

# Overview of Algorithms & Data Structures

Chaiyaporn Khemapatapan

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# Outline

## Lecture I

1. Motivation
2. Sorting algorithms
3. Complexity analysis
4. Linear data structures

## Lecture II

5. Nonlinear data structures
6. Abstract data types
7. Dijkstra's algorithm
8. Summary

# Attention!

Lecture aimed at non-computer scientists.

Focus is on explaining concepts,  
rather than technical correctness.

# 1. Motivation

# Motivation

- Algorithms
- Data Structures

1. **Everything** running on your computer is an **algorithm**
2. Analysing them is paramount to **writing, maintaining** and **improving** them
3. Several **tools** exist to help achieve this

# Motivation

- Algorithms
- Data Structures

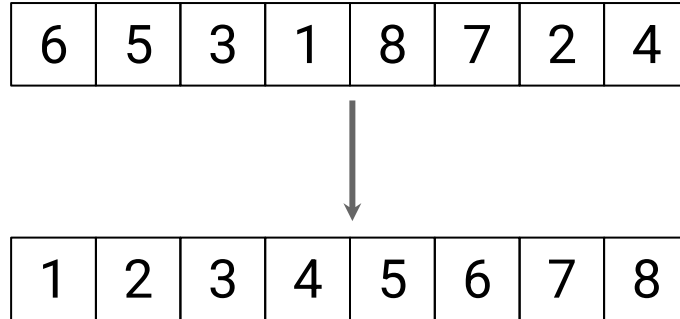
1. Data Structures define how data is **stored** in **RAM**
2. Many variations, each with **advantages** and **disadvantages**
3. Strongly coupled to algorithmic **complexity**

## 2. Sorting algorithms



# Sorting

- Suppose we have some **unsorted list**
- We want to make it **sorted**





# Insertion sort

- In pseudocode:

$i \leftarrow 1$   
• “Naive” sorting algorithm  
**while**  $i < \text{length}(A)$

$N$  steps

- One-by-one, take each

element and move it  
**while**  $j > 0$  and  $A[j-1] > A[j]$   $N$  steps

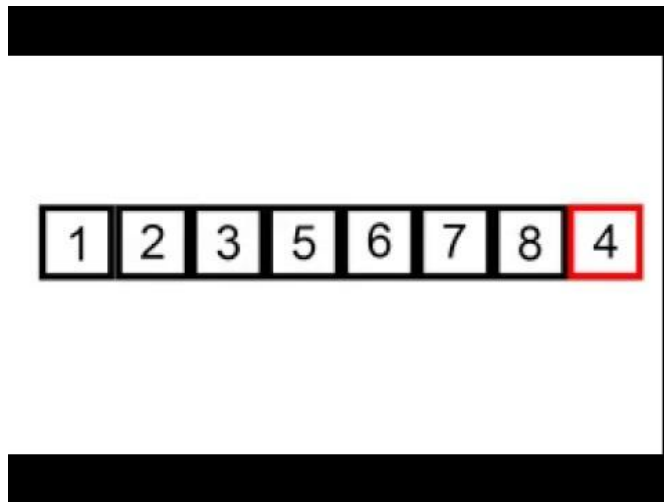
- When all elements have been

moved, list is sorted!

**end while**

$i \leftarrow i + 1$

**end while**



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Total:  $N^2$  steps

# Bubble sort

- Traverse the list, taking **pairs** of elements  **$N$  steps**
- **Swap** if order incorrect
- **Repeat**  $N$  times  **$N$  steps**
- Now it's **sorted!**

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Total:  **$N^2$  steps**

# Intermezzo: Divide and conquer

- Generic algorithm **strategy**
  - **Divide** the problem into smaller parts
  - **Solve** (conquer) the problem for each part
  - **Recombine** the parts
- Straightforward to **parallelise**
- Closely related to **map-reduce**
- Has been advocated by Caesar, Machiavelli, Napoleon...

# Merge sort

- Much **smarter** sort
  - **Split** the dataset into chunks
  - **Sort** each chunk
  - **Merge** the chunks back together
- Example of **divide-and-conquer**
- Splitting & sorting takes  **$\log_2(N)$  steps**
- Merging takes  **$N$  steps**

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Total:  **$N \log(N)$  steps**

# 3. Complexity analysis

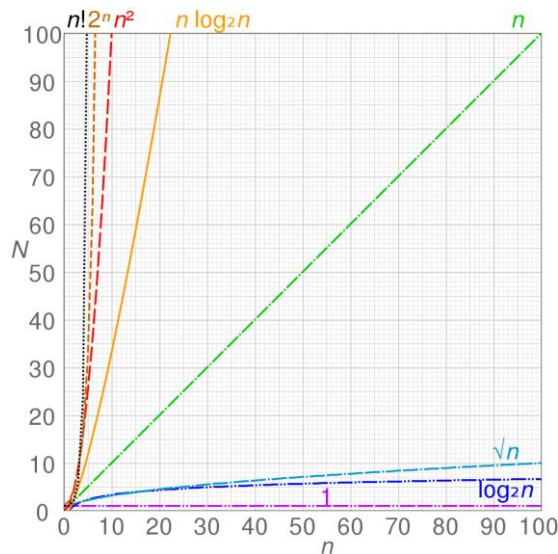


# Complexity – O

- We use “**Big-O**” notation
- Represents an **upper bound**
- Ignores **constants**
- Only shows **dominant term**
- In these examples: N is amount of **data**
- Linear:  **$O(N)$** 
  - $O(N) = O(2*N) = O(k*N + m)$   
for any constants  $k, m$
- Quadratic:  **$O(N^2)$** 
  - $O(N^2) = O(a*N^2 + b*N + c)$   
for any constants  $a, b, c$

# Complexity – O

- Roughly **three categories**, in **decreasing order**:
  - Exponential –  $O(k^N)$
  - Polynomial –  $O(N^k)$
  - Polylogarithmic –  $O(\log(N)^k)$
- This is an **abstraction**!
  - Does not *directly* relate to runtimes
  - A good  $O(N^2)$  algorithm may be faster than a bad  $O(\log(N))$  one
  - Depends on your input data!



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# Complexity – O

- Remember: it's an **upper bound**!
- And it's a **property** (not an equivalence relation)!
  - $O(N) = O(N^2)$
  - But  $O(N^2) \neq O(N)$
- $O(N) = O(2N)$
- $2N = O(N)$
- $N^2 + N = O(N^2)$
- $O(N^2) + O(N) = O(N^2)$
- $O(\log(N)) + O(N) = O(N)$
- $O(1) \rightarrow$  used for **constant time**



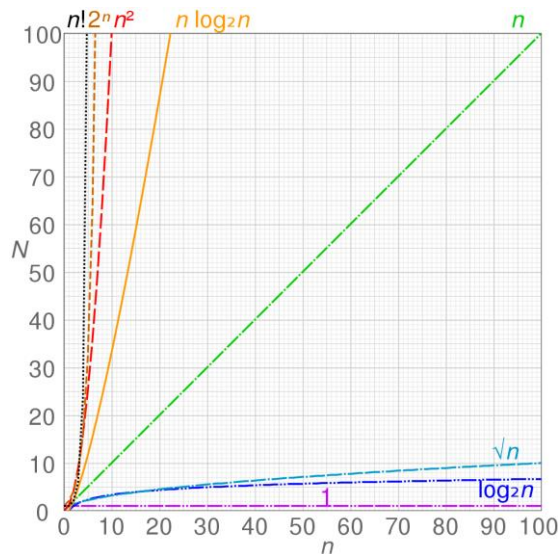
# Complexity – O

- Formal definition:

$$f(N) = O(g(N)) \leftrightarrow$$

$$\exists N_0, M : \forall N > N_0 : f(N) \leq M g(N)$$

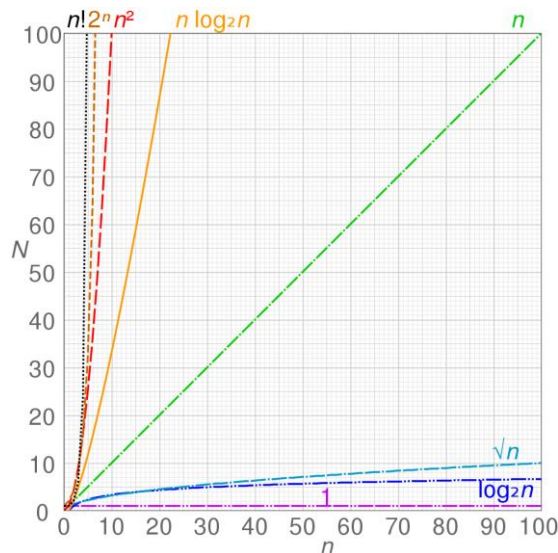
- In words: starting from  $N_0$ ,  $f$  is bounded from above by  $g$  (up to some constant  $M$ )



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# Complexity – $\Omega$

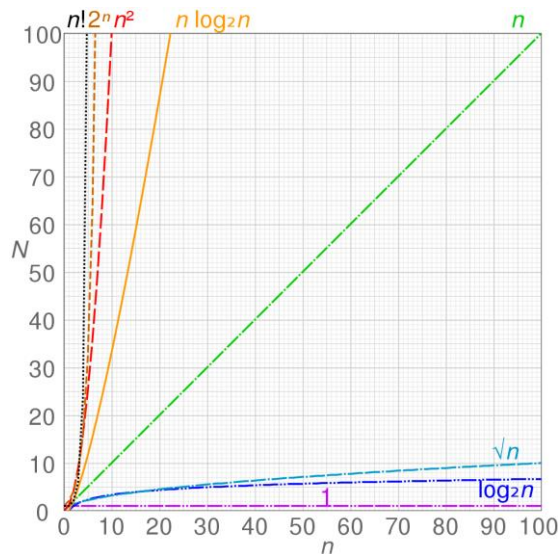
- $\Omega$ : **lower** bound
  - Formal definition:  
$$f(N) = \Omega(g(N)) \leftrightarrow$$
$$\exists N_0, M : \forall N > N_0 : f(N) \geq M g(N)$$
  - In words: starting from  $N_0$ ,  $f$  is bounded from below by  $g$  (up to some constant  $M$ )
- Examples:
  - $N = \Omega(\log(N))$
  - $N^2 = \Omega(N)$



