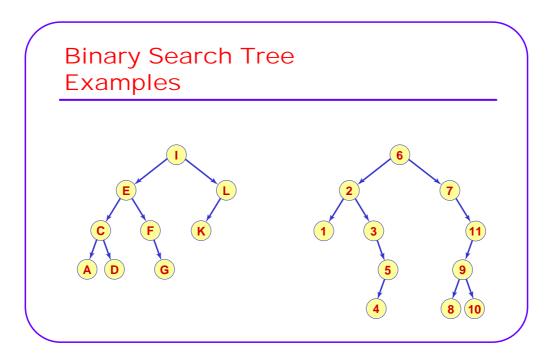
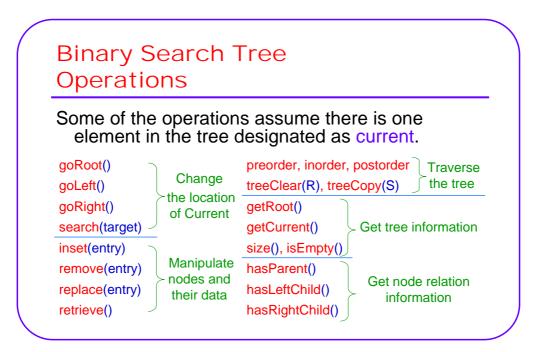
# Non-Linear Data Structures

**Binary Search Trees** 

# Binary Search Tree Definition

- A Binary Search Tree (BST) is a binary tree such that for each node, N, the following conditions are true:
  - 1. If L is any node in the left subtree of N, then the key of L is less than the key of N.
  - 2. If R is any node in the right subtree of N, then the key of R is greater than the key of N.
- The key values of the elements must be from an ordered set





# BST Implementation Using a Linked Structure

- Use the binary tree node generic class: BTNode<E>, as defined before.
- Use two BTNode<E> pointer variables:
  - root -- points to the root node of the tree.
  - current -- points to the current element node in the tree.
- Use one integer variable:
  - count -- stores the number of nodes in the tree.
- An empty BST is initialized by setting: root = current = null, and count = 0.
- The location of the current node is controlled by insert, remove, search, and the three go operations.



element

Right child

Left child

count

# BST Implementation as a Java Generic class Using a Linked Structure

# BST Implementation as a Java Generic class Using a Linked Structure

```
// MUTATOR METHODS:
public void goRoot() { if (count > 0) current = root; }
public void goLeft() { if (hasLeftChild()) current = current.getLeft(); }
public void goRight() { if (hasRightChild()) current = current.getRight();}
public void search(E target) { ... ... }
public void insert(E entry) { ... ... }
public boolean remove(E entry) { ... ... }
public void replace(E entry) { ... ... }
public void treeClear(BTNode<E> p) { ... ... }

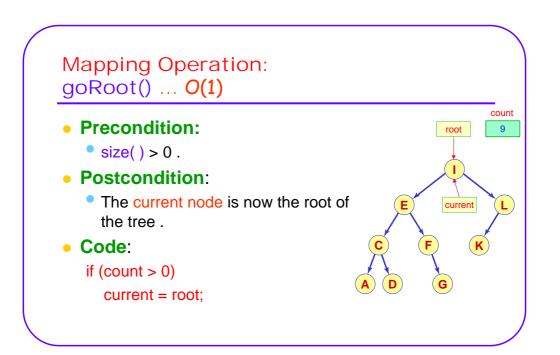
// PRIVATE HELPER FUNCTIONS:
private void createFirstNode(E entry) { ... ... }
private void addLeft(E entry) { ... ... }
private void addRight(E entry) { ... ... }
```

# BST Implementation as a Java Generic class Using a Linked Structure

```
// OBSERVER METHODS:
public E retrieve() {
    if (current != null)
        return current.getElement()
    else
        return null;
}

public boolean hasParent() { return current != root; }
public boolean hasLeftChild() { return current.getLeft() != null; }
public boolean hasRightChild() { return current.getRight() != null; }
public boolean isEmpty() { return (count == 0); }
public BTNode<E> getRoot() { return root; }
public BTNode<E> getCurrent() { return current; }
public int size() { return count; }
```

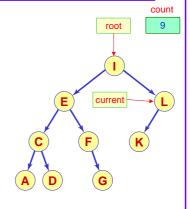
# Mapping Operation: goRoot() ... O(1) • Precondition: • size() > 0. • Postcondition: • The current node is now the root of the tree. • Code: if (count > 0) current = root;



# Mapping Operation: goLeft() ... O(1)

- Precondition:
  - hasLeftChild() returns true.
- Postcondition:
  - The current pointer has been shifted down to point to the left child of the original current node.
- Code:

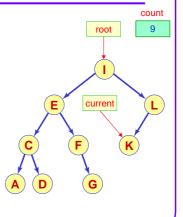
if (hasLeftChild())
 current = current.getLeft();



# Mapping Operation: goLeft() ... O(1)

- Precondition:
  - hasLeftChild() returns true.
- Postcondition:
  - The current pointer has been shifted down to point to the left child of the original current node.
- Code:

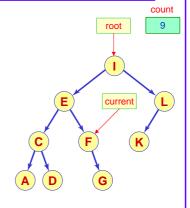
if (hasLeftChild())
 current = current.getLeft();



# Mapping Operation: goRight() ... O(1)

- Precondition:
  - hasRightChild() returns true.
- Postcondition:
  - The current pointer has been shifted down to point to the right child of the original current node.
- Code:

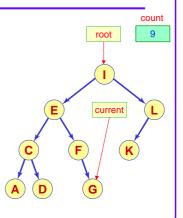
if (hasRightChild())
 current = current.getRight();



