 011101$

$K_5 = 011111\ 001110\ 110000\ 000111\ 111010\ 110101\ 001110\ 101000$

$K_6 = 011000\ 111010\ 010100\ 111110\ 010100\ 000111\ 101100\ 101111$

$K_7 = 111011\ 001000\ 010010\ 110111\ 111101\ 100001\ 100010\ 111100$

$K_8 = 111101\ 111000\ 101000\ 111010\ 110000\ 010011\ 101111\ 111011$

$K_9 = 111000\ 001101\ 101111\ 101011\ 111011\ 011110\ 011110\ 000001$

$K_{10} = 101100\ 011111\ 001101\ 000111\ 101110\ 100100\ 011001\ 001111$

$K_{11} = 001000\ 010101\ 111111\ 010011\ 110111\ 101101\ 001110\ 000110$

$K_{12} = 011101\ 010111\ 000111\ 110101\ 100101\ 000110\ 011111\ 101001$

$K_{13} = 100101\ 111100\ 010111\ 010001\ 111110\ 101011\ 101001\ 000001$

$K_{14} = 010111\ 110100\ 001110\ 110111\ 111100\ 101110\ 011100\ 111010$

$K_{15} = 101111\ 111001\ 000110\ 001101\ 001111\ 010011\ 111100\ 001010$

$K_{16} = 110010\ 110011\ 110110\ 001011\ 000011\ 100001\ 011111\ 110101$