Photo and Video Equipment Rental DBMS Final Report

12/03/2020

Introduction

In recent times, photo and video equipment rentals have seen significant growth leading to an increasingly competitive market. The focus of businesses in this market is to help customers find ideal equipment for a variety of use cases and provide them with a rental service for these products. One of the biggest challenges facing this approach is the requirement of a sophisticated database system needed to properly store and retrieve data in a secure and efficient manner. To accomplish this goal, a variety of entities are needed in order to properly track the distribution in various ways. This includes keeping a database of the client-rental relationship, the employee-customer relationship, as well as the company's product-distribution relationship.

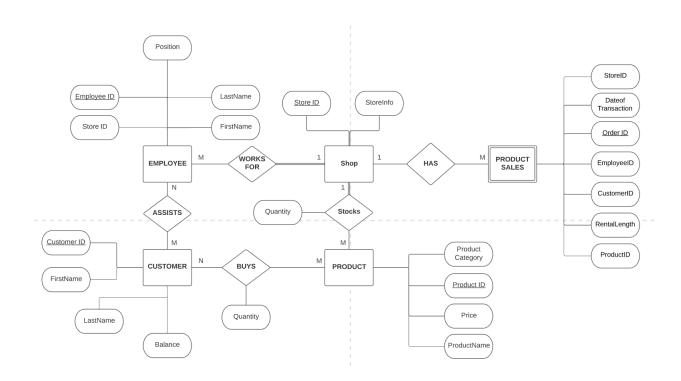
The first category that we will discuss is the side of the transaction that focuses on the customer. This category makes use of a few separate entities, mainly being the customer, product, and employee. For the customer, a unique primary key has been assigned alongside personal information, rental history, and balance owed attributes. This allows us to maintain a comprehensive understanding of each individual customer and their history. The goal of this relationship is to streamline the process for the customer. Tracking rental history allows us to easily rebook an order that was previously made. Tracking balance owed can also help build a positive relationship as a notification system can be implemented later on to help the customer keep track of their contracts and payments.

The second category will focus on the employee and customer relationship. This section makes use of the previously three mentioned entities, and adds another in the form of a product sale entity. The focus of this entity is to keep track of a unique order ID, which will include important information like the date of transaction, the customer ID, the employee who put the transaction through, as well as things like the length of the contract and the rental rate that was paid. This relationship allows the business to build an organized database of transactions which will be useful towards keeping track of a variety of attributes like sales volume. Another benefit is that if a customer has a problem, they will be able to be referred to the same employee that they made the transaction with, streamlining the process for both the business and the customer.

The final category is the product-distribution relationship. This category has a large breadth and can allow for use of new attributes for things like separate franchises, brands or manufacturers, product categories, as well as individual products themselves. The goal of this section is to be able to closely monitor distribution across various locations. This distribution includes what products each location will stock, as well as what products are currently in stock or sold out. This data is especially useful as it allows stores to monitor and adjust stock levels to better suit the needs of the individual customers in varying demographics. Manufacturers of the product will also be able to add new products, categories, and even make changes to the rates that they charge for rentals. As for the products themselves, essential attribute information is provided such as specifications, rental rate, and which stores currently have it in stock.

This highlights the general structure that we have followed when constructing the application. Following this we will present the final ER model, which has been structured to implement the features as highlighted above.

ER DIAGRAM DESIGN



RELATIONS & TABLES WITH DUMMY DATA

Shop

StoreID → StoreInformation

Employee

 $\begin{aligned} & \text{EmployeeID} \rightarrow \text{FirstName} \\ & \text{EmployeeID} \rightarrow \text{LastName} \\ & \text{EmployeeID} \rightarrow \text{Position} \end{aligned}$

 $\text{EmployeeID} \to \text{StoreID}$

Customer

Examples:

1	The Dundas Bay Location		
2	The Yorkdale location		

	♦ FIRSTNAME			
1	Samuel	Khan	Manager	1
2	Bolu	Zhang	Cashier	1
3	Carol	Sinclair	Cashier	1

1	724.96	Jonathan	(null)
2	765.97	Emily	Brown
3	815.93	Jonas	Rei

CustomerID → Balance

CustomerID → FirstName

CustomerID → LastName

Product

ProductID → ProductCategory

ProductID → ProductBrand

ProductID → ProductName

ProductID → Price

♦ PRODUCTID				♦ PRICE
1	Camera	Gucci	500D	133
2	Camera	Canon	The capturer	500
3	Camera	Nikon	the visualizer	450
4	Audio	Blue Yeti	Midnight blue microphone	75

3

1

2

3

4

1

1

1

1

ProductSale

OrderID → CustomerID

OrderID → ProductID

OrderID → DateOfTransaction

OrderID → StoreID	3 20-09-15	2	1	
OrderID → RentalLength				
OrderID \rightarrow EmployeeID				

