National University of Computer and Emerging Sciences, Lahore Campus

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Course Name:	Parallel and Distributing Computing	Course Code:	CS3006
Degree Program:	BS (CS)	Semester:	Spring 2023
Exam Duration:	60 Minutes	Total Marks:	30
Paper Date:	25/02/23	Weight	12.5
Exam Type:	Midterm I	Page(s):	4

Student : Name:_):	Roll No		Section:
Instruction:		Attempt all ques			be used but it should not be
		attached. If you	think some information is m	issing then assume	it and mention it clearly.
0	# 1.	[F 2 montes C O #	41		
<u>Ques</u>	stion # 1:	[5+2 marks, CLO #	寸		
From	the give	n options, select the	e best answer.		
i.	GPU is	an example of			
		Shared Memory SI			
		Distributed Memor	•		
		Shared Memory M			
	<u>d.</u>	_Distributed Memor	ry MIMD		
ii.			rk has switching node	es.	
		log p			
		$p/2 \times \log p$			
		$p \times log p$ p^2			
	<u>d.</u>	_P			
iii.	The	of a task	graph is usually a	of the	of the task-dependency graph:
	a.	edge-set; interaction	on; superset; nodes		
	b.	edge-set; mapping	; subset; edge-set		
	c.	edge-set; interaction	on; subset; edge-set		
	d.	edge-set; interacti	on; superset; edge-set		
iv.	The us	e of deco	mposition comes into play v	vhen we	
			t sure of the dependencies b		
		•	t sure of the size of the sear		ore
		•	ot sure of the size of the sea	•	
			avoid divide and conquer str		
V.	We car	n increase nerforma	nce of a uniprocessor by		
•.	a.	Increasing clock fre	-		
	b.	Increasing instructi	•		
	C.	Multithreading	, ,		
	d.	Any one from abou	ve can be used		
vi.		What is the maximu			uted in parallel and rest 30% in version of the program executing

a. 2.58b. 3.06

- c. 4.02
- d. 4.08

Question # 2: [08 marks, CLO # 1]

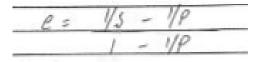
Assume a sequential program S has an execution time of S00 seconds. Further, assume that S_p is a parallel variant of S0. After an experimental evaluation over different number of processors, the following running times were achieved: -

Р	2	4	6	8
Execution Time (seconds)	294	203.84	176.62	164.92
Speedup				
Karp-Flatt Metric				

a) Calculate Speedups for each of the experimental configurations in the space provided below and then write your answers in the table above.

$$500/294 = 1.7006$$
 Speedup = T(S) / T(P) $500/203.84 = 2.4529$ $1/((F + (1-F)/p))$ $500/176.62 = 2.8309$ $500/164.92 = 3.0317$

b) Calculate the Karp-Flatt metric values in the space provided below and then write your answers in the table above. You also have to interpret the results of Karp-Flatt metric and write your opinion accordingly.



$$((1/1.70) - (1/2)) / (1 - 1/2) = 0.18$$
$$((1/2.45) - (1/4)) / (1 - 1/4) = 0.21$$
$$((1/2.83) - (1/6)) / (1 - 1/6) = 0.22$$
$$((1/3.03) - (1/8)) / (1 - 1/8) = 0.23$$

Since 'e' is increasing with p, it suggests that parallelization overhead is also contributing to the poor speedup. Hence, we need to reduce this overhead to improve speedups.

Question # 3: [10 marks, CLO # 1]

What is Fat Tree Routing? Draw a Fat tree having 4 levels with 8 leaf nodes. Provide a comparative analysis with a star topology with same number of nodes in terms of cost, diameter, bisection width and arc connectivity. You have to mention which routing is better in terms of each evaluation parameter.

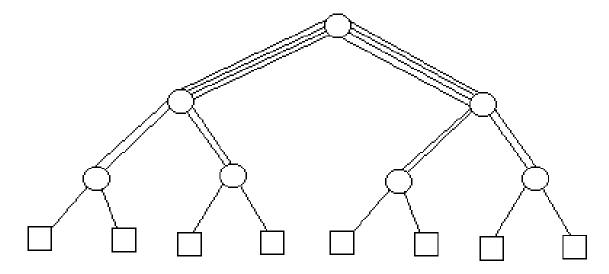
Answer

Cost, diameter, bisection width and arc connectivity

	diameter	bisection width	arc connectivity	Cost
Star	2	1	1	p-1
Fat Tree	2 log ((p+1)/2)	1	1	n x 2 ⁿ

p = number of nodes
n = number of levels

For the following tree, we have number of nodes 15 and levels 3, starting from 0.



In terms of above evaluation parameters, start has less cost and diameter however all traffic has to be passed through center node. On the other hand, fat tree manages the traffic well with distributed nodes and larger bandwidth at edges with large traffic.

(i) Show the output of the following program:

```
#include <iostream>
#include <omp.h>
using namespace std;
int main() {
                                                              At thread 0.0.0 iteration 0.1.2
  int nums[10] = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\};
                                                              At thread 1,1,1, iteration 3,4,5
  omp_set_num_threads(3);
                                                              At thread 2,2,2 iteration 6,7,8
#pragma omp parallel for schedule(static, 3)
                                                              At thread 0 iteration 9
  for (int j = 0; j < 10; j++) {
    nums[i] *= (i+3);
                                                              nums = [3,8,15.....]
    int x = omp get thread num();
    cout << "At thread: " << x << " iteration: ";
    cout << i << endl;
  }
  for (int i = 0; i < 10; i++) {
    cout << nums[i] << " ";
  cout << endl;
  return 0;
}
At thread: 0 iteration: 0
At thread: 0 iteration: 1
At thread: 0 iteration: 2
At thread: 1 iteration: 3
At thread: 1 iteration: 4
At thread: 1 iteration: 5
At thread: 2 iteration: 6
At thread: 2 iteration: 7
At thread: 2 iteration: 8
At thread: 0 iteration: 9
3 8 15 24 35 48 63 80 99 120
```

Semaphores

(ii) In general, what would you have to do if you are not allowed to use "OpenMP for construct" and you have to explicitly map the iterations to threads? Illustrate your answer with OpenMP code or pseudocode

We would have to make sure the iterations managed by different threads are mutually exclusive. We may use the thread-ID for this purpose (as a different value to initialize each threads iterations, and the number of threads as a separation or jump value between each iteration for a specific thread.