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## **Environment**

#### 1.1 Vimrc

## Data Structure

### 2.1 KD tree

```
/* kd_tree : finds the k-th closest point in O(kn^{1-\frac{1}{k}}). Usage : Stores the data in p[]. Call function init (n, k). Call min_kth (d, k). (or max_kth) (k is 1-based)
     2 Usage
 | based)
| Note : Switch to the commented code for Manhattan distance.
| Status : SPOJ-FAILURE Accepted.*/
| template <int MAXN = 200000, int MAXK = 2>
| struct kd_tree {
| int k, size;
| struct point { int data[MAXK], id; } p[MAXN];
| struct kd_node {
| int l, r; point p, dmin, dmax;
| kd_node (const point &rhs) : l (-1), r (-1), p (rhs) |
| , dmin (rhs), dmax (rhs) {}
| void merge (const kd_node &rhs, int k) {
| for (register int i = 0; i < k; ++i) {
| dmin.data[i] = std::max (dmax.data[i], rhs.dmin. data[i]); }
| long long min_dist (const point &rhs, int k) const {
| register long long ret = 0; |
| for (register int i = 0; i < k; ++i) {
| if (dmin.data[i] = rhs.data[i] & rhs.data[i] <= dmax.data[i] > continue; |
| ret += std::min (111 * (dmin.data[i] - rhs.data[i]) * (dmax. data[i] - rhs.data[i]) * (dmax. data[i] - rhs.data[i]) * (ret += std::max (0, rhs.data[i] - rhs.data[i]); |
| ret += std::max (0, dmin.data[i] - rhs.data[i]); |
| ret trun ret; |
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long ret = 0; |
| for (int i = 0; i < k; ++i) {
| int tmp = std::max (std::abs (dmin.data[i] - rhs.data[i]) |
| ret += std::max (std::abs (dmin.data[i] - rhs.data[i]) |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] |
| r
      3 Note
                                                           Switch to the commented code for Manhattan
                                            distance.
21
23 //
28
                                         | ]));
ret += 111 * tmp * tmp; }
  ret += std::max (std::abs (rhs.data[i] - dmax.
data[i]) + std::abs (rhs.data[i] - dmin.data[i]));
                   return ret; } tree[MAXN * 4];
struct result {
 long long dist; point d; result() {}
 result (const long long &dist, const point &d) :
    dist (dist), d (d) {}
 bool operator > (const result &rhs) const { return
 32
34
                35
 38
                             if ("tree[rt].r) tree[rt].merge (tree[tree[rt].r], k
                  if (~tree[rt].r) tree[rt].merge (tree[tree[rt].r], k
    ); }
std::priority_queue<result, std::vector<result>, std
    ::less <result>> heap_1;
std::priority_queue<result, std::vector<result>, std
    ::greater <result>> heap_r;
void _min_kth (const int &depth, const int &rt, const
    int &m, const point &d) {
    result tmp = result (sqrdist (tree[rt].p, d), tree[
        rt].p);
    if ((int)heap_1.size() < m) heap_1.push (tmp);
    else if (tmp < heap_1.top()) {
        heap_1.push (tmp); }
</pre>
 55
```

```
62
74
75
80
```

#### Splay 2.2

```
void push_down (int x) {
   if ( n[x].c[0]) push (n[x].c[0], n[x].t);
   if ( n[x].c[1]) push (n[x].c[1], n[x].t);
   if ( n[x].t = tag (); }
   void update (int x) {
      n[x].m = gen (x);
   if ( n[x].c[0]) n[x].m = merge (n[n[x].c[0]].m, n[x].
    m);
if (\tilde{n}[x].c[1]) n[x].m = merge (n[x].m, n[n[x].c[1]].
  11
12
```

## Link-cut tree

```
= u;

n[u].c[1] = v;

if (~v) n[v].f = u, n[v].p = -1;

update (u); u = n[v = u].p; }

splay (x); }
```

