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# **Environment**

#### 1.1 Vimrc

```
set ru nu ts=4 sts=4 sw=4 si sm hls is ar bs=2 mouse=a syntax on nm <F3> :vsplit %<.in <CR>
4 nm <F4> :!gedit % <CR>
5 au BufEnter *.cpp set cin
6 au BufEnter *.cpp nm <F5> :!time ./%< <CR>|nm <F7> :!
        gdb ./%< <CR>|nm <F8> :!time ./%< < %<.in <CR>|nm <F9> :!g++ % -o %< -g -std=gnu++14 -O2 -DLOCAL && size %< <CR>
 7 au BufEnter *.java nm <F5> :!time java %< <CR>|nm <F8>
:!time java %< < %<.in <CR>|nm <F9> :!javac % <CR
```

# Data Structure

#### 2.1 KD tree

```
/* kd_tree : finds the k-th closest point in O(kn^{1-\frac{1}{k}}). Usage : Stores the data in p[]. Call function init (n, k). Call min_kth (d, k). (or max_kth) (k is 1-based)
     2 Usage
 | based)
| Note : Switch to the commented code for Manhattan distance.
| Status : SPOJ-FAILURE Accepted.*/
| template <int MAXN = 200000, int MAXK = 2>
| struct kd_tree {
| int k, size;
| struct point { int data[MAXK], id; } p[MAXN];
| struct kd_node {
| int l, r; point p, dmin, dmax;
| kd_node (const point &rhs) : l (-1), r (-1), p (rhs) |
| , dmin (rhs), dmax (rhs) {}
| void merge (const kd_node &rhs, int k) {
| for (register int i = 0; i < k; ++i) {
| dmin.data[i] = std::max (dmax.data[i], rhs.dmin. data[i]); }
| long long min_dist (const point &rhs, int k) const {
| register long long ret = 0; |
| for (register int i = 0; i < k; ++i) {
| if (dmin.data[i] = rhs.data[i] & rhs.data[i] <= dmax.data[i] > continue; |
| ret += std::min (111 * (dmin.data[i] - rhs.data[i]) * (dmax. data[i] - rhs.data[i]) * (dmax. data[i] - rhs.data[i]) * (ret += std::max (0, rhs.data[i] - rhs.data[i]); |
| ret += std::max (0, dmin.data[i] - rhs.data[i]); |
| ret trun ret; |
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long ret = 0; |
| for (int i = 0; i < k; ++i) {
| int tmp = std::max (std::abs (dmin.data[i] - rhs.data[i]) |
| ret += std::max (std::abs (dmin.data[i] - rhs.data[i]) |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] |
| r
      3 Note
                                                           Switch to the commented code for Manhattan
                                            distance.
21
23 //
28
                                         | ]));
ret += 111 * tmp * tmp; }
  ret += std::max (std::abs (rhs.data[i] - dmax.
data[i]) + std::abs (rhs.data[i] - dmin.data[i]));
                   return ret; } tree[MAXN * 4];
struct result {
 long long dist; point d; result() {}
 result (const long long &dist, const point &d) :
    dist (dist), d (d) {}
 bool operator > (const result &rhs) const { return
 32
34
                35
 38
                             if ("tree[rt].r) tree[rt].merge (tree[tree[rt].r], k
                  if (~tree[rt].r) tree[rt].merge (tree[tree[rt].r], k
    ); }
std::priority_queue<result, std::vector<result>, std
    ::less <result>> heap_1;
std::priority_queue<result, std::vector<result>, std
    ::greater <result>> heap_r;
void _min_kth (const int &depth, const int &rt, const
    int &m, const point &d) {
    result tmp = result (sqrdist (tree[rt].p, d), tree[
        rt].p);
    if ((int)heap_1.size() < m) heap_1.push (tmp);
    else if (tmp < heap_1.top()) {
        heap_1.push (tmp); }</pre>
 55
```

```
62
74
75
80
```

#### Splay 2.2

```
void push_down (int x) {
   if ( n[x].c[0]) push (n[x].c[0], n[x].t);
   if ( n[x].c[1]) push (n[x].c[1], n[x].t);
   if ( n[x].t = tag (); }
   void update (int x) {
      n[x].m = gen (x);
   if ( n[x].c[0]) n[x].m = merge (n[n[x].c[0]].m, n[x].
    m);
if (\tilde{n}[x].c[1]) n[x].m = merge (n[x].m, n[n[x].c[1]].
  11
12
```

### Link-cut tree

```
= u;

n[u].c[1] = v;

if (~v) n[v].f = u, n[v].p = -1;

update (u); u = n[v = u].p; }

splay (x); }
```

#### Formula

#### Zellers congruence 3.1

```
/* Zeller's congruence : converts between a calendar
```

#### 3.2 Lattice points below segment

```
1 /* Euclidean-like algorithm : computes the sum of
     \sum_{i=0}^{n-1} \left[ \frac{a+bi}{m} \right] . \star /
```

# 3.3 Adaptive Simpson's method

```
/* Adaptive Simpson's method : integrates f in [1, r].

*/

struct simpson {

double area (double (*f) (double), double 1, double r

} {

double m = 1 + (r - 1) / 2;

return (f (1) + 4 * f (m) + f (r)) * (r - 1) / 6; }

double solve (double (*f) (double), double 1, double

r, double eps, double a) {

double m = 1 + (r - 1) / 2;

double left = area (f, 1, m), right = area (f, m, r)

if (fabs (left + right - a) <= 15 * eps) return left

+ right + (left + right - a) / 15.0;

return solve (f, 1, m, eps / 2, left) + solve (f, m,

r, eps / 2, right); }

double solve (double (*f) (double), double 1, double

r, double eps) {

return solve (f, 1, r, eps, area (f, 1, r)); };
```

#### 3.4 Neural network

```
1 /* Neural network : machine learning. */
2 template <int ft = 6, int n = 6, int MAXDATA = 100000>
3 struct network {
      double wp[n][ft], w[n], avg[ft + 1], sig[ft + 1], val
        network () {
std::mt19937_64 mt (time (0));
std::uniform_real_distribution <double> urdp (0, 2 *
sqrt (ft));
      network
        double
     void desc (double *x, double t, double eta) {
  double o = compute (x), del = (o - t); // * o * (1 -
        o)
for (int i = 0; i < n; ++i) for (int j = 0; j < ft;
21
     22
26
     dn; ++j)
sig[i] += (data[j][i] - avg[i]) * (data[j][i] - avg[i]);
for (int i = 0; i < ft + 1; ++i) sig[i] = sqrt (sig[i] / dn);
for (int i = 0; i < ft + 1; ++i) for (int j = 0; j < dn; ++j)
data[j][i] = (data[j][i] - avg[i]) / sig[i];
for (int cnt = 0; cnt < epoch; ++cnt) for (int test = 0; test < dn; ++test)
desc (data[test], data[test][ft], eta);
double predict (double *x) {
  for (int i = 0; i < ft; ++i) x[i] = (x[i] - avg[i])
  / sig[i];
  return compute (x) * sig[ft] + avg[ft]; }
</pre>
33
      return compute (x) * sig[ft] + avg[ft]; }
std::string to_string () {
  std::ostringstream os; os.precision (20); os << std</pre>
        42
        for (int i = 0; i < ft + 1; ++i) os << sig[i] << "_"
     return os.str (); }
void read (const std::string &str) {
    std::istringstream is (str);
    for (int i = 0; i < n; ++i) for (int j = 0; j < ft;
        ++j) is >> wp[i][j];
    for (int i = 0; i < n; ++i) is >> w[i];
    for (int i = 0; i < ft + 1; ++i) is >> avg[i];
    for (int i = 0; i < ft + 1; ++i) is >> sig[i]; } };
```

# 4 Number theory

# 4.1 Fast power module

```
/* Fast power module : x^n */ / int fpm (int x, int n, int mod) {
```

```
int ans = 1, mul = x; while (n) {
if (n & 1) ans = int (111 * ans * mul % mod);
mul = int (111 * mul * mul % mod); n >>= 1; }
return ans; }
```

## 4.2 Euclidean algorithm

```
/* Euclidean algorithm : solves for ax + by = gcd (a, b). */
void euclid (const long long &a, const long long &b, long long &x, long long &y) {
   if (b == 0) x = 1, y = 0;
   else euclid (b, a % b, y, x), y -= a / b * x; }
   6 long long inverse (long long x, long long m) {
      long long a, b; euclid (x, m, a, b); return (a % m + m) % m; }
```

#### 4.3 Discrete Fourier transform

### 4.4 Number theoretic transform

#### 4.5 Chinese remainder theorem

```
/* Chinese remainder theroem : finds positive integers
    x = out.first + k * out.second that satisfies x %
    in[i].second = in[i].first. */
2 struct crt {
3 long long fix (const long long &a, const long long &b
    ) { return (a % b + b) % b; }
```

