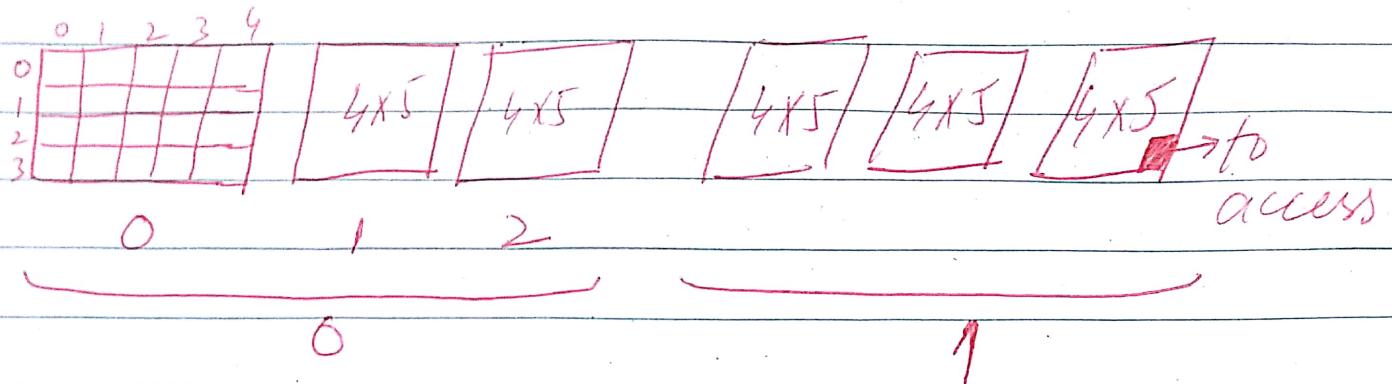




PART 01



Base Address + 1 (Block Size) + 2 (Square Size)
+ 3 (Row Size) + 4 (Int. Size)

$$\begin{aligned}
 &= 1 \left(\frac{240}{3 \times 4 \times 5 \times 4} \right) + 2 \left(\frac{4 \times 5 \times 4}{20} \right) + 3 \left(\frac{5 \times 4}{20} \right) + \frac{4 \times 4}{16} \\
 &= 460 + 16 \quad \underline{\quad 80 \quad} \quad \underline{\quad 60 \quad} \quad \underline{\quad 16 \quad} \\
 &= 476
 \end{aligned}$$

PART 02

```

bool anagrams (char *s, char *t) {
    if (strlen(s) != strlen(t))
        return false;
    int letters[256] = {0};
    for (int i = 0; i < strlen(s); i++) {
        letters[s[i]]++;
        letters[t[i]]--;
    }
    for (int i = 0; i < 256; i++)
        if (letters[i] != 0)
            return false;
    return true;
}

```



PART 03

Yes, we can generate the given sentence from the input data - using 2-Gram -

If we pick 3-Gram, the smallest possible gram with "Lahore" as last word, would be "is from Lahore" and we will not get our desired sentence -

PART 04

Only three values are out of order,
~~so the "Insertion Sort"~~ would be most efficient.

Variables	Array
$i=1 \quad j=1 \quad \text{temp} = 3$	3, 3, 4, 5, 5, 2, 3, 4, 10, 10
$i=2 \quad j=2 \quad \text{temp} = 4$	3, 3, 4, 5, 5, 2, 3, 4, 10, 10
$i=3, 4, \quad j=3, 4 \quad \text{temp} = 5, 5$	3, 3, 4, 5, 5, 2, 3, 4, 10, 10
$i=5, \quad j=5 \quad \text{temp} = 2$	3, 3, 4, 5, 5, <u>2</u> , 3, 4, 10, 10
$i=5, \quad j=4 \quad \text{temp} = 2$	3, 3, 4, 5, <u>5</u> , 2, 3, 4, 10, 10
$i=5, \quad j=3 \quad \text{temp} = 2$	3, 3, 4, <u>4</u> , 5, 2, 3, 4, 10, 10
$i=5, \quad j=2 \quad \text{temp} = 2$	3, 3, <u>3</u> , 4, 5, 2, 3, 4, 10, 10
$i=5, \quad j=1 \quad \text{temp} = 2$	3, <u>3</u> , 3, 4, 5, 2, 3, 4, 10, 10
$i=5, \quad j=0 \quad \text{temp} = 2$	<u>2</u> , 3, 3, 4, 5, 2, 3, 4, 10, 10
$i=6, \quad j=6 \quad \text{temp} = 3$	2, 3, 3, 4, 5, <u>2</u> , 3, 4, 10, 10
$i=6, \quad j=5, 4, 3, \quad \text{temp} = 3$	2, 3, 3, 3, 4, 5, <u>2</u> , 3, 4, 10, 10
$i=7, \quad j=7, 6, 5 \quad \text{temp} = 4$	2, 3, 3, 3, 4, <u>4</u> , 5, 2, 3, 4, 10, 10
$i=8 \quad \text{temp} = 10$	2, 3, 3, 3, 4, 4, 5, 5, 10, 10
$9 \quad \text{temp} = 10$	2, 3, 3, 3, 4, 4, 5, 5, 10, 10

