

# 2nd year paper

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## outline

1. Intro and Lit Review - Work on this section last. For the literature review I would focus on Walker (1981) and the Varian paper I gave you earlier this year.
2. Model - Define economic environment and assumptions on the utility functions/ production function/ etc. Use my draft as your main guide.
3. Generalized Lindahl Equilibrium - Define Generalized Lindahl Equilibria and provide an example (Rawlsian).
4. Generalized Walker Mechanism - define the mechanism and give your main result.
5. Discussion - I would talk about what you still need to do. For example, the other half of the mechanism result (all G. Lindahl Outcomes are achievable as NE of the game). Also, the set of NE is very large, we need to find a way to enforce a single outcome.

## To do

- Lit Review
- Reference work done jointly w/ Van Essen

## 1 Introduction

## 2 Model

A fairly simple economic environment will provide an adequate setting. There will be  $i = 1, \dots, n \geq 2$  individuals in the economy, each represented by a convex consumption set  $C_i = \mathbb{R}_i^2$  and an initial endowment of a private good  $\omega_i > 0$ . They consume bundles of a public good  $x$  and a private good  $y_i$ , and in doing so gain utility according to utility functions of the following form.

$$u_i(x, y_i) = \int_0^x v_i(t) dt + y_i \quad (1)$$

This functional form simplifies later analysis, as  $\frac{\partial u_i}{\partial x} = v_i(x) > 0$  and  $\frac{\partial^2 u_i}{\partial x^2} = v'_i(x) < 0$  for  $x \in \mathbb{R}_+$ .  $v_i$  is the the marginal rate of substitution between the public good and the private good for individual  $i$ .

The public good is produced in a constant returns to scale production process that uses the private good as an input: for each unit of the public good produced,  $c$  of the private good must be sacrificed. If  $\Omega = \sum_i \omega_i$ , then the most public good that can be produced is  $\frac{\Omega}{c}$ . An allocation in this economy is then an  $(n+1)$ -tuple  $(x, y_1, \dots, y_n) \in \mathbb{R}_+^{n+1}$ .

We shall assume that there is a unique Pareto optimal allocation in the economy, such that a unique level of provision of the public good is Pareto optimal. This is to say that there is a level of the public good that satisfies the Samuelson marginal condition:  $\sum_i v_i(x^{PO}) = c$  for  $x \in [0, \frac{\Omega}{c}]$ . To ensure this, we assume  $\sum_i v_i(0) > c$  and  $\sum_i v_i(\frac{\Omega}{c}) < c$ . In addition, we shall also assume that the initial endowments of the private good are sufficiently large for feasible provision of  $x^{PO}$ ,  $\omega_i \geq v_i(0)x^{PO}$ .

## 2.1 Classical Lindahl Equilibrium

An unregulated market cannot efficiently provide a public good. A firm cannot earn revenue by selling the right to consume a public good because it cannot prevent the consumption of a public good by those who have not paid for it. Consumers face a strong incentive not pay for the provision of a public good because once the public good is provided the consumption of it by one individual does not prevent the consumption of it by another. These are the usual non-rival and non-excludable characteristics of a public good. Some sort of market regulation is required to elicit and finance the efficient provision of the public good. Individuals can be taxed and production can be centralized by a government. One theoretical equilibrium that allows a Pareto optimal amount of the public good to be provided is a classical Lindahl equilibrium.

A classical Lindahl equilibrium uses individualized taxes, known as Lindahl prices, to raise revenue for the provision of a public good. It finds the socially optimal level of the public good and then charges each consumer their marginal valuation of the public good at the socially optimal level. With constant marginal costs of provision and downward sloping demand curves, a classical Lindahl equilibrium ensures that it is individually rational for individuals to pay the tax levied upon them and that the sum of such taxes exactly covers the provision of the public good.

For example, consider an instance of the economic setting outlined above with  $n = 2$ ,  $c = 5$ , and the following utility functions for the two individuals

### Classical Lindahl Equilibrium

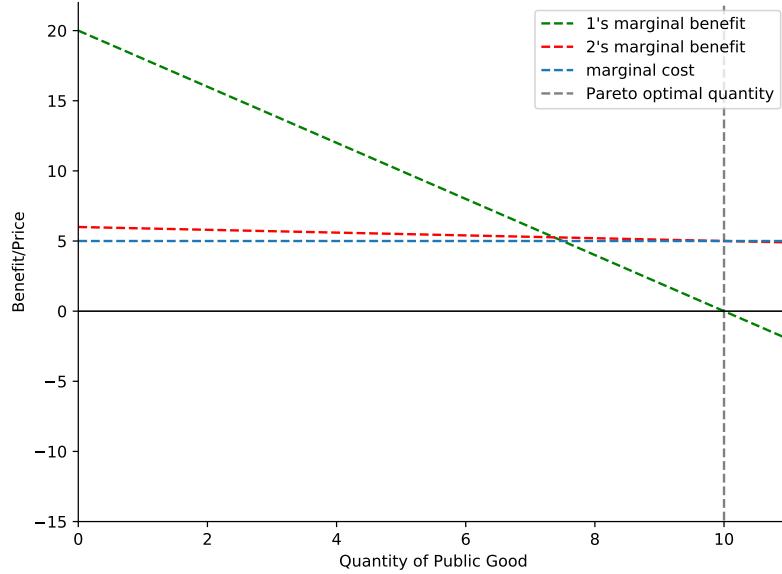


Figure 1: An example of a classical Lindahl Equilibrium for  $n = 2$  consumers. The vertical grey dotted line is located at the Pareto optimal level of the public good. Consumer 1 is charged nothing, while consumer 2 is charged 5 units of the private good.