#### 4.6 Linear Recurrence

#### 4.7 Berlekamp Massey algorithm

# 4.8 Baby step giant step algorithm

```
| /* Baby step giant step algorithm : Solves a^x = b \mod c in O(\sqrt{c}). */
| struct bsgs {
| int solve (int a, int b, int c) {
| std::unordered_map <int, int> b;
| int m = (int) sqrt ((double) c) + 1, res = 1;
| for (int i = 0; i < m; ++i) {
| if (bs.find (res) == bs.end ()) bs[res] = i;
| res = int (1LL * res * a % c);
| int mul = 1, inv = (int) inverse (a, c);
| for (int i = 0; i < m; ++i) mul = int (1LL * mul * inv % c);
| res = b % c;
| for (int i = 0; i < m; ++i) {
| if (bs.find (res) != bs.end ()) return i * m + bs[ res];
| res = int (1LL * res * mul % c); }
| return -1; };
```

#### 4.9 Miller Rabin primality test

#### 4.10 Pollard's Rho algorithm

# 5 Geometry

```
#define cd const double &
const double EPS = 1E-8, PI = acos (-1);
int sgn (cd x) { return x < -EPS ? -1 : x > EPS; }
int cmp (cd x, cd y) { return sgn (x - y); }
double sqr (cd x) { return x * x; }
```

#### 5.1 Point

#### **5.2** Line

#### 5.3 Circle

#### 5.4 Centers of a triangle

```
point incenter (cp a, cp b, cp c) {
   double p = dis (a, b) + dis (b, c) + dis (c, a);
   return (a * dis (b, c) + b * dis (c, a) + c * dis (a, b)) / p; }

   point circumcenter (cp a, cp b, cp c) {
      point p = b - a, q = c - a, s (dot (p, p) / 2, dot (q, q) / 2);
   return a + point (det (s, point (p.y, q.y)), det (point (p.x, q.x), s)) / det (p, q); }

   point orthocenter (cp a, cp b, cp c) { return a + b + c - circumcenter (a, b, c) * 2; }
```

#### 5.5 Fermat point

```
| /* Fermat point : finds a point P that minimizes | PA| + |PB| + |PC| . */
| point fermat_point (cp a, cp b, cp c) {
| if (a == b) return a; if (b == c) return b; if (c == a) return c;
| double ab = dis (a, b), bc = dis (b, c), ca = dis (c, a);
| double cosa = dot (b - a, c - a) / ab / ca;
| double cosb = dot (a - b, c - b) / ab / bc;
| double cosc = dot (b - c, a - c) / ca / bc;
| double sq3 = PI / 3.0; point mid;
| if (sgn (cosa + 0.5) < 0) mid = a;
| else if (sgn (cosb + 0.5) < 0) mid = b;
| else if (sgn (cosc + 0.5) < 0) mid = c;
| else if (sgn (det (b - a, c - a)) < 0) mid = line_intersect (line (a, b + (c - b).rot (sq3)), line (b, c + (a - c).rot (sq3)));
| else mid = line_intersect (line (a, c + (b - c).rot (sq3)), line (c, b + (a - b).rot (sq3)));
| return mid; }
```

## 5.6 Convex hull

#### 5.7 Half plane intersection

1 /\* Online half plane intersection : complexity O(n) each operation. \*/

```
std::vector <point> ret;
if (c.empty ()) return ret;
for (int i = 0; i < (int) c.size (); ++i) {
  int j = (i + 1) % (int) c.size ();
  if (turn_left (p.s, p.t, c[i])) ret.push_back (c[i]);</pre>
      if (two_side (c[i], c[j], p)) ret.push_back (
    line_intersect (p, line (c[i], c[j]))); }
return ret; }
        eturn ret; }
Offline half plane intersection : complexity
13 std::vector
                          <point> half_plane_intersect (std::vector
     16
     second.t - a.second.s, b.second.t - a.second.s)
    < 0;
else return cmp (a.first, b.first) < 0; });
h.resize (std::unique (g.begin (), g.end (), [] (
    const polar &a, const polar &b) { return cmp (a.
    first, b.first) == 0 }) - g.begin ());
for (int i = 0; i < (int) h.size (); ++i) h[i] = g[i]
    l second:</pre>
     24
25
     ret[++rear] = h[i]; }
while (rear - fore > 1 && !turn_left (ret[fore],
    line_intersect (ret[rear - 1], ret[rear])))
    rear;
     while (rear - fore > 1 && !turn_left (ret[rear],
    line_intersect (ret[fore], ret[fore + 1]))) ++
     if (rear - fore < 2) return std::vector <point> ();
std::vector <point> ans; ans.resize (rear + 1);
for (int i = 0; i < rear + 1; ++i) ans[i] =
    line_intersect (ret[i], ret[(i + 1) % (rear + 1)</pre>
31
```

#### 5.8 Nearest pair of points

```
/* Nearest pair of points : [l, r), need to sort p
first. */
double solve (std::vector <point> &p, int l, int r) {
   if (l + 1 >= r) return INF;
   int m = (l + r) / 2; double mx = p[m].x; std::vector
   <point> v;
double ret = std::min (solve(p, l, m), solve(p, m, r)
   );
for (int i = l; i < r; ++i)
   if (sqr (p[i].x - mx) < ret) v.push_back (p[i]);
sort (v.begin (), v.end (), [&] (cp a, cp b) { return
        a.y < b.y; });
for (int i = 0; i < v.size (); ++i)
   for (int j = i + 1; j < v.size (); ++j) {
      if (sqr (v[i].y - v[j].y) > ret) break;
      ret = min (ret, dis2 (v[i] - v[j])); }
   return ret; }
```

#### 5.9 Minimum circle

```
circle minimum_circle (std::vector <point> p) {
circle ret; std::random_shuffle (p.begin (), p.end ()
    );
}
for (int i = 0; i < (int) p.size (); ++i) if (!
    in_circle (p[i], ret)) {
    ret = circle (p[i], 0); for (int j = 0; j < i; ++j)
        if (!in_circle (p[j], ret)) {
        ret = make_circle (p[j], p[i]); for (int k = 0; k < j; ++k)
        if (!in_circle (p[k], ret)) ret = make_circle (p[i], p[j], p[j], p[k]);
}
return ret; }</pre>
```

#### 5.10 Intersection of a polygon and a circle

```
struct polygon_circle_intersect {
    double sector_area (cp a, cp b, const double &r) {
        double c = (2.0 * r * r - dis2 (a, b)) / (2.0 * r *
            r);
        return r * r * acos (c) / 2.0; }
    double da = (cp a, cp b, const double &r) {
        double dA = dot (a, a), dB = dot (b, b), dC =
            point_to_segment (point (), line (a, b));
        if (sgn (dA - r * r) <= 0 && sgn (dB - r * r) <= 0)
            return det (a, b) / 2.0;
        point tA = a.unit () * r, tB = b.unit () * r;
        if (sgn (dC - r) > 0) return sector_area (tA, tB, r)
        std::vector <point> ret = line_circle_intersect (
            line (a, b), circle (point (), r);
        if (sgn (dA - r * r) > 0 && sgn (dB - r * r) > 0)
```

#### 5.11 Union of circles

#### 6 Graph

#### 6.1 Hopcoft-Karp algorithm

# 6.2 Kuhn-Munkres algorithm

/\* Kuhn Munkres algorithm : weighted maximum ming algorithm for bipartition graphs with complexity  $O(N^3)$ .

#### 6.3 Blossom algorithm

```
_{1}^{\dagger}/\star Blossom algorithm : maximum match for general graph
       int *qhead, *qtail;
struct {
  int fa[MAXN];
  void init (int n) { for (int i = 1; i <= n; i++) fa[i
    ] = i;
  int find (int x) { if (fa[x] != x) fa[x] = find (fa[x]); return fa[x]; }
  void merge (int x, int y) { x = find (x); y = find (y); fa[x] = y; } ufs;
  void solve (int x, int y) {
  if (x == y) return;
  if (d[y] == 0) {
    solve (x, fa[fa[y]]); match[fa[y]] = fa[fa[y]];
    match[fa[fa[y]]] = fa[y];
  } else if (d[y] == 1) {
    solve (match[y], c1[y]); solve (x, c2[y]);
    match[c1[y]] = c2[y]; match[c2[y]] = c1[y]; }
  int lca (int x, int y, int root) {
    x = ufs.find (x); y = ufs.find (y);
    while (x != y && v[x] != 1 && v[y] != 0) {
      v[x] = 0; v[y] = 1;
      if (x != root) x = ufs.find (fa[x]);
      if (y != root) y = ufs.find (fa[y]);
      if (v[y] == 0) std::swap (x, y);
      for (int i = x; i != y; i = ufs.find (fa[i])) v[i] =
      v[y] = -1; return x; }
</pre>
        ufs.merge (i, b);
if (d[i] == 1) { c1[i] = x; c2[i] = y; *qtail++ = i
    ; } }
31
32
       37
                  38
           match[dest] = loc; return 1;
} else {
  fa[dest] = loc; fa[match[dest]] = dest;
  d[dest] = 1; d[match[dest]] = 0;
  *qtail++ = match[dest];
} else if (d[ufs.find (dest)] == 0) {
  int b = lca (loc, dest, root);
  contract (loc, dest, b); contract (dest, loc, b)
   ; } }
return 0; }
int solve (int n, const edge_list &e) {
  std::fill (fa, fa + n, 0); std::fill (c1, c1 + n, 0)
}
            std:'fill (c2, c2 + n, 0); std::fill (match, match + n, -1);
54
           std::fill (62, 62 ...,
    n, -1);
int re = 0; for (int i = 0; i < n; i++)
if (match[i] == -1) if (bfs (i, n, e)) ++re; else
    match[i] = -2;
return re; };</pre>
```

