

Diskretne strukture UNI

Vaje 3

1. (a) Pokaži, da tromestni veznik $A(p, q, r) \equiv r \Rightarrow (\neg p \wedge \neg q)$ predstavlja poln nabor veznikov.
- (b) Zaporedje izrazov A_n je definirano rekurzivno z

$$\begin{aligned} A_0 &= \neg p \\ A_n &= A(p, A_{n-1}, 1). \end{aligned}$$

Izračunaj A_{2019} .

a) $A(0,0,0) = 0 \Rightarrow (\neg 0 \wedge \neg 0) \sim 1$

$$A(1,1,1) = 1 \Rightarrow (\neg 1 \wedge \neg 1) \sim 1 \Rightarrow \neg 1 \sim 1 \Rightarrow 0 \sim 0$$

Ne obravnava konstant. Morda je poln.

$$\mathcal{P} = \{ \neg, \Rightarrow \} \quad \mathcal{A} = \{ A \}$$

Vezniki iz \mathcal{P} izrazimo z vezniki iz \mathcal{A} :

$$A(p, p, p) = p \Rightarrow \neg p \wedge \neg p \sim p \Rightarrow \neg p \sim \neg p \vee \neg p \sim \neg p$$

$$\bullet \underline{\neg p \sim A(p, p, p)}$$

$$\bullet \underline{p \Rightarrow q \sim p \Rightarrow q \wedge q \sim p \Rightarrow \neg(\neg q) \wedge \neg(\neg q) \sim A(\neg q, \neg q, p) \sim \underline{A(A(\neg q, \neg q, \neg q), A(\neg q, \neg q, \neg q), p)}}$$

$\Rightarrow \{A\}$ je poln nabor.

b) $A_0 = \neg p$

$$A(p, q, r) \equiv r \Rightarrow (\neg p \wedge \neg q)$$

$$A_n = A(p, A_{n-1}, 1).$$

$$A_1 = A(p, A_0, 1) = 1 \Rightarrow (\neg p \wedge \neg(\neg p)) \sim 1 \Rightarrow (\neg p \wedge p) \sim 1 \Rightarrow 0 \sim 0$$

$$A_2 = A(p, A_1, 1) = 1 \Rightarrow (\neg p \wedge \neg 0) \sim 1 \Rightarrow \neg p \sim \neg 1 \vee \neg p \sim 0 \vee \neg p \sim \neg p$$

$$A_3 = A(p, \underbrace{A_2}_{= A_0}, 1) = 0$$

$$A_4 = A(p, \underbrace{A_3}_{= A_1}, 1) = \neg p$$

\vdots

$$A_{2k} = \neg p \quad A_{2k+1} = 0$$

$$\underline{\underline{A_{2019} = 0}}$$

2. Naj bo A veznik $A(p, q, r) \equiv (p \vee q) \Rightarrow r$.

- (a) Kateri izmed naborov $\{A\}$, $\{A, 1\}$, $\{A, 0\}$, $\{A, \neg\}$ so polni?
 (b) Zaporedje izrazov A_n je definirano rekurzivno z

$$\begin{aligned} A_0 &= \neg p \\ A_1 &= \neg q \\ A_n &= A(p, q, A_{n-1} \wedge A_{n-2}) \end{aligned}$$

Izračunaj A_{2019} .

a) $A(0, 0, 0) = 0 \vee 0 \Rightarrow 0 \sim 0 \Rightarrow 0 \sim 1$

$A(1, 1, 1) = 1 \vee 1 \Rightarrow 1 \sim 1 \Rightarrow 1 \sim 1$

$\Rightarrow A$ ohranja enice $\Rightarrow \{A\}$ in $\{A, 1\}$ nista polna, ker ohranjata enice.

$\mathcal{N} = \{A, 0\}$, $\mathcal{P} = \{\neg, \Rightarrow\}$

\mathcal{P} izrazimo z \mathcal{N} :

• $\neg p \sim p \Rightarrow 0 \sim p \vee p \Rightarrow 0 \sim \underline{\underline{A(p, p, 0)}}$

• $p \Rightarrow q \sim p \vee p \Rightarrow q \sim \underline{\underline{A(p, p, q)}}$

$\Rightarrow \{A, 0\}$ je poln.