$$u_1(x, y_1) = \int_0^x (20 - 2t) dt + y_1$$

$$u_2(x, y_2) = \int_0^x (6 - 0.1t) dt + y_2$$

And so

$$v_1(x) = 20 - 2x$$

$$v_2(x) = 6 - 0.1x$$

This setting is shown graphically in figure 1. To construct a classical Lindahl equilibrium we would first find the Pareto optimal level of the public good,  $x^{PO}$ , via the Samuelson marginal condition,  $\sum_i v_i(x^{PO}) = c$ . Then, because the private good is a numeraire, the tax  $t_i$  levied upon each individual would be equal to their marginal rate of substitution evaluated at  $x^{PO}$ ,  $t_i = v_i(x^{PO})$ .

Here, the Pareto optimal amount of the public good is  $x^{PO} = 10$ . The marginal benefit of the public good at this level is 0 for consumer 1 and 5 for consumer 2, and so  $t_1 = 0$  and  $t_2 = 5$ .

We can see that the classical Lindahl equilibrium leaves both consumers better off than they would be if no public good was provided, so the outcome is indeed individually rational. It also raises tax revenue that covers the cost of providing the public good. But, it is also clear that consumer 1 is left far better off than consumer 2, even though consumer 2 incurs all of the costs.

This illustrates how a Lindahl equilibrium does not take into account the distribution of surplus among consumers when optimizing the provision of a public good. Even though the outcome is individually rational, some consumers may be left far better off than others. In order to address this flaw, we introduce the concept of a generalized Lindahl equilibrium.

### 3 Generalized Lindahl Equilibrium

A classical Lindahl equilibrium must be individually rational and raise tax revenue equal to the cost of provision of the public good. As stated in the previous section, the disregard of consumer surplus in such an equilibrium can lead to unfair outcomes. In this section, I explain how a generalization of the classical Lindahl result can allow a social planner to construct far more equitable outcomes.

The unjustness of the classical result can be thought of as a result of the linearity of prices faced by consumers. Consumers are charged one personalized tax across all margins of consumption of the public good; the surplus they attain at each margin is then the difference between their marginal valuations and these constant taxes. If those taxes were non-linear, a more equitable outcome could be achieved. A tax schedule could be created such that the equity of the benefits accruing to each consumer at each margin of provision of the public good were taken into account. This allowance for non-linear taxation *is* the generalization of the classical Lindahl equilibrium that constitutes the purpose of this paper.

A generalized Lindahl equilibrium sees each consumer receive a personalized tax schedule, instead of a single personalized tax. The outcome is still individually rational and covers the cost of provision, as in the classical case, but now a central planner or government can manipulate the personalized price schedules to distribute consumer surplus as they see fit.

**Definition 1.** *A profile of individualized price schedules induces a generalized Lindahl equilibrium if it is (1) individually rational for each individual to demand the Pareto optimal amount of the public good and (2) if the cost of providing the public good is exactly covered by taxes paid by each individual.*

Note that a generalized Lindahl equilibrium is agnostic about the specific distribution of consumer surplus. It merely requires that a tax schedule is administered. Consequently, there are many generalized Lindahl equilibria capable of providing the Pareto optimal level of the public good. The social planner can

chose an outcome that obtains a desired distribution from among all feasible generalized Lindahl equilibria. This will be explored further in the next section.

For now I shall discuss some necessary conditions for a generalized Lindahl equilibrium, using the economic setting described above for the exposition. A personalized tax schedule is defined as a function  $p_i : \mathbb{R} \rightarrow \mathbb{R}$ , such that the total tax revenue collected from individual  $i$  for a level of provision of the public good  $x$  is  $T_i(x) = \int_0^x p_i(t)dt$ . The budget constraint of a customer is then  $\int_0^x p_i(t)dt + y_i = \omega_i$ . If we solve for the level of the private good in terms of the level of the public good given their initial endowment,  $y_i(x)$ , and substitute this into their utility function, we can express their utility in terms of their initial allocation and the level of public good provided.

$$u_i(x, y_i(x)) = u_i \left( x, \omega_i - \int_0^x p_i(t)dt \right) = \int_0^x [v_i(t) - p_i(t)] dt + \omega_i \quad (2)$$

From this we get the first of two conditions for individual rationality:

$$\frac{\partial u_i(x, y_i(x))}{\partial x} \Big|_{x=x^{PO}} = 0 \rightarrow v_i(x^{PO}) = p_i(x^{PO}) \quad (3)$$

The prices charged at the margin for the public good in a generalized Lindahl equilibrium must be the Lindahl prices from the classical Lindahl equilibrium. The second condition for individual rationality requires that customers are indeed better off at the Pareto optimal outcome than they would be with zero public good provided.

$$u_i(x^{PO}, y_i(x^{PO})) \geq u_i(0, y_i(0)) \quad (4)$$

And finally, for the tax revenues to cover the cost of provision, we must have

$$\sum_i \int_0^{x^{PO}} p_i(t)dt = cx^{PO} \quad (5)$$

Equations 3, 4, and 5 must all hold for an equilibrium to qualify as a generalized Lindahl equilibrium. In the next section, they shall function as constraints on a social welfare problem facing a social planner who wants to optimally provide a public good, and who has CES preferences over the individual consumer surpluses it oversees.