1 20-09-15

2 20-09-16

Stocks

StoreID, ProductID → Quantity

Buys

CustomerID, ProductID → Quantity

	\$1	PRODUCTID	
1		1	5
1		2	15
1		3	12
	ID	♦ PRODUCTIE	0 ⊕ QUANTITY
	1		2 1
	1		4 3
	2		1 2

CROSS-TABLE RELATIONS

Employee-Shop

 $\mathsf{EmployeeID} \to \mathsf{StoreID}$

ProductSale-Product

OrderID → ProductID

ProductSale-Shop

ProductSale-Employee

OrderID → EmployeeID

ProductSale-Customer

OrderID → CustomerID

Product&Shop-Stocks

StoreID, ProductID → Stocks. Quantity

Product&Customer-Buys

CustomerID, ProductID → Buys. Quantity

Normalization of the below relation to BCNF using Bernstein's Algorithm

Our table was already in BCNF, so we have added tables 2 and 3 as shown below to our database in order to visualize the process of normalization.

R(ProductID,ProductCategory,ProductBrand,ProductName,Price,StoreID,Stocks,Buys,CustomerID)

```
FD's{ ProductID → ProductCategory, ProductBrand, ProductName, Price
                                                                                                                     1
         ProductName → ProductID, Price
                                                                                                                    2
         ProductID, ProductName → ProductCategory
                                                                                                                     3
         StoreID, ProductID → Stocks
                                                                                                                     4
         CustomerID, ProductID \rightarrow Buys }
                                                                                                                     5
FD = \{I \rightarrow CBNP, N \rightarrow IP, IN \rightarrow C, SiI \rightarrow S, CiI \rightarrow Bu\}
FD = \{I \rightarrow C, I \rightarrow B, I \rightarrow N, I \rightarrow P, N \rightarrow I, N \rightarrow P, IN \rightarrow C, Si\&I \rightarrow S, Ci&I \rightarrow Bu\}
I \rightarrow C: I+ = \{I,B,N,P\}
                                      We do not get C, so not redundant
                                      We do not get B, so not redundant
I \rightarrow B : I + = \{I,C,N,P\}
I \rightarrow N : I+ = \{I,C,B,P\}
                                      We do not get N, so not redundant
I \rightarrow P : I + = \{I,C,N,B,P\}
                                      We get P, transitive
```

 $N \rightarrow I$: $N+ = \{N,P,C,B\}$ We do not get N, so not redundant

 $N \rightarrow P: N+ = \{N,I,C,B,P\}$ We get P, transitive

 $IN \rightarrow C : IN+ = \{I,N,C\}$ We get C, so redundant

 $Si,I \rightarrow S: Si+= \{Si,I,S\}$ We do not get S, so not redundant $Ci,I \rightarrow Bu: Ci+= \{Ci,I,Bu\}$ We do not get Bu, so not redundant

2 is transitive and redundant(ProductID → ProductName , ProductName → ProductID) 3 is redundant

After removing redundancies,

$$FD = \{I \rightarrow C, I \rightarrow B, I \rightarrow N, I \rightarrow P, Si\&I \rightarrow S, Ci\&I \rightarrow Bu\}$$

All of the TRs are fully dependent, no partial dependencies.

 $ISiCi+ = \{I,Si,Ci,C,B,N,P\}$ this is the key.

R1(<u>ProductID</u>,ProductCategory,ProductBrand,ProductName,Price)

R2(ProductID, StoreID, Stocks)

R3(ProductID, CustomerID, Buys)

By looking at the updated FDs, we can see that the tables are in 3NF/BCNF.

3NF example from Assign 7:

Step 1: all tables are in the form where each cell contains a single value and each record is unique.

Step 2: all tables have a single column primary key in the form of a unique ID number.

Step 3: our table has no transitive functional dependencies.

For an example, we have added a transitive functional dependency to our Customer table

			♦ PERSONALINFORMATION
1	Mr.	724.96	Jonathan
2	Mrs	765.97	Emily Brown
3	Mr.	815.93	Jonas Rei

Here it has been split into a foreign key so it is 3NF

		BALANCE	♦ PERSONALINFORMATION
1	. 1	724.96	Jonathan
2	2	765.97	Emily Brown
3	1	815.93	Jonas Rei
Mr.	1		
Mrs.	2		

LIST OF RELATIONAL ALGEBRAS FOR ALL QUERIES

Relations

Shop(<u>StoreID</u>,storeInformation)

 $\textbf{Employee}(\underline{\text{EmployeeID}}, FirstName, LastName, Position, StoreID)$

Customer(CustomerID, Balance, FirstName, LastName)

Product(<u>ProductID</u>,ProductCategory,ProductBrand,ProductName,Price)

 $\label{lem:productSale} \textbf{ProductSale}(\underline{OrderID}, DateOfTransaction, RentalLength, EmployeeID, CustomerID, ProductID, StoreID)$

Buys(CustomerID, ProductID, Quantity)

Stocks(<u>StoreID</u>,<u>ProductID</u>,Quantity)

