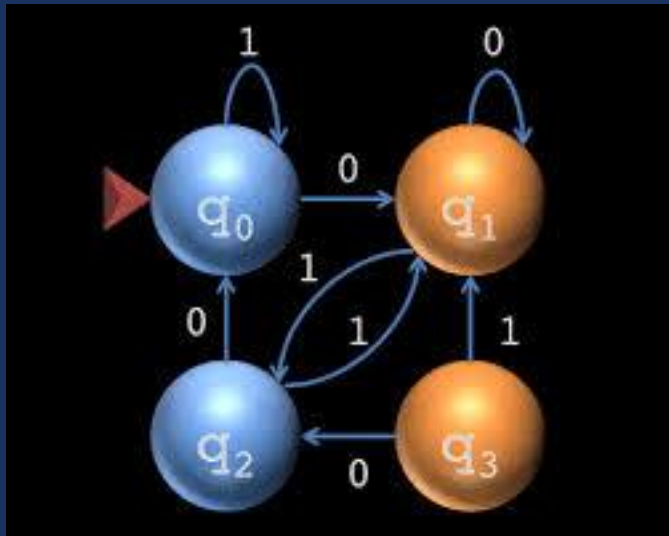


CS3005 –THEORY OF AUTOMATA

WEEK 1: INTRODUCTION TO FORMAL LANGUAGES & FINITE THEORY OF AUTOMATA



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MAJOR PARTS IN THEORY OF AUTOMATA

Foundation



```
graph TD; A[Foundation] --> B[Pushdown Automata Theory]; B --> C[Turing Theory];
```

Pushdown
Automata Theory

Turing Theory

DETAILED COURSE OUTLINE

Foundation - II

- Introduction to Theory of Automata.
- Languages.
- Recursive Definitions.
- Regular Expressions.
- Finite Automata.
- Transition Graphs.
- Kleene's Theorem.
- Finite Automata with Outputs.
- Regular Languages.
- Non-Regular Languages.
- Decidability

DETAILED COURSE OUTLINE

Pushdown Automata Theory - 7

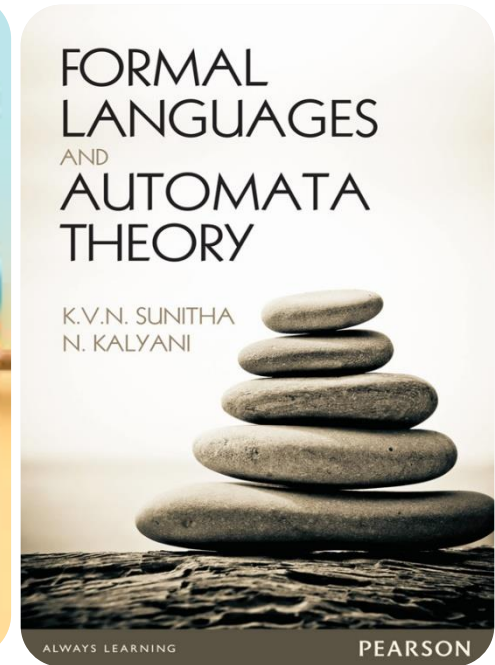
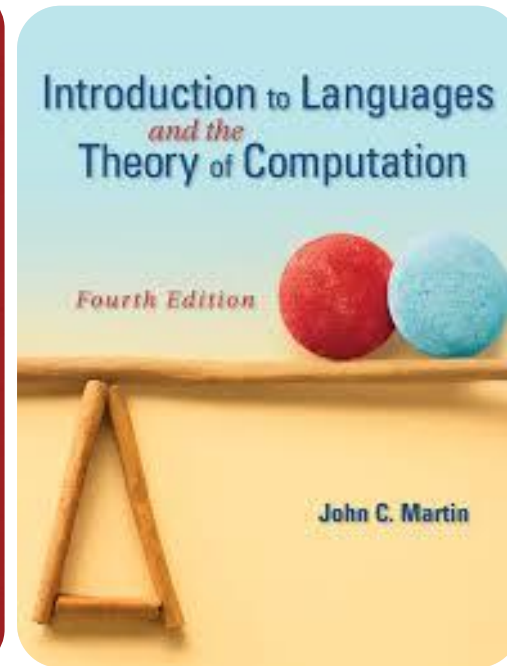
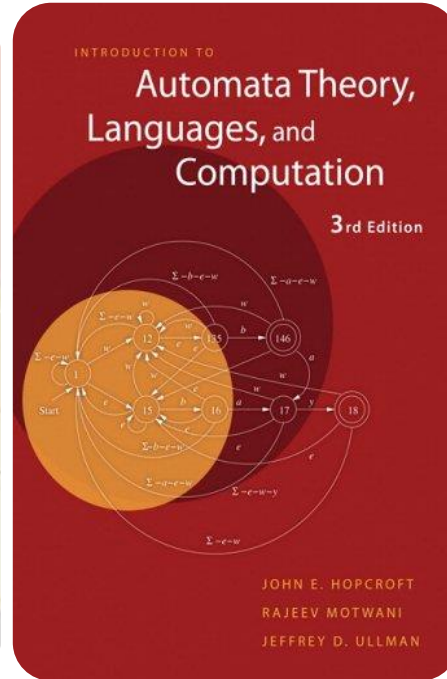
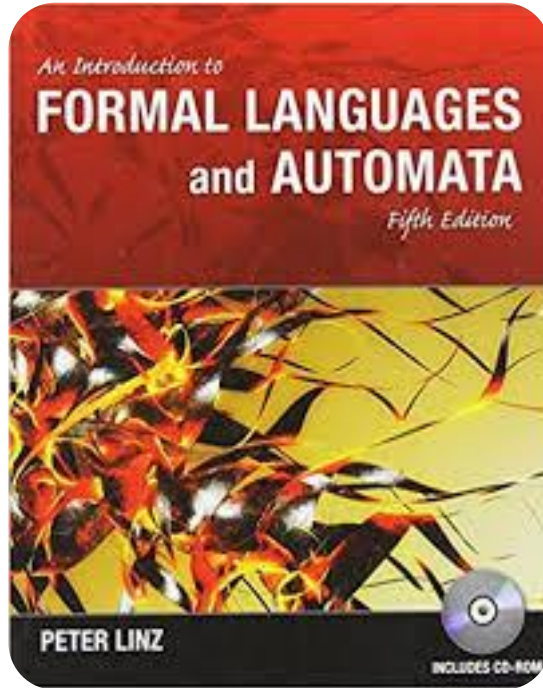
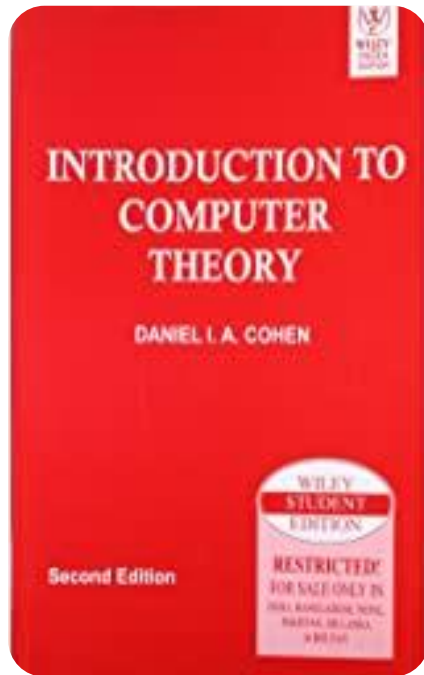
- Context Free Grammars.
- Grammatical Format.
- Pushdown Automata.
- $\text{CFG} = \text{PDA}$.
- Non-Context Free Languages.
- Context-Free Languages.
- Decidability

