

**Exercise 8.10:** Write and solve the constraints that are generated by the inter-procedural sign analysis for the program from Exercise 8.4 this time with context sensitivity using the call string approach with  $k = 1$ . (Even though this program does not need context sensitivity to be analyzed precisely, it illustrates the mechanism behind the call string approach.)

解答:

Constraint for entry node  $v$  of function  $\text{inc}(a)$ :

$$\llbracket v \rrbracket(c) = \bigcup_{\substack{w \in \text{pred}(v) \\ \wedge c = w \\ \wedge c' \in \text{Contexts}}} s_w^{c'}$$

其中  $s_w^{c'}$  表示从上下文  $c$  中节点  $w$  的调用创建的抽象状态:

$$s_w^{c'} = \begin{cases} \text{unreachable} & (\text{if } \llbracket w \rrbracket(c') = \text{unreachable}) \\ \perp [a \mapsto \text{eval}(\llbracket w \rrbracket(c'), 17)] & (\text{otherwise}) \end{cases}$$

$$s_w^{c'} = \begin{cases} \text{unreachable} & (\text{if } \llbracket w \rrbracket(c') = \text{unreachable}) \\ \perp [a \mapsto \text{eval}(\llbracket w \rrbracket(c'), 87)] & (\text{otherwise}) \end{cases}$$

Constraint for after-call node  $v$  labeled  $X = \square$ , with call node  $v'$  and exit node  $w \in \text{pred}(v)$ :

$$\llbracket v \rrbracket(c) = \begin{cases} \text{unreachable} & (\text{if } \llbracket v' \rrbracket(c) = \text{unreachable} \vee \llbracket w \rrbracket(v') = \text{unreachable}) \\ \llbracket v' \rrbracket(c)[X \mapsto \llbracket w \rrbracket(v'), (18)] & (\text{otherwise}) \end{cases}$$

$$\llbracket v \rrbracket(c) = \begin{cases} \text{unreachable} & (\text{if } \llbracket v' \rrbracket(c) = \text{unreachable} \vee \llbracket w \rrbracket(v') = \text{unreachable}) \\ \llbracket v' \rrbracket(c)[X \mapsto \llbracket w \rrbracket(v'), (88)] & (\text{otherwise}) \end{cases}$$