## Formula

### Zellers congruence 3.1

```
/* Zeller's congruence : converts between a calendar
```

## 3.2 Lattice points below segment

```
1 /* Euclidean-like algorithm : computes the sum of
     \sum_{i=0}^{n-1} \left[ \frac{a+bi}{m} \right] . \star /
```

```
if (b >= m) return (n - 1) * n / 2 * (b / m) + solve
    (n, a, b % m, m);
return solve ((a + b * n) / m, (a + b * n) % m, m, b)
```

#### Adaptive Simpson's method 3.3

```
|\mathbf{r}| / \star Adaptive Simpson's method : integrates f in [1, r].
  struct simpson {
  double area (double (*f) (double), double 1, double r
   double m = 1 + (r - 1) / 2;
return (f (1) + 4 * f (m) + f (r)) * (r - 1) / 6; }
double solve (double (*f) (double), double 1, double
    r, double eps, double a) {
    double m = 1 + (r - 1) / 2;
    double left = area (f, 1, m), right = area (f, m, r)
    :
```

#### 3.4Neural network

```
1 /* Neural network : machine learning. */
2 template <int ft = 6, int n = 6, int MAXDATA = 100000>
3 struct network {
      double wp[n][ft], w[n], avg[ft + 1], sig[ft + 1], val
        network () {
std::mt19937_64 mt (time (0));
std::uniform_real_distribution <double> urdp (0, 2 *
sqrt (ft));
      network
        double
      void desc (double *x, double t, double eta) {
  double o = compute (x), del = (o - t); // * o * (1 -
         o)
for (int i = 0; i < n; ++i) for (int j = 0; j < ft;
21
           wp[i][j] -= eta * del * val[i] * (1 - val[i]) * w[i
    ] * x[j];
22
      26
     dn; ++j)
sig[i] += (data[j][i] - avg[i]) * (data[j][i]
for (int i = 0; i < ft + 1; ++i) sig[i] = sqrt (sig[i] / dn);
for (int i = 0; i < ft + 1; ++i) for (int j = 0; j < dn; ++j)
data[j][i] = (data[j][i] - avg[i]) / sig[i];
for (int cnt = 0; cnt < epoch; ++cnt) for (int test = 0; test < dn; ++test)
desc (data[test], data[test][ft], eta);
double predict (double *x) {
  for (int i = 0; i < ft; ++i) x[i] = (x[i] - avg[i])
  / sig[i];
  return compute (x) * sig[ft] + avg[ft]; }
</pre>
33
      return compute (x) * sig[ft] + avg[ft]; }
std::string to_string () {
  std::ostringstream os; os.precision (20); os << std</pre>
        stat.ostringscrem = ,
    ::fixed;
for (int i = 0; i < n; ++i) for (int j = 0; j < ft;
    ++j) os << wp[i][j] << "_";
for (int i = 0; i < n; ++i) os << w[i] << "_";
for (int i = 0; i < ft + 1; ++i) os << avg[i] << "_"</pre>
42
         for (int i = 0; i < ft + 1; ++i) os << sig[i] << "_"
      return os.str (); }
void read (const std::string &str) {
    std::istringstream is (str);
    for (int i = 0; i < n; ++i) for (int j = 0; j < ft;
        ++j) is >> wp[i][j];
    for (int i = 0; i < n; ++i) is >> w[i];
    for (int i = 0; i < ft + 1; ++i) is >> avg[i];
    for (int i = 0; i < ft + 1; ++i) is >> sig[i]; } };
```

### Number theory 4

### 4.1 Fast power module

```
1 /* Fast power module : x^n */ 2 int fpm (int x, int n, int mod) {
```

```
int ans = 1, mul = x; while (n) {
  if (n & 1) ans = int (111 * ans * mul * mod);
  mul = int (111 * mul * mul * mod); n >>= 1; }
  return ans; }
```

