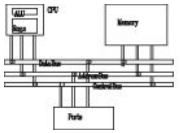
Conditional Processing

Computer Organization and Assembly Languages

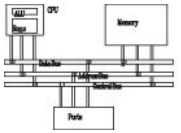
with slides by Kip Irvine

Announcements



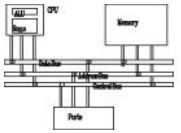
- Midterm exam: Room 103, 10:00am-12:00am next Thursday, open book, chapters 1-5.
- Assignment #2 is online.

Assignment #2 CRC32 checksum



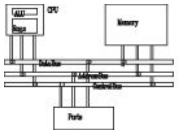
```
unsigned int crc32(const char* data,
                   size_t length)
{
  // standard polynomial in CRC32
  const unsigned int POLY = 0xEDB88320;
  // standard initial value in CRC32
  unsigned int reminder = 0xFFFFFFFF;
  for(size_t i = 0; i < length; i++){
    // must be zero extended
    reminder ^= (unsigned char)data[i];
    for(size_t bit = 0; bit < 8; bit++)
      if(reminder & 0x01)
        reminder = (reminder >> 1) ^ POLY;
      else
        reminder >>= 1;
  }
  return reminder ^ 0xFFFFFFFF;
}
```

Boolean and comparison instructions



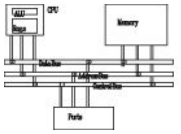
- CPU Status Flags
- AND Instruction
- OR Instruction
- XOR Instruction
- NOT Instruction
- Applications
- TEST Instruction
- CMP Instruction

Status flags - review



- The Zero flag is set when the result of an operation equals zero.
- The Carry flag is set when an instruction generates a result that is too large (or too small) for the destination operand.
- The Sign flag is set if the destination operand is negative, and it is clear if the destination operand is positive.
- The Overflow flag is set when an instruction generates an invalid signed result.
- Less important:
  - The Parity flag is set when an instruction generates an even number of 1 bits in the low byte of the destination operand.
  - The Auxiliary Carry flag is set when an operation produces a carry out from bit 3 to bit 4

NOT instruction



- Performs a bitwise Boolean NOT operation on a single destination operand
- Syntax: (no flag affected)

  NOT *destination*
- Example:

```
mov al, 11110000b
not al
```

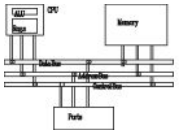
NOT 0 0 1 1 1 0 1 1

1 1 0 0 0 1 0 0 —— inverted

NOT

X	¬X
F	T
T	F

AND instruction



- Performs a bitwise Boolean AND operation between each pair of matching bits in two operands
- Syntax: (O=0,C=0,SZP)

  AND *destination, source*

- Example:

```
mov al, 00111011b
and al, 00001111b
```

```
  00111011
AND 00001111
  ---------
  00001011
```

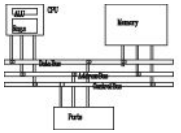
cleared — 0000 — 1011 — unchanged

bit extraction

AND

x	y	x ∧ y
0	0	0
0	1	0
1	0	0
1	1	1

OR instruction



- Performs a bitwise Boolean OR operation between each pair of matching bits in two operands
- Syntax: (O=0,C=0,SZP)

  OR *destination, source*
- Example:

```
mov dl, 00111011b
or  dl, 00001111b
```

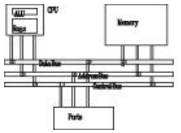
```
    00111011
OR  00001111
------------
    00111111
```

unchanged — 0 0 1 1 | 1 1 1 1 — set

OR

x	y	x ∨ y
0	0	0
0	1	1
1	0	1
1	1	1

XOR instruction



- Performs a bitwise Boolean exclusive-OR operation between each pair of matching bits in two operands
- Syntax: (O=0,C=0,SZP)

  XOR *destination, source*
- Example:

```
mov dl, 00111011b
xor dl, 00001111b
```

```
        00111011
XOR     00001111
        --------
        00110100
```

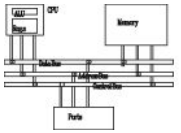
unchanged — inverted

XOR

x	y	x ⊕ y
0	0	0
0	1	1
1	0	1
1	1	0

XOR is a useful way to invert the bits in an operand and data encryption

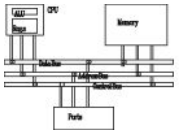
Applications (1 of 5)



- Task: Convert the character in AL to upper case.
- Solution: Use the AND instruction to clear bit 5.

```
mov al,'a'     ; AL = 01100001b
and al,11011111b ; AL = 01000001b
```

Applications (2 of 5)

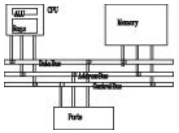


- Task: Convert a binary decimal byte into its equivalent ASCII decimal digit.
- Solution: Use the OR instruction to set bits 4 and 5.

```
mov al,6   ; AL = 00000110b
or  al,00110000b ; AL = 00110110b
```

The ASCII digit '6' = 00110110b

Applications (3 of 5)



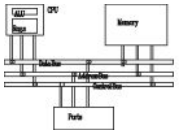
- Task: Turn on the keyboard CapsLock key
- Solution: Use the OR instruction to set bit 6 in the keyboard flag byte at 0040:0017h in the BIOS data area.

```
mov ax,40h                  ; BIOS segment
mov ds,ax
mov bx,17h                  ; keyboard flag byte
or BYTE PTR [bx],01000000b ; CapsLock on
```

This code only runs in Real-address mode, and it does not work under Windows NT, 2000, or XP.