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# Complexity – $\Theta$

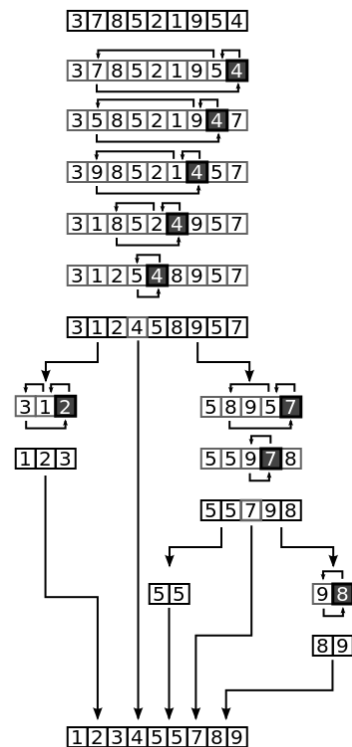
- $\Theta$ : **exact** bound
  - Formal definition:
$$f(N) = \Theta(g(N)) \leftrightarrow$$
$$f(N) = O(g(N)) \text{ and } f(N) = \Omega(g(N))$$
- Examples:
  - $N = \Theta(N)$
  - $2N = \Theta(N)$
  - $3N^2 + 5N + 1 = \Theta(N^2)$
  - $N \neq \Theta(N^2)$  (even though  $N = O(N^2)$ )



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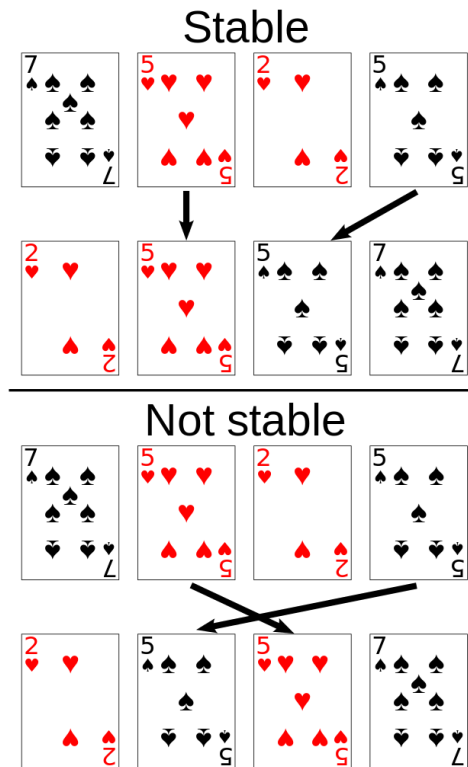
# One more sorting example: Quicksort

- Pick an element, called **pivot**
- **Partitioning**: reorder the array so that the pivot is in the correct place
- **Recursively** apply the above steps to the sub-arrays on either side of the pivot
- **Randomised-quicksort**: select the pivot **randomly**







# Stable sorting

- A sorting algorithm is **stable** iff it **conserves the order** of equal elements



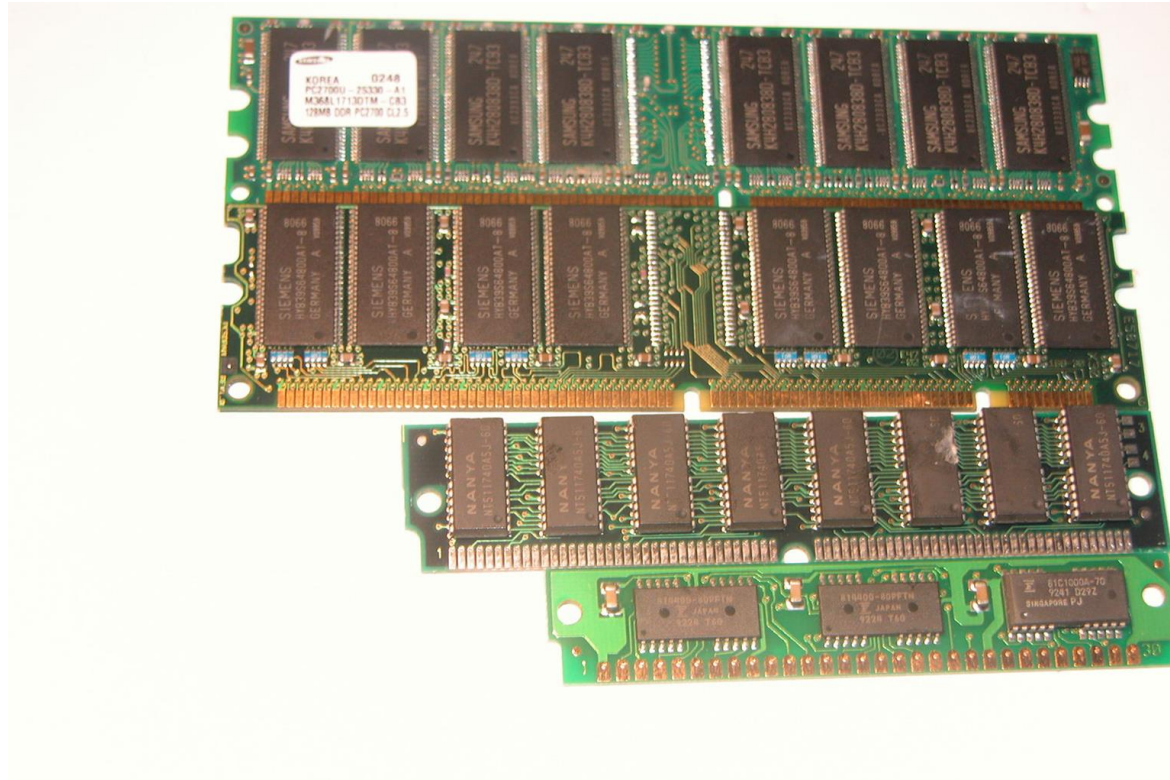
# Comparison of algorithms

Algorithm	Stable?	Complexity
Insertion sort		$O(N^2)$
Bubble sort		$O(N^2)$
Merge sort		$O(N \log(N))$
Quicksort		?

# 4. Linear data structures



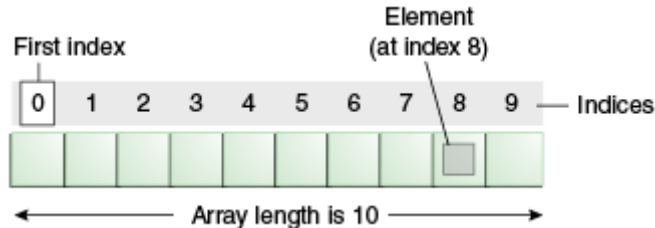
# Memory





# Arrays

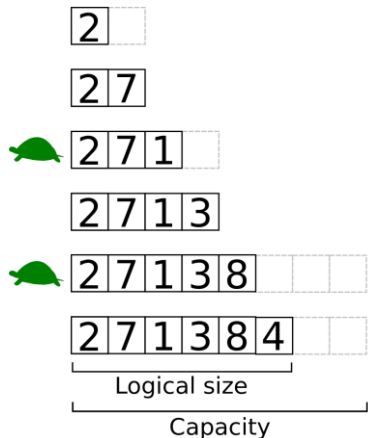
- **Linear, contiguous** list of data
- Accessible by **index**
- **Fixed-size**
  - $N * d$
- **Supported** by all major systems
- Back-insert/remove:  **$O(1)$**
- Random insert/remove:  **$O(N)$**
- Index-lookup:  **$O(1)$**
- Lookup:  **$O(N)$**



# Dynamic arrays

- **Linear, contiguous** list of data
- Accessible by **index**
- **Resizable**
- Back-insert/remove:  $O(1)^*$
- Random insert/remove:  $O(N)$
- Index-lookup:  $O(1)$
- Lookup:  $O(N)$

\*Amortised.



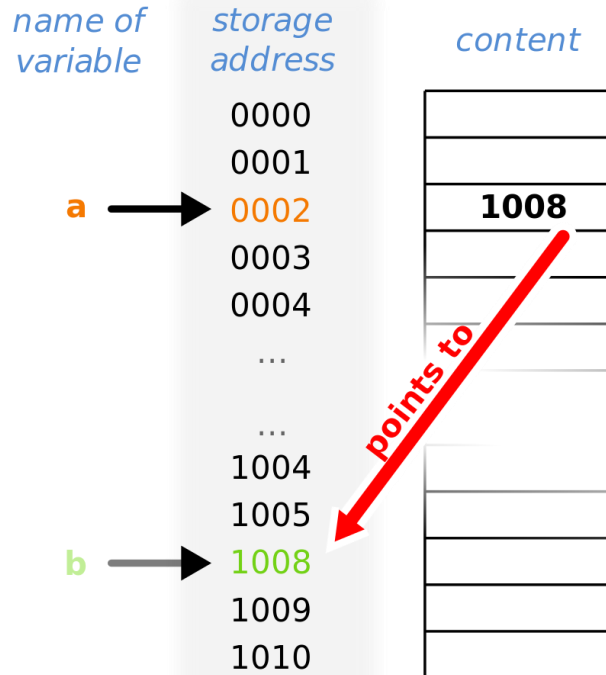
C++: **`std::vector`**

Python: **`list`**

C#: **`System.Collections.ArrayList`**

Java: **`java.util.ArrayList`**

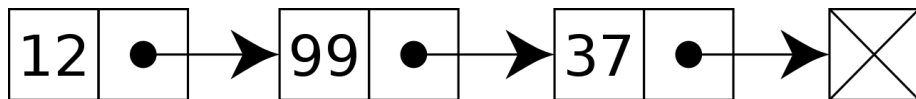
# Pointers



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# Linked list

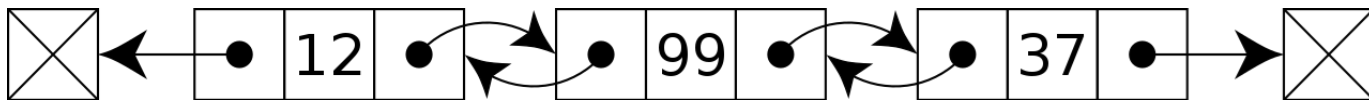
- **Linear, contiguous** list of data
- Accessible by **iteration**
- **Resizable**
- Back-insert/remove:  **$O(1)$**
- Random insert/remove:  **$O(N)$**
- Index-lookup:  **$O(N)$**
- Lookup:  **$O(N)$**



