

# RISC-V Single-Cycle CPU Simulation

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## **Abstract**

This document details the design, implementation, and verification of a 32-bit Single-Cycle RISC-V Processor. The simulation verifies the correct execution of arithmetic, logic, and memory instructions within a single clock cycle.

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# 1 Introduction

This document presents the design and simulation of a **32-bit Single-Cycle RISC-V Processor**, implemented and verified by students at the **Ghulam Ishaq Khan (GIK) Institute**.

## 1.1 Objective

The main objective is to ensure that the CPU executes instructions correctly in a **single clock cycle per instruction**. The design focuses on:

- **Data Path:** Correct movement of data between registers, ALU, and memory.
- **Control Unit:** Generation of accurate control signals (e.g., RegWrite, MemWrite, MemToReg).
- **Instruction Handling:** Execution of **I-type**, **R-type**, and **S-type** instructions.

## 2 Test Assembly Program

The program consists of 5 instructions designed to test arithmetic operations, memory access, and register usage.

Table 1: Test Assembly Program Execution Flow

Step	Address	Instruction	Machine Code	Description
1	0x00	addi x7, x0, 2	00200393	Adds 2 to x0 and stores in x7.
2	0x04	addi x6, x0, 0x123	12300313	Adds 0x123 (291) to x0; stores in x6.
3	0x08	sw x7, 0(x6)	00732023	Stores x7 value into memory at address x6.
4	0x0C	lw x8, 0(x6)	00032403	Loads data from memory address x6 into x8.
5	0x10	add x9, x8, x7	007404b3	Adds x8 + x7 and stores in x9 (Result: 4).

## 3 CPU Microarchitecture Overview

The CPU is composed of several modular Verilog components, each performing a specific role in the fetch-decode-execute cycle.

### 3.1 Instruction Memory

- Stores all program instructions.
- **Input:** Program Counter (PC).
- **Output:** 32-bit instruction word.

### 3.2 Program Counter (`register_flop`)

- Tracks the current instruction address.
- Updates on the rising edge of the clock.
- Resynchronous reset sets the PC to 0.

### 3.3 Register File (`regfile`)

- Contains 32 general-purpose registers (`x0–x31`).
- Supports dual-port read and single-port write.
- `x0` is hardwired to 0.

### 3.4 ALU (Arithmetic Logic Unit)

- Performs arithmetic (ADD, SUB) and logical operations.
- **Inputs:** Operand A, Operand B, ALU Control.
- **Output:** Calculation result and Zero flag.

### 3.5 Data Memory (`data_memory`)

- Simulates system RAM.
- **Store (SW):** Writes data when `mem_write` is high.
- **Load (LW):** Asynchronously outputs data based on address.

### 3.6 Control Unit

- Decodes the 7-bit opcode, `funct3`, and `funct7`.
- Generates control signals: `alu_src`, `mem_write`, `mem_to_reg`, etc.

## 4 Cycle-by-Cycle Execution

The following table details the internal state changes during the simulation.

Table 2: Cycle-by-Cycle Simulation Results

Cyc	Instruction	Stage	Activity	Result
1	addi x7, x0, 2	Decode	ALU adds $0 + 2$	x7 = 2
2	addi x6, x0, 0x123	Decode	ALU adds $0 + 291$	x6 = 291
3	sw x7, 0(x6)	Mem Wr	Mem write enabled at 0x123	M[291]=2
4	lw x8, 0(x6)	Mem Rd	Mem read enabled at 0x123	x8 = 2
5	add x9, x8, x7	Ex/WB	ALU adds $2 + 2$	x9 = 4

## 5 CPU Data Flow Diagram

The visual representation below illustrates the high-level connection of modules within the CPU.

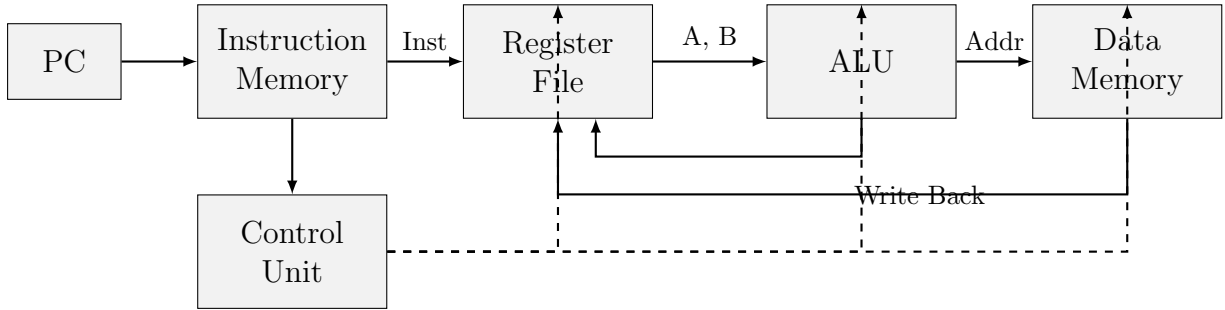


Figure 1: Abstract Data Path and Control Flow

## 6 Conclusion

The simulation validated that the **Single-Cycle RISC-V CPU** executes instructions correctly:

- Arithmetic and memory instructions were processed without stalls or hazards.
- The final register state **x9 = 4** confirms the correct interaction between the ALU, memory subsystem, and register file.
- The design successfully implements the core RISC-V Instruction Set Architecture subset.

## Future Enhancements

To further improve the design, the following features are proposed:

1. Implementation of branch instructions (BEQ, BNE) for control flow.
2. Extension of the ALU to support logic operations (AND, OR, XOR).
3. Integration of a visual timing diagram for signal verification.