Sharif University of Technology - Kolompeh - Notebook

Contents

1	Geo	metry 1	
	1.1	Rotating Calipers	
	1.2	Delaunay Triangulation $O(n^2)$	
	1.3	Convex Hull 3D	
	1.4	Half Plane Intersection	
	1.5	Useful Geo Facts	
2	String	\mathbf{ng} 5	
	2.1	Suffix Automata	
	2.2	Suffix Tree	
	2.3	Palindromic Tree	
_			
3		a structure 6	
	3.1	Treap	
	3.2	Treap Full	
	3.3	Link-cut Tree	
	3.4	Dynimic convex hull	
4	Cma	ph 8	
4	Gra j	Maximum matching - Edmond's blossom	
	4.1	Biconnected components	
	4.3	Flow - Dinic	
	4.4	Maximum weighted matching - Hungarian	
	4.5	Ear decomposition	
	4.6	Stoer-Wagner min cut $O(n^3)$	
	4.7	Directed minimum spanning tree $O(m \log n)$	
	4.8	Directed minimum spanning tree $O(nm)$	
	4.9	Dominator tree	
5	Con	abinatorics 14	
	5.1	LP simplex	
	5.2	FFT	
	5.3	NTT	
	5.4	Base Vectors in Z2	
	5.5	Gaussian Elimination	
	5.6	Stirling 1	
	5.7	Chinese remainder	
	5.8	Stirling 2	
	5.9	Popular LP	
	5.10	Duality of LP	
	5.11	Extended catalan	
	5.12	Find polynomial from it's points	
6	6 Constants 21		
O			
	$6.1 \\ 6.2$	Number of primes 21 Factorials 21	
	6.3	Powers of 3	
	6.4	C(2n,n)	
	6.5	Most divisor	
	0.5	10000 (1101001	

1 Geometry

1.1 Rotating Calipers

```
vector <pair<pt, pt>> get_antipodals(vector <pt> &p) {
        int n = sz(p);
        sort(p.begin(), p.end());
        vector <pt> U, L;
        for (int i = 0; i < n; i++) {
                while (sz(U) > 1 \&\& side(U[sz(U)-2], U[sz(U)-1], p[i])
                        U.pop_back();
                while (sz(L) > 1 \&\& side(L[sz(L)-2], L[sz(L)-1], p[i])
                         L.pop_back();
                U.pb(p[i]);
                L.pb(p[i]);
        vector <pair<pt, pt>> res;
        int i = 0, j = sz(L)-1;
        while (i+1 < sz(U) | | j > 0) {
                res.pb({U[i], L[j]});
                if (i+1 == sz(U))
                         j--;
                else if (j == 0)
                         <u>i</u>++;
                else if (cross(L[j]-L[j-1], U[i+1]-U[i]) >= 0)
                else
                         j--;
        return res;
```

1.2 Delaunay Triangulation $O(n^2)$

```
struct Delaunay{
       vector <pt> p;
       vector <pt> to;
       vector <int> nxt;
        int add_edge(pt q, int bef=-1){
                int cnt = sz(to);
                to.pb(q);
                nxt.pb(-1);
                if (bef !=-1) {
                        nxt[bef] = cnt;
                        to.pb(to[bef]);
                        nxt.pb(-1);
                return cnt;
        int before(int e) {
                int cur = e, last = -1;
                do{
                        last = cur;
                        cur = nxt[cur^1];
                }while (cur != e);
```

```
return last^1;
                                                                                     if (incircle(b, a, d, c)) return false;
                                                                                     return true:
                                                                             }
void easy_triangulate(){
        to.clear();
                                                                             void flip(int e) {
        nxt.clear();
                                                                                     int a = nxt[e], b = nxt[a],
        sort(p.begin(), p.end());
                                                                                             c = nxt[e^1], d = nxt[c];
        if (dir(p[0], p[1], p[2]) > 0)
                                                                                     nxt[d] = a;
                swap(p[1], p[2]);
                                                                                     nxt[b] = c;
        int to0 = add_edge(p[0]), to0c = add_edge(p[2]),
                                                                                     to[e] = to[c];
                to1 = add_edge(p[1]), to1c = add_edge(p[0]),
                                                                                     nxt[a] = e;
                                                                                     to[e^1] = to[a];
                to2 = add_edge(p[2]), to2c = add_edge(p[1]);
                                                                                     nxt[c] = e^1;
        nxt[to1] = to2; nxt[to2] = to0;
        nxt[to0] = to1; nxt[to0c] = to2c;
        nxt[to2c] = to1c; nxt[to1c] = to0c;
                                                                             void delaunay_triangulate(){
                                                                                     if (sz(to) == 0)
        int e = to0;
                                                                                             easy_triangulate();
        for (int i = 3; i < sz(p); i++) {
                                                                                     bool *mark = new bool[sz(to)];
                pt q = p[i];
                                                                                     fill(mark, mark + sz(to), false);
                                                                                     vector <int> bad;
                while (dir(q, to[e^1], to[e]) < 0 \mid | dir(q, to
                    [e^1], to [before(e)^1] < 0)
                                                                                     for (int e = 0; e < sz(to); e++) {
                                                                                             if (!mark[e/2] && !locally(e)){
                        e = nxt[e];
                vector <int> vis;
                                                                                                     bad.pb(e);
                while (dir(q, to[e^1], to[e]) > 0 \mid | dir(q, to
                                                                                                     mark[e/2] = true;
                    [e^1], to [before(e)^1]) > 0) {
                        vis.pb(e);
                        e = nxt[e];
                                                                                     while (sz(bad)) {
                                                                                             int e = bad.back();
                                                                                             bad.pop_back();
                int ex = add_edge(q, before(vis[0]));
                int last = ex^1;
                                                                                             mark[e/2] = false;
                for (int edge : vis) {
                                                                                             if (!locally(e)){
                        nxt[last] = edge;
                                                                                                      flip(e);
                        int eq = add_edge(q, edge);
                                                                                                      int to_check[4] = {nxt[e], nxt[nxt[e]]
                        nxt[edge] = eq;
                                                                                                          ]], nxt[e^1], nxt[nxt[e^1]]};
                        nxt[eq] = last;
                                                                                                      for (int i = 0; i < 4; i++)
                        last = eq^1:
                                                                                                              if (!mark[to_check[i]/2] && !
                nxt[ex] = last;
                nxt[last] = e;
bool incircle(pt a, pt b, pt c, pt d) {
        return a.z() * (b.x * (c.y - d.y) - c.x * (b.y - d.y)
            + d.x * (b.y - c.y)
                - b.z() * (a.x * (c.y - d.y) - c.x * (a.y - d.y)
                                                                             vector <tri> get_triangles() {
                    y) + d.x * (a.y - c.y))
                                                                                     vector <tri> res;
                + c.z() * (a.x * (b.v - d.v) - b.x * (a.v - d.
                                                                                     for (int e = 0; e < sz(to); e++) {
                    y) + d.x * (a.y - b.y))
                                                                                             pt a = to[e^1], b = to[e], c = to[nxt[e]];
                - d.z() * (a.x * (b.y - c.y) - b.x * (a.y - c.
                                                                                             if (dir(a, b, c) < 0) continue;</pre>
                    y) + c.x * (a.y - b.y)) > 0;
                                                                                             res.pb(tri(a, b, c));
                                                                                     return res;
bool locally(int e){
                                                                             Delaunay(vector <pt> p):p(p){}
        pt a = to[e^1], b = to[e], c = to[nxt[e]], d = to[nxt[e]]
            e^1]];
        if (dir(a, b, c) < 0) return true;</pre>
        if (dir(b, a, d) < 0) return true;</pre>
                                                                 1.3 Convex Hull 3D
        if (incircle(a, b, c, d)) return false;
```

locally(to_check[i])){

true;

bad.pb(to_check[i]);