### 3.1 Generalized Lindahl Equilibrium under a CES Social Welfare Function

We can imagine a society's preferences over the distribution of surplus being encoded into a marginal social welfare function. The maximization of the integral of this social welfare function, subject to the conditions necessary for a generalized Lindahl equilibrium stated in the previous section, is a theoretical

representation of how such a society could provide a public good in an equitable way.

In general, let's say that this marginal social welfare function is a function of the marginal consumer surpluses of each individual in a society:  $W : \mathbb{R}_+^n \rightarrow \mathbb{R}$  is differentiable, quasiconcave, increasing in each of its arguments, and  $W(\mathbf{0}) = 0$ . The restriction that  $W$  is defined on the non-negative elements of  $\mathbb{R}^n$  introduces a new constraint,  $v_i(t) - p_i(t) \geq 0 \forall t \in [0, x^{PO}]$ , which must hold for integration to be defined. The combination of this new constraint and  $v_i(t) - p_i(t) < 0 \forall t > x^{PO}$  provide a sufficient condition for individual rationality, ensuring that any optimal allocation of the public good will be individually rational. The problem facing a society is then

$$\max_{\{p_i(t)\}_{i=1}^n} \int_0^{x^{PO}} W(v_1(t) - p_1(t), \dots, v_n(t) - p_n(t)) dt \quad (6)$$

$$s.t. v_i(x^{PO}) = p_i(x^{PO}) \quad (7)$$

$$v_i(t) - p_i(t) \geq 0 \quad \forall t \in [0, x^{PO}] \quad (8)$$

$$v_i(t) - p_i(t) < 0 \quad \forall t > x^{PO} \quad (9)$$

$$\int_0^{x^{PO}} \sum_{i=1}^n p_i(x) dx = cx^{PO} \quad (10)$$

I will now solve a specific example of this problem in which society has CES preferences over individual marginal surplus.

$$\max_{\{p_i(t)\}_{i=1}^n} \int_0^{x^{PO}} \left[ \sum_{i=1}^n \beta_i (v_i(x) - p_i(x))^\alpha \right]^{\frac{1}{\alpha}} dx \quad (11)$$

$$s.t. v_i(x^{PO}) = p_i(x^{PO}) \quad (12)$$

$$\int_0^{x^{PO}} \sum_{i=1}^n p_i(x) dx = cx^{PO} \quad (13)$$

$\beta_i$  is an arbitrary weight for individual  $i$ 's surplus and  $-\infty > \alpha \leq 1$ . Equation 12 ensures that the level of the public good provided is the level demanded by each individual, and equation 13 ensures that tax revenue will equal the cost of provision. In this instance, 8 and 9 will not bind.

This is a calculus of variations problem. Throughout its solution any dependency on  $x$  is suppressed. The necessary conditions to solve it are 12, 13, and  $n$  Euler equations, one for each individual:

$$\lambda = -\beta_j (v_j - p_j)^{\alpha-1} \left[ \sum_{i=1}^n \beta_i (v_i - p_i)^\alpha \right]^{\frac{1-\alpha}{\alpha}} \quad \forall j = 1, \dots, n \quad (14)$$

$\lambda$  is the Lagrangian multiplier on the iso-perimetric constraint 13. If we set the Euler equations for individuals 1 and 2 equal to each other, we get

$$\left(\frac{\beta_2}{\beta_1}\right)^\sigma (v_1 - p_1) = v_2 - p_2$$

Where  $\sigma = \frac{1}{1-\alpha}$  is the elasticity of substitution between the surplus of any two individuals. I will define  $\Omega_{21} = (\beta_2/\beta_1)^\sigma$ . Substituting this in and rearranging, we get

$$p_2 - \Omega_{21}p_1 = v_2 - \Omega_{21}v_1$$

If we perform a similar manipulation among all  $n$  Euler conditions, we are left with a system of  $n - 1$  equations

$$\begin{aligned} p_2 - \Omega_{21}p_1 &= v_2 - \Omega_{21}v_1 \\ p_3 - \Omega_{32}p_2 &= v_3 - \Omega_{32}v_2 \\ &\vdots \\ p_n - \Omega_{n,n-1}p_{n-1} &= v_n - \Omega_{n,n-1}v_{n-1} \end{aligned}$$

One way to ensure that the iso-perimetric constraint holds is to assume that the taxes levied at each marginal provision of the public good are always equal to the marginal cost of production.

