Big Oh

- 2-7. [3] True or False?
 - (a) Is $2^{n+1} = O(2^n)$?
 - (b) Is $2^{2n} = O(2^n)$?
- 2-8. [3] For each of the following pairs of functions, either f(n) is in O(g(n)), f(n) is in $\Omega(g(n))$, or $f(n) = \Theta(g(n))$. Determine which relationship is correct and briefly explain why.
 - (a) $f(n) = \log n^2$; $g(n) = \log n + 5$

- (b) $f(n) = \sqrt{n}$; $g(n) = \log n^2$
- (c) $f(n) = \log^2 n$; $g(n) = \log n$
- (d) f(n) = n; $g(n) = \log^2 n$
- (e) $f(n) = n \log n + n; g(n) = \log n$
- (f) f(n) = 10; $q(n) = \log 10$
- (g) $f(n) = 2^n$; $g(n) = 10n^2$
- (h) $f(n) = 2^n$; $g(n) = 3^n$
- 2-9. [3] For each of the following pairs of functions f(n) and g(n), determine whether f(n) = O(g(n)), g(n) = O(f(n)), or both.
 - (a) $f(n) = (n^2 n)/2$, g(n) = 6n
 - (b) $f(n) = n + 2\sqrt{n}, g(n) = n^2$
 - (c) $f(n) = n \log n, g(n) = n\sqrt{n}/2$
 - (d) $f(n) = n + \log n, g(n) = \sqrt{n}$
 - (e) $f(n) = 2(\log n)^2$, $g(n) = \log n + 1$
 - (f) $f(n) = 4n \log n + n$, $g(n) = (n^2 n)/2$
- 2-10. [3] Prove that $n^3 3n^2 n + 1 = \Theta(n^3)$.
- 2-11. [3] Prove that $n^2 = O(2^n)$.
- 2-12. [3] For each of the following pairs of functions f(n) and g(n), give an appropriate positive constant c such that $f(n) \le c \cdot g(n)$ for all n > 1.
 - (a) $f(n) = n^2 + n + 1$, $g(n) = 2n^3$
 - (b) $f(n) = n\sqrt{n} + n^2$, $g(n) = n^2$
 - (c) $f(n) = n^2 n + 1$, $g(n) = n^2/2$
- 2-13. [3] Prove that if $f_1(n) = O(g_1(n))$ and $f_2(n) = O(g_2(n))$, then $f_1(n) + f_2(n) = O(g_1(n) + g_2(n))$.
- 2-14. [3] Prove that if $f_1(N) = \Omega(g_1(n))$ and $f_2(n) = \Omega(g_2(n))$, then $f_1(n) + f_2(n) = \Omega(g_1(n) + g_2(n))$.
- 2-15. [3] Prove that if $f_1(n) = O(g_1(n))$ and $f_2(n) = O(g_2(n))$, then $f_1(n) \cdot f_2(n) = O(g_1(n) \cdot g_2(n))$
- 2-16. [5] Prove for all $k \ge 1$ and all sets of constants $\{a_k, a_{k-1}, \dots, a_1, a_0\} \in R$, $a_k n^k + a_{k-1} n^{k-1} + \dots + a_1 n + a_0 = O(n^k)$
- 2-17. [5] Show that for any real constants a and b, b > 0 $(n+a)^b = \Theta(n^b)$
- 2-18. [5] List the functions below from the lowest to the highest order. If any two or more are of the same order, indicate which.

$$\begin{array}{llll} n & 2^n & n \lg n & \ln n \\ n-n^3+7n^5 & \lg n & \sqrt{n} & e^n \\ n^2+\lg n & n^2 & 2^{n-1} & \lg\lg n \\ n^3 & (\lg n)^2 & n! & n^{1+\varepsilon} \text{ where } 0 < \varepsilon < 1 \end{array}$$

2-19. [5] List the functions below from the lowest to the highest order. If any two or more are of the same order, indicate which.