## 6.4 Weighted blossom algorithm

```
1 /* Weighted blossom algorithm (vfleaking ver.) : maximum matching for general weighted graphs with complexity O(n^3).
```

```
: Set n to the size of the vertices. Run init () Set g[][].w to the weight of the edge. Run solve
 2 Usage :
 The first result is the answer, the second one is the number of matching pairs. Obtain the matching with match[].
 int w = 0): d(d), v(v), w(w), (, ),
int n, n_x;
edge g[MAXN * 2 + 1] [MAXN * 2 + 1];
int lab[MAXN * 2 + 1], match[MAXN * 2 + 1], slack[
MAXN * 2 + 1], st[MAXN * 2 + 1], pa[MAXN * 2 +
      12
23
      void augment (int u, int v) {
for (; ; ) {
  int xnv = st[match[u]]; set_match (u, v);
  if (!xnv) return; set_match (xnv, st[pa[xnv]]);
  u = st[pa[xnv]], v = xnv; }
int get_lca (int u, int v) {
  static int t = 0;
  for (++t; u || v; std::swap (u, v)) {
    if (u == 0) continue; if (vis[u] == t) return u;
    vis[u] = t; u = st[match[u]]; if (u) u = st[pa[u]];
    return 0: }

40
     63
        int xr = flower_irom[b][g[b][pa[b]].u], r-
b, xr);
for (int i = 0; i < pr; i += 2) {
  int xs = flower[b][i], xns = flower[b][i + 1];
  pa[xs] = g[xns][xs].u; S[xs] = 1, S[xns] = 0;
  slack[xs] = 0, set_slack(xns); q_push(xns);
} S[xr] = 1, pa[xr] = pa[b];
for {size_t i = pr + 1; i < flower[b].size (); ++i)</pre>
      int xs = flower[b][i]; S[xs] = -1, set_slack(xs); }
st[b] = 0; }
bool on_found_edge (const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1; int nu = st[match[v]];
    slack[v] = slack[nu] = 0; S[nu] = 0, q_push(nu);
} else if(S[v] == 0) {
  int lca = get_lca(u, v);
  if (!lca) return augment(u, v), augment(v, u), true
}
80
         else add_blossom(u, lca, v); }
return false; }
```

```
101
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            return false; }
std::pair <long long, int> solve () {
   memset (match + 1, 0, sizeof (int) * n); n_x = n;
   int n_matches = 0; long long tot_weight = 0;
   for (int u = 0; u <= n; ++u) st[u] = u, flower[u].
        clear();
   int w_max = 0;
   for (int u = 1; u <= n; ++u) for (int v = 1; v <= n;
        ++v) {
      flower_from[u][v] = (u == v ? u : 0); w_max = std::
            max (w_max, g[u][v].w); }
   for (int u = 1; u <= n; ++u) lab[u] = w_max;
   while (matching ()) ++n_matches;
   for (int u = 1; u <= n; ++u) if (match[u] && match[u] < u) tot_weight += g[u][match[u]].w;
   return std::make_pair (tot_weight, n_matches); }
void init () { for (int u = 1; u <= n; ++u) for (int v = 1; v <= n; ++v) g[u][v] = edge (u, v, 0); }
};</pre>
116
122
123
```

#### 6.5 Maximum flow

```
/* Sparse graph maximum flow : isap.*/
2 template <int MAXN = 1000, int MAXM = 100000>
3 struct isap {
4 struct flow_edge_list {
5 int size, begin[MAXN], dest[MAXM], next[MAXM], flow[
MAXM];
     20
25
se struct dinic {
struct flow_edge_list {
struct size, begin[MAXN], dest[MAXM], next[MAXM], flow[
MAXM];
     void clear (int n) { size = 0; std::fill (begin,
    begin + n, -1); }
flow_edge_list (int n = MAXN) { clear (n); }
```

```
42
48
50
51
      int ans = 0; n = n_; s = s_; dinic::t = t_;
     int ans = 0, in = in_, s = s_, dinfe..t = t_,
while (bfs (e)) {
  for (int i = 0; i < n; ++i) w[i] = e.begin[i];
  ans += dfs (e, s, INF); }
return ans; } };</pre>
```

#### 6.6Minimum cost flow

21

23

47

51

```
/* Sparse graph minimum cost flow : EK. */
template <int MAXN = 1000, int MAXM = 100000>
struct minimum_cost_flow {
   ans.second += num * e.cost[prev[i]]; } return ans; } ;;
/* Dense graph minimum cost flow : zkw. */
template <int MAXN = 1000, int MAXM = 100000>
  40
46
    MAXN];
int modlable() {
  int delta = INF;
  for (int i = 0; i < n; i++) {
    if (!visit[i] && slack[i] < delta) delta = slack[i]</pre>
      delta;
return 0; }
int dfs (cost_flow_edge_list &e, int x, int flow) {
  if (x == t) { tf += flow; tc += flow * (dis[s] - dis
       [t]); return flow; }
  visit[x] = 1; int left = flow;
  for (int i = e.begin[x]; ~i; i = e.next[i])
  if (e.flow[i] > 0 && !visit[e.dest[i]]) {
    int y = e.dest[i];
```

```
if (dis[y] + e.cost[i] == dis[x]) {
  int delta = dfs (e, y, std::min (left, e.flow[i])
              -= delta; e.flow[i ^ 1] += delta; left
    e.flów[i]
    -= delta;
if (!left) { visit[x] = false; return flow; }
```

## 6.7 Stoer Wagner algorithm

```
/* Stoer Wagner algorithm : Finds the minimum cut of
    an undirected graph. (1-based) */
template <int MAXN = 500>
struct stoer wagner {
    int n, edge[MAXN] [MAXN];
    int dist[MAXN],
    bool vis[MAXN], bin[MAXN];
    stoer_wagner () {
        memset (edge, 0, sizeof (edge));
        memset (bin, false, sizeof (bin)); }
    int contract (int &s, int &t) {
        memset (dist, 0, sizeof (dist));
        memset (vis, false, sizeof (vis));
        int i, j, k, mincut, maxc;
        for (i = 1; i <= n; i++) {
            k = -1; maxc = -1;
            for (j = 1; j <= n; j++)
            if (!bin[j] && !vis[j] && dist[j] > maxc) {
            k = j; maxc = dist[j]; }
            if (k == -1) return mincut;
            s = t; t = k; mincut = maxc; vis[k] = true;
            for (j = 1; j <= n; j++) if (!bin[j] && !vis[j])
            dist[j] += edge[k][j]; }
            return mincut; }
            int mincut, i, j, s, t, ans;
            for (mincut = INF, i = 1; i < n; i++) {
                 ans = contract (s, t); bin[t] = true;
            if (mincut > ans) mincut = ans;
            if (mincut = 0) return 0;
            for (j = 1; j <= n; j++) if (!bin[j])
            edge[s][j] = (edge[j][s] += edge[j][t]); }
            return mincut; };
}</pre>
1 /* Stoer Wagner algorithm : Finds the minimum cut of
```

#### 6.8 DN maximum clique

```
/* DN maximum clique : n <= 150 */
typedef bool BB[N]; struct Maxclique {
const BB *e; int pk, level; const float Tlimit;
struct Vertex { int i, d; Vertex (int i) : i(i), d(0)
{ };
typedef std::vector <Vertex> Vertices; Vertices V;
typedef std::vector <int> ColorClass; ColorClass QMAX,
O:
39
           ()) {
Q.push_back (R.back ().i); Vertices Rp; cut2 (R, Rp
           if ((int) Rp.size ()) {
  if((float) S[level].i1 / ++pk < Tlimit)
    degree_sort (Rp);
  color_sort (Rp); ++S[level].i1, ++level;
  expand_dyn (Rp); --level;</pre>
 42
```

```
57
58 BB e[N]; in
= true;
            int ans, sol[N]; for (...) e[x][y] = e[y][x]
set true;
set Maxclique mc (e, n); mc.mcqdyn (sol, ans); //0-based.
set for (int i = 0; i < ans; ++i) std::cout << sol[i] <</pre>
       std::endl;
```

