# Information Security CS3002 (Sections BDS-7A/B) Lecture 21

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# Types of Authentication

- There are two basic types of authentication: non-repudiable and repudiable.
- Repudiable Authentication involves factors, "what you know" and "what you have, "that can present problems to the authenticator because the information presented can be unreliable because such factors suffer from several well-known problems including the fact that possessions can be lost, forged, or easily duplicated.
- Non-repudiable Authentication involves characteristics whose proof of origin cannot be denied. Such characteristics include biometrics like iris patterns, retinal images, and hand geometry and they positively verify the identity of the individual

#### Authentication Mechanisms

In general, authentication takes one of the following three forms:

- Basic authentication involving a server:
  - The server maintains a user file of either passwords and usernames or some other useful piece of authenticating information. This information is always examined before authorization is granted.
- Challenge-response:
  - In which the server or any other authenticating system generates a challenge to the host requesting for authentication and expects a response.
- Centralized authentication:
  - In which a central server authenticates users on the network and in addition also authorizes and audits them

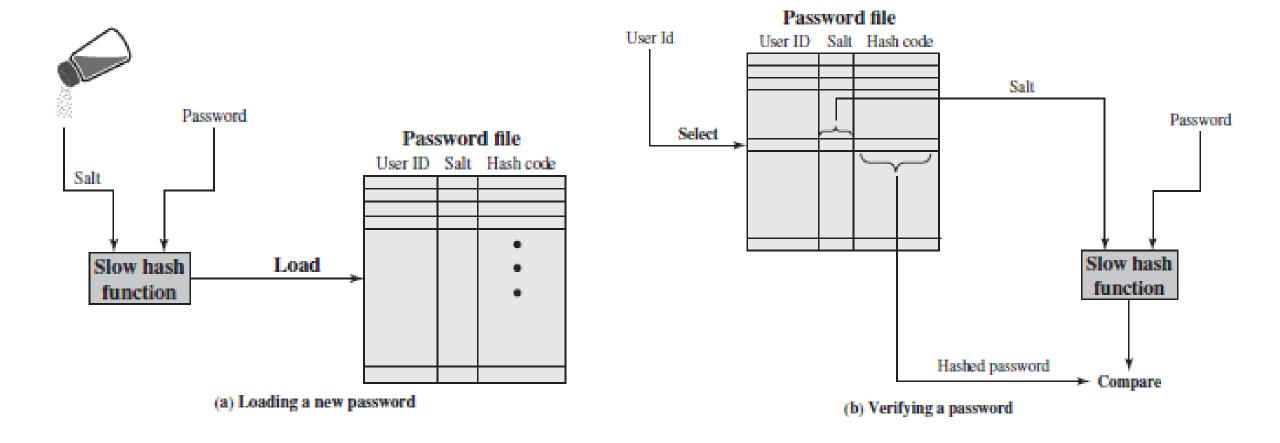
#### Password-based Authentication

- secret = the user password
- (client) create and transmit proof
  - F = I (the identity function)
  - i.e. proof = password (cleartext!)
- (server) verify the proof:
  - case #1: F = I (the identity function)
    - server knows all passwords in cleartext(!)
    - access control: proof = password ?
  - case #2: F = one-way hash (that is a digest)
    - server knows the passwords' digests, HUID
    - access control: f(proof) = HUID ?

#### **Passwords**

- Authentication based on alphanumeric characters or numbers
  - Benefits
    - Easy to remember (if only for one system)
  - Disadvantages
    - User-side password storage:
      - Post-It!
      - Client-side password manager or wallet
    - Password guessable (my son's name)
    - Password readable during transmission
    - Server-side password storage issues (hashing is must)
      - 35% passwords identified using dictionary attack
      - Use "salt"
    - Shoulder surfing
    - Using same password in multiple places