mark[to_check[i]/2] =

```
/*
  GETS:
  n->number of vertices
  you should use add_edge(u, v) and
  add pair of vertices as edges (vertices are 0..n-1)
  GIVES:
  output of edmonds() is the maximum matching in general graph
  match[i] is matched pair of i (-1 if there isn't a matched pair)
  0 (nh)
*/
#include<bits/stdc++.h>
using namespace std;
typedef pair<int,int> pii;
struct point{
  int X,Y,Z;
  point(int x=0, int y=0, int z=0) {
    X=x:
    Y=V:
    z=z;
  bool operator==(const point& rhs) const {
    return (rhs.X==this->X && rhs.Y==this->Y && rhs.Z==this->Z);
 bool operator<(const point& rhs) const {</pre>
    return rhs.X > this->X || (rhs.X == this->X && rhs.Y > this->Y) ||
         (rhs.X==this->X && rhs.Y==this->Y && rhs.Z>this->Z);
};
const int maxn=1000;
int n;
point P[maxn];
vector<point>ans:
queue<pii>Q;
set<pii>mark;
int cross2d(point p,point q) { return p.X*q.Y-p.Y*q.X; }
point operator - (point p, point q) { return point (p.X-q.X,p.Y-q.Y,p.Z-q.
    Z); }
int dot(point v,point u) { return u.X*v.X+u.Y*v.Y+u.Z*v.Z; }
point _cross(point u, point v) { return point(u.Y*v.Z-u.Z*v.Y, u.Z*v.X-u.
    X*v.Z,u.X*v.Y-u.Y*v.X);
point cross(point o, point p, point q) { return _cross(p-o, q-o);}
point shift(point p) { return point(p.Y,p.Z,p.X);}
point norm(point p)
  if(p.Y<p.X || p.Z<p.X) p=shift(p);</pre>
  if(p.Y<p.X) p=shift(p);</pre>
  return p;
int main()
  cin>>n;
  int mn=0;
  for (int i=0;i<n;i++) {</pre>
```

```
cin>>P[i].X>>P[i].Y>>P[i].Z;
  if(P[i]<P[mn]) mn=i;
int nx=(mn==0);
for (int i=0; i<n; i++)</pre>
  if(i!=mn && i!=nx && cross2d(P[nx]-P[mn],P[i]-P[mn])>0)
Q.push(pii(mn,nx));
while(!O.empty())
    int v=Q.front().first,u=Q.front().second;
    Q.pop();
    if (mark.find(pii(v,u))!=mark.end()) continue;
    mark.insert(pii(v,u));
    int p=-1;
    for (int q=0; q<n; q++)
      if(q!=v && q!=u)
        if (p==-1 || dot(cross(P[v],P[u],P[p]),P[q]-P[v])<0)</pre>
    ans.push_back(norm(point(v,u,p)));
    Q.push(pii(p,u));
    Q.push(pii(v,p));
sort(ans.begin(),ans.end());
ans.resize(unique(ans.begin(), ans.end())-ans.begin());
for(int i=0;i<ans.size();i++)</pre>
  cout << ans[i]. X << " " << ans[i]. Y << " " << ans[i]. Z << endl;</pre>
```

1.4 Half Plane Intersection

```
typedef int T;
typedef long long T2;
typedef long long T4; // maybe int128_t
const int MAXLINES = 100 * 1000 + 10;
const int INF = 20 * 1000 * 1000;
typedef pair<T, T> point;
typedef pair<point, point> line;
#define X first
#define Y second
#define A first
#define B second
// REPLACE ZERO WITH EPS FOR DOUBLE
point operator - (const point &a, const point &b) {
        return point(a.X - b.X, a.Y - b.Y);
T2 cross(point a, point b) {
        return ((T2)a.X * b.Y - (T2)a.Y * b.X);
bool cmp(line a, line b) {
        bool aa = a.A < a.B;</pre>
        bool bb = b.A < b.B;</pre>
```

```
if (aa == bb) {
                point v1 = a.B - a.A;
                point v2 = b.B - b.A;
                                                                             line dq[MAXLINES];
                if (cross(v1, v2) == 0)
                        return cross(b.B - b.A, a.A - b.A) > 0;
                                                                             vector<line> half plane(vector<line> lines) {
                else
                                                                                      lines.push back(line(point(INF, -INF), point(INF, INF)));
                                                                                      lines.push_back(line(point(-INF, INF), point(-INF, -INF)));
                        return cross (v1, v2) > 0;
                                                                                      lines.push_back(line(point(-INF, -INF), point(INF, -INF)));
        else
                                                                                      lines.push back(line(point(INF, INF), point(-INF, INF)));
                                                                                      sort(lines.begin(), lines.end(), cmp);
                return aa;
                                                                                     int ptr = 0;
                                                                                      for (int i = 0; i < lines.size(); i++)</pre>
bool parallel(line a, line b) {
                                                                                              if (i > 0 &&
        return cross(a.B - a.A, b.B - b.A) == 0;
                                                                                                      (lines[i-1].A < lines[i-1].B) == (lines[i-1].B)
                                                                                                          1.A < lines[i].B) &&</pre>
                                                                                                      parallel(lines[i - 1], lines[i]))
pair<T2, T2> alpha(line a, line b) {
                                                                                                      continue:
        return pair<T2, T2>(cross(b.A - a.A, b.B - b.A),
                                                                                              else
                                                 cross(a.B - a.A. b.B -
                                                                                                      lines[ptr++] = lines[i];
                                                                                     lines.resize(ptr):
                                                      b.A));
                                                                                      if (lines.size() < 2)</pre>
                                                                                              return lines;
bool fcmp(T4 flt, T4 flb, T4 f2t, T4 f2b) {
                                                                                      //print(lines);
        if (f1b < 0) {
                                                                                      int f = 0, e = 0;
               f1t *= -1;
                                                                                      dq[e++] = lines[0];
                f1b *= -1;
                                                                                      dq[e++] = lines[1];
                                                                                      for (int i = 2; i < lines.size(); i++) {</pre>
        if (f2b < 0) {
                                                                                              while (f < e - 1 && notin(dq[e - 2], dq[e - 1], lines[
                f2t *= -1;
                                                                                                      e--;
                f2b *= -1;
                                                                                              //print(vector<line>(dq + f, dq + e));
        return flt * f2b < f2t * f1b; // check with eps
                                                                                              if (e == f + 1) {
                                                                                                      T2 crs = cross(dq[f].B - dq[f].A, lines[i].B -
                                                                                                           lines[i].A) ;
bool check(line a, line b, line c) {
                                                                                                      if (crs < 0)
        bool crs = cross(c.B - c.A, a.B - a.A) > 0;
                                                                                                              return vector<line>();
        pair<T2, T2> a1 = alpha(a, b);
                                                                                                      else if (crs == 0 && cross(lines[i].B - lines[
        pair<T2, T2>a2=alpha(a, c);
                                                                                                          i].A, dq[f].B - lines[i].A) < 0)
        bool alp = fcmp(a1.A, a1.B, a2.A, a2.B);
                                                                                                              return vector<line>();
        return (crs ^ alp);
                                                                                              while (f < e - 1 \&\& notin(dq[f], dq[f + 1], lines[i]))
bool notin(line a, line b, line c) { // is intersection of a and b in
                                                                                              d\alpha[e++] = lines[i];
    ccw direction of c?
        if (parallel(a, b))
                                                                                      while (f < e - 1 \&\& notin(dq[e - 2], dq[e - 1], dq[f]))
                return false;
        if (parallel(a, c))
                                                                                      while (f < e - 1 \&\& notin(dq[f], dq[f + 1], dq[e - 1]))
                return cross(c.B - c.A, a.A - c.A) < 0;
                                                                                              f++;
        if (parallel(b, c))
                                                                                      vector<line> res;
                return cross(c.B - c.A, b.A - c.A) < 0;
                                                                                      res.resize(e - f);
        return ! (check(a, b, c) && check(b, a, c));
                                                                                      for (int i = f; i < e; i++)
                                                                                              res[i - f] = dq[i];
                                                                                      return res;
void print(vector<line> lines) {
        cerr << " " << endl; for (int i = 0; i < lines.size();</pre>
            i++)cerr << lines[i].A.X << " " << lines[i].A.Y << " -> "
            << lines[i].B.X << " " << lines[i].B.Y << endl;cerr << "
                                                                             int main() {
                                                                                      int n:
                                                                                      cin >> n:
                                                                                      vector<line> lines:
 << endl<< endl:
                                                                                      for (int i = 0; i < n; i++) {</pre>
```

