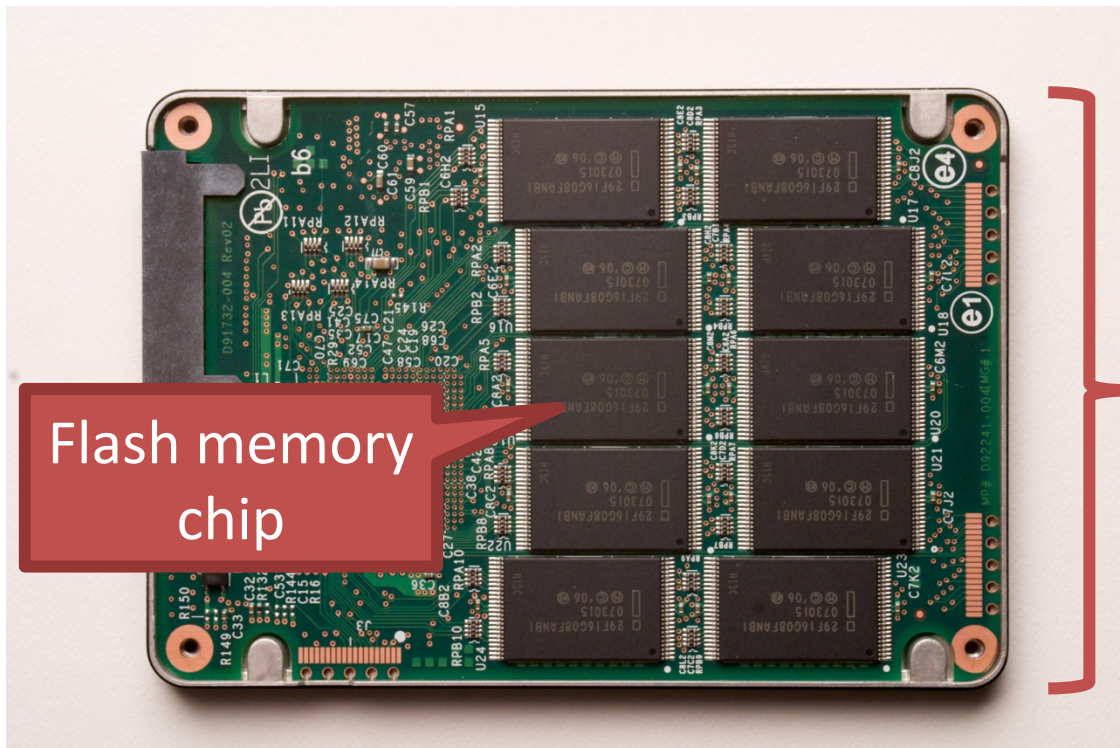


# Beyond Spinning Disks

- Hard drives have been around since 1956
  - The cheapest way to store large amounts of data
  - Sizes are still increasing rapidly
- However, hard drives are typically the slowest component in most computers
  - CPU and RAM operate at GHz
  - PCI-X and Ethernet are GB/s
- Hard drives are not suitable for mobile devices
  - Fragile mechanical components can break
  - The disk motor is extremely power hungry