# Using "salt" and hash: Fig 3.3, UNIX Password Scheme

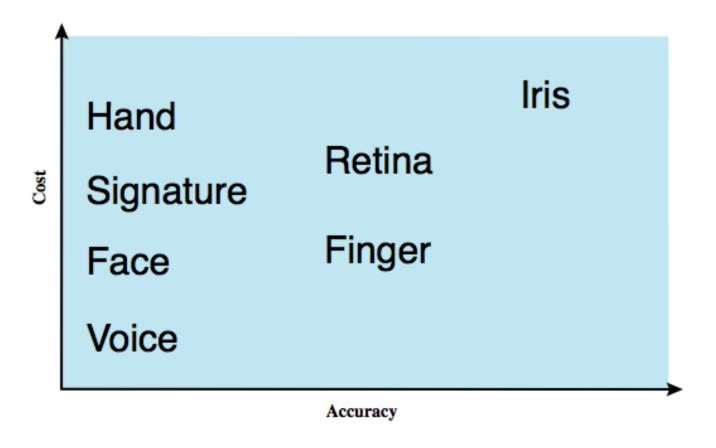


# Using "salt" and hash

- For each user ID:
  - Create/ask the password
  - Generate a random salt (should contain rarely used or control characters)
- Compute HP = hash(password | salt)
- Store the triples {UID, HP, salt<sub>UID</sub>}
- Advantages:
  - Prevents duplicate passwords from being visible in the password file (different HP for users having the same password)
  - Increases the difficulty of offline dictionary attacks
  - Nearly impossible to tell if a person used the same password on multiple systems

#### Biometric Authentication

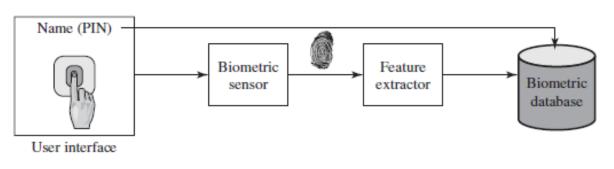
- Authenticate user based on one of their physical characteristics:
  - Facial
  - Fingerprint
  - Hand geometry
  - Retina pattern
  - Iris
  - Signature
  - Voice



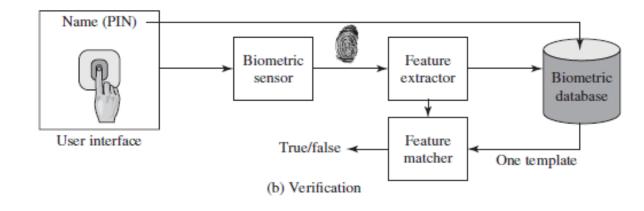
# Operation of a biometric system

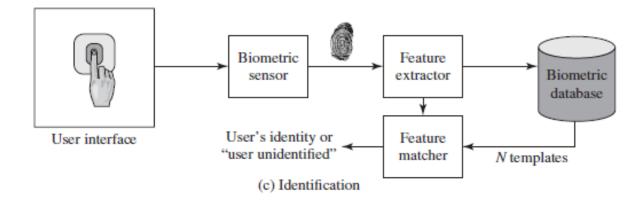
 Verification is analogous to user login via a smart card and a PIN

 Identification is biometric info but no IDs; system compares with stored templates



(a) Enrollment





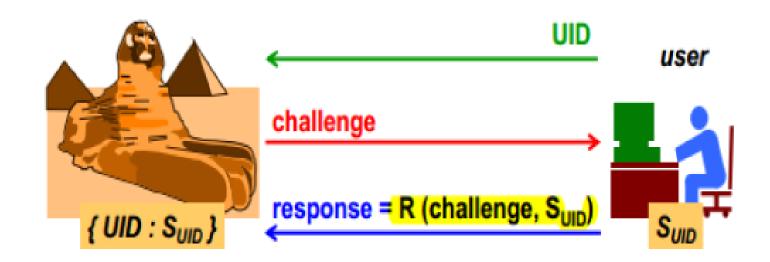
# Problems with biometric systems

- FAR = False Acceptance Rate
- FRR = False Rejection Rate
- FAR and FRR may be partly tuned but they heavily depend on the cost of the device

- variable biological characteristics:
  - Finger wound
  - Voice altered due to emotion or injury
  - retinal blood pattern altered due to alcohol or drug