1.5 Useful Geo Facts

```
Area of triangle with sides a, b, c: sqrt(S *(S-a)*(S-b)*(S-c)) where
    S = (a+b+c)/2

Area of equilateral triangle: s^2 * sqrt(3) / 4 where is side lenght

Pyramid and cones volume: 1/3 area(base) * height
Sphere volume: 4/3 pi r^3
Sphere area: 4 pi r^2

if p1, p2, p3 are points in circle, the center is
x = -((x2^2 - x1^2 + y2^2 - y1^2)*(y3 - y2) - (x2^2 - x3^2 + y2^2 - y3^2)*(y1 - y2)) / (2*(x1 - x2)*(y3 - y2) - 2*(x3 - x2)*(y1 - y2))
y = -((y2^2 - y1^2 + x2^2 - x1^2)*(x3 - x2) - (y2^2 - y3^3 + x2^2 - x3^2)*(x1 - x2)) / (2*(y1 - y2)*(x3 - x2) - 2*(y3 - y2)*(x1 - x2))
```

2 String

2.1 Suffix Automata

```
const int maxn = 2 e5 + 42; // Maximum amount of states
map < char , int > to [ maxn ]; // Transitions
int link [ maxn ]; // Suffix links
int len [ maxn ]; // Lengthes of largest strings in states
int last = 0; // State corresponding to the whole string
int sz = 1; // Current amount of states
void add_letter ( char c ) { // Adding character to the end
    int p = last ; // State of string s
   last = sz ++; // Create state for string sc
   len [ last ] = len [ p ] + 1;
    for (; to [ p ][ c ] == 0; p = link [ p ]) // (1)
        to [ p ][ c ] = last; // Jumps which add new suffixes
   if ( to [ p ][ c ] == last ) { // This is the first occurrence of
        c if we are here
        link [ last ] = 0;
        return ;
   int q = to [ p ][ c ];
   if ( len [ q ] == len [ p ] + 1) {
        link [ last ] = q;
        return ;
    // We split off cl from a here
    int cl = sz ++;
    to [cl] = to [q]; // (2)
```

2.2 Suffix Tree

```
#define fpos adla
const int inf = 1e9;
const int maxn = 1e4;
char s[maxn];
map<int, int> to[maxn];
int len[maxn], fpos[maxn], link[maxn];
int node, pos;
int sz = 1, n = 0;
int make_node(int _pos, int _len) {
    fpos[sz] = \_pos;
    len [sz] = _len;
    return sz++;
void go_edge() {
    while(pos > len[to[node][s[n - pos]]]) {
       node = to[node][s[n - pos]];
        pos -= len[node];
void add_letter(int c) {
    s[n++] = c;
    pos++;
    int last = 0;
    while(pos > 0) {
        go_edge();
        int edge = s[n - pos];
        int &v = to[node][edge];
        int t = s[fpos[v] + pos - 1];
        if(v == 0) {
            v = make_node(n - pos, inf);
            link[last] = node;
            last = 0;
        } else if(t == c) {
            link[last] = node;
            return;
        } else {
            int u = make_node(fpos[v], pos - 1);
            to[u][c] = make_node(n - 1, inf);
            to[u][t] = v;
            fpos[v] += pos - 1;
            len [v] -= pos - 1;
            v = u;
            link[last] = u;
            last = u;
        if(node == 0)
            pos--;
        else
            node = link[node];
```

```
int main() {
    len[0] = inf;
    string s;
    cin >> s;
    int ans = 0;
    for(int i = 0; i < s.size(); i++)
        add_letter(s[i]);
    for(int i = 1; i < sz; i++)
        ans += min((int)s.size() - fpos[i], len[i]);
    cout << ans << "\n";
}</pre>
```

2.3 Palindromic Tree

```
int n, last, sz;
void init() {
    s[n++] = -1;
    link[0] = 1;
    len[1] = -1;
    sz = 2;
int get_link(int v) {
    while (s[n - len[v] - 2] != s[n - 1]) v = link[v];
    return v;
void add_letter(int c) {
    s[n++] = c;
    last = get_link(last);
    if(!to[last][c]) {
        len [sz] = len[last] + 2;
        link[sz] = to[get_link(link[last])][c];
        to[last][c] = sz++;
    last = to[last][c];
```

3 Data structure

3.1 Treap

```
void insert (pitem & t, pitem it) {
    if (!t)
       t = it;
    else if (it->prior > t->prior)
        split (t, it->key, it->l, it->r), t = it;
       insert (it->key < t->key ? t->l : t->r, it);
void merge (pitem & t, pitem 1, pitem r) {
   if (!l || !r)
       t = 1 ? 1 : r;
    else if (l->prior > r->prior)
       merge (1->r, 1->r, r), t = 1;
       merge (r->1, 1, r->1), t = r;
void erase (pitem & t, int key) {
   if (t->key == key)
       merge (t, t->1, t->r);
    else
        erase (key < t->key ? t->l : t->r, key);
pitem unite (pitem l, pitem r) {
   if (!1 || !r) return 1 ? 1 : r;
   if (l->prior < r->prior) swap (l, r);
   pitem lt, rt;
   split (r, 1->key, lt, rt);
   1->1 = unite (1->1, lt);
   1->r = unite (1->r, rt);
   return 1;
```

3.2 Treap Full

```
typedef struct item * pitem;
struct item {
    int prior, value, cnt;
   bool rev;
   pitem 1, r;
};
int cnt (pitem it) {
    return it ? it->cnt : 0;
void upd_cnt (pitem it) {
   if (it)
       it->cnt = cnt(it->1) + cnt(it->r) + 1;
void push (pitem it) {
   if (it && it->rev) {
       it->rev = false;
       swap (it->1, it->r);
       if (it->1) it->1->rev ^= true;
       if (it->r) it->r->rev ^= true;
void merge (pitem & t, pitem l, pitem r) {
```

```
push (1);
   push (r);
   if (!l || !r)
       t = 1 ? 1 : r;
   else if (l->prior > r->prior)
        merge (1->r, 1->r, r), t = 1;
   else
        merge (r->1, 1, r->1), t = r;
   upd_cnt (t);
void split (pitem t, pitem & l, pitem & r, int key, int add = 0) {
   if (!t)
        return void( 1 = r = 0 );
   push (t);
   int cur_key = add + cnt(t->1);
   if (key <= cur_key)</pre>
        split (t->1, 1, t->1, key, add), r = t;
   else
        split (t->r, t->r, r, key, add + 1 + cnt(t->1)), l = t;
   upd_cnt (t);
void reverse (pitem t, int 1, int r) {
   pitem t1, t2, t3;
   split (t, t1, t2, 1);
   split (t2, t2, t3, r-l+1);
   t2->rev ^= true;
   merge (t, t1, t2);
   merge (t, t, t3);
void output (pitem t) {
   if (!t) return;
   push (t);
   output (t->1);
   printf ("%d ", t->value);
   output (t->r);
```

