## Lecture 9: Cryptography

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## Euclidean algorithms

Assume  $r_0 = a$  and  $r_1 = b$  with  $r_0 \ge r_1$ .

Euclidean finds gcd(a, b)

• 
$$r_0 = r_1 q_1 + r_2$$
:

$$r_{n-2} = r_{n-1}q_{n-1} + r_n$$

• 
$$r_{n-1} = r_n q_n + 0$$

Conclusion: 
$$gcd(a, b) = r_n$$

Extended Euclidean:  $as + bt = \gcd(a, b)$ Notation: \* = some number

• 
$$r_n = r_{n-2} + r_{n-1} *$$

• Use prev. equation to find  $r_{n-1}$ 

$$r_n = r_{n-2} + (r_{n-3} * + r_{n-2} *)(-q_{n-1})$$
  
 $r_n = r_{n-3} * + r_{n-2} *$ 

Keep doing this till first equation

$$r_n = r_0 s + r_1 t$$

### Modular inverses

$$a^{-1} \mod m =$$
 an integer  $b$  with  $ab \equiv 1 \pmod m$ . Note that 
$$a^{-1} \mod m \text{ exists} \Leftrightarrow \gcd(a,m) = 1$$

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① Compute gcd(a, m). If gcd(a, m) > 1,  $a^{-1} \mod m$  doesn't exist. If gcd(a, m) = 1, proceed to the next step.

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- ① Compute gcd(a, m). If gcd(a, m) > 1,  $a^{-1} \mod m$  doesn't exist. If gcd(a, m) = 1, proceed to the next step.
- 2 Find Bezout coefficients s and t of a and m:

$$as + mt = \gcd(a, m) = 1.$$

#### Conclusion:

$$a^{-1} \mod m = s$$
.



# Convert negative to positive in congruence

- In finding  $s = a^{-1} \mod m$  by Bezout coefficient, s might be negative. To convert s to positive, we add a suitable multiple of m.
- Examples

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- Examples

Since 
$$3 \cdot (-3) + 2 \cdot 5 = 1$$
, 
$$3^{-1} \mod 5 = (-3) \mod 5 = 2$$
 Since  $5 \cdot (-7) + 9 \times 4 = 1$ , 
$$5^{-1} \mod 9 = (-7) \mod 9 = 2$$

# Solving $ax \equiv b \pmod{m}$

- Find  $d = \gcd(a, m)$  (by factorizing a and m)
  - If  $d \nmid b$ , the equation has no solution.
  - If  $d \mid b$ , proceed to the next step.
- Write  $a = da_1, b = db_1, m = dm_1$ . Dividing all terms of  $ax \equiv b \pmod{m}$  by d, we obtain

$$a_1 x \equiv b_1 \pmod{m_1} \tag{1}$$

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**3** Multiplying both sides of (1) by  $a_2 = a_1^{-1} \mod m_1$ , we obtain

$$x \equiv b_1 a_2 \pmod{m_1}$$
.



#### Exercise 1

Let a,b,m be integers with  $\gcd(a,m) \nmid b$ . Prove that the equation

$$ax \equiv b \pmod{m}$$

has no solution.

#### Exercise 2

Consider integers a=252 and m=356.

(a) Let s,t be Bezout coefficients of a,m. What is s?

(A) -24 (B) 332 (C) -380 (D) Any A,B,C (E) None of A,B,C

- (b) For what  $b \in \mathbb{Z}$  does equation  $252x \equiv b \pmod{356}$  have solution?
- (c) Solve  $252x \equiv 12 \pmod{356}$ ?

(d) How many  $x \in \{0, 1, ..., 355\}$  satisfies  $252x \equiv 12 \pmod{356}$ ?

# What is cryptography?

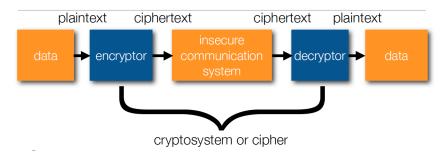
- The subject of transforming information so that it cannot be easily recovered without special knowledge
- Cryptography is a branch in cryptology which comprises of
  - **① Cryptography** How to design methods to hide information.
  - **2** Cryptanalysis How to break methods that hide information.

## Applications of cryptography

- Secrecy in communications: Military, spies, diplomats, banking (ATM cards, credit cards, PayPal)
- Integrity protection (ability to detect change in messages):
   Electronic form submission.

# Cryptography model

#### Model



### Conventional notations

- The plaintext is denoted by P
- The ciphertext is denoted by C
- Encryption

The sender encrypts  $P \to C$  and send C over the insecure channel.

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Decryption

Upon receiving C, the receiver decrypts  $C \to P$  to retrieve the information.

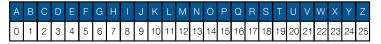
## Caesar cipher: Encryption

• Identify each letter with an integer between 0 and 25. Write out the plaintext P after this identification.

```
A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
O 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25
```

## Caesar cipher: Encryption

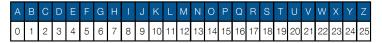
• Identify each letter with an integer between 0 and 25. Write out the plaintext P after this identification.



② Fix a key  $k \in \mathbb{Z}_{26} = \{0, 1, ..., 25\}$  (kept secret)

## Caesar cipher: Encryption

• Identify each letter with an integer between 0 and 25. Write out the plaintext P after this identification.



- ② Fix a key  $k \in \mathbb{Z}_{26} = \{0, 1, ..., 25\}$  (kept secret)
- lacktriangledown The plaintext P is encrypted to a cipher text C by

$$C = P + k \mod 26$$
.

