Timeline Of Learner Success

As learners engage in a blended eLearning ecosystem, they will build up a history of learning experiences. When that eLearning ecosystem adheres to a framework dedicated to supporting and understanding the learner, such as the Total Learning Architecture (TLA), it becomes possible to retell their story through data. One important aspect of that story is the learner's history of success.

1 Ideal Statements

In order to accurately portray a learner's timeline of success, there are a few requirements of the data produced by a Learning Record Provider (LRP). They are as follows:

- the learner must be uniquely and consistently identified across all LRPs
- learning activities which access a learner's understanding of material should report if the learner was successful or not
 - if the assessment is scored, the grade earned by the learner should be reported
 - if the assessment is scored, the minimum and maximum possible grade should be reported
- The learning activities must be uniquely and consistently identified across all LRPs
- The time at which a learner completed a learning activity must be recorded
 - The timestamp should contain an appropriate level of specificity.
 - ie. Year, Month, Day, Hour, Minute, Second, Timezone

1.1 statement parameters to utilize

The statement parameter locations here are written in JSONPath. This notation is also defined by the Z Schema JSONPath

- \$.timestamp
- \bullet \$.result.success
- \$.actor
- \bullet \$.verb.id

2 TLA Statement problems

The data collected at the TLA pilot run supports the following algorithm.

3 Algorithm

3.1 Summary

- 1. Query an LRS via a GET request to the statements endpoint using the parameters agent, since and until
- 2. Filter the results to the set of statements where:
 - \$.verb.id is one of:
 - http://adlnet.gov/expapi/verbs/passed
 - https://w3id.org/xapi/dod-isd/verbs/answered
 - http://adlnet.gov/expapi/verbs/completed
 - \$.result.success is true

3.2 Query an LRS via REST

How to query an LRS via a GET request to the Statements Resource $^{\rm 1~2}$

 $^{^1}$ S is the set of all statements parsed from the statements array within the HTTP response to the Curl request. It may be possible that multiple Curl requests are needed to retrieve all query results. If multiple requests are necessary, S is the result of concatenating the result of each request into a single set

 $^{^2}$ Querying an LRS will not be defined within the following Z specifications but the Results of the Query will be

3.3 xAPI Z Specifications

An xAPI statement(s) is only defined abstractly within the context of Z. A concrete definition for an xAPI statement(s) it outside the scope of this specification.

3.3.1 Basic Types

 $IFI ::= mbox \mid mbox_sha1sum \mid openid \mid account$

• Type unique to Agents and Groups, The concrete definition of the listed values is outside the scope of this specification

 $OBJECTTYPE ::= Agent \mid Group \mid SubStatement \mid StatementRef \mid Activity$

• A type which can be present in all activities as defined by the xAPI specification

 $INTERACTIONTYPE ::= true-false \mid choice \mid fill-in \mid long-fill-in \mid matching \mid performance \mid sequencing \mid likert \mid numeric \mid other$

• A type which represents the possible interaction Types as defined within the xAPI specification

INTERACTIONCOMPONENT ::= choices | scale | source | target | steps

- A type which represents the possible interaction components as defined within the xAPI specification
- the concrete definition of the listed values is outside the scope of this specification

 $CONTEXTTYPES ::= parent \mid grouping \mid category \mid other$

• A type which represents the possible context types as defined within the xAPI specification

[STATEMENT]

• Basic types for the results of querying an LRS

[AGENT, GROUP]

• Basic types for Agents and collections of Agents

3.3.2 Id Schema

• the schema *Id* introduces the component *id* which is a non-empty finite set of 1 value

3.3.3 Schemas for Agents and Groups

```
\begin{array}{c} Agent \\ agent : AGENT \\ objectType : OBJECTTYPE \\ name : \mathbb{F}_1 \ \# 1 \\ ifi : IFI \\ \\ objectType = Agent \\ agent = \{ifi\} \cup \mathbb{P}\{name, objectType\} \end{array}
```

• The schema *Agent* introduces the component *agent* which is a set consisting of an *ifi* and optionally an *objectType* and/or *name*

```
Member \_\_
Agent
member : \mathbb{F}_1
member = \{a : AGENT \mid \forall a : a_0...a_n \bullet a = agent\}
```

• The schema Member introduces the component member which is a set of objects a, where for every a within $a_0...a_n$, a is an agent

```
Group \\ Member \\ group: GROUP \\ objectType: OBJECTTYPE \\ ifi: IFI \\ \hline name: \mathbb{F}_1 \# 1 \\ objectType = Group \\ group = \{objectType, name, member\} \lor \{objectType, member\} \lor \\ \{objectType, ifi\} \cup \mathbb{P}\{name, member\}
```

• The schema *Group* introduces the component *group* which is of type *GROUP* and is a set of either *objectType* and *member* with optionaly *name* or *objectType* and *ifi* with optionally *name* and/or *member*

```
Actor \_\_\_\_
Agent
Group
actor : AGENT \lor GROUP
actor = agent \lor group
```

• The schema *Actor* introduces the component *actor* which is either an *agent* or *group*