# Mapping Operation: goRight() ... O(1)

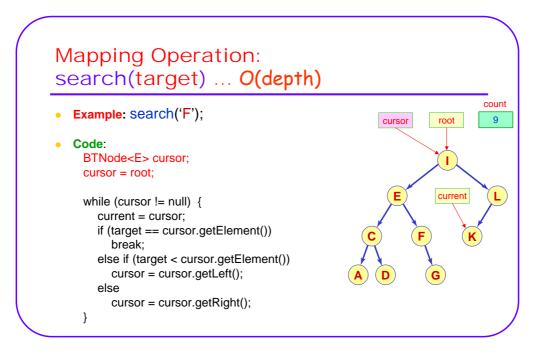
- Precondition:
  - hasRightChild() returns true.
- Postcondition:
  - The current pointer has been shifted down to point to the right child of the original current node.
- Code:

if (hasRightChild())
 current = current.getRight();

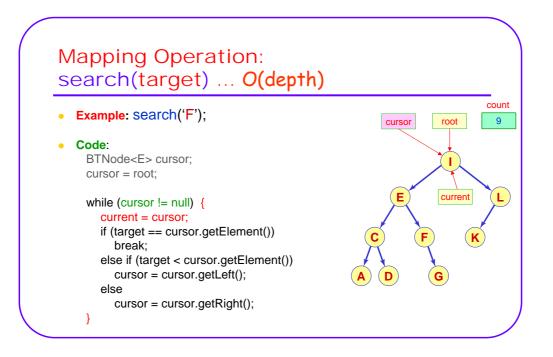


### Mapping Operation: search(target) ... O(depth) count **Precondition:** 9 root size() > 0.Postcondition: If the target element is in a node of the tree, then that node is made as the current node, otherwise, the current node is at the would be parent of the target Code: current BTNode<E> cursor; cursor = root; while (cursor != null) { current = cursor; if (target == cursor.getElement()) else if (target < cursor.getElement()) cursor = cursor.getLeft(); == and < are used here for brief, cursor = cursor.getRight(); Use compareTo() instead.

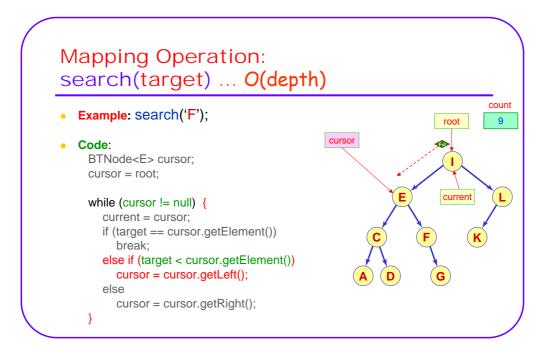
}



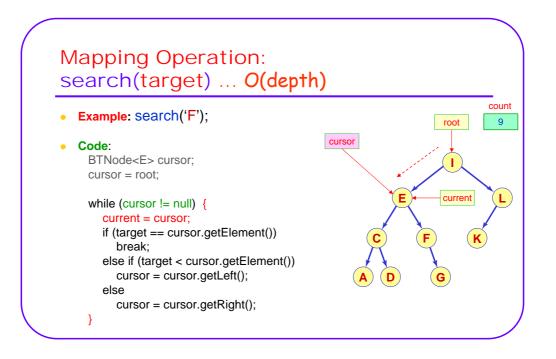
### Mapping Operation: search(target) ... O(depth) count Example: search('F'); 9 cursor root Code: BTNode<E> cursor; Ī cursor = root; current while (cursor != null) { current = cursor; if (target == cursor.getElement()) break; else if (target < cursor.getElement()) cursor = cursor.getLeft(); cursor = cursor.getRight();



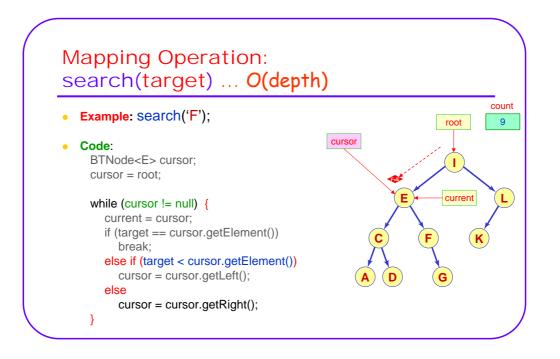
### Mapping Operation: search(target) ... O(depth) count Example: search('F'); 9 cursor root Code: BTNode<E> cursor; cursor = root; current while (cursor != null) { current = cursor; if (target == cursor.getElement()) break; else if (target < cursor.getElement())</pre> cursor = cursor.getLeft(); cursor = cursor.getRight();



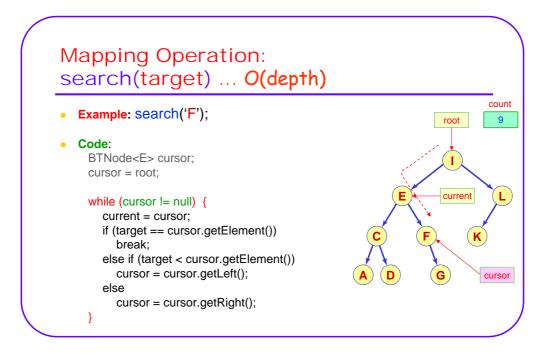
### Mapping Operation: search(target) ... O(depth) count Example: search('F'); 9 root cursor Code: BTNode<E> cursor; cursor = root; current while (cursor != null) { current = cursor; if (target == cursor.getElement()) break; else if (target < cursor.getElement()) cursor = cursor.getLeft(); cursor = cursor.getRight();



