

# Algorithms Notebook

Version 0.0.2

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# Chapter 1

# Introduction

Algorithms and data structures notes.

# Chapter 2

## Algorithm Analysis

### 2.1 Binary Logarithm

Logarithm is the inverse function to exponentiation which means that the logarithm of a number  $x$  to the base  $b$  is the exponent to which  $b$  must be raised to produce  $x$ . The binary logarithm is the power to which the number 2 must be raised to obtain the value  $n$ .

- The number of bits in the binary representation of a positive integer  $n$  is the integral part of  $\log_2(n) + 1$ .
- The binary logarithm also frequently appears in the analysis of algorithms because binary logarithms occur in the analysis of algorithms based on two-way branching. If a problem initially has  $n$  choices for its solution, and each iteration of the algorithm reduces the number of choices by a factor of two, then the number of iterations needed to select a single choice is again the integral part of  $\log_2(n)$ .

— [https://en.wikipedia.org/wiki/Binary\\_logarithm](https://en.wikipedia.org/wiki/Binary_logarithm)

### 2.2 Amortized Analysis

The motivation for amortized analysis is that looking at the worst-case run time can be too pessimistic. Instead, amortized analysis averages the running times of operations in a sequence over that sequence. Amortized

analysis is a useful tool that complements other techniques such as worst-case and average-case analysis.

Imagine a dynamic list backed by a fixed size array that doubles once capacity is reached and require  $n$  inserts into the new fixed size array.

n	# of inserts
1	1
2	1
3	1
...	1
n	n
...	n + 1

$$\frac{\# \text{ of inserts}}{n} \rightarrow \frac{n+n+1}{n} \rightarrow \frac{2n+1}{n} \rightarrow \frac{2n}{n} \rightarrow 2$$

— [https://en.wikipedia.org/wiki/Amortized\\_analysis](https://en.wikipedia.org/wiki/Amortized_analysis)

# Chapter 3

## Data Structures

### 3.1 Heaps

#### 3.1.1 Applications

- Heapsort
- Priority queue
- K-way merge

# Chapter 4

## Problem Solving Methods

### 4.1 Greedy Algorithms

#### 4.1.1 Overview

A greedy algorithm is any algorithm that follows the problem-solving heuristic of making the locally optimal choice at each stage. In many problems, a greedy strategy does not produce an optimal solution, but a greedy heuristic can yield locally optimal solutions that approximate a globally optimal solution in a reasonable amount of time.

— [https://en.wikipedia.org/wiki/Greedy\\_algorithm](https://en.wikipedia.org/wiki/Greedy_algorithm)

#### 4.1.2 Dijkstra's Shortest Path Algorithm

Dijkstra's shortest path algorithm is a search algorithm that finds the shortest path between a vertex and other vertices in a weighted graph.

### 4.2 Recursion

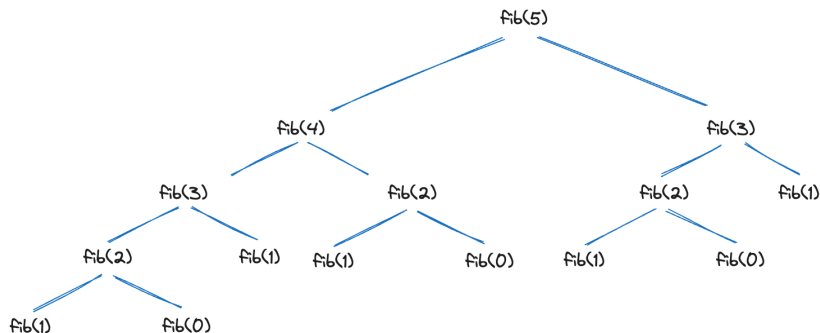
The power of recursion evidently lies in the possibility of defining an infinite set of objects by a finite statement. In the same manner, an infinite number of computations can be described by a finite recursive program, even if this program contains no explicit repetitions.

### 4.2.1 Fibonacci Sequence

$$F_0 = 0$$

$$F_1 = 1$$

$$F_n = F_{n-1} + F_{n-2} \quad \text{for } n > 1$$



$$F_n = F_{n-1} + F_{n-2}$$

$$F_5 = F_4 + F_3$$

$$F_5 = (F_3 + F_2) + (F_2 + F_1)$$

$$F_5 = (((F_2 + F_1) + (F_1 + F_0)) + ((F_1 + F_0) + F_1))$$

$$F_5 = (((((F_1 + F_0) + F_1) + (F_1 + F_0)) + ((F_1 + F_0) + F_1)))$$

$$F_5 = (((((1 + 0) + 1) + (1 + 0)) + ((1 + 0) + 1)))$$

$$F_5 = 5$$

```
export function fib(n: number): number {  
  if (n == 0 || n == 1) {  
    return n  
  }  
  return fib(n - 1) + fib(n - 2)  
}
```

## 4.3 Probability

<sup>1</sup><https://archive.org/details/algorithmsdatast00wirth/page/126>



# Chapter 5

# Appendix

## 5.1 Resources

- LeetCode
- Project Euler
- The Algorithm Design Manual
- Elements of Programming Interviews