

# OpenBox CS771

May 5, 2025

## 1 MATHEMATICAL DERIVATION

Starting with the upper signal equations for two PUF instances (denoted as PUFO and PUF1). For multiplexer index  $i = 1$ , the upper signal equations are:

**PUF0** ( $n = 0$ ):

$${}^0t_1^u = (1 - c_1) [{}^0t_0^u + {}^0p_1] + c_1 [{}^0t_0^l + {}^0s_1]$$

**PUF1** ( $n = 1$ ):

$${}^1t_1^u = (1 - c_1) [{}^1t_0^u + {}^1p_1] + c_1 [{}^1t_0^l + {}^1s_1]$$

We then do the subtraction of Upper Signals at MUX1 by defining the difference in upper signals at MUX1 as:

$$\Delta t_1^u = {}^1t_1^u - {}^0t_1^u$$

Substitute the equations for  ${}^1t_1^u$  and  ${}^0t_1^u$ :

$$\begin{aligned} \Delta t_1^u &= [(1 - c_1)({}^1t_0^u + {}^1p_1) + c_1({}^1t_0^l + {}^1s_1)] \\ &\quad - [(1 - c_1)({}^0t_0^u + {}^0p_1) + c_1({}^0t_0^l + {}^0s_1)] \\ &= (1 - c_1) \underbrace{({}^1t_0^u - {}^0t_0^u)}_{\Delta t_0^u} + c_1 \underbrace{({}^1t_0^l - {}^0t_0^l)}_{\Delta t_0^l} \\ &\quad + (1 - c_1) \underbrace{({}^1p_1 - {}^0p_1)}_{\epsilon_1^u} + c_1 \underbrace{({}^1s_1 - {}^0s_1)}_{\lambda_1^u + \epsilon_1^u} \end{aligned}$$

where:

- $\epsilon_1^u = {}^1p_1 - {}^0p_1$  (Intrinsic parameter mismatch)
- $\lambda_1^u = {}^1s_1 - {}^0s_1 - \epsilon_1^u$  (Net parameter mismatch)

After rearranging the terms we get the upper signal time difference at MUX1 as:

$$\Delta t_1^u = \Delta t_0^u + c_1 (-\Delta t_0^u + \Delta t_0^l) + \epsilon_1^u + c_1 \lambda_1^u$$

We then define the Delta Difference at MUX0 as:

$${}^1\Delta_0 - {}^0\Delta_0 = \underbrace{({}^1t_0^u - {}^0t_0^u)}_{\Delta t_0^u} - \underbrace{({}^1t_0^l - {}^0t_0^l)}_{\Delta t_0^l}$$

where:

- ${}^1\Delta_0 = {}^1t_0^u - {}^1t_0^l$  (Upper-lower difference for PUF1 at MUX0)

- ${}^0\Delta_0 = {}^0t_0^u - {}^0t_0^l$  (Upper-lower difference for PUF0 at MUX0)

The difference between upper-lower signal mismatches for PUF1 and PUF0 at MUX0:

$${}^1\Delta_0 - {}^0\Delta_0 = \Delta t_0^u - \Delta t_0^l$$

Now after elimination of  $\Delta t_0^l$ , We solve for  $\Delta t_0^l$ , From the first equation:

$$\Delta t_0^l = \Delta t_0^u - ({}^1\Delta_0 - {}^0\Delta_0)$$

Replace  $\Delta t_0^l$  in the MUX1 equation:

$$\Delta t_1^u = \Delta t_0^u + c_1 (-\Delta t_0^u + [\Delta t_0^u - ({}^1\Delta_0 - {}^0\Delta_0)]) + \epsilon_1^u + c_1 \lambda_1^u$$

$$\Delta t_1^u = \Delta t_0^u + c_1 (-{}^1\Delta_0 + {}^0\Delta_0) + \epsilon_1^u + c_1 \lambda_1^u$$

This equation no longer explicitly depends on  $\Delta t_0^l$

The delay at the  $i$ -th multiplexer is recursively defined as:

$$\Delta t_i^u = \Delta t_{i-1}^u + c_i (-{}^1\Delta_{i-1} + {}^0\Delta_{i-1}) + \epsilon_i^u + c_i \lambda_i^u$$

- $\Delta t_{i-1}^u$ : Delay from the previous multiplexer ( $i-1$ ).
- $c_i$ : Challenge bit for the  $i$ -th MUX ( $c_i \neq 0$ ).
- ${}^1\Delta_{i-1} - {}^0\Delta_{i-1}$ : Difference in upper-lower signal mismatches between PUF1 and PUF0 at MUX  $i-1$ .
- $\epsilon_i^u = {}^1p_i - {}^0p_i$ : Intrinsic mismatch in parameter  $p_i$ .
- $\lambda_i^u = {}^1s_i - {}^0s_i - \epsilon_i^u$ : Net mismatch between parameters  $s_i$  and  $p_i$ .

