

Example 19.1

Change the following IPv4 addresses from binary notation to dotted-decimal notation.

- a. 10000001 00001011 00001011 11101111
- b. 11000001 10000011 00011011 11111111

Solution

We replace each group of 8 bits with its equivalent decimal number (see Appendix B) and add dots for separation.

- a. 129.11.11.239
- b. 193.131.27.255

Example 19.2

Change the following IPv4 addresses from dotted-decimal notation to binary notation.

- a. 111.56.45.78
- b. 221.34.7.82

Solution

We replace each decimal number with its binary equivalent (see Appendix B).

- a. 01101111 00111000 00101101 01001110
- b. 11011101 00100010 00000111 01010010

Example 19.3

Find the error, if any, in the following IPv4 addresses.

- a. 111.56.045.78
- b. 221.34.7.8.20
- c. 75.45.301.14
- d. 11100010.23.14.67

Solution

- a. There must be no leading zero (045).
- b. There can be no more than four numbers in an IPv4 address.
- c. Each number needs to be less than or equal to 255 (301 is outside this range).
- d. A mixture of binary notation and dotted-decimal notation is not allowed.

IPv4 Address Rules (Reminder)

An IPv4 address:

- Consists of four decimal numbers (octets) separated by dots.
- Each octet must be between 0 and 255.
- Uses only dotted-decimal notation.
- No leading zeros are allowed in any octet.

a. 111.56.045.78

Error: Leading zero in one octet.

Explanation:

The octet 045 has a leading zero. IPv4 addresses do not allow leading zeros because they can cause ambiguity.

Correct form would be: 111.56.45.78

b. 221.34.7.8.20

Error: Too many octets.

Explanation:

An IPv4 address must contain exactly four numbers, but this address has five. Therefore, it is invalid.

Figure 19.2 Finding the classes in binary and dotted-decimal notation

| | First byte | Second byte | Third byte | Fourth byte |
|---------|---------------|----------------|---------------|----------------|
| Class A | 0 | | | |
| Class B | 10 | | | |
| Class C | 110 | | | |
| Class D | 1110 | | | |
| Class E | 1111 | | | |

a. Binary notation

| | First byte | Second byte | Third byte | Fourth byte |
|---------|---------------|----------------|---------------|----------------|
| Class A | 0-127 | | | |
| Class B | 128-191 | | | |
| Class C | 192-223 | | | |
| Class D | 224-239 | | | |
| Class E | 240-255 | | | |

b. Dotted-decimal notation

See b

Example 19.4

Find the class of each address.

- a. 00000001 00001011 00001011 11101111
- b. 11000001 10000011 00011011 11111111
- c. 14.23.120.8
- d. 252.5.15.111

Solution

- a. The first bit is 0. This is a class A address.
- b. The first 2 bits are 1; the third bit is 0. This is a class C address.
- c. The first byte is 14 (between 0 and 127); the class is A.
- d. The first byte is 252 (between 240 and 255); the class is E.

class A 0-127
class B 128-191
class C 192-223
class D 224-239
class E 240-255

c. 75.45.301.14

Error: Octet value out of range.

Explanation:

Each octet must be between 0 and 255.

The value 301 exceeds this limit, making the address invalid.

d. 11100010.23.14.67

Error: Mixing number formats.

Explanation:

The first octet 11100010 is written in binary, while the rest are in decimal.

IPv4 addresses must use only dotted-decimal notation throughout.

Mixing binary and decimal formats is not allowed.

Example 19.5

Figure 19.3 shows a block of addresses, in both binary and dotted-decimal notation, granted to a small business that needs 16 addresses.

Figure 19.3 A block of 16 addresses granted to a small organization

| | Block | Block |
|-------|------------------------------|--|
| First | 205.16.37.32 205.16.37.33 | 11001101 00010000 00100101 00100000 11001101 00010000 00100101 00100001 |
| Last | 205.16.37.47 | 11001101 00010000 00100101 00101111 |
| | a. Decimal | b. Binary |

16 Addresses

Example 19.6

A block of addresses is granted to a small organization. We know that one of the addresses is 205.16.37.39/28. What is the first address in the block?

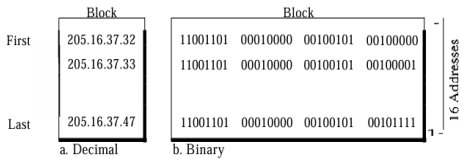
Solution

The binary representation of the given address is 11001101 00010000 00100101 00100111. If we set 32 - 28 rightmost bits to 0, we get 11001101 000100000100101 0010000 or 205.16.37.32. This is actually the block shown in Figure 19.3.

Last Address The last address in the block can be found by setting the 32 - *n* rightmost bits in the binary notation of the address to 1s.

The last address in the block can be found by setting the rightmost 32 - *n* bits to 1s.

Figure 19.3 A block of 16 addresses granted to a small organization



Example 19.7

Find the last address for the block in Example 19.6.

Solution

The binary representation of the given address is 11001101 000100000010010100100111. If we set 32 - 28 rightmost bits to 1, we get 11001101 00010000 0010010100101111 or 205.16.37.47. This is actually the block shown in Figure 19.3.

Number of Addresses The number of addresses in the block is the difference between the last and first address. It can easily be found using the formula 2^{32-n} .

The number of addresses in the block can be found by using the formula 2^{32-n} .

Example 19.8

Find the number of addresses in Example 19.6.

Solution

The value of *n* is 28, which means that number of addresses is 2^{32-28} or 16.

Example 19.9

Another way to find the first address, the last address, and the number of addresses is to represent the mask as a 32-bit binary (or 8-digit hexadecimal) number. This is particularly useful when we are writing a program to find these pieces of information. In Example 19.5 the /28 can be represented as 11111111 11111111 11111111 11110000 (twenty-eight 1s and four 0s). Find

- a. The first address
- h. The last address
- c. The number of addresses

Solution

- a. The first address can be found by ANDing the given addresses with the mask. ANDing here is done bit by bit. The result of ANDing 2 bits is 1 if both bits are 1s; the result is 0 otherwise.

Address: 11001101 00010000 00100101 00100111
Mask: 11111111 11111111 11111111 11110000
First address: 11001101 00010000 00100101 00100000

- b. The last address can be found by ORing the given addresses with the complement of the mask. ORing here is done bit by bit. The result of ORing 2 bits is 0 if both bits are 0s; the result is 1 otherwise. The complement of a number is found by changing each 1 to 0 and each 0 to 1.

Address: 11001101 00010000 00100101 00100111
Mask complement: 00000000 00000000 00000000 00001111
Last address: 11001101 00010000 00100101 00101111

- c. The number of addresses can be found by complementing the mask, interpreting it as a decimal number, and adding 1 to it.

Mask complement: 00000000 00000000 00000000 00001111
Number of addresses: 15 + 1 = 16

Key Concept (Exam Tip)

- **First address:** all host bits = 0
- **Last address:** all host bits = 1
- In a /28 block, addresses increase in steps of 16

Example 19.9

Another way to find the first address, last address, and number of addresses is to represent the subnet mask as a 32-bit binary number.

For /28, the mask is:

11111111 11111111 11111111 11100000

Copy code

(28 ones and 4 zeros)

Given address:

205.16.37.39

Copy code

Binary form of the address:

11001101 00010000 00100101 00100111

Copy code

c. Finding the number of addresses

The number of addresses is found by:

- 1. Taking the mask complement
- 2. Converting it to decimal
- 3. Adding 1

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Copy code

Mask complement (last octet): 00001111 = 15
15 + 1 = 16

Number of addresses = 16

a. Finding the first address

The first address is obtained by ANDing the address with the mask.

AND operation rule:

- 1 AND 1 = 1
- Otherwise = 0

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Copy code

Address: 11001101 00010000 00100101 00100111
Mask: 11111111 11111111 11111111 11100000

First address:11001101 00010000 00100101 00100000

Convert back to decimal:

205.16.37.32

Copy code

First address = 205.16.37.32

b. Finding the last address

The last address is found by ORing the address with the complement of the mask.

Mask complement

(Change 1 → 0 and 0 → 1)

less

Mask: 11111111 11111111 11111111 11100000
Mask complement: 00000000 00000000 00000000 00001111

OR operation rule:

- 0 OR 0 = 0
- Otherwise = 1

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Address: 11001101 00010000 00100101 00100111
Mask complement:00000000 00000000 00000000 00001111

Last address: 11001101 00010000 00100101 00101111

Convert back to decimal:

205.16.37.47

↓

Example 19.6

A block of IPv4 addresses is granted to a small organization. One address in the block is:

205.16.37.39/28

Find the first address in the block.

Step 1: Understand the prefix /28

- /28 means 28 bits are for the network.
- Remaining bits = $32 - 28 = 4$ bits for host addresses.
- A /28 block contains $2^4 = 16$ addresses.

Step 2: Convert the given IP address to binary

| Decimal | Binary |
|---------|----------|
| 205 | 11001101 |
| 16 | 00010000 |
| 37 | 00100101 |
| 39 | 00100111 |

So the full binary form is:

11001101 00010000 00100101 00100111

Example 19.7

Find the last address for the block in Example 19.6.

Given address: 205.16.37.39/28

Step 1: Understand the prefix /28

- /28 means 28 network bits
- Host bits = $32 - 28 = 4$
- Total addresses in the block = $2^4 = 16$

Step 2: Convert the given IP address to binary

| Decimal | Binary |
|---------|----------|
| 205 | 11001101 |
| 16 | 00010000 |
| 37 | 00100101 |
| 39 | 00100111 |

Binary representation:

11001101 00010000 00100101 00100111

Step 3: Set the host bits to 1

To find the last address, set the last 4 bits (host bits) to 1:

11001101 00010000 00100101 00101111

Step 3: Set the host bits to 0

Since /28, the last 4 bits are host bits.

Setting the last 4 bits to 0:

11001101 00010000 00100101 00100000

Step 4: Convert back to decimal

| Binary | Decimal |
|----------|---------|
| 11001101 | 205 |
| 00010000 | 16 |
| 00100101 | 37 |
| 00101111 | 47 |

✅ Final Answer

The last address in the block is:

205.16.37.47

Step 4: Convert back to decimal

| Binary | Decimal |
|----------|---------|
| 11001101 | 205 |
| 00010000 | 16 |
| 00100101 | 37 |
| 00100000 | 32 |

✅ Final Answer

The first address in the block is:

205.16.37.32