Artificial Intelligence

Constraint Satisfaction Problems





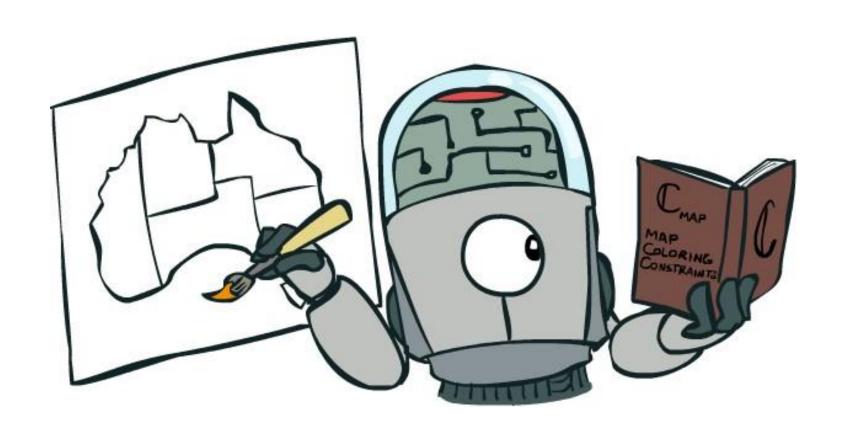
What is Search For?

 Assumptions about the world: a single agent, deterministic actions, fully observed state, discrete state space

- Planning: sequences of actions
 - The path to the goal is the important thing
 - Paths have various costs, depths
 - Heuristics give problem-specific guidance
- Identification: assignments to variables
 - The goal itself is important, not the path
 - All paths at the same depth (for some formulations)
 - CSPs are a specialized class of identification problems

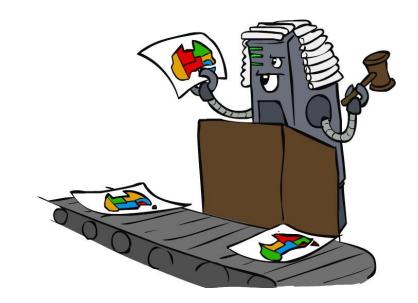


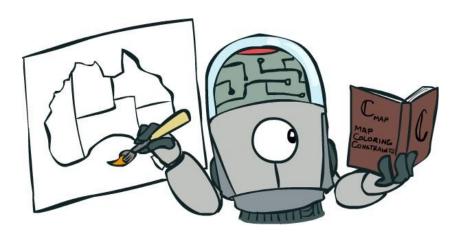
Constraint Satisfaction Problems



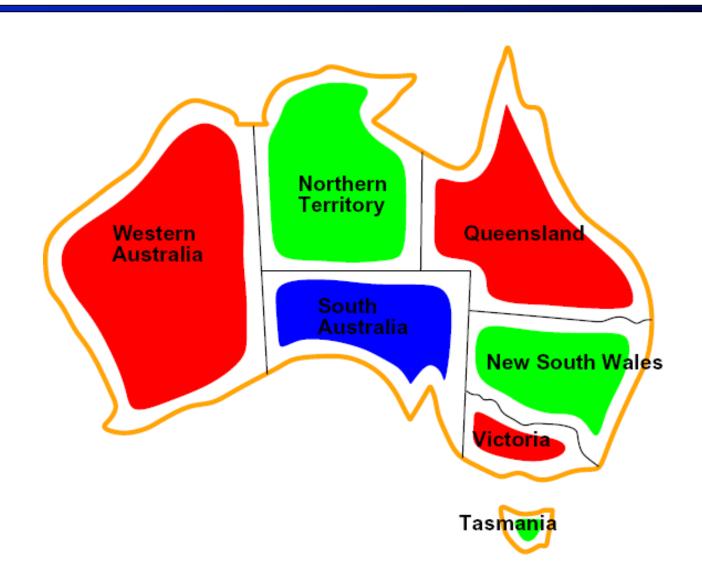
Constraint Satisfaction Problems

- Standard search problems:
 - State is a "black box": arbitrary data structure
 - Goal test can be any function over states
 - Successor function can also be anything
- Constraint satisfaction problems (CSPs):
 - A special subset of search problems
 - State is defined by variables X_i with values from a domain D (sometimes D depends on i)
 - Goal test is a set of constraints specifying allowable combinations of values for subsets of variables
- Simple example of a formal representation language
- Allows useful general-purpose algorithms with more power than standard search algorithms