# Solid State Drives

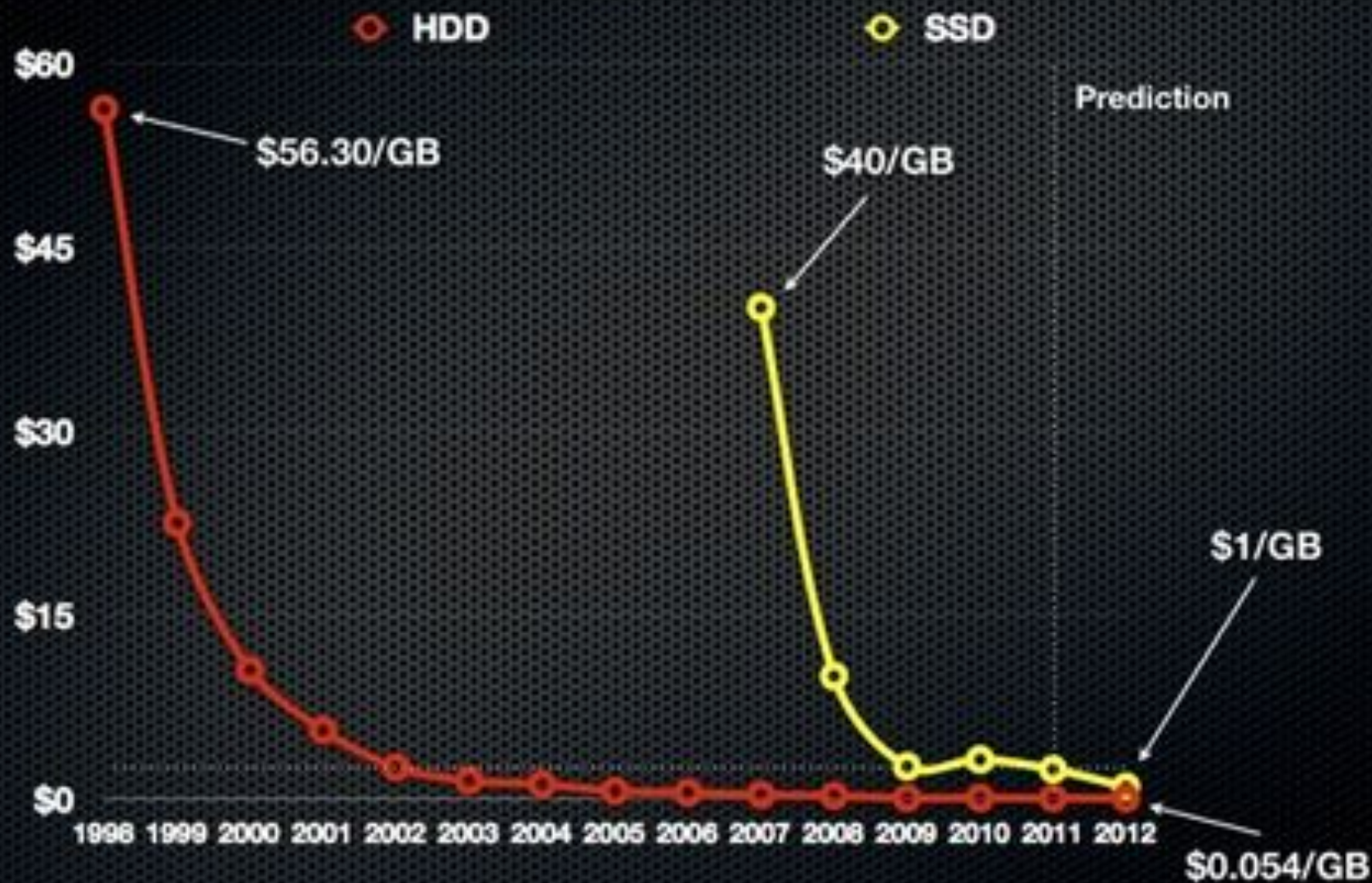
- NAND flash memory-based drives
  - High voltage is able to change the configuration of a floating-gate transistor
  - State of the transistor interpreted as binary data



# Advantages of SSDs

- More resilient against physical damage
  - No sensitive read head or moving parts
  - Immune to changes in temperature
- Greatly reduced power consumption
  - No mechanical, moving parts
- Much faster than hard drives
  - >500 MB/s vs ~200 MB/s for hard drives
  - No penalty for random access
    - Each flash cell can be addressed directly
    - No need to rotate or seek
  - Extremely high throughput
    - Although each flash chip is slow, they are RAIDed

# Average HDD and SSD prices in USD per gigabyte



# Challenges with Flash

- Flash memory is written in pages, but erased in blocks
  - Pages: 4 – 16 KB, Blocks: 128 – 256 KB
  - Thus, flash memory can become fragmented
  - Leads to the [write amplification](#) problem
- Flash memory can only be written a fixed number of times
  - Typically 3000 – 5000 cycles for MLC
  - SSDs use [wear leveling](#) to evenly distribute writes across all flash cells



