

Looping 100,000 times

```
COUNT1 EQU 0x20      ; Outer loop counter
COUNT2 EQU 0x21      ; Middle loop counter
COUNT3 EQU 0x22      ; Inner loop counter

        ORG 0x00
        GOTO MAIN

MAIN:
        MOVLW 0x55          ; Initialize PORTB = 55H
        MOVWF PORTB

        MOVLW d'10'          ; Load outer loop count = 10
        MOVWF COUNT1

LOOP1:
        MOVLW d'100'         ; Load middle loop count = 100
        MOVWF COUNT2

LOOP2:
        MOVLW d'100'         ; Load inner loop count = 100
        MOVWF COUNT3

LOOP3:
        COMF PORTB, F        ; Toggle PORTB
        DECFSZ COUNT3, F     ; Inner loop counter --
        GOTO LOOP3

        DECFSZ COUNT2, F     ; Middle loop counter --
        GOTO LOOP2

        DECFSZ COUNT1, F     ; Outer loop counter --
        GOTO LOOP1

        END
```

Low byte & High byte

```
L_BYTE EQU 0x05      ; Low byte of sum
H_BYTE EQU 0x06      ; High byte of sum

        ORG 0x00
        CLRF L_BYTE          ; Clear low byte
        CLRF H_BYTE          ; Clear high byte

        MOVLW 0x4C            ; W = 4Ch
        MOVWF 0x40             ; Store at 0x40
```

```

    MOVLW 0x3F      ; W = 3Fh
    MOVWF 0x41      ; Store at 0x41

    MOVLW 0xD3      ; W = D3h
    MOVWF 0x42      ; Store at 0x42

    MOVLW 0xA1      ; W = A1h
    MOVWF 0x43      ; Store at 0x43

        ; --- Add first number (0x40) ---
    MOVF 0x40, W
    ADDWF L_BYTE, F
    BNC N1
    INCF H_BYTE, F
N1:
        ; --- Add second number (0x41) ---
    MOVF 0x41, W
    ADDWF L_BYTE, F
    BNC N2
    INCF H_BYTE, F
N2:
        ; --- Add third number (0x42) ---
    MOVF 0x42, W
    ADDWF L_BYTE, F
    BNC N3
    INCF H_BYTE, F
N3:
        ; --- Add fourth number (0x43) ---
    MOVF 0x43, W
    ADDWF L_BYTE, F
    BNC OVER
    INCF H_BYTE, F
OVER:
    END

```

Set to swap even and odd bits of 56H and store in new file register

```
SRC    EQU 0x56
DST    EQU 0x57
TEMP1  EQU 0x20
TEMP2  EQU 0x21

ORG 0x00

; --- Extract even bits (mask = 01010101) ---
MOVF   SRC, W
ANDLW  0x55
MOVWF  TEMP1
RLF    TEMP1, F ; shift left → even bits move to odd positions

; --- Extract odd bits (mask = 10101010) ---
MOVF   SRC, W
ANDLW  0xAA
MOVWF  TEMP2
RRF    TEMP2, F ; shift right → odd bits move to even positions

; --- Combine ---
MOVF   TEMP1, W
IORWF  TEMP2, W
MOVWF  DST

END
```

a) Role of Stack and subroutines in PIC. What happens when the Stack overflows?

- **Stack role:**
 - The stack stores the return address whenever a subroutine (CALL) or interrupt is executed.
 - When the subroutine finishes (RETURN), the stored address is popped back so the program continues from the correct location.
- **Subroutines role:**
 - Subroutines allow code reuse by grouping instructions into callable blocks.
 - Helps reduce program size and improves modularity.
- **Stack overflow:**
 - If the stack becomes full and another CALL/interrupt occurs, it overwrites old return addresses.
 - This causes the program to return to the wrong place → **program crashes or behaves unpredictably.**

b) CALL vs RCALL for PIC18 with 4K program memory

- **CALL:**
 - Absolute call to any address in the entire program memory (useful for large memory space).
- **RCALL:**
 - Relative call (address specified as an offset from current location).
 - More compact instruction (uses fewer program words).
 - Suitable for small memory devices.

Since the chip has **only 4K program memory**, **RCALL** is more useful because it saves memory and is efficient for short-range subroutine calls.

Final short answers:

- a)** Stack stores return addresses for subroutines/interrupts. Overflow → overwriting → wrong return address → program error.
- b)** For a 4K ROM PIC18, **RCALL** is preferred (saves memory, efficient for small program space).