## 4.2 Euclidean algorithm

```
\frac{1}{2} /* Euclidean algorithm : solves for ax + by = gcd (a,
  b). */
void euclid (const long long &a, const long long &b,
    long long &x, long long &y) {
    if (b == 0) x = 1, y = 0;
    else euclid (b, a % b, y, x), y -= a / b * x; }
  long long inverse (long long x, long long m) {
  long long a, b; euclid (x, m, a, b); return (a % m +
     m) % m; }
```

## 4.3 Discrete Fourier transform

```
1 /* Discrete Fourier transform : the nafarious you-know
  -what thing.

2 Usage: call init for the suggested array size, and solve for the transform. (use f!=0 for the inverse
        template <int MAXN = 1000000>
 3 template <int MAXN = 1000000>
4 struct dft {
5    typedef std::complex <double> complex;
6    complex e[2][MAXN];
7    int init (int n) {
8        int len = 1;
9        for (; len <= 2 * n; len <<= 1);
10        for (int i = 0; i < len; ++i) {
11        e[0][i] = complex (cos (2 * PI * i / len), sin (2 * PI * i / len));
12        e[1][i] = complex (cos (2 * PI * i / len), -sin (2 * PI * i / len));
13        return len; }
14        void solve (complex *a, int n, int f) {</pre>
11
             return len; }
void solve (complex *a, int n, int f) {
  for (int i = 0, j = 0; i < n; ++i) {
    if (i > j) std::swap (a[i], a[j]);
    for (int t = n >> 1; (j ^= t) < t; t >>= 1); }
  for (int i = 2; i <= n; i <<= 1)
    for (int j = 0; j < n; j += i)
    for (int k = 0; k < (i >> 1); ++k) {
      complex A = a[j + k];
      complex B = e[f][n / i * k] * a[j + k + (i >> 1)
      ];
}
                 23
```

## Fast Walsh-Hadamard transform

```
1 /* Fast Walsh-Hadamard transform : binary operation
  /* Fast Walsh-Hadamara cransform. transform. */
void fwt (int *a, int n, int w) {
  for (int i = 1; i < n; i <<= 1)
    for(int j = 0; j < n; j += i << 1) {
      for(int k = 0; k < i; ++k) {
        int x = a[j + k], y = a[i + j + k];
        if (w) {
           }else{
/* xor : a[j + k] = x + y, a[i + j + k] = x - y,
    and : a[j + k] = x + y, or : a[i + j + k] = x
    + y; */
```

### Number theoretic transform 4.5

```
int mod, int prt) {
       inv[i][j] = (int) inverse (MOD[i], MOD[j]);
static int x[3];
for (int i = 0; i < 3; ++i) { x[i] = a[i];
    for (int j = 0; j < i; ++j) {
        int t = (x[i] - x[j] + MOD[i]) % MOD[i];</pre>
30
```

```
if (t < 0) t += MOD[i];
    x[i] = int (1LL * t * inv[j][i] % MOD[i]); }
int sum = 1, ret = x[0] % mod;
for (int i = 1; i < 3; ++i) {
    sum = int (1LL * sum * MOD[i - 1] % mod);
    ret += int (1LL * x[i] * sum % mod);
    if (ret >= mod) ret -= mod; }
return ret; };
```

## 4.6 Chinese remainder theorem

## 4.7 Linear Recurrence

## 4.8 Berlekamp Massey algorithm

## 4.9 Baby step giant step algorithm

## 4.10 Miller Rabin primality test

## 4.11 Pollard's Rho algorithm

## 5 Geometry

```
#define cd const double &
const double EPS = 1E-8, PI = acos (-1);
int sgn (cd x) { return x < -EPS ? -1 : x > EPS; }
int cmp (cd x, cd y) { return sgn (x - y); }
double sqr (cd x) { return x * x; }
```