DETAILED COURSE OUTLINE

Turing Theory - 7

- Turing Machines.
- Post Machines.
- Minsky's Theorem.
- Variations on the TM.
- TM Languages.
- The Chomsky Hierarchy.
- Computers.

BOOKS



MARKS DISTRIBUTION

- Mid term I & II – 30 Marks
- Final exam – 50 Marks
- Project – 10 Marks
- Quiz + Activities – 5 Marks (n-I)
- Assignments – 5 Marks (n-I)

FORMAL LANGUAGE & FINITE THEORY OF AUTOMATA

- When we call our study the Theory of Formal Languages, the word "formal" refers to the fact that all the rules for the language are explicitly stated in terms of what strings of symbols can occur. No liberties are tolerated, and no reference to any "deeper understanding" is required. Language will be considered solely as symbols on paper and not as expressions of ideas in the minds of humans.
- In this basic mod The term "formal" used here emphasizes that it is the form of the string of symbols we are interested in, not the meaning. We begin with only one finite set of fundamental units out of which we build structures. We shall call this the alphabet.
- A certain specified set of strings of characters from the alphabet will be called the language. Those strings that are permissible in the language we call words, The symbols in language is not communication among intellects, but a game of symbols with formal rules.

WHY STUDY FORMAL LANGUAGES & FINITE THEORY OF AUTOMATA?

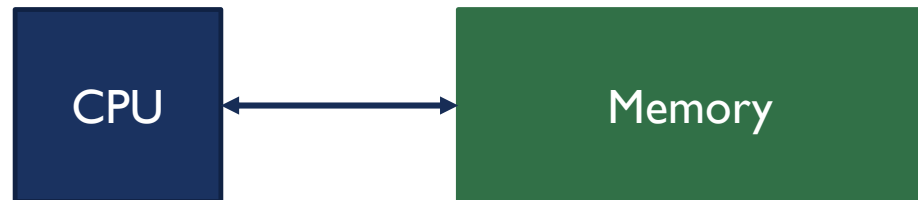
- As a computer professional you are dealing with logics. As we are dealing with high velocity, volume and multi type of file formatting we need high standards in computation machines.
- For the Processing of AI and Machine Learning, Deep Learning, Neural Networks we want to know how we are going to solve the problem and again our limitations come with hardware.
- Making a model in which we can able to understand that either the problem is solvable or unsolvable in given conditions or not.
- How we can compare that two different machines are performing same?
- How we are going to create new programming languages?
- Construction of new compilers to develop new programming languages.

