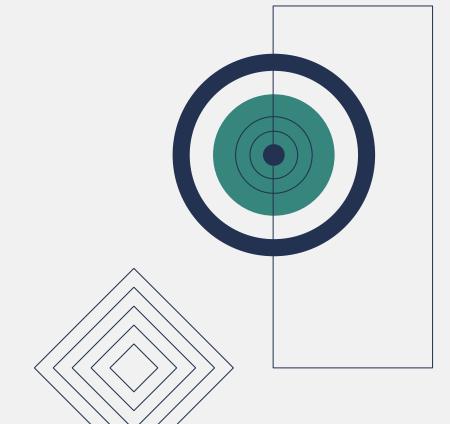


National University of Computer & Emerging Sciences

Incorporating Parallelism

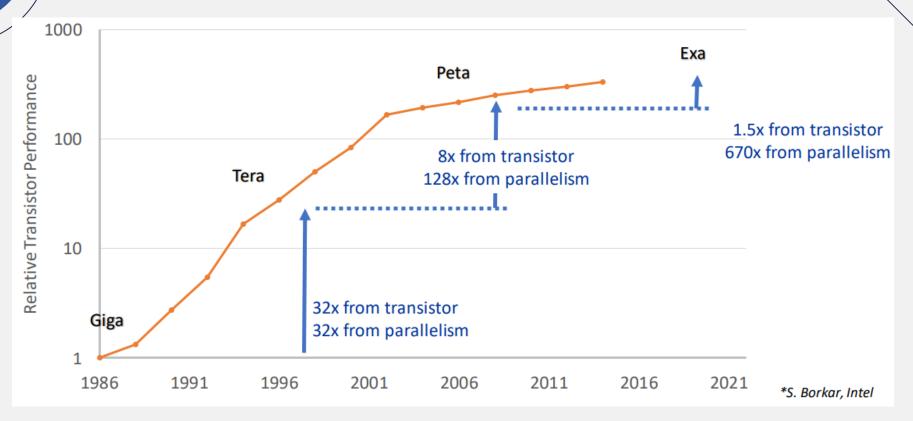


What is Parallelism

- Parallelism refers to the simultaneous execution of multiple tasks or instructions.
- Parallelism can be achieved in various ways, allowing multiple operations to be performed concurrently.

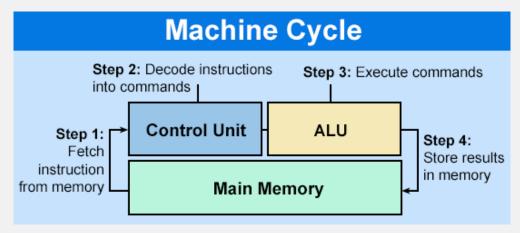


Performance from Parallelism



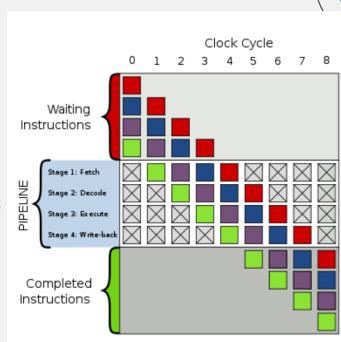
Instruction-Level Parallelism

- Involves executing multiple instructions (called Instruction Pipelining) at the same time.
- Instruction pipelining can reduce idle time in the processor
- Consider a machine cycle (having four-stage) of an Intel processor to execute one instruction.



Instruction-Level Parallelism

- This is a generic pipeline with four stages: Fetch, Decode, Execute, Writeback.
- At the top are the list of instructions waiting to be executed
- At the bottom are the list of instructions that have been completed
- In the middle are the instructions ins the pipeline.
- Pipelining improves the performance of a processor by overlapping the
 execution of multiple instructions.