## 5.1 Point

### 5.2 Line

## 5.3 Circle

```
11| std::vector <point> line_circle_intersect (cl a, cc b)
    if (cmp (point_to_line (b.c, a), b.r) > 0) return std
    ::vector <point> ();
double x = sqrt (sqr (b.r) - sqr (point_to_line (b.c,
a)));
return ret; }
//Counter-clockwise with respect of point Oa.
std::vector line> intangent (cc c1, cc c2) {
  point p = (b.c * a.r + a.c * b.r) / (a.r + b.r);
  std::vector pp = tangent (p, a), qq = tangent (p, b);
  if (pp.size () == 2 && qq.size () == 2) {
    ret.push_back (line (pp[0], qq[0]));
    ret.push_back (line (pp[1], qq[1]));
  return ret; }
```

## 5.4 Centers of a triangle

```
point incenter (cp a, cp b, cp c) {
   double p = dis (a, b) + dis (b, c) + dis (c, a);
   return (a * dis (b, c) + b * dis (c, a) + c * dis (a, b)) / p; }

   point circumcenter (cp a, cp b, cp c) {
      point p = b - a, q = c - a, s (dot (p, p) / 2, dot (q, q) / 2);
   return a + point (det (s, point (p.y, q.y)), det (point (p.x, q.x), s)) / det (p, q); }

   point orthocenter (cp a, cp b, cp c) { return a + b + c - circumcenter (a, b, c) * 2; }
```

## 5.5 Fermat point

```
| /* Fermat point : finds a point P that minimizes |PA| + |PB| + |PC|. */
2 point fermat_point (cp a, cp b, cp c) {
3 if (a == b) return a; if (b == c) return b; if (c == a) return c;
4 double ab = dis (a, b), bc = dis (b, c), ca = dis (c, a);
5 double cosa = dot (b - a, c - a) / ab / ca;
6 double cosb = dot (a - b, c - b) / ab / bc;
7 double cosc = dot (b - c, a - c) / ca / bc;
8 double sq3 = PI / 3.0; point mid;
9 if (sgn (cosa + 0.5) < 0) mid = a;
10 else if (sgn (cosb + 0.5) < 0) mid = b;
11 else if (sgn (cosc + 0.5) < 0) mid = c;
12 else if (sgn (det (b - a, c - a)) < 0) mid = 1 line_intersect (line (a, b + (c - b).rot (sq3)), line (b, c + (a - c).rot (sq3)));
13 else mid = line_intersect (line (a, c + (b - c).rot (sq3)));
14 return mid; }
```

### 5.6 Convex hull

```
1 //Counter-clockwise, with minimum number of points.
2 bool turn_left (cp a, cp b, cp c) { return sgn (det (b - a, c - a)) >= 0; }
3 std::vector <point> convex_hull (std::vector <point> a
    ) {
4 int cnt = 0; std::sort (a.begin (), a.end ());
5 std::vector <point> ret (a.size () << 1, point ());
6 for (int i = 0; i < (int) a.size (); ++i) {</pre>
```

## Half plane intersection

```
\mathbf{p}' /* Online half plane intersection : complexity O(n)
 ine p) {
std::vector <point> ret;
if (c.empty ()) return ret;
for (int i = 0; i < (int) c.size (); ++i) {
  int j = (i + 1) % (int) c.size ();
  if (turn_left (p.s, p.t, c[i])) ret.push_back (c[i])</pre>
      if ('wo_side (c[i], c[j], p)) ret.push_back (
    line_intersect (p, line (c[i], c[j]))); }
     return ret; }
* Offline half plane intersection : complexity
second.t - a.second.s, b.second.t - a.second.s)
< 0;
else return cmp (a.first, b.first) < 0; });
h.resize (std::unique (g.begin (), g.end (), [] (
    const polar &a, const polar &b) { return cmp (a.
    first, b.first) == 0 }) - g.begin ());
for (int i = 0; i < (int) h.size (); ++i) h[i] = g[i]</pre>
21
    24
      25
             fore:
     ret[++rear] = h[i]; }
while (rear - fore > 1 && !turn_left (ret[fore],
    line_intersect (ret[rear - 1], ret[rear]))) --
rear;
     while (rear - fore > 1 && !turn_left
28
            line_intersect (ret[fore], ret[fore + 1])))
     fore;
if (rear - fore < 2) return std::vector <point> ();
     fit (rear = lote < 2) return std:.vector <point> (),
std::vector <point> ans; ans.resize (rear + 1);
for (int i = 0; i < rear + 1; ++i) ans[i] =
    line_intersect (ret[i], ret[(i + 1) % (rear + 1)</pre>

]);
return ans; }
```

