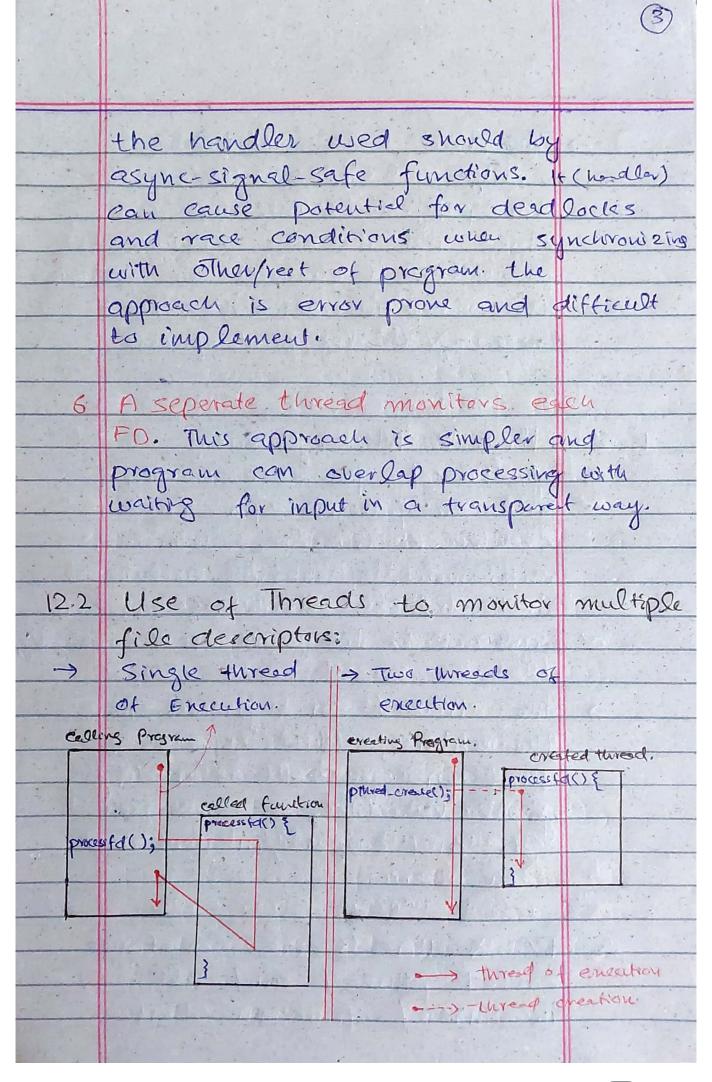
	CSE-302 15 Jan 2024/NO
	B: Unix System Programming
	System Programming Notes
	chapter 12: Unix/Posix
	Thréads.
->	One method of achieving parallelism
	is for multiple process to
	cooperate and synchronize throug
	shared memory and message
	passing
->	An alternative approach isses
	multiple threads of execution in a
	An alternative approach isses multiple threads of execution in a single address space.
10	
-	A motivating Problem: Monitoring.
	File descriptors:
->	Sin general approaches du
	Sin general approaches des monitor multiple file descriptos
	for input under POSIX are
	as follows.
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1	A seperat process (fork) monitors.
	each file descriptor.
	Problem since the child does't share
	any variable, we may use

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	and the state of the second met and a constate of the light of the second of the secon
100000	shared memory or message.
	passing to exenange information
2	Using select & system cell blocking
3	Using select & system call blocking using poll system call. I blocking
	Problem Once the blocking Call
	returns, the calling program
	handles each ready file descriptou
	In turn. Furthmore, the program
	can do no useful work.
4	Nou-blocking 1/0 with polling work
	well when the program has to
	do "usefull work" that it can
The contract of the contract o	perform between its intermittents
	cheeks to see if 1/0 is available.
	problems most problems are doftralt
	to structure in this way, and it
	timing for 1/0 cheek relative to
	useful work otherwir it can
	Dead to busy writing
5	POSIX asynchronous 1/0 can
	be used with or without signed.
	notification.
	Problems: If used with signals