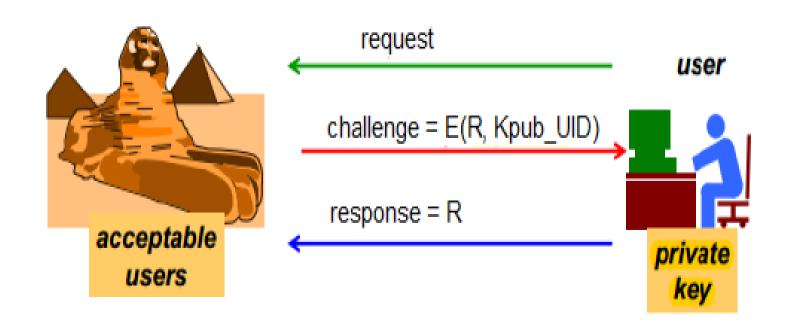
# Symmetric challenge-response authentication

- a challenge (typically a random nonce) is sent to the user ...
- ... who replies with the solution after a computation
- involving the shared secret and the challenge
- the server must know the secret in clear
- often R is a hash function (can't be encryption)

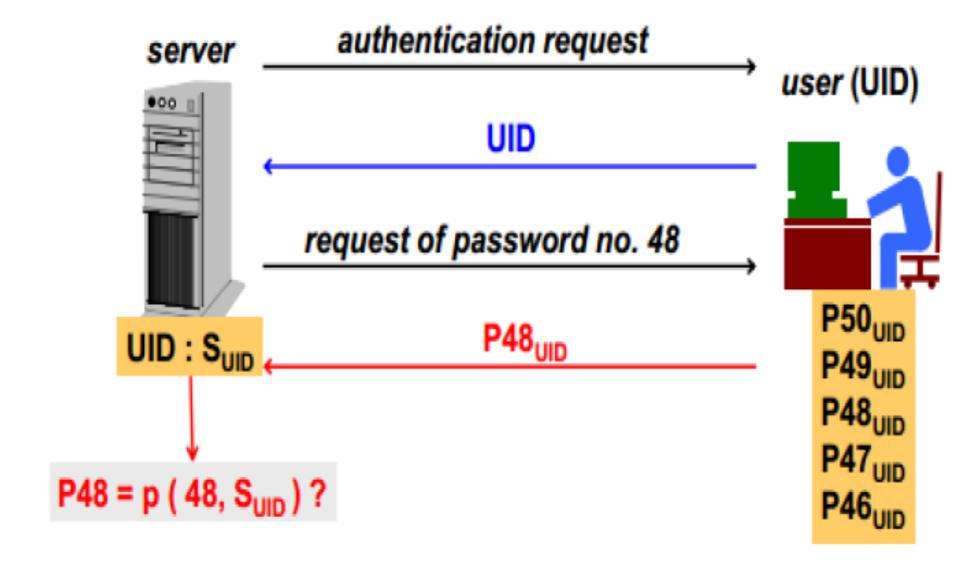


### Asymmetric challenge response system

- a random number R is encrypted with the user's public key ...
- and the user replies by sending R in clear thanks to its knowledge of the private key