## 5.8 Nearest pair of points

```
1 /* Nearest pair of points : [1, r), need to sort p
    first. */
    /* Nearest pair or points : [1, r), need to sort p
    first. */
double solve (std::vector <point> &p, int 1, int r) {
    if (1 + 1 >= r) return INF;
    int m = (1 + r) / 2; double mx = p[m].x; std::vector
    <point> v;
    double ret = std::min (solve(p, 1, m), solve(p, m, r))
        );
for (int i = 1; i < r; ++i)
    if (sqr (p[i].x - mx) < ret) v.push_back (p[i]);
sort (v.begin (), v.end (), [&] (cp a, cp b) { return
        a.y < b.y; });
for (int i = 0; i < v.size (); ++i)
for (int j = i + 1; j < v.size (); ++j) {
    if (sqr (v[i].y - v[j].y) > ret) break;
    ret = min (ret, dis2 (v[i] - v[j])); }
return ret; }
```

## 5.9 Minimum circle

```
1 circle minimum_circle (std::vector <point> p) {
2 circle ret; std::random_shuffle (p.begin (), p.end ()
```

## 5.10 Intersection of a polygon and a circle

```
struct polygon_circle_intersect {
double sector_area (cp a, cp b, const double &r) {
double c = (2.0 * r * r - dis2 (a, b)) / (2.0 * r *
      r);
return r * r * acos (c) / 2.0; }
double area (cp a, cp b, const double &r) {
    double dA = dot (a, a), dB = dot (b, b), dC =
        point_to_segment (point (), line (a, b));
    if (sgn (dA - r * r) <= 0 && sgn (dB - r * r) <= 0)
        return det (a, b) / 2.0;
    point tA = a.unit () * r, tB = b.unit () * r;
    if (sgn (dC - r) > 0) return sector_area (tA, tB, r);
```

### Union of circles

```
5.11
ang; }
 void addevent(cc a, cc b, std::vector <event> &evt,
    int &cnt) {
    double d2 = dis2 (a.c, b.c), d_ratio = ((a.r - b.r)
        * (a.r + b.r) / d2 + 1) / 2,
    p_ratio = sqrt (std::max (0., -(d2 - sqr(a.r - b.r)
        ) * (d2 - sqr(a.r + b.r)) / (d2 * d2 * 4)));
    point d = b.c - a.c, p = d.rot(PI / 2), q0 = a.c + d
        * d_ratio + p * p_ratio, q1 = a.c + d * d_ratio
        - p * p_ratio;
    double ang0 = atan2 ((q0 - a.c).y, (q0 - a.c).x),
        ang1 = atan2 ((q1 - a.c).x, (q1 - a.c).y);
    evt.emplace_back(q1, ang1, 1); evt.emplace_back(q0,
        ang0, -1); cnt += ang1 > ang0; }
bool same(cc a, cc b) { return sgn (dis (a.c, b.c))
        = 0 && sgn (a.r - b.r) == 0; }
bool overlap(cc a, cc b) { return sgn (a.r - b.r -
        dis (a.c, b.c)) >= 0; }
bool intersect(cc a, cc b) { return sgn (dis (a.c, b.
   void addevent(cc a, cc b, std::vector <event> &evt,
```