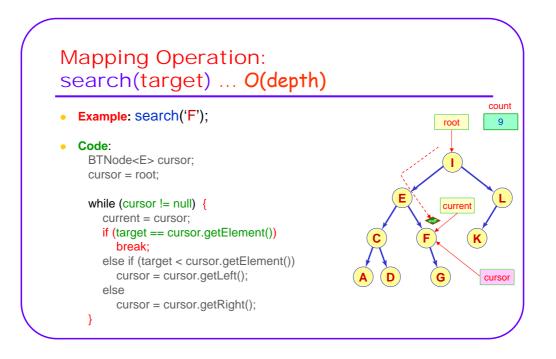
### Mapping Operation: search(target) ... O(depth) count Example: search('F'); 9 root cursor Code: BTNode<E> cursor; cursor = root; current while (cursor != null) { current = cursor; if (target == cursor.getElement()) break; else if (target < cursor.getElement())</pre> cursor = cursor.getLeft(); cursor = cursor.getRight();



```
Mapping Operation:
search(target) ... O(depth)
                                                                       count
  Example: search('F');
                                                                        9
                                                               root
  Code:
    BTNode<E> cursor;
    cursor = root;
                                                                current
    while (cursor != null) {
       current = cursor;
       if (target == cursor.getElement())
         break;
       else if (target < cursor.getElement())
         cursor = cursor.getLeft();
         cursor = cursor.getRight();
```

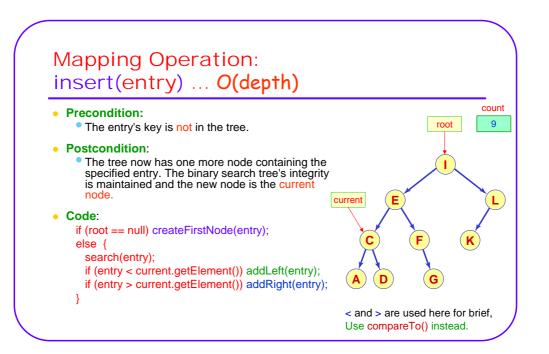


```
Mapping Operation:
search(target) ... O(depth)
                                                                       count
  Example: search('F');
                                                                        9
                                                               root
   Code:
    BTNode<E> cursor;
    cursor = root;
    while (cursor != null) {
                                                               current
       current = cursor;
       if (target == cursor.getElement())
       else if (target < cursor.getElement())
         cursor = cursor.getLeft();
         cursor = cursor.getRight();
```

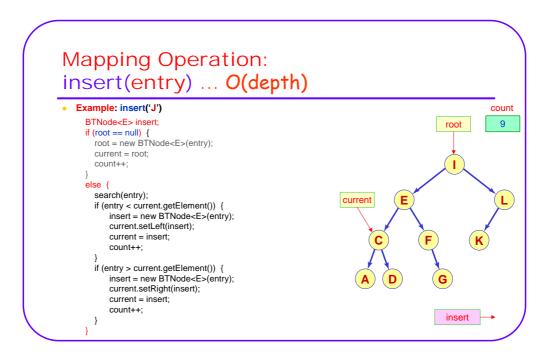


### Mapping Operation: search(target) ... O(depth) count Example: search('F'); 9 root Code: BTNode<E> cursor; cursor = root; while (cursor != null) { current current = cursor: if (target == cursor.getElement()) break; else if (target < cursor.getElement()) cursor = cursor.getLeft();

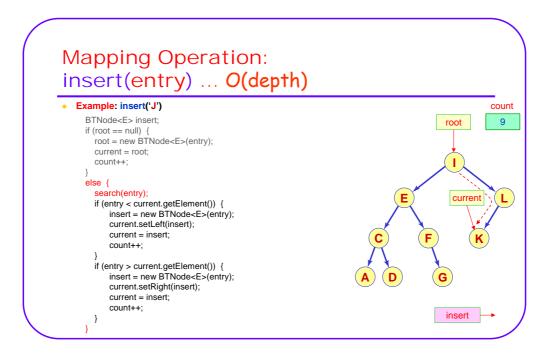
cursor = cursor.getRight();



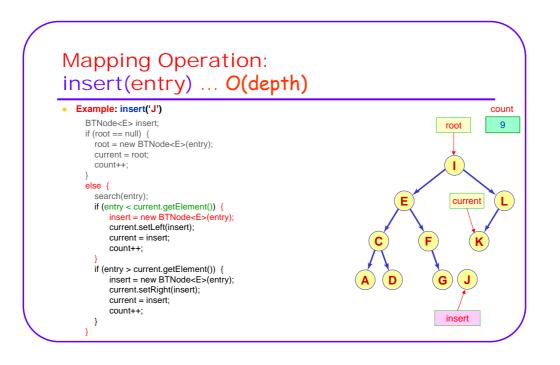
### Mapping Operation: insert(entry) ... O(depth) Code Expanded: count BTNode<E> insert; 9 root if (root == null) { root = new BTNode<E>(entry); current = root; count++; else { search(entry); if (entry < current.getElement()) { current insert = new BTNode<E>(entry); current.setLeft(insert); current = insert: count++; if (entry > current.getElement()) { insert = new BTNode<E>(entry); current.setRight(insert); current = insert; count++;