C++: `std::forward_list`

# Doubly Linked list

- Pointers **both ways**
- Uses **more memory**, but allows **iteration both ways**
- Back-insert/remove:  **$O(1)$**
- Random insert/remove:  **$O(N)$**
- Index-lookup:  **$O(N)$**
- Lookup:  **$O(N)$**



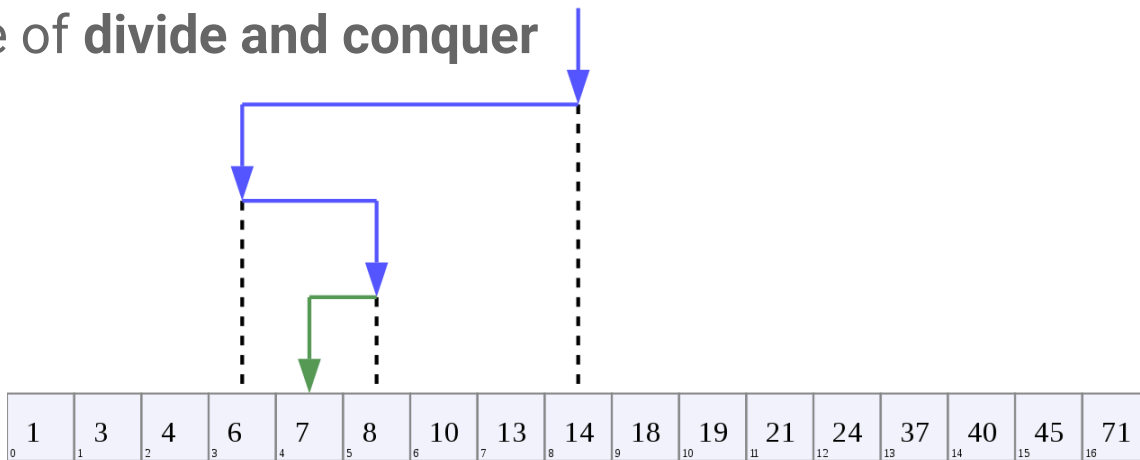
C++: **`std::list`**

C#: **`System.Collections.Generic.LinkedList`**

Java: **`java.util.LinkedList`**

# Binary search

- **Searches** a sorted linear data structure
- Takes  $\Theta(\log(N))$
- Example of **divide and conquer**



# Algorithms & Data Structures

## Lecture 2

L.J. Bel  
iCSC 2018

# Outline

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6. Abstract data types
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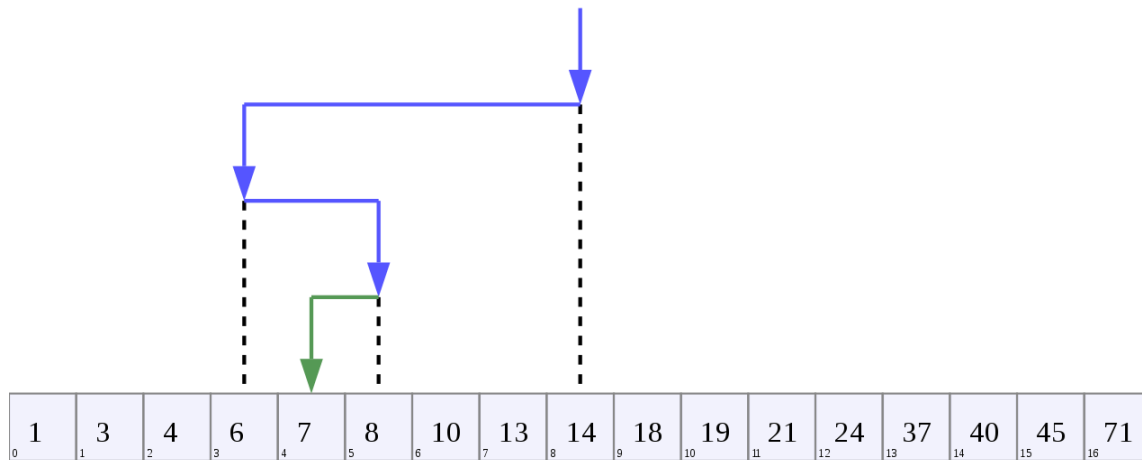


# 5. Nonlinear data structures



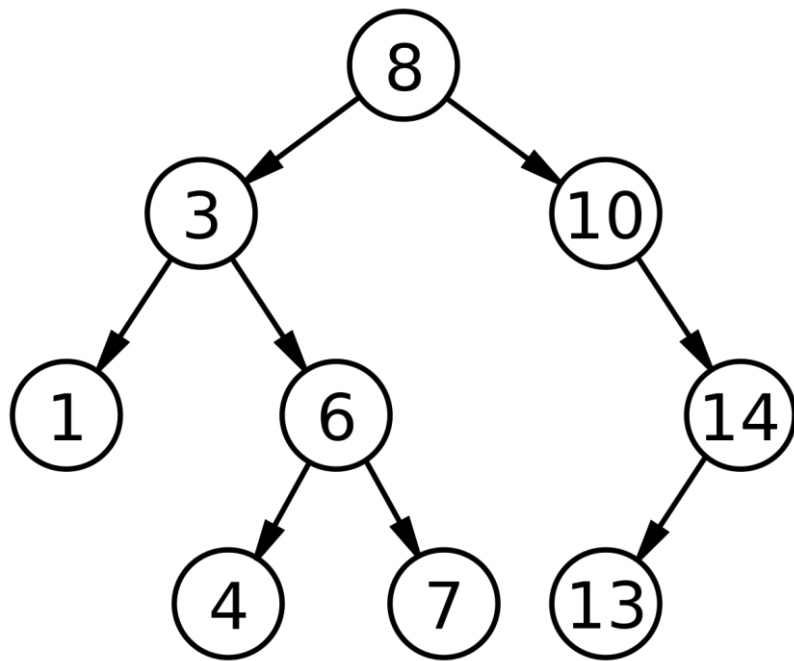
# Recall: Binary search

- **Searches** a sorted linear data structure
- Takes  $\Theta(\log(N))$
- ... let's use this as inspiration for a **data structure**!



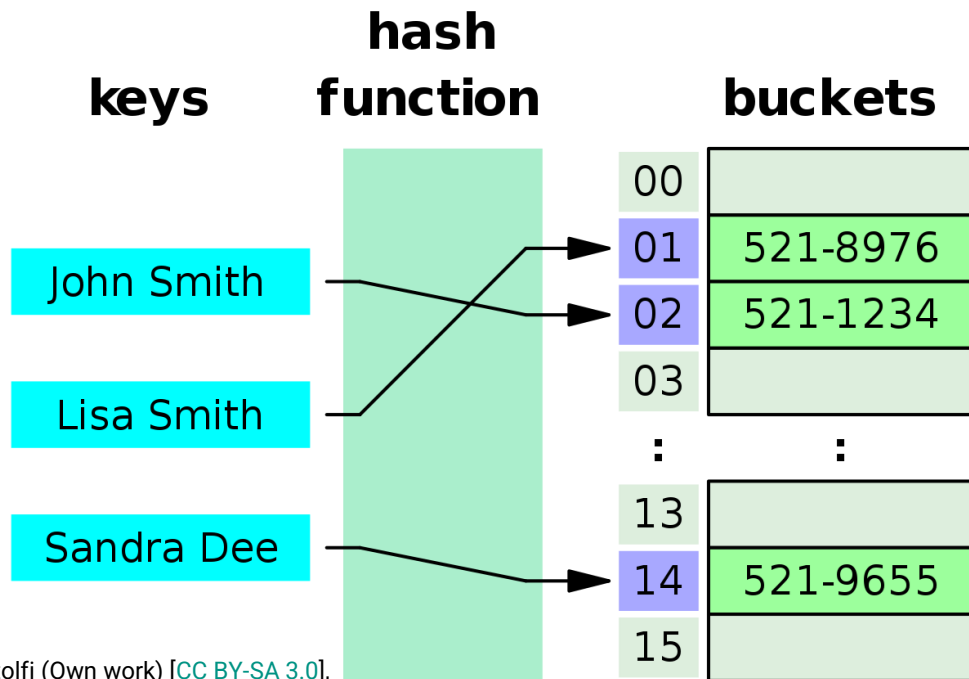
# Binary search trees

- **Tree** structure
- **Pointers** between nodes
  - To the right: only **larger**
  - To the left: only **smaller**
- Allows easy **sorted iteration**
- Search/insert/delete:  
all  **$O(\log(N))$**