#### One-time Passwords

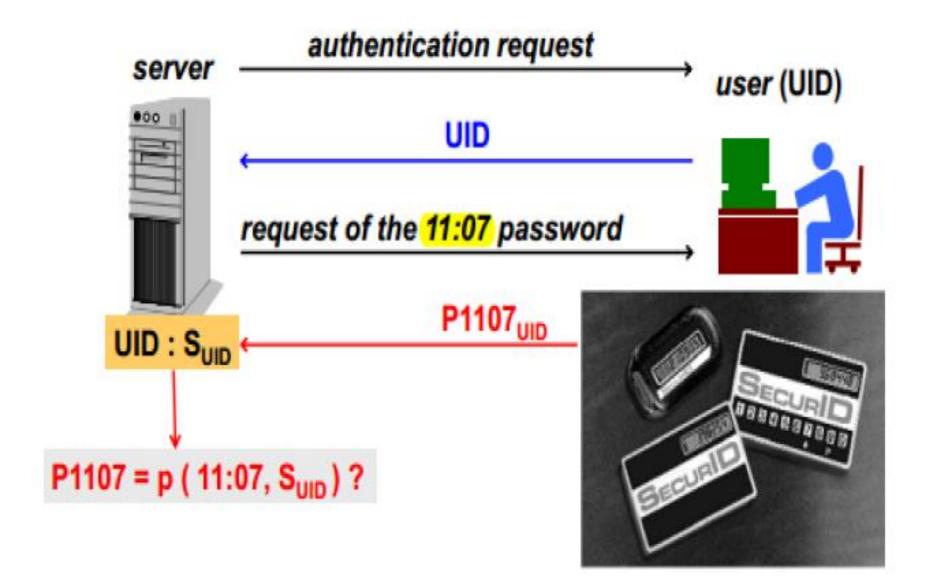


#### One-time Passwords





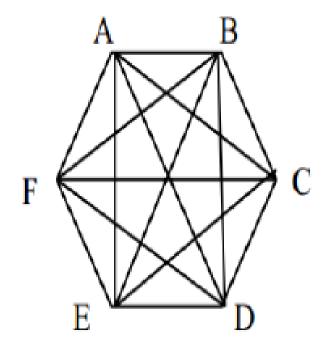
# Passwords (one-time token based)



#### SecurID: architecture

- Invented and patented by Security Dynamics
- Time-based synchronous OTP technique:
  - $P_{UID}(t) = h(S_{UID}, t)$
- The client sends:
  - user , PIN , token-code (seed, time)
- Based on user name and PIN the server verifies against three possible tokencodes:
  - TC-1, TC-0, TC+1
- More than once if there is a drift of more than one minute
- Wrong authentication attempts limited

# Secret keys for N-system network



- n system need n(n-1)/2 pairs of secret keys
- Each system remembers n-1 keys
- If a new system comes in n new key are generated
- If a system leaves, n-1 keys are removed

#### PKI and Certificate Authorities

- Certificate consists of:
  - A public key plus a User ID of the key owner
  - Signed by a third party trusted by community
  - Often govt/bank certificate authority (CA)
- Users obtain certificates from CA
  - Create keys & unsigned certificate, gives to CA, CA signs certificate & attaches signature, returns to user
- Other users can verify certificate
  - Checking signature on certificate using CA's public key

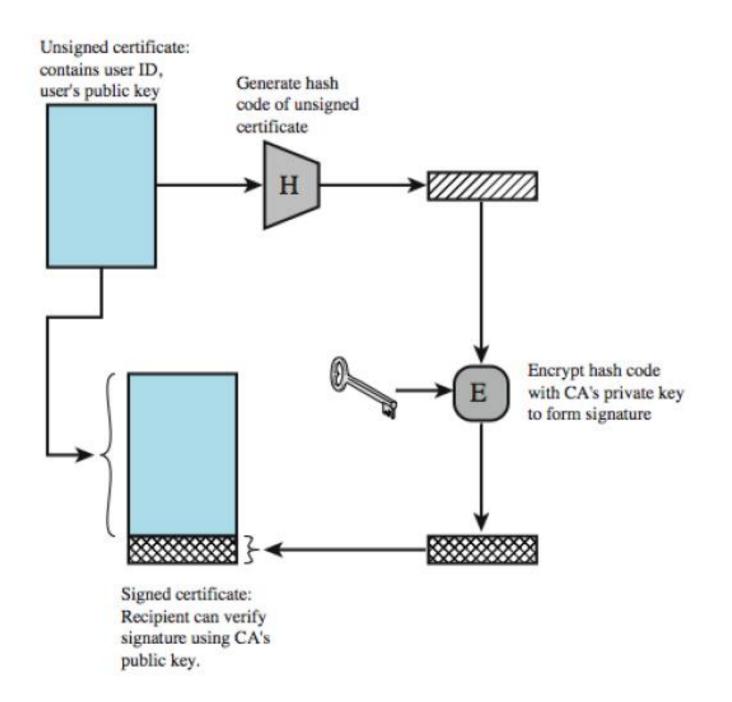
# Common Key Steps

- 1. User software creates a pair of keys: private and public
- 2. Clients prepares unsigned certificate that includes user ID and public key
- 3. User provides unsigned certificate to a CA
- 4. CA creates a signature:
  - I. Creates a hash code of the unsigned certificate
  - II. Encrypts the hash code with the CA's private key
- 5. CA attaches the signature to unsigned certificate to make signed certificate

