



Singularity

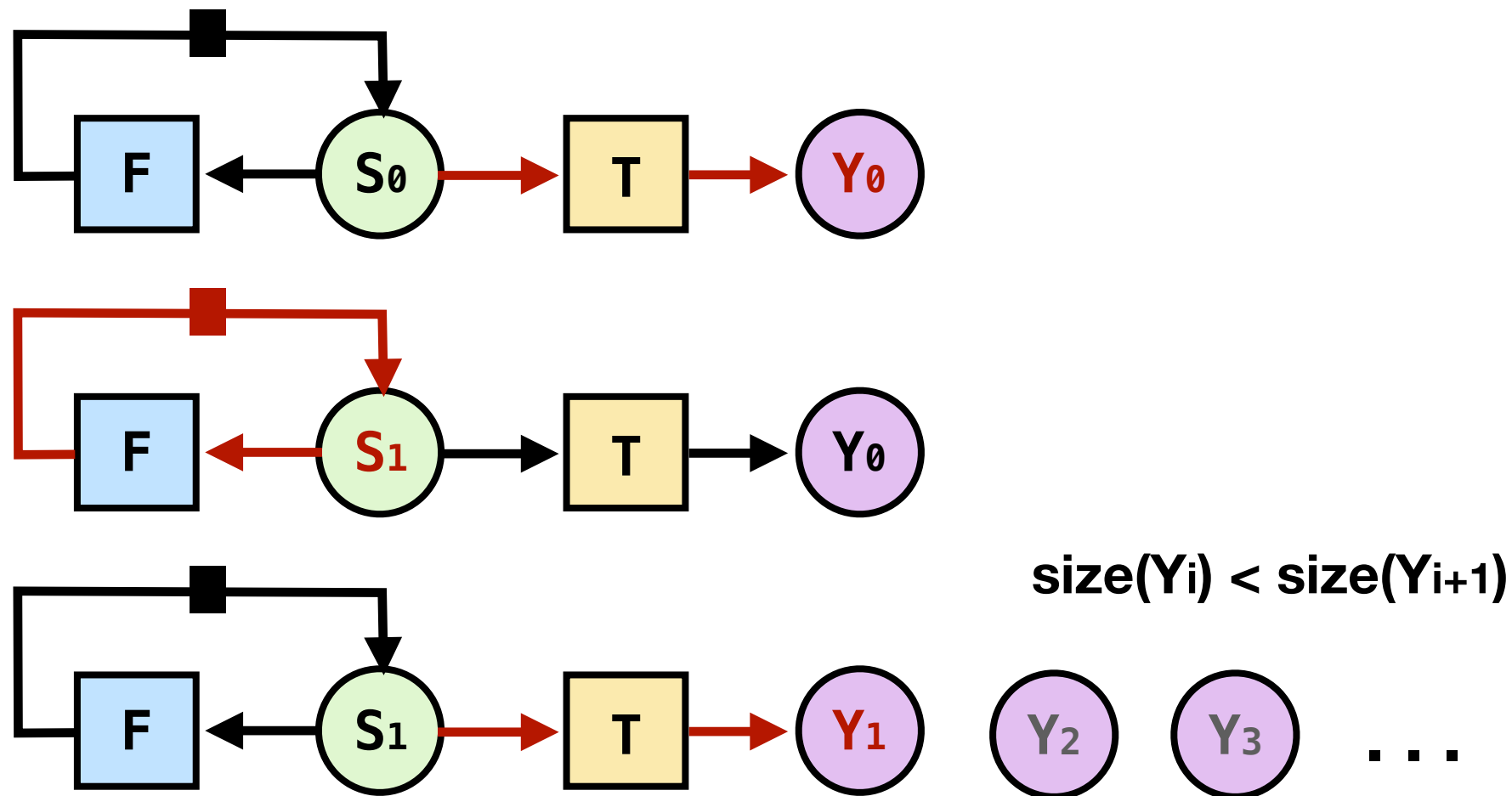
The closer your program gets to a singularity, the slower it runs.
— *general theory of relativity*