3.3.4 Verb Schema

```
 \begin{array}{c} -Verb \\ -Id \\ display, verb : \mathbb{F}_1 \\ \hline verb = \{id, display\} \vee \{id\} \end{array}
```

• The schema *Verb* introduces the component *verb* which is a set that consists of either *id* and the finite set *display* or just *id*

3.3.5 Object Schema

```
Extensions = \\ extensions, extensionVal : \mathbb{F}_1 \\ extensionId : \mathbb{F}_1 \#1 \\ \\ extensionS = \{e: (extensionId, extensionVal) \mid \forall i,j: e_i..e_j \bullet \\ (extensionId_i, extensionVal_i) \lor (extensionId_i, extensionVal_j) \land \\ (extensionId_j, extensionVal_i) \lor (extensionId_j, extensionVal_j) \land \\ extensionId_i \neq extensionId_j \}
```

- The schema Extensions introduces the component extensions which is a non-empty finite set that consists of ordered pairs of extensionId and extensionVal. Different extensionIds can have the same extensionVal but there can not be two identical extensionId values
- \bullet extension Id is a non-empty finite set with one value
- \bullet extensionVal is a non-empty finite set

```
\label{eq:contraction} \begin{split} & InteractionActivity \underline{\hspace{1cm}} \\ & interactionType : INTERACTIONTYPE \\ & correctResponsePattern : seq_1 \\ & interactionComponent : INTERACTIONCOMPONENT \\ \hline & interactionActivity = \{interactionType, correctReponsePattern, interactionComponent\} \lor \\ & \{interactionType, correctResponsePattern\} \end{split}
```

• The schema InteractionActivity introduces the component interactionActivity which is a set of either interactionType and correctResponsePattern or interactionType and correctResponsePattern and interactionComponent

```
\begin{tabular}{l} Log & Log
```

• The schema *Definition* introduces the component *definition* which is the non-empty, finite power set of *name*, *description*, *type*, *moreInfo* and *extensions*

```
Object ___
Id
Definition \\
Agent
Group
Statement
objectTypeA, objectTypeS, objectTypeSub, objectType:OBJECTTYPE
substatement: STATEMENT \\
object: \mathbb{F}_1
substatement = statement \\
objectTypeA = Activity
objectTypeS = StatementRef
objectTypeSub = SubStatement
objectType = objectTypeA \lor objectTypeS
object = \{id\} \lor \{id, objectType\} \lor \{id, objectTypeA, definition\}
         \vee \{id, definition\} \vee \{agent\} \vee \{group\} \vee \{objectTypeSub, substatement\}
         \vee \{id, objectTypeA\}
```

- The schema *Object* introduces the component *object* which is a nonempty finite set of either *id*, *id* and *objectType*, *id* and *objectTypeA* and *definition*, *agent*, *group*, or *substatement*
- The schema *Statement* and the corresponding component *statement* will be defined later on in this specification

3.3.6 Result Schema

```
Score = Score : \mathbb{F}_1
scaled, min, max, raw : \mathbb{Z}
scaled = \{n : \mathbb{Z} \mid -1.0 \le n \le 1.0\}
min = n < max
max = n > min
raw = raw = \{n : \mathbb{Z} \mid min \le n \le max\}
score = \mathbb{P}_1 \{scaled, raw, min, max\}
```

• The schema *Score* introduces the component *score* which is the non-empty powerset of *min*, *max*, *raw* and *scaled*

• The schema Result introduces the component result which is the nonempty power set of score, success, completion, response, duration and extensions

3.3.7 Context Schema

```
Instructor \\ Agent \\ Group \\ instructor: AGENT \lor GROUP \\ \\ instructor = agent \lor group
```

 \bullet The schema Instructor introduces the component instructor which can be ether an agent or a group

```
Team = Group
team : GROUP
team = group
```

• The schema Team introduces the component team which is a group

```
Context_{-}
Instructor
Team
Object
Extensions
registration, revision, platform, language: \mathbb{F}_1 \# 1
parentT, groupingT, categoryT, otherT: CONTEXTTYPES
contextActivities, statement: \mathbb{F}_1
statement = object \setminus (id, objectType, agent, group, definition)
parentT = parent
groupingT = grouping
categoryT = category
otherT = other
contextActivity = \{ca: object \setminus (agent, group, objectType, objectTypeSub, substatement)\}
contextActivityParent = (parentT, contextActivity)
contextActivityCategory = (categoryT, contextActivity)
contextActivityGrouping = (groupingT, contextActivity)
contextActivityOther = (otherT, contextActivity)
contextActivities = \mathbb{P}_1 \{ contextActivityParent, contextActivityCategory,
                        contextActivityGrouping, contextActivityOther\}
context = \mathbb{P}_1\{registration, instructor, team, contextActivities, revision, \}
              platform, language, statement, extensions \}
```

• The schema Context introduces the component context which is the nonempty powerset of registration, instructor, team, contextActivities, revision, platform, language, statement and extensions