#### 6.9 Dominator tree

```
/* Dominator tree : finds the immediate dominator (
    idom[]) of each node, idom[x] will be x if x does
    not have a dominator, and will be -1 if x is not
    reachable from s. */

template <int MAXN = 100000, int MAXM = 100000>

struct dominator tree /
if (dfn[p] > dfn[x]) { getfa (p); p = sdom[smin[p]]; }
  if (dfn[sdom[x]] > dfn[p]) sdom[x] = p; }
  tmp[sdom[x]].push (x); }
while (!tmp[x].empty ()) {
  int y = tmp[x].front (); tmp[x].pop (); getfa (y);
  if (x != sdom[smin[y]]) idom[y] = smin[y];
  else idom[y] = x; }
  for (int v : succ[x]) if (f[v] == x) fa[v] = x; }
  idom[s] = s; for (int i = 1; i < stamp; ++i) {
  int x = id[i]; if (idom[x] != sdom[x]) idom[x] =
      idom[idom[x]]; } };</pre>
```

#### String Suffix Array

```
/* Suffix Array : sa[i] - the beginning position of the ith smallest suffix, rk[i] - the rank of the suffix beginning at position i. height[i] - the longest common prefix of sa[i] and sa[i - 1]. */

template <int MAXN = 1000000, int MAXC = 26>

struct suffix array {
    for (int i = n - 1; i \ge 0; --i) sa[--sRank[x[i]]] =
      for (int i = n - 1; i >= 0; --i) sa[--sRank[x[y[i ]]]] = y[i];
std::swap (x, y); x[sa[0]] = 0; p = 1; y[n] = -1;
for (int i = 1; i < n; ++i)
20
```

44

46

```
x[sa[i]] = cmp (y, sa[i], sa[i - 1], d) ? p - 1 :
    p++; } }
void solve (int *a, int n) {
    a[n] = -1; doubling (a, n);
    for (int i = 0; i < n; ++i) rk[sa[i]] = i;
    int cur = 0;
    for (int i = 0; i < n; ++i)
    if (rk[i]) {
        if (cur) cur--:</pre>
23
                if (cur) cur--;
for (; a[i + cur] == a[sa[rk[i] - 1] + cur]; ++cur
                height[rk[i]] = cur; } };
```

#### 7.2 Suffix Automaton

```
/* Suffix automaton : head - the first state. tail -
the last state. Terminating states can be reached
via visiting the ancestors of tail. state::len -
the longest length of the string in the state.
state::right - 1 - the first location in the
string where the state can be reached. state::
parent - the parent link. state::dest - the
automaton link. */
template <int MAXN = 1000000, int MAXC = 26>
struct suffix_automaton {
struct state {
int len. right: state *parent. *dest[MAXC]:
                                  int len, right; state *parent, *dest[MAXC];
state (int len = 0, int right = 0) : len (len),
    right (right), parent (NULL) {
    memset (dest, 0, sizeof (dest)); }
} node_pool[MAXN * 2], *tot_node, *null = new state()
                                state *head, *tail;
void extend (int token) {
    state *p = tail;
    state *np = tail -> dest[token] ? null : new (
        tot_node++)    state (tail -> len + 1, tail -> len
        + 1);
    while (p && !p -> dest[token]) p -> dest[token] = np
        , p = p -> parent;
    if (!p) np -> parent = head;
    else {
        state *g = p -> dest[token]:
11
                                              if (!p' np' -> parent = head;
else {
    state *q = p -> dest[token];
    if (p -> len + 1 == q -> len) {
        np -> parent = q;
    } else {
        state *nq = new (tot_node++) state (*q);
        nq -> len = p -> len + 1;
        np -> parent = q -> parent = nq;
        while (p && p -> dest[token] == q) {
            p -> dest[token] == q) {
            p -> parent = nq;
            print = q -> parent = nq;
            print = q -> parent = nq;
            print = nq;

                                tail | np == null ? np -> parent : np; }
void init () {
  tot_node = node_pool;
  head = tail = new (tot_node++) state (); }
suffix_automaton () { init (); } };
```

#### 7.3 Palindromic tree

```
12
         now = now -> fail);
return now; }
bool extend (int token) {
  text[++size] = token; node *now = match (last);
  if (now -> child[token])
  return last = now -> child[token], false;
  last = now -> child[token] = new (tot_node++) node (
    now -> len + 2);
  if (now == odd) last -> fail = even;
  clast = now fail = even;
18
        if (now == odd) last -> fail = even;
else {
  now = match (now -> fail);
  last -> fail = now -> child[token]; }
  return true; }
void init() {
  text[size = 0] = -1; tot_node = node_pool;
  last = even = new (tot_node++) node (0); odd = new (
    tot_node++) node (-1);
  even -> fail = odd; }
palindromic_tree () { init (); } };
```

#### 7.4 Regular expression

```
std::string str = ("The_the_there");
std::regex pattern ("(th|Th)[\\w]*", std::
    regex_constants::optimize | std::regex_constants::
    ECMAScript);
std::smatch match; //std::cmatch for char *
s| std::regex_match (str, match, pattern);
```

```
| 10| for (std::sregex_iterator i = mbegin; i != mend; ++i)
```

#### Tips 8 8.1 Java

```
/* Java reference : References on Java IO, structures,
  doubleValue () / toPlainString () : converts to other
  types.

20 Arrays: Arrays.sort (T [] a); Arrays.sort (T [] a, int fromIndex, int toIndex); Arrays.sort (T [] a, int fromIndex, int toIndex, Comperator <? super T>
 int fromIndex, int toindex, Comperator <? super i>
comperator);
21 LinkedList <E> : addFirst / addLast (E) / getFirst /
getLast / removeFirst / removeLast () / clear () /
add (int, E) / remove (int) / size () / contains
/ removeFirstOccurrence / removeLastOccurrence (E)
22 ListIterator <E> listIterator (int index) : returns an
iterator :
 iterator :
23    E next / previous () : accesses and iterates.
24    hasNext / hasPrevious () : checks availablity.
25    nextIndex / previousIndex () : returns the index of a subsequent call.
26    add / set (E) / remove () : changes element.
27    PriorityQueue <E> (int initcap, Comparator <? super E> comparator) : add (E) / clear () / iterator () / peek () / poll () / size ()
28    TreeMap <K, V> (Comparator <? super K> comparator) :
        Map.Entry <K, V> ceilingEntry / floorEntry / higherEntry / lowerEntry (K): getKey / getValue () / setValue (V) : entries.
29    clear () / put (K, V) / get (K) / remove (K) / size ()
30    StringBuilder : StringBuilder (string) / append (int,
                       iterator :
 ()
30 StringBuilder: StringBuilder (string) / append (int, string, ...) / insert (int offset, ...) charAt (int) / setCharAt (int, char) / delete (int, int) / reverse () / length () / toString ()
31 String: String.format (String, ...) / toLowerCase / toUpperCase () */
32 /* Examples on Comparator:
33 public class Main {
34  public static class Point {
35  public int x; public int y;
             public int x; public int y;
public Point () {
  x = 0;
  y = 0; }
public Point (int xx, int yy) {
  x = xx;
                             yy; } };
c static class Cmp implements Comparator <Point>
              public int compare (Point a, Point b) {
         public static class Point implements Comparable <</pre>
             Point> {
public int x; public int y;
```

#### 8.2 Random numbers

```
std::mt19937_64 mt (time (0));
std::uniform_int_distribution <int> uid (1, 100);
std::uniform_real_distribution <double> urd (1, 100);
std::cout << uid (mt) << "\_" << urd (mt) << "\n";
```

#### 8.3 Read hack

#### 8.4 Stack hack

```
1 //C++
2 #pragma comment (linker, "/STACK:36777216")
3 //G++
4 int __size__ = 256 << 20;
5 char *_p__ = (char*) malloc(__size__) + __size__;
6 __asm__ ("movl_%0,_%%esp\n" :: "r"(_p__));</pre>
```

#### 8.5 Time hack

```
clock_t t = clock ();
z std::cout << 1. * t / CLOCKS_PER_SEC << "\n";
```

#### 8.6 Builtin functions

- \_builtin\_clz: Returns the number of leading 0-bits in x, starting at the most significant bit position. If x is 0, the result is undefined.
- undefined.

  \_builtin\_ctz: Returns the number of trailing 0-bits in x, starting at the least significant bit position. If x is 0, the result is undefined
- undefined.

  3. \_builtin\_clrsb: Returns the number of leading redundant sign bits in x, i.e. the number of bits following the most significant bit that are identical to it. There are no special cases for 0 or other values.
- other values.
  4. \_\_builtin\_popcount: Returns the number of 1-bits in x.
- 5. \_builtin\_parity: Returns the parity of x, i.e. the number of 1-bits in x modulo 2.
  6. \_builtin\_bswap16, \_builtin\_bswap32, \_builtin\_bswap64:
- b. \_\_bulltin\_bswap16, \_\_bulltin\_bswap32, \_\_bulltin\_bswap64:
  Returns x with the order of the bytes (8 bits as a group) reversed.