**BASE CASE (MUX0)**:- The initial delay depends only on intrinsic and parameter mismatches:

$$\Delta t_0^u = \epsilon_0^u + c_0 \lambda_0^u$$

**For  $i = 1$ :**

$$\Delta t_1^u = \Delta t_0^u + c_1 (-{}^1\Delta_0 + {}^0\Delta_0) + \epsilon_1^u + c_1 \lambda_1^u$$

Substitute  $\Delta t_0^u = \epsilon_0^u + c_0 \lambda_0^u$ :

$$\Delta t_1^u = \epsilon_0^u + c_0 \lambda_0^u + c_1 (-{}^1\Delta_0 + {}^0\Delta_0) + \epsilon_1^u + c_1 \lambda_1^u$$

**For  $i = 2$ :**

$$\Delta t_2^u = \Delta t_1^u + c_2 (-{}^1\Delta_1 + {}^0\Delta_1) + \epsilon_2^u + c_2 \lambda_2^u$$

Substitute  $\Delta t_1^u$  from above:

$$\Delta t_2^u = \epsilon_0^u + c_0 \lambda_0^u + \epsilon_1^u + c_1 \lambda_1^u + c_1 (-{}^1\Delta_0 + {}^0\Delta_0) + c_2 (-{}^1\Delta_1 + {}^0\Delta_1) + \epsilon_2^u + c_2 \lambda_2^u$$

**Generalized Delay Formula for MUX  $i$**

$$\Delta t_i^u = \sum_{k=1}^i c_k (-{}^1\Delta_{k-1} + {}^0\Delta_{k-1}) + \sum_{k=0}^i (\epsilon_k^u + c_k \lambda_k^u)$$

- $\sum_{k=1}^i c_k (-{}^1\Delta_{k-1} + {}^0\Delta_{k-1})$ : Cumulative signal adjustments.
- $\sum_{k=0}^i (\epsilon_k^u + c_k \lambda_k^u)$ : Total intrinsic + parameter mismatches.

### Definitions of $\Delta_{k-1}$ for simple arbiter PUF

For PUF1 ( $n = 1$ ):

$${}^1\Delta_{k-1} = \sum_{j=0}^{k-1} ({}^1w_j \cdot {}^1x_j) + {}^1\beta_{k-1}$$

For PUF0 ( $n = 0$ ):

$${}^0\Delta_{k-1} = \sum_{j=0}^{k-1} ({}^0w_j \cdot {}^0x_j) + {}^0\beta_{k-1}$$

Coefficient  $w_i$ :

$$\begin{aligned} {}^nw_0 &= {}^n\alpha_0, \\ {}^nw_i &= {}^n\alpha_i + {}^n\beta_{i-1} \quad (i > 0). \end{aligned}$$

Parameters  $\alpha_i$  and  $\beta_i$ :

$$\begin{aligned} {}^n\alpha_i &= \frac{1}{2} ({}^np_i - {}^nq_i + {}^nr_i - {}^ns_i), \\ {}^n\beta_i &= \frac{1}{2} ({}^np_i - {}^nq_i - {}^nr_i + {}^ns_i). \end{aligned}$$

where  $n = 0$  (PUF0) or  $n = 1$  (PUF1).