Ni g ider to kam ni kr raha ye code.

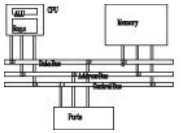
Applications (4 of 5)



- Task: Jump to a label if an integer is even.
- Solution: AND the lowest bit with a 1. If the result is Zero, the number was even.

```
mov ax,wordVal
and ax,1 ; low bit set?
jz  EvenValue  ; jump if Zero flag set
```

Applications (5 of 5)

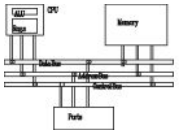


- Task: Jump to a label if the value in AL is not zero.
- Solution: OR the byte with itself, then use the JNZ (jump if not zero) instruction.

```
or  al,al
jnz IsNotZero  ; jump if not zero
```

ORing any number with itself does not change its value.

TEST instruction



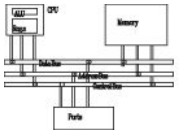
- Performs a nondestructive AND operation between each pair of matching bits in two operands
- No operands are modified, but the flags are affected.
- Example: jump to a label if either bit 0 or bit 1 in AL is set.

```
test al,00000011b
jnz  ValueFound
```

- Example: jump to a label if neither bit 0 nor bit 1 in AL is set.

```
test al,00000011b
jz   ValueNotFound
```

CMP instruction (1 of 3)



- Compares the destination operand to the source operand
  - Nondestructive subtraction of source from destination (destination operand is not changed)
- Syntax: (OSZCAP)

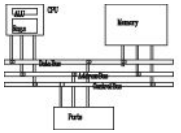
  CMP *destination, source*
- Example: destination == source

```
mov al,5
cmp al,5 ; Zero flag set
```

- Example: destination < source

```
mov al,4
cmp al,5 ; Carry flag set
```

CMP instruction (2 of 3)



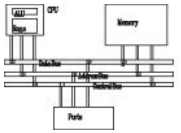
- Example: destination > source

```
mov al,6
cmp al,5 ; ZF = 0, CF = 0
```

(both the Zero and Carry flags are clear)

The comparisons shown so far were unsigned.

CMP instruction (3 of 3)



The comparisons shown here are performed with signed integers.

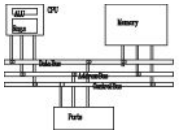
- Example: destination > source

```
mov al,5
cmp al,-2 ; Sign flag == Overflow flag
```

- Example: destination < source

```
mov al,-1
cmp al,5  ; Sign flag != Overflow flag
```

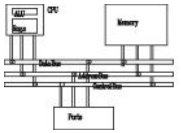
Setting and clearing individual flags



```
and al, 0      ; set Zero
or  al, 1      ; clear Zero
or  al, 80h      ; set Sign
and al, 7Fh      ; clear Sign
stc          ; set Carry
clc          ; clear Carry

mov al, 7Fh
inc al       ; set Overflow

or eax, 0      ; clear Overflow
```

Conditional jumps

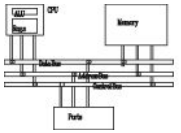
Conditional structures



- There are no high-level logic structures such as if-then-else, in the IA-32 instruction set. But, you can use combinations of comparisons and jumps to implement any logic structure.
- First, an operation such as CMP, AND or SUB is executed to modified the CPU flags. Second, a conditional jump instruction tests the flags and change the execution flow accordingly.

```
        CMP AL, 0
        JZ  L1
        :
L1:
```

J*cond* instruction



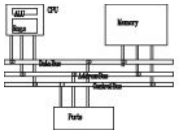
- A conditional jump instruction branches to a label when specific register or flag conditions are met

  `Jcond` *`destination`*

- Four groups: (some are the same)

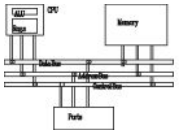
1. based on specific flag values
2. based on equality between operands
3. based on comparisons of unsigned operands
4. based on comparisons of signed operands

Jumps based on specific flags



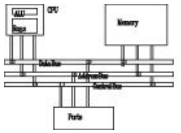
Mnemonic	Description	Flags
JZ	Jump if zero	ZF = 1
JNZ	Jump if not zero	ZF = 0
JC	Jump if carry	CF = 1
JNC	Jump if not carry	CF = 0
JO	Jump if overflow	OF = 1
JNO	Jump if not overflow	OF = 0
JS	Jump if signed	SF = 1
JNS	Jump if not signed	SF = 0
JP	Jump if parity (even)	PF = 1
JNP	Jump if not parity (odd)	PF = 0