<b>(4)</b>	
$\rightarrow$	The thread is "schedulable"
	entity" with its own value of
	the program counter, its own
A mines of the	Stack and its own scheduling
	parameters. It enecutes an independent
	stream of instructions, never
	reforning to the point of call.
	The first design of the state o
12-3	Thread Management:
	A thread package usually
	includes functions for three
	creation, destruction, schooluling,
	enforcement of mutual enclusion
	and conditional waiting, run-time
	system to manage threads transpanealy
	is theread A. Somers of
->	the entire address space of
	that process and can modify
	global variables, access open
	fals and cooperat/interfere with
	each others
	THE PROPERTY OF THE PARTY OF TH
->	POSIX unreads are sometimes
	called pthread and all
	the thread functions Start

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The second	All the water that a significant to the state of the stat
	with pthreads
->	Some posix threads mangment
	functions are:
#	POSIX function Description
112	Pesent Description
1	pthread cancel terminate another thread.
- 2,	" - create evente a turend.
3	" - detach sel thread to release resources.
4	- equal test two threads IDs for equility.
- +5	" - exit enit a-live of evithous enisting process.
6	" - Kill send a signal to a thread.
+7	" -join wait for a thread.
78	" _ self find out own thread 110.
->	Most pthreeds redurns oil
Park Clark	successful and nonzero error eadle
	if unsuccessful. (don't set errno).
->	Synopsis:
	#include < pinread.h>
	pthread_t pluread_self (void);
	Dilly and a second of the seco
	Princeds are refrenced by an ID
	of type pluread to A thread can find out its id by calling this
	find out its 14 by calling this
	tunction
	void * No errors are defined.

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->	Synopsis:
	#include (plhread. h)
	int pthread create (
	pruvead + * restrict thread;
	const prhread-attr_t * restrict attr,
- 11111	void*(*start_routine)(voia*),
	void* arg);
linese	This function creates a thread. by
(Miss	pointing to the 1D of newly created
	thread.
	attr: represent an attribute object
	that encapsulate the attribute of
	a thread. If NULL, the new thread
	has default attributes.
	Start-routine: the name of a function
	that the thread calle when it
	begin execution. It takes a single
7	parameters specified by arg, a
	pointer to void. It return a
	pointer to void, which is treated
	as an emit status by planeag-join.
	us O.
*	return monzero error code.
	The state of the s
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		<b>(</b>
$\rightarrow$	Synopsis	
	# include < pthread. 4>	
	int pthread_join (pthread)	t' thread,
	void ** value-ptr);	
	This function susspends the cal	Iling
v. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1.	thread untill the target three	q,
1	terminates.	
	thread: ID of terminating -	thread
	· value-ptre provides a Jocatio	n for
and the second	a pointer to the return sta	tus that
	the target thread passes t	o return.
Contract of the second	or 'pthread-exit'.	
7	15 0 *Us non-zero error co	12.
->	Synopsiso	
	#includ spihread hs	
	void pthread_enit (void	value ptr);
10 - 10 - 10 - 10 - 10 - 10 - 10 - 10 -	A call to this system call	causes
	only the calling thread to	
	terminate. A threog that	
	enecutes return from its top	
	level implicitly calls this f	
	value-ptr: this value is avail	lable
1	to a successful pluread-join	
	* Yord. Posix doe'f define any	

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The second secon	C HIN C China
	error for this function
-	The process can terminal by
	calling exit directly or by
	of the other process threads call
	of the other process threads call
A	all tureads are terminated.
	all threads vare terminated.
	n and the Court courses with a
7	A call to enit causes entire threed to terminale, a call
	to pthread_enit causes only
	The calling - Inread to terminate.
	Comment of the control of the contro
>	Parallelism is of two types
	· Data Level: in this type
	data is divided among
	. Task Jevel: tasks are
	divided anien multi-process/threads.
	Cold Cilian Cold Service Cold Cilian
	The state of the s
	gentlesh and thinkship thispier the
	The cold are the second of the
	Discheration with a college to the state of
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