Product  $x_j$ :

$$x_j = \prod_{k=j}^{a-1} (1 - 2c_k) \quad (\text{Product of challenge bits from MUX } j \text{ to } a-1).$$

After substituting the values the expression becomes:-

$$\Delta t_i^u = \underbrace{\sum_{a=1}^i c_a \left( \sum_{m=0}^{a-1} x_m \Delta w_m + \Delta \beta_{a-1} \right)}_{\text{Signal Adjustments}} + \underbrace{\sum_{a=0}^i (\epsilon_a^u + c_a \lambda_a^u)}_{\text{Intrinsic + Parameter Mismatches}}$$

**Key Terms:**

- $\Delta w_m = {}^0w_m - {}^1w_m$ : Difference in coefficients  $w_m$  between PUF0 and PUF1.
- $\Delta \beta_{a-1} = {}^0\beta_{a-1} - {}^1\beta_{a-1}$ : Difference in fixed parameters  $\beta_{a-1}$ .
- $\epsilon_a^u = {}^1p_a - {}^0p_a$ : Intrinsic mismatch in parameter  $p_a$ .
- $\lambda_a^u = {}^1s_a - {}^0s_a - \epsilon_a^u$ : Net mismatch between  $s_a$  and  $p_a$ .

Substituting  $i = 1$  into the formula:

$$\begin{aligned} \Delta t_1^u &= c_1 (x_0 \Delta w_0 + \Delta \beta_0) + \epsilon_0^u + c_0 \lambda_0^u + \epsilon_1^u + c_1 \lambda_1^u, \\ &= \text{Matches recursive derivation.} \end{aligned}$$

**Step 1: Substitute**  $x_m = \prod_{k=m}^{a-1} (1 - 2c_k)$   
 Replace  $x_m$  with the product term:

$$\Delta t_i^u = \underbrace{\sum_{a=1}^i c_a \left( \sum_{m=0}^{a-1} \left( \prod_{k=m}^{a-1} (1 - 2c_k) \right) \Delta w_m + \Delta \beta_{a-1} \right)}_{\text{Term 1}} + \underbrace{\sum_{a=0}^i \epsilon_a^u}_{\text{Term 2}} + \underbrace{\sum_{a=0}^i c_a \lambda_a^u}_{\text{Term 3}}$$

**Step 2: Expand Term 1** Expand the nested sums in Term 1:

$$\text{Term 1} = \sum_{a=1}^i \sum_{m=0}^{a-1} c_a \left( \prod_{k=m}^{a-1} (1 - 2c_k) \right) \Delta w_m + \sum_{a=1}^i c_a \Delta \beta_{a-1}$$

**Step 3: Separate  $c$ -Dependent Terms** Group all terms involving  $c$ :

$$\begin{aligned} \Delta t_i^u &= \underbrace{\sum_{a=0}^i c_a \lambda_a^u + \sum_{a=1}^i c_a \Delta \beta_{a-1}}_{\text{Linear Terms}} \\ &\quad + \underbrace{\sum_{a=1}^i \sum_{m=0}^{a-1} c_a \left( \prod_{k=m}^{a-1} (1 - 2c_k) \right) \Delta w_m}_{\text{Product Terms}} + \underbrace{\sum_{a=0}^i \epsilon_a^u}_b \end{aligned}$$

**Step 4: Define  $\psi^u$**  Construct the vector  $\psi^u$  to include both linear and product terms:

$$\psi^u(c) = \begin{bmatrix} c_0 \\ c_1 \\ \vdots \\ c_i \\ \sum_{m=0}^0 c_1 \prod_{k=m}^0 (1 - 2c_k) \\ \sum_{m=0}^1 c_2 \prod_{k=m}^1 (1 - 2c_k) \\ \vdots \\ \sum_{m=0}^{i-1} c_i \prod_{k=m}^{i-1} (1 - 2c_k) \end{bmatrix} \in \mathbb{R}^{(i+1) + \frac{i(i+1)}{2}}$$

where the product terms span all valid  $(a, m)$  pairs with  $0 \leq m \leq a - 1 \leq i - 1$ .

**Step 5: Construct Matrix  $Q$**  Matrix  $Q$  maps  $\psi^u$  to the delay formula:

$$Q = \begin{bmatrix} \underbrace{\lambda_0^u}_{\text{For } c_0} & \underbrace{\Delta \beta_0 + \lambda_1^u}_{\text{For } c_1} & \cdots & \underbrace{\Delta \beta_{i-1} + \lambda_i^u}_{\text{For } c_i} & \underbrace{\Delta w_0}_{\text{For } c_1(1-2c_0)} & \underbrace{\Delta w_0}_{\text{For } c_2(1-2c_0)(1-2c_1)} & \underbrace{\Delta w_1}_{\text{For } c_2(1-2c_1)} & \cdots \end{bmatrix}$$

Each  $\Delta w_m$  corresponds to a product term in  $\psi^u$ .