AUTOMATA

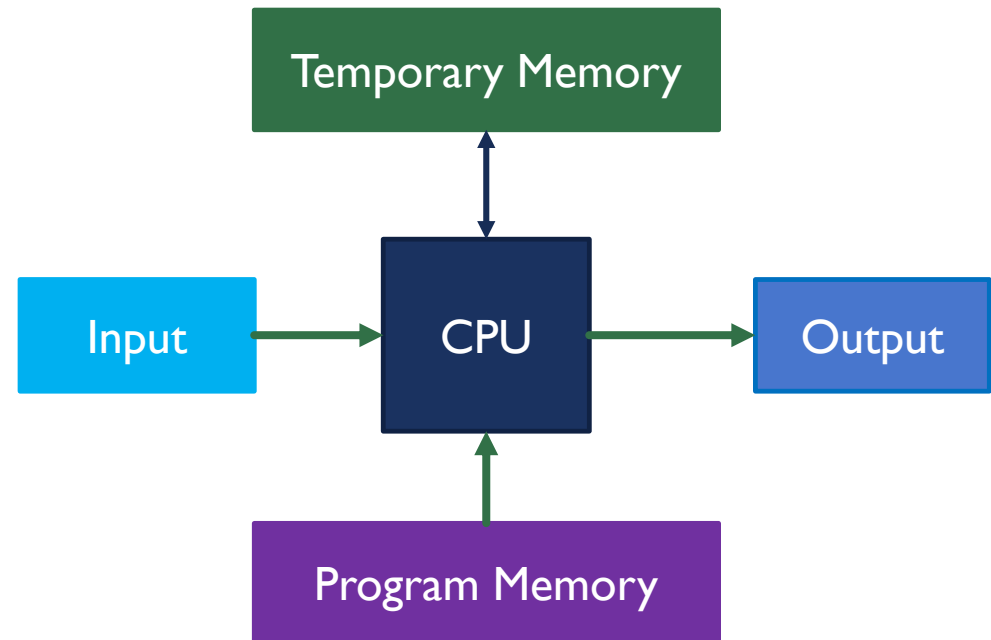
- The **word "automaton"** is the Latinization of **Greek** αὐτόματον, **automaton**, (neuter) "acting of one's own will".
- This **word** was first used by Homer to describe automatic door opening, or automatic movement of wheeled tripods.
- An automaton (Automata in plural) is an abstract self-propelled computing device which follows a predetermined sequence of operations automatically.
- An automaton with a finite number of states is called a **Finite Automaton (FA) or Finite State Machine (FSM)**.

COMPUTATION MODELS

A widely accepted model of computation

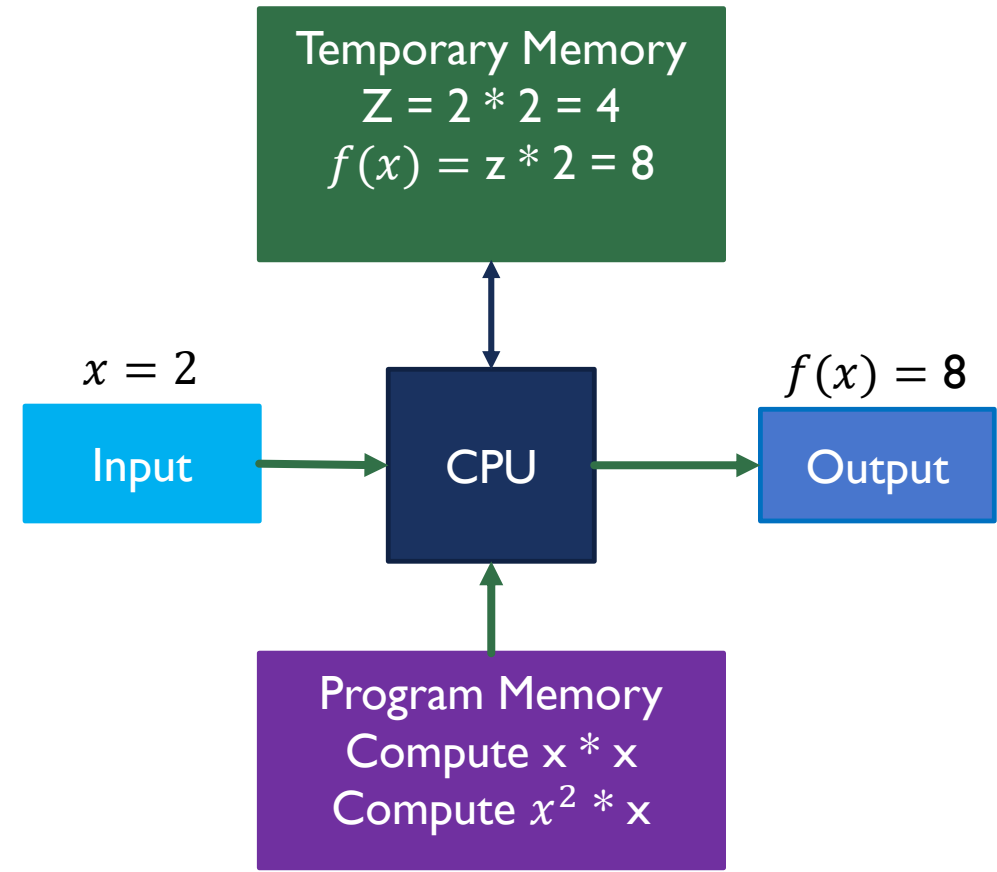
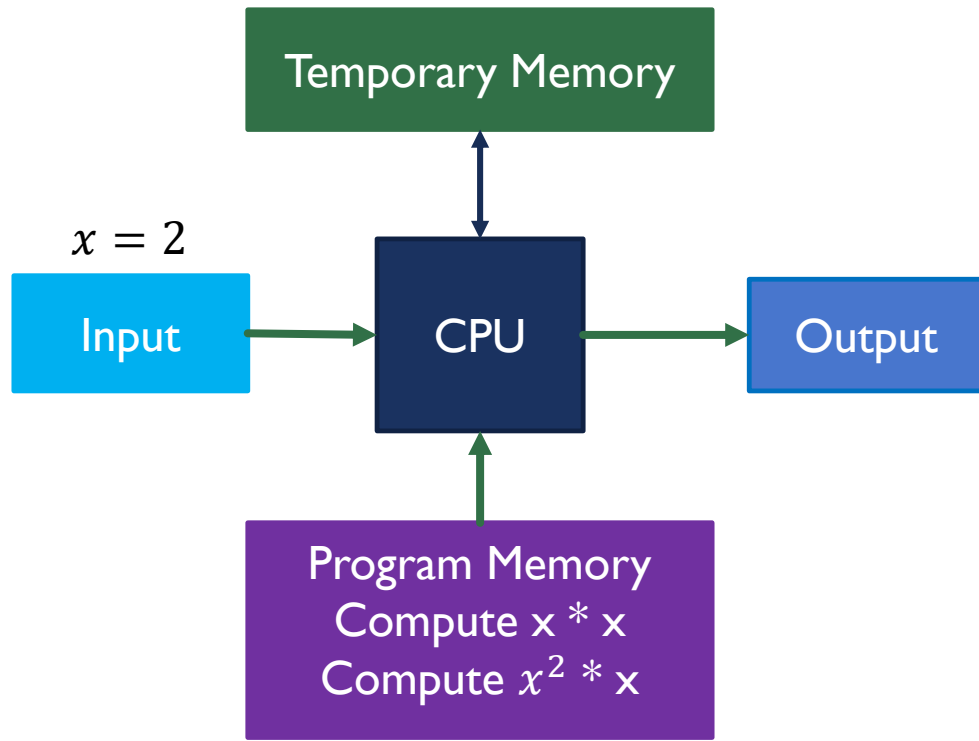