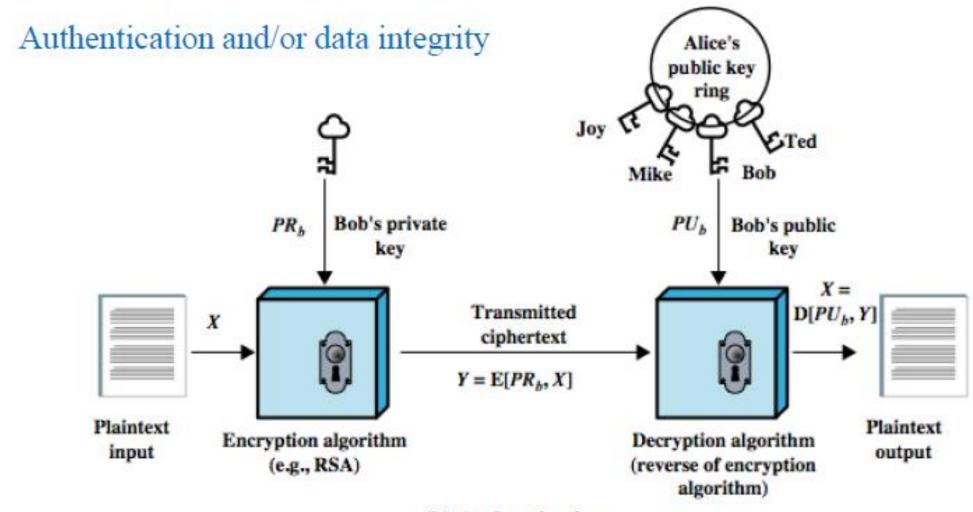
# Key Steps (continued)

- 6. CA returns the signed certificate to the client
- 7. Client may provide signed signature to other users
- 8. Any user may verify the certificate
  - i. Calculate the hash code of certificate (exclude signature)
  - ii. Decrypt signature using CA's public key
  - iii. Compare the two

# Public Key Certificates



# Public Key Authentication



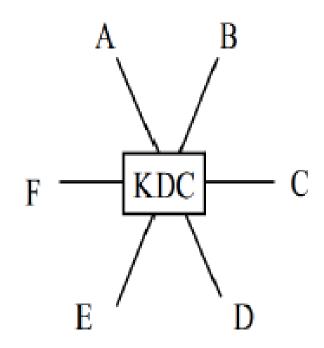
(b) Authentication

# Key Distribution Center (KDC)

- Each node is configured with KDC's key.
- KDC has all the keys.
- A  $\leftarrow \rightarrow$  B communication?
  - KDC sends a key K<sub>AB</sub> encrypted with A's key to A and B's key to B.



- If KDC is compromised, all systems are compromised
- KDC is single point of failure or performance bottleneck
- KDC has to be on-line all the time. Replication!



#### Kerberos

- Network authentication protocol
- Based on Trusted Third Party (TTP) KDC
- invented by MIT for project Athena



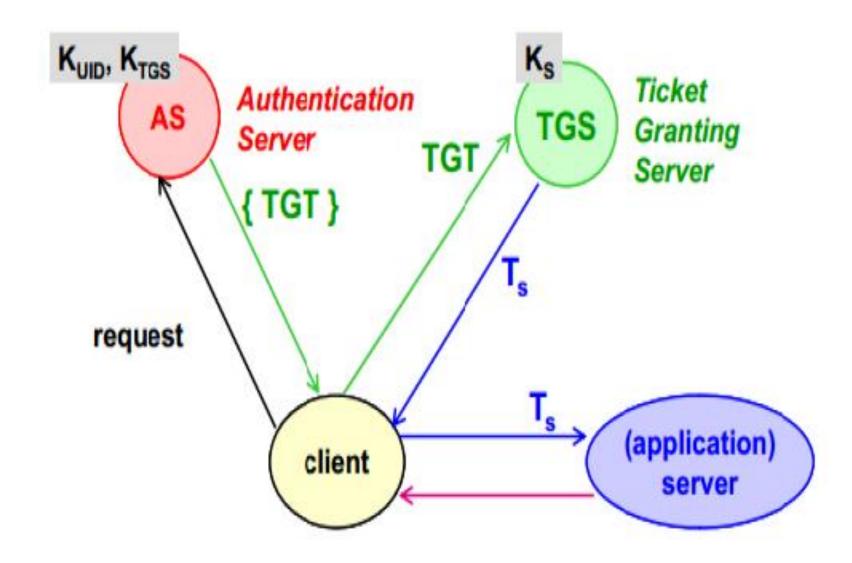
**Kerberos: The Network Authentication Protocol** 