  7 bitset is Find first () bitset x Find psy (id); bit
- 7. bitset::Find\_first(), bitset::Find\_next(idx): bitset built-in functions

# 8.7 Prufer sequence

In combinatorial mathematics, the Prufer sequence of a labeled tree is a unique sequence associated with the tree. The sequence for a tree on n vertices has length n-2.

One can generate a labeled tree's Prufer sequence by iteratively removing vertices from the tree until only two vertices remain. Specifically, consider a labeled tree T with vertices 1, 2, ..., n. At step i, remove

the leaf with the smallest label and set the ith element of the Prufer sequence to be the label of this leaf's neighbour.

One can generate a labeled tree from a sequence in three steps. The tree will have n+2 nodes, numbered from 1 to n+2. For each node set its degree to the number of times it appears in the sequence plus 1. Next, for each number in the sequence a[i], find the first (lowest-numbered) node, j, with degree equal to 1, add the edge (j, a[i]) to the tree, and decrement the degrees of j and a[i]. At the end of this loop two nodes with degree 1 will remain (call them u, v). Lastly, add the edge (u, v) to the tree.

The Prufer sequence of a labeled tree on n vertices is a unique sequence of length n-2 on the labels 1 to n- this much is clear. Somewhat less obvious is the fact that for a given sequence S of length n-2 on the labels 1 to n, there is a unique labeled tree whose Prufer sequence is S.

#### 8.8 Spanning tree counting

**Kirchhoff's Theorem**: the number of spanning trees in a graph G is equal to *any* cofactor of the Laplacian matrix of G, which is equal to the difference between the graph's degree matrix (a diagonal matrix with vertex degrees on the diagonals) and its adjacency matrix (a (0,1)-matrix with 1's at places corresponding to entries where the vertices are adjacent and 0's otherwise).

The number of edges with a certain weight in a minimum spanning tree is fixed given a graph. Moreover, the number of its arrangements can be obtained by finding a minimum spanning tree, compressing connected components of other edges in that tree into a point, and then applying Kirrchoff's theorem with only edges of the certain weight in the graph. Therefore, the number of minimum spanning trees in a graph can be solved by multiplying all numbers of arrangements of edges of different weights together.

#### 8.9 Mobius inversion

#### 8.9.1 Mobius inversion formula

$$[x=1] = \sum_{d|x} \mu(d)$$
 
$$x = \sum_{d|x} \mu(d)$$

#### 8.9.2 Gcd inversion

$$\begin{split} \sum_{a=1}^n \sum_{b=1}^n gcd^2(a,b) &= \sum_{d=1}^n d^2 \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [gcd(i,j) = 1] \\ &= \sum_{d=1}^n d^2 \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{t \mid gcd(i,j)} \mu(t) \\ &= \sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} [t|i] \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [t|j] \\ &= \sum_{t=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \lfloor \frac{n}{dt} \rfloor^2 \end{split}$$

The formula can be computed in O(nlogn) complexity. Moreover, let l=dt, then

$$\sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \lfloor \frac{n}{dt} \rfloor^2 = \sum_{l=1}^n \lfloor \frac{n}{l} \rfloor^2 \sum_{d|l} d^2 \mu(\frac{l}{d})$$

Let  $f(l) = \sum_{d|l} d^2 \mu(\frac{l}{d})$ . It can be proven that f(l) is multiplicative. Besides,  $f(p^k) = p^{2k} - p^{2k-2}$ .

Therefore, with linear sieve the formula can be computed in O(n) complexity.

#### 8.10 Numbers

#### 8.10.1 Binomial Coefficients

$$\binom{n}{k} = (-1)^k \binom{k-n-1}{k}, \quad \sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

$$\sum_{k=0}^n \binom{k}{m} = \binom{n+1}{m+1}$$

$$\sqrt{1+z} = 1 + \sum_{k=1}^\infty \frac{(-1)^{k-1}}{k \times 2^{2k-1}} \binom{2k-2}{k-1} z^k$$

$$\sum_{k=0}^r \binom{r-k}{m} \binom{s+k}{n} = \binom{r+s+1}{m+n+1}$$

$$C_{n,m} = \binom{n+m}{m} - \binom{n+m}{m-1}, n \ge m$$

$$\binom{n}{k} \equiv [n\&k = k] \pmod{2}$$

$$\binom{n_1 + \dots + n_p}{m} = \sum_{k_1 + \dots + k_p = m} \binom{n_1}{k_1} \dots \binom{n_p}{k_p}$$

#### 8.10.2Fibonacci Numbers

$$F(z) = \frac{z}{1 - z - z^2}$$

$$f_n = \frac{\phi^n - \hat{\phi}^n}{\sqrt{5}}, \phi = \frac{1 + \sqrt{5}}{2}, \hat{\phi} = \frac{1 - \sqrt{5}}{2}$$

$$\sum_{k=1}^n f_k = f_{n+2} - 1, \quad \sum_{k=1}^n f_k^2 = f_n f_{n+1}$$

$$\sum_{k=0}^n f_k f_{n-k} = \frac{1}{5}(n-1)f_n + \frac{2}{5}nf_{n-1}$$

$$\frac{f_{2n}}{f_n} = f_{n-1} + f_{n+1}$$

$$f_1 + 2f_2 + 3f_3 + \dots + nf_n = nf_{n+2} - f_{n+3} + 2]$$

$$\gcd(f_m, f_n) = f_{\gcd(m, n)}$$

$$f_n^2 + (-1)^n = f_{n+1}f_{n-1}$$

$$f_{n+k} = f_n f_{k+1} + f_{n-1}f_k$$

$$f_{2n+1} = f_n^2 + f_{n+1}^2$$

$$(-1)^k f_{n-k} = f_n f_{k-1} - f_{n-1}f_k$$

$$f_n = f_n f_{k-1} - f_{n-1}f_k$$

$$f_$$

#### 8.10.3 Lucas Numbers

$$L_0 = 2, L_1 = 1, L_n = L_{n-1} + L_{n-2} = \left(\frac{1+\sqrt{5}}{2}\right)^n + \frac{1-\sqrt{5}}{2}\right)^n$$
  
$$L(x) = \frac{2-x}{1-x-x^2}$$

#### 8.10.4 Catlan Numbers

$$c_1 = 1, c_n = \sum_{i=0}^{n-1} c_i c_{n-1-i} = c_{n-1} \frac{4n-2}{n+1} = \frac{\binom{2n}{n}}{n+1}$$
$$= \binom{2n}{n} - \binom{2n}{n-1}, c(x) = \frac{1-\sqrt{1-4x}}{2x}$$

#### 8.10.5 Stirling Cycle Numbers

Divide n elements into k non-empty cycles.

$$\begin{split} s(n,0) &= 0, s(n,n) = 1, s(n+1,k) = s(n,k-1) - ns(n,k) \\ & s(n,k) = (-1)^{n-k} {n \brack k} \\ & {n+1 \brack k} = n {n \brack k} + {n \brack k-1}, {n+1 \brack 2} = n! H_n \\ & x^{\underline{n}} = x(x-1)...(x-n+1) = \sum_k {n \brack k} (-1)^{n-k} x^k \\ & x^{\overline{n}} = x(x+1)...(x+n-1) = \sum_k {n \brack k} x^k \end{split}$$

#### 8.10.6 Stirling Subset Numbers

Divide n elements into k non-empty subsets

$${n+1 \choose k} = k {n \choose k} + {n \choose k-1}$$

$$x^n = \sum_k {n \choose k} x^{\underline{k}} = \sum_k {n \choose k} (-1)^{n-k} x^{\overline{k}}$$

$$m! {n \choose m} = \sum_k {m \choose k} k^n (-1)^{m-k}$$

For a fixed k, generating functions

$$\sum_{n=0}^{\infty} {n \brace k} x^{n-k} = \prod_{r=1}^{k} \frac{1}{1 - rx}$$

#### Motzkin Numbers

Draw non-intersecting chords between n points on a circle.