$$p_1 + p_2 + \dots + p_n = c \tag{15}$$

With this and the system of  $n - 1$  equations above, we now have a system of  $n$  equations that we can use to solve for closed-form solutions of individualized price schedules (the full solution is detailed in the appendix). Those solutions are

$$p_j = v_j - \frac{\beta_j^\sigma}{\sum_{i=1}^n \beta_i^\sigma} \left[ \sum_{i=1}^n (v_i) - c \right] \tag{16}$$

In this solution, the optimal personalized tax schedule has two parts. Individuals are taxed at each margin according to their private valuation at that margin,  $v_j$ . This baseline is then adjusted by what fraction of the marginal net social benefit ( $\sum_{i=1}^n (v_i) - c$ ) is redistributed back to them, as dictated by what share their social weights have in society ( $\beta_j^\sigma / \sum_{i=1}^n \beta_i^\sigma$ ). The redistributive nature of the tax schedules is made clear when the marginal net benefit earned under them is stated:

$$b_j \equiv v_j - p_j = \frac{\beta_j^\sigma}{\sum_{i=1}^n \beta_i^\sigma} \left[ \sum_{i=1}^n (v_i) - c \right]$$

The baseline outcome is that individuals are taxed such that all of their private surplus is taken away, and it is only the redistribution of net *social* benefit that leads to any individual being better off in the end.

For very small elasticities of substitution the complementarity of individuals' surplus in the eyes of society overrides their social weights. They are treated more or less equally. When  $\sigma = 0$ , consumer surpluses are perfect complements and we are left with the following tax schedules.

$$p_j = v_j - \frac{1}{n} \left[ \sum_{i=1}^n (v_i) - c \right]$$

When society distributes consumer surplus equally, individuals are left with equal shares of the net social benefit at each margin of provision of the public good. When  $\sigma = 1$  the original marginal social welfare function becomes a Cobb-Douglas function and the personalized tax schedules become

$$p_j = v_j - \frac{\beta_j}{\sum_{i=1}^n \beta_i} \left[ \sum_{i=1}^n (v_i) - c \right]$$

The weights now have an effect because individuals' surpluses are substitutable. An individual with a higher social weight now receives a larger share of the net social benefit at each margin because it is optimal for society to substitute their surplus for others'. As  $\sigma \rightarrow \infty$ , individuals' surpluses become perfect substitutes. Those with the highest social weight ( $j \in J_{max}$ , where  $J_{max} \equiv \{j : \beta_j = \max\{\beta_i\}_1^n\}$ ) receive all of the marginal net social benefit. Among these individuals, society is indifferent about the distribution of surplus. I shall assume that it is redistributed evenly among them. Those with lower weights receive no redistribution.

$$p_j = \begin{cases} v_j - \frac{1}{|J_{max}|} [\sum_{i=1}^n (v_i) - c] & \text{if } j \in J_{max} \\ v_j & \text{else} \end{cases}$$

Figure 2 plots some examples of 16 for various values of  $\sigma$ . I reuse the setting used to expose the unjustice of the classical Lindahl equilibrium in section 2.1: there are two consumers with  $v_1(x) = 20 - 2x$  and  $v_2(x) = 6 - 0.1x$ , the marginal cost of providing the public good is 5, and a Pareto optimal amount of it is 10.  $\beta_1$  is 10 and  $\beta_2$  is 1.

## 4 Generalized Walker Mechanism

As is often the case in the theory of public goods, one of the strongest assumptions made here is that private valuations of the public good are known by the social planner. In an effort to relax this assumption, I now describe a mechanism a social planner could implement in order to obtain a Lindahl equilibrium without knowledge of individuals' private valuations. When this mechanism is implemented, it induces a generalized Lindahl equilibrium. In general, a mechanism is defined by the messages agents are allowed to choose and the outcome functions that convert these messages into outcomes.

Price Schedules, Marginal Benefits, and Marginal Cost in a 2 Consumer Setting for Various Values of  $\sigma$

