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ULUCAY

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1-a) for open addressing Load factor must be 0.75

$$L_o = \frac{n}{s} \Rightarrow 0.75 = \frac{400}{\text{size}} \quad \frac{3}{4} = \frac{400}{\text{size}} \quad \text{size} = \frac{1600}{3}$$

$$\boxed{\text{size} = 533}$$

1-b) for chaining Load factor^{max} must be 3.0

$$L_c = \frac{400}{s} = 3 \quad s = \frac{400}{3} \quad \boxed{s = 134}$$

2) 4, 14, 24, 6

hash = %10

Linear probing

4	4
5	14
6	24
7	6

quadratic probing

4	4
5	14
6	6
7	
8	
9	24

- quadratic probing is increasing the randomize inserting

If we want to insert a couple of similar numbers in linear probing this cause a lot of collision because of their hash code. But when we turn it to quadratic probing we are distribute the numbers in table much more easily

3-) [3, 4, 7, 6, 9]

Y. M. G. S.

a) 3 | 4 7 6 9
 3 4 | 7 6 9
 3 4 7 | 6 9
 3 4 6 7 | 9
 3 4 6 7 9

b) 3 4 7 6 9
 3 4 7 6 9
 3 4 7 6 9
 3 4 7 6 9
 3 4 6 7 9

out of size
 $gap = 7/2.2$
 $gap = 3$
 $gap = 3/2.2$
 $gap = 1$

according to the implementation swap could be in 4th step

c) ③ 4 7 6 9
 - ③ ④ 6 7 9
 - ④ 6 7 9
 - ⑥ 7 9
 - ⑦ 9
 - ⑨

3 | 4 | 6 | 7 | 9