3.3 Link-cut Tree

```
Node x[N];
struct Node {
    int sz, label; /* size, label */
    Node *p, *pp, *l, *r; /* parent, path-parent, left, right pointers
    */
    Node() { p = pp = l = r = 0; }
};

void update(Node *x) {
    x->sz = 1;
    if(x->l) x->sz += x->l->sz;
    if(x->r) x->sz += x->r->sz;
}

void rotr(Node *x) {
    Node *y, *z;
    y = x->p, z = y->p;
```

```
if(y == z->1) z->1 = x;
        else z->r = x;
    x->pp = y->pp;
    y->pp = 0;
    update(y);
void rotl(Node *x) {
    Node *y, *z;
    y = x->p, z = y->p;
    if ((y->r = x->1)) y->r->p = y;
    x->1 = y, y->p = x;
    if((x->p = z))
        if(y == z->1) z->1 = x;
        else z \rightarrow r = x:
    x->pp = y->pp;
    y->pp = 0;
    update(y);
void splay(Node *x) {
    Node *y, *z;
    while (x->p) {
        v = x - > p;
        if(y->p == 0) {
             if(x == y->1) rotr(x);
             else rotl(x);
         else {
             z = y - > p;
             if(y == z->1) {
                 if(x == y->1) rotr(y), rotr(x);
                 else rotl(x), rotr(x);
             else { if(x == y->r) rotl(y), rotl(x);
                 else rotr(x), rotl(x);
    update(x);
Node *access(Node *x) {
    splav(x);
    if(x->r) {
        x \rightarrow r \rightarrow pp = x;
        x->r->p = 0;
        x->r = 0;
        update(x);
    Node *last = x;
    while(x->pp) {
        Node *y = x - pp;
        last = y;
        splay(y);
```

if ((y->1 = x->r)) y->1->p = y;

x->r = y, y->p = x;

if((x->p = z))

```
if(y->r) {
             y->r->pp = y;
             y \rightarrow r \rightarrow p = 0;
        y->r = x;
        x->p = v;
        x->pp = 0;
        update(y);
        splay(x);
    return last;
Node *root (Node *x) {
    access(x);
    while (x->1) x = x->1;
    splay(x);
    return x;
void cut(Node *x) {
    access(x);
    x - > 1 - > p = 0;
    x->1 = 0;
    update(x);
void link(Node *x, Node *y) {
    access(x);
    access(v);
    x->1 = y;
    y->p = x;
    update(x);
Node *lca(Node *x, Node *y) {
    access(x):
    return access(y);
int depth(Node *x) {
    access(x):
    return x->sz - 1;
void init(int n) {
    for (int i = 0; i < n; i++) {
        x[i].label = i;
        update(&x[i]);
```

3.4 Dynimic convex hull

```
const ld is_query = -(1LL << 62);
struct Line {
   ld m, b;
   mutable std::function<const Line *()> succ;
```

```
bool operator<(const Line &rhs) const {</pre>
        if (rhs.b != is_query) return m < rhs.m;</pre>
        const Line *s = succ();
        if (!s) return 0;
        ld x = rhs.m;
        return b - s -> b < (s -> m - m) * x;
};
struct HullDynamic : public multiset<Line> { // dynamic upper hull +
    max value query
    bool bad(iterator y) {
        auto z = next(y);
        if (y == begin())
            if (z == end()) return 0;
            return y->m == z->m && y->b <= z->b;
        auto x = prev(y);
        if (z == end()) return y->m == x->m && y->b <= x->b;
        return (x-b-y-b) * (z-m-y-m) >= (y-b-z-b) * (y-m)
            - x->m):
    void insert line(ld m, ld b) {
        auto y = insert({m, b});
        y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
        if (bad(v)) {
            erase(y);
            return;
        while (next(y) != end() \&\& bad(next(y))) erase(next(y));
        while (y != begin() && bad(prev(y))) erase(prev(y));
    ld best(ld x) {
        auto l = *lower_bound((Line) {x, is_query});
        return 1.m * x + 1.b:
};
```

4 Graph

4.1 Maximum matching - Edmond's blossom

```
/*
GETS:
n->number of vertices
you should use add_edge(u,v) and
add pair of vertices as edges (vertices are 0..n-1)
(note: please don't add multiple edge)

GIVES:
output of edmonds() is the maximum matching in general graph
match[i] is matched pair of i (-1 if there isn't a matched pair)

O(mn^2)
*/
```

```
#include <bits/stdc++.h>
using namespace std;
struct struct_edge{int v;struct_edge* nxt;};
typedef struct_edge* edge;
const int MAXN=500;
struct Edmonds
        struct_edge pool[MAXN*MAXN*2];
        edge top=pool,adj[MAXN];
        int n, match[MAXN], qh, qt, q[MAXN], father[MAXN], base[MAXN];
        bool ing[MAXN], inb[MAXN];
        void add_edge(int u,int v)
                top->v=v,top->nxt=adj[u],adj[u]=top++;
                top->v=u,top->nxt=adj[v],adj[v]=top++;
        int LCA(int root,int u,int v)
                static bool inp[MAXN];
                memset(inp, 0, sizeof(inp));
                while(1)
                         inp[u=base[u]]=true;
                         if (u==root) break;
                         u=father[match[u]];
                while(1)
                         if (inp[v=base[v]]) return v;
                         else v=father[match[v]];
        void mark blossom(int lca.int u)
                while (base[u]!=lca)
                         int v=match[u];
                         inb[base[u]]=inb[base[v]]=true;
                         u=father[v];
                         if (base[u]!=lca) father[u]=v;
        void blossom_contraction(int s,int u,int v)
                int lca=LCA(s,u,v);
                memset(inb, 0, sizeof(inb));
                mark_blossom(lca,u);
                mark_blossom(lca, v);
                if (base[u]!=lca)
                         father[u]=v;
                if (base[v]!=lca)
                         father[v]=u;
                for (int u=0; u < n; u++)
                        if (inb[base[u]])
                                 base[u]=lca;
```

```
if (!inq[u])
                                  inq[q[++qt]=u]=true;
int find augmenting path(int s)
        memset(inq, 0, sizeof(inq));
        memset(father,-1,sizeof(father));
        for (int i=0;i<n;i++) base[i]=i;</pre>
        inq[q[qh=qt=0]=s]=true;
        while (qh<=qt)</pre>
                 int u=q[qh++];
                 for (edge e=adj[u];e;e=e->nxt)
                         int v=e->v;
                         if (base[u]!=base[v] && match[u]!=v)
                                  if (v==s || (match[v]!=-1 &&
                                      father[match[v]]!=-1))
                                          blossom contraction(s,
                                               u, v);
                                  else if (father[v]==-1)
                                          father[v]=u;
                                          if (match[v] == -1)
                                                   return v;
                                          else if (!inq[match[v
                                              ]])
                                                   inq[q[++qt] =
                                                       match[v]]=
                                                       true;
        return -1;
int augment_path(int s,int t)
        int u=t, v, w;
        while (u!=-1)
                 v=father[u];
                w=match[v];
                match[v]=u;
                match[u]=v;
                u=w;
        return t!=-1;
int edmonds()
        int matchc=0;
        memset (match, -1, sizeof (match));
        for (int u=0;u<n;u++)</pre>
                if (match[u] ==-1)
                         matchc+=augment_path(u,
```

```
find_augmenting_path(u));
return matchc;
};
```

4.2 Biconnected components

```
vector<int> adj[maxn];
bool vis[maxn];
int dep[maxn], par[maxn], lowlink[maxn];
vector<vector<int> > comp;
stack<int> st;
void dfs(int u, int depth = 0, int parent = -1)
        vis[u] = true;
        dep[u] = depth;
        par[u] = parent;
        lowlink[u] = depth;
        st.push(u):
        for (int i = 0; i < adj[u].size(); i++)</pre>
                int v = adj[u][i];
                if (!vis[v])
                        dfs(v, depth + 1, u);
                        lowlink[u] = min(lowlink[u], lowlink[v]);
                else
                        lowlink[u] = min(lowlink[u], dep[v]);
        if (lowlink[u] == dep[u] - 1)
                comp.push_back(vector<int>());
                while (st.top() != u)
                        comp.back().push_back(st.top());
                        st.pop();
                comp.back().push_back(u);
                st.pop();
                comp.back().push_back(par[u]);
void bicon(int n)
        for (int i = 0; i < n; i++)
                if (!vis[i])
                        dfs(i);
```