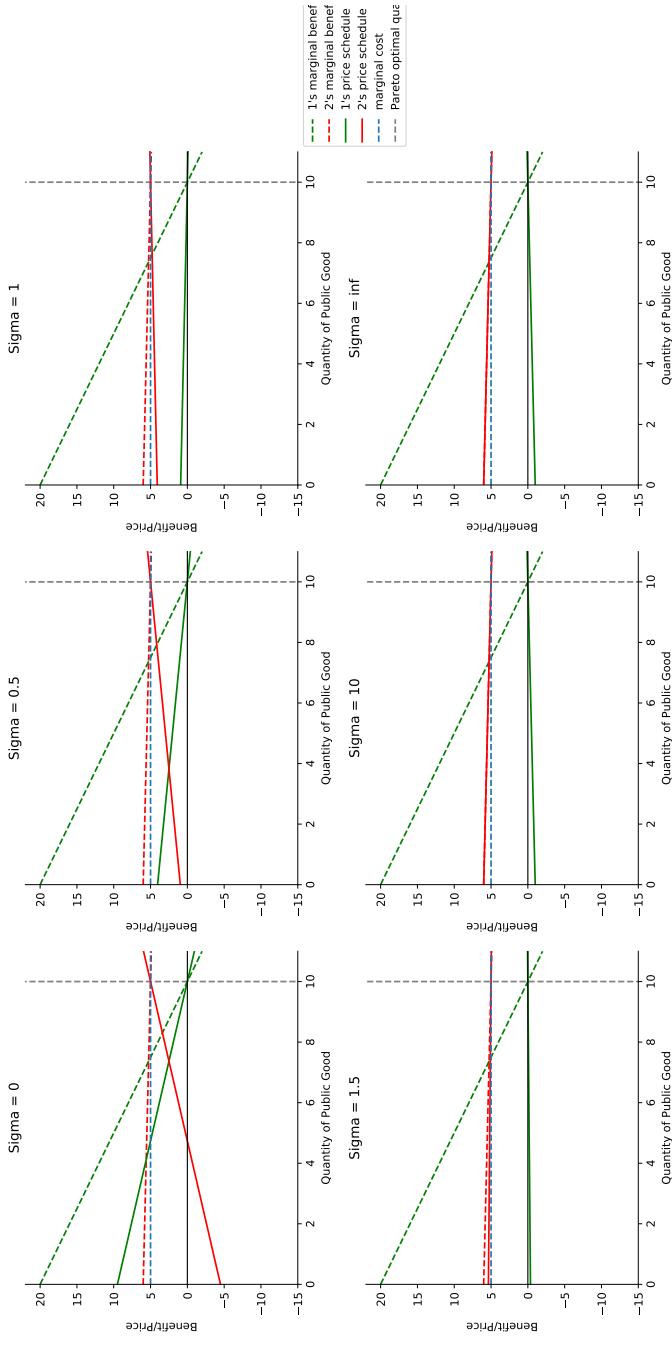


Figure 2: Above are several representations of price schedules leading to generalized Lindahl equilibria for a setting with two individuals and various values of  $\sigma$ , the elasticity of substitution between the surpluses of any two individuals. The marginal benefit functions are  $v_1(x) = 20 - 2x$ ,  $v_2(x) = 6 - x/10$ . Social welfare weights are  $\beta_1 = 10$  and  $\beta_2 = 1$ . Marginal cost of providing the public good is 5, and the Pareto optimal amount of it is 10.

In our environment (with a slight change of notation) an outcome is specified by a level of the public good  $x$  and a set of personalized tax schedules for each individual  $\tau_i$  such that each individual pays in total  $T_i(x) = \int_0^x \tau_i(t)dt$ . The messages sent by individuals  $i = 1, \dots, n$  are continuous functions  $p_i : \mathbb{R} \rightarrow \mathbb{R}$ . The list of all of these messages is a message profile  $\mathbf{p} = (p_1, \dots, p_n)$ . There are  $n + 1$  outcome functions, one to determine the provision of the public good,  $\mathcal{X} : \prod_i^n C(\mathbb{R}) \rightarrow \mathbb{R}$ , and  $n$  to determine what tax schedules are to be administered,  $\mathcal{T}_i : \prod_i^n C(\mathbb{R}) \rightarrow C(\mathbb{R})$ . The functional forms of these outcome functions are

$$\mathcal{X}(\mathbf{p}) = \arg \max_{t \geq 0} \int_o^t \sum_{i=1}^n p_i(m) dm \quad (17)$$

$$\mathcal{T}_i(\mathbf{p}) = \frac{c}{n} + p_{i+2} - p_{i+1} \quad (18)$$

Note that the subscript  $i + j$  is understood to be a modulo operation,  $n + j = j$ , and  $c$  is once again the marginal cost of the public good (this follows the mechanism design in Walker (1981)).

This mechanism defines a game form: when preferences are added to it a game is induced. The outcome functions 17 and 18 each are designed to influence what strategies are used to satisfy these preferences in such a game. First, when all other messages  $p_{-j}$  are taken as given, 17 allows any individual message  $p_j$  to demand a certain level of the public good: no matter what function results from  $\sum_{i \neq j} p_i(m)$ , the addition of  $p_j$  can create an integrand  $\sum_{i=1}^n p_i(m)$  such that  $\mathcal{X}(\mathbf{p})$  produces any outcome  $x$  an individual  $j$  desires. Second, 18 states that the tax schedule that will be administered to any individual does not depend on their own message. In tandem, these two outcome functions mean that a utility optimizing agent playing a game induced by this mechanism only has one margin of optimization, the public good. I now add the preferences from the economic setting described in section 2.1 in order to induce and analyze the resulting game.

