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## **Environment**

#### 1.1 Vimrc

```
set ru nu ts=4 sts=4 sw=4 si sm hls is ar bs=2 mouse=a syntax on nm <F3> :vsplit %<.in <CR>
nm <F4> :!gedit % <CR>
s au BufEnter *.cpp set cin
au BufEnter *.cpp nm <F5> :!time ./%< <CR>|nm <F7> :!
    gdb ./%< <CR>|nm <F8> :!time ./%< < %<.in <CR>|nm <F9> :!g++ % -o %< -g -std=gnu++14 -O2 -DLOCAL && size %< <CR>
```

## Data Structure

## 2.1 KD tree

```
/* kd_tree : finds the k-th closest point in O(kn^{1-\frac{1}{k}}). Usage : Stores the data in p[]. Call function init (n, k). Call min_kth (d, k). (or max_kth) (k is 1-based)
 2 Usage
3 Note
                     Switch to the commented code for Manhattan
               distance.
21
23 //
28
               | ]));
ret += 111 * tmp * tmp; }
  ret += std::max (std::abs (rhs.data[i] - dmax.
data[i]) + std::abs (rhs.data[i] - dmin.data[i]));
       return ret; } tree[MAXN * 4];
struct result {
long long dist; point d; result() {}
result (const long long &dist, const point &d) :
    dist (dist), d (d) {}
bool operator > (const result &rhs) const { return
32
34
     35
38
          if ("tree[rt].r) tree[rt].merge (tree[tree[rt].r], k
      if ("tree[rt].r) tree[rt].merge (tree[tree[tree]...,
); }
std::priority_queue<result, std::vector<result>, std
::less <result>> heap_1;
std::priority_queue<result, std::vector<result>, std
::greater <result>> heap_r;
void _min_kth (const int &depth, const int &rt, const
    int &m, const point &d) {
  result tmp = result (sqrdist (tree[rt].p, d), tree[
    rt].p);
if ((int)heap_l.size() < m) heap_l.push (tmp);
else if (tmp < heap_l.top()) {
  heap_l.pop();
  heap_l.push (tmp); }</pre>
55
```

```
62
74
75
80
```

#### Splay 2.2

```
void push_down (int x) {
   if ( n[x].c[0]) push (n[x].c[0], n[x].t);
   if ( n[x].c[1]) push (n[x].c[1], n[x].t);
   if ( n[x].t = tag (); }
   void update (int x) {
      n[x].m = gen (x);
   if ( n[x].c[0]) n[x].m = merge (n[n[x].c[0]].m, n[x].
   if ("n[x].c[1]) n[x].m = merge (n[x].m, n[n[x].c[1]].
  11
12
```

#### 2.3Link-cut tree

```
void access (int x)
int u = x, v = -1;
while (u != -1) {
  = u;

n[u].c[1] = v;

if (~v) n[v].f = u, n[v].p = -1;

update (u); u = n[v = u].p; }

splay (x); }
```

## Formula

#### Zellers congruence 3.1

```
/* Zeller's congruence : converts between a calendar
```

## 3.2 Lattice points below segment

```
1 /* Euclidean-like algorithm : computes the sum of
     \sum_{i=0}^{n-1} \left[ \frac{a+bi}{m} \right] . \star /
```

```
if (b >= m) return (n - 1) * n / 2 * (b / m) + solve
    (n, a, b % m, m);
return solve ((a + b * n) / m, (a + b * n) % m, m, b)
```

#### 3.3 Adaptive Simpson's method

```
|\mathbf{r}| / \star Adaptive Simpson's method : integrates f in [1, r].
  struct'simpson {
  double area (double (*f) (double), double 1, double r
   double m = 1 + (r - 1) / 2;
return (f (1) + 4 * f (m) + f (r)) * (r - 1) / 6; }
double solve (double (*f) (double), double 1, double
    r, double eps, double a) {
    double m = 1 + (r - 1) / 2;
    double left = area (f, 1, m), right = area (f, m, r)
```

#### 3.4Neural network

```
1 /* Neural network : machine learning. */
2 template <int ft = 6, int n = 6, int MAXDATA = 100000>
3 struct network {
    network
     for (int i = 0; i < ft + 1; ++i) avg[i] = sig[i] =
    0; }
double compute (double *x) {
  for (int i = 0; i < n; ++i) {
    val[i] = 0; for (int j = 0; j < ft; ++j) val[i] +=
        wp[i][j] * x[j];
    val[i] = 1 / (1 + exp (-val[i])); }
double res = 0; for (int i = 0; i < n; ++i) res +=
    val[i] * w[i];
/ return 1 / (1 + exp (-res));
return res; }
roid desc (double *x, double t, double eta) {</pre>
    double
   void desc (double *x, double t, double eta) {
  double o = compute (x), del = (o - t); // * o * (1 -
     o)
for (int i = 0; i < n; ++i) for (int j = 0; j < ft;
21
       22
    26
                  ++j)
     33
    return compute (x) * sig[ft] + avg[ft]; }
std::string to_string () {
  std::ostringstream os; os.precision (20); os << std</pre>
     42
     for (int i = 0; i < ft + 1; ++i) os << sig[i] << "_"
   return os.str (); }
void read (const std::string &str) {
    std::istringstream is (str);
    for (int i = 0; i < n; ++i) for (int j = 0; j < ft;
        ++j) is >> wp[i][j];
    for (int i = 0; i < n; ++i) is >> w[i];
    for (int i = 0; i < ft + 1; ++i) is >> avg[i];
    for (int i = 0; i < ft + 1; ++i) is >> sig[i]; } };
```

## Number theory 4

#### 4.1Fast power module

```
1 /* Fast power module : x^n */ 2 int fpm (int x, int n, int mod) {
```

```
int ans = 1, mul = x; while (n) {
  if (n & 1) ans = int (111 * ans * mul % mod);
  mul = int (111 * mul * mul % mod); n >>= 1; }
  return ans; }
     long long mod (long long x, long long y, long long
mod) {
mod) {

long long t = (x * y - (long long) ((long double) x /
mod * y + 1E-3) * mod) % mod;

return t < 0 ? t + mod : t; }

long long long long n, long long
mod) {
       long long ans = 1, mul = x; while (n) {
  if (n & 1) ans = mul_mod (ans, mul, mod);
  mul = mul_mod (mul, mul, mod); n >>= 1; ]
  return ans; }
```

## 4.2 Euclidean algorithm

```
\frac{1}{2} /* Euclidean algorithm : solves for ax + by = gcd (a,
  b). */
void euclid (const long long &a, const long long &b,
   long long &x, long long &y) {
   if (b == 0) x = 1, y = 0;
   else euclid (b, a % b, y, x), y -= a / b * x; }
  long long inverse (long long x, long long m) {
  long long a, b; euclid (x, m, a, b); return (a % m +
     m) % m; }
```

## 4.3 Discrete Fourier transform

```
1 /* Discrete Fourier transform : the nafarious you-know
     -what thing.

e : call init for the suggested array size, and solve for the transform. (use f!=0 for the inverse
template <int MAXN = 1000000>
11
18
```

## Fast Walsh-Hadamard transform

```
/* Fast Walsh-Hadamard transform : binary operation
        transform. */
 / ( a[j + k] = (x + y) / 2, a[i + j + k] = (x
- y) / 2, and : a[j + k] = x - y, or : a[i +
| + k] = y - x; */
       }else{
/* xor : a[j + k] = x + y, a[i + j + k] = x - y,
    and : a[j + k] = x + y, or : a[i + j + k] = x
    + y; */
} } }
```

## Number theoretic transform

```
if (f == 1) {
int rev = fpm (n, mod - 2, mod);
```

```
for (int i = 0; i < n; ++i) a[i] = int (111 * a[i]
    * rev % mod); }
int crt (int *a, int mod) {
  static int inv[3][3];
for (int i = 0; i < 3; ++i) for (int j = 0; j < 3;
    ++i)</pre>
24
                      inv[i][j] = (int) inverse (MOD[i], MOD[j]);
static int x[3];
for (int i = 0; i < 3; ++i) { x[i] = a[i];
    for (int j = 0; j < i; ++j) {
        int t = (x[i] - x[j] + MOD[i]) % MOD[i];
        if (t < 0) t += MOD[i];
        x[i] = int (1LL * t * inv[j][i] % MOD[i]); }
int sum = 1, ret = x[0] % mod;
for (int i = 1; i < 3; ++i) {
    sum = int (1LL * sum * MOD[i - 1] % mod);
    ret += int (1LL * x[i] * sum % mod);
    if (ret >= mod) ret -= mod; }
return ret; } };
                        return ret; } };
```

## 4.6 Chinese remainder theorem

```
/* Chinese remainder theroem : finds positive integers
    x = out.first + k * out.second that satisfies x %
    in[i].second = in[i].first. */
struct crt {
    false;
n *= (in[i].first - out.first) / divisor;
n = fix (n, in[i].second);
out.first += out.second * n;
out.second *= in[i].second / divisor;
out.first = fix (out.first, out.second); }
return true; } };
```

#### 4.7Linear Recurrence

```
1 /* Linear recurrence : finds the n-th element of a
 linear recurrence.
2 Usage : vector <int> - first n terms, vector <int> - transition function, calc (k) : the kth term mod
              MOD.
    Example : In : {2, 1}, {2, 1} : a_1 = 2, a_2 = 1, a_n = 2a_{n-1} + a_{n-2}, Out : calc (3) = 5, calc (10007) = 959155122 (MOD 1E9+7) */
 result.erase(result.begin() + n + 1, result.end());
      return result; }
linear_rec (const std::vector <int> &first, const std
::vector <int> &trans) : first(first), trans(
trans) {
        ::vector <int> &trans) : first(first), trans(
    trans) {
    n = first.size(); std::vector <int> a(n + 1, 0); a
        [1] = 1; bin.push_back(a);
    for (int i = 1; i < LOG; ++i) bin.push_back(add(bin
        [i - 1], bin[i - 1])); }
int solve (int k) {
    std::vector <int> a(n + 1, 0); a[0] = 1;
    for (int i = 0; i < LOG; ++i) if (k >> i & 1) a =
        add(a, bin[i]);
    int ret = 0;
    for (int i = 0; i < n; ++i) if ((ret += (long long)
        a[i + 1] * first[i] % MOD) >= MOD) ret -= MOD;
    return ret; } };
19
```

#### Berlekamp Massey algorithm 4.8

```
/* Berlekamp Massey algorithm : Complexity: O(n^2)
Requirement: const MOD, inverse(int)
Input: the first elements of the sequence
Output: the recursive equation of the given sequence
Example In: {1, 1, 2, 3}
Example Out: {1, 1000000006, 1000000006} (MOD = 1e9+7)
| struct berlekamp-massey {
| struct Poly { std::vector <int> a; Poly() { a.clear()
         Poly (std::vector <int> &a) : a (a) {}
int length () const { return a.size(); }
Poly move (int d) { std::vector <int> na (d, 0);
na.insert (na.end (), a.begin (), a.end ());
return Poly (na); }
int calc(std::vector <int> &d, int pos) { int ret =
0;
600
              for (int i = 0; i < (int) a.size (); ++i) {
```

```
if ((ret += 1LL * d[pos - i] * a[i] % MOD) >= MOD)
  ret -= MOD; } }
return ret; }
comprator - (c
```