Jumps based on equality



Mnemonic	Description
JE	Jump if equal (*leftOp* = *rightOp*)
JNE	Jump if not equal (*leftOp* ≠ *rightOp*)
JCXZ	Jump if CX = 0
JECXZ	Jump if ECX = 0

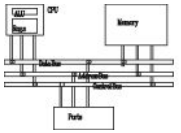
Jumps based on unsigned comparisons



Mnemonic	Description
JA	Jump if above (if *leftOp* > *rightOp*)
JNBE	Jump if not below or equal (same as JA)
JAE	Jump if above or equal (if *leftOp* >= *rightOp*)
JNB	Jump if not below (same as JAE)
JB	Jump if below (if *leftOp* < *rightOp*)
JNAE	Jump if not above or equal (same as JB)
JBE	Jump if below or equal (if *leftOp* <= *rightOp*)
JNA	Jump if not above (same as JBE)

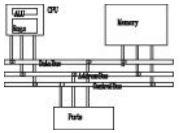
$>\geqq<\leqq$

Jumps based on signed comparisons



Mnemonic	Description
JG	Jump if greater (if *leftOp* > *rightOp*)
JNLE	Jump if not less than or equal (same as JG)
JGE	Jump if greater than or equal (if *leftOp* >= *rightOp*)
JNL	Jump if not less (same as JGE)
JL	Jump if less (if *leftOp* < *rightOp*)
JNGE	Jump if not greater than or equal (same as JL)
JLE	Jump if less than or equal (if *leftOp* <= *rightOp*)
JNG	Jump if not greater (same as JLE)

Examples



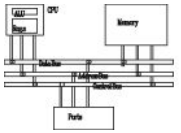
- Compare unsigned AX to BX, and copy the larger of the two into a variable named Large

```
    mov Large,bx
    cmp ax,bx
    jna Next
    mov Large,ax
Next:
```

- Compare signed AX to BX, and copy the smaller of the two into a variable named Small

```
    mov Small,ax
    cmp bx,ax
    jnl Next
    mov Small,bx
Next:
```