**Proposition 1.** *At any Nash equilibrium induced by our mechanism, any equilibrium outcome  $(\vec{\tau}, x)$  coincides with a generalized Lindahl outcome and is Pareto optimal.*

*Proof.* To be a generalized Lindahl outcome, each individual must be given a personal tax schedule, the cost of providing the public good must be covered, and the allocation  $(\vec{\tau}, x)$  must be individually rational and Pareto optimal. If the budget constraint of a player is  $\int_0^x p_i(t)dt + y_i = \omega_i$  and their utility function is  $u_i(x, y_i)$ , then each player faces the following utility maximization problem:

$$\begin{aligned} & \max_{p_i(x)} u_i \left( x, \omega_i - \int_0^x \tau_i(t)dt \right) \\ & s.t. \quad x = \arg \max_{t \geq 0} \int_o^t \sum_{i=1}^n p_i(m) dm \end{aligned}$$

or

$$\max_{p_i(x)} u_i \left( x, \omega_i - \int_0^x \left[ \frac{c}{n} + p_{i+2}(t) - p_{i+1}(t) \right] dt \right) \quad (19)$$

$$s.t. x = \arg \max_{t \geq 0} \int_0^t \sum_{i=1}^n p_i(m) dm \quad (20)$$

The problem can be simplified further, as a choice over  $x$ . The objective function does not *directly* depend on an individual's choice of message  $p_i$ , but the amount of the public good provided does, via the constraint 20: if a player understands  $x$  is chosen according to 20, then they understand that the following first order condition  $\sum_{i=1}^n p_i(x) = 0$  must hold. Each player takes all other messages  $p_{-i}$  as given, and submits a  $p_i$  such that the maximizing argument of 20 is the  $x$  they deem optimal. Only by this level of  $x$  does  $p_i$  affect their utility; the problem can be rewritten as:

$$\max_x u_i \left( \omega_i - \int_0^x \left[ \frac{c}{n} + p_{i+2}(t) - p_{i+1}(t) \right] dt, x \right) \quad (21)$$

The first order condition for utility maximization is:

$$\frac{u_{i,x}(y(x), x)}{u_{i,y_i}(y(x), x)} = \frac{c}{n} + p_{i+2}(x) - p_{i+1}(x) \quad (22)$$

Equation 22 states that in a Nash equilibrium a best responding individual submits a message  $p_i$  such that the level of the public good the message implicitly demands equates their marginal rate of substitution to their personalized tax at the margin demanded. If we assume that  $u_i$  is quasi-concave (we can't, sign/magnitude of  $p_{i+2} - p_{i+1}$  matters) then 22 defines a global maximum, and we can state that a Nash equilibrium is individually rational for all players to accept.

Note that the best response message  $p_i$  is not unique. 22 only applies at a single point, the level of the public good provided. This is true because the only feature of  $p_i$  relevant to utility maximization is where it dictates  $\sum_{i=1}^n p_i(x) = 0$ . Consequently, this is also the only feature relevant to the achievement of a Nash equilibrium. Therefore, a best response  $p_i$  could be any integrable function that satisfies the point constraint 22.

The sum of all  $n$  first order conditions is

$$\sum_{i=1}^n \frac{u_{i,x}(y(x), x)}{u_{i,y_i}(y(x), x)} = c \quad (23)$$

The modulo dependence of the personalized taxes creates a cancellation that leads to the sum of all marginal rates of substitution equalling the marginal cost of provision: the outcome is Pareto optimal.

The final qualification left to prove is that the cost of provision is covered. This is true via the construction of the personalized tax outcome function. No matter what messages are submitted, the sum of all tax schedules  $\tau_i$  equals marginal cost at every single margin,  $\sum_{i=1}^n \tau_i = c$ .

Any Nash equilibrium of the game induced by our mechanism is a generalized Lindahl outcome: the outcome provides individualized tax schedules, covers the cost of provision of the public good, and is Pareto optimal and individually rational. This concludes the proof  $\square$

**Proposition 2.** *Any generalized Lindahl outcome  $(\vec{\tau}, x)$  can be achieved in a Nash equilibrium of a game induced by the mechanism by a correspondence of  $\mathbf{p}$ .*

*Proof.* Suppose we have a generalized Lindahl outcome  $(\vec{\tau}, x)$ . Note that this means that  $x$  is the Pareto optimal amount of the public good, so it is a best response for each individual to submit a message such that the outcome functions 17 and 18 provide  $x$ . With this in mind, if we show that a profile of messages  $\mathbf{p}$  is capable of generating an outcome  $(\vec{\tau}, x)$  via 17 and 18, we have shown that the generalized Lindahl equilibrium that led to 17 and 18 is a Nash equilibrium in a game induced by the mechanism.

This proof again follows Walker (1981). From 17 and 18, we can say we have the following system of  $n$  equations:

$$\begin{aligned} x &= \arg \max_{t \geq 0} \int_o^t \sum_{i=1}^n p_i(m) dm \\ \tau_n - \frac{c}{n} &= p_2 - p_1 \\ \tau_1 - \frac{c}{n} &= p_3 - p_2 \\ \tau_2 - \frac{c}{n} &= p_4 - p_3 \\ &\vdots \\ \tau_{n-2} - \frac{c}{n} &= p_n - p_{n-1} \end{aligned}$$

And now we are aided by the fact that a correspondence of  $\mathbf{p}$  is capable of producing a Nash equilibrium. The above system has  $n$  unknowns  $p_i$ . To define a unique profile  $\mathbf{p}$  defined across all margins of the public good, we would need  $n$  linearly independent equations that provide structure across said margins. This system only has  $n - 1$  such equations: the functions  $\tau_i - \frac{c}{n} = p_{i+2} - p_{i+1}$ . The  $n^{th}$  equation,  $x = \arg \max_{t \geq 0} \int_o^t \sum_{i=1}^n p_i(m) dm$ , only has information on a single margin, the last margin provided. But, this is all the information that is needed to construct a correspondence of profiles  $\mathbf{p}$  that is capable of achieving a

Nash equilibrium. Thus any generalized Lindahl equilibrium can be achieved in a Nash equilibrium of a game induced by the mechanism by a correspondence of  $\mathbf{p}$ .  $\square$

There is no generalized Lindahl equilibrium the mechanism cannot create, and every equilibrium of the mechanism is a generalized Lindahl equilibrium.