## 4.9 Baby step giant step algorithm

```
ı /* Baby step giant step algorithm : Solves a^x = b \mod c
     /* Baby step giant step algorithm : Solves a^x = b \mod a in O(\sqrt{c}). */ struct bsgs {
 int solve (int a, int b, int c) {
   std::unordered_map <int, int> bs;
   int m = (int) sqrt ((double) c) + 1, res = 1;
   for (int i = 0; i < m; ++i) {
    if (bs.find (res) == bs.end ()) bs[res] = i;
    res = int (1LL * res * a * c); }
   int mul = 1, inv = (int) inverse (a, c);
   for (int i = 0; i < m; ++i) mul = int (1LL * mul * inv * c);
   res = b * c;
   for (int i = 0; i < m; ++i) {
      if (bs.find (res) != bs.end ()) return i * m + bs[
      res];
              res = iii (1LL * res * mul % c); } return -1; } };
```

## 4.10 Miller Rabin primality test

```
_{1} /* Miller Rabin : tests whether a certain integer is
```

## 4.11 Pollard's Rho algorithm

```
19
 if (rem > 1) ans.push_back (rem); }
return ans; } ;;
```

## 5 Geometry

```
#define cd const double &
const double EPS = 1E-8, PI = acos (-1);
int sgn (cd x) { return x < -EPS ? -1 : x > EPS; }
int cmp (cd x, cd y) { return sgn (x - y); }
double sqr (cd x) { return x * x; }
```

## 5.1 Point

```
#define cp const point &

struct point {

double x, y;

explicit point (cd x = 0, cd y = 0) : x (x), y (y) {}

int dim () const { return sgn (y) == 0 ? sgn (x) < 0

: sgn (y) < 0; }

point unit () const { double 1 = sqrt (x * x + y * y)

//counter-clockwise

point rot90 () const { return point (-y, x); }

//clockwise

point rot90 () const { return point (y, -x); }

point rot (cd t) const {

double c = cos (t), s = sin (t);

return point (x * c - y * s, x * s + y * c); };

bool operator == (cp a, cp b) { return cmp (a.x, b.x)

== 0 && cmp (a.y, b.y) == 0; }

bool operator < (cp a, cp b) { return (cmp (a.x, b.x)

== 0) ? cmp (a.y, b.y) < 0 : cmp (a.x, b.x) < 0; }

point operator + (cp a, cp b) { return point (-a.x, -a.y); }

point operator + (cp a, cp b) { return point (a.x + b. x, a.y + b.y); }

point operator * (cp a, cd b) { return point (a.x - b. x, a.y - b.y); }

point operator / (cp a, cd b) { return point (a.x + b, a.y * b); }

point operator / (cp a, cd b) { return point (a.x + b, a.y * b); }

couble dot (cp a, cp b) { return a.x * b.x + a.y * b.y ; }

double dot (cp a, cp b) { return a.x * b.x + a.y * b.y ; }

double dis2 (cp a, cp b = point ()) { return sqrt (dis2 (a, b)); }

double dis (cp a, cp b = point ()) { return sqrt (dis2 (a, b)); }
```

## 5.2 Line

## 5.3 Circle

```
#define cc const circle &
struct circle {
   point c; double r;
   explicit circle (point c = point (), double r = 0) :
        c (c), r (r) {} };
   bool operator == (cc a, cc b) { return a.c == b.c &&
        cmp (a.r, b.r) == 0; }
```

```
if (cmp (point_to_line (b.c, a), b.r) > 0) return std
    ::vector <point> ();
double x = sqrt (sqr (b.r) - sqr (point_to_line (b.c,
    a)));
return std::vector <point> ({a.c + r * x - r.rot90 () * h, a.c + r * x + r.rot90 () * h}); }

//Counter-clockwise with respect of point a.

make_circle (a, b.c); return circle_intersect (p, b); }

//Counter-clockwise with respect of point (p, b); }
      33//Counter-clockwise with respect of point O_a.
34|std::vector <line> extangent (cc a, cc b) {
35| std::vector <line> ret;
                      if (pp.size () == 2 && qq.size () == 2) {
   if (cmp (a.r, b.r) < 0) std::swap (pp[0], pp[1]),
      std::swap (qq[0], qq[1]);
   ret.push_back (line (pp[0], qq[0]));
   ret.push_back (line (pp[1], qq[1])); }
return ret; }
//Counter-clockwise with recount of the content 
                return ret; }
//Counter-clockwise with respect of point Oa.
std::vector std::vector std::vector + a.c + b.r) / (a.r + b.r);
std::vector pp = tangent (p, a), qq = tangent (p, b);
if (pp.size () == 2 && qq.size () == 2) {
    ret.push_back (line (pp[0], qq[0]));
    ret.push_back (line (pp[1], qq[1])); }
return ret; }
```

## 5.4 Centers of a triangle

## 5.5 Fermat point

#### 5.6Convex hull

```
//Counter-clockwise, with minimum number of points.
bool turn_left (cp a, cp b, cp c) { return sgn (det (b - a, c - a)) >= 0; }
std::vector <point> convex_hull (std::vector <point> a
```

## 5.7 Half plane intersection

```
\mathbf{1}| /* Online half plane intersection : complexity O(n)
   each operation. */
std::vector <point> cut (const std::vector<point> &c,
    line p) {
    std::vector <point> ret;
}
              inne p) {
std::vector <point> ret;
if (c.empty ()) return ret;
for (int i = 0; i < (int) c.size (); ++i) {
  int j = (i + 1) % (int) c.size ();
  if (turn_left (p.s, p.t, c[i])) ret.push_back (c[i])</pre>
if (two_side (c[i], c[j], p)) ret.push_back (
    line_intersect (p, line (c[i], c[j]))); }
return ret; }
/* Offline half plane intersection : complexity
10 /* Offilme half product in the control of the co
std::vector < 
<h)</h>
                                                                        <point> half_plane_intersect (std::vector
              < 0:
             else return cmp (a.first, b.first) < 0; });
h.resize (std::unique (g.begin (), g.end (), [] (
    const polar &a, const polar &b) { return cmp (a.
    first, b.first) == 0 }) - g.begin ());
for (int i = 0; i < (int) h.size (); ++i) h[i] = g[i
    ].second;
int forcon con = -1; std::vegtor < line ret (b.</pre>
               int fore = 0, rear = -1; std::vector <line> ret (h.
    size (), line ());
for (int i = 0; i < (int) h.size (); ++i) {
   while (fore < rear && !turn_left (h[i],
        line_intersect (ret[rear - 1], ret[rear]))) --
   rear;
   while (fore < rear ff !turn_left (h[i])</pre>
                     25
               inte_intersect (ret[rore], ret[rore + 1])) ++
fore;
if (rear - fore < 2) return std::vector <point> ();
std::vector <point> ans; ans.resize (rear + 1);
for (int i = 0; i < rear + 1; ++i) ans[i] =
    line_intersect (ret[i], ret[(i + 1) % (rear + 1)
    1).</pre>
              return ans; }
```

## 5.8 Nearest pair of points

```
1 /* Nearest pair of points : [1, r), need to sort p
    first. */
1 /* Nearest pair of points : [1, r], need to sort p
    first. */
2 double solve (std::vector <point> &p, int 1, int r) {
3    if (1 + 1 >= r) return INF;
4    int m = (1 + r) / 2; double mx = p[m].x; std::vector <point> v;
5    double ret = std::min (solve(p, 1, m), solve(p, m, r)).
         for (int i = 1; i < r; ++i)
  if (sqr (p[i].x - mx) < ret) v.push_back (p[i]);
  sort (v.begin (), v.end (), [&] (cp a, cp b) { return
      a.y < b.y; });
  for (int i = 0; i < v.size (); ++i)
  for (int j = i + 1; j < v.size (); ++j) {
    if (sqr (v[i].y - v[j].y) > ret) break;
    ret = min (ret, dis2 (v[i] - v[j])); }
  return ret; }
```

#### 5.9 Minimum circle

```
circle minimum_circle (std::vector <point> p) {
  circle ret; std::random_shuffle (p.begin (), p.end ()
  for (int i = 0; i < (int) p.size (); ++i) if (!
   in_circle (p[i], ret)) {</pre>
```

## 5.10 Intersection of a polygon and a circle

```
struct polygon_circle_intersect {
  double sector_area (cp a, cp b, const double &r) {
    double c = (2.0 * r * r - dis2 (a, b)) / (2.0 * r *
        return r * r * acos (c) / 2.0; }
double area (cp a, cp b, const double &r) {
    double dA = dot (a, a), dB = dot (b, b), dC =
        point_to_segment (point (), line (a, b));
    if (sgn (dA - r * r) <= 0 && sgn (dB - r * r) <= 0)
        return det (a, b) / 2.0;
    point tA = a.unit () * r, tB = b.unit () * r;
    if (sgn (dC - r) > 0) return sector_area (tA, tB, r)
       10
```

```
5.11
         Union of circles
void addevent(cc a, cc b, std::vector <event> &evt,
    i+cnt;
for (int j = 0; j < C; ++j) if (j != i && !same (c[
    i], c[j]) && overlap (c[j], c[i])) ++cnt;
for (int j = 0; j < C; ++j) if (j != i && !overlap
    (c[j], c[i]) && !overlap (c[i], c[j]) &&
    intersect (c[i], c[j]))
addevent (c[i], c[j], evt, cnt);
if (evt.empty ()) area[cnt] += PI * c[i].r * c[i].r</pre>
    area[cnt] += ang * c[i].r * c[i].r / 2 - sin(ang) * c[i].r * c[i].r / 2; } } };
```

## 5.12 3D point

15

23

```
#define cp3 const point3 &
struct point3 {
double x, y, z;
explicit point3 (cd x = 0, cd y = 0, cd z = 0) : x (x
), y (y), z (z) {};
point3 operator + (cp3 a, cp3 b) { return point3 (a.x
+ b.x, a.y + b.y, a.z + b.z);
point3 operator - (cp3 a, cp3 b) { return point3 (a.x
- b.x, a.y - b.y, a.z - b.z);
point3 operator * (cp3 a, cd b) { return point3 (a.x *
b, a.y * b, a.z * b); }
point3 operator / (cp3 a, cd b) { return point3 (a.x /
b, a.y / b, a.z / b); }
```

```
| double dot (cp3 a, cp3 b) { return a.x * b.x + a.y * b | .y + a.z * b.z; } | 10 point3 det (cp3 a, cp3 b) { return point3 (a.y * b.z - a.z * b.y, -a.x * b.z + a.z * b.x, a.x * b.y - a. y * b.x); } | 11 double dis2 (cp3 a, cp3 b = point3 ()) { return sqr (a .x - b.x) + sqr (a.y - b.y) + sqr (a.z - b.z); } | 12 double dis (cp3 a, cp3 b = point3 ()) { return sqr ( a dis2 (a, b)); } | 13 //right-handed, if x+ -> y+ is right-handed | 14 point3 rotate(cp3 p, cp3 axis, double w) { | 15 double x = axis.x, y = axis.y, z = axis.z; | 16 double x = axis.x, y = axis.y, z = axis.z; | 17 double a[4][4]; memset(a, 0, sizeof (a)); | 18 a[3][3] = 1; | 19 a[0][0] = ((y * y + z * z) * cosw + x * x) / s; | 19 a[0][0] = ((y * y + z * z) * cosw + x * x) / s; | 19 a[0][1] = x * x * (1 - cosw) / s - y * sinw / ss; | 10 a[1][1] = (x * x + z * z) * cosw + y * y) / s; | 18 a[1][1] = ((x * x + z * z) * cosw + y * y) / s; | 19 a[1][1] = ((x * x + z * z) * cosw + y * y) / s; | 19 a[2][0] = x * z * (1 - cosw) / s - x * sinw / ss; | 19 a[2][0] = x * z * (1 - cosw) / s - x * sinw / ss; | 19 a[2][0] = (x * x * y * (1 - cosw) / s - x * sinw / ss; | 19 a[2][0] = (x * x * y * (1 - cosw) / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x * sinw / ss; | 10 cosw / s - x
```