CSP Examples



Example: Map Coloring

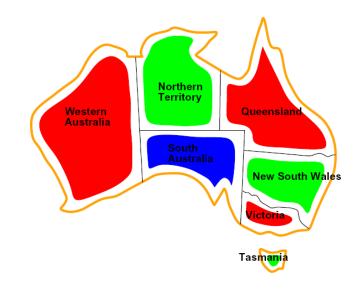
- Variables: WA, NT, Q, NSW, V, SA, T
- Domains: D = {red, green, blue}
- Constraints: adjacent regions must have different colors

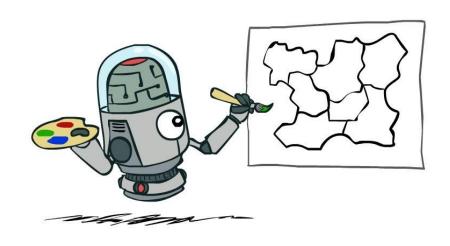
Implicit: $WA \neq NT$

Explicit: $(WA, NT) \in \{(red, green), (red, blue), \ldots\}$

Solutions are assignments satisfying all constraints, e.g.:

{WA=red, NT=green, Q=red, NSW=green, V=red, SA=blue, T=green}





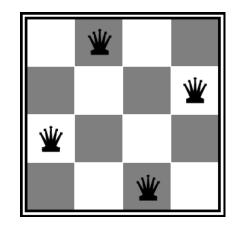
Example: N-Queens

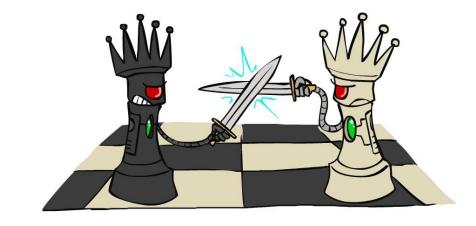
Formulation 1:

• Variables: X_{ij}

■ Domains: {0, 1}

Constraints





$$\forall i, j, k \ (X_{ij}, X_{ik}) \in \{(0,0), (0,1), (1,0)\}$$

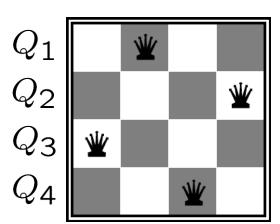
 $\forall i, j, k \ (X_{ij}, X_{kj}) \in \{(0,0), (0,1), (1,0)\}$
 $\forall i, j, k \ (X_{ij}, X_{i+k,j+k}) \in \{(0,0), (0,1), (1,0)\}$
 $\forall i, j, k \ (X_{ij}, X_{i+k,j-k}) \in \{(0,0), (0,1), (1,0)\}$

$$\sum_{i,j} X_{ij} = N$$

Example: N-Queens

Formulation 2:

- lacktriangle Variables: Q_k
- Domains: $\{1, 2, 3, ... N\}$



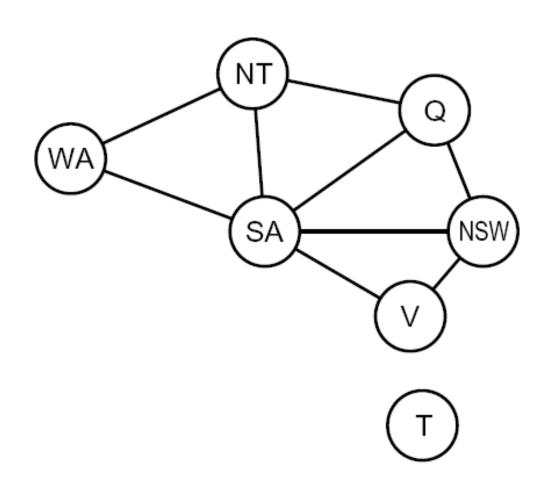
Constraints:

Implicit: $\forall i, j$ non-threatening (Q_i, Q_j)

Explicit: $(Q_1, Q_2) \in \{(1,3), (1,4), \ldots\}$

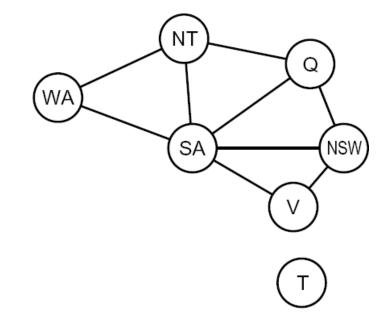
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Constraint Graphs



Constraint Graphs

- Binary CSP: each constraint relates (at most) two variables
- Binary constraint graph: nodes are variables, arcs show constraints
- General-purpose CSP algorithms use the graph structure to speed up search. E.g., Tasmania is an independent subproblem!