Pick n numbers  $k_1,k_2,...,k_n \in \{-1,0,1\}$  so that  $\sum_i^a k_i (1 \le a \le n)$  is non-negative and the sum of all numbers is 0.

$$M_{n+1} = M_n + \sum_{i=0}^{n-1} M_i M_{n-1-i} = \frac{(2n+3)M_n + 3nM_{n-1}}{n+3}$$

$$M_n = \sum_{i=0}^{\lfloor \frac{n}{2} \rfloor} {n \choose 2k} Catlan(k)$$

$$M(X) = \frac{1 - x - \sqrt{1 - 2x - 3x^2}}{2x^2}$$

#### 8.10.8 **Eulerian Numbers**

Permutations of the numbers 1 to n in which exactly k elements are

#### 8.10.9 Harmonic Numbers

Sum of the reciprocals of the first n natural numbers.

$$\sum_{k=1}^{n} H_k = (n+1)H_n - n$$

$$\sum_{k=1}^{n} kH_k = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}$$

$$\sum_{k=1}^{n} {n \choose m} H_k = {n+1 \choose m+1} (H_{n+1} - \frac{1}{m+1})$$

#### 8.10.10 Pentagonal Number Theorem

$$\prod_{n=1}^{\infty} (1-x^n) = \sum_{n=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2}$$

$$p(n) = p(n-1) + p(n-2) - p(n-5) - p(n-7) + \cdots$$

$$f(n,k) = p(n) - p(n-k) - p(n-2k) + p(n-5k) + p(n-7k) - \cdots$$

#### 8.10.11 Bell Numbers

Divide a set that has exactly n elements

$$B_n = \sum_{k=1}^n {n \brace k}, \quad B_{n+1} = \sum_{k=0}^n {n \brack k} B_k$$
$$B_{p^m+n} \equiv mB_n + B_{n+1} \pmod{p}$$
$$B(x) = \sum_{n=0}^\infty \frac{B_n}{n!} x^n = e^{e^x - 1}$$

#### 8.10.12 Bernoulli Numbers

$$B_n = 1 - \sum_{k=0}^{n-1} {n \choose k} \frac{B_k}{n-k+1}$$

$$G(x) = \sum_{k=0}^{\infty} \frac{B_k}{k!} x^k = \frac{1}{\sum_{k=0}^{\infty} \frac{x^k}{(k+1)!}}$$

$$\sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k n^{m-k+1}$$

#### 8.10.13 Sum of Powers

$$\begin{split} \sum_{i=1}^n i^2 &= \frac{n(n+1)(2n+1)}{6}, \quad \sum_{i=1}^n i^3 = (\frac{n(n+1)}{2})^2 \\ &\sum_{i=1}^n i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} \\ &\sum_{i=1}^n i^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12} \end{split}$$

#### 8.10.14 Sum of Squares

Denote  $r_k(n)$  the ways to form n with k squares. If:

$$n = 2^{a_0} p_1^{2a_1} \cdots p_r^{2a_r} q_1 b_1 \cdots q_s b_s$$

where 
$$p_i \equiv 3 \mod 4$$
,  $q_i \equiv 1 \mod 4$ , then 
$$r_2(n) = \begin{cases} 0 & \text{if any } a_i \text{ is a half-integer} \\ 4 \prod_{i=1}^{r} (b_i + 1) & \text{if all } a_i \text{ are integers} \end{cases}$$

 $r_3(n) > 0$  when and only when n is not  $4^a(8b+7)$ .

# Appendix 9.1 Calculus table

#### $(\operatorname{arcsec} x)' = \frac{1}{x\sqrt{1-x^2}}$ $\left(\frac{u}{v}\right)' = \frac{u'v - uv'}{v^2}$ $(a^x)' = (\ln a)a^x$ $(\tanh x)' = \operatorname{sech}^2 x$ $(\tan x)' = \sec^2 x$ $(\coth x)' = -\operatorname{csch}^2 x$ $(\cot x)' = \csc^2 x$ $(\operatorname{sech} x)' = -\operatorname{sech} x \tanh x$ $(\sec x)' = \tan x \sec x$ $(\operatorname{csch} x)' = -\operatorname{csch} x \operatorname{coth} x$ $(\csc x)' = -\cot x \csc x$ $(\operatorname{arcsinh} x)' = \frac{1}{\sqrt{1+x^2}}$ $(\arcsin x)' = \frac{1}{\sqrt{1-x^2}}$ $(\operatorname{arccosh} x)' = \frac{1}{\sqrt{x^2 - 1}}$ $(\arccos x)' = -\frac{1}{\sqrt{1-x^2}}$ $(\arctan x)' = \frac{1}{1-x^2}$ $(\operatorname{arctanh} x)' = \frac{1}{1-x^2}$ $(\arctan x)' = \frac{1}{1+x^2}$ $(\operatorname{arccoth} x)' = \frac{1}{x^2 - 1}$ $(\operatorname{arccot} x)' = -\frac{1}{1+x^2}$ $(\operatorname{arccsch} x)' = -\frac{1}{|x|\sqrt{1+x^2}}$ $(\operatorname{arccsc} x)' = -\frac{1}{x\sqrt{1-x}}$

# **9.1.1** $ax + b \ (a \neq 0)$

1.  $\int \frac{x}{ax+b} dx = \frac{1}{a^2} (ax + b - b \ln |ax + b|) + C$ 

2. 
$$\int \frac{x^2}{ax+b} dx = \frac{1}{a^3} \left( \frac{1}{2} (ax+b)^2 - 2b(ax+b) + b^2 \ln|ax+b| \right) + C$$

 $(\operatorname{arcsech} x)' = -\frac{1}{x\sqrt{1-x^2}}$ 

3.  $\int \frac{\mathrm{d}x}{x(ax+b)} = -\frac{1}{b} \ln \left| \frac{ax+b}{x} \right| + C$ 

4.  $\int \frac{\mathrm{d}x}{x^2(ax+b)} = -\frac{1}{bx} + \frac{a}{b^2} \ln \left| \frac{ax+b}{x} \right| + C$ 

5.  $\int \frac{x}{(ax+b)^2} dx = \frac{1}{a^2} \left( \ln|ax+b| + \frac{b}{ax+b} \right) + C$ 

6.  $\int \frac{x^2}{(ax+b)^2} dx = \frac{1}{a^3} \left( ax + b - 2b \ln|ax+b| - \frac{b^2}{ax+b} \right) + C$ 

7.  $\int \frac{\mathrm{d}x}{x(ax+b)^2} = \frac{1}{b(ax+b)} - \frac{1}{b^2} \ln \left| \frac{ax+b}{x} \right| + C$ 

#### **9.1.2** $\sqrt{ax+b}$

1.  $\int \sqrt{ax+b} dx = \frac{2}{3a} \sqrt{(ax+b)^3} + C$ 

7.  $\int \frac{\mathrm{d}x}{x^2 \sqrt{ax+b}} = -\frac{\sqrt{ax+b}}{bx} - \frac{a}{2b} \int \frac{\mathrm{d}x}{x\sqrt{ax+b}}$ 

8.  $\int \frac{\sqrt{ax+b}}{x} dx = 2\sqrt{ax+b} + b \int \frac{dx}{x\sqrt{ax+b}}$ 9.  $\int \frac{\sqrt{ax+b}}{x^2} dx = -\frac{\sqrt{ax+b}}{x} + \frac{a}{2} \int \frac{dx}{\sqrt{ax+b}}$ 

# **9.1.3** $x^2 \pm a^2$

1.  $\int \frac{\mathrm{d}x}{x^2 + a^2} = \frac{1}{a} \arctan \frac{x}{a} + C$ 

2.  $\int \frac{\mathrm{d}x}{(x^2+a^2)^n} = \frac{x}{2(n-1)a^2(x^2+a^2)^{n-1}} + \frac{2n-3}{2(n-1)a^2} \int \frac{\mathrm{d}x}{(x^2+a^2)^{n-1}}$ 

3.  $\int \frac{\mathrm{d}x}{x^2 - a^2} = \frac{1}{2a} \ln \left| \frac{x - a}{x + a} \right| + C$ 

3. 
$$\int \frac{1}{x^2 - a^2} = \frac{1}{2a} \ln \left| \frac{1}{x + a} \right| + C$$
9.1.4  $ax^2 + b$   $(a > 0)$ 

1. 
$$\int \frac{dx}{ax^2 + b} = \begin{cases} \frac{1}{\sqrt{ab}} \arctan \sqrt{\frac{a}{b}}x + C & (b > 0) \\ \frac{1}{2\sqrt{-ab}} \ln \left| \frac{\sqrt{ax} - \sqrt{-b}}{\sqrt{ax} + \sqrt{-b}} \right| + C & (b < 0) \end{cases}$$
2. 
$$\int \frac{x}{ax^2 + b} dx = \frac{1}{2a} \ln \left| ax^2 + b \right| + C$$
3. 
$$\int \frac{x^2}{ax^2 + b} dx = \frac{x}{a} - \frac{b}{a} \int \frac{dx}{ax^2 + b}$$
4. 
$$\int \frac{dx}{x(ax^2 + b)} = \frac{1}{2b} \ln \frac{x^2}{|ax^2 + b|} + C$$
5. 
$$\int \frac{dx}{x^2(ax^2 + b)} = -\frac{1}{bx} - \frac{b}{a} \int \frac{dx}{ax^2 + b}$$

$$\begin{array}{l}
0. \quad \int \frac{x}{x^{3}(ax^{2}+b)} = \frac{2b^{2} \ln \frac{x}{x^{2}} - \frac{2bx^{2}}{2b(ax^{2}+b)} + \frac{1}{2b} \int \frac{dx}{ax^{2}+b} \\
7. \quad \int \frac{dx}{(ax^{2}+b)^{2}} = \frac{x}{2b(ax^{2}+b)} + \frac{1}{2b} \int \frac{dx}{ax^{2}+b} \\
9.1.5 \quad ax^{2} + bx + c \quad (a > 0) \\
1. \quad \frac{dx}{ax^{2}+bx+c} = \begin{cases} \frac{2}{\sqrt{4ac-b^{2}}} \arctan \frac{2ax+b}{\sqrt{4ac-b^{2}}} + C \quad (b^{2} < 4ac) \\ \frac{1}{\sqrt{b^{2}-4ac}} \ln \left| \frac{2ax+b-\sqrt{b^{2}-4ac}}{2ax+b+\sqrt{b^{2}-4ac}} \right| + C \quad (b^{2} > 4ac) \\
2. \quad \int \frac{x}{ax^{2}+bx+c} dx = \frac{1}{2a} \ln |ax^{2} + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^{2}+bx+c} \\
0.1.6 \quad \sqrt{x^{2}+a^{2}} \quad \sqrt{a^{2}+a^{2}} \quad (a > 0)
\end{cases}$$