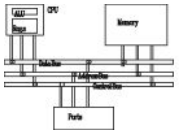
Examples

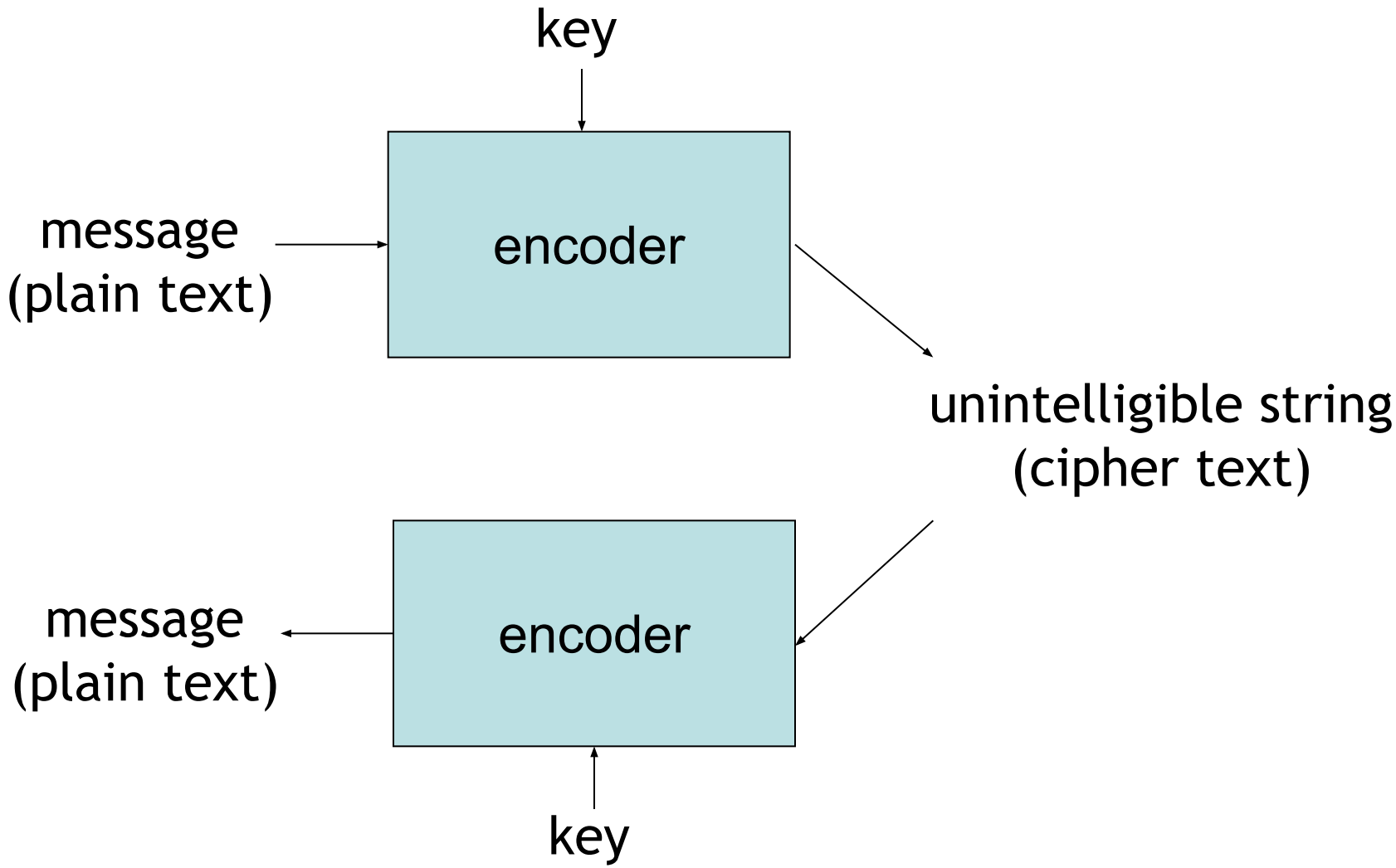


- Find the first even number in an array of unsigned integers

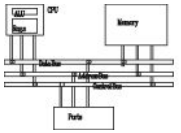
```
.date
intArray DWORD 7,9,3,4,6,1
.code
...
        mov  ebx, OFFSET intArray
        mov  ecx, LENGTHOF intArray
L1:     test DWORD PTR [ebx], 1
        jz   found
        add  ebx, 4
        loop L1
...
```

String encryption





Encrypting a string

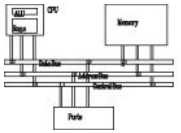


```
KEY = 239
.data
buffer BYTE BUFMAX DUP(0)
bufSize DWORD ?
.code
    mov ecx,bufSize  ; loop counter
    mov esi,0 ; index 0 in buffer
L1:
    xor buffer[esi],KEY ; translate a byte
    inc esi    ; point to next byte
    loop L1
```

```
Message: Attack at dawn.

Cipher text: «¢¢Äîä-Ä¢-ïÄÿü-Gs

Decrypted: Attack at dawn.
```

Conditional loops

LOOPZ and LOOPE



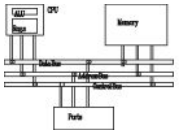
- Syntax:

  LOOPE *destination*

  LOOPZ *destination*

- Logic:
  - ECX ← ECX – 1
  - if ECX > 0 and ZF=1, jump to *destination*
- The destination label must be between -128 and +127 bytes from the location of the following instruction
- Useful when scanning an array for the first element that meets some condition.

LOOPNZ and LOOPNE



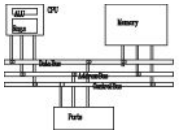
- Syntax:

  LOOPNZ *destination*

  LOOPNE *destination*

- Logic:
  - ECX ← ECX – 1;
  - if ECX > 0 and ZF=0, jump to *destination*

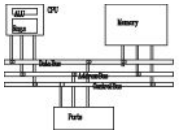
LOOPNZ example



The following code finds the first positive value in an array:

```
.data
array SWORD -3,-6,-1,-10,10,30,40,4
sentinel SWORD 0
.code
    mov esi,OFFSET array
    mov ecx,LENGTHOF array
next:
    test WORD PTR [esi],8000h  ; test sign bit
    pushfd ; push flags on stack
    add esi,TYPE array
    popfd  ; pop flags from stack
    loopnz next   ; continue loop
    jnz quit  ; none found
    sub esi,TYPE array  ; ESI points to value
quit:
```

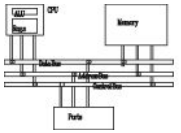
Your turn



Locate the first nonzero value in the array. If none is found, let ESI point to the sentinel value:

```
.data
array SWORD 50 DUP(?)
sentinel SWORD 0FFFFh
.code
    mov esi,OFFSET array
    mov ecx,LENGTHOF array
L1: cmp WORD PTR [esi],0    ; check for zero









quit:
```

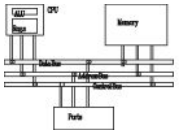
Solution



```
.data
array SWORD 50 DUP(?)
sentinel SWORD 0FFFFh
.code
    mov esi,OFFSET array
    mov ecx,LENGTHOF array
L1: cmp WORD PTR [esi],0    ; check for zero
    pushfd ; push flags on stack
    add esi,TYPE array
    popfd  ; pop flags from stack
    loope next    ; continue loop
    jz quit    ; none found
    sub esi,TYPE array  ; ESI points to value
quit:
```