Combine  $Q$ ,  $\psi^u$ , and  $b$ :

$$\Delta t_i^u = Q^u \cdot \psi^u(c) + b^u \quad \text{where} \quad \left\{ \begin{array}{l} \psi^u(c) = \begin{bmatrix} c_0 \\ c_1 \\ \vdots \\ c_i \\ \frac{\sum_{m=0}^0 c_1 \prod_{k=m}^0 (1-2c_k)}{\sum_{m=0}^1 c_2 \prod_{k=m}^1 (1-2c_k)} \\ \vdots \\ \frac{\sum_{m=0}^{i-1} c_i \prod_{k=m}^{i-1} (1-2c_k)}{\sum_{m=0}^i c_{i+1} \prod_{k=m}^i (1-2c_k)} \end{bmatrix} \in \mathbb{R}^{(i+1) + \frac{i(i+1)}{2}}, \\ Q^u = [\lambda_0^u \quad \Delta\beta_0 + \lambda_1^u \quad \cdots \quad \Delta\beta_{i-1} + \lambda_i^u \quad \Delta w_0 \quad \Delta w_0 \quad \Delta w_1 \quad \cdots], \\ b^u = \sum_{a=0}^i \epsilon_a^u. \end{array} \right.$$

### Step 6: Final Expression

Combine  $P^l$ ,  $\psi^l$ , and  $b^l$ :

$$\Delta t_i^l = P^l \cdot \psi^l(c) + b^l \quad \text{where} \quad \left\{ \begin{array}{l} \psi^l(c) = \begin{bmatrix} c_0 \\ c_1 \\ \vdots \\ c_i \\ \frac{\sum_{m=0}^0 c_1 \prod_{k=m}^0 (1-2c_k)}{\sum_{m=0}^1 c_2 \prod_{k=m}^1 (1-2c_k)} \\ \vdots \\ \frac{\sum_{m=0}^{i-1} c_i \prod_{k=m}^{i-1} (1-2c_k)}{\sum_{m=0}^i c_{i+1} \prod_{k=m}^i (1-2c_k)} \end{bmatrix} \in \mathbb{R}^{(i+1) + \frac{i(i+1)}{2}}, \\ P^l = [\lambda_0^l \quad \Delta\beta_0 + \lambda_1^l \quad \cdots \quad \Delta\beta_{i-1} + \lambda_i^l \quad \Delta w_0 \quad \Delta w_0 \quad \Delta w_1 \quad \cdots], \\ b^l = \sum_{a=0}^i \epsilon_a^l. \end{array} \right.$$

### Step 7: General XOR Formula for $K = 2$

The generalized XOR function is:

$$f = \frac{1 + (-1)^{K+1} \prod_i \text{sign}(\mathbf{w}_i^\top \mathbf{x})}{2}$$

Substituting  $K = 2$ , we get:

$$f = \frac{1 - \text{sign}(\Delta t_i^u) \cdot \text{sign}(\Delta t_i^l)}{2}$$

### Step 8: Matching to Linear Sign Function Form

We aim to express the XOR output as:

$$f = \frac{1 + \text{sign}(\tilde{\mathbf{W}}^\top \tilde{\phi}(\mathbf{c}) + \tilde{b})}{2}$$

To match:

$$\frac{1 - \text{sign}(\Delta t_i^u) \cdot \text{sign}(\Delta t_i^l)}{2} = \frac{1 + \text{sign}(-\Delta t_i^u \cdot \Delta t_i^l)}{2} = \frac{1 + \text{sign}(\tilde{\mathbf{W}}^\top \tilde{\phi}(\mathbf{c}) + \tilde{b})}{2}$$

So we define:

$$\tilde{\mathbf{W}}^\top \tilde{\phi}(\mathbf{c}) + \tilde{b} = -\Delta t_i^u \cdot \Delta t_i^l$$

Hence:

$$\tilde{\phi}(\mathbf{c}) = [\psi_j \psi_k, \psi_l]^\top, \quad \tilde{\mathbf{W}} = \text{constructed via expansion of } -(Q^u \cdot \psi(\mathbf{c}) + b^u)(P^l \cdot \psi(\mathbf{c}) + b^l), \quad \tilde{b} = -b^u b^l$$