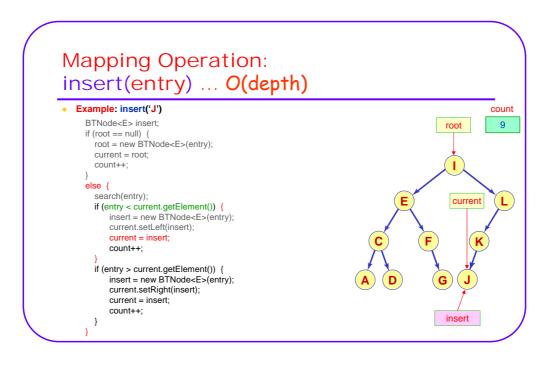
### Mapping Operation: insert(entry) ... O(depth) Example: insert('J') count BTNode<E> insert; 9 root if (root == null) { root = new BTNode<E>(entry); current = root; count++; else { search(entry); if (entry < current.getElement()) {</pre> current insert = new BTNode<E>(entry); current.setLeft(insert); current = insert: count++; if (entry > current.getElement()) { insert = new BTNode<E>(entry); current.setRight(insert); current = insert;



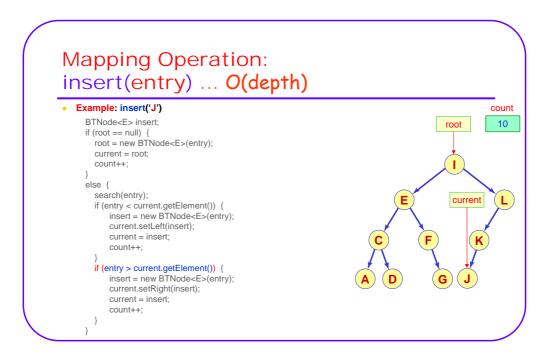
### Mapping Operation: insert(entry) ... O(depth) Example: insert('J') count BTNode<E> insert; 9 root if (root == null) { root = new BTNode<E>(entry); current = root; count++; else { search(entry); if (entry < current.getElement()) { current insert = new BTNode<E>(entry); current.setLeft(insert); current = insert: count++; if (entry > current.getElement()) { insert = new BTNode<E>(entry); current.setRight(insert); current = insert;



### Mapping Operation: insert(entry) ... O(depth) Example: insert('J') count BTNode<E> insert; 9 root if (root == null) { root = new BTNode<E>(entry); current = root; count++; else { search(entry); if (entry < current.getElement()) { current insert = new BTNode<E>(entry); current.setLeft(insert); current = insert: count++; if (entry > current.getElement()) { insert = new BTNode<E>(entry); current.setRight(insert); current = insert; insert



### Mapping Operation: insert(entry) ... O(depth) Example: insert('J') count BTNode<E> insert; 10 root if (root == null) { root = new BTNode<E>(entry); current = root; count++; else { search(entry); if (entry < current.getElement()) { current insert = new BTNode<E>(entry); current.setLeft(insert); current = insert: count++; if (entry > current.getElement()) { insert = new BTNode<E>(entry); current.setRight(insert); current = insert; count++; insert



# Mapping Operation: remove(entry) ... O(depth)

- Precondition:
  - size() > 0.
- Postcondition:
  - If entry is in the tree then it has been deleted from it. The tree now has one node less than before. The root of the tree is now the current node, the return value is true, and the binary search tree's integrity is maintained. If entry is not in the tree then nothing happens and the return value is false.

# Mapping Operation: remove(entry) ... O(depth)

Algorithm:

This is the most complex operation on a BST. Two cases exist:

Case1, If either subtree of the node to be deleted is empty:

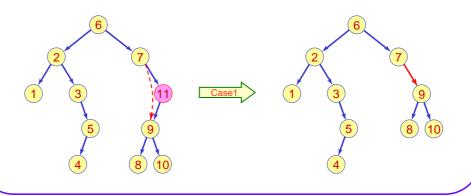
 Replace the pointer from its parent with the pointer to its non-empty subtree.

Case2, If both subtrees of the node, N, to be deleted are non-empty:

- Find the rightmost node, R, in the left subtree of N.
   This is the node with the maximum key value on the left subtree of N, and it is guaranteed to have at least one non-empty subtree.
- 2. Move the contents of R to N.
- 3. Replace the pointer from R's parent with the pointer to R's left subtree.

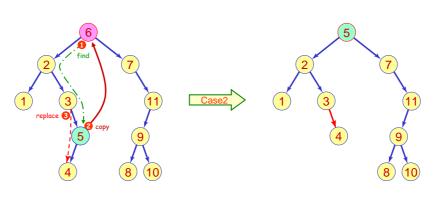
# Mapping Operation: remove(entry) ... O(depth)

• Example 1: remove node with key value = 11.



# Mapping Operation: remove(entry) ... O(depth)

• Example 2: remove node with key value = 6.



## Mapping Operation: remove(entry) ... O(depth)

Code:

```
// This code depends on the findParent(entry) helper function, which is used to give
// access to the parent of the node containing the given entry. It has the following
// postcondition: If the entry is found in a node of the tree, then current points at that
// node and its parent link is returned. If the found node is the root, then both current
// and the returned link point at the root node. If the entry is not found in the tree, then
// current is set at the would be parent of entry and the value null is returned.
BTNode<E> parent, temp;
if (root == null) return false;
                                        // no nodes to delete
temp = current;
parent = findParent(entry);
                                        // find the entry and its parent node
if (parent == null) {
                                        // the entry was not found in the tree
  current = temp;
                                        // restore current
  return false;
```