Without ILP Example

Cycle	Instruction Fetch (IF)	Instruction Decode (ID)	Instruction Execute (EX)	Instruction Write- back (WB)		
1	I1 IF					
2		I1 ID				
3			I1 EX			
4				I1 WB		
5	12 IF					
6		I2 ID				
7			I2 EX			
8				I2 WB		
9	13 IF					
10		I3 ID				
11			I3 EX			
12				I3 WB		
13	14 IF					
14		I4 ID				
15			I4 EX			

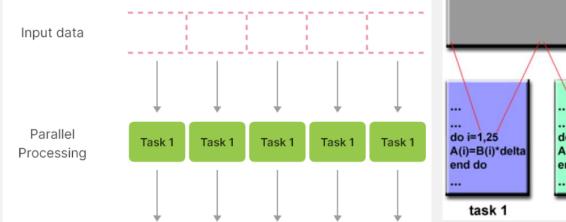
With ILP Example

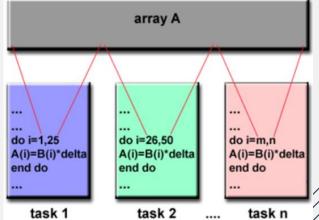
Cycle	Instruction Fetch (IF)	Instruction Decode (ID)	Instruction Execute (EX)	Instruction Write- back (WB)
1	I1 IF			
2	I2 IF	I1 ID		
3	13 IF	I2 ID	I1 EX	
4	14 IF	I3 ID	I2 EX	I1 WB
5		I4 ID	I3 EX	I2 WB
6			I4 EX	I3 WB
7				I4 WB

Data Parallelism

Threads involved in doing similar things, but at different locations.

- The data set is organized into a common structure, such as an array.
- A set of tasks work collectively on that structure; however, each task works on a different region.
- Tasks perform the same operation on their portion of data, for example, "multiply every array element by some value".









- Involves dividing a program into smaller tasks that can be executed independently.
- program:
- ...
- if CPU="a" then do task "A"
- else if CPU="b" then do task "B"
- end if
- . . .
- end program
- Task Parallelism means concurrent execution of the different task on multiple computing cores or threads using the same or different data.

Threads involved in doing different things.



Task-Level Parallelism (TLP)

- Code executed by CPU "a":
- program:
- do task "A"
- end program
- Code executed by CPU "b":
- program:
- do task "B"
- end program
- In the general case, different execution core/threads communicate with one another as they work.



- This type of parallelism is effective when the nature of the work doesn't require strict predictability in the decomposition process.
- Load balancing becomes crucial in task parallelism to ensure that tasks are evenly distributed across the available processing resources.
- Examples:
 - Scientific simulations, where different parts of a simulation can be executed concurrently.

Parallel vs Distributed Systems

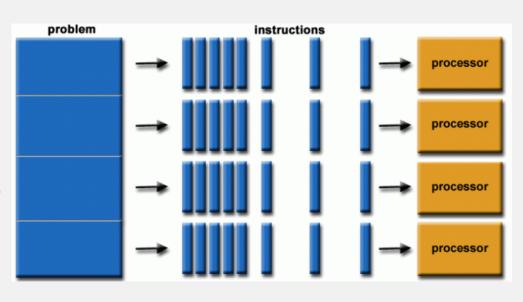


Motivation

- Parallel computing and distributed computing are two fundamental paradigms used to improve
 - computational speed,
 - efficiency, and
 - o scalability.
- Both approaches involve multiple processors working together, but they differ in architecture, communication, and execution strategies.

Parallel Computing Systems

 Parallel computing refers to the simultaneous execution of multiple tasks using multiple processing elements within a single system to solve a problem faster.



Key Characteristics

- Uses multiple processors (or cores) within a single system.
- Shared memory or distributed memory.
- Tasks are divided into subtasks, which execute concurrently.
- Low-latency communication.
- Requires tight synchronization among processors.
- **Example Systems** (e.g., Multicore Systems, high-performance computers, supercomputers etc).



Shared Memory Architecture

- All processors have access to a common memory.
- Uses synchronization mechanisms (e.g., locks, semaphores).

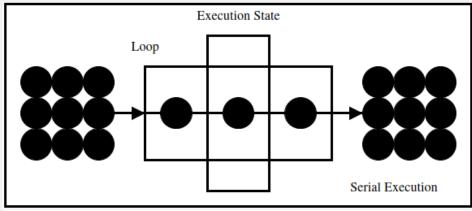
Distributed Memory Architecture

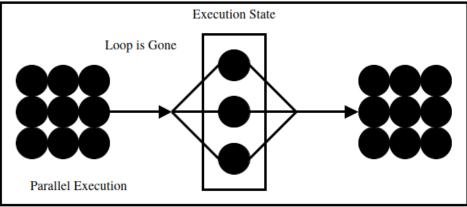
- Each processor has its own local memory.
- Uses message passing for inter-processor communication.