- Named after Greek mythological character "Cerberus" Three headed dog protecting the entrance of Hades
- Used by popular operating systems and servers
- Protect against eavesdropping and firewall limitation to users and replay attacks

#### Kerberos overview

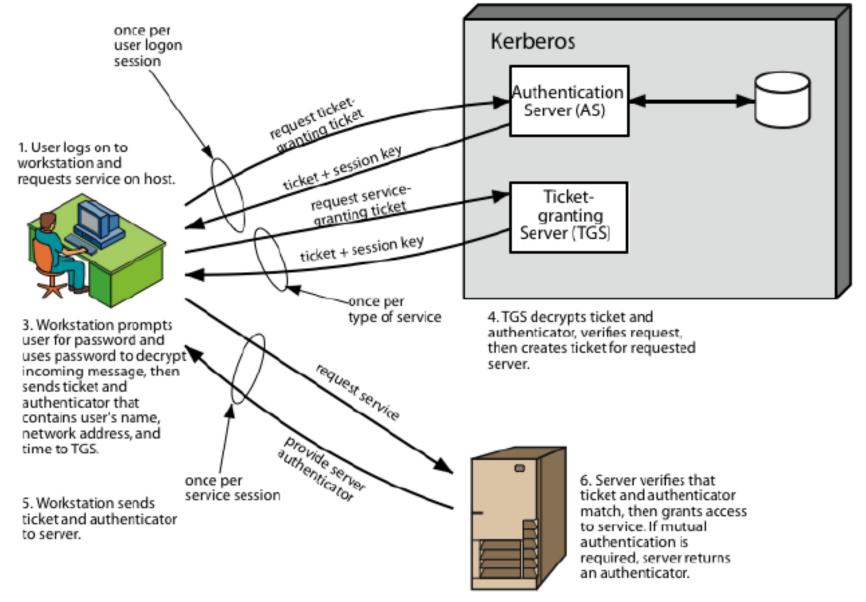
- Authentication server authenticates a user to a specific service in the network
- TGS: Ticket Granting Server, grants ticket to the user
- Authentication server and TGS can be the same system. They work as a single unit.
- Application Server provides the service to the user
- The client/user, Auth. Server & TGS, Application server are the 3 heads of Kerberos!