## 5.13 3D line

## 5.14 3D convex hull

# 6 Graph

```
template <int MAXN = 100000, int MAXM = 100000>
struct edge_list {
int size, begin[MAXN], dest[MAXM], next[MAXM];
```

## 6.1 Hopcoft-Karp algorithm

```
/* Hopcoft-Karp algorithm : unweighted maximum matching for bipartition graphs with complexity O(m\sqrt{n}). */

template <int MaxN = 100000, int MaxM = 100000>
struct hopcoft_karp {

using edge_list = std::vector <int> [MaxN];

bool dfs (edge_list <MaxN, MaxM> &e, int x) {

for (int i = e.begin[x]; ~i; i = e.next[i]) {

int y = e.dest[i], w = my[y];

if (!~w || (lv[x] + 1 == lv[w] && dfs (e, w))) {

mx[x] = y; my[y] = x; return true; } }

lv[x] = -1; return false; }

int solve (edge_list <MaxN, MaxM> &e, int n, int m) {

std::fill (mx, mx + n, -1); std::fill (my, my + m, -1); for (int ans = 0; ;) {

std::vector <int> q;

for (int i = 0; i < n; ++i)

if (mx[i] == -1) {

lv[i] = 0; q.push_back (i);

} else lv[i] = -1;

for (int head = 0; head < (int) q.size(); ++head) {

int x = q[head];

for (int i = e.begin[x]; ~i; i = e.next[i]) {

int y = e.dest[i], w = my[y];

if (~w && lv[w] < 0) { lv[w] = lv[x] + 1; q.

push_back (w); } }

int d = 0; for (int i = 0; i < n; ++i) if (!~mx[i] && dfs (e, i)) ++d;

if (d == 0) return ans; else ans += d; } };
```

## 6.2 Kuhn-Munkres algorithm

```
/* Kuhn Munkres algorithm: weighted maximum ming algorithm for bipartition graphs with complexity O(N^3).

Note: The graph is 1-based. */

template <int MAXN = 500>

template <int MAXN = 500>

template <int MAXN];

struct kuhn_munkres {
   int n, w[MAXN] [MAXN], lx[MAXN], ly[MAXN], m[MAXN],
        way[MAXN], sl[MAXN];

bool u[MAXN];

void hungary(int x) {
   m[0] = x; int j0 = 0;
   std::fill (sl, sl + n + 1, INF); std::fill (u, u + n + 1, false);

do {
   u[j0] = true; int i0 = m[j0], d = INF, j1 = 0;
   for (int j = 1; j <= n; ++j)
   if (u[j] == false) {
      int cur = -w[i0][j] - lx[i0] - ly[j];
      if (sl[j] < d) { d = sl[j]; j1 = j; } }

   for (int j = 0; j <= n; ++j) {
      if (u[j]) { lx[m[j]] += d; ly[j] -= d; }
      if (u[j]) { lx[m[j]] += d; ly[j] -= d; }
      if (u[j]) { lx[m[j]] += d; ly[j] -= d; }
      int j1 = way[j0]; m[j0] = m[j1]; j0 = j1;
   } while (j0); }

int solve() {
   for (int i = 1; i <= n; ++i) m[i] = lx[i] = ly[i] = way[i] = 0;
   for (int i = 1; i <= n; ++i) hungary (i);
   int sum = 0; for (int i = 1; i <= n; ++i) sum += w[m [i]][i];
   return sum; } ;
```

## 6.3 Blossom algorithm

## 6.4 Weighted blossom algorithm

```
1 /* Weighted blossom algorithm (vfleaking ver.) : maximum matching for general weighted graphs with
                                                                                                                                    80
 complexity O(n^3). 2 Usage : Set n to the size of the vertices. Run init () . Set g[][].w to the weight of the edge. Run solve
 The first result is the answer, the second one is the number of matching pairs. Obtain the matching with match[].
 match[].
4 Note: 1-based. */
5 struct weighted_blossom {
6 static const int INF = INT_MAX, MAXN = 400;
7 struct edge{ int u, v, w; edge (int u = 0, int v = 0, int w = 0): u(u), v(v), w(w) {} };
                                                                                                                                    92
      int n, n_x;
edge g[MAXN * 2 + 1][MAXN * 2 + 1];
int lab[MAXN * 2 + 1], match[MAXN * 2 + 1], slack[
MAXN * 2 + 1], st[MAXN * 2 + 1], pa[MAXN * 2 +
     1];
int flower_from[MAXN * 2 + 1][MAXN + 1], S[MAXN * 2 +
1], vis[MAXN * 2 + 1];
std::vector <int> flower[MAXN * 2 + 1]; std::queue <
    99
13
                                                                                                                                  100
                                                                                                                                  101
19
21
                                                                                                                                  113
23
29
      void augment (int u, int v) {
  for (; ; ) {
    int xnv = st[match[u]]; set_match (u, v);
    if (!xnv) return; set_match (xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv; } }
```

```
int get_lca (int u, int v) {
  static int t = 0;
  for (++t; u || v; std::swap (u, v)) {
   if (u == 0) continue; if (vis[u] == t) return u;
   vis[u] = t; u = st[match[u]]; if (u) u = st[pa[u]];
  }
  return 0; }

37
                  return 0; }
void add_blossom (int u, int lca, int v) {
  int b = n + 1; while (b <= n_x && st[b]) ++b;
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca]; flower[b].clear ();
  flower[b].push_back (lca);
  for (int x = u, y; x != lca; x = st[pa[y]]) {
    flower[b].push_back (x), flower[b].push_back (y =
        st[match[x]]), q_push (y); }
  std::reverse (flower[b].begin () + 1, flower[b].end
    ());
  for (int x = v, y; x != lca; x = st[pa[y]]) {
                           int xs = flower[b][i]; S[xs] = -1, set_slack(xs); }
st[b] = 0; }
bool on_found_edge (const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1; int nu = st[match[v]];
    slack[v] = slack[nu] = 0; S[nu] = 0, q_push(nu);
} else if(S[v] == 0) {
  int lca = get_lca(u, v);
  if (!lca) return augment(u, v), augment(v, u), true
                for (int x = 1; x \lefta n_x, v.a, z \lefta n_x, \lefta n_x, z \lefta n_x, \lefta n_x
                   return false; }
std::pair <long long, int> solve () {
  memset (match + 1, 0, sizeof (int) * n); n_x = n;
  int n_matches = 0; long long tot_weight = 0;
  for (int u = 0; u <= n; ++u) st[u] = u, flower[u].
     clear();
  int w_max = 0;</pre>
                            for (int u = 1; u <= n; ++u) for (int v = 1; v <= n; ++v) {
                          flower_from[u][v] = (u == v ? u : 0); w_max = std::
    max (w_max, g[u][v].w); }
for (int u = 1; u <= n; ++u) lab[u] = w_max;
while (matching ()) ++n_matches;
for (int u = 1; u <= n; ++u) if (match[u] && match[u] < u) tot_weight += g[u][match[u]].w;
```

22

```
return std::make_pair (tot_weight, n_matches); }
void init () { for (int u = 1; u <= n; ++u) for (int
    v = 1; v <= n; ++v) g[u][v] = edge (u, v, 0); }</pre>
```

#### 6.5Maximum flow

```
/* Sparse graph maximum flow : isap.*/
2 template <int MAXN = 1000, int MAXM = 100000>
 13
dflow; e.flow[cur[p] 1] += dflow; }
else {
  int mindist = n + 1;
  for (int i = e.begin[u]; ~i; i = e.next[i])
  if (e.flow[i] && mindist > d[e.dest[i]]) {
    mindist = d[e.dest[i]]; cur[u] = i; }
  if (!--gap[d[u]]) return maxflow;
  gap[d[u] = mindist + 1]++; u = pre[u]; } }
  /* Dense graph maximum flow : dinic. */
  template <int MAXN = 1000, int MAXM = 100000>
  struct dinic {
    struct flow edge list {
     38
          int ans = 0;
                                        n = n_; s = s_; dinic::t = t_;
          int ans = 0; n = n_; s = s_; dinic::t = t_;
while (bfs (e)) {
  for (int i = 0; i < n; ++i) w[i] = e.begin[i];
  ans += dfs (e, s, INF); }
return ans; } };</pre>
```

## 6.6 Minimum cost flow

```
/* Sparse graph minimum cost flow : EK. */
2 template <int MAXN = 1000, int MAXM = 100000>
3 struct minimum_cost_flow {
4  struct cost_flow_edge_list {
5  int size, begin[MAXN], dest[MAXM], next[MAXM], cost[MAXM], flow[MAXM];
6  void clear (int n) { size = 0; std::fill (begin, begin + n, -1); }
7  cost_flow_edge_list (int n = MAXN) { clear (n); }
8  void add_edge (int u, int v, int c, int f) {
9  dest[size] = v; next[size] = begin[u]; cost[size] = c; flow[size] = f; begin[u] = size++;
10  dest[size] = u; next[size] = begin[v]; cost[size] = -c; flow[size] = 0; begin[v] = size++; };
11  int n, s, t, prev[MAXN], dist[MAXN], occur[MAXN];
12  bool augment (cost_flow_edge_list &e) {
13   std::vector <int> queue;
14   std::fill (dist, dist + n, INF); std::fill (occur, occur + n, 0);
15   dist[sl= 0; cover[sl= true; queue push back (sl= cover[sl= true;
```

```
for (int head = 0; head < (int)queue.size(); ++head)</pre>
                             int x = queue[head];
                                int x = queue[head];
for (int i = e.begin[x]; ~i; i = e.next[i]) {
  int y = e.dest[i];
  if (e.flow[i] && dist[y] > dist[x] + e.cost[i]) {
    dist[y] = dist[x] + e.cost[i]; prev[y] = i;
    if (!occur[y]) {
        occur[y] = true; queue.push_back (y); } }
cccur[x] = false; }
                              for
   20
                 if (!occur[y]) {
   occur[y] = true; queue.push_back (y); } }
occur[x] = false; }
return dist[t] < INF; }
std::pair <int, int> solve (cost_flow_edge_list &e,
   int n_, int s_, int t_) {
   n = n_; s = s_; t = t_; std::pair <int, int> ans =
        std::make_pair (0, 0);
while (augment (e)) {
   int num = INF;
   24
                            int num = INF;
for (int i = t; i != s; i = e.dest[prev[i] ^ 1]) {
   num = std::min (num, e.flow[prev[i]]); }
                            ans.first += num;
for (int i = t; i != s; i = e.dest[prev[i] ^ 1
e.flow[prev[i]] -= num; e.flow[prev[i] ^ 1] +
in the struct of the stru
                                  ans.second += num * e.cost[prev[i]]; } }
                       slack[i] = INF; }
if (delta == INF) return 1;
for (int i = 0; i < n; i++) if (visit[i]) dis[i] +=
    delta;
return 0; }</pre>
                 return 0; }
int dfs (cost_flow_edge_list &e, int x, int flow) {
  if (x == t) { tf += flow; tc += flow * (dis[s] - dis
      [t]); return flow; }
  visit[x] = 1; int left = flow;
  for (int i = e.begin[x]; ~i; i = e.next[i])
    if (e.flow[i] > 0 && !visit[e.dest[i]]) {
    int v = e dest[i]:
                                  int
                                      int y = e.dest[i];
int y = e.dest[i];
if (dis[y] + e.cost[i] == dis[x]) {
  int delta = dfs (e, y, std::min (left, e.flow[i])
    );
    ;
}
                                       e.flow[i] -= delta; e.flow[i ^ 1] += delta; left
                                       -= delta;
if_(!left) { visit[x] = false; return flow; }
                 do { do {
   std::fill (visit + 1, visit + t + 1, 0);
} while (dfs (e, s, INF)); } while (!modlable ());
return std::make_pair (tf, tc);
};
```