Hybrid Architecture

Combination of shared and distributed memory models.

Serial VS parallel Execution





Challenges in Parallel Systems

Synchronization Issues

- Ensuring correct execution order of parallel tasks.
- Dealing with race conditions by using locks and barriers.

Load Balancing

- Efficiently distributing tasks across multiple processors.
- Dynamic vs. static load balancing strategies.

Memory Management

- Managing shared memory access efficiently.
- Avoiding bottlenecks in memory bandwidth.

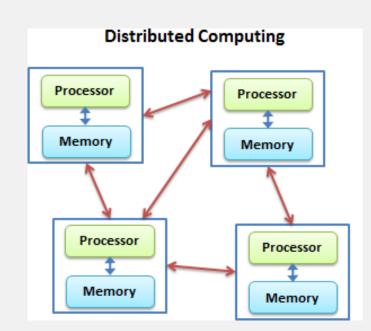
Programming Complexity

- Writing efficient parallel algorithms.
- Adapting sequential programs to parallel execution.



 Distributed computing involves multiple independent computers (nodes) working together over a network to achieve a common goal.

 Goal: Scalability, fault tolerance, and distributed workload.



Key Characteristics

Multiple autonomous computers (nodes) connected via a network.

• Each node has its **own memory and processing** power.

Tasks are distributed and executed independently.

• Communication occurs via message passing.

• Used for scalability, fault tolerance, and decentralized control.

Types of Distributed Computing

Cluster Computing

 A group of tightly coupled homogeneous computers acting as a single system. e.g. Google Cloud Cluster, AWS Clusters.

Grid Computing

 Uses geographically distributed computers (same/different) connected via a network. Example: SETI@Home, CERN Grid.

Cloud Computing

 Provides scalable computing resources on demand via the internet. e.g. Amazon AWS, Microsoft Azure, Google Cloud.

Challenges in Distributed Systems

Network Latency and Bandwidth

- Communication delays due to network overhead.
- Limited bandwidth causing bottlenecks.

Fault Tolerance and Reliability

Handling node failures and system crashes.

Data Consistency

Ensuring data consistency across distributed nodes.

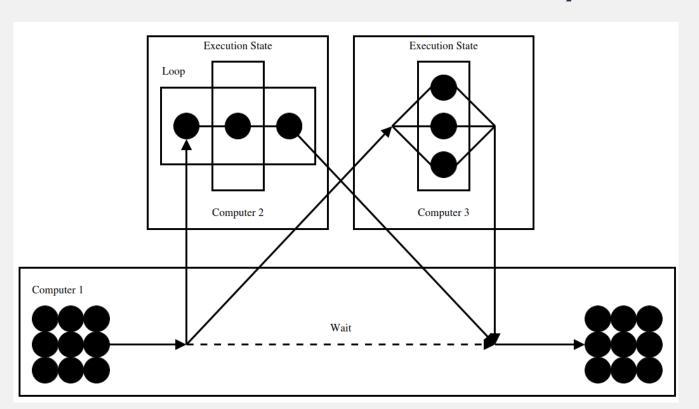
Security and Privacy Concerns

Protecting data during transmission and storage.

Resource Management

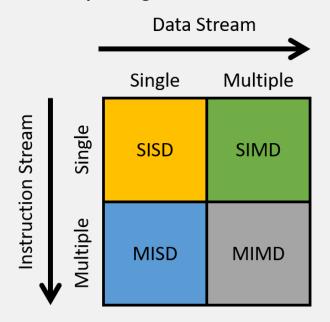
Efficient allocation of computing and storage resources.

Execution on a Distributed System



Flynn's Taxonomy of Architecture

 Flynn's taxonomy is commonly used to describe the ecosystem of modern computing architecture

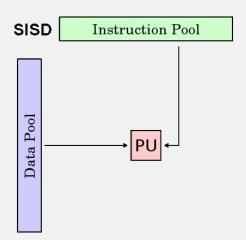


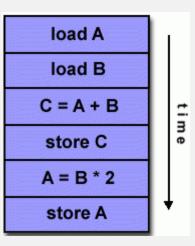
A **stream** in this context is a flow of data or instructions.

A **single stream** issues one element per time unit, similar to a queue. In contrast, **multiple streams** typically issue many elements per time unit (think of multiple queues).