Convert C back to a string of letters.



# Caesar cipher: Decryption

• Since  $C = P + k \mod 26$ , we have

$$P = C - k \mod 26$$
.

## Example 1

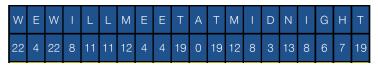
Encrypt the following text using Caesar cipher with the key k=11.

WEWILLMEETATMIDNIGHT.



## Example 1 solution

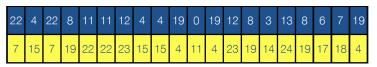
• Convert the plaintext into integers between 0 and 25.



$$P = (22, 4, 22, 8, 11, 11, 12, 4, 4, 19, 0, 19, 12, 8, 3, 13, 8, 6, 7, 19)$$

# Example 1 solution

• Add each integer by 11 and compute the result modulo 26.

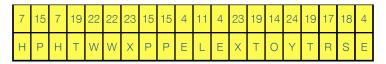


$$C = P + 11 \mod 26 = (7, 15, 7, 19, 22, 22, 23, 5, 15, 4, 11, \dots)$$



## Example 1 solution

#### Convert the integers back to letters

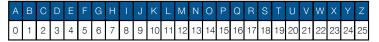


**HPHTWWXPPELEXTOYTRSE** 

#### Encryption: $P \rightarrow C$

lacksquare Identify each letter with an integer between 0 and 25.

Write out the plaintext P after this identification.



- **2** Fix a key  $k \in \mathbb{Z}_{26} = \{0, 1, \dots, 25\}.$
- lacksquare The plaintext P is encrypted to a cipher text C by

$$C = P + k \mod 26$$
.

**Decryption:**  $C \rightarrow P$ 

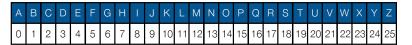
$$P = C - k \mod 26$$
.



### Exercise 3

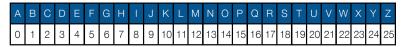
Using the Caesar cipher with the key k=6,  $\mbox{encrypt}$  the plaintext

#### **SITISGREAT**



#### Exercise 4

**Decrypt** the ciphertext ZNK KGXRE HOXJ MKZY ZNK CUXS using Caesar cipher with the key k=6.



## How to break Caesar cipher?

The Caesar cipher is easy to break, as the key space is small.

- There are 26 possible values  $0, 1, \ldots, 25$  for the key k.
- 2 Try all possible values.

## Affine cipher: Encryption

Identify each letter with an integer between 0 and 25.
 Write out the plaintext P after this identification.

### Affine cipher: Encryption

- Identify each letter with an integer between 0 and 25. Write out the plaintext P after this identification.
- ② Choose  $a, b \in \mathbb{Z}_{26}$  such that gcd(a, 26) = 1.

$$a \in \{1, 3, 5, 7, 9, 11, 15, 17, 19, 21, 23, 25\}$$

## Affine cipher: Encryption

- Identify each letter with an integer between 0 and 25.
   Write out the plaintext P after this identification.
- ② Choose  $a, b \in \mathbb{Z}_{26}$  such that gcd(a, 26) = 1.

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lacktriangledown The plaintext P is encrypted to a ciphertext C by

$$C = aP + b \mod 26$$
.



## Affine cipher: Decryption

• Decryption: Since  $C = aP + b \mod 26$ , we have

$$P = a^{-1}(C - b) \mod 26.$$

• Remark.  $a^{-1} \mod 26$  exists because gcd(a, 26) = 1.

# Remarks on affine cipher

• When a=1, the affine cipher becomes Caesar cipher.

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- When a = 1, the affine cipher becomes Caesar cipher.
- The key space for affine cipher is

$$K = \{(a, b) : a, b \in \mathbb{Z}_{26}, \gcd(a, 26) = 1\}.$$

# Key space of affine cipher

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• There are 12 choices for a

$$a \in \{1, 3, 5, 7, 9, 11, 15, 17, 19, 21, 23, 25\}$$

• There are 26 choices for b

$$b \in \{0, 1, \dots, 25\}$$

# Key space of affine cipher

$$K = \{(a, b) : a, b \in \mathbb{Z}_{26}, \gcd(a, 26) = 1\}.$$

• There are 12 choices for a

$$a \in \{1, 3, 5, 7, 9, 11, 15, 17, 19, 21, 23, 25\}$$

• There are 26 choices for b

$$b \in \{0, 1, \dots, 25\}$$

• The key space is  $|K| = 12 \cdot 26 = 312$ .



## Example 2

Encrypt the plaintext SITISTHEBEST using affine cipher with the key (a,b)=(3,13).

# Summary on affine cipher

#### Encryption: $P \rightarrow C$

- Identify each letter with an integer between 0 and 25. Write out the plaintext P after this identification.
- ② Choose a key k = (a, b) with  $a \in \mathbb{Z}_{26}^*, b \in \mathbb{Z}_{26}$ .
- lacktriangledown The plaintext P is encrypted to a cipher text C by

$$C = aP + b \mod 26$$
.

#### **Decryption:** $C \rightarrow P$

$$P = a^{-1}(C - b) \mod 26.$$



#### Exercise 5

In this exercise, we decrypt the ciphertext AXG using affine cipher with key  $(a,b)=(7,3)\,$ 

(a) Write out the formula which gives **decryption rule** for affine cipher with the key (a, b) = (7, 3).

### Exercise 5

(b) Find  $7^{-1} \mod 26$  and decrypt AXG using affine cipher with key (a,b)=(7,3).