# Mapping Operation: remove (entry) ... O(depth)

```
// At this point, the node to be deleted is found in the tree. current is at the found node
// and parent is at its parent node. If entry is at the root then both pointers are at the root.
// Two cases exist: 1: The node has at least one null child.
                       2. The node has two non-null children.
if (current.getRight() == null) {
                                                    // Case1, with right null child.
  // Attach its left child to its parent.
  if (current == parent) root = current.getLeft();
  if (current == parent.getLeft()) parent.setLeft(current.getLeft());
  if (current == parent.getRight()) parent.setRight(current.getLeft());
  // Delete the node.
  if (current != null)
      { current.setLeft(null); current.setRight(null); current.setElement(null); }
  current = root;
  count--;
  return true;
```

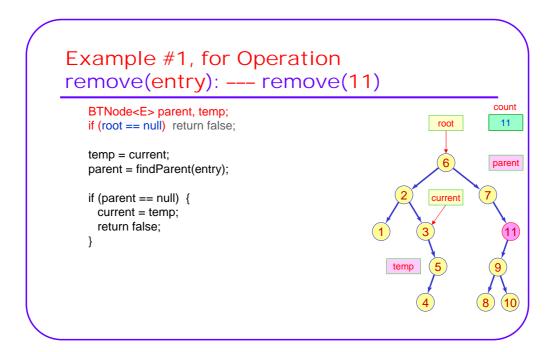
# Mapping Operation: remove(entry) ... O(depth)

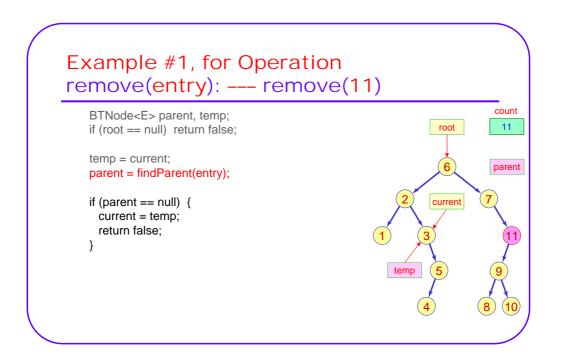
# Mapping Operation: remove(entry) ... O(depth)

```
// Case2: The node has two non-null children.
// Find the node with the maximum element value in the left subtree,
// This maximum node is guaranteed to have at least one null child.
parent = current; temp = current.getLeft();
while (temp.getRight() != null) { parent = temp; temp = temp.getRight(); }
// The maximum node is found. Copy its contents to subtree's root node.
current.setElement(temp.getElement());
// Attach its left child to its parent node.
if (temp == parent.getLeft()) parent.setLeft(temp.getLeft());
if (temp == parent.getRight()) parent.setRight(temp.getLeft());
// Delete the maximum node.
if (temp != null)
  { temp.setLeft(null); temp.setRight(null); temp.setElement(null); temp = null; }
current = root;
count--;
return true
```

# The helper function: findParent(target)

```
BTNode<E> parent = root;
int where = 0;
                                                             // used to trace current direction
current = root;
while (current != null) {
  if (c.compare(target, current.getElement()) == 0) return parent;
                                                                         // target is found.
  else if (c.compare(target, current.getElement()) < 0) {
                                                                         // check left subtree.
     current = current.getLeft();
                                                                         // move current left
                                                             // move parent following current
     if (where == -1) parent = parent.getLeft();
     if (where == +1) parent = parent.getRight();
     where = -1;
                                                             // record that current moved left
                                                                         // check right subtree.
  else {
     current = current.getRight();
                                                                         // move current right
     if (where == -1) parent = parent.getLeft();
                                                              // move parent following current
     if (where == +1) parent = parent.getRight();
     where = +1;
                                                             // record that current moved right
return null;
```



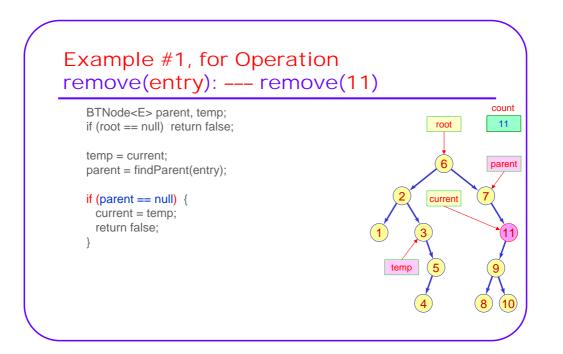


```
Example #1, for Operation remove(entry): --- remove(11)

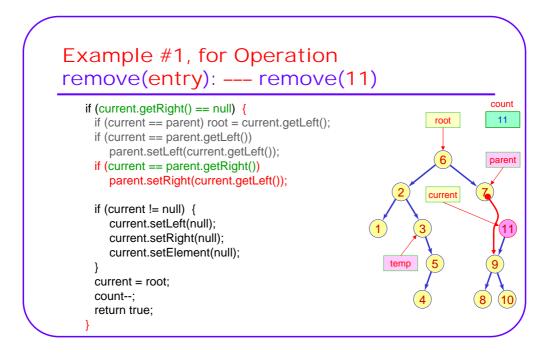
BTNode<E> parent, temp; if (root == null) return false; temp = current; parent = findParent(entry); if (parent == null) { current = temp; return false; }

if (parent == null) { 2 current = temp; return false; }

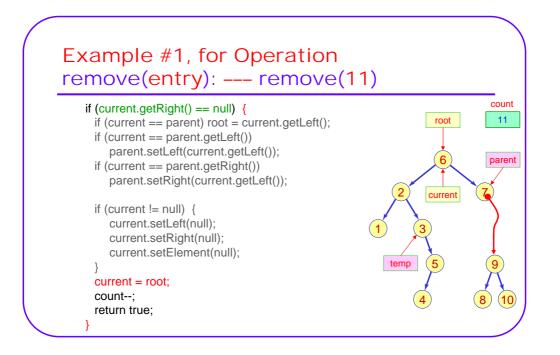
1 3 11
```