# Write Amplification

G moved to new block by the garbage collector

Cleaned block can now be rewritten

Block X			
K	D	G	C'
L	E	A'	D'
C	F	B'	E'

Block Y			
G	C''	F''	J
A''	D''	H	A'''
B''	E''	I	B'''

- Once all pages have been written, valid pages must be consolidated to free up space
- **Write amplification**: a write triggers garbage collection/compaction
  - One or more blocks must be read, erased, and rewritten before the write can proceed

# Garbage Collection

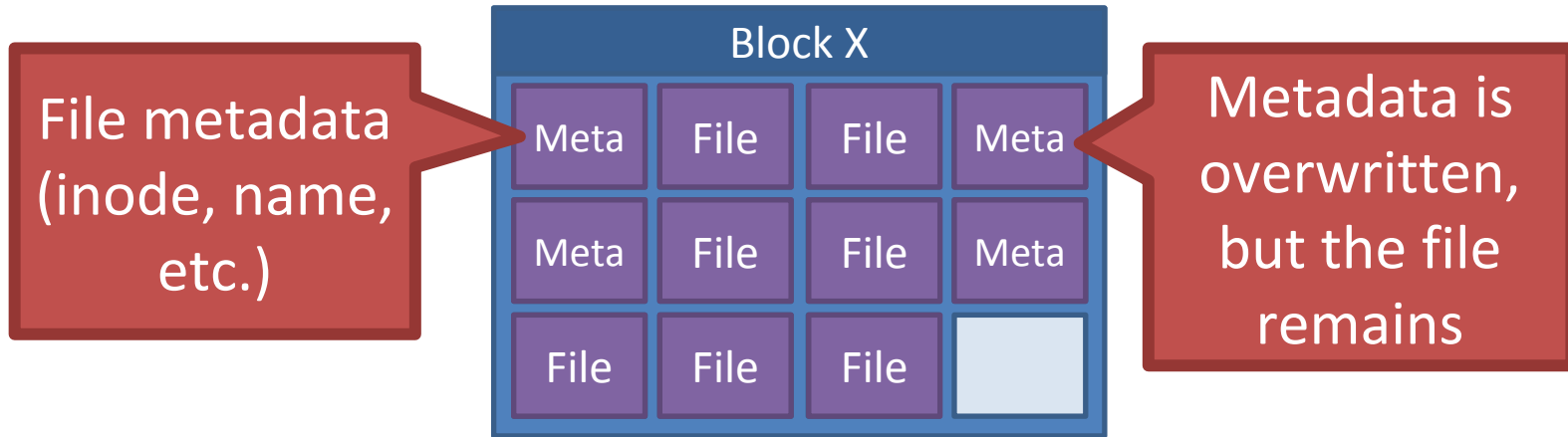
- Garbage collection (GC) is vital for the performance of SSDs
- Older SSDs had fast writes up until all pages were written once
  - Even if the drive has lots of “free space,” each write is amplified, thus reducing performance
- Many SSDs over-provision to help the GC
  - 240 GB SSDs actually have 256 GB of memory
- Modern SSDs implement background GC
  - However, this doesn't always work correctly

# The Ambiguity of Delete

- Goal: the SSD wants to perform background GC
  - But this assumes the SSD knows which pages are invalid
- Problem: most file systems don't actually delete data
  - On Linux, the “delete” function is `unlink()`
  - Removes the file meta-data, but not the file itself