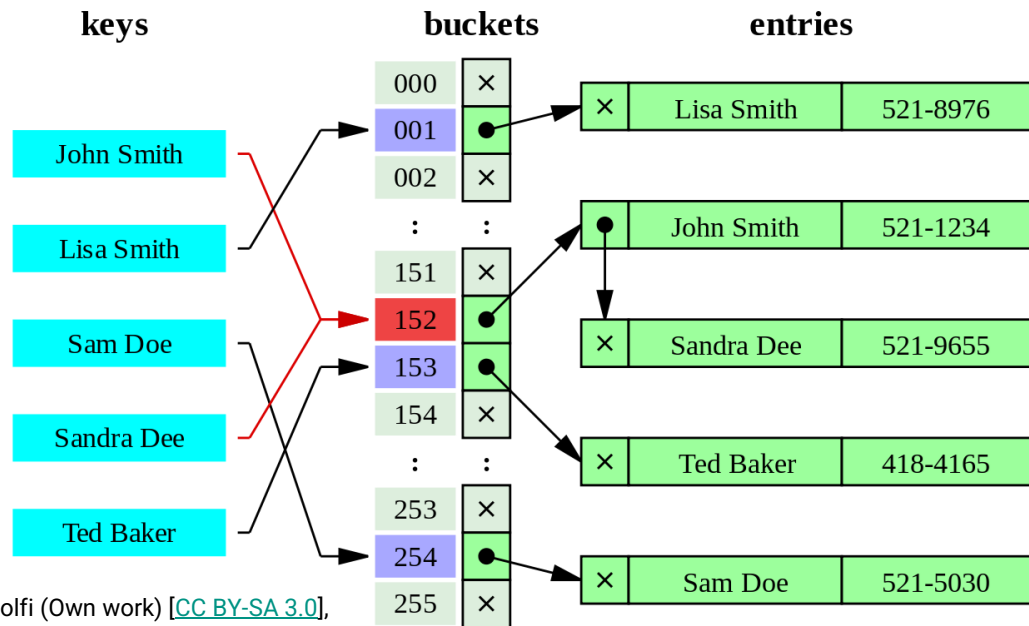
# Hash tables

- Idea: create **buckets** numbered 1 to B
- For each item, **compute** in which bucket it belongs
- **Put** the item in that bucket
- Search/insert/delete: all  **$O(1)$**



# Hash tables

- Problem: **clashing** hashes!
- Solution: replace entry with **linked list** (chaining)
- New problem: **load factor** can become **too high**!
- Solution: **copy** to new table with **more buckets**



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# Comparing data structures

Data structure → Operation ↓	Dynamic array	Linked list	Binary search tree	Hash table
Lookup	$O(N)$	$O(N)$	$O(\log(N))$	$O(1)$
Indexed lookup	$O(1)$	$O(N)$	N/A	N/A
Back-insert	$O(1)^*$	$O(1)$	$O(\log(N))$	$O(1)^*$
Random insert	$O(N)$	$O(N)$	N/A	N/A
Remove	$O(N)$	$O(N)$	$O(\log(N))$	$O(1)^*$

\*Amortised.

# 6. Abstract data types



# Why “Abstract”?

- Abstract Data Type (**ADT**) does **not** define a real data structure
  - Only defines an **interface**
  - **Implemented** using one of the “real” data structures
- Usually **limits** operations compared to actual DS
- Enhances **flexibility**

Related to several **core programming principles**:

- Program against the **interface**, not the **implementation**!
- Use **high cohesion, loose coupling**
- **Separate the concerns**



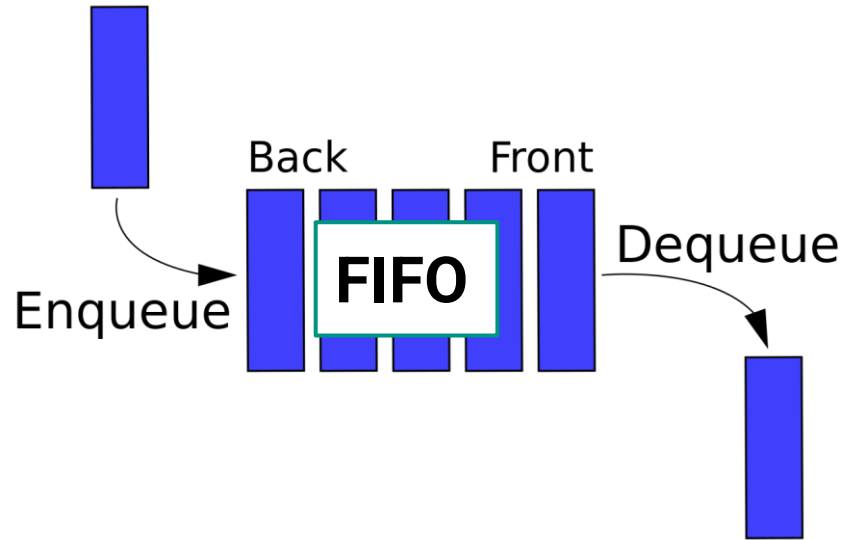
# Queue

Operations:

- **Enqueue:** add item to beginning of queue
- **Dequeue:** retrieve and remove item from end of queue

Typical underlying data structure:

- **Linked list**
- **Dynamic array**



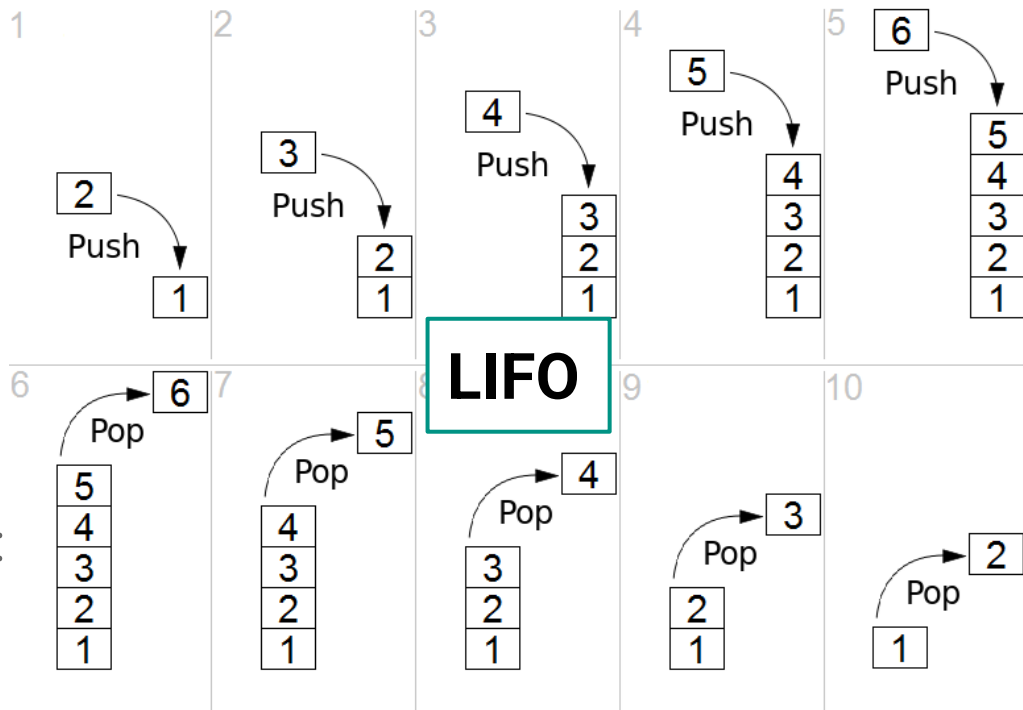
# Stack

Operations:

- **Push:** add item to top of stack
- **Pop:** retrieve and remove item from top of stack

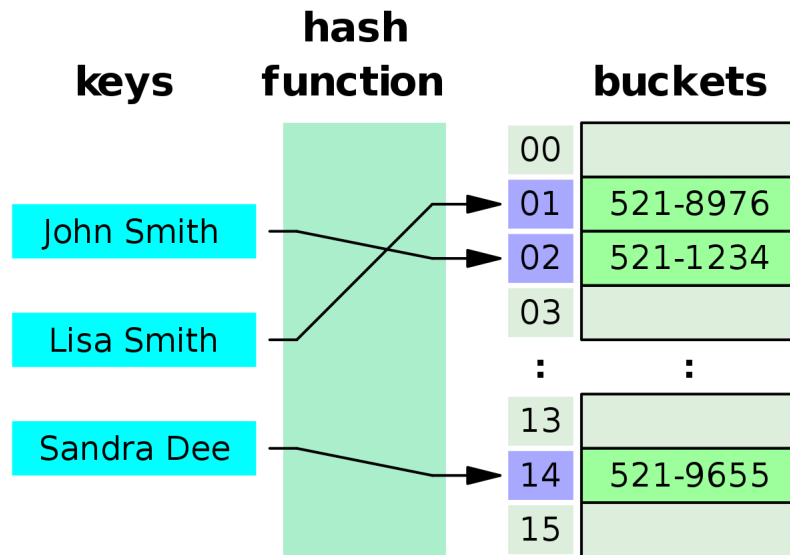
Typical underlying data structure:

- **Linked list**
- **Dynamic array**



# Map

- Map: dataset that maps (associates) **keys** to **values**
- Keys are **unique** (values need not be)
- Values can be **retrieved** by key
- Not indexed...
  - ...although an array could be seen as a map with **integer keys!**

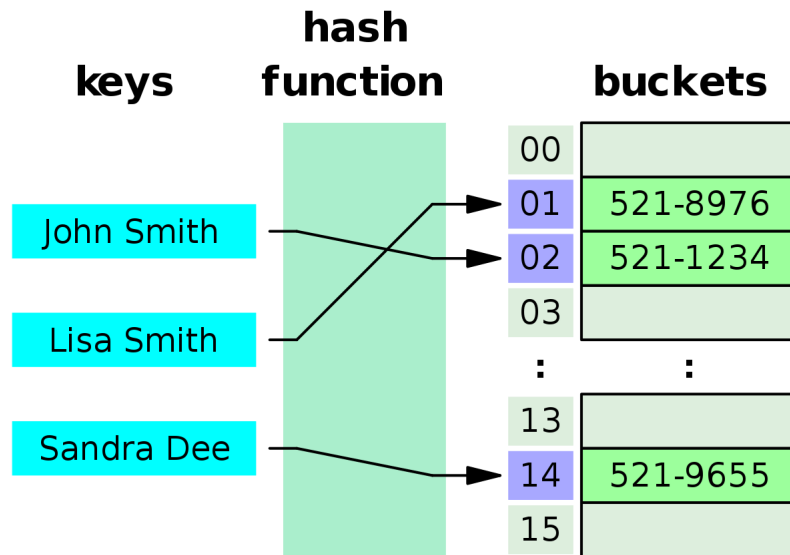


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# Map

Operations:

- **Lookup:** retrieve value for a key
- **Insert:** add key-value pair
- **Replace:** replace value for a specified key
- **Remove:** remove key-value pair



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# Map

Typical implementations:

- Binary Search Tree
  - Requires **sortable keys**
  - Can do **indexed/range** queries!
  - Fast with many **insertions**
- Hash Table
  - Generally very **fast**
  - **Space-efficient**
  - Need to keep **load factor** under control...