## 6.7 Stoer Wagner algorithm

```
int solve () {
   int mincut, i, j, s, t, ans;
   for (mincut = INF, i = 1; i < n; i++) {
    ans = contract (s, t); bin[t] = true;
    if (mincut > ans) mincut = ans;
    if (mincut == 0) return 0;
    for (j = 1; j <= n; j++) if (!bin[j])
    edge[s][j] = (edge[j][s] += edge[j][t]); }
   return mincut; } };</pre>
```

## 6.8 DN maximum clique

```
typedef std::vector <Vertex> Vertices; Vertices V;
typedef std::vector <int> ColorClass; ColorClass QMAX,
Q;
std::vector <ColorClass> C;
static bool desc_degree (const Vertex &vi,const Vertex &vi) { return vi.d > vj.d; }
yoid init_colors (Vertices &v) {
const int max_degree = v[0].d;
if for (int i = 0; i < (int) v.size(); ++i) v[i].d = std
::min (i, max_degree) + 1; }
void set_degrees (Vertices &v) {
if for (int i = 0, j; i < (int) v.size (); ++i)
if for (v[i].d = j = 0; j < (int) v.size (); ++i)
if for (v[i].d = j = 0; j < (int) v.size (); ++j)
if v[i].d += e[v[i].i][v[j].i]; }
if std::vector <StepCount { int i1, i2; StepCount(): i1 (0), i2
(0) { };
std::vector <StepCount> S;
if bool cut1 (const int pi, const ColorClass &A) {
if (e[pi][A[i]]) return true; return false; }
if (e[A.back().i][A]i].i]) B.push_back(A[i].i); }
if (e[A.back().i][A]i].i]) B.push_back(A[i].i); }
if (e[A.back().i][A]i].i]) B.push_back(A[i].i); }
if (e[A.back().i][A]i].i]) B.push_back(A[i].i); }
if (int i = 0; i < (int) A.size () + 1, 1);
column for (int i = 0; i < (int) R.size (); ++i) {
int pi = R[i].i, k = 1; while (cut1(pi, C[k])) ++k;
if (k > maxno) maxno = k, C[maxno + 1].clear();
column for (int i = 0; i < (int) R.size (); ++i)
if (j > 0) R[j - 1].d = 0;
for (int k = min_k; k <= maxno; ++k)
if (j > 0) R[j - 1].d = 0;
for (int i = 0; i < (int) C[k].size (); ++i)
if (int) R.size () if + slevel - 1].i1 - slevel
li2;
stlevel].i2 = S[level - 1].i1;
while ((int) R.size ()) {
location for int i = 0; i < (int) C[k].size (); ++i)</pre>
                            39
                                            Q.push_back (R.back ().i); Vertices Rp; cut2 (R, Rp
  Q.push_back (R.back ().i); Vertices Rp; cut2 (R, Rp
);
if ((int) Rp.size ()) {
    if ((float) S[level].il / ++pk < Tlimit)
        degree_sort (Rp);

    color_sort (Rp); ++S[level].il, ++level;
    expand_dyn (Rp); --level;
} else if ((int) Q.size () > (int) QMAX.size ())
    QMAX = Q;

Q.pop_back (); } else return; R.pop_back (); } }

Q.pop_back (); } else return; R.pop_back (); } {
    void mcqdyn (int *maxclique, int &sz) {
    set_degrees (V); std::sort(V.begin (), V.end (),
        desc_degree); init_colors (V);
    for (int i = 0; i < (int) V.size () + 1; ++i) S[i].il
        = S[i].i2 = 0;
    expand_dyn (V); sz = (int) QMAX.size ();
    for (int i = 0; i < (int) QMAX.size (); i++)
        maxclique[i] = QMAX[i]; }

zvoid degree_sort (Vertices & R) {
    set_degrees(R); std::sort(R.begin(), R.end(),
        desc_degree);
    Maxclique (const BB *conn, const int sz, const float
        tt = .025) : pk (0), level (1), Tlimit (tt) {
        for (int i = 0; i < sz; i++) V.push_back (Vertex (i));
    set = conn, C.resize (sz + 1), S.resize (sz + 1); };
    set BB e[N]; int ans, sol[N]; for (...) e[x][y] = e[y][x]</pre>
     BB e[N]; int ans, sol[N]; for (...) e[x][y] = e[y][x] = true;
     Maxclique mc (e, n); mc.mcqdyn (sol, ans); //0-based.
for (int i = 0; i < ans; ++i) std::cout << sol[i] << std::endl;
```

## 6.9 Dominator tree

## 6.10 Tarjan

```
| /* Tarjan : strongly-connected components. */
| template <int MAXN = 1000000>
| struct tarjan {
| int comp[MAXN], size;
| int dfn[MAXN], ind, low[MAXN], ins[MAXN], stk[MAXN],
| stks;
| void dfs (const edge_list <MAXN, MAXM> &e, int i) {
| dfn[i] = low[i] = ind++;
| ins[i] = 1; stk[stks++] = i;
| for (int x = e.begin[i]; "x; x = e.next[x]) {
| int j = e.dest[x]; if (!"dfn[j]) {
| dfs (j);
| if (low[i] > low[j]) low[i] = low[j];
| if (low[j] > dfn[i]); //vertex-biconnected
| if (low[j] > dfn[i]); //vertex-biconnected
| if (low[j] > dfn[i]); //edge-biconnected
| } else if (ins[j] && low[i] > dfn[j])
| low[i] = dfn[j]; }
| if (dfn[i] == low[i]) { //strongly-connected
| for (int j = -1; j!= i;
| j = stk[--stks], ins[j] = 0, comp[j] = size);
| ++size; }
| void solve (const edge_list <MAXN, MAXM> &e, int n) {
| size = ind = stks = 0;
| std::fill (dfn, dfn + n, -1);
| for (int i = 0; i < n; ++i) if (!"dfn[i])
| dfs (e, i); } ;</pre>
```

## 7 String 7.1 Manacher

## 7.2 Suffix Array

## Suffix Automaton

```
/* Suffix automaton : head - the first state. tail -
the last state. Terminating states can be reached
via visiting the ancestors of tail. state::len -
the longest length of the string in the state.
state::right - 1 - the first location in the
string where the state can be reached. state::
parent - the parent link. state::dest - the
automaton link. */
template <int MAXN = 1000000, int MAXC = 26>
struct suffix automaton f
   3 struct suffix_automaton {
4 struct state {
               int len, right; state *parent, *dest[MAXC];
state (int len = 0, int right = 0) : len (len),
    right (right), parent (NULL) {
    memset (dest, 0, sizeof (dest)); }
} node_pool[MAXN * 2], *tot_node, *null = new state()
               state *head, *tail;
void extend (int token) {
  state *p = tail;
  state *np = tail -> dest[token] ? null : new (
     tot_node++) state (tail -> len + 1, tail -> len
     + 1);
12
                    while (p && !p -> dest[token]) p -> dest[token] = np
if (!p) np -> parent;
else {
    state *q = p -> dest[token];
    if (p -> len + 1 == q -> len) {
        np -> parent = q;
    } else {
        state *nq = new (tot_node++) state (*q);
        np -> parent = q;
    } else {
        state *nq = new (tot_node++) = q;
        np -> parent = q -> parent = nq;
        while (p && p -> dest[token] == q) {
            p -> dest[token] == q) {
            p -> dest[token] = nq, p = p -> parent;
        }
    }
tail = np == null ? np -> parent : np; }
roid init () {
              tail = np == null ? np -> parent : np; }
void init () {
  tot_node = node_pool;
  head = tail = new (tot_node++) state (); }
suffix_automaton () { init (); } };
```

## 7.4 Palindromic tree

```
/* Palindromic tree : extend () - returns whether the tree has generated a new node. odd, even - the root of two trees. last - the node representing the last char. node::len - the palindromic string length of the node. */
2 template <int MAXN = 1000000, int MAXC = 26>
12
      now = now -> fail);
return now; }
bool extend (int token) {
  text[++size] = token; node *now = match (last);
  if (now -> child[token])
  return last = now -> child[token], false;
  last = now -> child[token] = new (tot_node++) node (
    now -> len + 2);
  if (now == odd) last -> fail = even;
else /
18
```

## 7.5 Regular expression

```
std::string str = ("The_the_there");
std::regex pattern ("(th|Th)[\\w]*",
                           std::
std::regex_match (str, match, pattern);
match = *i
match.length () gives length and match.position ()
    gives position of the match. */ }
```