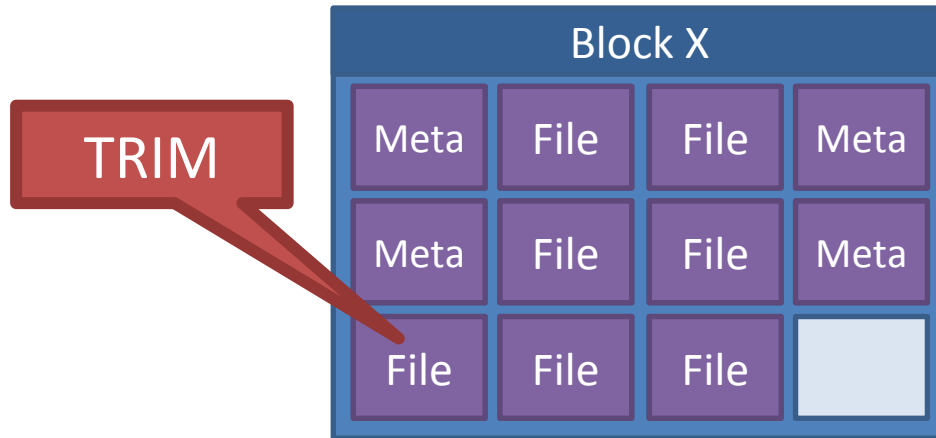
# Delete Example



1. File is written to SSD
  2. File is deleted
  3. The GC executes
    - 9 pages look valid to the SSD
    - The OS knows only 2 pages are valid
- Lack of explicit delete means the GC wastes effort copying useless pages
  - Hard drives are not GCed, so this was never a problem

# TRIM

- New SATA command TRIM (SCSI – UNMAP)
  - Allows the OS to tell the SSD that specific LBAs are invalid, may be GCed



- OS support for TRIM
  - Win 7, OSX Snow Leopard, Linux 2.6.33, Android 4.3
- Must be supported by the SSD firmware

# Wear Leveling

- Recall: each flash cell wears out after several thousand writes
- SSDs use **wear leveling** to spread writes across all cells
  - Typical consumer SSDs should last ~5 years

W

If the GC runs now, page G must be copied

Examples

Wait as long as possible before garbage collecting

Block X			
K	D	G	C'
L	E	A'	D'
C	F	B'	E'

Block Y			
F'	C''	F''	G'
A''	D''	H	A'''
B''	E''	I	B'''

Blocks with long lived data receive less wear

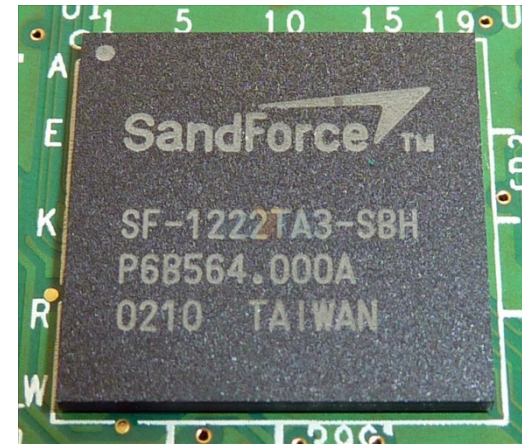
Block X			
M*	D	G	J
N*	E	H	K
O*	F	I	L

Block Y			
A	D	G	J
B	E	H	K
C	F	I	L

SSD controller periodically swap long lived data to different blocks

# SSD Controllers

- SSDs are extremely complicated internally
- All operations handled by the SSD controller
  - Maps LBAs to physical pages
  - Keeps track of free pages, controls the GC
  - May implement background GC
  - Performs wear leveling via data rotation
- Controller performance is crucial for overall SSD performance



# Flavors of NAND Flash Memory

## Multi-Level Cell (MLC)

- Multiple bits per flash cell
  - For two-level: 00, 01, 10, 11
  - 2, 3, and 4-bit MLC is available
- Higher capacity and cheaper than SLC flash
- Lower throughput due to the need for error correction
- 3000 – 5000 write cycles
- Consumes more power

**Consumer-grade drives**

## Single-Level Cell (SLC)

- One bit per flash cell
  - 0 or 1
- Lower capacity and more expensive than MLC flash
- Higher throughput than MLC
- 10000 – 100000 write cycles

**Expensive, enterprise drives**