

Banker's Algorithm

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Example

- **Considering a system with five processes T_0 through T_4 and three resources of type A, B, C.**
- **Resource type A has 10 instances, B has 5 instances and type C has 7 instances.**
- **Suppose at time t_0 following snapshot of the system has been taken**

Example of Banker's Algorithm

- 5 threads T_0 through T_4 ;
3 resource types:
 A (10 instances), B (5 instances), and C (7 instances)
- Snapshot at time T_0 :

	<u>Allocation</u>	<u>Max</u>	<u>Available</u>
	$A \ B \ C$	$A \ B \ C$	$A \ B \ C$
T_0	0 1 0	7 5 3	3 3 2
T_1	2 0 0	3 2 2	
T_2	3 0 2	9 0 2	
T_3	2 1 1	2 2 2	
T_4	0 0 2	4 3 3	

Need Matrix

- The content of the matrix ***Need*** is defined to be ***Max – Allocation***

	<u><i>Need</i></u>
	<i>A B C</i>
T_0	7 4 3
T_1	1 2 2
T_2	6 0 0
T_3	0 1 1
T_4	4 3 1

Safety Algo - 2

- For $i=1$, $\text{Need}_1 = (1,2,2)$
- $\text{Finish}[1] = \text{false} \ \&\& \ \text{Need}_1 (1,2,2) \leq \text{Work} (3,3,2)$
- So T_1 is kept in safe sequence

	<u>Need</u>		
	A	B	C
T_0	7	4	3
T_1	1	2	2
T_2	6	0	0
T_3	0	1	1
T_4	4	3	1

Safety Algo - 3

- $Work = Work + Allocation_1$
- $Work = (3,3,2) + (2,0,0) = (5,3,2)$
- $Finish[5] = \{false, true, false, false, false\}$

	<u>Need</u>		
	A	B	C
T_0	7	4	3
T_1	1	2	2
T_2	6	0	0
T_3	0	1	1
T_4	4	3	1

	<u>Allocation</u>			<u>Max</u>			<u>Available</u>		
	A	B	C	A	B	C	A	B	C
T_0	0	1	0	7	5	3	3	3	2
T_1	2	0	0	3	2	2			
T_2	3	0	2	9	0	2			
T_3	2	1	1	2	2	2			
T_4	0	0	2	4	3	3			

Safety Algo - 4

- For $i=2$, $\text{Need}_2 = (6,0,0)$
- $\text{Finish}[2] = \text{false} \ \&\& \ \text{Need}_2 (6,0,0) > \text{Work} (5,3,2)$
- So T_2 must wait

	<u>Need</u>		
	A	B	C
T_0	7	4	3
T_1	1	2	2
T_2	6	0	0
T_3	0	1	1
T_4	4	3	1

	<u>Allocation</u>			<u>Max</u>			<u>Available</u>		
	A	B	C	A	B	C	A	B	C
T_0	0	1	0	7	5	3	3	3	2
T_1	2	0	0	3	2	2			
T_2	3	0	2	9	0	2			
T_3	2	1	1	2	2	2			
T_4	0	0	2	4	3	3			

Safety Algo - 5

- For $i=3$, $\text{Need}_3 = (0,1,1)$
- $\text{Finish}[3] = \text{false} \ \&\& \ \text{Need}_3 (0,1,1) \leq \text{Work} (5,3,2)$
- So T_3 is kept in safe sequence

	<u>Need</u>		
	A	B	C
T_0	7	4	3
T_1	1	2	2
T_2	6	0	0
T_3	0	1	1
T_4	4	3	1

	<u>Allocation</u>			<u>Max</u>			<u>Available</u>		
	A	B	C	A	B	C	A	B	C
T_0	0	1	0	7	5	3	3	3	2
T_1	2	0	0	3	2	2			
T_2	3	0	2	9	0	2			
T_3	2	1	1	2	2	2			
T_4	0	0	2	4	3	3			

Safety Algo - 6

- $Work = Work + Allocation_3$
- $Work = (5,3,2) + (2,1,1) = (7,4,3)$
- $Finish[5] = \{false, true, false, true, false\}$

	<u>Need</u>		
	A	B	C
T_0	7	4	3
T_1	1	2	2
T_2	6	0	0
T_3	0	1	1
T_4	4	3	1

	<u>Allocation</u>			<u>Max</u>			<u>Available</u>		
	A	B	C	A	B	C	A	B	C
T_0	0	1	0	7	5	3	3	3	2
T_1	2	0	0	3	2	2			
T_2	3	0	2	9	0	2			
T_3	2	1	1	2	2	2			
T_4	0	0	2	4	3	3			