Observations

- **Many complexity vulnerabilities can be triggered by inputs with special *patterns***
 - a sorted array
 - a graph with negative-weight loops
 - objects with same hash value
 - . . .
- **Fuzzing large inputs with such properties can be challenging because of the huge search space**
 - the space of possible inputs grows exponentially with size
- **Instead of concrete inputs, we can search for *input patterns***
- **We developed *Singularity*, a pattern fuzzer for complexity vulnerabilities**

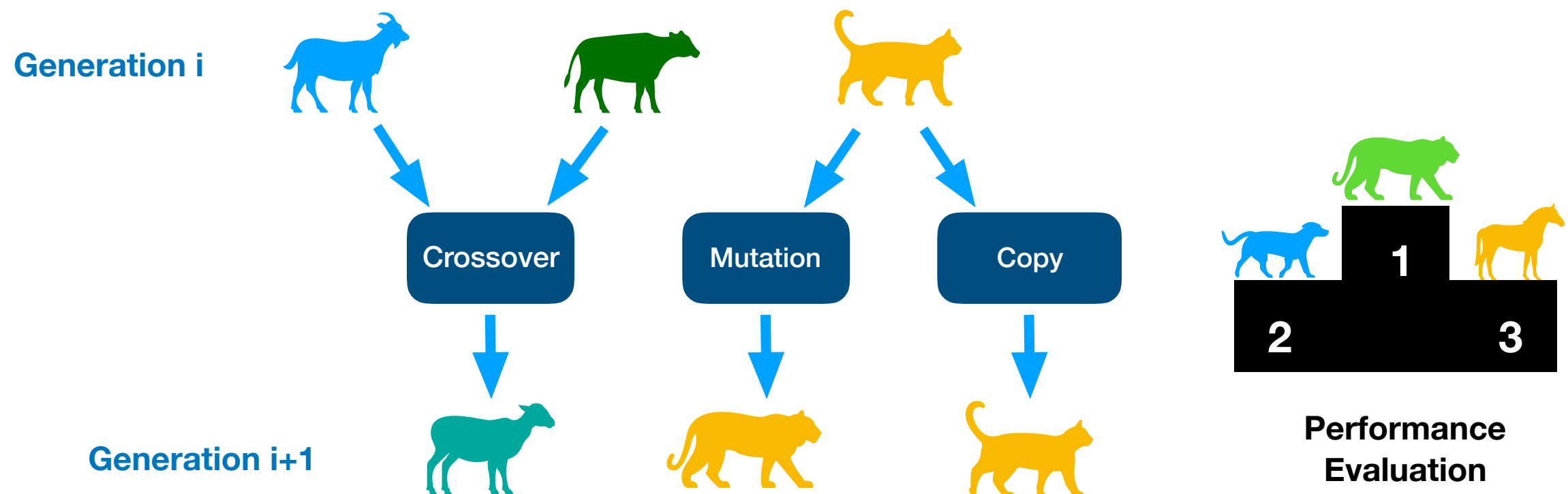
Input Patterns

- **We can represent input patterns as programs**
 - Starting state **S**₀, Transformer **T**, Updater **F**
 - One pattern can be used to construct inputs of *arbitrarily large size*
 - Represent **S**₀, **T**, **F** as DSL expressions and search for them instead
 - **S** and **Y** can be tuples containing multiple values



Resource Usage Maximization

- **Search for patterns that can maximize resource usage**
 - To evaluate a pattern, we use it to generate an input of some specified size, and measure the corresponding resource usage (running time, space usage, e.t.c.)
 - A pattern with the highest resource usage will be able to trigger any potential complexity vulnerability
 - To search for patterns with maximal resource usage, we use genetic programming as the optimization algorithm



Solving SearchableBlog

- In the recent engagement, Singularity helped us solve the *SearchableBlog* problem (Q27)
 - We need to find a vector **V** that maximizes the time used by the following code

```
Vector v1 = MatrixRoutines.processVector(V, 0.999);  
Matrix A = MatrixRoutines.concatenateColumn(ajMatrix, v1);  
estimateStationaryVector(A, bVector, 1.0E-7);
```

- *estimateStationaryVector* is a numerical iterative algorithm whose running time has complex dependency on the initial vector **V** and the given matrix *ajMatrix*
- Singularity found the correct pattern within 5min and used it to construct following vector, whose length is 730!

