

I hereby pledge on my honor that I will strictly adhere to academic integrity codes and the work done on this examination is solely my own and I will not receive/give my help from/to anybody or source during this examination.

1.) for ArrayList: (Assuming i starts 1)
 2) $sum = 0 \Rightarrow \Theta(1)$ (Otherwise it will be infinite loop) $= O(\infty)$
 $for(int i=1; i < myList.size(); i *= 2) \Rightarrow \Theta(\log n)$
 $sum += myList.get(i); \Rightarrow \Theta(1)$
 This method $T(n)$ for ArrayList is constant

$$T(n) = \underbrace{\Theta(1)}_1 + (\underbrace{\Theta(\log(n))}_2 * \underbrace{\Theta(1)}_3)$$

- We can ignore 1 according to the following rule

$$T(n) = \max(T_1(n), T_2(n))$$

- in for loop i is increasing $i = i * 2 \Rightarrow (\log n)$

$$\boxed{T(n) = \Theta(\log(n))}$$

2) for LinkedList

$sum = 0 \Rightarrow \Theta(1)$ ignore

$for(int i=1; i < myList.size(); i *= 2) \Rightarrow \Theta(\log n)$
 $sum += myList.get(i);$

This is working $\Theta(n)$ time for LinkedList

$$T_n = \Theta_1 + (\Theta(\log n) * \Theta(n)) \quad \boxed{T(n) = \Theta(n \log n)}$$

2-) int foo(A l) ?

if (l == l.size()) return 0; $\Rightarrow \Theta(1)$

int sum = 0; $\Rightarrow \Theta(1)$

int x = l.get(l.size()-1); $\Rightarrow \Theta(1)$ constant for ArrayList

for (int i = 0; i < l.size(); ++i) $\Rightarrow \underline{\Theta(n)}$

sum += x % i $\Rightarrow \Theta(1)$

l.remove(l.size()-1); $\Rightarrow \Theta(1)$ because deleting last element
not shifting array for each time

sum += foo(l); $\Rightarrow \underline{\Theta(n)}$ it will call l.size() times $\Theta(n)$

$$\frac{n \cdot (n+1)}{2} = n^2 \text{ (ignore const)}$$

$$T(n) = T(n-1) + \Theta(n)$$

$$T(n-1) = T(n-2) + \Theta(n-1)$$

$$T(n) = \Theta(n) * \Theta(n)$$

$$\boxed{T(n) = \Theta(n^2)}$$

3-)

a) $O(n \log n)$ ✓ \Rightarrow Because it will run max length of list it is $O(n)$ we can accept

$\Theta(n)$ ✓ \rightarrow for this function it is more detailed answer it would be this, because it will process l.size() time in while loop

$\Omega(n^2)$ ✗ \rightarrow we cannot accept this because omega notation must be lower than $\Theta(n)$ notation $n^2 < n$ this is wrong

b) $O(n \log n)$ ✓ \rightarrow it will complexity for the take on index

$\square \rightarrow \square \xrightarrow{\leftarrow} \square$ max it will be $O(\frac{n}{2})$

so $n \log n > \frac{n}{2}$ we can accept it because of O notation

$\Theta(n)$ ✓ \rightarrow This is also correct $\Theta(\frac{n}{2}) = \Theta(n)$ we can ignore constant

$\Omega(n^2)$ ✗ \rightarrow We can't accept this because $n^2 < n/2$ is false for omega notation