# Kerberos high-level view

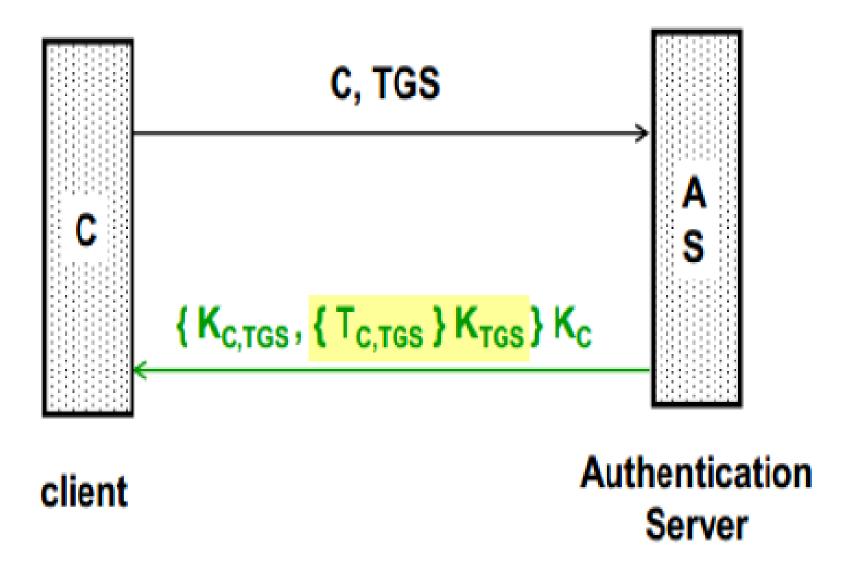


# Kerberos Protocol

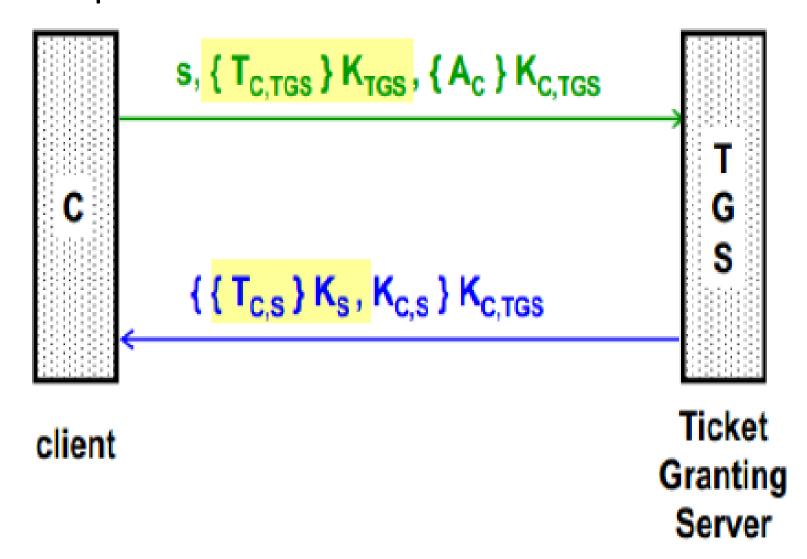
 AS verifies user's access right in database, creates ticket-granting ticket and session key. Results are encrypted using key derived from user's password.



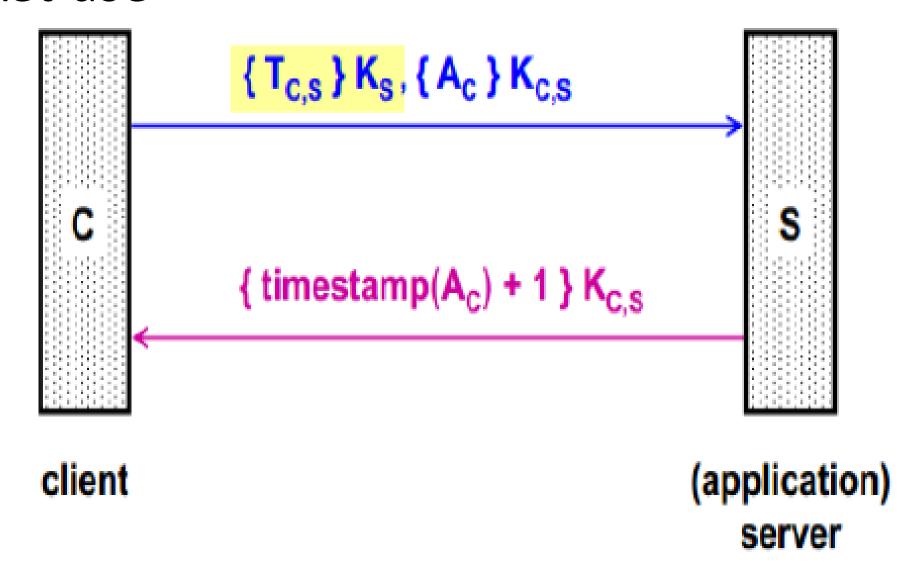
# TGT Request



# Ticket Request



#### Ticket use



# Kerberos: Component Breakdown

- K<sub>C, TGS</sub>:
  - This is the session key that is shared between the client and the TGS. It is generated by the Authentication Server (AS) and included in the Ticket-Granting Ticket (TGT) that the client receives during the initial authentication process.
- $\{T_{C,TGS}\}$   $K_{TGS}$ :
  - This is the encrypted message that contains the client's identity (TC) and the name of the desired service (TGS), encrypted using the TGS's secret key (KTGS). The TGS's secret key is known only to the TGS and the Kerberos Key Distribution Center (KDC).
- {K<sub>C,TGS</sub>, {T<sub>C,TGS</sub>} K<sub>TGS</sub>}:
  - This is the message that combines the session key and the encrypted ticket information. It is encrypted using the session key  $K_{C,TGS}$  that is shared between the client and the TGS.
- K<sub>C</sub>:
  - This is the final layer of encryption that is applied to the entire message using the session key Kc that is shared between the client and the service. The final encrypted message is then sent to the TGS.

# Other Authentication Systems

- OATH (open authentication)
  - Interoperability of authentication systems based on OTP, both symmetric and asymmetric
- SSO (single sign-on)
  - Single credential multiple services