The different components of memory



PROCESSING THE MODEL

Example: $f(x) = x^3$



AUTOMATON

Temporary Memory

Automaton

CPU

Program Memory

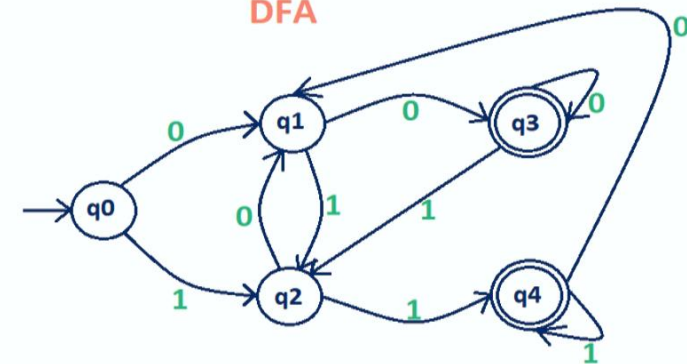
Input

Output

Temporary Memory

Automaton

DFA

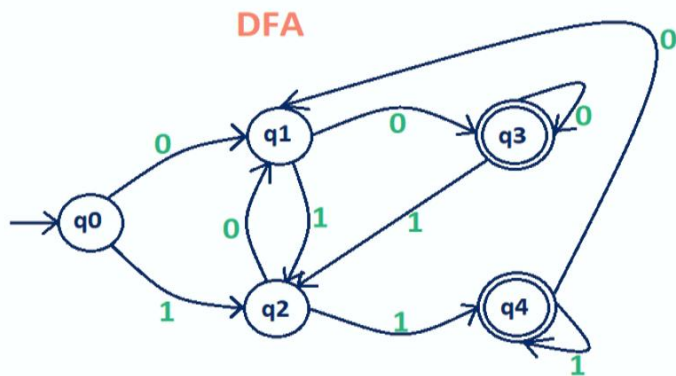


Output

CPU + Program Memory = States + Transitions

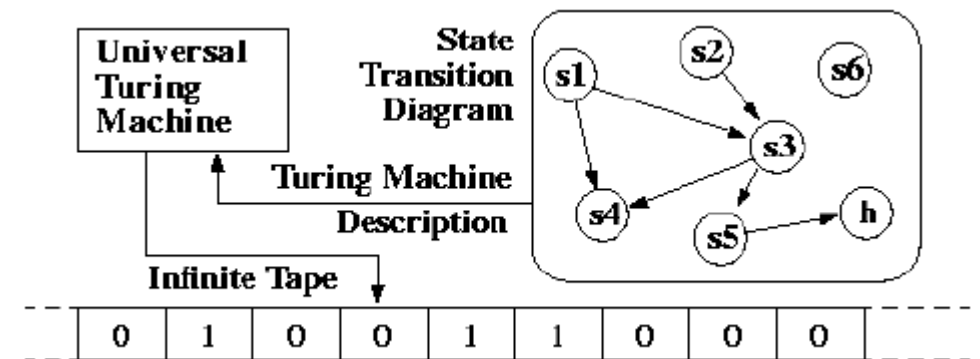
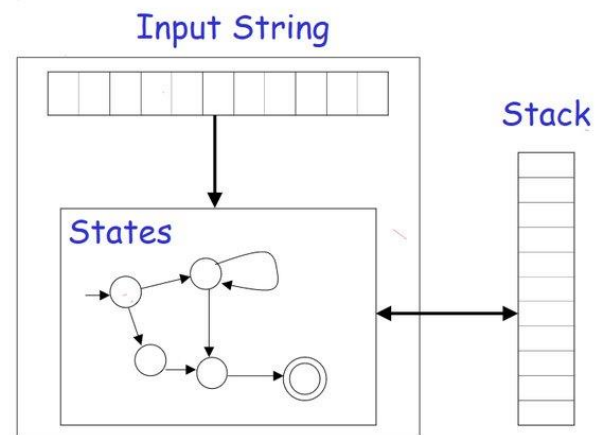
DIFFERENT KINDS OF AUTOMATA