4.3 Flow - Dinic

```
const int MAXN = ???; //XXX
const int MAXE = ???????; //XXX

int from[MAXE], to[MAXE], cap[MAXE], prv[MAXE], head[MAXN], pt[MAXN],
    ec;
```

```
void addEdge(int u, int v, int uv, int vu = 0) {
        from[ec] = u, to[ec] = v, cap[ec] = uv, prv[ec] = head[u],
            head[u] = ec++;
        from[ec] = v, to[ec] = u, cap[ec] = vu, prv[ec] = head[v],
            head[v] = ec++;
int lv[MAXN], q[MAXN];
bool bfs(int source, int sink){
        memset(lv, 63, sizeof(lv));
        int h = 0, t = 0;
        lv[source] = 0;
        q[t++] = source;
        while (t-h) {
                int v = q[h++];
                for (int e = head[v]; ~e; e = prv[e])
                        if (cap[e] && lv[v] + 1 < lv[to[e]]){</pre>
                                 lv[to[e]] = lv[v] + 1;
                                 a[t++] = to[e]:
        return lv[sink] < 1e8;</pre>
int dfs(int v, int sink, int f = 1e9) {
        if (v == sink || f == 0)
                return f;
        int ret = 0;
        for (int &e = pt[v]; ~e; e = prv[e])
                if (lv[v]+1 == lv[to[e]]) {
                         int x = dfs(to[e], sink, min(f, cap[e]));
                         cap[e] -= x;
                        cap[e^1] += x;
                        ret += x;
                        f -= x;
                        if (!f)
                                 break:
        return ret;
int dinic(int source, int sink){
        int ret = 0;
        while (bfs(source, sink)){
                memcpy(pt, head, sizeof(head));
                ret += dfs(source, sink);
        return ret;
```

4.4 Maximum weighted matching - Hungarian

```
const int N = 2002;
const int INF = 1e9;
int hn, weight[N][N];
int x[N], y[N];
int hungarian() // maximum weighted perfect matching
{
```

```
int n = hn;
int p, q;
vector<int> fx(n, -INF), fy(n, 0);
fill(x, x + n, -1);
fill(y, y + n, -1);
for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j)
                fx[i] = max(fx[i], weight[i][j]);
for (int i = 0; i < n; ) {</pre>
        vector<int> t(n, -1), s(n+1, i);
        for (p = 0, q = 1; p < q \&\& x[i] < 0; ++p) {
                int k = s[p];
                for (int j = 0; j < n \&\& x[i] < 0; ++j)
                        if (fx[k] + fy[j] == weight[k][j] && t
                             [i] < 0)
                                 s[q++] = y[j], t[j] = k;
                                 if (y[j] < 0) // match found!
                                         for (int p = j; p >=
                                             0; j = p
                                                 y[i] = k = t[i]
                                                      ], p = x[k]
                                                      ], x[k] =
                                                      j;
        if (x[i] < 0) {
                int d = INF;
                for (int k = 0; k < q; ++k)
                        for (int j = 0; j < n; ++j)
                                 if (t[j] < 0) d = min(d, fx [s
                                     [k]] + fy[j] - weight[s]k
                                     ]][j]);
                for (int j = 0; j < n; ++j) fy[j] += (t[j] <
                    0? 0: d);
                for (int k = 0; k < q; ++k) fx[s[k]] -= d;
        } else ++i;
int ret = 0;
for (int i = 0; i < n; ++i) ret += weight[i][x[i]];</pre>
return ret;
```

4.5 Ear decomposition

- 1- Find a spanning tree of the given graph and choose a root for the tree.
- 2- Determine, for each edge uv that is not part of the tree, the distance between the root and the lowest common ancestor of \boldsymbol{u} and \boldsymbol{v} .
- 3- For each edge uv that is part of the tree, find the corresponding "master edge", a non-tree edge wx such that the cycle formed by adding wx to the tree passes through uv and such that, among such edges, w and x have a lowest common ancestor that is as close to the root as possible (with ties broken by edge identifiers).
- 4- Form an ear for each non-tree edge, consisting of it and the tree edges for which it is the master, and order the ears by their master edges' distance from the root (with the same tie-breaking rule).

4.6 Stoer-Wagner min cut $O(n^3)$

```
const int N = -1, MAXW = -1;
int g[N][N], v[N], w[N], na[N];
bool a[N];
int minCut( int n ) // initialize q[][] before calling!
    for ( int i = 0; i < n; i++ ) v[i] = i;
    int best = MAXW * n * n;
    while (n > 1)
        // initialize the set A and vertex weights
        a[v[0]] = true;
        for ( int i = 1; i < n; i++ )
            a[v[i]] = false;
            na[i - 1] = i;
            w[i] = g[v[0]][v[i]];
        // add the other vertices
        int prev = v[0];
        for( int i = 1; i < n; i++ )</pre>
            // find the most tightly connected non-A vertex
            int zj = -1;
            for( int j = 1; j < n; j++ )</pre>
                if( !a[v[j]] && (zj < 0 \mid \mid w[j] > w[zj]))
                    zj = j;
            // add it to A
            a[v[zj]] = true;
            // last vertex?
            if(i == n - 1)
                // remember the cut weight
                best = min(best, w[zj]);
                // merge prev and v[zj]
                for ( int j = 0; j < n; j++ )
                    q[v[j]][prev] = q[prev][v[j]] += q[v[zj]][v[j]];
                v[zj] = v[--n];
                break;
            prev = v[zj];
            // update the weights of its neighbors
            for ( int j = 1; j < n; j++ ) if ( [a[v[j]]] )
                w[j] += q[v[zj]][v[j]];
    return best:
```

4.7 Directed minimum spanning tree $O(m \log n)$

```
/*
        GETS:
                call make_graph(n) at first
                you should use add_edge(u,v,w) and
                add pair of vertices as edges (vertices are 0..n-1)
        GIVES:
                output of dmst(v) is the minimum arborescence with
                    root v in directed graph
                (INF if it hasn't a spanning arborescence with root v)
        O(mlogn)
*/
#include <bits/stdc++.h>
using namespace std;
const int INF = 2e7;
struct MinimumAborescense
        struct edge {
                int src, dst, weight;
        struct union_find {
                vector<int> p:
                union_find(int n) : p(n, -1) { };
                bool unite(int u, int v) {
                        if ((u = root(u)) == (v = root(v))) return
                            false;
                        if (p[u] > p[v]) swap(u, v);
                        p[u] += p[v]; p[v] = u;
                        return true;
                bool find(int u, int v) { return root(u) == root(v); }
                int root(int u) { return p[u] < 0 ? u : p[u] = root(p[</pre>
                    ul); }
                int size(int u) { return -p[root(u)]; }
        };
        struct skew heap {
                struct node {
                        node *ch[2];
                        edge key;
                        int delta;
                } *root;
                skew_heap() : root(0) { }
                void propagate(node *a) {
                        a->key.weight += a->delta;
                        if (a->ch[0]) a->ch[0]->delta += a->delta;
                        if (a->ch[1]) a->ch[1]->delta += a->delta;
                        a \rightarrow delta = 0;
                node *merge(node *a, node *b) {
                        if (!a || !b) return a ? a : b;
                        propagate(a); propagate(b);
                        if (a->key.weight > b->key.weight) swap(a, b);
                        a - ch[1] = merge(b, a - ch[1]);
```

```
swap(a->ch[0], a->ch[1]);
                return a;
        void push(edge key) {
                node *n = new node();
                n \rightarrow ch[0] = n \rightarrow ch[1] = 0;
                n->key = key; n->delta = 0;
                root = merge(root, n);
        void pop() {
                propagate(root);
                node *temp = root;
                root = merge(root->ch[0], root->ch[1]);
        edge top() {
                propagate(root);
                return root->key;
        bool empty() {
                return !root;
        void add(int delta) {
                root->delta += delta;
        void merge(skew_heap x) {
                root = merge(root, x.root);
};
vector<edge> edges;
void add_edge(int src, int dst, int weight) {
        edges.push_back({src, dst, weight});
int n:
void make_graph(int _n) {
        n = _n;
        edges.clear();
int dmst(int r) {
        union_find uf(n);
        vector<skew_heap> heap(n);
        for (auto e: edges)
                heap[e.dst].push(e);
        double score = 0;
        vector<int> seen(n, -1);
        seen[r] = r;
        for (int s = 0; s < n; ++s) {
                vector<int> path;
                for (int u = s; seen[u] < 0;) {</pre>
                        path.push_back(u);
                         seen[u] = s;
                         if (heap[u].empty()) return INF;
                         edge min_e = heap[u].top();
                         score += min_e.weight;
                         heap[u].add(-min_e.weight);
                        heap[u].pop();
                         int v = uf.root(min_e.src);
```