Conditional structures

Block-structured IF statements

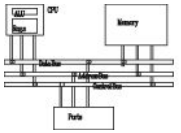


Assembly language programmers can easily translate logical statements written in C++/Java into assembly language. For example:

```
if( op1 == op2 )
  X = 1;
else
  X = 2;
```

```
    mov eax,op1
    cmp eax,op2
    jne L1
    mov X,1
    jmp L2
L1: mov X,2
L2:
```

Example

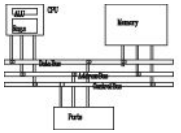


Implement the following pseudocode in assembly language. All values are unsigned:

```
if( ebx <= ecx )
{
  eax = 5;
  edx = 6;
}
```

```
      cmp ebx,ecx
      ja  next
      mov eax,5
      mov edx,6
next:
```

Example

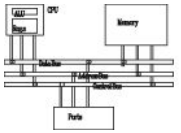


Implement the following pseudocode in assembly language. All values are 32-bit signed integers:

```
if( var1 <= var2 )
  var3 = 10;
else
{
  var3 = 6;
  var4 = 7;
}
```

```
    mov eax,var1
    cmp eax,var2
    jle L1
    mov var3,6
    mov var4,7
    jmp L2
L1: mov var3,10
L2:
```

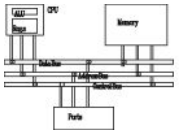
Compound expression with AND



- When implementing the logical AND operator, consider that HLLs use short-circuit evaluation
- In the following example, if the first expression is false, the second expression is skipped:

```
if (al > bl) AND (bl > cl)
  X = 1;
```

Compound expression with AND

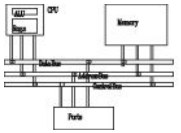


```
if (al > bl) AND (bl > cl)
  X = 1;
```

This is one possible implementation . . .

```
    cmp al,bl ; first expression...
    ja  L1
    jmp next
L1:
    cmp bl,cl ; second expression...
    ja  L2
    jmp next
L2:    ; both are true
    mov X,1   ; set X to 1
next:
```

Compound expression with AND

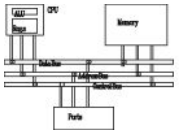


```
if (al > bl) AND (bl > cl)
  X = 1;
```

But the following implementation uses 29% less code by reversing the first relational operator. We allow the program to "fall through" to the second expression:

```
    cmp al,bl ; first expression...
    jbe next  ; quit if false
    cmp bl,cl ; second expression...
    jbe next  ; quit if false
    mov X,1   ; both are true
next:
```

Your turn . . .



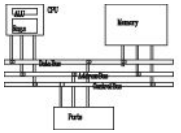
Implement the following pseudocode in assembly language. All values are unsigned:

```
if( ebx <= ecx
 && ecx > edx )
{
  eax = 5;
  edx = 6;
}
```

```
        cmp ebx,ecx
        ja  next
        cmp ecx,edx
        jbe next
        mov eax,5
        mov edx,6
next:
```

(There are multiple correct solutions to this problem.)

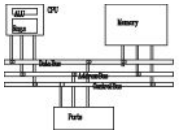
Compound Expression with OR



- In the following example, if the first expression is true, the second expression is skipped:

```
if (al > bl) OR (bl > cl)
  X = 1;
```

Compound Expression with OR

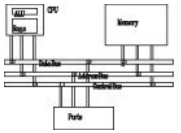


```
if (al > bl) OR (bl > cl)
  X = 1;
```

We can use "fall-through" logic to keep the code as short as possible:

```
    cmp al,bl ; is AL > BL?
    ja  L1 ; yes
    cmp bl,cl ; no: is BL > CL?
    jbe next  ; no: skip next statement
L1: mov X,1   ; set X to 1
next:
```

WHILE Loops

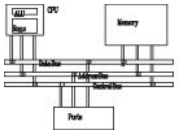


A WHILE loop is really an IF statement followed by the body of the loop, followed by an unconditional jump to the top of the loop. Consider the following example:

```
while( eax < ebx)
    eax = eax + 1;
```

```
_while:
        cmp eax,ebx   ; check loop condition
       jae _endwhile ; false? exit loop
       inc eax    ; body of loop
       jmp _while ; repeat the loop
_endwhile:
```

Your turn . . .