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### 8 Tips 8.1 Java

```
/* Java reference : References on Java IO, structures, etc. */
import java.io.*;
import java.lang.*;
import java.math.*;
import java.util.*;
/* Common usage:
      | /* Common usage:
| Scanner in = new Scanner (System.in);
| Scanner in = new Scanner (new BufferedInputStream (
| 7 | Scanner in = new Scanner (System.in); |
| 8 | Scanner in = new Scanner (new BufferedInputStream (
| System.in); |
| 9 | in.nextInt () / in.nextBigInteger () / in. |
| nextBigDecimal () / in.nextDouble () |
| 10 | in.nextLine () / in.hasNext () |
| 11 | System.out.print (...); |
| 12 | System.out.printf (...); |
| 13 | System.out.printf (...); |
| 14 | BigInteger : BigInteger.valueOf (int) / abs / negate () / max / min / add / subtract / multiply / divide / remainder (BigInteger) / gcd (BigInteger) / modInverse (BigInteger mod) / modPow ( |
| BigInteger ex, BigInteger mod) / pow (int ex) / |
| not () / and / or / xor (BigInteger) / shiftLeft / |
| shiftRight (int) / compareTo (BigInteger) / |
| intValue () / longValue () / toString (int radix) / isProbablePrime (int certainty) / |
| nextProbablePrime (int certainty) / |
| nextProbablePrime () |
| BigDecimal : consists of a BigInteger value and a | scale. The scale is the number of digits to the | right of the decimal point. |
| divide (BigDecimal) : exact divide. |
| divide (BigDecimal) : exact divide. |
| divide (BigDecimal) : exact divide. |
| HALF_EVEN / HALF_UP / UNNECESSARY / UP. |
| BigDecimal setScale (int newScale, RoundingMode | roundingMode) : returns a BigDecimal with newScale | |
| 19 | doubleValue () / toPlainString () : converts to other | |
    doubleValue () / toPlainString () : converts to other
                                     types.

ys: Arrays.sort (T [] a); Arrays.sort (T [] a,
int fromIndex, int toIndex); Arrays.sort (T [] a,
int fromIndex, int toIndex, Comperator <? super T>
   20 Arrays
  int fromindex, int toindex, Comperator <? super T>
    comperator);
21 LinkedList <E> : addFirst / addLast (E) / getFirst /
    getLast / removeFirst / removeLast () / clear () /
    add (int, E) / remove (int) / size () / contains
    / removeFirstOccurrence / removeLastOccurrence (E)
22 ListIterator <E> listIterator (int index) : returns an
 iterator :
    E next / previous () : accesses and iterates.
    hasNext / hasPrevious () : checks availablity.
    nextIndex / previousIndex () : returns the index of a
        subsequent call.
    add / set (E) / remove () : changes element.
    PriorityQueue <E> (int initcap, Comparator <? super E>
        comparator) : add (E) / clear () / iterator () /
        peek () / poll () / size ()

28 TreeMap <K, V> (Comparator <? super K> comparator) :
    Map.Entry <K, V> ceilingEntry / floorEntry /
        higherEntry / lowerEntry (K): getKey / getValue ()
        / setValue (V) : entries.

29 clear () / put (K, V) / get (K) / remove (K) / size ()
30 StringBuilder : StringBuilder (string) / append (int,
                                         iterator
   StringBuilder: StringBuilder (string) / append (int, string, ...) / insert (int offset, ...) charAt (int) / setCharAt (int, char) / delete (int, int) / reverse () / length () / toString ()

String: String: String.format (String, ...) / toLowerCase / toUpperCase () */
  31 String: String.format (Str
toUpperCase () */
32 /* Examples on Comparator :
33 public class Main {
                public class Main {
  public static class Point {
    public int x; public int y;
    public Point () {
      x = 0;
      y = 0; }
  public Point (int xx, int yy) {
      x = xx;
      y = yy; } };
  public static class Cmp implements Comparator <Point>
                  > (c);
return; } };
              */
/* or
                 /* OF :
public static class Point implements Comparable <
    Point> {
    public int x; public int y;
    public Point () {
                       public Point () {
  x = 0;
  y = 0; }
public Point (int xx, int yy) {
  x = xx;
  y = yy; }
public int compareTo (Point p) {
  if (x < p.x) return -1;
  if (x == p.x) {</pre>
```

```
if (y < p.y) return -1;
if (y == p.y) return 0; }
return 1; }
public boolean equalTo (Point p) {
return (x == p.x && y == p.y); }
public int hashCode () {
  return x + y; } };</pre>
67
68
69
70
71
72
  */
//Faster IO :
```

#### 8.2 Random numbers

```
std::mt19937_64 mt (time (0));
std::uniform_int_distribution <int> uid (1, 100);
std::uniform_real_distribution <double> urd (1, 100);
std::cout << uid (mt) << "_" << urd (mt) << "\n";
```

#### Read hack 8.3

```
__always_inline__, __artificial__))
_int next_int () {
const int SIZE = 110000; static char buf[SIZE + 1];
```

#### Stack hack 8.4

```
//C++
""
#pragma comment (linker, "/STACK:36777216")
//G++
int __size__ = 256 << 20;</pre>
           _size__ = 256 << 20;
*_p_ = (char*) malloc(__size__) + __size__;
__ ("movl_%0,_%%esp\n" :: "r"(_p__));
```

#### 8.5 Time hack

```
clock_t t = clock ();
std::cout << 1. * t / CLOCKS_PER_SEC << "\n";</pre>
```

## Builtin functions

- \_builtin\_clz: Returns the number of leading 0-bits in x, starting at the most significant bit position. If x is 0, the result is
- undefined.
  \_builtin\_ctz: Returns the number of trailing 0-bits in x, starting at the least significant bit position. If x is 0, the result is
- undefined.
  \_builtin\_clrsb: Returns the number of leading redundant sign bits in x, i.e. the number of bits following the most significant bit that are identical to it. There are no special cases for 0 or
- other values.
  \_builtin\_popcount: Returns the number of 1-bits in x.
  \_builtin\_parity: Returns the parity of x, i.e. the number of 1-bits in x modulo 2.
  \_builtin\_bswap16, \_builtin\_bswap32, \_builtin\_bswap64: Returns x with the order of the bytes (8 bits as a group) reversed.
- bitset::Find\_first(), bitset::Find\_next(idx): bitset built in functions.

  Prufer sequence

In combinatorial mathematics, the Prufer sequence of a labeled tree is a unique sequence associated with the tree. The sequence for a tree on n vertices has length n-2.

One can generate a labeled tree's Prufer sequence by iteratively removing vertices from the tree until only two vertices remain. Specifically, consider a labeled tree T with vertices 1, 2, ..., n. At step i, remove the leaf with the smallest label and set the ith element of the Prufer sequence to be the label of this leaf's neighbour.

One can generate a labeled tree from a sequence in three steps. The

tree will have n+2 nodes, numbered from 1 to n+2. For each node set its degree to the number of times it appears in the sequence plus 1. Next, for each number in the sequence a[i], find the first (lowestnumbered) node, j, with degree equal to 1, add the edge (j, a[i]) to the tree, and decrement the degrees of j and a[i]. At the end of this loop two nodes with degree 1 will remain (call them u, v). Lastly, add the edge (u, v) to the tree.

The Prufer sequence of a labeled tree on n vertices is a unique sequence of length n-2 on the labels 1 to n - this much is clear. Somewhat less obvious is the fact that for a given sequence S of length n-2 on the labels 1 to n, there is a unique labeled tree whose Prufer sequence

 ${}^{\text{is }S.}_{f 8.8}$ Spanning tree counting

Kirchhoff's Theorem: the number of spanning trees in a graph G is equal to any cofactor of the Laplacian matrix of G, which is equal to the difference between the graph's degree matrix (a diagonal matrix with vertex degrees on the diagonals) and its adjacency matrix (a (0,1)matrix with 1's at places corresponding to entries where the vertices are adjacent and 0's otherwise).

The number of edges with a certain weight in a minimum spanning tree is fixed given a graph. Moreover, the number of its arrangements can be obtained by finding a minimum spanning tree, compressing connected components of other edges in that tree into a point, and then applying Kirrchoff's theorem with only edges of the certain weight in the graph. Therefore, the number of minimum spanning trees in a graph can be solved by multiplying all numbers of arrangements of edges of different weights together.

#### 8.9 Mobius inversion

#### 8.9.1 Mobius inversion formula

$$[x = 1] = \sum_{d|x} \mu(d)$$
$$x = \sum_{d|x} \mu(d)$$

#### 8.9.2Gcd inversion

$$\begin{split} \sum_{a=1}^{n} \sum_{b=1}^{n} gcd^{2}(a,b) &= \sum_{d=1}^{n} d^{2} \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [gcd(i,j) = 1] \\ &= \sum_{d=1}^{n} d^{2} \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{t \mid gcd(i,j)} \mu(t) \\ &= \sum_{d=1}^{n} d^{2} \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} [t \mid i] \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [t \mid j] \\ &= \sum_{d=1}^{n} d^{2} \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \lfloor \frac{n}{dt} \rfloor^{2} \end{split}$$

The formula can be computed in O(nlogn) complexity. Moreover, let l = dt, then

$$\sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \big\lfloor \frac{n}{dt} \big\rfloor^2 = \sum_{l=1}^n \big\lfloor \frac{n}{l} \big\rfloor^2 \sum_{d|l} d^2 \mu(\frac{l}{d})$$

Let  $f(l) = \sum_{d|l} d^2 \mu(\frac{l}{d})$ . It can be proven that f(l) is multiplicative. Besides,  $f(p^k) = p^{2k} - p^{2k-2}$ .

Therefore, with linear sieve the formula can be computed in O(n)complexity.

#### 8.10 **2-SAT**

In terms of the implication graph, two literals belong to the same strongly connected component whenever there exist chains of implications from one literal to the other and vice versa. Therefore, the two literals must have the same value in any satisfying assignment to the given 2-satisfiability instance. In particular, if a variable and its negation both belong to the same strongly connected component, the instance cannot be satisfied, because it is impossible to assign both of these literals the same value. As Aspvall et al. showed, this is a necessary and sufficient condition: a 2-CNF formula is satisfiable if and only if there is no variable that belongs to the same strongly connected component as its

This immediately leads to a linear time algorithm for testing satisfiability of 2-CNF formulae: simply perform a strong connectivity analysis on the implication graph and check that each variable and its negation belong to different components. However, as Aspvall et al. also showed, it also leads to a linear time algorithm for finding a satisfying assignment, when one exists. Their algorithm performs the following steps:

Construct the implication graph of the instance, and find its strongly connected components using any of the known linear-time algorithms for strong connectivity analysis.

Check whether any strongly connected component contains both a variable and its negation. If so, report that the instance is not satisfiable

and halt.
Construct the condensation of the implication graph, a smaller graph that has one vertex for each strongly connected component, and an edge from component i to component j whenever the implication graph contains an edge uv such that u belongs to component i and v belongs to component j. The condensation is automatically a directed acyclic graph and, like the implication graph from which it was formed, it is skew-symmetric.

Topologically order the vertices of the condensation. In practice this may be efficiently achieved as a side effect of the previous step, as

components are generated by Kosaraju's algorithm in topological order and by Tarjan's algorithm in reverse topological order.

For each component in the reverse topological order, if its variables do not already have truth assignments, set all the literals in the component to be true. This also causes all of the literals in the complementary component to be set to false.

## 8.11 Interesting numbers

## 8.11.1 Binomial Coefficients

$$\binom{n}{k} = (-1)^k \binom{k-n-1}{k}, \quad \sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

$$\sum_{k=0}^n \binom{k}{m} = \binom{n+1}{m+1}$$

$$\sqrt{1+z} = 1 + \sum_{k=1}^\infty \frac{(-1)^{k-1}}{k \times 2^{2k-1}} \binom{2k-2}{k-1} z^k$$

$$\sum_{k=0}^r \binom{r-k}{m} \binom{s+k}{n} = \binom{r+s+1}{m+n+1}$$

$$C_{n,m} = \binom{n+m}{m} - \binom{n+m}{m-1}, n \ge m$$

$$\binom{n}{k} \equiv [n\&k = k] \pmod{2}$$

$$\binom{n_1 + \dots + n_p}{m} = \sum_{k_1 + \dots + k_n = m} \binom{n_1}{k_1} \dots \binom{n_p}{k_p}$$

## 8.11.2 Fibonacci Numbers

$$F(z) = \frac{z}{1 - z - z^2}$$

$$f_n = \frac{\phi^n - \hat{\phi}^n}{\sqrt{5}}, \phi = \frac{1 + \sqrt{5}}{2}, \hat{\phi} = \frac{1 - \sqrt{5}}{2}$$

$$\sum_{k=1}^n f_k = f_{n+2} - 1, \quad \sum_{k=1}^n f_k^2 = f_n f_{n+1}$$

$$\sum_{k=0}^n f_k f_{n-k} = \frac{1}{5}(n-1)f_n + \frac{2}{5}nf_{n-1}$$

$$\frac{f_{2n}}{f_n} = f_{n-1} + f_{n+1}$$

$$f_1 + 2f_2 + 3f_3 + \dots + nf_n = nf_{n+2} - f_{n+3} + 2]$$

$$\gcd(f_m, f_n) = f_{\gcd(m, n)}$$

$$f_n^2 + (-1)^n = f_{n+1}f_{n-1}$$

$$f_{n+k} = f_n f_{k+1} + f_{n-1}f_k$$

$$f_{2n+1} = f_n^2 + f_{n+1}^2$$

$$(-1)^k f_{n-k} = f_n f_{k-1} - f_{n-1}f_k$$

$$\begin{cases} f_r, & m \bmod 4 = 0; \\ (-1)^{r+1}f_{n-r}, & m \bmod 4 = 1; \\ (-1)^n f_r, & m \bmod 4 = 2; \\ (-1)^{r+1+n}f_{n-r}, & m \bmod 4 = 3. \end{cases}$$