## 5 Discussion

## 6 Appendix

The solution will take place as follows: I will solve for  $p_1$  in terms of  $p_2, v_1$  and  $v_2$ , and then substitute  $p_1$  into 15; I will iterate this process two more times, substituting equations for  $p_2$  and  $p_3$  into 15; at that point a pattern will emerge that will allow me to extrapolate a closed-form solution for to  $p_n$ , and then for  $p_{n-1}$ ; finally, I will use the pattern observable in the solutions for  $p_n$  and  $p_{n-1}$  solution to deduce a closed form solution for  $p_i$  in general.

As stated, first solve for  $p_1$

$$p_1 = \frac{-v_2}{\Omega_{21}} + v_1 + \frac{p_2}{\Omega_{21}}$$

Substitute into 15

$$\left( \frac{-v_2}{\Omega_{21}} + v_1 + \frac{p_2}{\Omega_{21}} \right) + p_2 + p_3 + \dots + p_n = c$$

Combine terms, keeping prices and marginal benefits organized, and note that  $\Omega_{ij} = \Omega_{ji}^{-1}$ . We get our first instance of a pattern.

$$(1 + \Omega_{12}) p_2 + p_3 + \dots + p_n + v_1 - \Omega_{12} v_2 = c$$

Iterating, solve for  $p_2$  in terms of  $p_3, v_2$  and  $v_3$  using the equation relating  $p_2$  to  $p_3$  from our system of equalities. Substitute into above.

$$(1 + \Omega_{12}) \left( \frac{-v_3}{\Omega_{32}} + v_2 + \frac{p_3}{\Omega_{32}} \right) + p_3 + p_4 + \dots + p_n + v_1 - \Omega_{12} v_2 = c$$

Combine terms again, and note that  $\Omega_{Ki}/\Omega_{Kj} = \Omega_{ij}$ . A nice cancellation occurs. We get a second instance of the pattern.

$$(1 + \Omega_{23} + \Omega_{13}) p_3 + p_4 + \dots + p_n + v_1 + v_2 - (\Omega_{23} + \Omega_{13}) v_3 = c$$

Iterate one last time, substituting for  $p_3$ .

$$(1 + \Omega_{23} + \Omega_{13}) \left( \frac{-v_4}{\Omega_{43}} + v_3 + \frac{p_4}{\Omega_{43}} \right) + p_4 + p_5 + \dots + p_n + v_1 + v_2 - (\Omega_{23} + \Omega_{13})v_3 = c$$

Combine terms, and we get a third instance of the pattern.

$$(1 + \Omega_{34} + \Omega_{24} + \Omega_{14}) p_4 + p_5 + \dots + p_n + v_1 + v_2 + v_3 - (\Omega_{34} + \Omega_{24} + \Omega_{14})v_4 = c$$

Let's quickly look at all three instances of the pattern together.

$$(1 + \Omega_{12}) p_2 + p_3 + \dots + p_n + v_1 - \Omega_{12}v_2 = c$$

$$(1 + \Omega_{23} + \Omega_{13}) p_3 + p_4 + \dots + p_n + v_1 + v_2 - (\Omega_{23} + \Omega_{13})v_3 = c$$

$$(1 + \Omega_{34} + \Omega_{24} + \Omega_{14}) p_4 + p_5 + \dots + p_n + v_1 + v_2 + v_3 - (\Omega_{34} + \Omega_{24} + \Omega_{14})v_4 = c$$

By extrapolation, and noting that  $\Omega_{ii} = 1$ ,  $p_n$  must be implicitly defined by

$$p_n \sum_{i=1}^n \Omega_{in} + \sum_{i=1}^{n-1} v_i - v_n \sum_{i=1}^{n-1} \Omega_{in} = c$$

Solving for  $p_n$ , we get

$$p_n = \left( \sum_{i=1}^n \Omega_{ni} \right)^{-1} \left[ c - \sum_{i=1}^{n-1} v_i + v_n \sum_{i=1}^{n-1} \Omega_{ni} \right] \quad (24)$$

To get a closed form solution for any  $p_i$ , I'll first use 6 to solve for  $p_{n-1}$  using the equation from our system of equations that relates  $p_{n-1}$  to  $p_n$ . Then with closed form solutions for both  $p_n$  and  $p_{n-1}$ , a pattern emerges that I use to solve for the general closed form solution for  $p_i$ .