- 2-20. [5] Find two functions f(n) and g(n) that satisfy the following relationship. If no such f and g exist, write "None."
 - (a) f(n) = o(g(n)) and $f(n) \neq \Theta(g(n))$
 - (b) $f(n) = \Theta(g(n))$ and f(n) = o(g(n))
 - (c) $f(n) = \Theta(q(n))$ and $f(n) \neq O(q(n))$
 - (d) $f(n) = \Omega(g(n))$ and $f(n) \neq O(g(n))$
- 2-21. [5] True or False?
 - (a) $2n^2 + 1 = O(n^2)$
 - (b) $\sqrt{n} = O(\log n)$
 - (c) $\log n = O(\sqrt{n})$
 - (d) $n^2(1+\sqrt{n}) = O(n^2 \log n)$
 - (e) $3n^2 + \sqrt{n} = O(n^2)$
 - (f) $\sqrt{n} \log n = O(n)$
 - (g) $\log n = O(n^{-1/2})$
- 2-22. [5] For each of the following pairs of functions f(n) and g(n), state whether f(n) = O(g(n)), $f(n) = \Omega(g(n))$, $f(n) = \Theta(g(n))$, or none of the above.
 - (a) $f(n) = n^2 + 3n + 4$, g(n) = 6n + 7
 - (b) $f(n) = n\sqrt{n}, q(n) = n^2 n$
 - (c) $f(n) = 2^n n^2$, $g(n) = n^4 + n^2$
- 2-23. [3] For each of these questions, briefly explain your answer.
 - (a) If I prove that an algorithm takes $O(n^2)$ worst-case time, is it possible that it takes O(n) on some inputs?
 - (b) If I prove that an algorithm takes $O(n^2)$ worst-case time, is it possible that it takes O(n) on all inputs?
 - (c) If I prove that an algorithm takes $\Theta(n^2)$ worst-case time, is it possible that it takes O(n) on some inputs?

- (d) If I prove that an algorithm takes $\Theta(n^2)$ worst-case time, is it possible that it takes O(n) on all inputs?
- (e) Is the function $f(n) = \Theta(n^2)$, where $f(n) = 100n^2$ for even n and $f(n) = 20n^2 n \log_2 n$ for odd n?
- 2-24. [3] For each of the following, answer yes, no, or can't tell. Explain your reasoning.
 - (a) Is $3^n = O(2^n)$?
 - (b) Is $\log 3^n = O(\log 2^n)$?
 - (c) Is $3^n = \Omega(2^n)$?
 - (d) Is $\log 3^n = \Omega(\log 2^n)$?
- 2-25. [5] For each of the following expressions f(n) find a simple g(n) such that $f(n) = \Theta(g(n))$.
 - (a) $f(n) = \sum_{i=1}^{n} \frac{1}{i}$.
 - (b) $f(n) = \sum_{i=1}^{n} \lceil \frac{1}{i} \rceil$.
 - (c) $f(n) = \sum_{i=1}^{n} \log i$.
 - (d) $f(n) = \log(n!)$.
- 2-26. $\it [5]$ Place the following functions into increasing asymptotic order.

$$f_1(n) = n^2 \log_2 n, f_2(n) = n(\log_2 n)^2, f_3(n) = \sum_{i=0}^n 2^i, f_4(n) = \log_2(\sum_{i=0}^n 2^i).$$

2-27. [5] Place the following functions into increasing asymptotic order. If two or more of the functions are of the same asymptotic order then indicate this.

$$f_1(n) = \sum_{i=1}^n \sqrt{i}, f_2(n) = (\sqrt{n}) \log n, f_3(n) = n \sqrt{\log n}, f_4(n) = 12n^{\frac{3}{2}} + 4n,$$

- 2-28. [5] For each of the following expressions f(n) find a simple g(n) such that $f(n) = \Theta(g(n))$. (You should be able to prove your result by exhibiting the relevant parameters, but this is not required for the homework.)
 - (a) $f(n) = \sum_{i=1}^{n} 3i^4 + 2i^3 19i + 20.$
 - (b) $f(n) = \sum_{i=1}^{n} 3(4^{i}) + 2(3^{i}) i^{19} + 20.$
 - (c) $f(n) = \sum_{i=1}^{n} 5^{i} + 3^{2i}$.
- 2-29. [5] Which of the following are true?
 - (a) $\sum_{i=1}^{n} 3^i = \Theta(3^{n-1}).$
 - (b) $\sum_{i=1}^{n} 3^i = \Theta(3^n)$.
 - (c) $\sum_{i=1}^{n} 3^i = \Theta(3^{n+1}).$
- 2-30. [5] For each of the following functions f find a simple function g such that $f(n) = \Theta(g(n))$.
 - (a) $f_1(n) = (1000)2^n + 4^n$.
 - (b) $f_2(n) = n + n \log n + \sqrt{n}$.
 - (c) $f_3(n) = \log(n^{20}) + (\log n)^{10}$.
 - (d) $f_4(n) = (0.99)^n + n^{100}$.

2-31. [5] For each pair of expressions (A,B) below, indicate whether A is O, o, Ω , ω , or Θ of B. Note that zero, one or more of these relations may hold for a given pair; list all correct ones.

	A	B
(a)	n^{100}	2^n
(b)	$(\lg n)^{12}$	\sqrt{n}
(c)	\sqrt{n}	$n^{\cos(\pi n/8)}$
(d)	10^n	100^{n}
(e)	$n^{\lg n}$	$(\lg n)^n$
(f)	$\lg(n!)$	$n \lg n$