## 1 Z Notation Introduction

The following subsections provide a high level overview of select properties of Z Notation based on "The Z Notation: A Reference Manual" by J. M. Spivey. A copy of this reference manual can be found at dave/docs/z/Z-notation reference manual.pdf. In many cases, definitions will be pulled directly from the reference manual and when this occurs, the relevant page number(s) will be included. For a proper introduction with tutorial examples, see chapter 1, "Tutorial Introduction" from pages 1 to 23. For the LaTeX symbols used to write Z, see the reference document found at dave/docs/z/zed-csp-documentation.pdf.

### 1.1 Decorations

The following decorations are used through this document and are taken directly from the reference manual. For a complete summary of the Syntax of Z, see chapter 6, Syntax Summary, startingon page 142.

```
[indicates final state of an operation]
[indicates input to an operation]
[indicates output of an operation]
[indicates output of an operation]
[indicates the schema results in a change to the state space]
[indicates the schema does not result in a change to the state space]
[indicates output of the left schema is input to the right schema]
```

## 1.2 Types

Objects have a type which characterizes them and distinguish them from other kinds of objects.

- Basic types are sets of objects which have no internal structure of interest meaning the concrete definition of the members is not relevant, only their shared type.
- Free types are used to describe (potentially nested and/or recursive) sets
  of objects. In the most simple case, a free type can be an enumeration of
  constants.

Within the xAPI Formal Specification, both of these types are used to describe the Inverse Functional Identifier property.

- Introduction of the basic types MBOX,  $MBOX\_SHA1SUM$ , OPENID and ACCOUNT allows the specification to talk about these constraints within the xAPI specification without defining their exact structure
- The free type IFI is defined as one of the above basic types meaning an object of type IFI is of type MBOX or  $MBOX\_SHA1SUM$  or OPENID or ACCOUNT.

Types can be composed together to form composite types and thus complex objects.

```
[MBOX, MBOX\_SHA1SUM, OPENID, ACCOUNT] \\ IFI ::= MBOX \mid MBOX\_SHA1SUM \mid OPENID \mid ACCOUNT \\
```

Within the xAPI Formal Specification, IFI is used within the definition of an agent as presented in the schema Agent.

```
\begin{array}{c} Agent \\ agent : AGENT \\ objectType : OBJECTTYPE \\ name : \mathbb{F}_1 \# 1 \\ ifi : IFI \\ \\ objectType = Agent \\ agent = \{ifi\} \cup \mathbb{P}\{name, objectType\} \end{array}
```

See section 2.2, pages 28 to 34, and chapter 3, pages 42 to 85, for more information about Schemas and the Z Language.

### 1.3 Sets

A collection of elements that all share a type. A set is characterized solely by which objects are members and which are not. Both the order and repetition of objects are ignored. Sets are written in one of two ways:

- listing their elements
- by a property which is characteristic of the elements of the set.

such that the following law from page 55 holds for some object y

$$y \in \{x_1, ..., x_n\} \iff y = x_1 \vee ... \vee y = x_n$$

## 1.4 Ordered Pairs

Two objects (x, y) where x is paired with y. An n-tuple is the pairing of n objects together such that equality between two n-tuple pairs is given by the law from page 55

$$(x_1, ..., x_n) = (y_1, ..., y_n) \iff x_1 = y_1 \land ... \land x_n = y_n$$

When ordered pairs are used with respect to application (as seen on page 60)

$$fx \Rightarrow f(x) \iff (x,y) \in f$$

which states that f(x) is defined if and only if there is a unique value y which result from fx Additionally, application associates to the left

$$fxy \Rightarrow (fx)y \Rightarrow (f(x), y)$$

meaning f(x) results in a function which is then applied to y.

## 1.5 Sequences

A collection of elements where their ordering matters such that

$$\langle a_1, ..., a_n \rangle \Rightarrow \{1 \mapsto a_1, ..., n \mapsto a_n\}$$

as seen on page 115. Additionally, iseq is used to describe a sequence whose members are distinct.