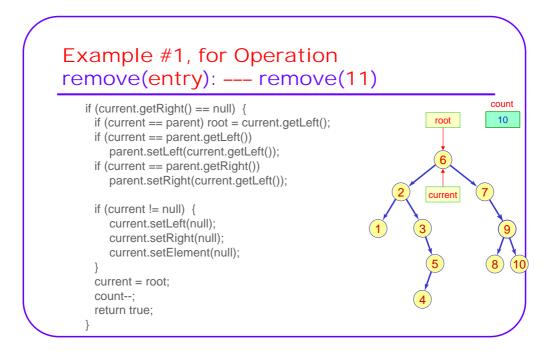
### Example #1, for Operation remove(entry): --- remove(11) count if (current.getRight() == null) { 11 root if (current == parent) root = current.getLeft(); if (current == parent.getLeft()) parent.setLeft(current.getLeft()); parent 6 if (current == parent.getRight()) parent.setRight(current.getLeft()); current if (current != null) { current.setLeft(null); current.setRight(null); current.setElement(null); temp current = root; count--; return true;



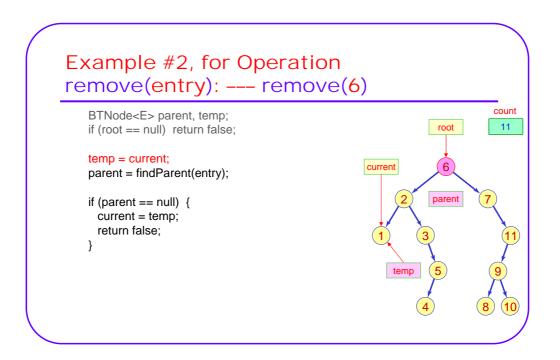
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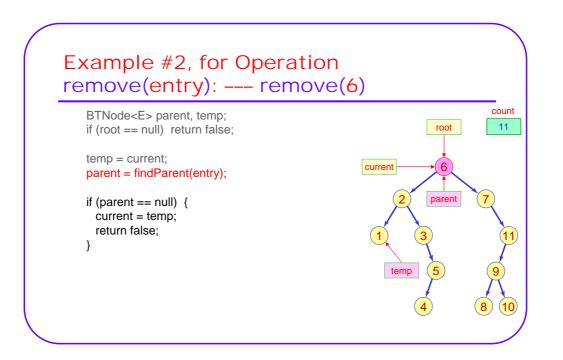
### Example #1, for Operation remove(entry): --- remove(11) count if (current.getRight() == null) { 10 root if (current == parent) root = current.getLeft(); if (current == parent.getLeft()) parent.setLeft(current.getLeft()); parent if (current == parent.getRight()) parent.setRight(current.getLeft()); current if (current != null) { current.setLeft(null); current.setRight(null); current.setElement(null); temp current = root; count--; return true;



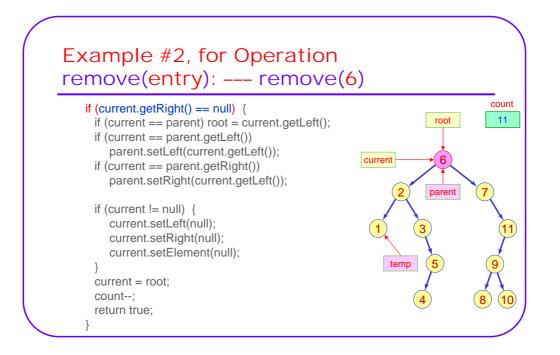
```
Example #2, for Operation
remove(entry): --- remove(6)
                                                                      count
   BTNode<E> parent, temp;
                                                                      11
                                                            root
   if (root == null) return false;
   temp = current;
                                                current
                                                             \left[ 6\right]
   parent = findParent(entry);
                                                           parent
   if (parent == null) {
     current = temp;
     return false;
                                                     temp
```



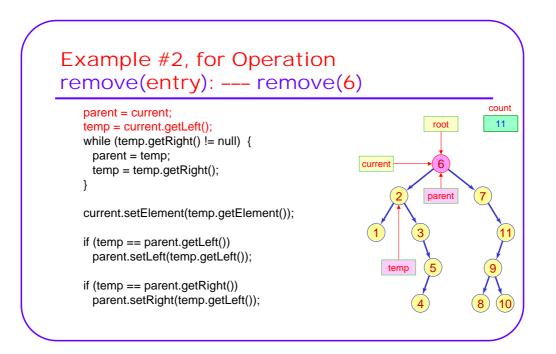
```
Example #2, for Operation
remove(entry): --- remove(6)
                                                                     count
   BTNode<E> parent, temp;
                                                                      11
                                                            root
   if (root == null) return false;
   temp = current;
                                                current
                                                             \left[ 6\right]
   parent = findParent(entry);
                                                           parent
   if (parent == null) {
     current = temp;
     return false;
                                                    temp
```



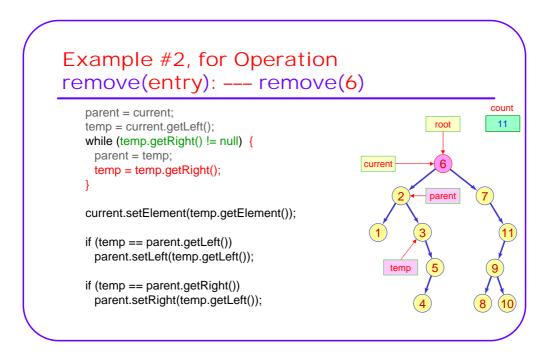
### Example #2, for Operation remove(entry): --- remove(6) count BTNode<E> parent, temp; 11 root if (root == null) return false; temp = current; current 6 parent = findParent(entry); if (parent == null) { parent current = temp; return false; temp