Example: Cryptarithmetic

Variables:

$$F T U W R O X_1 X_2 X_3$$

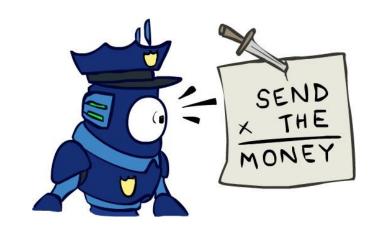
Domains:

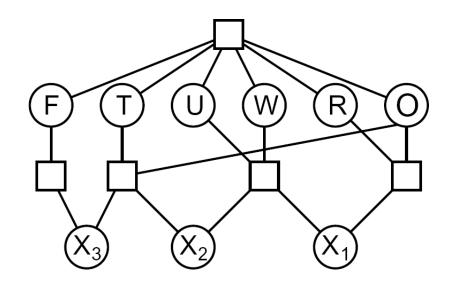
$$\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$$

Constraints:

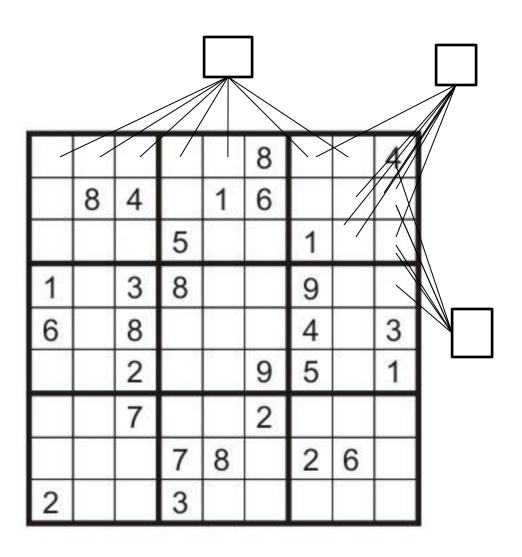
$$O + O = R + 10 \cdot X_1$$

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Example: Sudoku



- Variables:
 - Each (open) square
- Domains:
 - **•** {1,2,...,9}
- Constraints:

9-way alldiff for each column

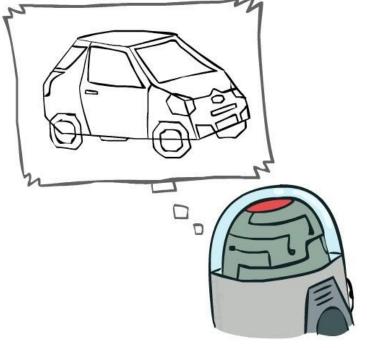
9-way alldiff for each row

9-way alldiff for each region

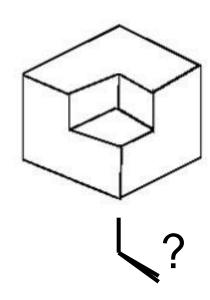
(or can have a bunch of pairwise inequality constraints)

Example: The Waltz Algorithm

- The Waltz algorithm is for interpreting line drawings of solid polyhedra as 3D objects
- An early example of an AI computation posed as a CSP







Approach:

- Each intersection is a variable
- Adjacent intersections impose constraints on each other
- Solutions are physically realizable 3D interpretations

Varieties of CSPs and Constraints



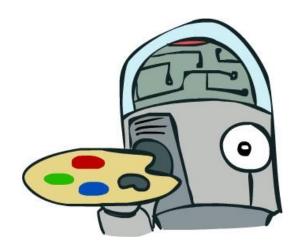
Varieties of CSPs

Discrete Variables

- Finite domains
 - Size d means $O(d^n)$ complete assignments
 - E.g., Boolean CSPs, including Boolean satisfiability (NP-complete)
- Infinite domains (integers, strings, etc.)
 - E.g., job scheduling, variables are start/end times for each job
 - Linear constraints solvable, nonlinear undecidable

Continuous variables

- E.g., start/end times for Hubble Telescope observations
- Linear constraints solvable in polynomial time by LP methods (see cs170 for a bit of this theory)