## 1.6 Bags

A collection of elements where the number of times an element appears in the collection is meaningful.

$$[a_1, ..., a_n] \Rightarrow \{a_1 \mapsto c_1, ..., a_n \mapsto c_n\}$$

As described on page 124 (replacing c with k), each element  $a_i$  appears  $c_i$  times in the list  $a_1, ..., a_n$  such that the number of occurances of  $a_i$  within bag A is returned by

$$count A a_i \equiv A \# a_i$$

# 1.7 Maps

This document introduces a named subcategory of sets, map of the free type KV, which are akin to sequences and bags. To enumerate the members of a map,  $\langle\langle ... \rangle\rangle$  is used but should not be confused with  $d_i \langle\langle E_i[T] \rangle\rangle$  within a Free Type definition. The distinction between the two usages is context dependent but in general, if  $\langle\langle ... \rangle\rangle$  is used outside of a constructor declaration within a Free Type definition, it should be assumed to represent a map.

$$KV ::= base \mid associate \langle \langle KV \times X \times Y \rangle \rangle$$

where

base [is a constant which is the empty  $KV \Rightarrow \langle \langle \rangle \rangle$ ] associate [is a constructor and is inferred to be an injection]

The full enumeration of all properties, constraints and functions specific to a map with type KV will be defined elsewhere but associate can be understood to (in the most basic case) operate as follows.

$$associate(base, x_i, y_i) = \langle \langle (x_i, y_i) \rangle \rangle \Rightarrow \langle \langle x_i \mapsto y_i \rangle \rangle$$

The enumeration of a map was chosen to be  $\langle ... \rangle$  as a map is a collection of injections such that if M is the result of  $associate(base, x_i, y_i)$  from above then

$$atKey(M, x_i) = y_i \iff x_i \mapsto y_i \land (x_i, y_i) \in M$$

# 1.8 Select Operations

The follow are defined in Chapter 4 (The Mathematical Tool-kit) within the reference manual and are used extensively throughout this document. In many cases, the functions listed here will serve as Operations in the context of Primitives and Algorithms.

#### 1.8.1 Functions

```
 \begin{array}{lll} & \rightarrow & [\text{relate each } x \in X \text{ to at most one } y \in Y, \text{ page } 105] \\ & \rightarrow & [\text{relate each } x \in X \text{ to exactly one } y \in Y, \text{ page } 105] \\ & \mapsto & [\text{map different elements of } x \text{ to different } y, \text{ page } 105] \\ & \mapsto & [\rightarrowtail \text{ that are also } \rightarrow, \text{ page } 105] \\ & \mapsto & [X \rightarrow Y \text{ where whole of } Y \text{ is the range, page } 105] \\ & \mapsto & [X \rightarrow Y \text{ whole of } X \text{ as domain and whole of } Y \text{ as range, page } 105] \\ & \mapsto & [\text{map } x \in X \text{ one-to-one with } y \in Y, \text{ page } 105] \\ & \times & Y = & \{f: X \leftrightarrow Y \mid (\forall x: X; y1, y2: Y \bullet \\ & & (x \mapsto y_1 \in f \ \land \ (x \mapsto y_2) \in f \Rightarrow y_1 = y_2))\} \\ & X \rightarrow Y = & \{f: X \rightarrow Y \mid \text{dom } f = X\} \\ & X \rightarrowtail Y = & \{f: X \rightarrow Y \mid (\forall x_1, x_2: \text{dom } f \bullet f(x_1) = f(x_2) \Rightarrow x_1 = x_2)\} \\ & X \rightarrowtail Y = & (X \rightarrowtail Y) \cap (X \rightarrow Y) \\ \end{array}
```