Type of Automata	Memory Type
Finite Automata	No Temporary Memory
Push Down Automata	Stack
Turing Machines	Random Access Memory



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Pushdown Automaton -- PDA



1/20/2025

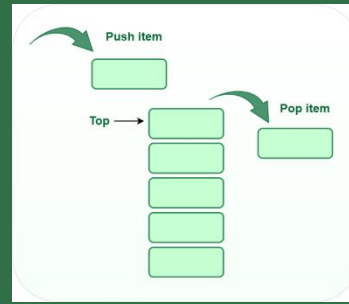
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COMPARISON OF PDA WITH TM

Feature	Push Down Automata	Turing Machine
Tape	Has only one tape	Has infinite tape
Memory	Has a stack	Has infinite memory
Computational Power	Less powerful	More powerful
Acceptance of Languages	Accept Context Free Languages	Accepts all Languages
Type of Automata	Non-Deterministic	Deterministic

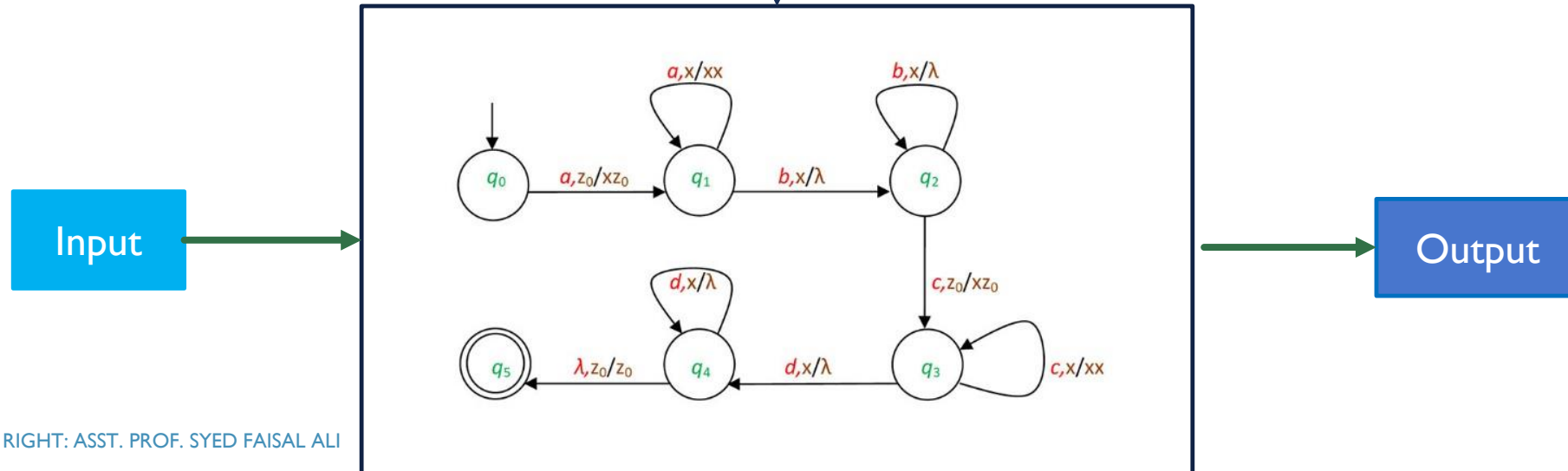
PUSH DOWN AUTOMATA (PDA)

Temporary Memory



Example: Parsers for Programming Languages (medium computing power)