C++: **std::map**

C#: **System.Collections.Generic  
SortedSet**

Java: **java.util.TreeMap**

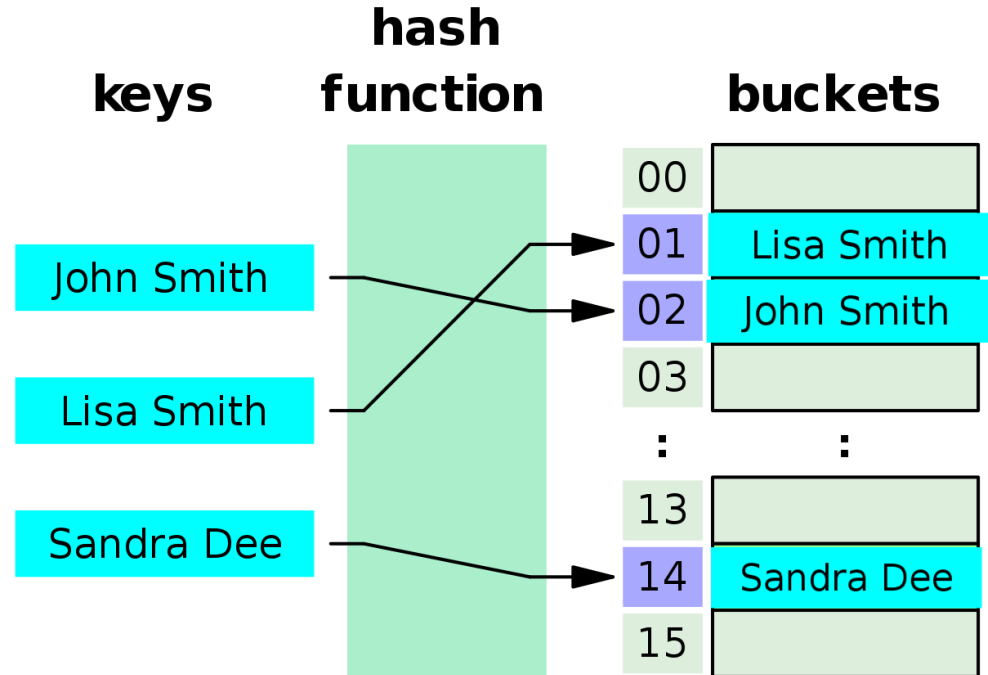
Python: **dict**

C++: **std::unordered\_map**

Java: **java.util.HashMap**

# Set

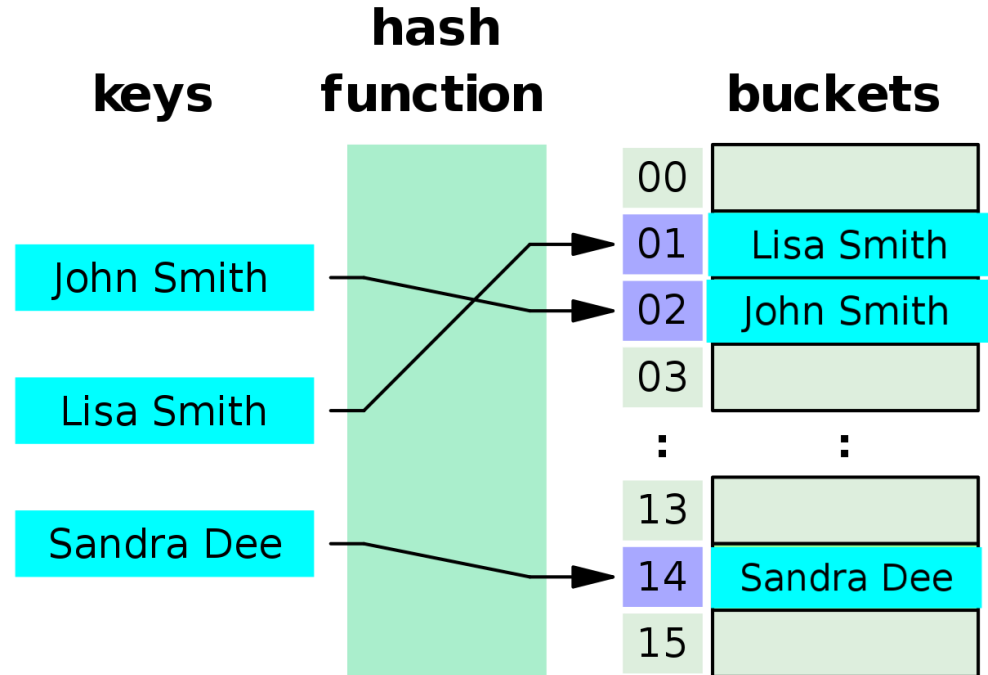
- Set: dataset that contains certain **values**
- No **ordering**, no **multiplicity**
- A value is either **present** or **not**



# Set

Operations:

- **Contains:** check whether a value is present
- **Add:** add a value
- **Remove:** remove a value



# Set

Typical implementations:

- Binary Search Tree
- Hash Table
- *Bloom filter*

C++: **std::set**

C#: **System.Collections.Generic.SortedSet**

Java: **java.util.TreeSet**

Python: **set** (and **frozenset**)

C++: **std::unordered\_set**

C#: **System.Collections.Generic.HashSet**

Java: **java.util.HashSet**



# Comparing ADTs

Abstract Data Type → Operation ↓	Queue	Stack	Map	Set
Lookup	N/A*	N/A*	By key	Contains
Add	Enqueue	Push	Key + value	Add
Replace	N/A	N/A	By key	N/A
Remove	Dequeue	Pop	By key	Remove

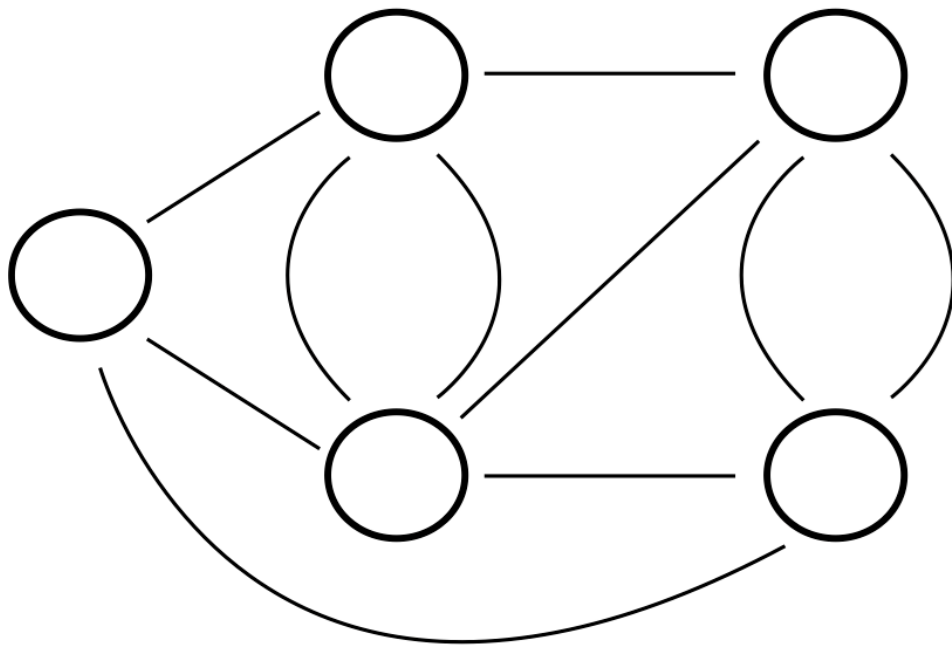
\*Only by removing element  
(some may support *peek*)