#### 1.8.2 Ordered Pairs, Maplet and Composition of Relations

 $X \twoheadrightarrow Y == \{ f : X \rightarrowtail Y \mid \operatorname{ran} f = Y \}$   $X \twoheadrightarrow Y == (X \twoheadrightarrow Y) \cap (X \rightarrow Y)$   $X \rightarrowtail Y == (X \twoheadrightarrow Y) \cap (X \rightarrowtail Y)$ 

```
first [returns the first element of an ordered pair, page 93] second [returns the second element of an ordered pair, page 93] 

→ [maplet is a graphic way of expressing an ordered pair, page 95] 
dom [set of all x \in X related to atleast one y \in Y by R, page 96] 
ran [set of all y \in Y related to atleast one x \in X by R, page 96] 
[The composition of two relationships, page 97] 

o [The backward composition of two relationships, page 97]
```

```
[X, Y] = first : X \times Y \to X
second : X \times Y \to Y
\forall x : X; y : Y \bullet
first(x, y) = x \land
second(x, y) = y
```

```
[X, Y] = \frac{[X, Y]}{\operatorname{dom} : (X \leftrightarrow Y) \to \mathbb{P} X}
\operatorname{ran} : (X \leftrightarrow Y) \to \mathbb{P} Y
\forall R : X \leftrightarrow Y \bullet
\operatorname{dom} R = \{x : X; y : Y \mid x\underline{R}y \bullet x\} \land
\operatorname{ran} R = \{x : X; y : Y \mid x\underline{R}y \bullet y\}
```

### 1.8.3 Numeric

succ [the next natural number, page 109]
.. [set of integers within a range, page 109]
# [number of members of a set, page 111]
min [smallest number in a set of numbers, page 113]
max [largest number in a set of numbers, page 113]

```
succ : \mathbb{N} \to \mathbb{N}
\underline{\quad \dots : \mathbb{Z} \times \mathbb{Z} \to \mathbb{P} \mathbb{Z}}
\forall n : \mathbb{N} \bullet succ(n) = n + 1
foralla, b : \mathbb{Z} \bullet
a ... b = \{ k : \mathbb{Z} \mid a \le k \le b \}
```

```
\begin{aligned} & \min: \mathbb{P}_1 \, \mathbb{Z} \to \mathbb{Z} \\ & \max: \mathbb{P}_1 \, \mathbb{Z} \to \mathbb{Z} \\ \\ & \min = \{ S: \mathbb{P}_1 \, \mathbb{Z}; m: \mathbb{Z} \, | \\ & m \in S \, \wedge \, (\forall n: S \bullet m \leq n) \bullet S \mapsto m \} \\ & \max = \{ S: \mathbb{P}_1 \, \mathbb{Z}; m: \mathbb{Z} \, | \\ & m \in S \, \wedge \, (\forall n: S \bullet m \geq n) \bullet S \mapsto m \} \end{aligned}
```

### 1.8.4 Sequences

```
[concatenation of two sequences, page 116]
                                             [reverse a sequence, page 116]
rev
                                     [first element of a sequence, page 117]
head
last
                                     [last element of a sequence, page 117]
                 [all elements of a sequence except for the first, page 117]
tail
front
                  [all elements of a sequence except for the last, page 117]
          [sub seq based on provided indices, order maintained, page 118]
       [sub seq based on provided condition, order maintained, page 118]
squash
             [compacts a fn of positive integers into a sequence, page 118]
^{\prime}
                              [flatten seq of seqs into single seq, page 121]
disjoint
                [pairs of sets in family have empty intersection, page 122]
partition
                      [union of all pairs of sets = the family set, page 122]
```

```
[X] = [X]
- ^ - : \operatorname{seq} X \times \operatorname{seq} X \to \operatorname{seq} X
rev : \operatorname{seq} X \to \operatorname{seq} X
\forall s, t : \operatorname{seq} X \bullet
s ^ t = s \cup \{ n : \operatorname{dom} t \bullet n + \#s \mapsto t(n) \}
\forall s : \operatorname{seq} X \bullet
revs = (\lambda n : \operatorname{dom} s \bullet s(\#s - n + 1))
```

### 1.8.5 Bags

```
[X] = count : bag X \rightarrow (X \rightarrow \mathbb{N})
-\#_{-} : bag X \times X \rightarrow \mathbb{N}
-\otimes_{-} : \mathbb{N} \times bag X \rightarrow bag X
\forall B : bag X \bullet
countB = (\lambda x : X \bullet 0) \oplus B
\forall x : X; B : bag x \bullet
B \# x = count B x
\forall n : \mathbb{N}; B : bag X; x : X \bullet
(n \otimes B) \# x = n * (B \# x)
```

```
[X] = \underbrace{items : \operatorname{seq} X \to \operatorname{bag} X}
\forall s : \operatorname{seq} X; x : X \bullet 
(items s) \# x = \# \{ i : \operatorname{dom} s \, | \, s(i) = x \}
```