Push Down Automaton

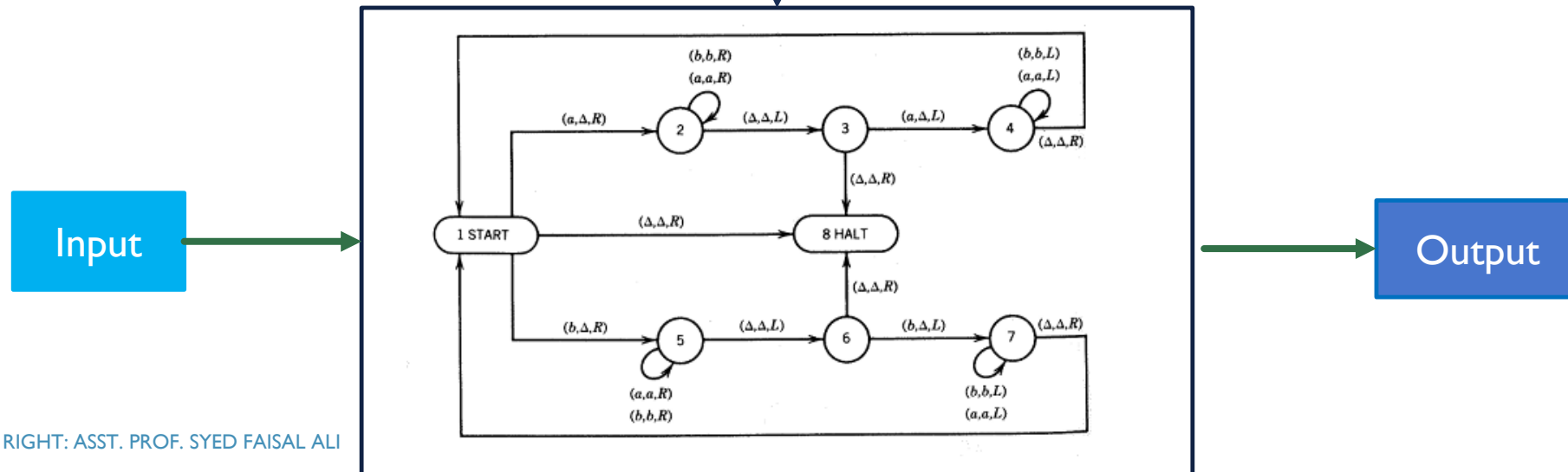


TURING MACHINE (TM)

Random Access Memory

Example: Any algorithm (highest known computing power)

Turing Machine



POWER OF AUTOMATA

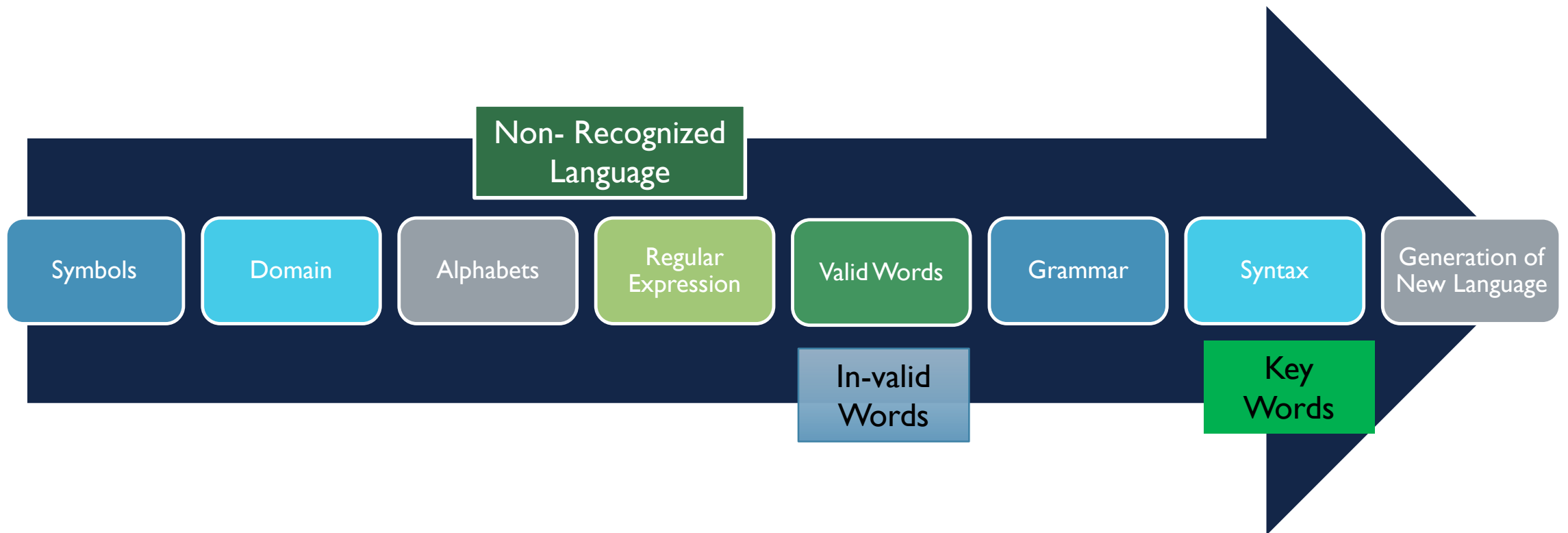
Simple Problems	More Complex Problems	Hardest Problems
Finite Automata <	Push Down Automata <	Turing Machine
Less Power		More Power

Solve more computational problems

THE MOST POWERFUL COMPUTATION MODEL

- Turing Machine is the most powerful known computation model
- **Question:** Can Turing Machine solve all computation problems?
- **Answer:** No (There are unsolvable problems too)

WHAT WE DO IN AUTOMATA AS COMPUTER SCIENTIST?