Safety Algo - 7

- For $i=4$, $\text{Need}_3 = (4,3,1)$
- $\text{Finish}[4] = \text{false} \ \&\& \ \text{Need}_4 (4,3,1) \leq \text{Work} (7,4,3)$
- So T_4 is kept in safe sequence

	<u>Need</u>		
	A	B	C
T_0	7	4	3
T_1	1	2	2
T_2	6	0	0
T_3	0	1	1
T_4	4	3	1

	<u>Allocation</u>			<u>Max</u>			<u>Available</u>		
	A	B	C	A	B	C	A	B	C
T_0	0	1	0	7	5	3	3	3	2
T_1	2	0	0	3	2	2			
T_2	3	0	2	9	0	2			
T_3	2	1	1	2	2	2			
T_4	0	0	2	4	3	3			

Safety Algo - 8

- $Work = Work + Allocation_4$
- $Work = (7,4,3) + (0,0,2) = (7,4,5)$
- $Finish[5] = \{false, true, false, true, true\}$

	<u>Need</u>			
	A	B	C	
T_0	7	4	3	
T_1	1	2	2	
T_2	6	0	0	
T_3	0	1	1	
T_4	4	3	1	

	<u>Allocation</u>			<u>Max</u>			<u>Available</u>		
	A	B	C	A	B	C	A	B	C
T_0	0	1	0	7	5	3	3	3	2
T_1	2	0	0	3	2	2			
T_2	3	0	2	9	0	2			
T_3	2	1	1	2	2	2			
T_4	0	0	2	4	3	3			

Safety Algo - 9

- For $i=0$, $\text{Need}_0 = (7,4,3)$
- $\text{Finish}[0] = \text{false} \ \&\& \ \text{Need}_0 (7,4,3) \leq \text{Work} (7,4,5)$
- So T_0 is kept in safe sequence

	<u>Need</u>		
	A	B	C
T_0	7	4	3
T_1	1	2	2
T_2	6	0	0
T_3	0	1	1
T_4	4	3	1

	<u>Allocation</u>			<u>Max</u>	<u>Available</u>
	A	B	C	A B C	A B C
T_0	0	1	0	7 5 3	3 3 2
T_1	2	0	0	3 2 2	
T_2	3	0	2	9 0 2	
T_3	2	1	1	2 2 2	
T_4	0	0	2	4 3 3	

Safety Algo - 10

- $Work = Work + Allocation_0$
- $Work = (7,4,5) + (0,1,0) = (7,5,5)$
- $Finish[5] = \{true, true, false, true, true\}$

	<u>Need</u>		
	A	B	C
T_0	7	4	3
T_1	1	2	2
T_2	6	0	0
T_3	0	1	1
T_4	4	3	1

	<u>Allocation</u>			<u>Max</u>			<u>Available</u>		
	A	B	C	A	B	C	A	B	C
T_0	0	1	0	7	5	3	3	3	2
T_1	2	0	0	3	2	2			
T_2	3	0	2	9	0	2			
T_3	2	1	1	2	2	2			
T_4	0	0	2	4	3	3			

Safety Algo - 11

- For $i=2$, $Need_2 = (6,0,0)$
- $Finish[2] = \text{false} \ \&\& \ Need_2 (6,0,0) \leq Work (7,5,5)$
- So T_2 is kept in safe sequence

	<u>Need</u>		
	A	B	C
T_0	7	4	3
T_1	1	2	2
T_2	6	0	0
T_3	0	1	1
T_4	4	3	1

	<u>Allocation</u>	<u>Max</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 5 3	3 3 2
T_1	2 0 0	3 2 2	
T_2	3 0 2	9 0 2	
T_3	2 1 1	2 2 2	
T_4	0 0 2	4 3 3	

Safety Algo - 12

- $Work = Work + Allocation_2$
- $Work = (7,5,5) + (3,0,2) = (10,5,7)$
- $Finish[5] = \{true, true, true, true, true\}$

	<u>Need</u>		
	A	B	C
T_0	7	4	3
T_1	1	2	2
T_2	6	0	0
T_3	0	1	1
T_4	4	3	1

	<u>Allocation</u>			<u>Max</u>			<u>Available</u>		
	A	B	C	A	B	C	A	B	C
T_0	0	1	0	7	5	3	3	3	2
T_1	2	0	0	3	2	2			
T_2	3	0	2	9	0	2			
T_3	2	1	1	2	2	2			
T_4	0	0	2	4	3	3			

