Parallel and Distributed Computing CS3006

Lecture 17

MPI-IV

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- MPI provides its own optimized implementations for most of the collective operations that we performed in CH#4
- These operations are called collective as all the processes must have a call to collective functions
- Every collective operation take a communicator (such as MPI_COMM_WORLD) as argument
 - all the processes within that communicator must have a corresponding call to the operation

Barrier synchronization operation

The barrier synchronization operation is performed in MPI using:

```
int MPI_Barrier(MPI_Comm comm)
```

The call to the MPI_barrier returns only all the processes in the group have called this function

The one-to-all broadcast:

Buffer of the source process is copied to the buffers of other processes

The all-to-one reduction operation

- Dual of one-to-all broadcast
- Every process including target provides sendbuf for its value that is to be used for the reduction
- After the reduction, reduced value is stored in recybuf of target process
- Every process must also provide recybuf, though it may not be target of the reduction

```
int MPI_Reduce(void *sendbuf, void *recvbuf, int count,
    MPI_Datatype datatype, MPI_Op op, int target,
    MPI_Comm comm)
```

Here MPI_Op is MPI defined set of operations for reduction

The all-to-one reduction operation

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Operation	Meaning	Datatypes	
MPI_MAX	Maximum	C integers and floating point	
MPI_MIN	Minimum	C integers and floating point	
MPI_SUM	Sum	C integers and floating point	
MPI_PROD	Product	C integers and floating point	
MPI_LAND	Logical AND	C integers	
MPI_BAND	Bit-wise AND	C integers and byte	
MPI_LOR	Logical OR	C integers	
MPI_BOR	Bit-wise OR	C integers and byte	
MPI_LXOR	Logical XOR	C integers	
MPI_BXOR	Bit-wise XOR	C integers and byte	
MPI_MAXLOC	max-min value-location	Data-pairs	
MPI_MINLOC	min-min value-location	Data-pairs	

MPI_MAXLOC and MPI_MINLOC

- The operation MPI_MAXLOC combines pairs of values (v_i, l_i) and returns the pair (v, l) such that v is the maximum among all v_i 's and l is the corresponding l_i (if there are more than one, it is the smallest among all these l_i 's).
- lacktriangle MPI MINLOC does the same, except for minimum value of v_i .

An example use of the MPI_MINLOC and MPI_MAXLOC operators.

MPI_MAXLOC and MPI_MINLOC

MPI datatypes for data-pairs used with the MPI_MAXLOC and MPI MINLOC reduction operations.

MPI Datatype	C Datatype	
MPI_2INT	pair of ints	
MPI_SHORT_INT	short and int	
MPI_LONG_INT	long and int	
MPI_LONG_DOUBLE_INT	long double and int	
MPI_FLOAT_INT	float and int	
MPI_DOUBLE_INT	double and int	

The All-Reduce operation

- MPI_AllReduce is used when the result of the reduction operation is needed by all processes
- Equal to All-to-one reduction followed by one-to-all broadcast

- After Allreduce operation, recybuf of all the processes contain reduced value
- Note: no target for reduction is given

Prefix(scan) operation

- Recall 4.3-for prefix-sum: After the operation, every process has sum of the buffers of the previous processes and its own.
- MPI_Scan() is MPI primitive for the prefix operations.
- All the operators that can be used for reduction can also be used for the scan operation
- If buffer is an array of elements, then recvbuf is also an array containing element-wise prefix at each position.

Exclusive-Prefix (Exscan) operation

- Exclusive-prefix-sum: After the operation, every process has sum of the buffers of the previous processes excluding its own.
- MPI_Exscan() is MPI primitive for the exclusive-prefix operations.
- The recybuf of first process is remains unchanged as there is no process before it. Some MPI distributions place identity value for the given associative operator.

- Recall section 4.4: After the **Gather** operation, a single target process accumulates[concatenates] buffers of all the other processes without any reduction operator.
- Each process sends element(s) in its sendbuf to the target process.
- Total number of elements to be sent by each process must be same. This number is specified in sendcount and is equal to recvcount.
- On the target, recybuf stores elements sent by all the processes in rank order. Elements received at target by process#i, will be stored starting from (i*sendcount)th index of recybuf.