STRING & LANGUAGES

- The mathematical study of the “Theory of Computation” begins by understanding the Mathematics of strings of symbols.
- *Alphabet: It is defined as a finite set of symbols.*
- **Example :** Roman alphabet $\{a, b, \dots, z\}$.
- “Binary Alphabet” $\{0, 1\}$ is pertinent to the theory of computation.
- **String:** A “string” over an alphabet is a finite sequence of symbols from that alphabet, which is usually written next to one another and not separated by commas.
- (i) If $\Sigma_a = \{0, 1\}$ then 001001 is a string over Σ_a .
- (ii) If $\Sigma_b = \{a, b, c, \dots, z\}$ then axyrpqstcd is a string over Σ_b .
- **Length of String:** The “length” of a string is its length as a sequence. The length of a string w is written as $|w|$.
- **Example:** $|10011| = 5$

STRING & LANGUAGES

- **Empty String:** The string of zero length is called the “empty string”. This is denoted by ϵ_0 , or ε_0 or λ (Epsilon / Lambda).
- The empty string plays the role of 0 in a number system.
- **Reverse String:** If $w = w_1 w_2 w_3 \dots w_n$ where each $w_i \in \Sigma$, the reverse of w is $w_n w_{n-1} w_{n-2} \dots w_1$
- **Substring:** z is a substring of w if z appears consecutively within w .
- As an example, ‘deck’ is a substring of ‘abcdeckabcjkl’.
- **Concatenation:** Assume a string x of length m and string y of length n , the concatenation of x and y is written xy , which is the string obtained by appending y to the end of x , as in $x_1 x_2 x_3 \dots x_m y_1 y_2 y_3 \dots y_n$.

SUFFIX , PREFIX AND LEXICOGRAPHICAL ORDER

- To concatenate a string with itself many times we use the “superscript” notation.

$$\overbrace{xx \dots x}^k = x^k$$

- **Suffix:** If $w = xv$ for some x , then v is a suffix of w .
- **Prefix:** If $w = vy$ for some y , then v is a prefix of w .
- **Lexicographic ordering:** The Lexicographic ordering of strings is the same as the dictionary ordering, except that shorter strings precede longer strings.

LANGUAGES

- A *language* is a subset of Σ^* for some alphabet Σ .
- **Example:** The set of strings of 0's and 1's with no two consecutive 1's.
- $L = \{\epsilon, 0, 1, 00, 01, 10, 000, 001, 010, 100, 101, 0000, 0001, 0010, 0100, 0101, 1000, 1001, 1010, \dots\}$
- A super set can be more complex and have all the members in the Universal set of 0 and 1.
- $\{0,1\}^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, 001, \dots\}$
- **Subtlety:** 0 as a string, 0 as a symbol look the same.

GENERATING LANGUAGE

- Let us consider some simple examples of languages. If we start with an alphabet having only one letter, the letter x,

$$\Sigma = \{x\}$$

- We can define a language by saying that any non empty string of alphabet characters is a word.

$$L = \{x \ xx \ xxx \ xxxx \dots\}$$

- Or to write this in an alternate form:

$$L = \{x^n \text{ for } n = 1 \ 2 \ 3 \dots\}$$

MAN, GOAT, WOLF AND CABBAGE

- Now consider a riddle in which there is a man, a goat, a wolf, and a cabbage. He has to go across the river by boat but he can only take one thing at a time. If he takes the cabbage but leaves the goat and the wolf together, the wolf will eat the goat. If he takes the wolf but leaves the goat and cabbage together, then the goat will eat the cabbage. How can he get the goat, the wolf and the cabbage across the river?