Varieties of Constraints

Varieties of Constraints

Unary constraints involve a single variable (equivalent to reducing domains), e.g.:

$$SA \neq green$$

Binary constraints involve pairs of variables, e.g.:

$$SA \neq WA$$

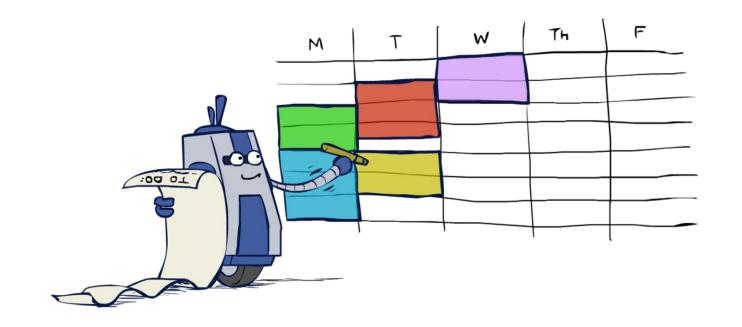
 Higher-order constraints involve 3 or more variables: e.g., cryptarithmetic column constraints



- Preferences (soft constraints):
 - E.g., red is better than green
 - Often representable by a cost for each variable assignment
 - Gives constrained optimization problems
 - (We'll ignore these until we get to Bayes' nets)

Real-World CSPs

- Scheduling problems: e.g., when can we all meet?
- Timetabling problems: e.g., which class is offered when and where?
- Assignment problems: e.g., who teaches what class
- Hardware configuration
- Transportation scheduling
- Factory scheduling
- Circuit layout
- Fault diagnosis
- ... lots more!



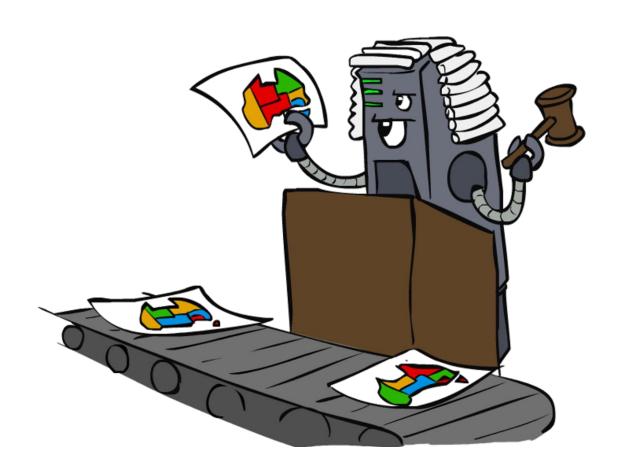
Many real-world problems involve real-valued variables...

Solving CSPs



Standard Search Formulation

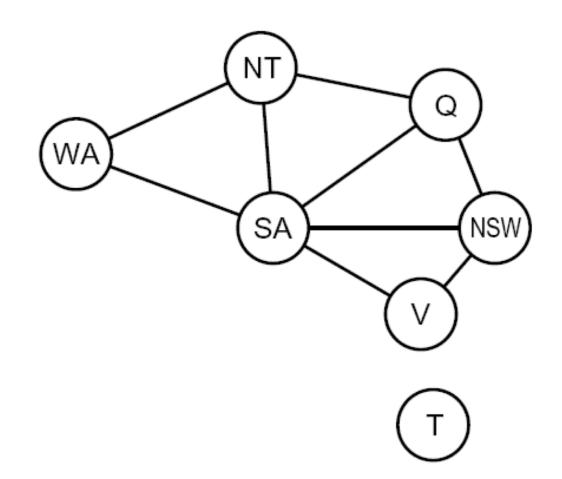
- Standard search formulation of CSPs
- States defined by the values assigned so far (partial assignments)
 - Initial state: the empty assignment, {}
 - Successor function: assign a value to an unassigned variable
 - Goal test: the current assignment is complete and satisfies all constraints
- We'll start with the straightforward, naïve approach, then improve it