[0,0,0,0,0,0,...,0,1]

- Initial state: empty vector **[]**
- Updater: **λv. append(v, 0)**
- Transformer: **λv. append(v, 1)**
- First few elements:

[1], [0,1], [0,0,1], [0,0,0,1], ..

Solving Braidit

- **Singularity also helped us solve the *Braidit* problem (Q17 Q22)**
 - To trigger the time or space vulnerability, we needed to construct inputs that can cause **Weave.isEquivalent** method run a long time
 - The inputs correspond to two strings (**S1**, **S2**) and one integer (**n**)
 - We let Singularity maximize the running time of the following code

```
val w1 = new Weave(S1, n)
val w2 = new Weave(S2, n)
w1.isEquivalent(w2)
```

- It found an efficient pattern within 5min, which gives the following inputs for string length less than 50

```
S1 = "thYPGxofWNEvmdULCtkbSJArIZQHypgX0FwneVMDuLcTKBsja"
S2 = "qykYPGxofWNEvmdULCtkbSJArIZQHypgX0FwneVMDuLcTKBsja"
n = 27
```

- Note that the two strings are not equivalent but share a common part. Patterns like this can be challenging for other non-pattern-based fuzzers.

How it works

- Suppose we want to maximize the running time of a QuickSort implementation which always select the middle element in the array as pivot

```
def quick_sort(xs: list):  
    if length(xs) <= 1:  
        return xs  
    pivot = xs[length(xs) // 2]  
    left, middle, right = [], [], []  
    for x in xs:  
        if x==pivot:  
            middle.append(x)  
        elif x<pivot:  
            left.append(x)  
        else:  
            right.append(x)  
    left = quick_sort(left)  
    right = quick_sort(right)  
    return concat(left, middle, right)
```

How it works

- **Suppose we want to maximize the running time of a QuickSort implementation which always select the middle element in the array as pivot**

- One way to achieve this is using the following pattern:

`[], [1,0], [3,1,0,2], [5,3,1,0,2,4], ...`

- Which can be represented as
 - initial state: `[]`
 - transformer: **identity**
 - updater: **`$\lambda x.append(prepend(length(x) + 1, x), length(x))$`**
- To find the desired pattern, Singularity starts with a population of randomly-generated input patterns
- For simplicity, assume init states and transformers are already given
- It is very unlikely that Singularity will discover the correct updater in the this first generation, but it will discover useful sub-optimal updaters. e.g.

```
f1 =  $\lambda x.append(x, length(x))$ 
f2 =  $\lambda x.prepend(length(x), x)$ 
..
```


How it works

- Desired updater:

$\lambda x. \text{append}(\text{prepend}(\text{length}(x) + 1, x), \text{length}(x))$

- Updaters found in first population:

```
f1 =  $\lambda x. \text{append}(x, \text{length}(x))$   
f2 =  $\lambda x. \text{prepend}(\text{length}(x), x)$   
..
```

- For the next population, we randomly pick input patterns with high resource usage from the previous population
- f1 and f2 are likely to be selected because they have higher than average resource usage
- We then use these input patterns to generate a new population by combining them using genetic operators, such as mutation and crossover
- For example, we can obtain the following updater f3 from f1 and f2 by using the crossover operation

$f3 = \lambda x. \text{append}(\text{prepend}(\text{length}(x), x), \text{length}(x))$

- f3 has higher resource usage than f1 and f2, but still sub-optimal

How it works


- Desired updater:

$\lambda x. \text{append}(\text{prepend}(\text{length}(x) + 1, x), \text{length}(x))$

- Updaters found:

```
f1 =  $\lambda x. \text{append}(x, \text{length}(x))$   
f2 =  $\lambda x. \text{prepend}(\text{length}(x), x)$   
f3 =  $\lambda x. \text{append}(\text{prepend}(\text{length}(x), x), \text{length}(x))$   
..
```

- We continue the process of generating new populations and monitor their average performance
- In general, average performance will keep increasing over generations
- At some point, Singularity will generate the desired updater by mutating the sub-expression in f3:



```
f3 =  $\lambda x. \text{append}(\text{prepend}(\text{length}(x), x), \text{length}(x))$   
f4 =  $\lambda x. \text{append}(\text{prepend}(\text{length}(x)+1, x), \text{length}(x))$ 
```

- Hence, we found the pattern and can use it to construct an attack vector of arbitrarily large size