

HW07

BY ZIHAN ZHOU

1 WILD

Case $r \cong \text{Wild}$:

to show :

$\forall cs: \text{char list}, k: \text{char list} \rightarrow \text{bool} \{ \text{match wild } cs \hookrightarrow \text{true} \mid \exists p, s, p \circ s \cong cs \wedge p \in L(_) \wedge ks \hookrightarrow \text{true} \}$

```
match Wild cs k =  
  (case cs of  
    nil => false  
    | c'::cs' => k cs')
```

Since $\text{match Wild } cs \ k \cong \text{true}$

then cs is for sure not nil .

so cs must has form $c'::cs'$, and $k \ cs' \hookrightarrow \text{true}$

let $p = [c']$, $s=cs'$,

then $cs=p@s$, and $p \in L(_)$, and $k \ s \hookrightarrow \text{true}$

2 BOTH

Case $r \sim = \text{Both}(r1, r2)$: To show:

$\forall cs: \text{char list}, k: (\text{char list} \rightarrow \text{bool}) \{ \text{match}(\text{Both}(r1, r2)) \ cs' \ k \hookrightarrow \text{true} \mid \exists p, s: \text{char list}, p \circ s \cong cs' \wedge p \in L(r1 \cap r2) \wedge ks \hookrightarrow \text{true} \}$

two IHs:

IH1: $\forall cs_1: \text{char list}, k_1: (\text{char list} \rightarrow \text{bool}) \{ \text{match}(r1) \ cs_1 \ k_1 \hookrightarrow \text{true} \mid \exists p_1, s_1: \text{char list}, p_1 \circ s_1 \cong cs_1 \wedge p_1 \in L(r1) \wedge k_1 \ s_1 \hookrightarrow \text{true} \}$

IH2: $\forall cs_2: \text{char list}, k_2: (\text{char list} \rightarrow \text{bool}) \{ \text{match}(r2) \ cs_2 \ k_2 \hookrightarrow \text{true} \mid \exists p_2, s_2: \text{char list}, p_2 \circ s_2 \cong cs_2 \wedge p_2 \in L(r2) \wedge k_2 \ s_2 \hookrightarrow \text{true} \}$

```
| match (Both(r1,r2)) cs k =  
  match r1 cs (fn u1=>  
    match r2 cs (fn u2=>  
      u1=u2 andalso k(u1) andalso k(u2)))
```

we know

1. $\text{match}(\text{Both}(r1,r2)) \ cs \ k \hookrightarrow \text{true}$

so

```
match r1 cs (fn u1=>  
  match r2 cs (fn u2=>  
    u1=u2 andalso k(u1) andalso k(u2))) \hookrightarrow \text{true}
```

by IH1:

let $cs_1 = cs$,
 let $k_1 = (\text{fn } u1 = >$
 match $r2$ cs ($\text{fn } u2 = >$
 $u1 = u2$ andalso $k(u1)$ andalso $k(u2)))$
 then we have
 2. $cs_1 = p_1 \circ u_1$
 3. $p_1 \in L(r1)$
 4. ($\text{fn } u1 = >$
 match $r2$ cs ($\text{fn } u2 = >$
 $u1 = u2$ andalso $k(u1)$ andalso $k(u2)))$ $u1 \hookrightarrow \text{true}$
 5. match $r2$ cs ($\text{fn } u2 = >$
 $u1 = u2$ andalso $k(u1)$ andalso $k(u2)) \hookrightarrow \text{true}$
 apply IH2 on 5.
 let $cs_2 = cs$
 let $k_2 = (\text{fn } u2 = >$
 $u1 = u2$ andalso $k(u1)$ andalso $k(u2))$
 then we have
 6. $cs_2 = p_2 \circ u_2$
 7. $p_2 \in L(r2)$
 8. ($\text{fn } u2 = >$
 $u1 = u2$ andalso $k(u1)$ andalso $k(u2))$ $u2 \hookrightarrow \text{true}$
 9. $u1 = u2$ andalso $k(u1)$ andalso $k(u2) \hookrightarrow \text{true}$
 by lemma 1
 because $u1 = u2$ and $p1@u1 = p2@u2 = cs$, so $p1 = p2$
 So we find
 $p = p1 = p2$
 $p \in L(r1) \cap L(r2) \Rightarrow p \in L(r1 \cap r2)$
 $u = u1 = u2$
 $cs' = cs = p@u$
 $k\ u = \text{true}$
 QED