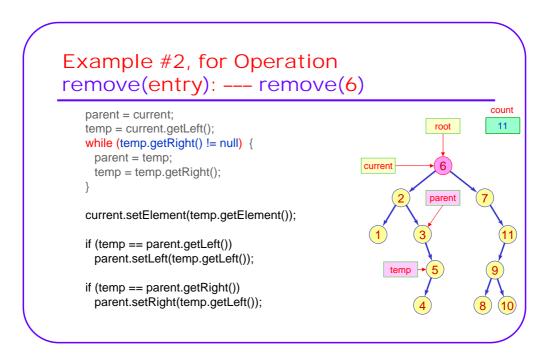
### Example #2, for Operation remove(entry): --- remove(6) count if (current.getLeft() == null) { 11 root if (current == parent) root = current.getRight(); if (current == parent.getLeft()) parent.setLeft(current.getRight()); current 6 if (current == parent.getRight()) parent.setRight(current.getRight()); parent if (current != null) { current.setLeft(null); current.setRight(null); current.setElement(null); temp current = root; count--; return true;



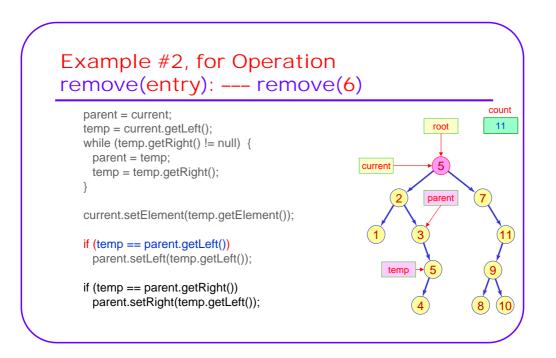
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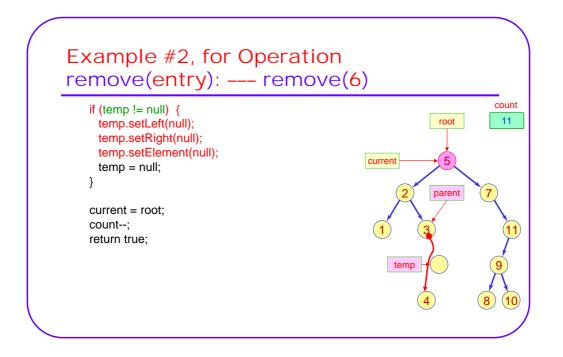
### Example #2, for Operation remove(entry): --- remove(6) count parent = current; 11 root temp = current.getLeft(); while (temp.getRight() != null) { parent = temp; current 6 temp = temp.getRight(); parent current.setElement(temp.getElement()); if (temp == parent.getLeft()) parent.setLeft(temp.getLeft()); temp if (temp == parent.getRight()) parent.setRight(temp.getLeft());



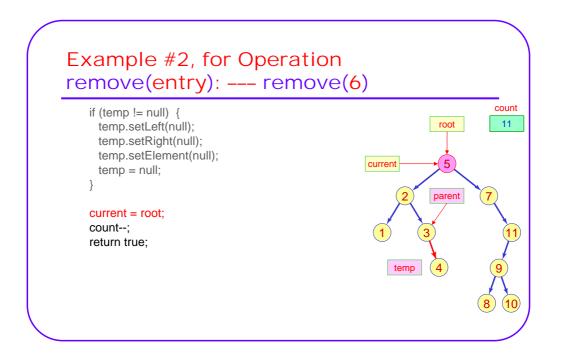
### Example #2, for Operation remove(entry): --- remove(6) count parent = current; 11 root temp = current.getLeft(); while (temp.getRight() != null) { parent = temp; current 5 temp = temp.getRight(); parent current.setElement(temp.getElement()); if (temp == parent.getLeft()) parent.setLeft(temp.getLeft()); temp → if (temp == parent.getRight()) parent.setRight(temp.getLeft());



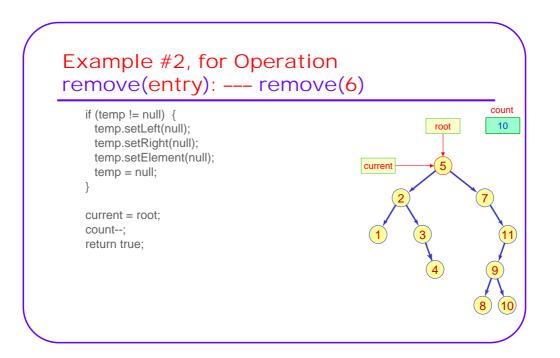
### Example #2, for Operation remove(entry): --- remove(6) count parent = current; 11 temp = current.getLeft(); root while (temp.getRight() != null) { parent = temp; current [5] temp = temp.getRight(); parent current.setElement(temp.getElement()); if (temp == parent.getLeft()) parent.setLeft(temp.getLeft()); temp if (temp == parent.getRight()) parent.setRight(temp.getLeft());



```
Example #2, for Operation
remove(entry): --- remove(6)
                                                                  count
   if (temp != null) {
                                                                   11
                                                         root
    temp.setLeft(null);
    temp.setRight(null);
    temp.setElement(null);
                                              current
                                                          [5]
    temp = null;
                                                        parent
   current = root;
   count--;
   return true;
```