4.8 Directed minimum spanning tree O(nm)

```
/*
        GETS:
                call make_graph(n) at first
                you should use add_edge(u,v,w) and
                add pair of vertices as edges (vertices are 0..n-1)
        GIVES:
                output of dmst(v) is the minimum arborescence with
                     root v in directed graph
                 (-1 if it hasn't a spanning arborescence with root v)
        O(mn)
#include <bits/stdc++.h>
using namespace std;
const int INF = 2e7;
struct MinimumAborescense
        int n;
        struct edge {
                int src, dst;
                int weight;
        vector<edge> edges;
        void make_graph(int _n) {
                n=\underline{n};
                edges.clear():
        void add_edge(int u, int v, int w) {
                edges.push_back({u, v, w});
        int dmst(int r) {
```

```
int N = n;
for (int res = 0; ;) {
        vector<edge> in(N, {-1,-1,(int)INF});
        vector<int> C(N, -1);
        for (auto e: edges)
                if (in[e.dst].weight > e.weight)
                        in[e.dst] = e;
        in[r] = \{r, r, 0\};
        for (int u = 0; u < N; ++u) { // no comming
            edge ==> no aborescense
                if (in[u].src < 0) return -1;</pre>
                res += in[u].weight;
        vector<int> mark(N, -1); // contract cycles
        int index = 0:
        for (int i = 0; i < N; ++i) {
                if (mark[i] != -1) continue;
                int u = i:
                while (mark[u] == -1) {
                        mark[u] = i;
                        u = in[u].src;
                if (mark[u] != i || u == r) continue;
                for (int v = in[u].src; u != v; v = in
                     [v].src) C[v] = index;
                C[u] = index++;
        if (index == 0) return res; // found
            arborescence
        for (int i = 0; i < N; ++i) // contract</pre>
                if (C[i] == -1) C[i] = index++;
        vector<edge> next;
        for (auto &e: edges)
                if (C[e.src] != C[e.dst] && C[e.dst]
                        next.push_back({C[e.src], C[e.
                             dst], e.weight - in[e.dst
                             ].weight});
        edges.swap(next);
        N = index; r = C[r];
```

4.9 Dominator tree

};

```
tree[i].clear();
                 sdom[i] = idom[i] = dsu[i] = label[i] = i;
                arr[i] = -1;
        cnt = 0;
void add_edge(int u, int v)
        adj[u].push_back(v);
void dfs(int v)
        arr[v] = cnt;
        rev[cnt] = v;
        cnt++;
        for (int i = 0; i < adj[v].size(); i++)</pre>
                int u = adj[v][i];
                if (arr[u] == -1)
                         dfs(u);
                         par[arr[u]] = arr[v];
                 radj[arr[u]].push_back(arr[v]);
int find(int v, int x = 0)
        if (dsu[v] == v)
                return (x ? -1 : v);
        int u = find(dsu[v], x + 1);
        if (u < 0)
                return v;
        if (sdom[label[dsu[v]]] < sdom[label[v]])</pre>
                label[v] = label[dsu[v]];
        dsu[v] = u;
        return (x ? u : label[v]);
void merge(int u, int v)
        dsu[v] = u;
void build(int root)
        dfs(root);
        int n = cnt;
        for (int v = n - 1; v >= 0; v--)
                 for (int i = 0; i < radj[v].size(); i++)</pre>
                         int u = radj[v][i];
                         sdom[v] = min(sdom[v], sdom[find(u)]);
                if (v > 0)
                         bucket[sdom[v]].push_back(v);
                for (int i = 0; i < bucket[v].size(); i++)</pre>
                         int u = bucket[v][i];
                         int w = find(u);
                         if (sdom[u] == sdom[w])
                                 idom[u] = sdom[u];
```

5 Combinatorics

5.1 LP simplex

```
// Two-phase simplex algorithm for solving linear programs of the form
11
       maximize
                    C^T X
11
       subject to Ax \le b
11
                    x >= 0
11
// INPUT: A -- an m x n matrix
          b -- an m-dimensional vector
          c -- an n-dimensional vector
          x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
           above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
  int m, n;
```

```
VI B, N;
VVD D;
LPSolver(const VVD &A, const VD &b, const VD &c) :
 m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
  for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] =
      A[i][i];
  for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1; D[i][n + i]
       1 = b[i]; 
  for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
 N[n] = -1; D[m + 1][n] = 1;
void Pivot(int r, int s) {
 double inv = 1.0 / D[r][s];
 for (int i = 0; i < m + 2; i++) if (i != r)
    for (int j = 0; j < n + 2; j++) if (j != s)
     D[i][j] = D[r][j] * D[i][s] * inv;
  for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;
 for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;</pre>
 D[r][s] = inv;
 swap(B[r], N[s]);
bool Simplex(int phase) {
 int x = phase == 1 ? m + 1 : m;
 while (true) {
    int s = -1;
    for (int j = 0; j \le n; j++) {
      if (phase == 2 \&\& N[j] == -1) continue;
      if (s == -1 \mid | D[x][\dot{j}] < D[x][s] \mid | D[x][\dot{j}] == D[x][s] && N[\dot{j}]
           < N[s]) s = j;
    if (D[x][s] > -EPS) return true;
    int r = -1;
    for (int i = 0; i < m; i++) {
      if (D[i][s] < EPS) continue;</pre>
      if (r == -1 \mid \mid D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s]
        (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r][s]) && B[i] <
             B[r]) r = i;
    if (r == -1) return false:
    Pivot(r, s);
DOUBLE Solve(VD &x) {
 int r = 0;
  for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1]) r = i;
 if (D[r][n + 1] < -EPS) {
    Pivot(r, n);
    if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -</pre>
        numeric limits<DOUBLE>::infinity();
    for (int i = 0; i < m; i++) if (B[i] == -1) {
      int s = -1;
      for (int j = 0; j \le n; j++)
        if (s == -1 \mid | D[i][j] < D[i][s] \mid | D[i][j] == D[i][s] && N[
            j] \langle N[s] \rangle s = j;
      Pivot(i, s);
```

```
if (!Simplex(2)) return numeric limits<DOUBLE>::infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
    return D[m][n + 1];
};
int main() {
  const int m = 4;
  const int n = 3:
  DOUBLE _A[m][n] = {
    \{6, -1, 0\},
    \{-1, -5, 0\},
    { 1, 5, 1 },
    \{-1, -5, -1\}
  DOUBLE _b[m] = { 10, -4, 5, -5 };
  DOUBLE _{c[n]} = \{ 1, -1, 0 \};
  VVD A(m);
  VD b(\underline{b}, \underline{b} + m);
  VD c(c, c+n);
  for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);</pre>
  LPSolver solver(A, b, c);
  VD x;
  DOUBLE value = solver.Solve(x);
  cerr << "VALUE: " << value << endl; // VALUE: 1.29032</pre>
  cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
  for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
  cerr << endl:
  return 0;
```

5.2 FFT

```
const int LG = 20; // IF YOU WANT TO CONVOLVE TWO ARRAYS OF LENGTH N
    AND M CHOOSE LG IN SUCH A WAY THAT 2^LG > n + m
const int MAX = 1 << LG;

struct point
{
    double real, imag;
    point(double _real = 0.0, double _imag = 0.0)
    {
        real = _real;
        imag = _imag;
    }
};
point operator + (point a, point b)
{
    return point(a.real + b.real, a.imag + b.imag);
}
point operator - (point a, point b)
{
    return point(a.real - b.real, a.imag - b.imag);
}</pre>
```

```
point operator * (point a, point b)
        return point(a.real * b.real - a.imag * b.imag, a.real * b.
            imag + a.imag * b.real);
void fft(point *a, bool inv)
        for (int mask = 0; mask < MAX; mask++)</pre>
                int rev = 0;
                for (int i = 0; i < LG; i++)
                        if ((1 << i) & mask)
                                 rev = (1 << (LG - 1 - i));
                if (mask < rev)</pre>
                        swap(a[mask], a[rev]);
        for (int len = 2; len <= MAX; len *= 2)</pre>
                double ang = 2.0 * M_PI / len;
                if (inv)
                        ang *=-1.0;
                point wn(cos(ang), sin(ang));
                for (int i = 0; i < MAX; i += len)
                         point w(1.0, 0.0);
                         for (int j = 0; j < len / 2; j++)
                                 point t1 = a[i + j] + w * a[i + j +
                                     len / 21;
                                 point t2 = a[i + j] - w * a[i + j +
                                     len / 2];
                                 a[i + j] = t1;
                                 a[i + j + len / 2] = t2;
                                 w = w * wn;
        if (inv)
                for (int i = 0; i < MAX; i++)
                         a[i].real /= MAX;
                        a[i].imaq /= MAX;
```

5.3 NTT

```
b >>= 1;
                 a = 1LL * a * a % MOD;
        return ans;
void ntt(int *a, bool inv)
        for (int mask = 0; mask < MAX; mask++)</pre>
                 int rev = 0;
                 for (int i = 0; i < LG; i++)</pre>
                         if ((1 << i) & mask)
                                 rev = (1 << (LG - 1 - i));
                 if (mask < rev)</pre>
                         swap(a[mask], a[rev]);
        for (int len = 2; len <= MAX; len *= 2)</pre>
                 int wn = bpow(ROOT, MAX / len);
                 if (inv)
                         wn = bpow(wn, MOD - 2);
                 for (int i = 0; i < MAX; i += len)</pre>
                         int w = 1;
                         for (int j = 0; j < len / 2; j++)
                                  int l = a[i + j];
                                  int r = 1LL * w * a[i + j + len / 2] %
                                  a[i + j] = (l + r);
                                  a[i + j + len / 2] = 1 - r + MOD;
                                  if (a[i + j] >= MOD)
                                          a[i + j] -= MOD;
                                  if (a[i + j + len / 2] >= MOD)
                                          a[i + j + len / 2] -= MOD;
                                  w = 1LL * w * wn % MOD;
        if (inv)
                 int x = bpow(MAX, MOD - 2);
                 for (int i = 0; i < MAX; i++)</pre>
                         a[i] = 1LL * a[i] * x % MOD;
```

5.4 Base Vectors in Z2

```
for (int i=B-1; i>=0; --i) if (x >> i & 1)
            a[i] = x;
                         for(int j = i - 1; j >= 0; j--) if(a[j] >> i & 5.6 Stirling 1
                              1) a[j] = x;
            return;
    int size()
        int cnt = 0;
        for(int i=0; i<B; ++i) if(a[i]) ++cnt;</pre>
        return cnt:
} ;
```

5.5Gaussian Elimination

```
const int N = 505, MOD = 1e9 + 7;
typedef vector <ll> vec;
11 pw(ll a, ll b) {
        if(!b)
                return 1:
        11 x = pw(a, b/2);
        return x * x % MOD * (b % 2 ? a : 1) % MOD;
11 inv(11 x) { return pw(x, MOD - 2); }
bool solve() {
        int n = in():
        vector <vec> matrix(n);
        for (int i = 0; i < n; i++)
                for (int j = 0; j < n; j++) {
                        matrix[i].push_back((in() % MOD + MOD) % MOD);
                }
        11 \text{ res} = 1;
        for (int i = 0; i < n; i++) {
                int ind = -1;
                for(int row = i; row < n; row++)</pre>
                         if (matrix[row][i])
                                 ind = row;
                if(ind == -1) {
                         res = 0;
                         break;
                if(i != ind)
                         res = (MOD - res) %MOD;
                matrix[i].swap(matrix[ind]);
                res = res * matrix[i][i] % MOD;
                11 inverse = inv(matrix[i][i]);
                for(int row = i + 1; row < n; row++) {</pre>
                         11 z = matrix[row][i] * inverse % MOD;
                         for (int j = 0; j < n; j++)
                                 matrix[row][j] = (matrix[row][j] % MOD
                                       - matrix[i][i]*z % MOD + MOD) %
                                      MOD;
        cout << res << endl;
```

```
#include <bits/stdc++.h>
using namespace std:
typedef long long 11;
#define pb push_back
const int mod = 998244353;
const int root = 15311432:
const int root_1 = 469870224;
const int root_pw = 1 << 23;</pre>
const int N = 400004;
vector<int> v[N];
11 modInv(ll a, ll mod = mod) {
        11 \times 0 = 0, \times 1 = 1, \times 0 = 0, \times 1 = 0;
        while(r1){
                 11 q = r0 / r1;
                 x0 -= q * x1; swap(x0, x1);
                 r0 = q * r1; swap(r0, r1);
        return x0 < 0 ? x0 + mod : x0;
void fft (vector<int> &a, bool inv) {
        int n = (int) a.size();
        for (int i=1, j=0; i<n; ++i) {</pre>
                 int bit = n >> 1;
                 for (; j>=bit; bit>>=1)
                         j -= bit;
                 j += bit;
                 if (i < j)
                         swap (a[i], a[j]);
        for (int len=2; len<=n; len<<=1) {</pre>
                 int wlen = inv ? root 1 : root;
                 for (int i=len; i<root_pw; i<<=1)</pre>
                         wlen = int (wlen * 111 * wlen % mod);
                 for (int i=0; i<n; i+=len) {</pre>
                         int w = 1;
                         for (int j=0; j<len/2; ++j) {</pre>
                                  int u = a[i+j], v = int (a[i+j+len/2]
                                        * 111 * w % mod);
                                  a[i+j] = u+v < mod ? u+v : u+v-mod;
                                  a[i+j+len/2] = u-v >= 0 ? u-v : u-v+
                                      mod:
                                  w = int (w * 111 * wlen % mod);
        if(inv) {
                 int nrev = modInv(n, mod);
                 for (int i=0; i<n; ++i)</pre>
                         a[i] = int (a[i] * 111 * nrev % mod);
        }
```

```
void pro(const vector<int> &a, const vector<int> &b, vector<int> &res)
        vector<int> fa(a.begin(), a.end()), fb(b.begin(), b.end());
        while (n < (int) max(a.size(), b.size())) n <<= 1;</pre>
        n <<= 1;
        fa.resize (n), fb.resize (n);
        fft(fa, false), fft (fb, false);
        for (int i = 0; i < n; ++i)
                fa[i] = 1LL * fa[i] * fb[i] % mod;
        fft (fa, true);
        res = fa;
int S(int n, int r) {
        int nn = 1;
        while(nn < n) nn <<= 1;
        for(int i = 0; i < n; ++i) {
                v[i].push back(i);
                v[i].push_back(1);
        for(int i = n; i < nn; ++i) {
                v[i].push_back(1);
        for (int j = nn; j > 1; j >>= 1) {
                int hn = j >> 1;
                for(int i = 0; i < hn; ++i) {
                        pro(v[i], v[i + hn], v[i]);
        return v[0][r];
int fac[N], ifac[N], inv[N];
void prencr() {
        fac[0] = ifac[0] = inv[1] = 1;
        for(int i = 2; i < N; ++i)
                inv[i] = mod - 1LL * (mod / i) * inv[mod % i] % mod;
        for (int i = 1; i < N; ++i) {
                fac[i] = 1LL * i * fac[i - 1] % mod;
                ifac[i] = 1LL * inv[i] * ifac[i - 1] % mod;
int C(int n, int r) {
        return (r \ge 0 \&\& n \ge r)? (1LL * fac[n] * ifac[n - r] % mod
            * ifac[r] % mod) : 0;
int main(){
        prencr();
        int n, p, q;
        cin >> n >> p >> q;
```

```
ll ans = C(p + q - 2, p - 1);
ans *= S(n - 1, p + q - 2);
ans %= mod;
cout << ans;</pre>
```

5.7 Chinese remainder

```
long long GCD(long long a, long long b) { return (b == 0) ? a : GCD(b,
inline long long LCM(long long a, long long b) { return a / GCD(a, b)
inline long long normalize(long long x, long long mod) { x %= mod; if
    (x < 0) x += mod; return x; }
struct GCD_type { long long x, y, d; };
GCD_type ex_GCD(long long a, long long b)
    if (b == 0) return {1, 0, a};
    GCD_type pom = ex_GCD(b, a % b);
    return {pom.y, pom.x - a / b * pom.y, pom.d};
int testCases;
int t:
long long r[N], n[N], ans, lcm;
int main()
    cin >> t;
    for(int i = 1; i \le t; i++) cin >> r[i] >> n[i], normalize(r[i], n
    ans = r[1];
    lcm = n[1];
    for (int i = 2; i \le t; i++)
        auto pom = ex_GCD(lcm, n[i]);
        int x1 = pom.x;
        int d = pom.d;
        if((r[i] - ans) % d != 0) return cerr << "No solutions" <<</pre>
        ans = normalize(ans + x1 * (r[i] - ans) / d % (n[i] / d) * lcm
            , lcm * n[i] / d);
        lcm = LCM(lcm, n[i]); // you can save time by replacing above
            lcm * n[i] /d by <math>lcm = lcm * n[i] / d
    cout << ans << " " << lcm << endl:
    return 0:
```

5.8 Stirling 2

$$\left\{\begin{array}{c} \mathbf{n} \\ \mathbf{k} \end{array}\right\} = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} \binom{k}{j} j^{n}$$

5.9 Popular LP

BellmanFord:

maximize X_n

 $X_1 = 0$

and for each edge (v->u) and weight w:

 $X_u - X_v \le w$

Flow:

maximize Σf_{out} (where out is output edges of vertex 1)

for each vertex (except 1 and n):

 $\Sigma f_{in} - \Sigma f_{out} = 0$ (where in is input edges of v and out is output edges of v)

Dijkstra(IP):

minimize $\Sigma z_i * w_i$

for each edge (v->u) and weight w:

 $0 \le z_i \le 1$

and for each ST-cut which vertex 1 is in S and vertex n is in T:

 $\Sigma z_e \geq 1$ (for each edge e from S to T)

5.10 Duality of LP

primal: Maximize c^Tx subject to $Ax \leq b, x \geq 0$ dual: Minimize b^Ty subject to $A^Ty \geq c, y \geq 0$

5.11 Extended catalan

number of ways for going from 0 to A with k moves without going to -B:

$$\binom{k}{\frac{A+k}{2}} - \binom{k}{\frac{2B+A+k}{2}}$$

5.12 Find polynomial from it's points

$$P(x) = \sum_{i=1}^{n} y_i \prod_{j=1, j \neq i}^{n} \frac{x - x_j}{x_i - x_j}$$

Useful formulas

 $\binom{n}{k} = \frac{n!}{k!(n-k)!}$ objects out of n- number of ways to choose k

 $\binom{n+k-1}{k-1}$ — number of ways to choose k objects out of n with repetitions $\binom{n}{n}$ — Stirling numbers of the first kind; number of

permutations of n elements with k cycles ${n+1\brack m}=n{n\brack m}+{n\brack m-1}$ ${n \brack m}$ — Stirling numbers of the first kind; number of

$${\binom{n+1}{m}} = n {\binom{n}{m}} + {\binom{n}{m-1}}$$
$$(x)_n = x(x-1)\dots x - n + 1 = \sum_{k=0}^n (-1)^{n-k} {\binom{n}{k}} x^k$$

of partitions of set $1, \ldots, n$ into k disjoint subsets. ${n+1 \brace m} = k \begin{Bmatrix} n \end{Bmatrix} + \begin{Bmatrix} n \cr k-1 \end{Bmatrix}$ ${n \brace m} - ext{Stirling numbers of the second kind; number}$

$$\begin{Bmatrix} n+1 \\ m \end{Bmatrix} = k \begin{Bmatrix} n \\ k \end{Bmatrix} + \begin{Bmatrix} n \\ k-1 \end{Bmatrix}$$

$$\sum_{k=0}^{n} {n \brace k} (x)_k = x^n$$

$$C_n = \frac{1}{n+1} {2n \choose n}$$
 — Catalan numbers $C(x) = \frac{1-\sqrt{1-4x}}{2x}$

If
$$a_n = \sum_{k=0}^{n} {n \choose k} b_k$$
, then $b_n = \sum_{k=0}^{n} (-1)^{n-k} {n \choose k} a_k$

•
$$a = (1, x, x^2, ...), b = (1, (x+1), (x+1)^2, ...)$$

•
$$a_i = i^k, b_i = {n \brace i} i!$$

Burnside's lemma

shifts of array, rotations and symmetries of $n \times n$ matrix, ...) Let G be a group of action on set X (Ex.: cyclic

action f that transforms x to y: f(x) = y. Call two objects x and y equivalent if there is an

The number of equivalence classes then can be calculated as follows: $C = \frac{1}{|G|} \sum_{f \in G} |X^f|$, where X^f

is the set of fixed points of $f: X^f = \{x | f(x) = x\}$

Generating functions

 $a_0, a_1, \dots, a_n, \dots$ is $A(x) = \sum_{i=1}^{\infty} a_i x^i$ Ordinary generating function (o.g.f.) for sequence Exponential generating function (e.g.f.)

sequence $a_0, a_1, \dots, a_n, \dots$ is $A(x) = \sum_{n=0}^{\infty} a_i x^i$

 $B(x) = A'(x), b_{n-1} = n \cdot a_n$

with FFT using $\widetilde{a_n} = \frac{a_n}{n!}$) $c_n = \sum$ $c_n =$ $\sum_{k=0}^{n} a_k b_{n-k} \text{ (o.g.f. convolution)}$ $\sum_{k} {n \choose k} a_k b_{n-k}$ (e.g.f. convolution, compute

General linear recurrences

also can compute all a_n with Divide-and-Conquer algorithm in $O(n\log^2 n)$. If $a_n =$ $\sum_{k=1} b_k a_{n-k}, \text{ then } A(x) =$

Inverse polynomial modulo x'

Given A(x), find B(x) ; $A(x)B(x) = 1 + x^l \cdot Q(x)$ for some Q(x) such

- 1. Start with $B_0(x) = \frac{1}{a_0}$
- 2. Double $B_{k+1}(x) = (-B_k(x)^2 A(x) + 2B_k(x)) \mod x^{2^{k+1}}$

Fast subset convolution

Given array a_i of size 2^k , calculate $b_i =$

for
$$b = 0..k-1$$

for $i = 0..2^k-1$
if $(i & (1 << b)) != 0:$
 $a[i + (1 << b)] += a[i]$

Hadamard transform

size $2 \times 2 \times \ldots \times 2$, calculate FFT of that array: Treat array a of size 2^k as k-dimentional array

6 Constants

6.1 Number of primes

30: 10 60: 17 100: 25 1000: 168 10000: 1229 100000: 9592 1000000: 78498 10000000: 664579

6.2 Factorials

1: 1 2: 2 3: 6 4: 24 5: 120 6: 720 7: 5040 8: 40320 9: 362880 10: 362880 11: 39916800 12: 479001600 13: 6227020800 14: 87178291200 15: 1307674368000

6.3 Powers of 3

1: 3 2: 9 3: 27 4: 81 5: 243 6: 729 7: 2187 8: 6561 9: 19683 10: 59049 11: 177147 12: 531441 13: 1594323 14: 4782969 15: 14348907 16: 43046721 17: 129140163 18: 387420489 19: 1162261467 20: 3486784401

$6.4 \quad C(2n,n)$

1: 2

2: 6 3: 20 4: 70 5: 252 6: 924 7: 3432 8: 12870 9: 48620 10: 184756 11: 705432 12: 2704156 13: 10400600 14: 40116600 15: 155117520

6.5 Most divisor

<= 1e2: 60 with 12 divisors <= 1e3: 840 with 32 divisors <= 1e4: 7560 with 64 divisors <= 1e5: 83160 with 128 divisors <= 1e6: 720720 with 240 divisors <= 1e7: 8648640 with 448 divisors <= 1e8: 73513440 with 768 divisors <= 1e9: 735134400 with 1344 divisors <= 1e10: 6983776800 with 2304 divisors <= 1e11: 97772875200 with 4032 divisors <= 1e12: 963761198400 with 6720 divisors <= 1e13: 9316358251200 with 10752 divisors <= 1e14: 97821761637600 with 17280 divisors <= 1e15: 866421317361600 with 26880 divisors <= 1e16: 8086598962041600 with 41472 divisors <= 1e17: 74801040398884800 with 64512 divisors <= 1e18: 897612484786617600 with 103680 divisors