$\mathcal{N} = \{A, \neg\}$, $\mathcal{P} = \{\neg, \Rightarrow\}$

\mathcal{P} izrazimo z \mathcal{N} :

• $\neg p \sim \neg p$

• $p \Rightarrow q \sim p \vee p \Rightarrow q \sim \underline{\underline{A(p, p, q)}}$

$\Rightarrow \{A, \neg\}$ je poln.

$A(p, q, r) \equiv (p \vee q) \Rightarrow r$.

$A_0 = \neg p$

$A_1 = \neg q$

$A_n = A(p, q, A_{n-1} \wedge A_{n-2})$

b) $A_0 = \neg p$

$A_1 = \neg q$

$A_2 = A(p, q, A_1 \wedge A_0) = A(p, q, \neg q \wedge \neg p) = p \vee q \Rightarrow \neg p \wedge \neg q \sim p \vee q \Rightarrow \neg(p \vee q) \sim \neg(p \vee q) \vee \neg(p \vee q) \sim$
 $\sim \neg(p \vee q) \sim \neg p \wedge \neg q$

$A_3 = A(p, q, A_2 \wedge A_1) = A(p, q, \neg p \wedge \neg q \wedge \neg q) = A(p, q, \neg p \wedge \neg q) = \neg p \wedge \neg q$

$A_n = A(p, q, A_3 \wedge A_2) = A(p, q, \neg p \wedge \neg q \wedge \neg p \wedge \neg q) = A(p, q, \neg p \wedge \neg q) = \neg p \wedge \neg q$

\vdots

$A_{2019} = \neg p \wedge \neg q$

3. Veznik A je definiran s predpisom $A(p, q, r) \equiv (p \wedge q) \vee (\neg p \wedge \neg r)$.

- (a) Samo z veznikom A zapiši izraze 1 , $p \wedge q$ in $p \Rightarrow q$.
 (b) Kateri izmed naborov $\{A\}$, $\{A, 1\}$, $\{A, 0\}$, $\{A, \Rightarrow\}$, $\{A, \vee\}$ so polni?
 (c) Zaporedje izrazov I_n je definirano rekurzivno s predpisi

$$\begin{aligned} I_0 &= \neg p \\ I_1 &= p \\ I_n &= A(I_{n-1}, I_{n-2}, I_{n-2}) \end{aligned}$$

Izračunaj I_{2019}

a) $\underline{A(p, p, p)} = (p \wedge p) \vee (\neg p \wedge \neg p) \sim p \vee \neg p \sim \underline{\underline{1}}$

$$A(p, q, r) \equiv \underbrace{(p \wedge q)}_{p \wedge q} \vee \underbrace{(\neg p \wedge \neg r)}_{0 \sim \neg p \wedge 0 \sim \neg p \wedge \neg 1}$$

$$A(p, 2, 1) \sim p \wedge 2 \vee \neg p \wedge \neg 1 \sim p \wedge 2 \vee 0 \sim \underline{p \wedge 2} \sim \underline{A(p, 2, A(p, p, p))}$$

$$A(p, q, r) \equiv (p \wedge q) \vee (\neg p \wedge \neg r) \sim \neg(p \vee \neg) \vee p \wedge q \sim \underbrace{p \vee \neg}_{p} \Rightarrow \underbrace{p \wedge q}_q$$

$$A(p, 2, p) \sim p \vee p \Rightarrow p \wedge 2 \sim p \Rightarrow p \wedge 2 \sim \neg p \vee (p \wedge 2) \sim \underbrace{(\neg p \vee p)}_1 \wedge (\neg p \vee 2) \sim \neg p \vee 2 \sim p \Rightarrow 2$$

$$\underline{p \Rightarrow 2} \sim \underline{A(p, 2, p)}^*$$

a) $A(1, 1, 1) \sim 1 \Rightarrow \{A\}, \{A, 1\}, \{A, \Rightarrow\}$ niso polni, ker (vsi vezniki) ohranjajo enice.