## 8.11.3 Lucas Numbers

$$L_0 = 2, L_1 = 1, L_n = L_{n-1} + L_{n-2} = \left(\frac{1+\sqrt{5}}{2}\right)^n + \frac{1-\sqrt{5}}{2}\right)^n$$
$$L(x) = \frac{2-x}{1-x-x^2}$$

## 8.11.4 Catlan Numbers

$$c_1 = 1, c_n = \sum_{i=0}^{n-1} c_i c_{n-1-i} = c_{n-1} \frac{4n-2}{n+1} = \frac{\binom{2n}{n}}{n+1}$$
$$= \binom{2n}{n} - \binom{2n}{n-1}, c(x) = \frac{1-\sqrt{1-4x}}{2x}$$

## 8.11.5 Stirling Cycle Numbers

Divide n elements into k non-empty cycles.

$$\begin{split} s(n,0) &= 0, s(n,n) = 1, s(n+1,k) = s(n,k-1) - ns(n,k) \\ & s(n,k) = (-1)^{n-k} {n \brack k} \\ & {n+1 \brack k} = n {n \brack k} + {n \brack k-1}, {n+1 \brack 2} = n! H_n \\ & x^{\underline{n}} = x(x-1)...(x-n+1) = \sum_k {n \brack k} (-1)^{n-k} x^k \\ & x^{\overline{n}} = x(x+1)...(x+n-1) = \sum_k {n \brack k} x^k \end{split}$$

## 8.11.6 Stirling Subset Numbers

Divide n elements into k non-empty subsets.

$${n+1 \choose k} = k {n \choose k} + {n \choose k-1}$$

$$x^n = \sum_k {n \choose k} x^{\underline{k}} = \sum_k {n \choose k} (-1)^{n-k} x^{\overline{k}}$$

$$m! {n \choose m} = \sum_k {m \choose k} k^n (-1)^{m-k}$$

For a fixed k, generating functions

$$\sum_{n=0}^{\infty} {n \brace k} x^{n-k} = \prod_{r=1}^{k} \frac{1}{1 - rx}$$

## 8.11.7 Motzkin Numbers

Draw non-intersecting chords between n points on a circle.

Pick n numbers  $k_1, k_2, ..., k_n \in \{-1, 0, 1\}$  so that  $\sum_i^a k_i (1 \le a \le n)$  is non-negative and the sum of all numbers is 0.

$$M_{n+1} = M_n + \sum_{i=0}^{n-1} M_i M_{n-1-i} = \frac{(2n+3)M_n + 3nM_{n-1}}{n+3}$$

$$M_n = \sum_{i=0}^{\lfloor \frac{n}{2} \rfloor} {n \choose 2k} Catlan(k)$$

$$M(X) = \frac{1 - x - \sqrt{1 - 2x - 3x^2}}{2x^2}$$

## 8.11.8 Eulerian Numbers

Permutations of the numbers 1 to n in which exactly k elements are greater than the previous element.

## 8.11.9 Harmonic Numbers

Sum of the reciprocals of the first n natural numbers.

$$\sum_{k=1}^{n} H_k = (n+1)H_n - n$$

$$\sum_{k=1}^{n} kH_k = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}$$

$$\sum_{k=1}^{n} {n \choose m} H_k = {n+1 \choose m+1} (H_{n+1} - \frac{1}{m+1})$$

## 8.11.10 Pentagonal Number Theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = \sum_{n=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2}$$

$$= p(n-1) + p(n-2) - p(n-5) - p(n-7) + \cdots$$

 $p(n) = p(n-1) + p(n-2) - p(n-5) - p(n-7) + \cdots$  $f(n,k) = p(n) - p(n-k) - p(n-2k) + p(n-5k) + p(n-7k) - \cdots$ 

## 8.11.11 Bell Numbers

Divide a set that has exactly n elements.

$$B_{n} = \sum_{k=1}^{n} {n \brace k}, \quad B_{n+1} = \sum_{k=0}^{n} {n \brack k} B_{k}$$

$$B_{p^{m}+n} \equiv mB_{n} + B_{n+1} \pmod{p}$$

$$B(x) = \sum_{n=0}^{\infty} \frac{B_{n}}{n!} x^{n} = e^{e^{x} - 1}$$

## 8.11.12 Bernoulli Numbers

$$B_n = 1 - \sum_{k=0}^{n-1} {n \choose k} \frac{B_k}{n-k+1}$$

$$G(x) = \sum_{k=0}^{\infty} \frac{B_k}{k!} x^k = \frac{1}{\sum_{k=0}^{\infty} \frac{x^k}{(k+1)!}}$$

$$\sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k n^{m-k+1}$$

## **8.11.13** Sum of Powers

$$\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}, \quad \sum_{i=1}^{n} i^3 = (\frac{n(n+1)}{2})^2$$
$$\sum_{i=1}^{n} i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

$$\sum_{i=1}^{n} i^{5} = \frac{n^{2}(n+1)^{2}(2n^{2}+2n-1)}{12}$$

## 8.11.14 Sum of Squares

Denote  $r_k(n)$  the ways to form n with k squares. If:

$$n = 2^{a_0} p_1^{2a_1} \cdots p_r^{2a_r} q_1 b_1 \cdots q_s b_s$$

where  $p_i \equiv 3 \mod 4$ ,  $q_i \equiv 1 \mod 4$ , then

$$r_2(n) = \begin{cases} 0 & \text{if any } a_i \text{ is a half-integer} \\ 4 \prod_{i=1}^{r} (b_i + 1) & \text{if all } a_i \text{ are integers} \end{cases}$$

 $r_3(n) > 0$  when and only when n is not  $4^a(8b+7)$ .

## 8.11.15 Derangement

$$D_1 = 0, D_2 = 1, D_n = n! \left(\frac{1}{0!} - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + \frac{(-1)^n}{n!}\right)$$
$$D_n = (n-1)(D_{n-1} + D_{n-2})$$

## 8.11.16 Tetrahedron Volume

If U, V, W, u, v, w are lengths of edges of the tetrahedron (first three form a triangle; u opposite to U and so on)

$$V = \frac{\sqrt{4u^2v^2w^2 - \sum_{cyc} u^2(v^2 + w^2 - U^2)^2 + \prod_{cyc} (v^2 + w^2 - U^2)}}{12}$$

## Appendix Calculus table

$\left(\frac{u}{v}\right)' = \frac{u'v - uv'}{v^2}$	(arcson x)' = 1
$(a^x)' = (\ln a)a^x$ $(\tan x)' = \sec^2 x$	$(\operatorname{arcsec} x)' = \frac{1}{x\sqrt{1-x^2}}$ $(\tanh x)' = \operatorname{sech}^2 x$ $(\coth x)' = -\operatorname{csch}^2 x$
$(\cot x)' = \csc^2 x$ $(\sec x)' = \tan x \sec x$ $(\csc x)' = -\cot x \csc x$	$(\operatorname{sech} x)' = -\operatorname{sech} x \tanh x$ $(\operatorname{csch} x)' = -\operatorname{csch} x \coth x$
$(\arcsin x)' = \frac{1}{\sqrt{1-x^2}}$ $(\arccos x)' = -\frac{1}{\sqrt{1-x^2}}$	$(\operatorname{arcsinh} x)' = \frac{1}{\sqrt{1+x^2}}$ $(\operatorname{arccosh} x)' = \frac{1}{\sqrt{x^2-1}}$
$(\arctan x)' = \frac{1}{1+x^2}$ $(\operatorname{arccot} x)' = -\frac{1}{1+x^2}$	$(\operatorname{arctanh} x)' = \frac{1}{1-x^2}$ $(\operatorname{arccoth} x)' = \frac{1}{x^2-1}$
$(\operatorname{arccsc} x)' = -\frac{1}{x\sqrt{1-x^2}}$	$(\operatorname{arcsech} x)' = -\frac{1}{ x \sqrt{1+x^2}}$ $(\operatorname{arcsech} x)' = -\frac{1}{x\sqrt{1-x^2}}$

## **9.1.1** $ax + b \ (a \neq 0)$

- 1.  $\int \frac{x}{ax+b} dx = \frac{1}{a^2} (ax+b-b \ln |ax+b|) + C$
- 2.  $\int \frac{x^2}{ax+b} dx = \frac{1}{a^3} \left( \frac{1}{2} (ax+b)^2 2b(ax+b) + b^2 \ln|ax+b| \right) + C$
- 3.  $\int \frac{\mathrm{d}x}{x(ax+b)} = -\frac{1}{b} \ln \left| \frac{ax+b}{x} \right| + C$
- 4.  $\int \frac{\mathrm{d}x}{x^2(ax+b)} = -\frac{1}{bx} + \frac{a}{b^2} \ln \left| \frac{ax+b}{x} \right| + C$
- 5.  $\int \frac{x}{(ax+b)^2} dx = \frac{1}{a^2} \left( \ln|ax+b| + \frac{b}{ax+b} \right) + C$
- $6. \int \frac{x^2}{(ax+b)^2} dx = \frac{1}{a^3} \left( ax + b 2b \ln|ax + b| \frac{b^2}{ax+b} \right) + C$   $7. \int \frac{dx}{x(ax+b)^2} = \frac{1}{b(ax+b)} \frac{1}{b^2} \ln\left|\frac{ax+b}{x}\right| + C$

## **9.1.2** $\sqrt{ax+b}$

- 1.  $\int \sqrt{ax+b} \mathrm{d}x = \frac{2}{3a} \sqrt{(ax+b)^3} + C$
- 2.  $\int x\sqrt{ax+b}dx = \frac{2}{15a^2}(3ax-2b)\sqrt{(ax+b)^3} + C$

- $2. \int x\sqrt{ax+b} dx = \frac{2}{15a^2} (3ax-2b)\sqrt{(ax+b)^3} + C$   $3. \int x^2\sqrt{ax+b} dx = \frac{2}{105a^3} (15a^2x^2 12abx + 8b^2)\sqrt{(ax+b)^3} + C$   $4. \int \frac{x}{\sqrt{ax+b}} dx = \frac{2}{3a^2} (ax-2b)\sqrt{ax+b} + C$   $5. \int \frac{2}{\sqrt{ax+b}} dx = \frac{2}{15a^3} (3a^2x^2 4abx + 8b^2)\sqrt{ax+b} + C$   $6. \int \frac{dx}{x\sqrt{ax+b}} = \begin{cases} \frac{1}{\sqrt{b}} \ln \left| \frac{\sqrt{ax+b} \sqrt{b}}{\sqrt{ax+b} + \sqrt{b}} \right| + C \quad (b > 0) \\ \frac{2}{\sqrt{-b}} \arctan \sqrt{\frac{a-b}{b}} + C \quad (b < 0) \end{cases}$   $7. \int \frac{dx}{x^2\sqrt{ax+b}} = -\frac{\sqrt{ax+b}}{bx} \frac{a}{2b} \int \frac{dx}{x\sqrt{ax+b}}$   $8. \int \frac{\sqrt{ax+b}}{x} dx = 2\sqrt{ax+b} + b \int \frac{dx}{x\sqrt{ax+b}}$