### Example #2, for Operation remove(entry): --- remove(6) count if (temp != null) { 10 root temp.setLeft(null); temp.setRight(null); temp.setElement(null); current [5] temp = null; parent current = root; count--; return true; temp



# Mapping Operation: replace(entry) ... O(depth)

- Precondition:
  - size() > 0.
- Postcondition:
  - The data at the <u>current</u> node has been replaced with the new entry, and it has been moved to its proper location in the tree so that the binary search tree's integrity is maintained.
- Code:

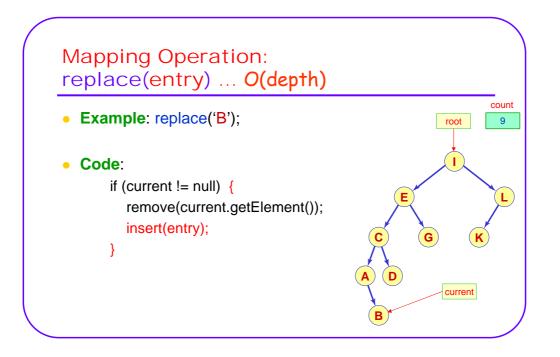
```
if (current != null) {
   remove(current.getElement());
   insert(entry);
}
```

# Mapping Operation: replace(entry) ... O(depth) • Example: replace('B'); • Code: if (current != null) { remove(current.getElement()); insert(entry); }

```
Mapping Operation:
replace(entry) ... O(depth)

• Example: replace('B');

• Code:
    if (current != null) {
        remove(current.getElement());
        insert(entry);
    }
```

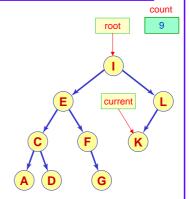


# Mapping Operation: retrieve() ... O(1)

- Precondition:
  - size() > 0.
- Postcondition:
  - The return value is the element from the current node.
- Code:

if (current != null)
 return current.getElement();

Example: retrieve() > returns 'K'.



# Mapping Operation: Tree Traversal print(p, depth) ... O(n)

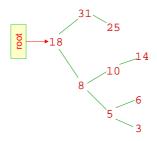
- Precondition:
  - p is a pointer to a node in a binary tree or null to indicate the empty tree.
  - If the pointer is not null, then depth is the level of the node at p.
- Postcondition:
  - If p is non-null, then the elements of node p and all its descendants have been written out, using a backward in-order traversal. Each node is indented four times its depth, so to resemble the tree shape.
- Code:

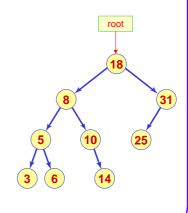
```
 \begin{aligned} & \text{String s} = \text{``''}; \\ & \text{if (p} \sqsubseteq \text{null)} \left\{ \\ & \text{print(p.getRight(), depth+1);} \\ & \text{for (int } \sqsubseteq 0; \text{i} < \text{depth; i++) } \text{s} + \text{``'}; \\ & \text{System.out.println(s+p.getElement());} \\ & \text{/'} \text{ Print the node contents } \\ & \text{print(p.getLeft(), depth+1);} \\ & \text{/'} \text{ Do the left branch last} \end{aligned}
```



# Mapping Operation: Tree Traversal print(p, depth) ... O(n)

- Example: print (root,0);
- Result:





# Mapping Operation: treeClear(p) ... O(n)

- Precondition:
  - p is the root pointer of a binary tree which may be null for the empty tree.
- Postcondition:
  - All nodes at the root or below have been returned to the heap, and root have been set to null. The tree is now empty.
- Code:

```
if (p!= null) {
    treeClear( p.getLeft( ) );
    treeClear( p.getRight( ) );
    if (p == root) {
        root = null;
        current == null;
        count = 0;
    }
    p.element = null;
    p := null;
```



# BST Implementation Performance Issues

- Most of the operations of the BST are O(1).
- All traversal operations are O(n).
- The most important operation of the BST is the search operation. We define the following:
  - The search length of a node:
    - The number of nodes examined to find the node.
  - The search length of a tree:
    - The average of the search lengths for all of its nodes
- The time complexity of the search operation depends on the height of the tree rather than its size.
- There are three more operations that need to find a node before proceeding, these are: insert, remove, and replace. Their time complexity depends on the time complexity of the search operation.

# BST Implementation Performance Issues (cont.)

- Binary trees have many shapes, but can be grouped into one of the following categories:
  - 1. The degenerate binary tree: Where each node has exactly one child, except the only leaf node.
  - 2. The random binary tree: Which is formed by adding elements in random order to an empty BST.
  - 3. The minimum height binary tree: as given before.
  - 4. The full binary tree: as given before.



# BST Implementation Performance Issues (cont.)

### Search lengths for binary trees of size n:

Tree Type	Average Search Length	Maximum Search Length of any node
Degenerate	½ (n+1)	n
Random	1.4 log <sub>2</sub> (n+1)	n
Minimum Height	log <sub>2</sub> (n+1)	log <sub>2</sub> (n+1)
Full	log <sub>2</sub> (n+1)	log <sub>2</sub> (n+1)

The search operation for a random BST is O(log<sub>2</sub> n)

# Binary Search Tree Application in Sorting a Linear List -- tree sort

Beside its use in searching, a BST can also be used for sorting:

### Tree Sort Algorithm:

Given a linear list of items, A, produce a sorted list of them, B.

- 1. Form a binary search tree by inserting the items from A, one by one, into an empty BST. O(n log n)
- 2. Use the inorder traversal operation on the BST to get the sorted list, B. O(n)

### Example:

