

# Data Communication & Networks

## Chapter 3: Data Link Layer

# Data Link Layer

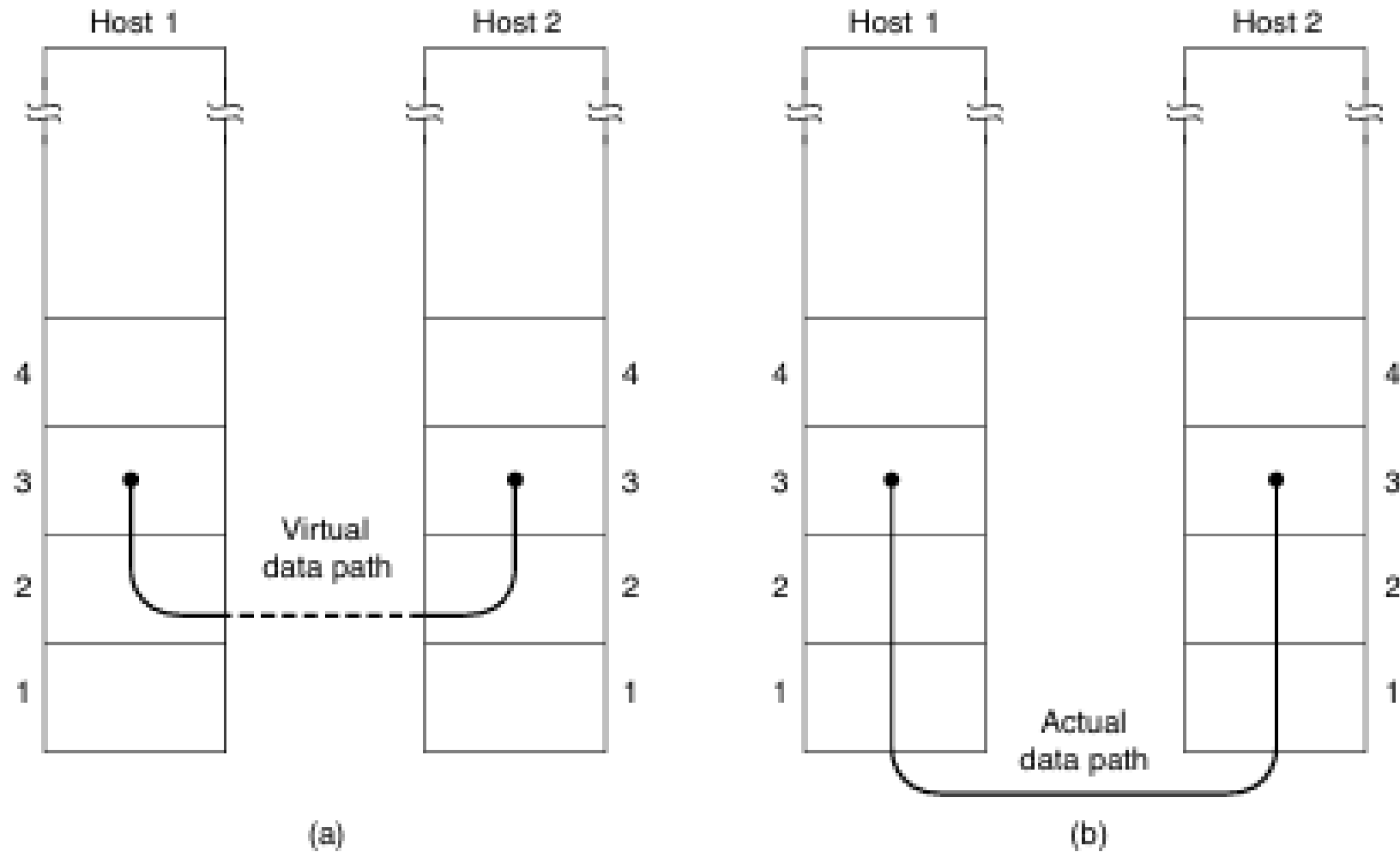
## Data Link Layer Design Issues

- Network layer services
- Framing
- Error control
- Flow control

# Data link Layer

- Physical layer delivers bits of information to and from data link layer. The functions of Data Link Layer are:
  1. Providing a well-defined service interface to the network layer.
  2. Dealing with transmission errors.
  3. Regulating the flow of data so that slow receivers are not swamped by fast senders.
- Data Link layer
  - Takes the packets from Physical layer, and
  - Encapsulates them into **frames**

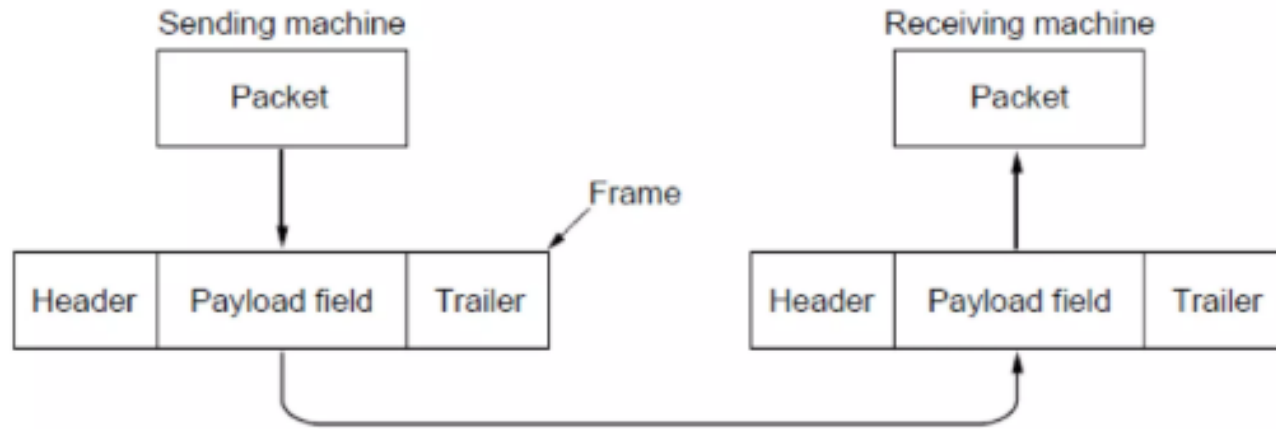
# Services Provided to Network Layer



**Figure 3-2.** (a) Virtual communication. (b) Actual communication.

# Packets and Frames

## Packets and Frames



Relationship between packets and frames.

- Header: Source address/ Destination address
- Trailer: Error detecting/correcting codes

# Services Offered

## Possible Services Offered



1. Unacknowledged connectionless service.
2. Acknowledged connectionless service.
3. Acknowledged connection-oriented service.

# Services Offered

## Unacknowledged Connectionless Service

- It consists of having the source machine send independent frames to the destination machine without having the destination machine acknowledge them.
- Example: Ethernet, Voice over IP, etc. in all the communication channel where real time operation is more important than quality of transmission.

# Services Offered

## Acknowledged Connectionless Service

- Each frame send by the Data Link layer is acknowledged and the sender knows if a specific frame has been received or lost.
- Typically the protocol uses a specific time period that if has passed without getting acknowledgment it will re-send the frame.
- This service is useful for commutation when an unreliable channel is being utilized (e.g., 802.11 WiFi).
- Network layer does not know frame size of the packets and other restriction of the data link layer. Hence it becomes necessary for data link layer to have some mechanism to optimize the transmission.



# Services Offered

## Acknowledged Connection Oriented Service

- Source and Destination establish a connection first.
- Each frame sent is numbered
  - Data link layer guarantees that each frame sent is indeed received.
  - It guarantees that each frame is received only once and that all frames are received in the correct order.
- Examples:
  - Satellite channel communication,
  - Long-distance telephone communication, etc.

# Framing

## Framing

- Transmission of the data link layer starts with breaking up the bit stream
  - into discrete frames
  - Computation of a checksum for each frame, and
  - Include the checksum into the frame before it is transmitted.
- Receiver computes its checksum error for a receiving frame and if it is different from the checksum that is being transmitted will have to deal with the error.
- Framing is more difficult than one could think!

# Framing

## Framing Methods

1. Byte count.
2. Flag bytes with byte stuffing.
3. Flag bits with bit stuffing.
4. Physical layer coding violations.

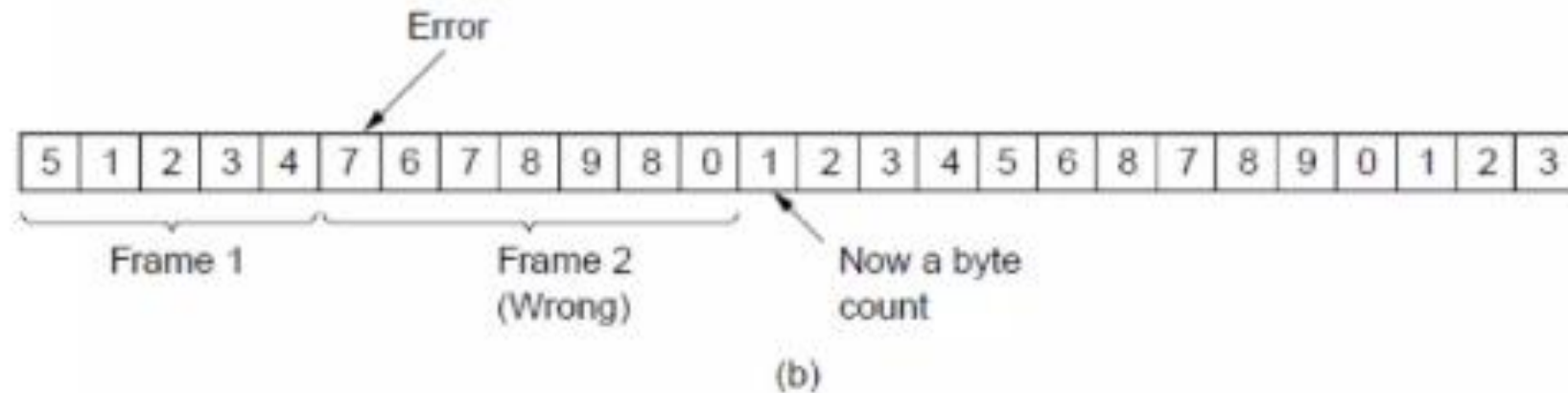
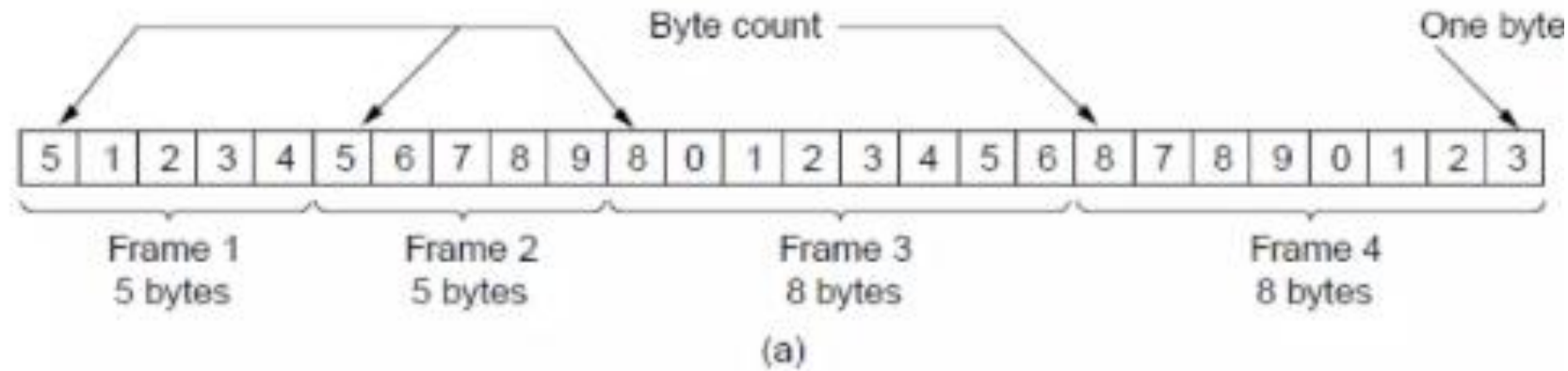
# Framing Method 1

## Byte Count Framing Method



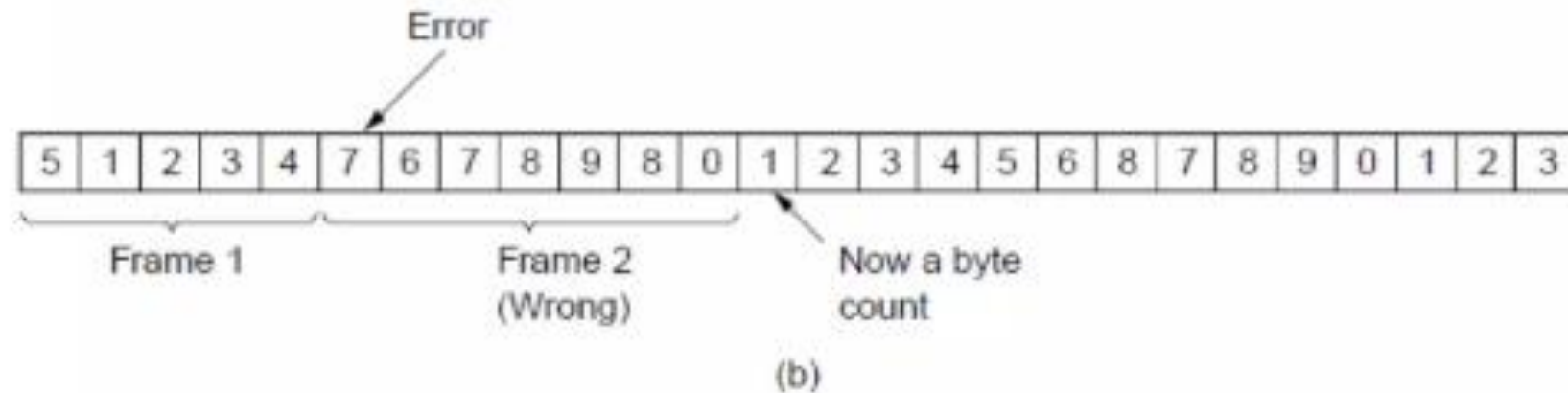
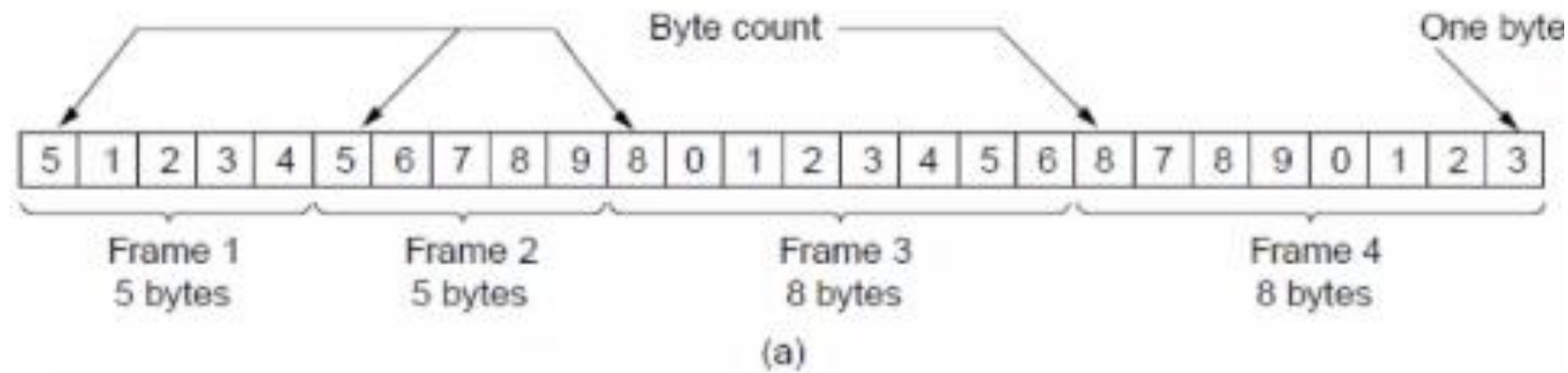
- It uses a field in the header to specify the number of bytes in the frame.
- Once the header information is being received it will be used to determine end of the frame.
- See figure in the next slide:
- Trouble with this algorithm is that when the count is incorrectly received the destination will get out of synch with transmission.
  - Destination may be able to detect that the frame is in error but it does not have a means (in this algorithm) how to correct it.

# Framing Method 1



A byte stream. (a) Without errors. (b) With one error.

# Framing Method 1



A byte stream. (a) Without errors. (b) With one error.



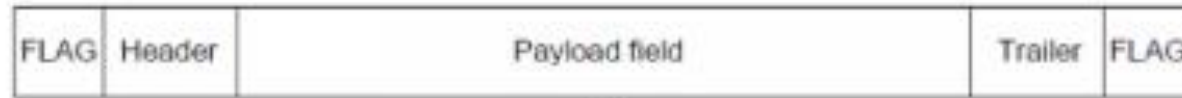
# Framing Method 2

## Flag Bytes with Byte Staffing Framing Method

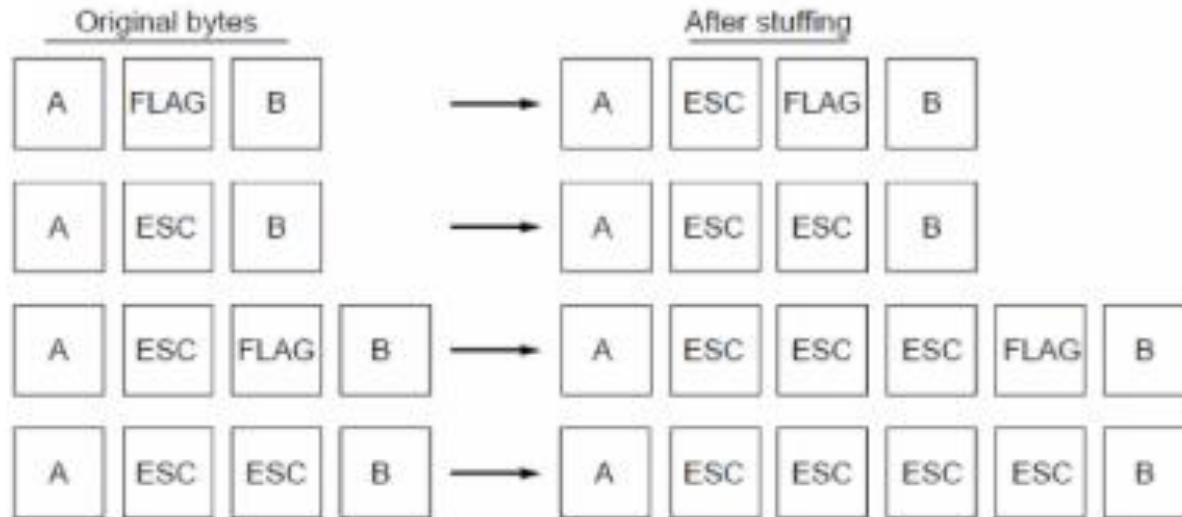
- This method gets around the boundary detection of the frame by having each appended by the frame start and frame end special bytes.
- If they are the same (beginning and ending byte in the frame) they are called **flag byte**.
- In the next slide figure this byte is shown as FLAG.
- If the actual data contains a byte that is identical to the FLAG byte (e.g., picture, data stream, etc.) the convention that can be used is to have escape character inserted just before the "FLAG" character.

# Framing Method 2

## Framing (2)



(a)



(b)

- a) A frame delimited by flag bytes.
- b) Four examples of byte sequences before and after byte stuffing.



# Bit Stuffing

- This method achieves the same thing as Byte Stuffing method by using Bits (1) instead of Bytes (8 Bits).
- It was developed for High-level Data Link Control (HDLC) protocol.
- Each frame begins and ends with a special bit pattern:
  - 01111110 or 0x7E <- Flag Byte
  - Whenever the sender's data link layer encounters five consecutive 1s in the data it automatically stuffs a 0 bit into the outgoing bit stream.
  - USB uses bit stuffing.

# Framing Method 3: Bit Stuffing

## Framing (3)

(a) 0 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 0 0 1 0

(b) 0 1 1 0 1 1 1 1 1 0 1 1 1 1 0 1 1 1 1 0 1 0 0 1 0

Stuffed bits

(c) 0 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 0 0 1 0

# Data Link Layer

## Data Link Layer Design Issues

- Network layer services ✓
- Framing ✓
- Error control
- Flow control

# Error Control: Error Detection and Error Correction

- Error Detection
- Error correction

# Error Detection: Error Detecting Codes

Linear, systematic block codes

1. Parity.
2. Checksums.
3. Cyclic Redundancy Checks (CRCs).

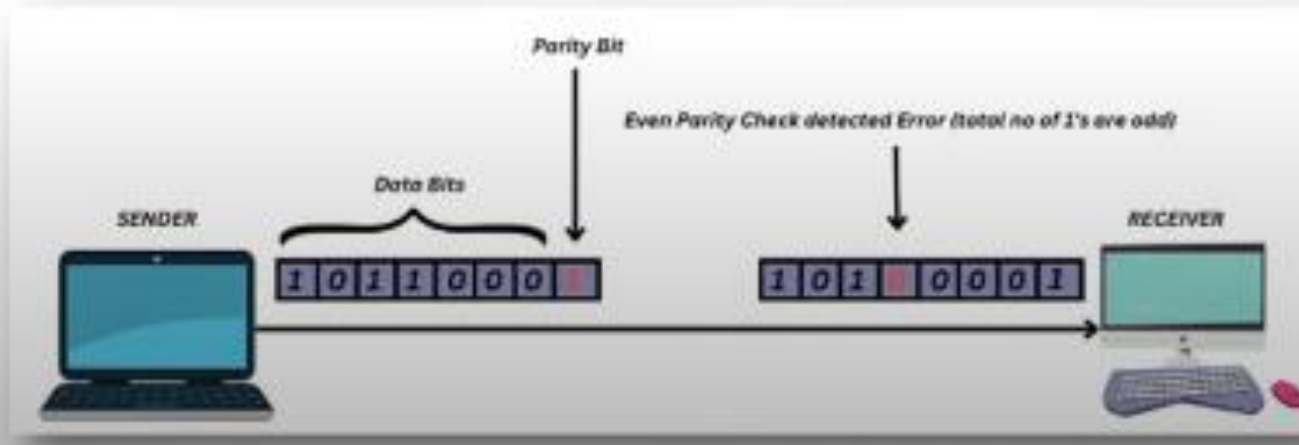
# Error Detection: Method 1 (Parity Bits)

- The parity bit is extra bit added to the message from sender side, to help in error detection at the receiver side.
- Also called as Data-word bits or Codeword.

1. **Even Parity:** If number of 1s is even in data bits, parity bit value is 0.

2. **Odd Parity:** If number of 1s is odd in data bits, parity bit value is 1.

Data Bits							Parity Bits
1	0	1	0	1	0	1	0
1	0	1	0	0	0	1	1



# Error Detection: Method 2 (Checksum)

- It detects all errors involving an odd number of bits as well even number of bits.
- In checksum, at each word is transmitted it is added to previously send words & sum is retained at transmitter.

Sender



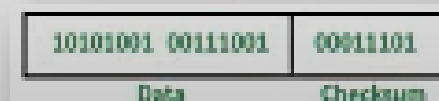
Subunit 1 : 1 0 1 0 1 0 0 1  
Subunit 2 : 0 0 1 1 1 0 0 1  
Sum : 1 1 1 0 0 0 1 0  
Checksum : 0 0 0 1 1 1 0 1

(Compliment of 1's)

Receiver



Subunit 1 : 1 0 1 0 1 0 0 1  
Subunit 2 : 0 0 1 1 1 0 0 1  
Sum : 1 1 1 0 0 0 1 0  
Checksum : 0 0 0 1 1 1 0 1

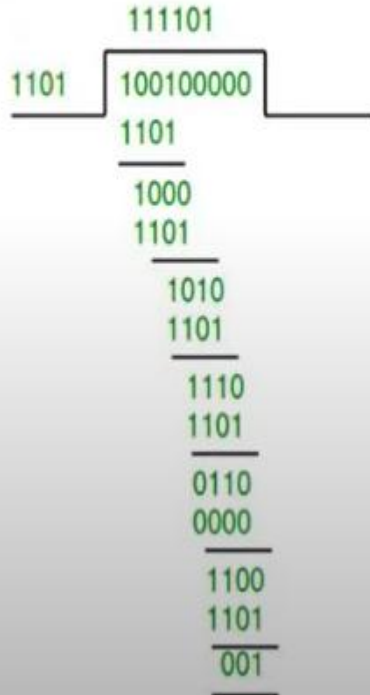


# Error Detection: Method 3 (Cyclic Redundancy Check)

- CRC is a method of detecting accidental changes/errors in the communication channel.
- CRC uses Generator Polynomial which is available on both sender and receiver side.

## Example:

Sender



Data Words sent : 1 0 0 1 0 0

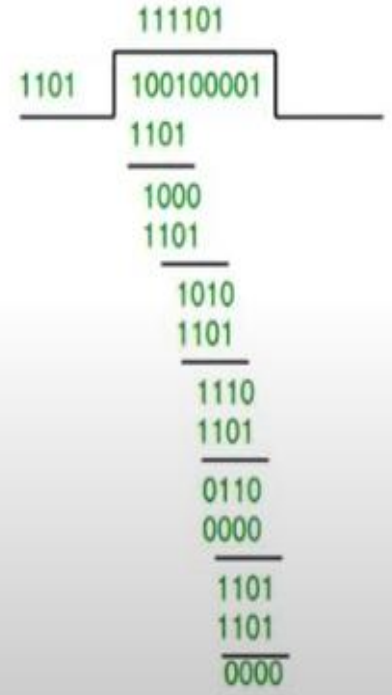
Key: 1 1 0 1

Reminder : 0 0 1

Data Sent : 1 0 0 1 0 0 0 0 1



Receiver



Data Words sent : 1 0 0 1 0 0

Key: 1 1 0 1

Reminder : 0 0 1

Data Sent : 1 0 0 1 0 0 0 0 1



# Error Detection & Error Correction: Hamming Code

- Hamming code Done in Lecture. See Lecture Notes