## **9.1.6** $\sqrt{x^2 + a^2}$ (a > 0)

1. 
$$\int \frac{\mathrm{d}x}{\sqrt{x^2 + a^2}} = \operatorname{arsh} \frac{x}{a} + C_1 = \ln(x + \sqrt{x^2 + a^2}) + C$$

2.  $\int \frac{dx}{\sqrt{(x^2+a^2)^3}} = \frac{x}{a^2\sqrt{x^2+a^2}} + C$ 

3.  $\int \frac{x}{\sqrt{x^2 + a^2}} dx = \sqrt{x^2 + a^2} + C$ 

6.  $\int \frac{x^2}{\sqrt{(x^2 + a^2)^3}} dx = -\frac{x}{\sqrt{x^2 + a^2}} + \ln(x + \sqrt{x^2 + a^2}) + C$ 

7.  $\int \frac{\mathrm{d}x}{x\sqrt{x^2+a^2}} = \frac{1}{a} \ln \frac{\sqrt{x^2+a^2}-a}{|x|} + C$ 

9.  $\int \sqrt{x^2 + a^2} dx = \frac{x}{2} \sqrt{x^2 + a^2} + \frac{a^2}{2} \ln(x + \sqrt{x^2 + a^2}) + C$ 

10.  $\int \sqrt{(x^2 + a^2)^3} dx = \frac{x}{8} (2x^2 + 5a^2) \sqrt{x^2 + a^2} + \frac{3}{8} a^4 \ln(x + \sqrt{x^2 + a^2}) + C$ 

11.  $\int x\sqrt{x^2+a^2}\,\mathrm{d}x = \frac{1}{3}\sqrt{(x^2+a^2)^3} + C$ 

12.  $\int x^2 \sqrt{x^2 + a^2} \, dx = \frac{x}{8} (2x^2 + a^2) \sqrt{x^2 + a^2} - \frac{a^4}{8} \ln(x + \sqrt{x^2 + a^2}) + C$ 

13.  $\int \frac{\sqrt{x^2 + a^2}}{x} dx = \sqrt{x^2 + a^2} + a \ln \frac{\sqrt{x^2 + a^2} - a}{|x|} + C$ 

14.  $\int \frac{\sqrt{x^2 + a^2}}{2} dx = -\frac{\sqrt{x^2 + a^2}}{x} + \ln(x + \sqrt{x^2 + a^2}) + C$ 

#### **9.1.7** $\sqrt{x^2-a^2}$ (a>0)

1.  $\int \frac{\mathrm{d}x}{\sqrt{x^2 - a^2}} = \frac{x}{|x|} \operatorname{arch} \frac{|x|}{a} + C_1 = \ln \left| x + \sqrt{x^2 - a^2} \right| + C$ 

6.  $\int \frac{x^2}{\sqrt{(x^2 - a^2)^3}} dx = -\frac{x}{\sqrt{x^2 - a^2}} + \ln|x + \sqrt{x^2 - a^2}| + C$ 7.  $\int \frac{dx}{x\sqrt{x^2 - a^2}} = \frac{1}{a} \arccos \frac{a}{|x|} + C$ 

 $x^{2}\sqrt{x^{2} - a^{2}}$ 9.  $\int \sqrt{x^{2} - a^{2}} dx = \frac{x}{2}\sqrt{x^{2} - a^{2}} - \frac{a^{2}}{2} \ln|x + \sqrt{x^{2} - a^{2}}| + C$ 10.  $\int \sqrt{(x^{2} - a^{2})^{3}} dx = \frac{x}{8}(2x^{2} - 5a^{2})\sqrt{x^{2} - a^{2}} + \frac{3}{8}a^{4} \ln|x + \sqrt{x^{2} - a^{2}}| + C$ 11.  $\int x\sqrt{x^{2} - a^{2}} dx = \frac{1}{3}\sqrt{(x^{2} - a^{2})^{3}} + C$ 

12.  $\int x^2 \sqrt{x^2 - a^2} dx = \frac{x}{8} (2x^2 - a^2) \sqrt{x^2 - a^2} - \frac{a^4}{8} \ln|x + \sqrt{x^2 - a^2}| + C$ 

13.  $\int \frac{\sqrt{x^2 - a^2}}{x} dx = \sqrt{x^2 - a^2} - a \arccos \frac{a}{|x|} + C$ 

14.  $\int \frac{\sqrt{x^2 - a^2}}{x^2} dx = -\frac{\sqrt{x^2 - a^2}}{x} + \ln|x + \sqrt{x^2 - a^2}| + C$ 

# 9.1.8 $\sqrt{a^2 - x^2}$ (a > 0)1. $\int \frac{dx}{\sqrt{a^2 - x^2}} = \arcsin \frac{x}{a} + C$

2.  $\frac{dx}{\sqrt{(a^2 - x^2)^3}} = \frac{x}{a^2 \sqrt{a^2 - x^2}} + C$ 3.  $\int \frac{x}{\sqrt{a^2 - x^2}} dx = -\sqrt{a^2 - x^2} + C$ 

7.  $\int \frac{\mathrm{d}x}{x\sqrt{a^2-x^2}} = \frac{1}{a} \ln \frac{a-\sqrt{a^2-x^2}}{|x|} + C$ 

9.  $\int \sqrt{a^2 - x^2} dx = \frac{x}{2} \sqrt{a^2 - x^2} + \frac{a^2}{2} \arcsin \frac{x}{a} + C$ 

10.  $\int \sqrt{(a^2 - x^2)^3} dx = \frac{x}{8} (5a^2 - 2x^2) \sqrt{a^2 - x^2} + \frac{3}{8} a^4 \arcsin \frac{x}{a} + C$ 

11.  $\int \sqrt{x^2 - x^2} \, dx = -\frac{1}{3} \sqrt{(a^2 - x^2)^3 + C}$ 12.  $\int x^2 \sqrt{a^2 - x^2} \, dx = \frac{x}{8} (2x^2 - a^2) \sqrt{a^2 - x^2} + \frac{a^4}{8} \arcsin \frac{x}{a} + C$ 13.  $\int \frac{\sqrt{a^2 - x^2}}{x} \, dx = \sqrt{a^2 - x^2} + a \ln \frac{a - \sqrt{a^2 - x^2}}{|x|} + C$ 

# **9.1.9** $\sqrt{\pm ax^2 + bx + c}$ (a > 0)

1.  $\int \frac{\mathrm{d}x}{\sqrt{ax^2 + bx + c}} = \frac{1}{\sqrt{a}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$ 

2.  $\int \sqrt{ax^2 + bx + c} dx = \frac{2ax + b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac - b^2}{8\sqrt{a^3}} \ln|2ax| + b + \frac{4ac - b^2}{8\sqrt{a^3}} \ln|2ax| + b$ 

 $\frac{2\sqrt{a}\sqrt{ax^2+bx+c}|+C}{3. \int \frac{x}{\sqrt{ax^2+bx+c}}dx = \frac{1}{a}\sqrt{ax^2+bx+c} - \frac{b}{2\sqrt{a^3}} \ln|2ax+b+2\sqrt{a}\sqrt{ax^2+bx+c}|+$ 

6.  $\int \frac{b^2 + 4ac}{8\sqrt{a^3}} \arcsin \frac{\frac{2ax - b}{\sqrt{b^2 + 4ac}} + C}{\sqrt{c + bx - ax^2}} dx = -\frac{1}{a}\sqrt{c + bx - ax^2} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{\frac{2ax - b}{\sqrt{b^2 + 4ac}} + C}{\sqrt{b^2 + 4ac}}$ 

# **9.1.10** $\sqrt{\pm \frac{x-a}{x-b}}$ & $\sqrt{(x-a)(x-b)}$

1.  $\int \sqrt{\frac{x-a}{x-b}} dx = (x-b)\sqrt{\frac{x-a}{x-b}} + (b-a)\ln(\sqrt{|x-a|} + \sqrt{|x-b|}) + C$ 

2.  $\int \sqrt{\frac{x-a}{b-x}} dx = (x-b)\sqrt{\frac{x-a}{b-x}} + (b-a)\arcsin\sqrt{\frac{x-a}{b-x}} + C$ 