PART 05

<i>n</i>	<i>i</i>	newn	A
8	1	0	5 6 3 1 8 7 2 4
11	2	0	5 <u>3 6</u> 1 8 7 2 4
11	3	2	5 <u>3 1</u> <u>6</u> 8 7 2 4
11	4	3	5 3 1 6 8 7 2 4
11	5	3	5 3 1 <u>6 7</u> <u>8</u> 2 4
11	6	5	5 3 1 6 7 <u>2</u> <u>8</u> 4
11	7	6	5 3 1 6 7 <u>2</u> <u>4</u> 8
7	1	0	<u>3 5</u> 1 6 7 2 <u>4</u> 8
11	2	1	<u>3 1</u> <u>5</u> 6 7 2 4 8
11	3	2	<u>3 1</u> <u>5</u> 6 7 2 4 8
11	4	2	<u>3 1</u> <u>5</u> 6 7 2 4 8
11	5	2	<u>3 1</u> <u>5</u> 6 <u>2</u> <u>2</u> 4 8
11	6	5	<u>3 1</u> <u>5</u> 6 <u>2</u> <u>4</u> 7 8
6	1	0	1 3 5 6 2 4 7 8
11	2	0	1 3 5 <u>6</u> 2 4 7 8
11	3	0	1 3 5 6 2 4 7 8
11	4	0	1 3 5 <u>2</u> <u>6</u> 4 7 8
11	5	4	1 3 5 <u>2</u> <u>4</u> <u>6</u> 7 8
5	1	0	1 3 5 2 4 6 7 8
11	2	0	1 3 5 2 4 6 7 8
11	3	0	1 3 <u>2</u> <u>5</u> 4 6 7 8
11	4	3	1 3 2 4 5 6 7 8
4	1	0	1 3 2 4 <u>5</u> <u>6</u> <u>7</u> <u>8</u>
11	2	0	<u>1</u> 2 3 4 <u>5</u> <u>6</u> <u>7</u> <u>8</u>
11	3	2	<u>1</u> 2 3 4 <u>5</u> <u>6</u> <u>7</u> <u>8</u>
2	1	0	1 2 3 4 <u>5</u> <u>6</u> <u>7</u> <u>8</u>

In case that multiple elements at the end of the list are sorted, then n_{new} will keep the index of last swap. For next iteration of while loop, the n is initialized from n_{new} , i.e. last swap index and results in the improvement as many elements are ignored from the end.

$O(n^2)$ no improvement here (worst case)