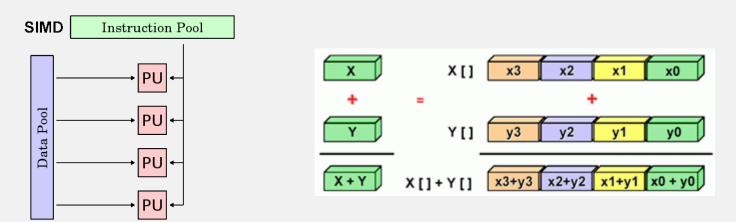
- A serial (non-parallel) computer
- A single compute unit processing a single instruction at a time. Most commercially available processors prior to the mid-2000s were SISD machines.





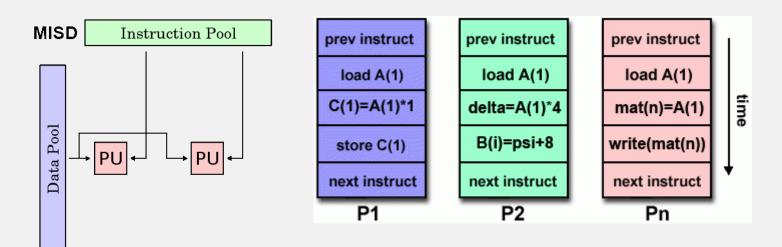
Single Instruction, Multiple Data (SIMD)

- A type of parallel computer, which execute the same instruction on multiple data simultaneously in lockstep fashion.
- All instructions are placed into a queue, while data is distributed among different compute units. The most well-known example of the SIMD architecture is the graphics processing unit.



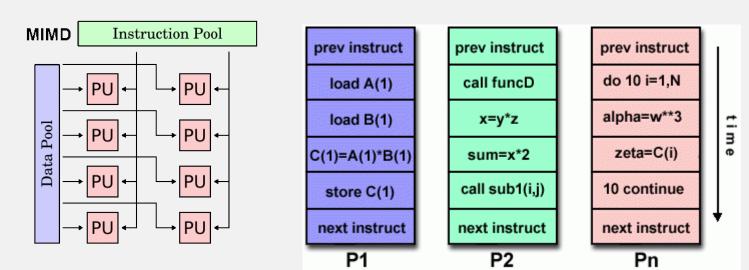
Multiple Instruction, Single Data (MISD)

 A type of parallel computer which have multiple instruction units performing on a single data stream. MISD systems were typically designed for incorporating fault tolerance in mission-critical systems, such as the flight control programs for NASA shuttles.



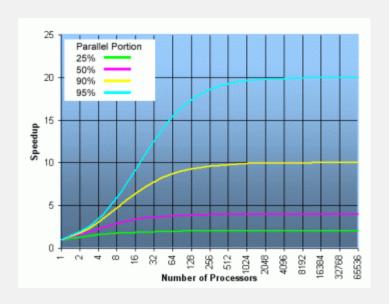
Multiple Instruction, Multiple Data (MIMD)

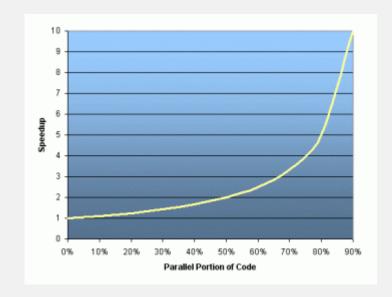
 A type of parallel computer which can work on multiple instructions or multiple data streams. Since nearly all modern computers use multicore CPUs, most are classified as MIMD machines.



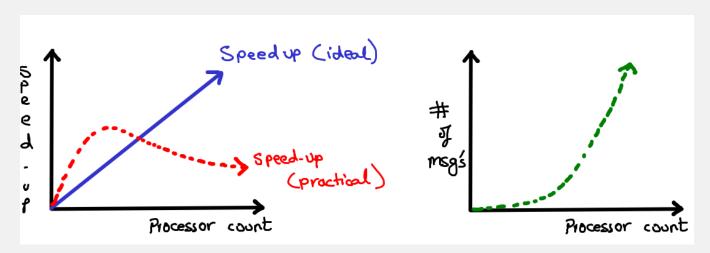
Amdahl's Law

 Amdahl's Law states that potential program speedup is defined by the fraction of code that can be parallelized (P).





