

1 Rate of Completions

As learners engage in activities supported by a learning ecosystem, they will build up a history of learning experiences. When the digital resources of that learning ecosystem adhere to a framework dedicated to supporting and understanding the learner, such as the Total Learning Architecture (TLA), it becomes possible to retell their learning story through data and data visualization. One important aspect of that story is the rate of completion of the various digital resources within the learning ecosystem.

1.1 Initialization

$init(state)$ sets up an empty KV within $state$ for the Algorithm to update at each $step$

$$init(state) = state_0$$

where

$$state_0 = associate(state, < state, completions >, <>) \iff atKey(state, < state, completions >) = nil$$

otherwise

$$state_0 = state$$

such that if

$$state = < a \mapsto b >$$

then

$$state_0 = < a \mapsto b, state \mapsto completions \mapsto <> >$$

1.2 Relevant?

$relevant?(state, stmt)$ determines if $stmt$ is valid for use within $step$ of $rateOfCompletions$ and does so by looking into various $k \rightarrow v$ within $stmt$. The following Primitives are used as the *body* of $relevant?(state, stmt)$

- is the Object of the Statement an Activity?

$$activityType = atKey(stmt, < object, objectType >)$$

$$activity?(activityType) = true \iff activityType = Activity \vee activityType = nil$$

- is the Verb indicative of a completion event?

$$verbId = atKey(stmt, < verb, id >)$$

$$completionVerb?(verbId) = true$$

$$\iff$$

$$verbId = http : //adlnet.gov/expapi/verbs/passed$$

$$\begin{aligned}
& \vee \\
& verbId = https : //w3id.org/xapi/dod - isd/verbs/answered \\
& \vee \\
& verbId = http : //adlnet.gov/expapi/verbs/completed
\end{aligned}$$

- does the *stmt* indicate completion using Result?

$$\begin{aligned}
result &= atKey(stmt, < result, completion >) \\
resultCompletion &= true \iff result = true
\end{aligned}$$

such that the body of *relevant?* contains

$$p_a(stmt) = activity?(atKey(stmt, < object, objectType >))$$

and

$$p_v(stmt) = completionVerb?(atKey(stmt, < verb, id >))$$

and

$$p_r(stmt) = resultCompletion(atKey(stmt, < result, completion >))$$

which are used to form higher level Primitives

$$p_{continue}(stmt) = stmt \iff p_a(stmt) = true$$

and

$$p_{completed?}(stmt) = stmt \iff p_v(stmt) = true \vee p_r(stmt) = true$$

which results in a final Primitive $p_{return?}$

$$p_{return?}(stmt) = object?(p_{completed?}(p_{continue}(stmt)))$$

which defines the *body* of *relevant?*

$$relevant?(stmt) = p_{return?}(stmt) \Rightarrow object?(p_{completed?}(p_{continue}(stmt)))$$

and can be summarized as

$$\begin{aligned}
& relevant?(state, stmt) = true \\
& \iff \\
& activity?(activityType) = true \\
& \wedge \\
& completionVerb?(verbId) = true \vee resultCompletion = true
\end{aligned}$$

1.3 Accept?

rateOfCompletions does not require further boolean logic to determine if *stmt* and *state* can be passed to *step*

$$accept?(state, stmt) = object?(stmt)$$

which should always return true assuming valid xAPI Statements are passed to *rateOfCompletions*

1.4 Step

1.4.1 summary

step(state, stmt) updates *state* to include

$$$.object.id \mapsto \langle domain, statementCount, name \rangle$$

where

$$domain \mapsto \langle start, end \rangle$$

$$statementCount \mapsto \mathbb{R}$$

$$name \mapsto \langle $.object.definition.name \rangle$$

at

$$\langle state, completions, $.object.id \rangle$$

1.4.2 processing

step starts by extracting the relevant information from *stmt*

- *currentTime*

$$currentTime = atKey(stmt, timestamp)$$

- *name*

$$name_{stmt} = atKey(stmt, \langle object, definition, name \rangle)$$

- *objectId*

$$objectId = atKey(stmt, \langle object, id \rangle)$$

which allows for the previous *state* to be resolved using *objectId*

- *domain*

$$domain_{state} = atKey(state, \langle state, completions, objectId, domain \rangle)$$

$$start_{state} = first(domain_{state})$$

$$end_{state} = last(domain_{state})$$

- *statementCount*

$$statementCount_{state} = atKey(state, < state, completions, objectId, statementCount >)$$

- *name*

$$name_{state} = atKey(state, < state, completions, objectId, name >)$$

so that the previous state can be used along side the information parsed from *stmt*

- does $start_{state}$ need to be updated to *currentTime*?

where

$$inSeconds_{stmt} = isoToUnixEpoch(currentTime)$$

$$inSeconds_{start} = isoToUnixEpoch(start_{state}) \iff start_{state} \neq nil$$

such that

$$start(state, stmt) = currentTime$$

$$\iff$$

$$start_{state} = nil$$

$$\vee$$

$$inSeconds_{stmt} \leq inSeconds_{start}$$

otherwise

$$start(state, stmt) = start_{state}$$

- does end_{state} need to be updated to *currentTime*?