$$\mathcal{N} = \{A, 0\} \quad \mathcal{P} = \{\neg, \Rightarrow\}$$

\mathcal{P} izrazimo z \mathcal{N} :

$$\bullet \underline{p \Rightarrow 2} \sim \underline{A(p, 2, p)}^*$$

$$\bullet \underline{\neg p} \sim p \Rightarrow 0 \sim \underline{A(p, 0, p)} \quad \{A, 0\} \text{ je polni nabor.}$$

$$\mathcal{N} = \{A, \vee\} \quad \mathcal{P} = \{\neg, \Rightarrow\}$$

\mathcal{P} izrazimo z \mathcal{N} :

$$\bullet \underline{p \Rightarrow 2} \sim \underline{A(p, 2, p)}^*$$

$$\bullet \underline{\neg p} \sim p \vee p \quad \{A, \vee\} \text{ je polni nabor.}$$

$$\textcircled{d} \quad \left. \begin{array}{l} I_0 = \neg p \\ I_1 = p \end{array} \right\}$$

$$I_n = A(I_{n-1}, I_{n-2}, I_{n-2})$$

$$A(p, q, r) \equiv (p \wedge q) \vee (\neg p \wedge \neg r)$$

$$I_2 = A(I_1, I_0, I_0) = A(p, \neg p, \neg p) = (p \wedge \neg p) \vee (\neg p \wedge \neg \neg p) \sim p \wedge \neg p \sim 0$$

$$I_3 = A(I_2, I_1, I_1) = A(0, p, p) = (0 \wedge p) \vee (\neg 0 \wedge \neg p) \sim 0 \vee \neg p \sim \neg p \quad \left. \begin{array}{l} \\ \\ \end{array} \right\}$$

$$I_4 = A(I_3, I_2, I_2) = A(\neg p, 0, 0) = (\neg p \wedge 0) \vee (\neg \neg p \wedge \neg 0) \sim 0 \vee p \sim p$$

$$I_5 = I_2 = 0$$

$$I_6 = I_3 = \neg p$$

⋮

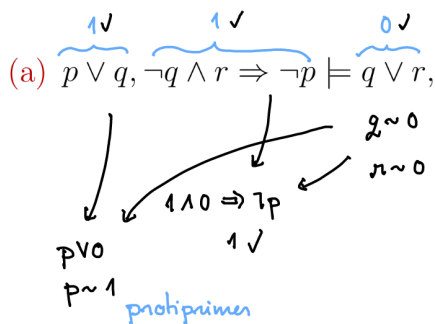
$$I_{3k} = \neg p \longleftarrow 2019 : 3 = 673, 0 \text{ ost.}$$

$$I_{3k+1} = p \qquad I_{2019} = \neg p$$

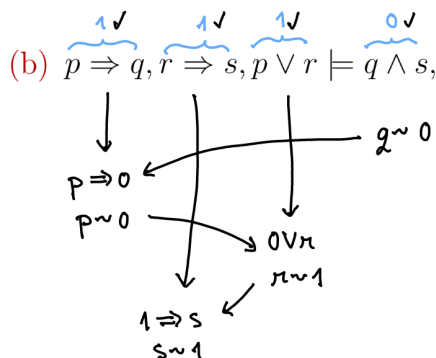
$$I_{3k+2} = 0$$

4. Kateri od naslednjih sklepov so pravilni?

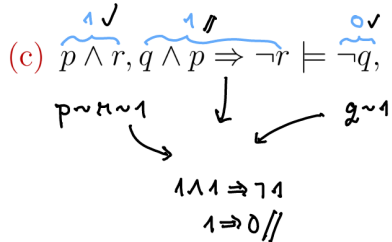
- (a) $p \vee q, \neg q \wedge r \Rightarrow \neg p \models q \vee r,$
 (b) $p \Rightarrow q, r \Rightarrow s, p \vee r \models q \wedge s,$
 (c) $p \wedge r, q \wedge p \Rightarrow \neg r \models \neg q,$
 (d) $p \Rightarrow q, p \vee s, q \Rightarrow r, s \Rightarrow t, \neg r \models t,$
 (e) $p \Rightarrow q, p \wedge s, q \wedge r \Rightarrow t, s \Rightarrow r \models t,$
 (f) $p \Leftrightarrow q, \neg p, \neg(q \Rightarrow r) \vee t, s \vee t \Rightarrow r \models r \wedge \neg p,$



Pri $p \sim 1, q \sim 0$ sta predpostavki pravilni, zaključek pa napačen. Sklep je napačen.



Pri $p \sim q \sim 0, r \sim s \sim 1$ so predpostavke pravilne, zaključek pa napačen. Sklep je napačen.