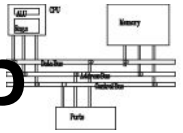


Implement the following loop, using unsigned 32-bit integers:

```
while( ebx <= val1)
{
    ebx = ebx + 5;
    val1 = val1 - 1
}
```

```
_while: cmp ebx,val1     ; check loop condition
        ja  _endwhile  ; false? exit loop
        add ebx,5        ; body of loop
        dec val1
        jmp while          ; repeat the loop
_endwhile:
```

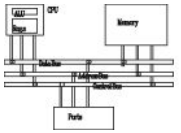
Example: IF statement nested in a loop



```
while(eax < ebx)
{
  eax++;
  if (ebx==ecx)
    X=2;
  else
    X=3;
}
```

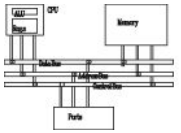
```
_while:   cmp   eax, ebx
          jae   _endwhile
          inc   eax
          cmp   ebx, ecx
          jne   _else
          mov   X, 2
          jmp   _while
_else:    mov   X, 3
          jmp   _while
_endwhile:
```

Table-driven selection



- Table-driven selection uses a table lookup to replace a multiway selection structure (switch-case statements in C)
- Create a table containing lookup values and the offsets of labels or procedures
- Use a loop to search the table
- Suited to a large number of comparisons

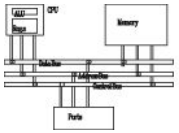
Table-driven selection



Step 1: create a table containing lookup values and procedure offsets:

```
.data
CaseTable BYTE 'A'   ; lookup value
    DWORD Process_A  ; address of procedure
    EntrySize = ($ - CaseTable)
    BYTE 'B'
    DWORD Process_B
    BYTE 'C'
    DWORD Process_C
    BYTE 'D'
    DWORD Process_D

NumberOfEntries = ($ - CaseTable) / EntrySize
```

Table-driven selection

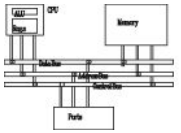


Step 2: Use a loop to search the table. When a match is found, we call the procedure offset stored in the current table entry:

```
    mov ebx,OFFSET CaseTable ; point EBX to the table
    mov ecx,NumberOfEntries ; loop counter

L1:cmp al,[ebx]; match found?
    jne L2; no: continue
    call NEAR PTR [ebx + 1] ; yes: call the procedure
    jmp L3; and exit the loop
L2:add ebx,EntrySize ; point to next entry
    loop L1   ; repeat until ECX = 0

L3:
```

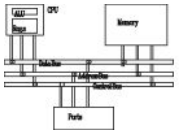
required for procedure pointers

Application: finite-state machines



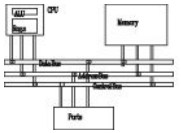
- A finite-state machine (FSM) is a graph structure that changes state based on some input. Also called a state-transition diagram.
- We use a graph to represent an FSM, with squares or circles called nodes, and lines with arrows between the circles called edges (or arcs).
- A FSM is a specific instance of a more general structure called a directed graph (or digraph).
- Three basic states, represented by nodes:
  - Start state
  - Terminal state(s)
  - Nonterminal state(s)

Finite-state machines

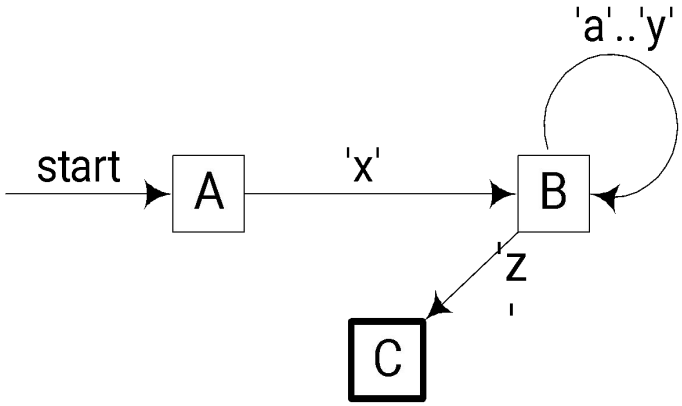


- Accepts any sequence of symbols that puts it into an accepting (final) state
- Can be used to recognize, or validate a sequence of characters that is governed by language rules (called a regular expression)

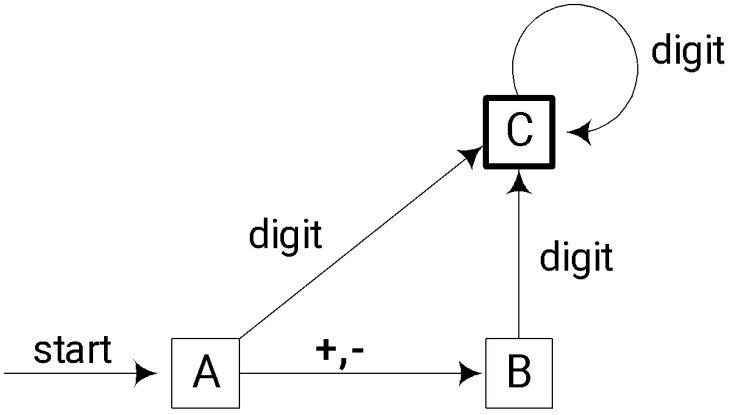
FSM Examples



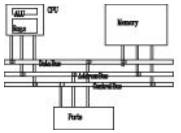
- FSM that recognizes strings beginning with 'x', followed by letters 'a'..'y', ending with 'z':