Safe Sequence

- The system is in a safe state since the sequence $\langle T_1, T_3, T_4, T_2, T_0 \rangle$ satisfies safety criteria

T_1 Request (1,0,2)

T_1 checks

Step1: Request₁ ≤ Need₁ i.e., (1,0,2) ≤ (1,2,2) and

Step2: Request₁ ≤ Available i.e., (1,0,2) ≤ (3,3,2)

According to Step3 in Resource-Request Algorithm for Process T_i

Available = Available – Request_i;

Allocation_i = Allocation_i + Request_i;

Need_i = Need_i – Request_i;

We have

(3,3,2) – (1,0,2) = (2,3,0) Available

(2,0,0) + (1,0,2) = (3,0,2) Allocation

(1,2,2) – (1,0,2) = (0,2,0) Need

T_1 Request (1,0,2)

- Determine whether this new state is safe. Execute Safety Algo

Safety Algo - 1

- $m=3, n=5$
- $Work = Available = (2,3,0)$
- $Finish[5] = \{false, false, false, false, false\}$
- For $i=0$, $Need_0 = (7,4,3)$
- $Finish[0] = false \ \&\& \ Need_0 (7,4,3) > Work (2,3,0)$
- So T_0 has to wait

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 2

- For $i=1$, $\text{Need}_1 = (0,2,0)$
- $\text{Finish}[1] = \text{false} \ \&\& \ \text{Need}_1 (0,2,0) \leq \text{Work} (2,3,0)$
- So T_1 is kept in safe sequence

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 3

- $Work = Work + Allocation_1$
- $Work = (2,3,0) + (3,0,2) = (5,3,2)$
- $Finish[5] = \{false, true, false, false, false\}$

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 4

- For $i=2$, $\text{Need}_2 = (6,0,0)$
- $\text{Finish}[2] = \text{false}$ && $\text{Need}_2 (6,0,0) > \text{Work} (5,3,2)$
- So T_2 must wait

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 5

- For $i=3$, $\text{Need}_3 = (0,1,1)$
- $\text{Finish}[3] = \text{false} \ \&\& \ \text{Need}_3 (0,1,1) \leq \text{Work} (5,3,2)$
- So T_3 is kept in safe sequence

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 6

- $Work = Work + Allocation_3$
- $Work = (5,3,2) + (2,1,1) = (7,4,3)$
- $Finish[5] = \{false, true, false, true, false\}$

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 7

- For $i=4$, $\text{Need}_4 = (4,3,1)$
- $\text{Finish}[4] = \text{false} \ \&\& \ \text{Need}_4 (4,3,1) \leq \text{Work} (7,4,3)$
- So T_4 is kept in safe sequence

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 8

- $Work = Work + Allocation_4$
- $Work = (7,4,3) + (0,0,2) = (7,4,5)$
- $Finish[5] = \{false, true, false, true, true\}$

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 9

- For $i=0$, $\text{Need}_0 = (7,4,3)$
- $\text{Finish}[0] = \text{false}$ && $\text{Need}_0 (7,4,3) \leq \text{Work} (7,4,5)$
- So T_0 is kept in safe sequence

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	<i>A B C</i>	<i>A B C</i>	<i>A B C</i>
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 10

- $Work = Work + Allocation_0$
- $Work = (7,4,5) + (0,1,0) = (7,5,5)$
- $Finish[5] = \{true, true, false, true, true\}$

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 11

- For $i=2$, $\text{Need}_2 = (6,0,0)$
- $\text{Finish}[2] = \text{false} \ \&\& \ \text{Need}_2 (6,0,0) \leq \text{Work} (7,5,5)$
- So T_2 is kept in safe sequence

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safety Algo - 12

- $Work = Work + Allocation_2$
- $Work = (7,5,5) + (3,0,2) = (10,5,7)$
- $Finish[5] = \{true, true, true, true, true\}$

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
T_0	0 1 0	7 4 3	2 3 0
T_1	3 0 2	0 2 0	
T_2	3 0 2	6 0 0	
T_3	2 1 1	0 1 1	
T_4	0 0 2	4 3 1	

Safe Sequence

- The system is in a safe state since the sequence $\langle T_1, T_3, T_4, T_0, T_2 \rangle$ satisfies safety criteria