Ne najdemo protiprimerov. Sklep je pravilen. Dokaži:

1. $p \wedge r$
2. $q \wedge p \Rightarrow \neg r$
3. p
4. r
5. $\neg(q \wedge p)$
6. $\neg q \vee \neg p$
7. $\neg q$

} predp.

$P_0(1)$

$P_0(1)$

MT(2,4)

$\sim 5.$

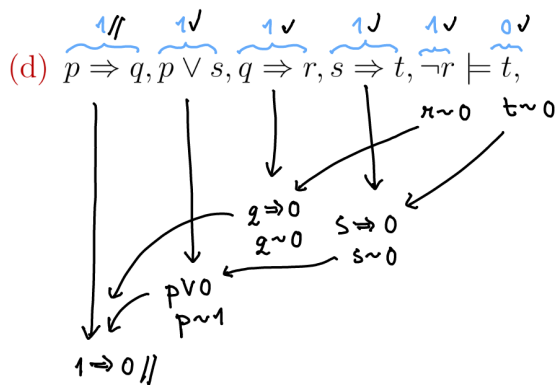
DS(6,3)

$\vdash \neg q$

POENOSTAVITEV P_0
 $A \wedge B \models A$

MODUS TOLLENS
 $A \Rightarrow B, \neg B \models \neg A$

DISJUNKTIVNI SILOGIZEM
 $A \vee B, \neg A \models B$



Dokaz:

- | | | |
|----------------------|-------------------|------------|
| 1. $p \Rightarrow q$ | } predp. | $\vdash t$ |
| 2. $p \vee s$ | | |
| 3. $q \Rightarrow r$ | | |
| 4. $s \Rightarrow t$ | | |
| 5. $\neg r$ | | |
| 6. $\neg q$ | MT(3,5) | |
| 7. $\neg p$ | MT(1,6) | |
| 8. s | DS(2,7) | |
| 9. t | MP(4,8) | |

MODUS PONENS
 $A \Rightarrow B, A \vdash B$

(e) $p \Rightarrow q, p \wedge s, q \wedge r \Rightarrow t, s \Rightarrow r \models t,$

- | | | |
|-------------------------------|-------------------|------------|
| 1. $p \Rightarrow q$ | } predp. | $\vdash t$ |
| 2. $p \wedge s$ | | |
| 3. $q \wedge r \Rightarrow t$ | | |
| 4. $s \Rightarrow r$ | | |
| 5. p | PO(2) | |
| 6. s | PO(2) | |
| 7. r | MP(4,6) | |
| 8. q | MP(1,5) | |
| 9. $q \wedge r$ | Zd(7,8) | |
| 10. t | MP(3,9) | |

ZDRUŽITEV
 $A, B \vdash A \wedge B$

★ $A \Leftrightarrow B \sim (A \Rightarrow B) \wedge (B \Rightarrow A)$

(f) $p \Leftrightarrow q, \neg p, \neg(q \Rightarrow r) \vee t, s \vee t \Rightarrow r \models r \wedge \neg p.$

- | | | |
|---|-------------------|--------------------------|
| 1. $p \Leftrightarrow q$ ★ | } predp. | $\vdash r \wedge \neg p$ |
| 2. $\neg p$ | | |
| 3. $\neg(q \Rightarrow r) \vee t$ | | |
| 4. $s \vee t \Rightarrow r$ | | |
| 5. $(p \Rightarrow q) \wedge (q \Rightarrow p)$ | $\sim 1.$ ★ | |
| 6. $p \Rightarrow q$ | PO(5) | |
| 7. $q \Rightarrow p$ | PO(5) | |
| 8. $\neg(\neg q \vee r) \vee t$ | $\sim 3.$ | |
| 9. $\neg q$ | MT(7,2) | |
| 10. $\neg q \vee r$ | PR(9,4) | |

PRIDRUŽITEV \neg
 $A \vdash A \vee B$

11. t $DS(8, 10)$
12. $s \vee t$ $Pr(11, s)$
13. r $MP(4, 12)$
14. $r \wedge 7p$ $Zd(13, 2)$

HIPOTETIČNI SILOGIZEM

$A \Rightarrow B, B \Rightarrow C \models A \Rightarrow C$