# 7. Dijkstra's algorithm



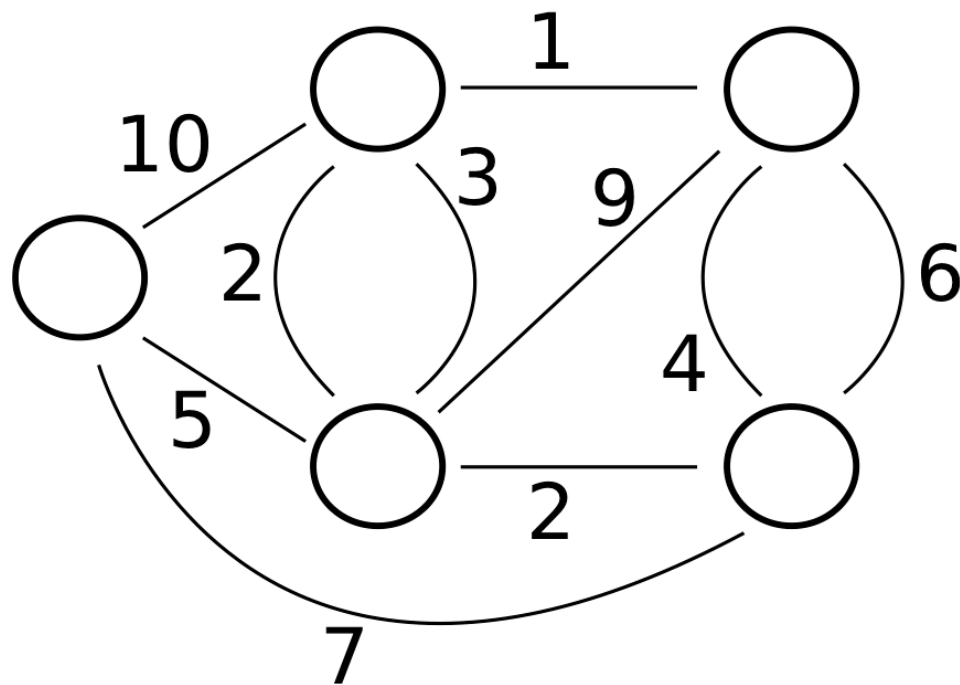
# Graphs

- Another data structure!
- Consists of **vertices** ( $V$ ) and **edges** ( $E$ )



# Graphs

- Another data structure!
- Consists of **vertices** ( $V$ ) and **edges** ( $E$ )
- Edges may carry a **weight**

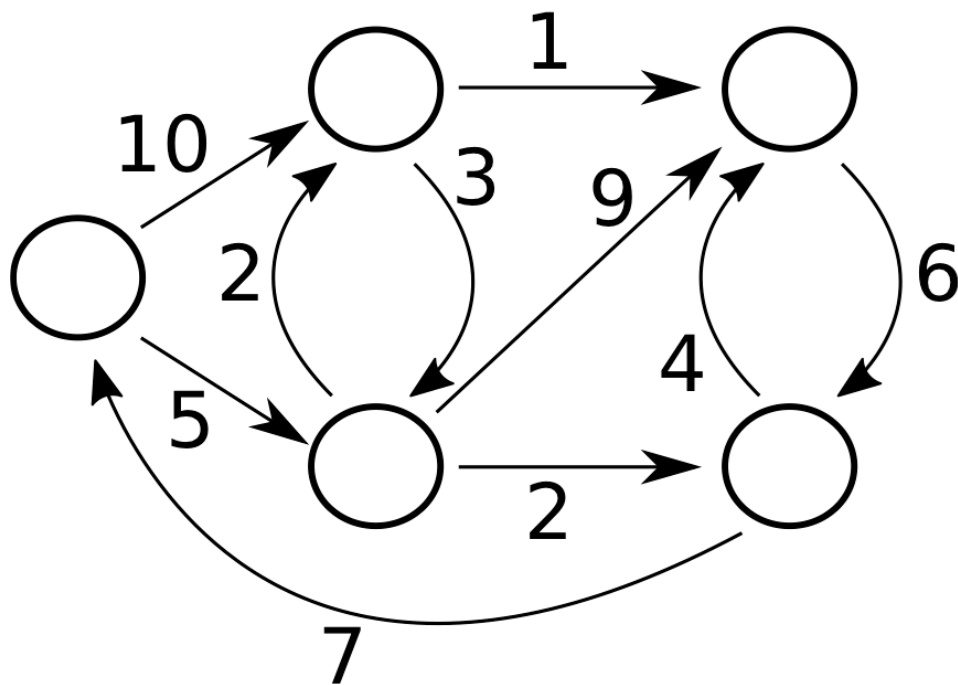


# Graphs

- Another data structure!
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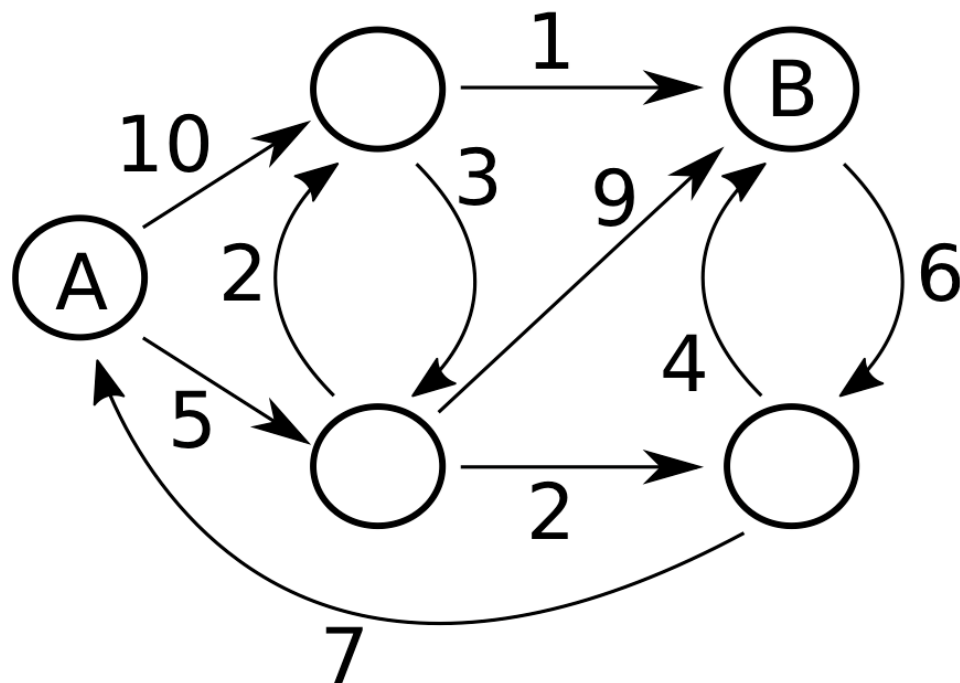
Directed graph:

- Edges are **directed**



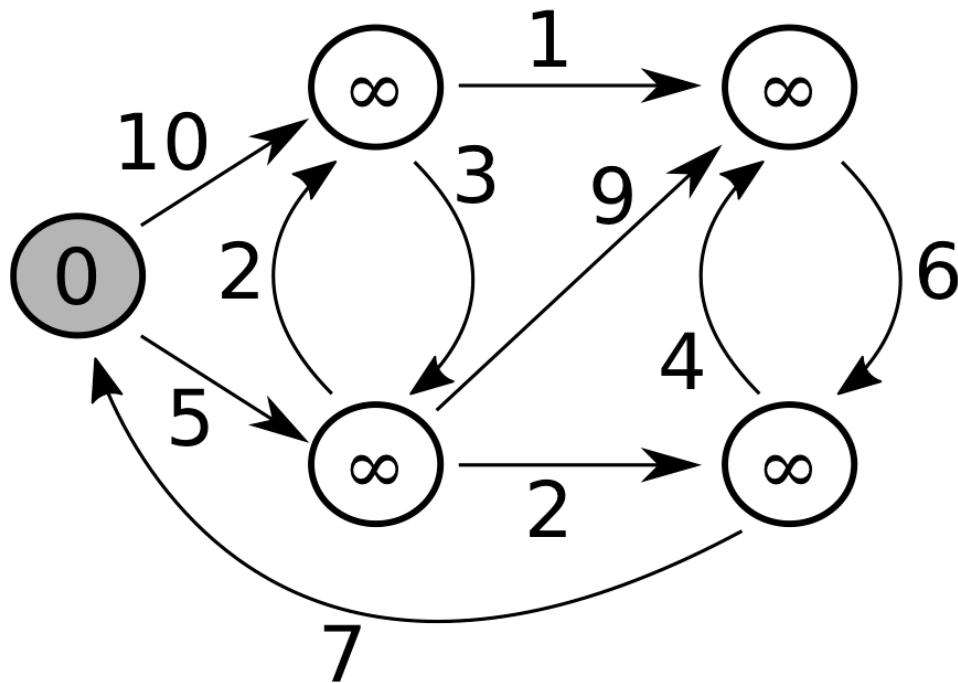
# Pathfinding

- Problem: find **shortest path** from A to B
- Shortest is defined as **lowest total edge weights**



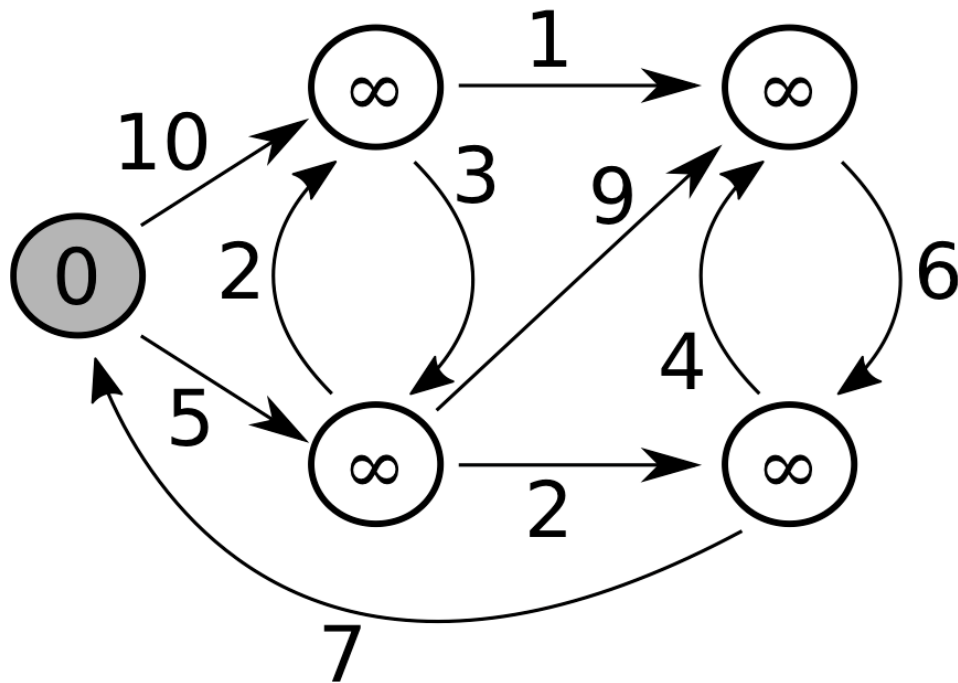
# Dijkstra's algorithm

- Algorithm to obtain **shortest path** from a given vertex to **any other** vertex
- Example of **greedy** algorithm
- Initially: set shortest-path **estimates** to **0** for start vertex and  $\infty$  for the others