Queries and their RA

--Selects distinct dates where products were sold

SELECT

DISTINCT dateoftransaction

FROM ProductSale

ORDER BY dateoftransaction ASC

 $\Pi_{date of transaction} Products ale$

--Shows most popular products sold

SELECT productid, COUNT(*) AS "Times Sold" FROM productsale GROUP BY productid HAVING COUNT(*) > 1 ORDER BY COUNT(*) DESC

 $\Pi_{\textit{productid}, \textit{count}(^*)}(\sigma_{\textit{count}(^*)>1} \rho_{\textit{"Times Sold"/count}(^*)} \textit{Productsale})$

--Shows the orders fulfilled by an employee on specified dates

SELECT orderid, dateoftransaction

FROM productsale

WHERE employeeid = 003

AND (dateoftransaction = '2020-09-16' OR dateoftransaction = '2020-09-15')

 $\Pi_{orderid,\, date of transaction}(\sigma_{employeeid\,=\,003\,\Lambda\,(date of transaction\,=\,2020-09-16\,\text{V}\,\, date of transaction\,=\,2020-09-15}\,Properties and the properties of the prop$

--Selects all the managers

SELECT DISTINCT employeeid, position, storeid, FirstName, LastName FROM Employee

WHERE position = 'Manager'

ORDER BY storeid DESC

 $\Pi_{employeeid, position, storeid, firstname, lastname}$ ($\sigma_{position = 'manager'}$ Employeeid, position = 'manager'

--Shows how many customers are in the database

SELECT 'Unique Customers', count(customerid)

FROM Customer;

 $\Pi_{'Unique\ customers',\ count(customerid)} Customer$

--Selects products with a given budget and category

SELECT*

FROM Product

WHERE ProductCategory = 'Camera'

AND price < 250

 $\sigma_{product category = 'Camera' \ \Lambda \ price < 250} Product$

--Shows store exclusive items

SELECT productid

FROM stocks

WHERE storeid = 2

MINUS

SELECT productid

FROM stocks

WHERE storeid = 1

ORDER BY productid

$$\Pi_{productid} \left(\sigma_{storeid=2} Stocks \right) - \Pi_{productid} \left(\sigma_{storeid=1} Stocks \right)$$

--Joins the PRODUCT table with STOCKS table to show each product and the stock at each store

SELECT product.productid,product.ProductName,product.product.producteategory, product.productbrand,product.price,stocks.quantity,stocks.storeid FROM product
FULL JOIN stocks
ON product.productid = stocks.productid
ORDER BY storeid

Includes = product.productid,product.ProductName,product.productcategory,
product.productbrand,product.price,stocks.quantity,stocks.storeid

$$\Pi_{Includes}(\mathit{Product} \bowtie_{\mathit{product.productid} = \mathit{stocks.productid}} \mathit{Stocks})$$

--Joins the CUSTOMER table with BUYS table to display products purchased by each customer

 $SELECT\ customer. First Name, customer. Last Name, Customer. balance, buys. product id,$

buys.quantity

FROM Customer

LEFT JOIN buys

ON buys.customerid = customer.customerid

ORDER BY customer.customerid;

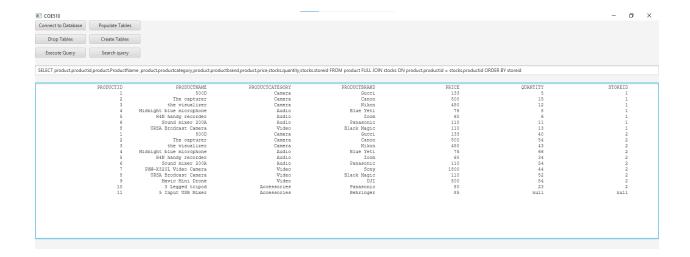
Includes = customer.FirstName,customer.LastName,Customer.balance,buys.productid,buys.quantity

GUI PROJECT

We made a java gui application that can connect to the application, drop tables, create tables, populate tables, and execute queries. A screenshot of the gui is shown below. First, the user must click "Connect to Database" to open the connection to the server. If no data is stored, the user can click "Drop Tables" followed by "Create Tables" followed by "Populate Tables" to initialize dummy data. From there, the user can type a search query or an execution query as desired in the text bar. The result will be displayed in the text field below.



The application is submitted as a jar file so it can be run on Windows OS, along with a zip file of the Netbeans project. Below is a screenshot of it executing an advanced query and displaying results in table form.



Product sale before deletion:



Deletion Query:



Product sale after deletion:



Closing Thoughts

In conclusion, the database achieved all of our design goals. Looking towards the future, more advanced systems could be layered on top of the database to provide increased functionality. For example, an email notification server could be useful to help customers keep on top of their rental due dates. Integration with a website could also be useful in order for customers to be able to search stock before coming to the store. Finally, more prepared statements and text fields could be useful for database functionality in order to streamline use for less experienced users.