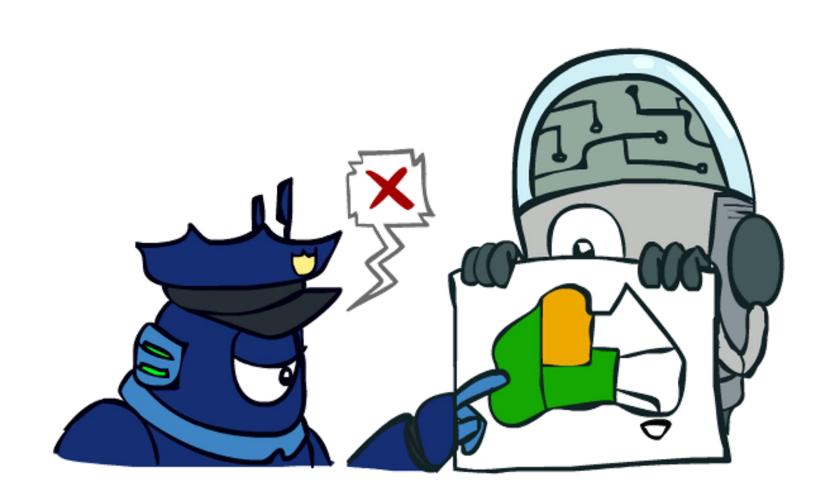
Search Methods

What would BFS do?

What would DFS do?

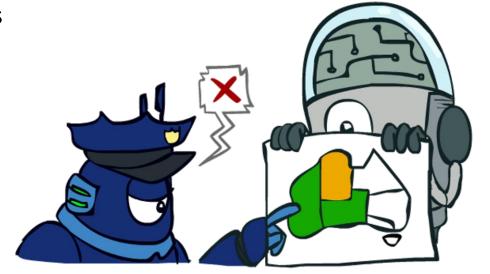


Backtracking Search

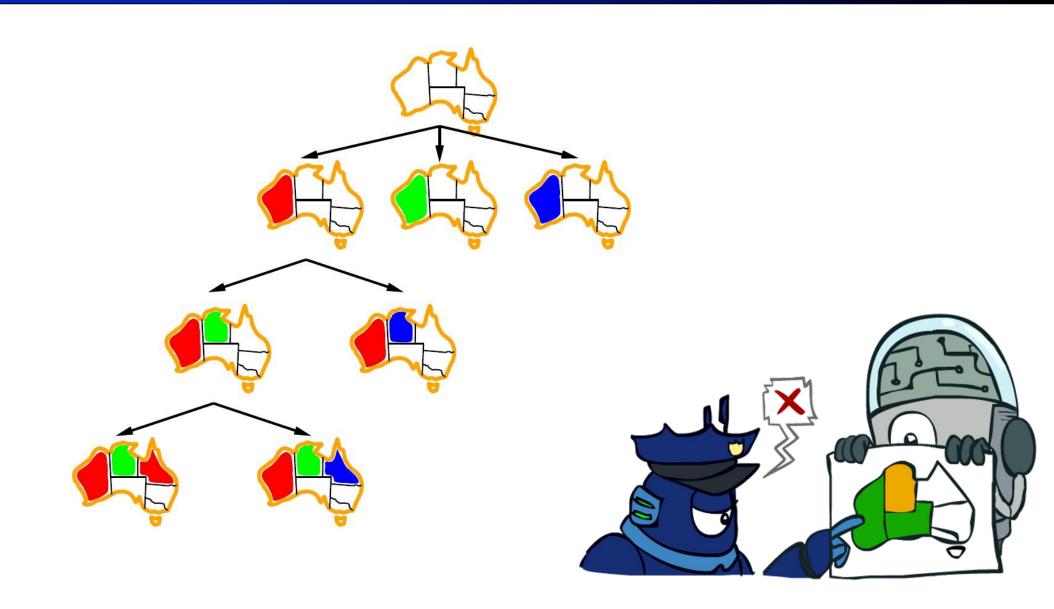


Backtracking Search

- Backtracking search is the basic uninformed algorithm for solving CSPs
- Idea 1: One variable at a time
 - Variable assignments are commutative, so fix ordering
 - I.e., [WA = red then NT = green] same as [NT = green then WA = red]
 - Only need to consider assignments to a single variable at each step
- Idea 2: Check constraints as you go
 - I.e. consider only values which do not conflict with previous assignments
 - Might have to do some computation to check the constraints
 - "Incremental goal test"
- Depth-first search with these two improvements is called backtracking search (not the best name)
- Can solve n-queens for $n \approx 25$



Backtracking Example



Backtracking Search

```
function BACKTRACKING-SEARCH(csp) returns solution/failure
  return Recursive-Backtracking({ }, csp)
function Recursive-Backtracking (assignment, csp) returns soln/failure
   if assignment is complete then return assignment
   var \leftarrow \text{Select-Unassigned-Variable}(\text{Variables}[csp], assignment, csp)
  for each value in Order-Domain-Values (var, assignment, csp) do
       if value is consistent with assignment given Constraints [csp] then
           add \{var = value\} to assignment
           result \leftarrow Recursive-Backtracking(assignment, csp)
           if result \neq failure then return result
           remove \{var = value\} from assignment
  return failure
```

- Backtracking = DFS + variable-ordering + fail-on-violation
- What are the choice point?