PART 06

sorted	i	list	
false		5, 6, 3, 1, 8, 7, 2, 4	Initial
true	1	5, <u>3, 6</u> , 1, 8, 7, 2, 4	
false	3	5, 3, <u>6, 1</u> , 8, 7, 2, 4	
false	5	5, 3, 6, <u>1, 8, 2, 7</u> , 4	
false	0	<u>3, 5</u> , 6, 1, 8, 2, 7, 4	
false	2	3, 5, <u>1, 6</u> , 8, 2, 7, 4	
false	4	3, 5, 1, 6, <u>2, 8</u> , 7, 4	
false	6	3, 5, 1, 6, <u>2, 8, 4, 7</u>	<u>1st While</u>
true	1	3, <u>1, 5</u> , 6, 2, 8, 4, 7	
false	3	3, 1, 5, <u>2, 6</u> , 8, 4, 7	
false	5	3, 1, 5, 2, 6, <u>4, 8, 7</u>	
false	0	<u>1, 3, 5</u> , 2, 6, 4, 8, 7	
false	2	<u>1, 3, 2, 5</u> , 6, 4, 8, 7	
false	4	<u>1, 3, 2, 5, 4</u> , 6, 8, 7	
false	6	<u>1, 3, 2, 5, 4, 6, 7, 8</u>	<u>2nd While</u>
true	1	<u>1, 2, 3, 5, 4, 6, 7, 8</u>	
false	3	<u>1, 2, 3, 4, 5</u> , 6, 7, 8	
false	5	<u>1, 2, 3, 4, 5, 6, 7, 8</u>	
false	0	<u>1, 2, 3, 4, 5, 6, 7, 8</u>	
false	2	<u>1, 2, 3, 4, 5, 6, 7, 8</u>	
false	4, 6	<u>1, 2, 3, 4, 5, 6, 7, 8</u>	<u>3rd While</u>
true	0	"	Sorted
"	3	"	
"	5	"	
"	0	"	
"	2, 4, 6	"	
"	0	"	



Interesting Bubble Sort moves largest and second largest elements towards the end of the list, hence sorts early -

Secondly, it can run two separate sortings on two separate machines (distribu-

$O(n^2)$ still the worst case



PART 07

swapped	start	end	i	a
true	0	7		5, 6, 3, 1, 8, 7, 2, 4
false	0	7	0	5, 6, 3, 1, 8, 7, 2, 4
false	0	7	1	5, <u>3, 6</u> , 1, 8, 7, 2, 4
true	0	7	2	5, <u>3, 1, 6</u> , 8, 7, 2, 4
"	0	7	3	5, <u>3, 1, 6</u> , 8, 7, 2, 4
"	0	7	4	5, <u>3, 1, 6</u> , <u>7, 8</u> , 2, 4
"	0	7	5	5, <u>3, 1, 6</u> , 7, 2, 8, 4
"	0	7	6	5, <u>3, 1, 6</u> , 7, 2, 4, 8 ←
false	0	6	5	5, <u>3, 1, 6</u> , <u>7, 2, 4</u> , 8
"	0	6	4	5, <u>3, 1, 6</u> , <u>2, 7, 4</u> , 8
true	0	6	3	5, <u>3, 1, 2, 6</u> , <u>7, 4, 8</u>
"	0	6	2	5, <u>3, 1, 2, 6</u> , <u>7, 4, 8</u>
"	0	6	1	5, <u>1, 3</u> , <u>2, 6</u> , <u>7, 4, 8</u> ←
"	0	6	0	1, <u>5, 3</u> , <u>2, 6</u> , <u>7, 4, 8</u>
false	1	6	1	1, <u>3, 5</u> , <u>2, 6</u> , <u>7, 4, 8</u> .
true	1	6	2	1, <u>3, 2, 5</u> , <u>6, 7, 4, 8</u>
"	1	6	3	1, <u>3, 2, 5, 6</u> , <u>7, 4, 8</u>
"	1	6	4	1, <u>3, 2, 5, 6</u> , <u>7, 4, 8</u>
"	1	6	5	1, <u>3, 2, 5, 6, 4</u> , <u>7, 8</u> ←
false	1	5	4	1, <u>3, 2, 5, 4, 6</u> , <u>7, 8</u>
true	1	5	3	1, <u>3, 2, 4, 5</u> , <u>6, 7, 8</u>
"	1	5	2	1, <u>3, 2, 4, 5, 6</u> , <u>7, 8</u>
"	1	5	1	1, <u>2, 3, 4, 5, 6</u> , <u>7, 8</u> ←
false	2	5	2, 3, 4	1, <u>2, 3, 4, 5, 6</u> , <u>7, 8</u> .



Cocktail sort is better as it moves (down moving) smaller elements towards the start in every reverse loop, hence terminates early.
(Turtles Vs. Rabbits concept)

$O(n^2)$ is still the worst case -