where

$$inSeconds_{stmt} = isoToUnixEpoch(currentTime)$$

$$inSeconds_{end} = isoToUnixEpoch(end_{state}) \iff end_{state} \neq nil$$

such that

$$end(state, stmt) = currentTime$$

$$\iff$$

$$end_{state} = nil$$

$$\vee$$

$$inSeconds_{stmt} \geq inSeconds_{end}$$

otherwise

$$end(state, stmt) = end_{state}$$

- what should *statementCount* be?

$$nStmts(state) = 1 \iff statementCount_{state} = 0 \vee nil$$

\vee

$$nStmts(state) = 1 + statementCount_{state} \iff statementCount_{state} \geq 1$$

- do we need to add a new *name*?

$$allNames(state, stmt) = append(name_{state}, name_{stmt}, count(name_{state}))$$

\iff

$$name_{stmt} \notin name_{state}$$

otherwise

$$allNames(state, stmt) = name_{state}$$

which allows for the following primitives to be defined

$$p_{start}(state, stmt) = start(state, stmt)$$

$$p_{end}(state, stmt) = end(state, stmt)$$

$$p_{stmtCount}(state, stmt) = nStmts(state)$$

$$p_{names}(state, stmt) = allNames(state, stmt)$$

and establish relevant paths into *state*

$$K_{domain} = \langle state, completions, objectId, domain \rangle$$

$$K_{stmtCount} = \langle state, completions, objectId, statementCount \rangle$$

$$K_{names} = \langle state, completions, objectId, name \rangle$$

which are used within higher level primitives concerned with updating *state*

$$p_{updateStart}(state, stmt)$$

\equiv

$$associate(state, K_{domain}, append(remove(domain_{state}, 0), p_{start}(state, stmt), 0))$$

and

$$p_{updateEnd}(state, stmt)$$

\equiv

$$associate(state, K_{domain}, append(remove(domain_{state}, 1), p_{end}(state, stmt), 1))$$

and

$$p_{updatedCount}(state, stmt)$$

\equiv

$$associate(state, K_{stmtCount}, p_{stmtCount}(state, stmt))$$

and

$$\begin{aligned} & p_{updatedNames}(state, stmt) \\ & \equiv \\ & associate(state, K_{names}, p_{names}(state, stmt)) \end{aligned}$$

such that *body* of *step* is defined as

$$step(state, stmt) = p_{updateNames}(p_{updateCount}(p_{updateEnd}(p_{updateStart}(state, stmt), stmt), stmt), stmt)$$

where

$$state' = p_{updateStart}(state, stmt)$$

and

$$state'' = p_{updateEnd}(state', stmt)$$

and

$$state''' = p_{updateCount}(state'', stmt)$$

such that

$$step(state, stmt) = p_{updateNames}(state''', stmt)$$

1.5 Result

The only *opts* used by *rateOfCompletions* is *timeUnit*

$$timeUnit = second \vee minute \vee hour \vee day \vee month \vee year$$

and will default to *day* if not passed to *rateOfCompletions*

$$result(state) = result(state, < timeUnit \mapsto day >)$$

which is passed to *rateOf* along with the arguments parsed from *state*

$$unit = atKey(opts, timeUnit)$$

$$allCompletions(state) = atKey(state, < state, completions >)$$

such that

$$\forall k_n : i..n..j \in allCompletions(state)$$

the following primitives are called each iteration

$$getCount(state, k_n) = atKey(allCompletions(state), < k_n, statementCount >)$$

$$getStart(state, k_n) = atKey(allCompletions(state), < k_n, domain, start >)$$

$$getEnd(state, k_n) = atKey(allCompletions(state), < k_n, domain, end >)$$

$$getName(state, k_n) = atKey(allCompletions(state), < k_n, name >)$$

which allows for

$$rate_n(state, k_n, unit) = rateOf(getCount(state, k_n), getStart(state, k_n), getEnd(state, k_n), unit)$$

such that

$$value_n(state, k_n, unit) = \langle x_n, y_n \rangle$$

where

$$name_n(state, k_n) = first(getName(state, k_n))$$

$$x_n = x \mapsto name_n(state, k_n) \iff name_n(state, k_n) \neq nil$$

otherwise

$$x_n = x \mapsto k_n$$

and

$$y_n = y \mapsto rate_n(state, k_n, unit)$$

such that

$$value_n(state, k_n, unit) = \langle name_n(state, k_n), rate_n(state, k_n, unit) \rangle$$

and

$$value(state, unit) = \forall k_n : i..n..j \in allCompletions(state) \exists! value_n(state, k_n, unit) = \langle x_n, y_n \rangle$$

$$\Rightarrow$$

$$value(state, unit) = \langle value_i(state, k_i, unit)..value_n(state, k_n, unit)..value_j(state, k_j, unit) \rangle$$

which allows the body of *result* to be defined using

$$unit = atKey(opts, timeUnit)$$

$$K_{store} = \langle state, completions, values, unit \rangle$$

so that *result* returns an updated *state* with the rate of completions per *unit* located at K_{store}

$$result(state, opts) = associate(state, K_{store}, value(state, unit))$$