

0.0.1 Associate At

The Primitive *associateAt* establishes a relationship between $k?_j$ and $v?$ at the nesting $k?_i \dots k?_{j-1}$ within $m!$

$$\frac{\begin{array}{l} k? = \langle k?_i \dots k?_j \rangle \\ (k?_j, m?_{k?_j}) \in m? \vee (k?_j, m?_{k?_j}) \notin m? \\ (k?_j, m?_{k?_j}) \notin m! \\ (k?_j, v?) \in m! \end{array}}{m! = \text{associateAt}(m?, k?, v?)}$$

This implies that any existing mapping at $k?_j \in m?$ will be overwritten by *associateAt* but an existing mapping is not a precondition. The following helper Operation *getFirstKey* is introduced to help with the recursive nature of *associateAt*.

$$\frac{\begin{array}{l} \text{GetFirstKey}[KV, \text{Collection}] \\ m? : KV \\ k? : \text{Collection} \\ v! : V \\ \text{getFirstKey}_- : KV \times \text{Collection} \twoheadrightarrow V \end{array}}{v! = \text{getFirstKey}(m?, k?) \bullet v! = \text{atKey}(m?, \text{head}(k?))}$$

This allows for the recursive aspect of *associateAt* to be defined as $\langle \text{getFirstKey}_-, \text{recur}_- \rangle^{\#k?-1}$

$$\frac{\begin{array}{l} \text{AssociateAt}[KV, \text{Collection}, V] \\ \text{GetFirstKey}, \text{Recur} \\ m?, m! : KV \\ k? : \text{Collection} \bullet \forall k?_n \in k?_{\langle i \dots j \rangle} \mid k_n : K \\ v? : V \\ \text{associateAt}_- : KV \times \text{Collection} \times V \twoheadrightarrow KV \end{array}}{\begin{array}{l} \text{associateAt} = \langle \langle \text{atFirstKey} \text{ recur}_- \rangle^{\#k?-1}, \text{associate}_- \rangle \\ m! = \text{associateAt}(m?, k?, v?) \bullet \\ \quad \forall n : i \dots j - 1 \in k? \bullet j = \text{first}(\text{last}(k?)) \Rightarrow \text{first}(j, k?_j) \mid \exists_1 v_n \bullet \\ \quad \text{let } c_n == \text{tail}(k?)^{n-i} \\ \quad \quad v_i == \text{getFirstKey}(m?, c_n) \Rightarrow \\ \quad \quad \quad c_n = k? \iff n = i \bullet v_i = \text{atKey}(m?, \text{head}(k?)) \\ \quad \quad v_n = \text{recur}(v_i, c_n, \text{getFirstKey}_-)^{j-1} \\ \quad \quad v_{j-1} == \text{getFirstKey}(v_n, (j-1 \upharpoonright k?)) \iff n = j-2 \\ \quad \quad v_j == \text{associate}(v_{j-1}, \text{last}(k?), v?) \Rightarrow \langle \langle k?_j \mapsto v? \rangle \rangle \cup v_{j-1} \triangleleft k?_j \\ = (v_j \cup v_{n-\text{succ}(1)} \triangleleft k?_{n-\text{succ}(0)})^{j-1} \bullet n \leq j-1 \Rightarrow \\ \quad v_{j-4} \cup (v_{j-3} \cup (v_j \cup v_{j-2} \triangleleft k?_{j-1}) \triangleleft k?_{j-2}) \triangleleft k?_{j-3} \end{array}}$$

No other mappings which exist at the various levels of depth will be altered