Ideal VS Practical Speed Up



- Why ideal speedup is not possible:
- data transfer (through message exchanges), I/O bottlenecks, race conditions, dependencies, load balancing, deadlocks, synchronization, node failures, lazy programmers, . . .

Amdahl's Law

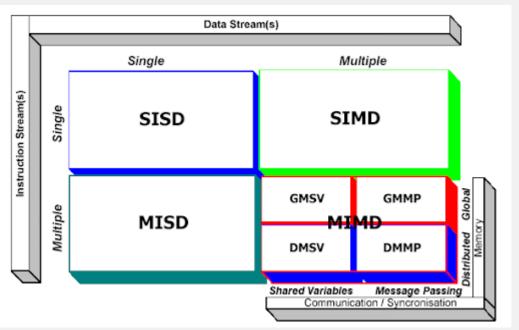
Speedup = 1/(1-p)

- if none of the code can be parallelized, P = 0 and the speedup = 1 (no speedup).
- If all of the code is parallelized, P = 1 and the speedup is infinite (in theory).
- If 50% of the code can be parallelized, maximum speedup = 2, meaning the code will run twice as fast.
- Introducing the number of processors performing the parallel fraction of work, the relationship can be modeled by:

where P = parallel fraction, N = number of processors and S = serial fraction.

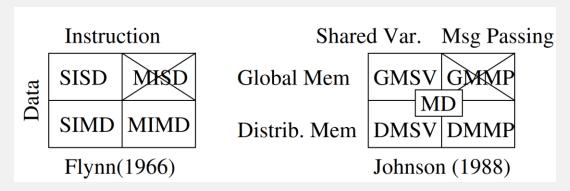
Johnson Taxonomy (Classification)

• Johnson further refined the MIMD model by using the memory system and communications as criteria:



Johnson Taxonomy (Classification)

- Shared Memory MIMD (tightly coupled multiprocessor systems)
 - GMSV (Global/Shared Memory, Shared Variables)
 - GMMP (Global/Shared Memory, Message Passing)
- **Distributed Memory MIMD** (loosely coupled multicomputer systems)
 - DMSV (Distributed Memory, Shared Variable)
 - DMMP (Distributed Memory, Message Passing)



Shared-Variable vs Message-Passing

Shared-Variable Technique

- Common variables are shared among threads/cores
- Synchronization becomes very important between threads.
- Threads need to work together but not interfering

Message-Passing Technique

- In the case of Distributed Computing System there is no concept of shared memory.
- In this model, the only way to communicate information with another node is to send data over the network. We say that nodes are passing messages to each other.
- Unlike threads, synchronization happens automatically as part of the message passing process.

Johnson Taxonomy

- Global Memory-Shared Variables (GMSV) -- Shared memory machines
- Distributed Memory-Message Passing (DMMP) -- Message passing machines
- Distributed Memory-Shared Variables (DMSV) -- Hybrid machines Memory is physically distributed but appears logically shared

Johnson's Taxonomy

- **GMSV** Symmetric Multi-Processors (OpenMP, PThreads)
- **DMMP** Distributed systems (If simulated: Omnet++, NS2/3, etc.), NUMA based Cluster systems (OpenMPI, RPC), GPU Computing (OpenCL, CUDA), HPC based systems
- **DMSV** Cloud Model: Memory is distributed but access is through same address space (Hadoop, MapReduce)

Applications and Examples

General Purpose

- Word processor (spell check, auto save, . . .)
- Spreadsheet (automatic background recalculations after cell edits)
- Operating system (User threads, Kernel threads, . . .)
- Server (multi client handling
- GUI's (Event listeners)
- Video games (rendering of animation, taking care of physics,
 AI)
- Web browser (tabs, multi-segment download accelerators, . . .)
- O ...

Applications and Examples

Scientific Purposes

- Mathematical problems (Ordinary differential equations, partial differential equations, linear algebra, numerical analysis)
- Climate modelling and weather prediction
- Fluid dynamics (video games, simulation of dam bursts, tsunamies, etc.)
- Computational Biology, Finance, Physics, Chemistry
- Machine learning, statistical methods
- Image processing (medical images, spectral images, satellite images)
- Finding Aliens, Finding cure for Alzheimers, Cancers, . . .