First, solve for  $p_{n-1}$ . We know that

$$p_{n-1} = \frac{-v_n}{\Omega_{n,n-1}} + v_{n-1} + \frac{p_n}{\Omega_{n,n-1}}$$

Substitute in 6. Step by step simplification:

$$\begin{aligned}
p_{n-1} &= \frac{-v_n}{\Omega_{n,n-1}} + v_{n-1} + \frac{1}{\Omega_{n,n-1}} \left( \sum_{i=1}^n \Omega_{in} \right)^{-1} \left[ c - \sum_{i=1}^{n-1} v_i + v_n \sum_{i=1}^{n-1} \Omega_{in} \right] \\
&= -\Omega_{n-1,n} v_n + v_{n-1} + \left( \sum_{i=1}^n \Omega_{n,n-1} \Omega_{in} \right)^{-1} \left[ c - \sum_{i=1}^{n-1} v_i + v_n \sum_{i=1}^{n-1} \Omega_{in} \right] \\
&= -\Omega_{n-1,n} v_n + v_{n-1} + \left( \sum_{i=1}^n \Omega_{i,n-1} \right)^{-1} \left[ c - \sum_{i=1}^{n-1} v_i + v_n \sum_{i=1}^{n-1} \Omega_{in} \right] \\
&= \left( \sum_{i=1}^n \Omega_{i,n-1} \right)^{-1} \left[ -\Omega_{n-1,n} v_n \sum_{i=1}^n \Omega_{i,n-1} + v_{n-1} \sum_{i=1}^n \Omega_{i,n-1} + c - \sum_{i=1}^{n-1} v_i + v_n \sum_{i=1}^{n-1} \Omega_{in} \right] \\
&= \left( \sum_{i=1}^n \Omega_{i,n-1} \right)^{-1} \left[ -v_n \sum_{i=1}^n \Omega_{in} + v_{n-1} \sum_{i=1}^n \Omega_{i,n-1} + c - \sum_{i=1}^{n-1} v_i + v_n \sum_{i=1}^{n-1} \Omega_{in} \right] \\
&= \left( \sum_{i=1}^n \Omega_{i,n-1} \right)^{-1} \left[ -v_n \sum_{i=1}^n \Omega_{in} + v_{n-1} \sum_{i=1}^n \Omega_{i,n-1} + c - \sum_{i=1}^{n-1} v_i + v_n \sum_{i=1}^{n-1} \Omega_{in} \right] \\
&= \left( \sum_{i=1}^n \Omega_{i,n-1} \right)^{-1} \left[ v_{n-1} \sum_{i=1}^n \Omega_{i,n-1} + c - \sum_{i=1}^{n-1} v_i + v_n \left( \sum_{i=1}^{n-1} \Omega_{in} - \sum_{i=1}^n \Omega_{in} \right) \right]
\end{aligned}$$

At this point, we need to break steps down term by term. Note that the term in the parenthesis within the brackets equals -1, and that

$$v_{n-1} \sum_{i=1}^n \Omega_{i,n-1} = v_{n-1} \left( 1 + \sum_{i \neq n-1} \Omega_{i,n-1} \right) = v_{n-1} + v_{n-1} \sum_{i \neq n-1} \Omega_{i,n-1}$$

The lone  $v_{n-1}$  will cancel with the final term in the  $\sum_{i=1}^{n-1} v_i$  summation. This leaves us with

$$p_{n-1} = \left( \sum_{i=1}^n \Omega_{i,n-1} \right)^{-1} \left[ v_{n-1} \sum_{i \neq n-1} \Omega_{i,n-1} + c - \sum_{i=1}^{n-2} v_i - v_n \right]$$

Or

$$p_{n-1} = \left( \sum_{i=1}^n \Omega_{i,n-1} \right)^{-1} \left[ c - \sum_{i \neq n-1} v_i + v_{n-1} \sum_{i \neq n-1} \Omega_{i,n-1} \right] \quad (25)$$

Our closed form solution to  $p_{n-1}$ . Finally let's compare 7 to 6.

$$p_{n-1} = \left( \sum_{i=1}^n \Omega_{i,n-1} \right)^{-1} \left[ c - \sum_{i \neq n-1} v_i + v_{n-1} \sum_{i \neq n-1} \Omega_{i,n-1} \right]$$

$$p_n = \left( \sum_{i=1}^n \Omega_{in} \right)^{-1} \left[ c - \sum_{i=1}^{n-1} v_i + v_n \sum_{i=1}^{n-1} \Omega_{in} \right]$$

From the above pattern, it's clear that the closed form solution for any  $p_j$  is

$$p_j = \left( \sum_{i=1}^n \left( \frac{\beta_i}{\beta_j} \right)^\sigma \right)^{-1} \left[ c - \sum_{i \neq j} v_i + v_j \sum_{i \neq j} \left( \frac{\beta_i}{\beta_j} \right)^\sigma \right] \quad (26)$$

Inside the brackets, add  $0 = v_j \left( \frac{\beta_j}{\beta_j} \right)^\sigma - v_j$  and rearrange to get:

$$p_j = v_j - \frac{\beta_j^\sigma}{\sum_{i=1}^n \beta_i^\sigma} \left[ \sum_{i=1}^n (v_i) - c \right] \quad (27)$$