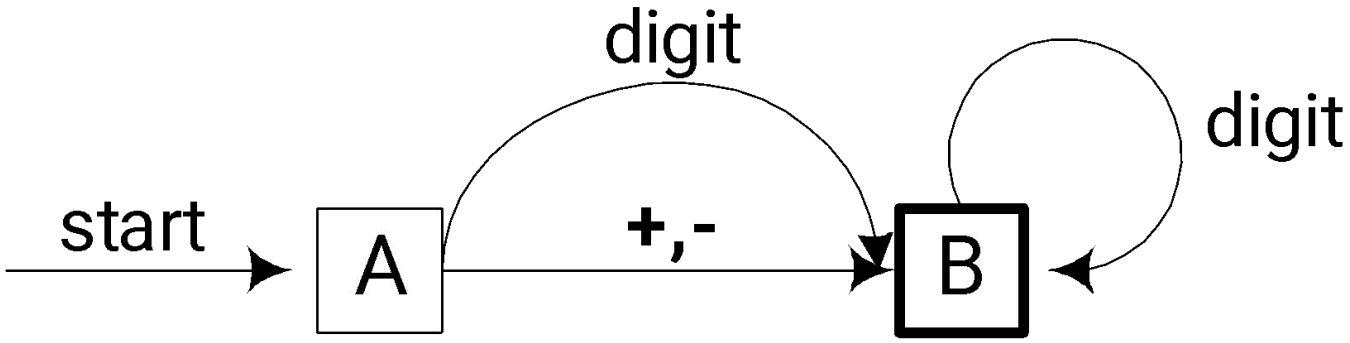
- FSM that recognizes signed integers:



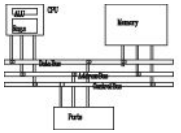
Your turn . . .

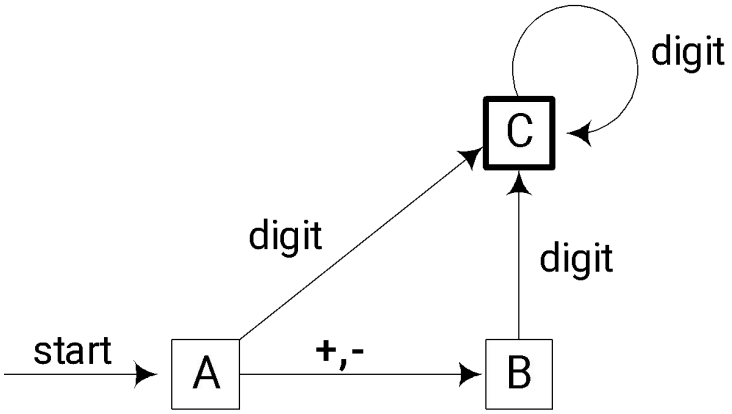


- Explain why the following FSM does not work as well for signed integers as the one shown on the previous slide:



Implementing an FSM

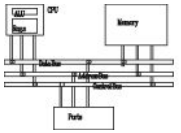




The following is code from State A in the Integer FSM:

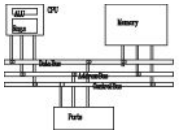
```
StateA:
   call Getnext      ; read next char into AL
   cmp al,'+'        ; leading + sign?
   je StateB         ; go to State B
   cmp al,'-'        ; leading - sign?
   je StateB         ; go to State B
   call IsDigit      ; ZF = 1 if AL = digit
   jz StateC         ; go to State C
   call DisplayErrorMsg ; invalid input found
   jmp Quit
```

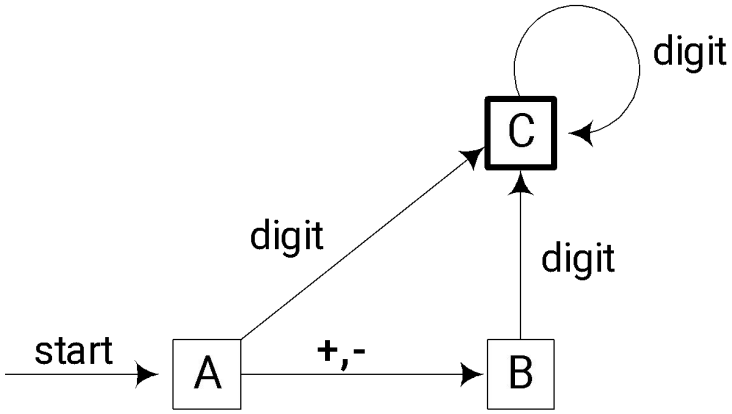
Isdigit



```
Isdigit PROC
    cmp al,'0'
    jb    L1
    cmp al,'9'
    ja L1
    test ax,0
L1: ret
Isdigit ENDP
```