Final Result

$$f = \frac{1 + \text{sign}(-\Delta t_i^u \cdot \Delta t_i^l)}{2} = \frac{1 + \text{sign}(\tilde{\mathbf{W}}^\top \tilde{\phi}(\mathbf{c}) + \tilde{b})}{2}$$

Computation of  $\phi(\mathbf{c})$  from  $\psi(\mathbf{c})$

Let  $\psi(\mathbf{c}) = [\psi_1, \psi_2, \dots, \psi_n]^T$ . We compute the product:

$$\Delta t_i^u \cdot \Delta t_i^l = (Q^u \cdot \psi + b^u)(P^l \cdot \psi + b^l)$$

Expanding the terms:

$$\begin{aligned} \Delta t_i^u \cdot \Delta t_i^l &= (Q^u \cdot \psi)(P^l \cdot \psi) + b^u(P^l \cdot \psi) + b^l(Q^u \cdot \psi) + b^u b^l \\ &= \sum_{j=1}^n \sum_{k=1}^n Q_j^u P_k^l \psi_j \psi_k + b^u \sum_{k=1}^n P_k^l \psi_k + b^l \sum_{j=1}^n Q_j^u \psi_j + b^u b^l \end{aligned}$$

Therefore, the feature map  $\tilde{\phi}(\mathbf{c})$  is:

$$\tilde{\phi}(\mathbf{c}) = \begin{bmatrix} \psi_1 \psi_1 \\ \psi_1 \psi_2 \\ \vdots \\ \psi_j \psi_k \\ \vdots \\ \psi_n \psi_n \\ \psi_1 \\ \vdots \\ \psi_n \end{bmatrix} \in \mathbb{R}^{[n(n+1)/2] + n}$$

## 2 DIMENSIONALITY

It includes all second-order pairwise products and all first-order terms of  $\psi(\mathbf{c})$ .

For  $i = 7$ :

$$\dim(\psi^u) = (i+1) + \frac{i(i+1)}{2} = 8 + \frac{7 \cdot 8}{2} = 8 + 28 = 36$$

$$\text{Let } n = \dim(\psi^u) = 36$$

$$\dim(\tilde{\phi}) = \frac{n(n+1)}{2} + n = \frac{36 \cdot 37}{2} + 36 = 666 + 36 = \boxed{702}$$

### 3 Kernel SVM

To replicate our model using a kernel SVM without explicit feature mapping, a polynomial kernel of degree 3 would be sufficient:

$$K(x, x') = (\gamma x^\top x' + r)^3$$

Choose  $\gamma = 1$ ,  $r = 0$ . This kernel generates all 3rd, 2nd, and 1st-order interactions, matching our  $\phi(x)$  structure.

### 4 Inversion Simple Arbiter

We have an arbitrary buffer with 4 delays  $P, Q, R$ , and  $S$ , where the direct delays are  $P$  and  $Q$ , and the cross delays are  $R$  and  $S$ . The task is to determine the non-negative delays given the models of the 64-bit buffers. In total, there are  $64 \times 4 = 256$  variables, but only 65 equations (64 weights and one bias). This leaves us with 164 variables and only 65 equations.

The idea is to set all values of  $R$  and  $S$  to 0, and all values of  $Q$  to 0 except for  $q_{64}$ . This reduces the problem to 64 values of  $P$  and one value of  $Q$ , totaling 65 variables. Now, with 65 equations and 65 variables, we can solve for all 64 values of  $P$  and the 64<sup>th</sup> value of  $Q$ .

However, some of these values may be negative, which violates the constraint that delays must be non-negative ( $d \geq 0$ ). To address this we can do two things **1.** We add a constant to all  $P, Q, R$ , and  $S$  values to ensure they become non-negative. This constant is the absolute value of the largest negative number among the  $P$  values and  $q_{64}$ . **2.** We add a constant to only the  $P$  and  $Q$  values to ensure they become non-negative (others are already non-negative with 0 value). Similarly, this constant will also be the absolute value of the largest negative number among the  $P$  values and  $q_{64}$ .

**Note:** Since all the  $W$ s have equal number of + and - sign delays ( $p, q, r, s$ , variable), so adding a constant i.e.  $p_i + \epsilon_i, q_i + \epsilon_i, r_i + \eta_i, s_i + \eta_i$ , in delays won't affect the  $W$ s. Our model remains the same.