3.3.8 Timestamp and Stored Schema

```
Timestamp \\ timestamp : \mathbb{F}_1 \# 1 Stored \\ stored : \mathbb{F}_1 \# 1
```

• The schema *Timestamp* and *stored* introduce the components *timestamp* and *stored* respectively. Each are non-empty finite sets containing one value

3.3.9 Attachements Schema

3.3.10 Statement and Statements Schema

```
Statement \\ Id \\ Actor \\ Verb \\ Object \\ Result \\ Context \\ Timestamp \\ Stored \\ Attachements \\ statement, \$: STATEMENT \\ \\ statement = \{actor, verb, object, stored\} \cup \\ \mathbb{P}\{\mathrm{id}, result, context, timestamp, attachments\} \\ \$ \sim statement
```

- The schema *Statement* introduces the component *statement* which consists of the components *actor*, *verb*, *object* and *stored* and the optional components *id*, *result*, *context*, *timestamp*, and/or *attachments*
- The schema Statement also binds the component statement to the variable \$ so that JSONPath can be used within Operation schemas which require reaching into a statement. This is accomplished by using the . (select) notation starting at \$ (root) and navigating into subsequent components of the statement

```
Statements \_
Statement
statements : \mathbb{F}_1
statements = \{s : statement\}
```

• The schema *Statements* introduces the component *statements* which is a non-empty finite set of components *statement*

3.4 Timeline Learner Success Z Specifications

The following Z Schemas define the system state, initialization and operations necessary to perform the Timeline Learner Success Algorithm

3.4.1 Timeline Leaner Success System State

```
Timeline Learner Success \\ S_{extra}, S_{completion}, S_{success}, S_{failure} : \mathbb{P} S \\ \\ S_{extra} \cup S_{completion} = S \\ S_{extra} \cap S_{completion} = \{\} \\ S_{success} \cup S_{failure} = S_{completion} \\ S_{success} \cap S_{failure} = \{\} \\
```

- The sets $S_{extra}, S_{completion}, S_{success}, S_{failure}$ are the powerset of S
- The union of sets S_{extra} and $S_{completion}$ is equal to the complete set of statements S
- No values are shared between the sets S_{extra} and $S_{completion}$
- The union of sets $S_{success}$ and $S_{failure}$ is equal to the set $S_{completion}$
- No values are shared between the sets $S_{success}$ and $S_{failure}$

3.4.2 Initial State of Timeline Learner Success System

```
InitTimelineLearnerSuccess TimelineLearnerSuccess S_{extra} = \{\} S_{completion} = \{\} S_{success} = \{\} S_{failure} = \{\}
```

• The sets S_{extra} , $S_{completion}$, $S_{success}$, $S_{failure}$ are all initially empty

3.4.3 Filter for Completion

 \bullet The var $V_{completion}$ has a value of one of the above IRIs and is of type Verb

- The updated set $S'_{completion}$ is the set of all statements s where $V_{completion}$ is in s and s is in S
- the updated set S'_{extra} is the set of all statements s where $V_{completion}$ is not in s and s is in s

3.4.4 Filter for Success

```
ResultSuccessTrue
R_{successful}: Result
R_{successful} = true
R_{successful} \neq false
```

• The var $R_{successful}$ has a value of true but not false and is of the type Result

```
Filter For Success \\ \Delta Time line Learner Success \\ s_{completion}: Statement \\ \\ s_{completion} \in S_{completion} \\ S'_{success} = \{s_{completion}: Statement \mid R_{successful} \in s_{completion}\} \\ S'_{failure} = \{s_{completion}: Statement \mid R \not\in s_{completion}\} \\
```

- The set $s_{completion}$ is of type Statement and is in the set $S_{completion}$
- The updated set $S'_{success}$ is the set of all statements $s_{completion}$ where $R_{successful}$ is in $s_{completion}$
- The updated set $S'_{failure}$ is the set of all statements $s_{completion}$ where $R_{successful}$ is not in $s_{completion}$

3.4.5 Return

```
Return
\Xi Timeline Learner Success
S_{success}!: Statements
S_{success}!= S_{success}
```

 \bullet The returned variable $S_{success}!$ is equal to the current state of variable $S_{success}$

3.5 Pseudocode

Algorithm 1: Timeline of Learner Success

```
Input: S
Result: S_{success}
while S is not empty do
     for each Statement s in S
     if s.verb.id = V_{completion} then
      | add s to S_{completion}
     else
      | add s to S_{extra}
     \quad \text{end} \quad
\quad \mathbf{end} \quad
 \begin{array}{l} \textbf{while} \ S_{completion} \ is \ not \ empty \ \textbf{do} \\ | \ \text{for each Statement} \ s_{completion} \ \text{in} \ S_{completion} \end{array} 
     if s_{completion}.result.success = R_{success} then
          add s_{completion} to S_{success}
     else
      | add s_{completion} to S_{failure}
     end
end
```

3.6 Result JSON Schema

JSON schema describing the returned data structure

3.7 Visualization Description

description of the associated visualization in english

3.8 Visualization prototype

This section will be updated to a prototype viz