- Gatherv
 - Each process can have different message length.
 - Recvcounts[i] = Total elements to be received by ith processing node
 - Displs[i]= starting index in recybuf to store message received from ith process.

- Gathery (Displs calculation example)
 - Let each process have elements one more than their rank.
 - Then calculation of displs[] at target is calculated as:

PO	P1	P2	Р3
32	12, 15	4,9,14	20,23,27,31

Recvcounts	1	1	2	3	4
displs	0	2	<u>0+1=1</u>	1+2=3	3+3=6

- MPI_Allgather
- Same as All-to-All broadcast described in section 4.2
- Every process serve as target for the gather

- Note: No target for gather
- Unlike MPI_Gather, it gathers sendbufs of all the processes in recybufs of all the processes.

MPI_Gather and its variants

MPI_Allgatherv

- Here every process will have to supply the valid calculated arrays of recvcounts and displs.
- Furthermore, it is also necessary for all the processors to provide a recybuf [an array] of sufficient size to store all the elements of all the processes.

```
intial values at source::0::=83 86 77 15 93 35 86 92
rank=0: Received:83 86
rank=1: Received:77 15
rank=2: Received:93 35
rank=3: Received:86 92
```

MPI_scatter

Scatters data stored in sendbuf of source process between all the processes as discussed in ch#4.

Sendcount and recvcount should be the same and represent total elements to be given to each process.

```
intial values at source::0::=83 86
                                        77
                                                 15
                                                         93
                                                                 35
                                                                          86
                                                                                                  21
rank=0: Received:83
rank=1: Received:86
                        77
rank=2: Received:15
                        93
                                35
rank=3: Received:86
                        92
                                49
                                        21
```

MPI_scatterv

- Here sendcounts is an array of P size such that its ith index contains number of elements to be sent to ith process.
- displs[i] indicates the index in sendbuf from which sendcounts[i] values are to be sent to ith process.

- Values of sendbuf and sendcounts at all processes except the source are ignored but, you have to provide pointers though pointing to nothing.
- Every process will have to calculate its own recvcount
- Challenge: Write a program that scatters even rows of a 2P X P array to P processes. Assume that we are using 1d array to simulate 2D data.

MPI_AlltoAll

- This routine is used to perform operation known as alltoall personalized communication in CH#4.
- Each process has P messages, one for each process.
 - parallel matrix transpose operations.

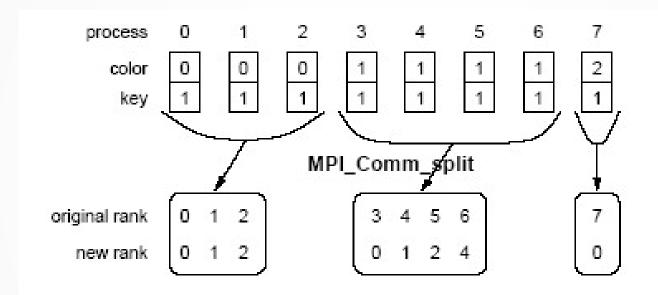
- Here sendbuf is of size p*message size for each process.
- Size of receive buffer is equal to sendbuf.
- Sendcount and recvcount have same integer value representing elements to be sent to each process and elements to be received from each process, respectively.
- Read and implement vector variant for alltoall personalized communication

Groups and Communicators

- In many parallel algorithms, communication operations need to be restricted to certain subsets of processes.
- MPI provides mechanisms for partitioning the group of processes that belong to a communicator into subgroups each corresponding to a different communicator.
- The simplest such mechanism is:

This operation groups processors by color and sorts resulting groups on the key (i.e., highest key→ highest new_rank).

Groups and Communicators



Using MPI_Comm_split to split a group of processes in a communicator into subgroups.

Questions



References

1. Kumar, V., Grama, A., Gupta, A., & Karypis, G. (2017). *Introduction to parallel computing*. Redwood City, CA: Benjamin/Cummings.