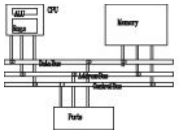
Your turn

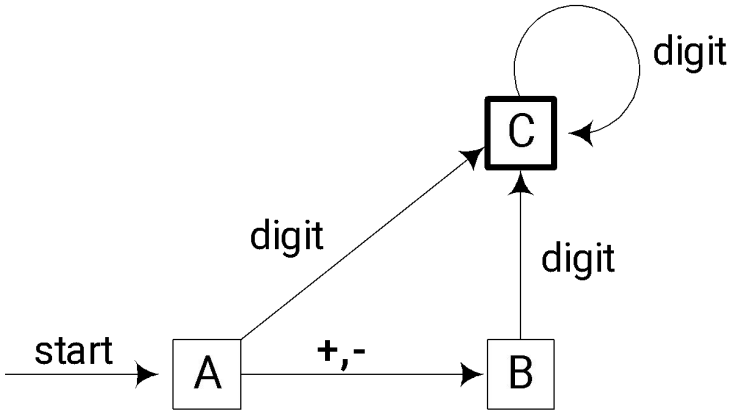





```
StateB:
  call  Getnext       ; read next char into AL
  call  Isdigit       ; ZF = 1 if AL is a digit
  jz    StateC
  call  DisplayErrorMsg  ; invalid input found
  jmp   Quit
```

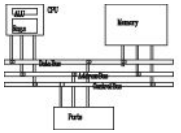
Implementing an FSM






```
StateC:
  call  Getnext    ; read next char into AL
  jz    Quit       ; quit if Enter pressed
  call  Isdigit    ; ZF = 1 if AL is digit
  jz    StateC
  cmp   AL,ENTER_KEY ; Enter key pressed?
  je    Quit; yes: quit
  call  DisplayErrorMsg ; no: invalid input
  jmp   Quit
```

Finite-state machine example

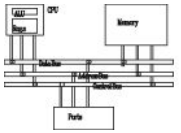


- [*sign*]*integer*.[*integer*][*exponent*]

  sign → {+|-}

  exponent → E[{+|-}]integer

High-level directives



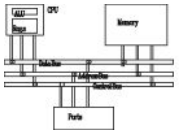
- .IF, .ELSE, .ELSEIF, and .ENDIF can be used to create block-structured IF statements.
- Examples:

```
.IF eax > ebx
   mov edx,1
.ELSE
   mov edx,2
.ENDIF
```

```
.IF eax > ebx && eax > ecx
   mov edx,1
.ELSE
   mov edx,2
.ENDIF
```

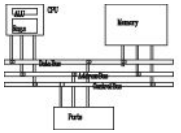
- MASM generates "hidden" code for you, consisting of code labels, CMP and conditional jump instructions.

Relational and logical operators



Operator	Description
expr1 == *expr2*	Returns true when *expression1* is equal to *expr2*.
expr1 != *expr2*	Returns true when *expr1* is not equal to *expr2*.
expr1 > *expr2*	Returns true when *expr1* is greater than *expr2*.
expr1 >= *expr2*	Returns true when *expr1* is greater than or equal to *expr2*.
expr1 < *expr2*	Returns true when *expr1* is less than *expr2*.
expr1 <= *expr2*	Returns true when *expr1* is less than or equal to *expr2*.
! *expr*	Returns true when *expr* is false.
expr1 && *expr2*	Performs logical AND between *expr1* and *expr2*.
expr1 || *expr2*	Performs logical OR between *expr1* and *expr2*.
expr1 & *expr2*	Performs bitwise AND between *expr1* and *expr2*.
CARRY?	Returns true if the Carry flag is set.
OVERFLOW?	Returns true if the Overflow flag is set.
PARITY?	Returns true if the Parity flag is set.
SIGN?	Returns true if the Sign flag is set.
ZERO?	Returns true if the Zero flag is set.

MASM-generated Code



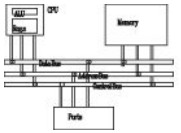
```
.data
val1   DWORD 5
result DWORD ?
.code
mov eax,6
.IF eax > val1
mov result,1
.ENDIF
```

Generated code:

```
    mov eax,6
    cmp eax,val1
    jbe @C0001
    mov result,1
@C0001:
```

MASM automatically generates an unsigned jump (JBE).

.REPEAT directive

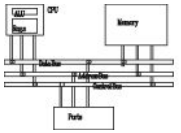


Executes the loop body before testing the loop condition associated with the .UNTIL directive.

Example:

```
; Display integers 1 - 10:

mov eax,0
.REPEAT
    inc eax
    call WriteDec
    call Crlf
.UNTIL eax == 10
```

.WHILE directive



Tests the loop condition before executing the loop body The .ENDW directive marks the end of the loop.

Example:

```
; Display integers 1 - 10:

mov eax,0
.WHILE eax < 10
    inc eax
    call WriteDec
    call Crlf
.ENDW
```