3.  $\int \frac{\mathrm{d}x}{\sqrt{(x-a)(b-x)}} = 2\arcsin\sqrt{\frac{x-a}{b-x}} + C \ (a < b)$ 

4.  $\int \sqrt{(x-a)(b-x)} dx = \frac{2x-a-b}{4} \sqrt{(x-a)(b-x)} + \frac{(b-a)^2}{4} \arcsin \sqrt{\frac{x-a}{b-x}} + C$ 

# 9.1.11 Triangular function

1.  $\int \tan x dx = -\ln|\cos x| + C$ 2.  $\int \cot x dx = \ln|\sin x| + C$ 

3.  $\int \sec x dx = \ln \left| \tan \left( \frac{\pi}{4} + \frac{x}{2} \right) \right| + C = \ln \left| \sec x + \tan x \right| + C$ 

4.  $\int \csc x \, dx = \ln \left| \tan \frac{x}{2} \right| + C = \ln \left| \csc x - \cot x \right| + C$ 

 $\int \sec^2 x \, \mathrm{d}x = \tan x + C$ 

5.  $\int \sec^{2} x dx = \tan x + C$ 6.  $\int \csc^{2} x dx = -\cot x + C$ 7.  $\int \sec x \tan x dx = \sec x + C$ 8.  $\int \csc x \cot x dx = -\csc x + C$ 9.  $\int \sin^{2} x dx = \frac{x}{2} - \frac{1}{4} \sin 2x + C$ 10.  $\int \cos^{2} x dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$ 

$$\begin{array}{ll} 11. & \int \sin^n x \, \mathrm{d} x = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x \, \mathrm{d} x \\ 12. & \int \cos^n x \, \mathrm{d} x = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x \, \mathrm{d} x \\ 13. & \frac{\mathrm{d} x}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^n - 1} x + \frac{n-2}{n-1} \int \frac{\mathrm{d} x}{\sin^n - 2} x \\ 14. & \frac{\mathrm{d} x}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{\mathrm{d} x}{\cos^n - 2} x \\ 15. & \int \cos^m x \sin^n x \, \mathrm{d} x \\ & = \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^m x \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+n} \int \cos^m x \sin^{n-2} x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+n} \int \cos^m x \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^m x \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^m x \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^m x \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^{m+1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^m x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^m x \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^m x \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^m x \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{m+n} \cos^m x \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x \, \mathrm{d} x \\ & = -\frac{1}{2(a+b)} \sin^n x$$

## 9.1.12 Inverse triangular function (a > 0)

$$\begin{array}{l} 1. \quad \int \arcsin \frac{x}{a} \mathrm{d}x = x \arcsin \frac{x}{a} + \sqrt{a^2 - x^2} + C \\ 2. \quad \int x \arcsin \frac{x}{a} \mathrm{d}x = (\frac{x^2}{2} - \frac{a^2}{4}) \arcsin \frac{x}{a} + \frac{x}{4} \sqrt{x^2 - x^2} + C \\ 3. \quad \int x^2 \arcsin \frac{x}{a} \mathrm{d}x = \frac{x^3}{3} \arcsin \frac{x}{a} + \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 - x^2} + C \\ 4. \quad \int \arccos \frac{x}{a} \mathrm{d}x = x \arccos \frac{x}{a} - \sqrt{a^2 - x^2} + C \\ 5. \quad \int x \arccos \frac{x}{a} \mathrm{d}x = (\frac{x^2}{2} - \frac{a^2}{4}) \arccos \frac{x}{a} - \frac{x}{4} \sqrt{a^2 - x^2} + C \\ 6. \quad \int x^2 \arccos \frac{x}{a} \mathrm{d}x = \frac{x}{3} \arccos \frac{x}{a} - \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 - x^2} + C \\ 7. \quad \int \arctan \frac{x}{a} \mathrm{d}x = x \arctan \frac{x}{a} - \frac{a}{2} \ln(a^2 + x^2) + C \\ 8. \quad \int x \arctan \frac{x}{a} \mathrm{d}x = \frac{1}{2} (a^2 + x^2) \arctan \frac{x}{a} - \frac{a}{2} x + C \\ 9. \quad \int x^2 \arctan \frac{x}{a} \mathrm{d}x = \frac{x^3}{3} \arctan \frac{x}{a} - \frac{a}{6} x^2 + \frac{a^3}{6} \ln(a^2 + x^2) + C \end{array}$$

## 9.1.13 Exponential function

```
1. \int a^x dx = \frac{1}{\ln a} a^x + C
2. \int e^{ax} dx = \frac{1}{a} a^{ax} + C
\begin{aligned} 2. & \int e^{ax} \, dx = \frac{1}{a} a^{ax} + C \\ 3. & \int x e^{ax} \, dx = \frac{1}{a^2} (ax - 1) a^{ax} + C \\ 4. & \int x^n e^{ax} \, dx = \frac{1}{a} x^n e^{ax} - \frac{n}{a} \int x^{n-1} e^{ax} \, dx \\ 5. & \int x a^x \, dx = \frac{x}{\ln a} a^x - \frac{1}{(\ln a)^2} a^x + C \\ 6. & \int x^n a^x \, dx = \frac{1}{\ln a} x^n a^x - \frac{n}{\ln a} \int x^{n-1} a^x \, dx \\ 7. & \int e^{ax} \sin bx \, dx = \frac{1}{a^2 + b^2} e^{ax} (a \sin bx - b \cos bx) + C \\ 6. & \int x^n \cos bx \, dx = \frac{1}{a^2 + b^2} e^{ax} (b \sin bx + a \cos bx) + C \end{aligned}
  \frac{n(n-1)b^2}{a^2+b^2n^2} \int e^{ax} \sin^{n-2} bx dx
10. \int e^{ax} \cos^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) +
             \frac{n(n-1)b^2}{a^2+b^2n^2} \int e^{ax} \cos^{n-2} bx dx
```

#### 9.1.14 Logarithmic function

	$\int \ln x  \mathrm{d}x = x \ln x - x + C$
	$\int \frac{\mathrm{d}x}{x \ln x} = \ln \left  \ln x \right  + C$
3.	$\int x^n \ln x dx = \frac{1}{n+1} x^{n+1} (\ln x - \frac{1}{n+1}) + C$
	$\int (\ln x)^n dx = x(\ln x)^n - n \int (\ln x)^{n-1} dx$
5.	$\int x^m (\ln x)^n dx = \frac{1}{m+1} x^{m+1} (\ln x)^n - \frac{n}{m+1} \int x^m (\ln x)^{n-1} dx$

#### 9.2Regular expression

#### 9.2.1Special pattern characters

Characters	Description
	Not newline
\t	Tab (HT)
\n	Newline (LF)
\v	Vertical tab (VT)
\f	Form feed (FF)
\r	Carriage return (CR)
\cletter	Control code
\xhh	ASCII character
\uhhhh	Unicode character
\0	Null
\int	Backreference
\d	Digit
\D	Not digit
\s	Whitespace
\S	Not whitespace
\ W	Word (letters, numbers and the underscore)
\W	Not word
\character	Character
[class]	Character class
[^class]	Negated character class

#### 9.2.2Quantifiers

Characters	Times
*	0 or more
+	1 or more
?	0 or 1
{int}	int
{int,}	int or more
{min,max}	Between min and max

By default, all these quantifiers are greedy (i.e., they take as many characters that meet the condition as possible). This behavior can be overridden to ungreedy (i.e., take as few characters that meet the condition as possible) by adding a question mark (?) after the quantifier.

#### 9.2.3Groups

Characters	Description
(subpattern)	Group with backreference
(?:subpattern)	Group without backreference

#### 9.2.4Assertions

Characters	Description
^	Beginning of line
\$	End of line
\b	Word boundary
\B	Not a word boundary
(?=subpattern)	Positive lookahead
(?!subpattern)	Negative lookahead

#### 9.2.5Alternative

A regular expression can contain multiple alternative patterns simply by separating them with the separator operator (|): The regular expression will match if any of the alternatives match, and as soon as 9.2.6 Character classes

Class	Description
[:alnum:]	Alpha-numerical character
[:alpha:]	Alphabetic character
[:blank:]	Blank character
[:cntrl:]	Control character
[:digit:]	Decimal digit character
[:graph:]	Character with graphical representation
[:lower:]	Lowercase letter
[:print:]	Printable character
[:punct:]	Punctuation mark character
[:space:]	Whitespace character
[:upper:]	Uppercase letter
[:xdigit:]	Hexadecimal digit character
[:d:]	Decimal digit character
[:w:]	Word character
[:s:]	Whitespace character

Please note that the brackets in the class names are additional to those opening and closing the class definition. For example:

[[:alpha:]] is a character class that matches any alphabetic character. [abc[:digit:]] is a character class that matches a, b, c, or a

[^[:space:]] is a character class that matches any character ex-