- 9.  $\int \frac{\sqrt{ax+b}}{x^2} dx = -\frac{\sqrt{ax+b}}{x} + \frac{a}{2} \int \frac{dx}{x\sqrt{ax+b}}$

## **9.1.3** $x^2 \pm a^2$

- 1.  $\int \frac{\mathrm{d}x}{x^2 + a^2} = \frac{1}{a} \arctan \frac{x}{a} + C$
- 2.  $\int \frac{dx}{(x^2+a^2)^n} = \frac{x}{2(n-1)a^2(x^2+a^2)^{n-1}} + \frac{2n-3}{2(n-1)a^2} \int \frac{dx}{(x^2+a^2)^{n-1}}$
- 3.  $\int \frac{\mathrm{d}x}{x^2 a^2} = \frac{1}{2a} \ln \left| \frac{x a}{x + a} \right| + C$

- 9.1.4  $ax^{2} + b \ (a > 0)$ 1.  $\int \frac{dx}{ax^{2} + b} = \begin{cases} \frac{1}{\sqrt{ab}} \arctan \sqrt{\frac{a}{b}}x + C & (b > 0) \\ \frac{1}{2\sqrt{-ab}} \ln \sqrt{\frac{ax \sqrt{-b}}{\sqrt{ax + \sqrt{-b}}}} + C & (b < 0) \end{cases}$ 2.  $\int \frac{x}{ax^{2} + b} dx = \frac{1}{2a} \ln |ax^{2} + b| + C$ 

  - $3. \int \frac{x^2}{ax^2 + b} dx = \frac{x}{a} \frac{b}{a} \int \frac{dx}{ax^2 + b}$   $4. \int \frac{dx}{x(ax^2 + b)} = \frac{1}{2b} \ln \frac{x^2}{ax^2 + b} + C$   $5. \int \frac{dx}{x^2(ax^2 + b)} = -\frac{1}{bx} \frac{a}{b} \int \frac{dx}{ax^2 + b}$

7. 
$$\int \frac{dx}{(ax^2+b)^2} = \frac{x}{2b(ax^2+b)} + \frac{1}{2b} \int \frac{dx}{ax^2+b}$$

## **9.1.5** $ax^2 + bx + c \ (a > 0)$

1. 
$$\frac{dx}{ax^{2} + bx + c} = \begin{cases}
\frac{2}{\sqrt{4ac - b^{2}}} & \arctan \frac{2ax + b}{\sqrt{4ac - b^{2}}} + C & (b^{2} < 4ac) \\
\frac{1}{\sqrt{b^{2} - 4ac}} & \ln \left| \frac{2ax + b - \sqrt{b^{2} - 4ac}}{2ax + b + \sqrt{b^{2} - 4ac}} \right| + C & (b^{2} > 4ac) \\
2. \int \frac{x}{ax^{2} + bx + c} dx = \frac{1}{2} \ln |ax^{2} + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^{2} + bx + c}
\end{cases}$$

## **9.1.6** $\sqrt{x^2 + a^2}$ (a > 0)

- 1.  $\int \frac{\mathrm{d}x}{\sqrt{x^2 + a^2}} = \operatorname{arsh} \frac{x}{a} + C_1 = \ln(x + \sqrt{x^2 + a^2}) + C$

- 4.  $\int \frac{x}{\sqrt{(x^2 + a^2)^3}} dx = -\frac{1}{\sqrt{x^2 + a^2}} + C$ 5.  $\int \frac{x^2}{\sqrt{x^2 + a^2}} dx = \frac{x}{2} \sqrt{x^2 + a^2} \frac{a^2}{2} \ln(x + \sqrt{x^2 + a^2}) + C$
- 6.  $\int \frac{x^2}{\sqrt{(x^2+a^2)^3}} dx = -\frac{x}{\sqrt{x^2+a^2}} + \ln(x+\sqrt{x^2+a^2}) + C$

- 9.  $\int \sqrt{x^2 + a^2} dx = \frac{x}{2} \sqrt{x^2 + a^2} + \frac{a^2}{2} \ln(x + \sqrt{x^2 + a^2}) + C$ 10.  $\int \sqrt{(x^2 + a^2)^3} dx = \frac{x}{8} (2x^2 + 5a^2) \sqrt{x^2 + a^2} + \frac{3}{8} a^4 \ln(x + \sqrt{x^2 + a^2}) + C$
- 11.  $\int x\sqrt{x^2+a^2}dx = \frac{1}{3}\sqrt{(x^2+a^2)^3} + C$
- 11.  $\int x\sqrt{x^2 + a^2} \, dx = \frac{1}{8}(2x^2 + a^2)\sqrt{x^2 + a^2} \frac{a^4}{8}\ln(x + \sqrt{x^2 + a^2}) + C$ 13.  $\int \frac{\sqrt{x^2 + a^2}}{x} \, dx = \sqrt{x^2 + a^2} + a \ln \frac{\sqrt{x^2 + a^2} a}{|x|} + C$
- 14.  $\int \frac{\sqrt{x^2 + a^2}}{2} dx = -\frac{\sqrt{x^2 + a^2}}{2} + \ln(x + \sqrt{x^2 + a^2}) + C$

# **9.1.7** $\sqrt{x^2 - a^2}$ (a > 0)

- 1.  $\int \frac{\mathrm{d}x}{\sqrt{x^2 a^2}} = \frac{x}{|x|} \operatorname{arch} \frac{|x|}{a} + C_1 = \ln \left| x + \sqrt{x^2 a^2} \right| + C$
- 2.  $\int \frac{dx}{\sqrt{(x^2 a^2)^3}} = -\frac{x}{a^2 \sqrt{x^2 a^2}} + C$ 3.  $\int \frac{x}{\sqrt{x^2 a^2}} dx = \sqrt{x^2 a^2} + C$

- 6.  $\int \frac{x^2}{\sqrt{(x^2 a^2)^3}} dx = -\frac{x}{\sqrt{x^2 a^2}} + \ln|x + \sqrt{x^2 a^2}| + C$ 7.  $\int \frac{dx}{x\sqrt{x^2 a^2}} = \frac{1}{a} \arccos \frac{a}{|x|} + C$

- 9.  $\int \sqrt{x^2 a^2} dx = \frac{x}{2} \sqrt{x^2 a^2} \frac{a^2}{2} \ln|x + \sqrt{x^2 a^2}| + C$ 10.  $\int \sqrt{(x^2 a^2)^3} dx = \frac{x}{8} (2x^2 5a^2) \sqrt{x^2 a^2} + \frac{3}{8} a^4 \ln|x + \sqrt{x^2 a^2}| + C$ 11.  $\int x \sqrt{x^2 a^2} dx = \frac{1}{3} \sqrt{(x^2 a^2)^3} + C$
- 12.  $\int x^2 \sqrt{x^2 a^2} dx = \frac{x}{8} (2x^2 a^2) \sqrt{x^2 a^2} \frac{a^4}{8} \ln|x + \sqrt{x^2 a^2}| + C$
- 13.  $\int \frac{\sqrt{x^2 a^2}}{x} dx = \sqrt{x^2 a^2} a \arccos \frac{a}{|x|} + C$
- 14.  $\int \frac{\sqrt{x^2 a^2}}{2} dx = -\frac{\sqrt{x^2 a^2}}{x} + \ln|x + \sqrt{x^2 a^2}| + C$

# 9.1.8 $\sqrt{a^2 - x^2} (a > 0)$ 1. $\int \frac{dx}{\sqrt{a^2 - x^2}} = \arcsin \frac{x}{a} + C$ 2. $\frac{dx}{\sqrt{(a^2 - x^2)^3}} = \frac{x}{a^2 \sqrt{a^2 - x^2}} + C$ 3. $\int \frac{x}{\sqrt{a^2 - x^2}} dx = -\sqrt{a^2 - x^2} + C$

- 6.  $\int \frac{x^2}{\sqrt{(a^2 x^2)^3}} dx = \frac{x}{\sqrt{a^2 x^2}} \arcsin \frac{x}{a} + C$
- 7.  $\int \frac{\mathrm{d}x}{x\sqrt{a^2 x^2}} = \frac{1}{a} \ln \frac{a \sqrt{a^2 x^2}}{|x|} + C$
- 9.  $\int \sqrt{a^2 x^2} dx = \frac{x}{2} \sqrt{a^2 x^2} + \frac{a^2}{2} \arcsin \frac{x}{a} + C$
- 10.  $\int \sqrt{(a^2 x^2)^3} \, dx = \frac{x}{8} (5a^2 2x^2) \sqrt{a^2 x^2} + \frac{3}{8} a^4 \arcsin \frac{x}{a} + C$
- 11.  $\int \sqrt{u^2 x^2} \, dx = -\frac{1}{3} \sqrt{(a^2 x^2)^3} + C$ 12.  $\int x^2 \sqrt{a^2 x^2} \, dx = \frac{1}{8} (2x^2 a^2) \sqrt{a^2 x^2} + \frac{a^4}{8} \arcsin \frac{x}{a} + C$ 13.  $\int \frac{\sqrt{a^2 x^2}}{x} \, dx = \sqrt{a^2 x^2} + a \ln \frac{a \sqrt{a^2 x^2}}{|x|} + C$
- 14.  $\int \frac{\sqrt{a^2 x^2}}{a^2} dx = -\frac{\sqrt{a^2 x^2}}{x} \arcsin \frac{x}{a} + C$

- 9.1.9  $\sqrt{\pm ax^2 + bx + c}$  (a > 0)1.  $\int \frac{dx}{\sqrt{ax^2 + bx + c}} = \frac{1}{\sqrt{a}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$ 
  - $\sqrt{\frac{ax + ca + c}{4a}}$ 2.  $\int \sqrt{ax^2 + bx + c} dx = \frac{2ax + b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac b^2}{8\sqrt{a^3}} \ln |2ax + b| + \frac{ac b^2}{8\sqrt{a^3}} \ln |2ax +$  $2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$

3. 
$$\int \frac{x}{\sqrt{ax^2 + bx + c}} \, dx = \frac{1}{a} \sqrt{ax^2 + bx + c} - \frac{b}{2\sqrt{a^3}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + \frac{C}{\sqrt{ax^2 + bx + c}} + \frac{C}{\sqrt{ax^2 + bx + c}} + \frac{C}{\sqrt{ax^2 + bx - ax^2}} = -\frac{1}{\sqrt{a}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$
5. 
$$\int \sqrt{c + bx - ax^2} \, dx = \frac{2ax - b}{4a} \sqrt{c + bx - ax^2} + \frac{b^2 + 4ac}{8\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$
6. 
$$\int \frac{x}{\sqrt{c + bx - ax^2}} \, dx = -\frac{1}{a} \sqrt{c + bx - ax^2} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$
9.1.10 
$$\sqrt{\pm \frac{x - a}{x - b}} \, dx = (x - b)\sqrt{\frac{x - a}{x - b}} + (b - a) \ln(\sqrt{|x - a|} + \sqrt{|x - b|}) + C$$
2. 
$$\int \sqrt{\frac{x - a}{b - x}} \, dx = (x - b)\sqrt{\frac{x - a}{b - x}} + (b - a) \arcsin \sqrt{\frac{x - a}{b - x}} + C$$
3. 
$$\int \frac{dx}{\sqrt{(x - a)(b - x)}} = 2 \arcsin \sqrt{\frac{x - a}{b - x}} + C (a < b)$$
4. 
$$\int \sqrt{(x - a)(b - x)} \, dx = \frac{2x - a - b}{b - x} \sqrt{(x - a)(b - x)} + \frac{(b - a)^2}{4} \arcsin \sqrt{\frac{x - a}{b - x}} + C$$