# Dijkstra's algorithm

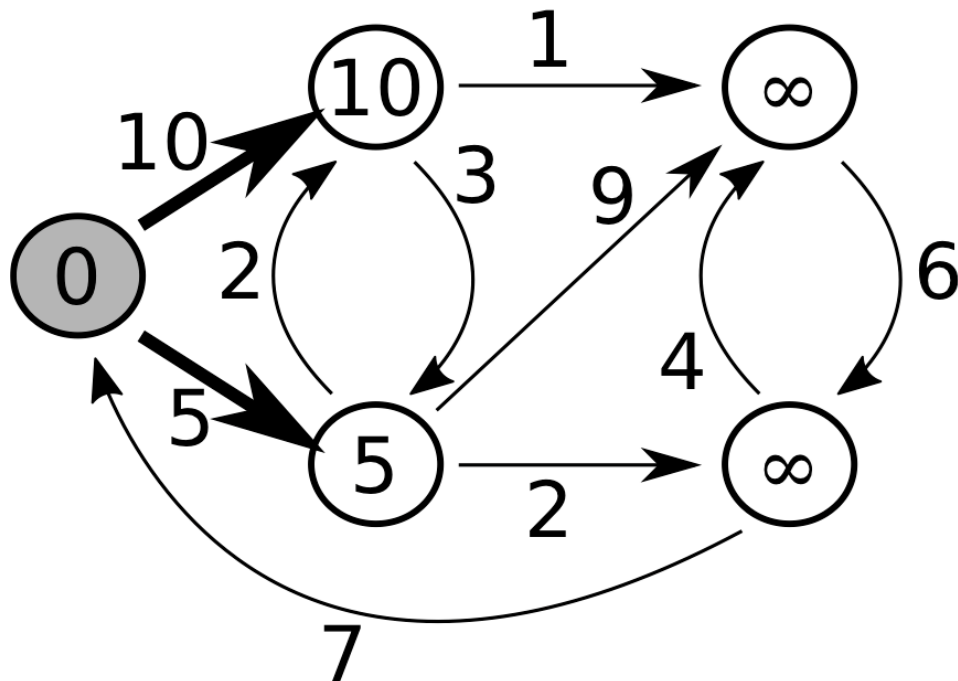
- Repeat the following:
  - Select unvisited vertex with **lowest** estimate
  - Look at paths to **unvisited** nodes
  - Update estimates if **lower** than previous estimate