Here is the generalized model for the delay after passing through 64 responses:

$$\Delta_{64} = w_1 \cdot x_1 + w_2 \cdot x_2 + \dots + w_{64} \cdot x_{64} + b$$

The generalized weight  $w_i$  in terms of all delays is given by:

$$\begin{aligned} w_1 &= \alpha_1 \\ w_i &= \alpha_i + \beta_{i-1} \quad (\text{for } i = 2, 3, \dots, 64) \\ b &= \beta_{64} \end{aligned}$$

The delays are defined as:

$$\begin{aligned} d_i &\stackrel{\text{def}}{=} (1 - 2c_i) \\ \alpha_i &\stackrel{\text{def}}{=} (p_i - q_i + r_i - s_i)/2 \\ \beta_i &\stackrel{\text{def}}{=} (p_i - q_i - r_i + s_i)/2 \end{aligned}$$

### Demonstration with a 2-bit Response

For simplicity, consider a 2-bit response where all  $Q$  and  $S$  values are set to 0, except for  $q_2$  and all  $P_s$ . This leaves us with 3 equations and 3 variables,  $p_1, p_2$ , and  $q_2$ .

$$\begin{aligned} w_1 &= \alpha_1 = \frac{(p_1 - q_1 + r_1 - s_1)}{2} \\ w_2 &= \alpha_2 + \beta_1 = \frac{(p_2 - q_2 + r_2 - s_2)}{2} + \frac{(p_1 - q_1 - r_1 + s_1)}{2} \\ w_3 &= b = \beta_2 = \frac{(p_2 - q_2 - r_2 + s_2)}{2} \end{aligned}$$

The system can be represented in matrix form as:

$$\begin{bmatrix} w_1 \\ w_2 \\ w_3 \end{bmatrix} = \begin{bmatrix} 1 & 0 & -1 & 0 & 1 & 0 & -1 & 0 \\ 1 & 1 & -1 & -1 & -1 & 1 & 1 & 1 \\ 0 & 1 & 0 & -1 & 0 & -1 & 0 & 1 \end{bmatrix} \begin{bmatrix} p_1 \\ p_2 \\ q_1 \\ q_2 \\ r_1 \\ r_2 \\ s_1 \\ s_2 \end{bmatrix}$$

Simplifying, we get:

$$w_1 = p_1$$

$$w_2 = p_2$$

$$w_3 - p_2 = q_2$$

Suppose we obtain  $p_1 = -2$ ,  $p_2 = 1$ , and  $q_2 = -3$ . To ensure all values are non-negative, we can do two things as stated above.

1. Add 3 (the absolute value of the largest negative number, which is  $-3$ ) to all  $P, Q, R$ , and  $S$  values. The final adjusted values will be:

$$\begin{array}{cccccccc} p_1 & p_2 & q_1 & q_2 & r_1 & r_2 & s_1 & s_2 \\ 1 & 4 & 0 & 3 & 3 & 3 & 3 & 3 \end{array}$$

2. Same as above, but just add 3 (the absolute value of the largest negative number, which is  $-3$ ) only to all  $P$  and  $Q$  values as all other ( $R$  and  $S$ ) values are already non-negative (kept 0 in the starting itself). Then the final adjusted values will be:

$$\begin{array}{cccccccc} p_1 & p_2 & q_1 & q_2 & r_1 & r_2 & s_1 & s_2 \\ 1 & 4 & 0 & 3 & 0 & 0 & 0 & 0 \end{array}$$

For our solution, we have chosen the second option.

## 5 CODE for ML-PUF

## 6 CODE for Invertor-PUF

## 7 EXPERIMENTS ON LinearSVC and Logistic Regression

### SVM Loss Comparison

Loss Function	Training Time (s)	Test Accuracy (%)
Hinge Loss	2.33	95.75
Squared Hinge Loss	2.22	99.19

Table 1: Comparison of Hinge and Squared Hinge loss functions for SVM.

### Effect of Regularization Parameter (C) on Accuracy and Time



<b>C Value</b>	<b>Model</b>	<b>Training Time (s)</b>	<b>Test Accuracy (%)</b>
0.0001 (Low)	LinearSVC	0.197	83.06
	LogisticRegression	0.269	69.19
0.01 (Medium)	LinearSVC	0.862	100.00
	LogisticRegression	0.782	96.75
10 (High)	LinearSVC	24.387	100.00
	LogisticRegression	0.934	100.00

Table 2: Effect of different regularization strengths (C values) on SVM and Logistic Regression.