# 9.1.11 Triangular function

- 1.  $\int \tan x dx = -\ln|\cos x| + C$ 2.  $\int \cot x dx = \ln|\sin x| + C$
- $\int \sec x \, dx = \ln \left| \tan \left( \frac{\pi}{4} + \frac{x}{2} \right) \right| + C = \ln \left| \sec x + \tan x \right| + C$
- 4.  $\int \csc x \, \mathrm{d}x = \ln \left| \tan \frac{x}{2} \right| + C = \ln \left| \csc x \cot x \right| + C$

- 5.  $\int \sec^2 x dx = \tan x + C$ 6.  $\int \csc^2 x dx = -\cot x + C$ 7.  $\int \sec x \tan x dx = \sec x + C$ 8.  $\int \csc x \cot x dx = -\csc x + C$
- 9.  $\int \sin^2 x \, dx = \frac{x}{2} \frac{1}{4} \sin 2x + C$
- 10.  $\int \cos^2 x \mathrm{d}x = \frac{x}{2} + \frac{1}{4} \sin 2x + C$

- 11.  $\int \sin^n x dx = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x dx$ 12.  $\int \cos^n x dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x dx$ 13.  $\frac{dx}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\sin^{n-2} x}$ 14.  $\frac{dx}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x}$ 15.

$$\begin{split} & \int \cos^m x \sin^n x dx \\ &= \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x dx \\ &= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x dx \end{split}$$

- 16.  $\int \sin ax \cos bx dx = -\frac{1}{2(a+b)} \cos(a+b)x \frac{1}{2(a-b)} \cos(a-b)x + C$
- 17.  $\int \sin ax \sin bx dx = -\frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$
- 18.  $\int \cos ax \cos bx dx = \frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$

18. 
$$\int \cos ax \cos bx dx = \frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$$
19. 
$$\int \frac{dx}{a+b \sin x} = \begin{cases} \frac{2}{\sqrt{a^2-b^2}} \arctan \frac{a \tan \frac{x}{2} + b}{\sqrt{a^2-b^2}} + C & (a^2 > b^2) \\ \frac{1}{\sqrt{b^2-a^2}} \ln \left| \frac{a \tan \frac{x}{2} + b - \sqrt{b^2-a^2}}{a \tan \frac{x}{2} + b + \sqrt{b^2-a^2}} \right| + C & (a^2 < b^2) \end{cases}$$
20. 
$$\int \frac{dx}{a+b \cos x} = \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan \left(\sqrt{\frac{a-b}{a+b}} \tan \frac{x}{2}\right) + C & (a^2 < b^2) \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln \left| \frac{\tan \frac{x}{2} + \sqrt{\frac{a+b}{b-a}}}{\tan \frac{x}{2} - \sqrt{\frac{a+b}{b-a}}} \right| + C & (a^2 < b^2) \end{cases}$$
21. 
$$\int \frac{dx}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{ab} \arctan \left(\frac{b}{a} \tan x\right) + C$$
22. 
$$\int \frac{dx}{a^2 \cos^2 x - b^2 \sin^2 x} = \frac{1}{2ab} \ln \left| \frac{b \tan x}{b \tan x} \right| + C$$
23. 
$$\int x \sin ax dx = \frac{1}{a^2} \sin ax - \frac{1}{a} x \cos ax + C$$
24. 
$$\int x^2 \sin ax dx = -\frac{1}{a^2} \cos ax + \frac{2}{a^2} x \sin ax + \frac{2}{a^3} \cos ax + C$$
25. 
$$\int x \cos ax dx = \frac{1}{a^2} \cos ax + \frac{1}{a} x \sin ax + C$$
26. 
$$\int x^2 \cos ax dx = \frac{1}{a} x^2 \sin ax + \frac{2}{a^2} x \cos ax - \frac{2}{a^3} \sin ax + C$$
27. 
$$\ln x \cos ax dx = \frac{1}{a} x^2 \sin ax + \frac{2}{a^2} x \cos ax - \frac{2}{a^3} \sin ax + C$$
28. 
$$\ln x \cos ax dx = \frac{1}{a} \cos ax + \frac{2}{a^2} \cos ax - \frac{2}{a^3} \sin ax + C$$
29. 
$$\ln x \cos ax dx = \frac{1}{a} \cos ax + \frac{2}{a^2} x \cos ax - \frac{2}{a^3} \sin ax + C$$
21. 
$$\ln x \cos ax dx = \frac{1}{a} \cos ax + \frac{2}{a^2} \cos ax - \frac{2}{a^3} \sin ax + C$$
21. 
$$\ln x \cos ax dx = \frac{1}{a} \cos ax + \frac{2}{a^2} x \cos ax - \frac{2}{a^3} \sin ax + C$$
22. 
$$\ln x \cos ax dx = \frac{1}{a} \cos ax + \frac{2}{a^2} x \cos ax - \frac{2}{a^3} \sin ax + C$$

## 9.1.12 Inverse triangular function (a > 0)

- 1.12 Inverse triangular function (a > 0)1.  $\int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2 x^2} + C$ 2.  $\int x \arcsin \frac{x}{a} dx = (\frac{x^2}{2} \frac{a^2}{4}) \arcsin \frac{x}{a} + \frac{x}{4} \sqrt{x^2 x^2} + C$ 3.  $\int x^2 \arcsin \frac{x}{a} dx = \frac{x^3}{3} \arcsin \frac{x}{a} + \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 x^2} + C$ 4.  $\int \arccos \frac{x}{a} dx = x \arccos \frac{x}{a} \sqrt{a^2 x^2} + C$ 5.  $\int x \arccos \frac{x}{a} dx = (\frac{x^2}{2} \frac{a^2}{4}) \arccos \frac{x}{a} \frac{x}{4} \sqrt{a^2 x^2} + C$ 6.  $\int x^2 \arccos \frac{x}{a} dx = \frac{x^3}{3} \arccos \frac{x}{a} \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 x^2} + C$ 7.  $\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} \frac{a}{2} \ln(a^2 + x^2) + C$ 8.  $\int x \arctan \frac{x}{a} dx = \frac{1}{2} (a^2 + x^2) \arctan \frac{x}{a} \frac{a}{2} x + C$ 9.  $\int x^2 \arctan \frac{x}{a} dx = \frac{x^3}{3} \arctan \frac{x}{a} \frac{a}{6} x^2 + \frac{a^3}{6} \ln(a^2 + x^2) + C$ 1.13 Exponential function

## 9.1.13 Exponential function

- 1.  $\int a^x dx = \frac{1}{\ln a} a^x + C$ 2.  $\int e^{ax} dx = \frac{1}{a} a^{ax} + C$

- 2.  $\int e^{ax} dx = \frac{1}{a} a^{ax} + C$ 3.  $\int x e^{ax} dx = \frac{1}{a^2} (ax 1) a^{ax} + C$ 4.  $\int x^n e^{ax} dx = \frac{1}{a} x^n e^{ax} \frac{n}{a} \int x^{n-1} e^{ax} dx$ 5.  $\int x a^x dx = \frac{x}{\ln a} a^x \frac{1}{(\ln a)^2} a^x + C$ 6.  $\int x^n a^x dx = \frac{1}{\ln a} x^n a^x \frac{n}{\ln a} \int x^{n-1} a^x dx$ 7.  $\int e^{ax} \sin bx dx = \frac{1}{a^2 + b^2} e^{ax} (a \sin bx b \cos bx) + C$
- 8.  $\int e^{ax} \cos bx dx = \frac{1}{a^2 + b^2} e^{ax} (b \sin bx + a \cos bx) + C$
- 9.  $\int e^{ax} \sin^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \sin^{n-1} bx (a \sin bx nb \cos bx) +$  $\frac{a^2}{a^2+b^2n^2} \int e^{ax} \sin^{n-2} bx dx$
- 10.  $\int e^{ax} \cos^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) +$

 $\tfrac{n(n-1)b^2}{a^2+b^2n^2}\int \mathrm{e}^{ax}\cos^{n-2}bx\mathrm{d}x$ 

## 9.1.14 Logarithmic function

- 1.  $\int \ln x \, dx = x \ln x x + C$ 2.  $\int \frac{dx}{x \ln x} = \ln |\ln x| + C$ 3.  $\int x^n \ln x \, dx = \frac{1}{n+1} x^{n+1} (\ln x \frac{1}{n+1}) + C$ 4.  $\int (\ln x)^n \, dx = x (\ln x)^n n \int (\ln x)^{n-1} \, dx$ 5.  $\int x^m (\ln x)^n \, dx = \frac{1}{m+1} x^{m+1} (\ln x)^n \frac{n}{m+1} \int x^m (\ln x)^{n-1} \, dx$

#### 9.2Regular expression

## 9.2.1 Special pattern characters

Characters	Description
	Not newline
\t	Tab (HT)
\n	Newline (LF)
\A	Vertical tab (VT)
\f	Form feed (FF)
\r	Carriage return (CR)
\cletter	Control code
\xhh	ASCII character
\uhhhh	Unicode character
\0	Null
\int	Backreference
\d	Digit
\D	Not digit
\s	Whitespace
\S	Not whitespace
/W	Word (letters, numbers and the underscore)
\W	Not word
\character	Character
[class]	Character class
[^class]	Negated character class
[^class]	Negated character class

#### 9.2.2Quantifiers

Characters	Times
*	0 or more
+	1 or more
?	0 or 1
{int}	int
{int,}	int or more
{min,max}	Between min and max

By default, all these quantifiers are greedy (i.e., they take as many characters that meet the condition as possible). This behavior can be overridden to ungreedy (i.e., take as few characters that meet the condition as possible) by adding a question mark (?) after the quantifier.

## 9.2.3 Groups

Characters	Description
(subpattern)	Group with backreference
(?:subpattern)	Group without backreference

#### 9.2.4 Assertions

Characters	Description
^	Beginning of line
\$	End of line
\b	Word boundary
\B	Not a word boundary
(?=subpattern)	Positive lookahead
(?!subpattern)	Negative lookahead

## Alternative

A regular expression can contain multiple alternative patterns simply by separating them with the separator operator (|): The regular expression will match if any of the alternatives match, and as soon as

# one does. 9.2.6 Character classes

Class	Description
[:alnum:]	Alpha-numerical character
[:alpha:]	Alphabetic character
[:blank:]	Blank character
[:cntrl:]	Control character
[:digit:]	Decimal digit character
[:graph:]	Character with graphical representation
[:lower:]	Lowercase letter
[:print:]	Printable character
[:punct:]	Punctuation mark character
[:space:]	Whitespace character
[:upper:]	Uppercase letter
[:xdigit:]	Hexadecimal digit character
[:d:]	Decimal digit character
[:w:]	Word character
[:s:]	Whitespace character

Please note that the brackets in the class names are additional to those opening and closing the class definition. For example:

[[:alpha:]] is a character class that matches any alphabetic character.

- acter. [abc[:digit:]] is a character class that matches a, b, c, or a
- digit.

  [^[:space:]] is a character class that matches any character ex-