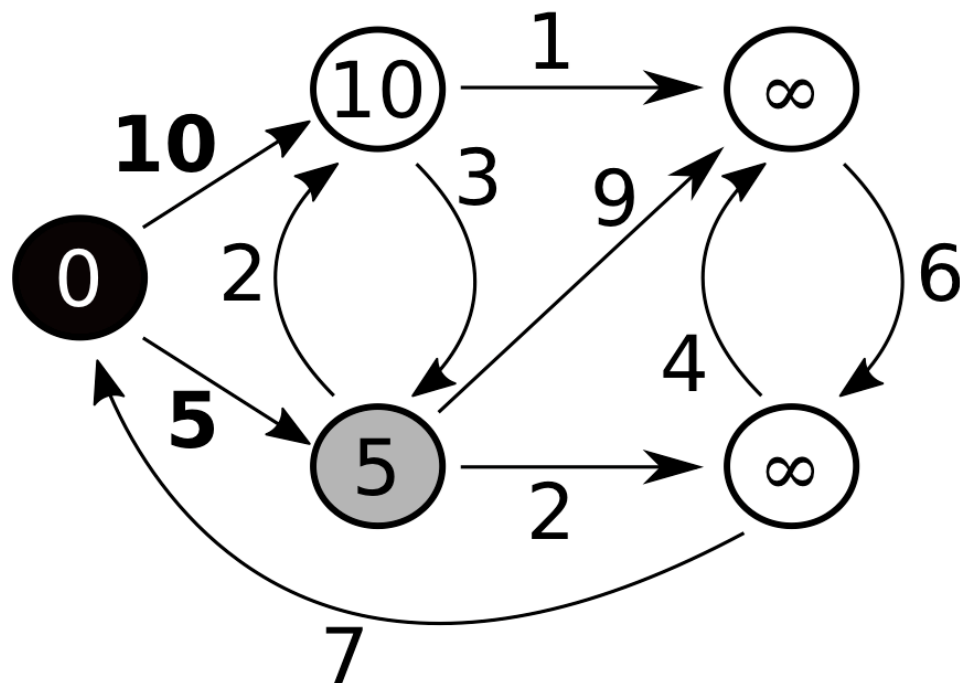
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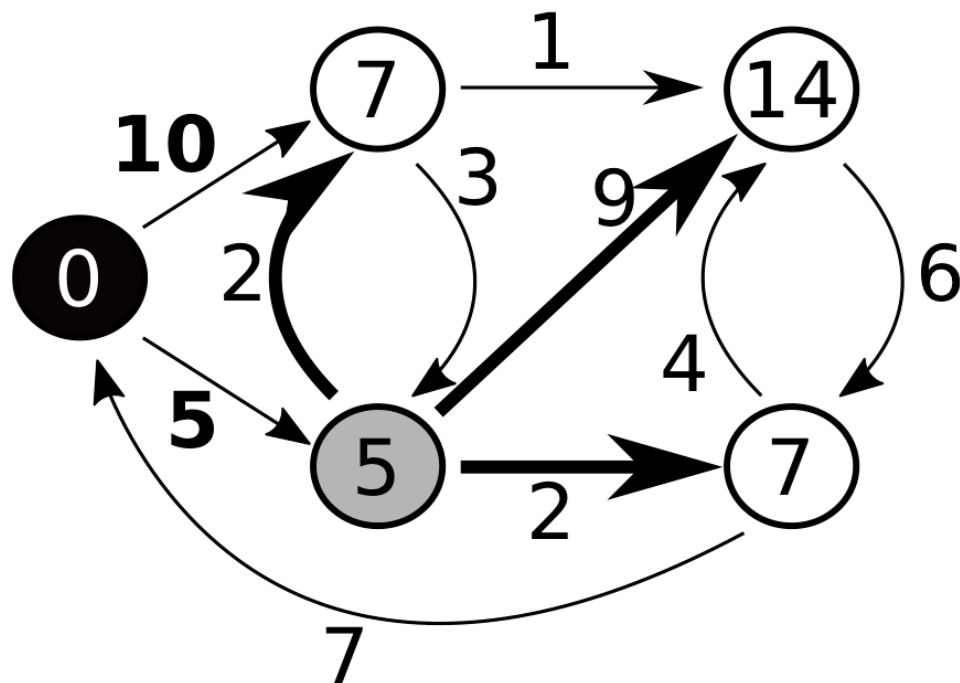
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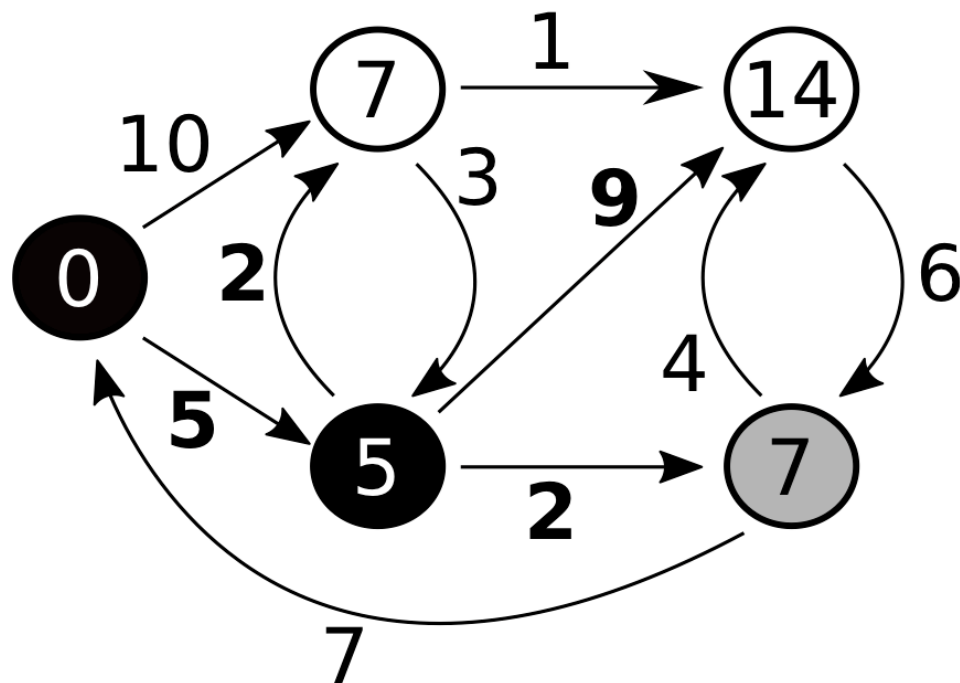
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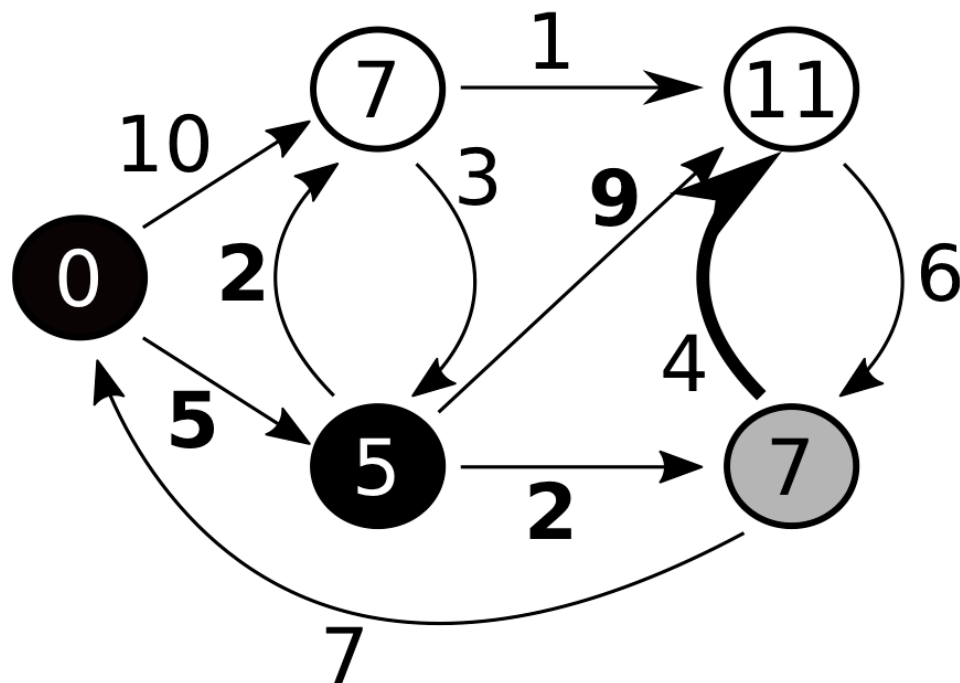
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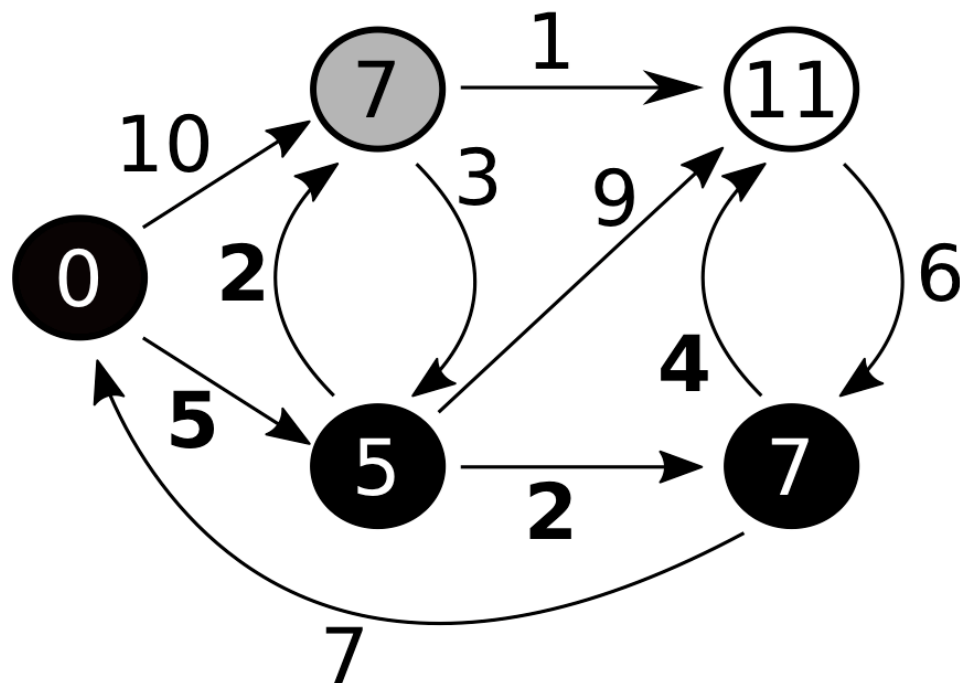
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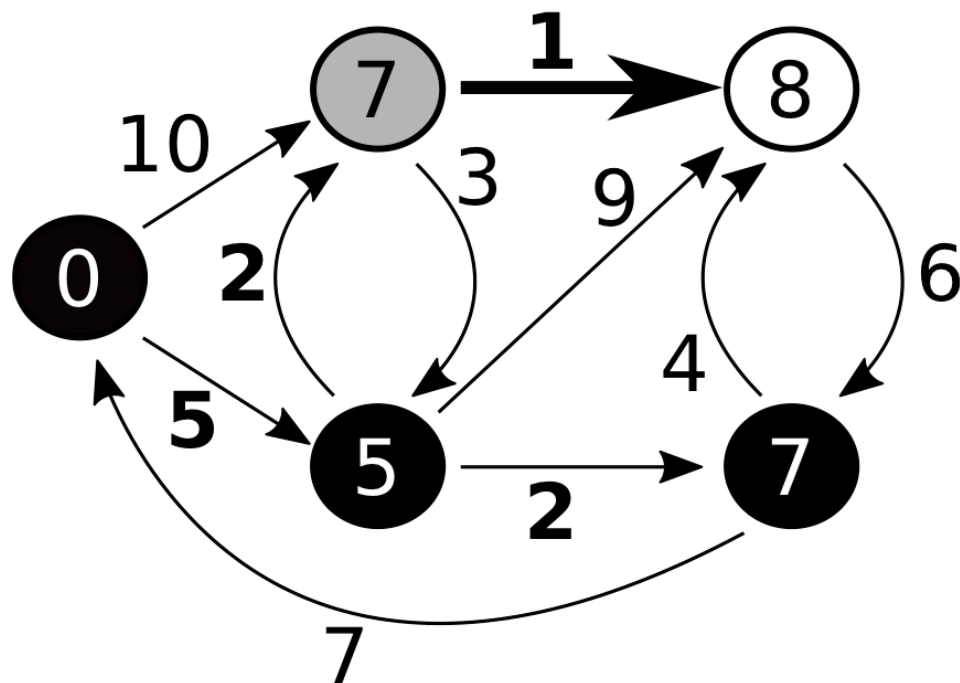
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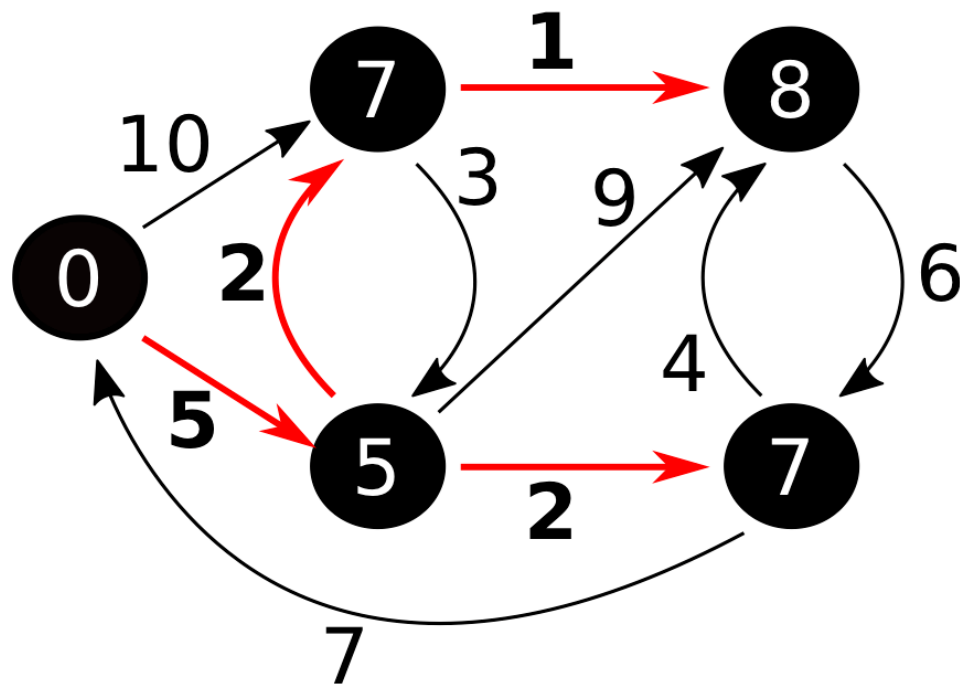
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# Dijkstra's algorithm

- Repeat the following:
  - Select unvisited vertex with **lowest** estimate
  - Look at paths to **unvisited** nodes
  - Update estimates if **lower** than previous estimate
- Shortest paths indicated in red
- Complexity:  $O(E + V \log V)$





# 8. Summary



# Summary

- Concepts

- Divide and conquer
- Complexity
  - $O, \Theta, \Omega$
- (Un)stable sorting
- Pointers

# Summary

- Concepts
- **Sorting algorithms**

- Insertion sort
- Bubble sort
- Merge sort
- Quicksort

# Summary

- Concepts
- Sorting algorithms
- Data structures

- Arrays
- Dynamic arrays
- (Doubly) linked lists
- Binary search trees
- Hash tables
- (Directed) graphs

# Summary

- Concepts
  - Sorting algorithms
  - Data structures
  - Abstract data types
- Queues
  - Stacks
  - Maps
  - Sets

# Summary

- Concepts
  - Sorting algorithms
  - Data structures
  - Abstract data types
  - Algorithms
- Binary search
  - Dijkstra's algorithm