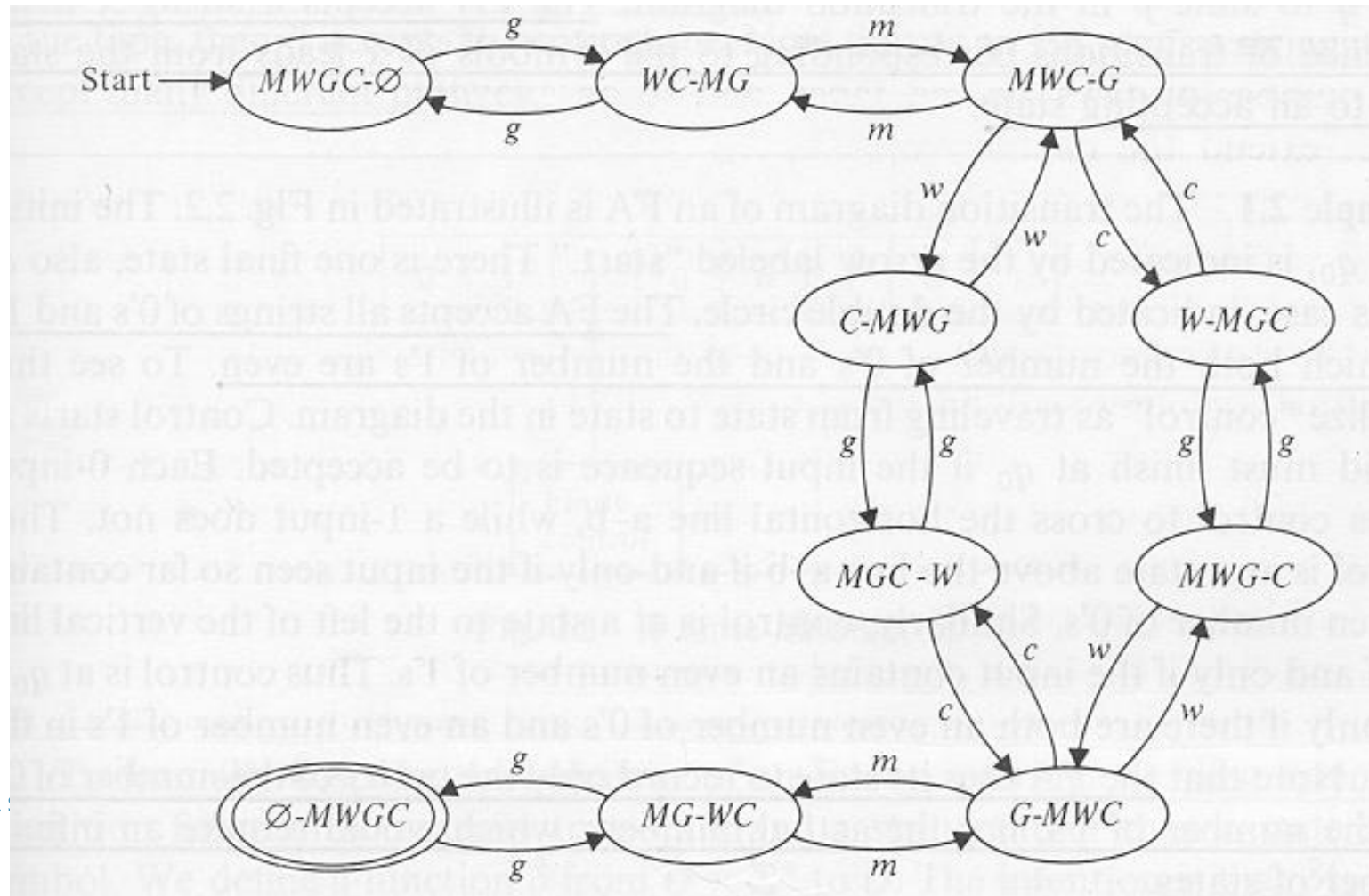
SOLUTION

- The man first takes the goat across the river. He leaves the goat and returns to get the wolf. He leaves the wolf on the other side to but brings back the goat. He then takes the cabbage and leaves the goat. When he is on the other side he reaves the cabbage with wolf. The wolf will not eat the cabbage. The man goes back to pick up the goat and finally returns to the other side where the wolf and cabbage are waiting.
- Now the question is how you are going to solve this riddle using automata?

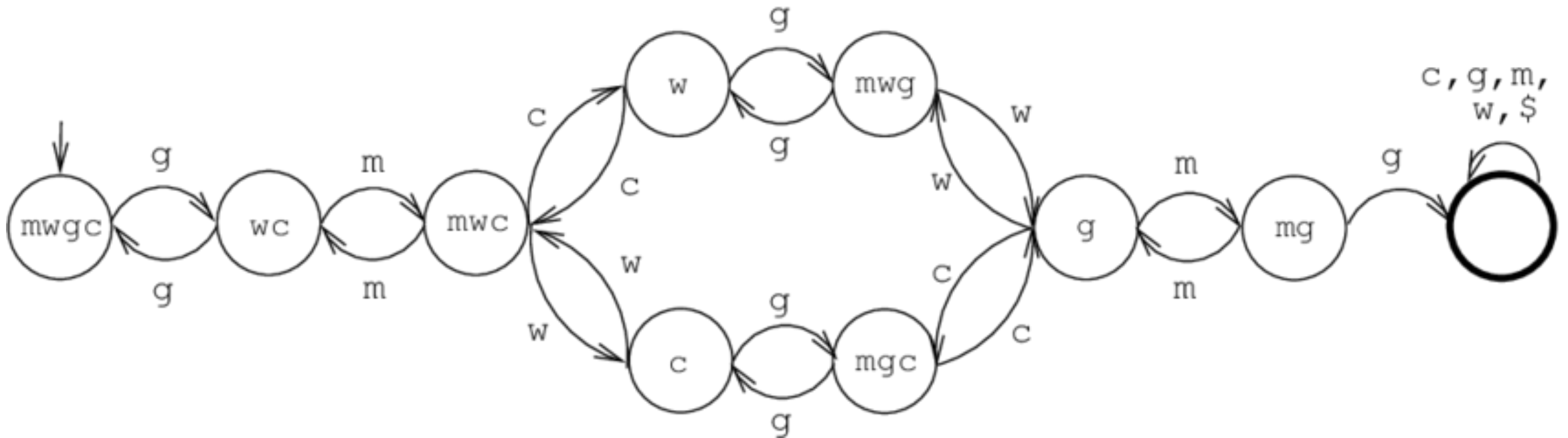
SOLVE THIS RIDDLE USING AUTOMATA

- Step 1 : Find all the elements which are in the riddle.
- Step 2: Find the power set for the riddle.
- Step 3: Find all the constraints in your riddle.
- Step 4: Now identify your initial and final conditions.
- Step 6: Starting from initial condition try to find how we are going to connect to next step.
- Step 7: Continue step 6 until you reach final state.

TRANSITION DIAGRAM FOR MAN, WOLF, GOAT AND CABBAGE



GENERALIZED AUTOMATA FOR THE RIDDLE



READING ASSIGNMENT

- Read what you have been taught today.
- Download book and read Chapter 1 & Chapter 2 from Daniel I A Cohen book.
- Update your python tool so we can start doing programming.
- Find one real problem and try to solve your problem by using automata.



END OF LECTURE

