

Tutorial - 2

Name - Mushtaq Agha
Section - 4
U-Roll No. - 2016857

Ans-1

```
void fun(int n) {
    int j = 1, i = 0;
    while (i < n) {
        i = i + j;
        j = j + 1;
    }
}
```

$j = 1$ $i = 0 + 1$
 $j = 2$ $i = 0 + 1 + 2$
 $j = 3$ $i = 0 + 1 + 2 + 3$
 \vdots
 Loop ends when $i \geq n$
 $0 + 1 + 2 + 3 + \dots + n \geq n$
 $\frac{k(k+1)}{2} \geq n$
 $k^2 \geq n$
 $k \geq \sqrt{n}$

$$\Rightarrow O(\sqrt{n})$$

Ans-2

Recurrence Relⁿ for Fibonacci Series

$$T(n) = T(n-1) + T(n-2) \quad T(0) = T(1) = 1$$

• if $T(n-1) \approx T(n-2)$

$$T(n) \approx 2T(n-2)$$

Put $n = n-2$

$$T(n-2) \approx 2T(n-4)$$

$$T(n) \approx 2(2T(n-4)) = 4T(n-4)$$

Put $n = n-4$

$$T(n-4) \approx 2T(n-6)$$

$$T(n) \approx 4(2T(n-6)) = 8T(n-6)$$

$$\vdots$$

$$T(n) = 2^k T(n-2k)$$

$$T(n) = 2^n T(n-2k)$$

$$2^n T(0)$$

$$O(2^n)$$

OR

$$\therefore n-2k = 0$$

$$n = 2k$$

$$n-2k = 1$$

$$n = 1 + 2k$$

$$\frac{n-1}{2} = k$$

Ans 3 complexity $n \lg n, n^3, \lg \lg n$

2. $O(n \lg n)$

```
for (int i = 0; i < n; i++) {
    for (int j = 1; j < n; j = j * 2) {
        // O(1)
    }
}
```

• $O(n^3)$

```
for (i = 0 to n; i++)
{
    for (j = 0 to n; j++) {
        for (k = 0 to n; k++) {
            // O(1)
        }
    }
}
```

• $O(\lg \lg n)$:
 for (i = 1 to n i = i * 2) {
 for (j = 1 to n j = j * 2) {
 // O(1)
 }
 }

Ans 4

$$T(n) = T(n/4) + T(n/2) + (n^2)$$

We know that

$$T\left(\frac{n}{2}\right) > T\left(\frac{n}{4}\right)$$

$$\text{So } T(n) \geq 2T\left(\frac{n}{2}\right) + (n^2)$$

applying Master theorem $T(n) = aT\left(\frac{n}{b}\right) + f(n)$

$$a = 2, b = 2, f(n) = n^2$$

$$c = \log_b a = \log_2 2 = 1$$

Compare n^c with n^2
 $f(n) > n^c$ So $T(n) = \Theta(n^2)$

Ans 5

```
int fun( int n) {
    for (int i=1; i<=n; i++) {
        for (int j=1; j<=n; j++) {
            // O(1) 333
        }
    }
}
```

So total complexity
 $= O(n^2 + n^4 + n^6 + \dots)$
 $= O(n^2)$

$i=1$ — n times
 $j=1$
 $j=2$
 $j=3$
 \vdots

$i=2$ — $j=1$ Loop ends when $j > n$
 $j=3$
 $j=5$
 $j=7$ $1+3+5+7 > n$
 \vdots $k > \frac{n}{2}$
 $- n$ times

$i=3$ — $j=1$ $1+4+7 > n$
 $j=4$ $k > \frac{n}{3}$
 $j=7$
 $i=4$ $k > \frac{n}{4}$
 \vdots
 $i=n$

Ans 8

```
for (i=2; i<=n; i = pow(i, k)) {
    // O(1) 3
}
```

Complexity $\text{Pow}(i, k) \rightarrow O(\log N) = \log(k)$

$i=2$
 $i=2^k$
 $i=2^{k^2}$
 \vdots
 $i=2^{k^m}$

Loop ends when $i > n$
 $2^{k^m} > n$

$\log(2^{k^m}) > \log n$
 $k^m > \log n$

$\log(k^m) > \log(\log n)$

$m \log k > \log(\log n)$

$m > \log(\log n)$

$T(n) = O(\log(\log n))$ $\log(k)$

26

$$a) 100 < \lg n < \sqrt{n} < n < \lg(\lg(n)) < n \lg n < \lg n!$$

$$< n^2 < \lg^2 n < 2^n < 2^n < 4^n$$

$$b) 1 < \sqrt{\lg n} < \lg n < 2 \lg n < \lg_2 n < n < 2n < 4n < \lg^2 n$$

$$< n \lg n < \lg n! < 2n! < n^2 < 2 \times 2^n$$

$$c) a! < \lg_2 n < \lg_2 n < n \lg_2 n < n \lg_2 n < \lg n! < n! < 8n$$

$$< 8n^2 < 7n^3 < 8^{2n}$$