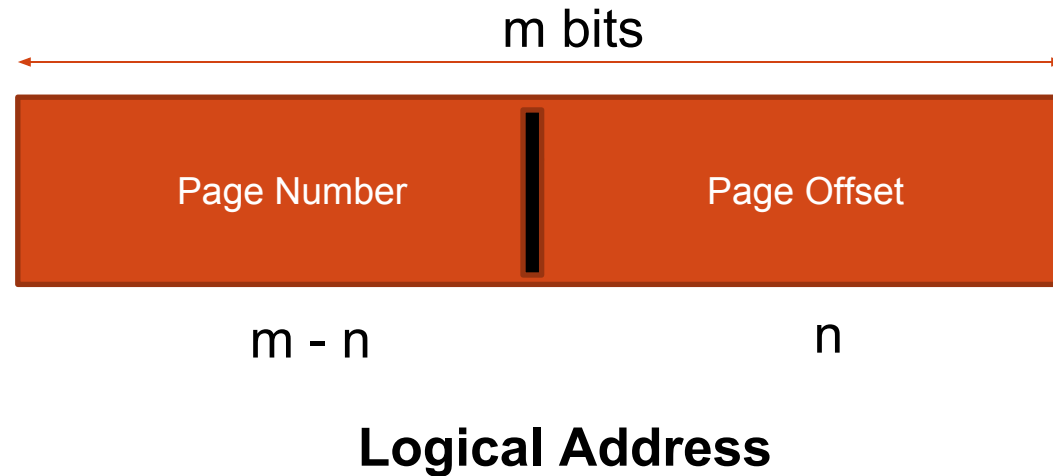


PAGING

■ **Rules:**

- If we have M bit processor, then logical address will also of M bits long.
- Let the size of logical address space is 2^m and page size is 2^n (bytes or words), then the high order ($m-n$) bits of logical address designate the page number (p) and lower order bits (n) designate the page of offset.



- **Example:**

- Consider the logical address space of **8** pages of **1024** words mapped onto physical memory of **32** frames

- **1. How many bits are there in logical address:**

Given page size = 1024 words = 1k

Page size = $2^{10} \Leftrightarrow 2^n$ so, $n = 10$

Logical Address space = $8 \times 1K = 8K = 2^3 \times 2^{10} = 2^{13}$ so, $m = 13$

$p = 3 (m - n)$

$d = 10 (n)$

There are **13 bits** in logical address space.

- **2. How many bits are there in physical address**

Physical address space = $32 \times 1K = 32K = 2^5 \times 2^{10} = 2^{15}$

There are **15 bits** in physical address.

- **Example:**

- Given **32-bit** processor, page size = 1024 bytes

- **Find:**

- **1. size of logical address :**

32 bits (as 32-bit processor)

- **2. No. of bits to represent page number and offset:**

As page size = 1024 bytes = 2^{10} (2^n) so, **n = 10**

No. of bits used to represent page no. = $m - n = 32 - 10 = \mathbf{22 \text{ bits}}$

- **3. Max size of logical address space:**

2^{32} bytes = $2^2 \times 2^{30} = 4\text{GBytes}$

- **4. Max number of pages in logical address space:**

Space = $2^{m-n} = 2^{22} = 2^2 \times 2^{20} = 4\text{MBytes} = 4 \text{ Million}$

- **5. Max length of page table of a process :**

4 M entries (max no. of pages in process)

- ***Note to remember:***

- $1K = 2^{10} \text{ Bytes}$

- $1MB = 1024KB$

$= 2^{10} \times 2^{10}$

$= 2^{20} \text{ Bytes}$

- $1 \text{ GB} = 1024 \text{ MB}$

$= 2^{30} \text{ Bytes}$

❑ **Example:** Assume a logical address space of 16 pages of 1024 words, each mapped into a physical memory of 32 frames.

- No. of bits needed for $p = \lceil \log 16 \rceil = 2^4 = 4 \text{ bits}$
- No. of bits needed for $f = \lceil \log 32 \rceil = 2^5 = 5 \text{ bits}$
- No. of bits needed for $d = 2048 \text{ bytes} = 2^{11}$
- Logical address size = $|p| + |d| = 4 + 11 = 15 \text{ bits}$
- Physical address size = $|f| + |d| = 5 + 11 = 16 \text{ bits}$

❑ **Page Table Size**—Page table size = $NP * PTES$, (assuming a byte size page table entry)

- where NP is the number of pages in the process address space and $PTES$ is the page table entry size (equal to $|f|$)

- 16-bit logical address
- 2^{16}
- 8K page size
- Maximum pages in a process address space =
- $2^{16}/8K = 8192$ pages
- $|d| = \log_2 8K = 13$ bits
- $|p| = 16 - 13 = 3$ bits

▪ **Example:**

- Logical address = 32-bit
- Process address space = 2^{32} B = 4 GB
- Main memory = RAM = 512 MB
- Page size = 4K
- Maximum pages in a process address space = $2^{32}/4K=1M$
- $|d| = 32-20=12\text{bits}$
- $|p| = \text{ceiling}[\log 1M]=20\text{bits}$
- No. of frames = $512\text{MB}/4K = 128K$
- $|f| = \text{ceiling}[\log 128K]=17\text{bits}$
- Physical address = $|f|+|d|=17+12=29\text{bits}$