Improving Backtracking

- General-purpose ideas give huge gains in speed
- Ordering:
 - Which variable should be assigned next?
 - In what order should its values be tried?
- Filtering: Can we detect inevitable failure early?
- Structure: Can we exploit the problem structure?

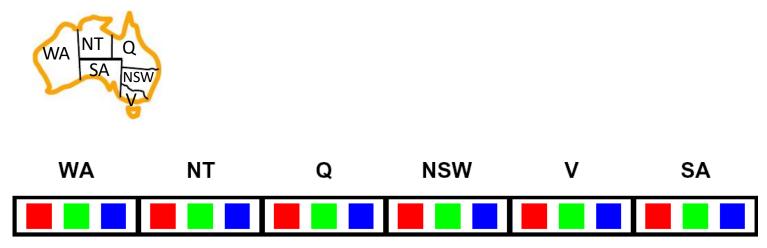


Filtering



Filtering: Forward Checking

- Filtering: Keep track of domains for unassigned variables and cross off bad options
- Forward checking: Cross off values that violate a constraint when added to the existing assignment



Filtering: Constraint Propagation

 Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:

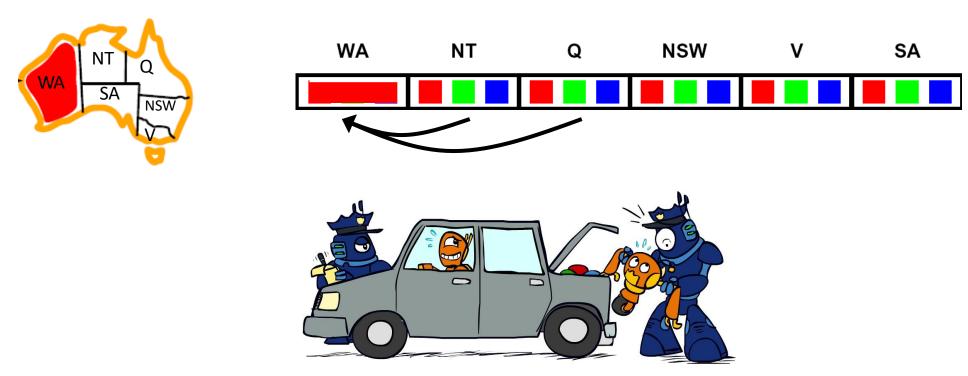




- NT and SA cannot both be blue!
- Why didn't we detect this yet?
- Constraint propagation: reason from constraint to constraint

Consistency of A Single Arc

An arc X → Y is consistent iff for every x in the tail there is some y in the head which could be assigned without violating a constraint

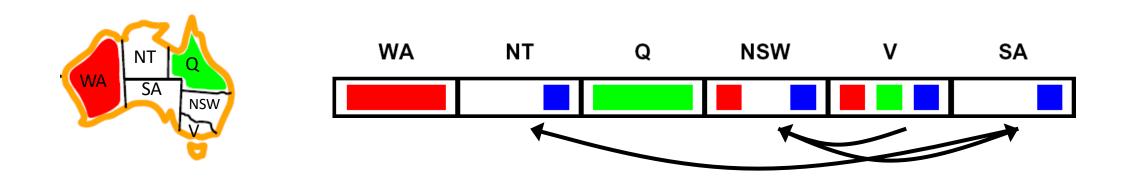


Delete from the tail!

Forward checking: Enforcing consistency of arcs pointing to each new assignment

Arc Consistency of an Entire CSP

A simple form of propagation makes sure all arcs are consistent:



- Important: If X loses a value, neighbors of X need to be rechecked!
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment
- What's the downside of enforcing arc consistency?

Remember: Delete from the tail!

Limitations of Arc Consistency

- After enforcing arc consistency:
 - Can have one solution left
 - Can have multiple solutions left
 - Can have no solutions left (and not know it)

• Arc consistency still runs inside a backtracking search!