#### 6 Graph

## 6.1 Hopcoft-Karp algorithm

```
1 /* Hopcoft-Karp algorithm : unweighted maximum
matching for bipartition graphs with complexity
\begin{array}{c} O(m\sqrt{n}). \  \, \star/\\ \text{2 template} < \text{int MAXN} = 100000, \ \text{int MAXM} = 100000>\\ \text{3 struct hopcoft\_karp} \{\\ \text{4 using edge\_list} = \text{std}::\text{vector} < \text{int}> [\text{MAXN}];\\ \text{5 int mx}[\text{MAXN}], \ \text{my}[\text{MAXM}], \ \text{lv}[\text{MAXN}];\\ \text{6 bool dfs (edge\_list} < \text{MAXN}, \ \text{MAXM}> \&e, \ \text{int} \ \text{x}) \{\\ \text{7 for (int} \ y : e[x]) \{\\ \text{8 int} \ \text{w} = \text{my}[y];\\ \text{9 if } (!^{\infty} \ \text{w} \ | \ (\text{lv}[x] + 1 == \text{lv}[\text{w}] \&\& \ \text{dfs} \ (e, \ \text{w}))) \{\\ \text{10 } \ \text{mx}[x] = y; \ \text{my}[y] = x; \ \text{return true}; \ \}\\ \text{11 } \ \text{lv}[x] = -1; \ \text{return false}; \ \}\\ \text{12 int solve (edge\_list} \&e, \ \text{int} \ \text{n, int} \ \text{m}) \{\\ \text{13 } \ \text{std}::fill \ (\text{mx}, \ \text{mx} + \text{n, } -1); \ \text{std}::fill \ (\text{my}, \ \text{my} + \text{m, } -1);\\ \text{14 } \ \text{for (int} \ \text{ans} = 0: : ) \ \}\\ \end{array}
                                    for (int ans = 0; ; ) {
  std::vector <int> q;
```

```
for (int i = 0; i < n; ++i)
if (mx[i] == -1) {
  lv[i] = 0; q.vush_back (i);
} else lv[i] = -1;
for (int head = 0; head < (int) q.size(); ++head) {
  int restricted.</pre>
17
18
                                                                                                       int solve
          int x = q[head];
for (int y : e[x]) { int w = my[y]; if (~w && lv[w
        ] < 0) {</pre>
```

## 6.2 Kuhn-Munkres algorithm

```
/* Kuhn Munkres algorithm : weighted maximum ming algorithm for bipartition graphs with complexity O(N^3).
| O(N°).
| Note: The graph is 1-based. */
| template <int MAXN = 500>
| struct kuhn_munkres {
| int n, w[MAXN] [MAXN], lx[MAXN], ly[MAXN], m[MAXN],
| way[MAXN]; | bool u[MAXN];
| void hungary(int x) {
| m[0] = x; int j0 = 0;
| std::fill (sl, sl + n + 1, INF); std::fill (u, u + n + 1, false);
| do {
        11
```

## 6.3 Blossom algorithm

```
template <int MAXN = 500, int MAXM = 250000>
struct blossom {
  using edge_list = std::vector <int> [MAXN];
  int match[MAXN], d[MAXN], fa[MAXN], c1[MAXN], c2[MAXN]
  ], v[MAXN], q[MAXN];
  int *qhead, *qtail;
  struct {
   int fa[MAXN]
 1 /* Blossom algorithm : maximum match for general graph
    10
11
       ror (int i = x; i != y; i = urs.rind (ra[i])) V[i] = -1;
v[y] = -1; return x; }
roid contract (int x, int y, int b) {
for (int i = urs.find (x); i != b; i = urs.find (fa[i])) {
urs.rind (ra[i]) }
    31
36
39
              } else {
  fa[dest] = loc; fa[match[dest]] = dest;
  d[dest] = 1; d[match[dest]] = 0;
  *qtail++ = match[dest];
} else if (d[ufs.find (dest)] == 0) {
  int b = lca (loc, dest, root);
```

```
contract (loc, dest, b); contract (dest, loc, b)
return 0; } }}
return 0; }
int solve (int n, const edge_list &e) {
std::fill (fa, fa + n, 0); std::fill (c1, c1 + n, 0)
return re; } };
```

## 6.4 Weighted blossom algorithm

```
/* Weighted blossom algorithm (vfleaking ver.) : maximum matching for general weighted graphs with
                    complexity O(n^3).
   Usage: Set n to the size of the vertices. Run init ()
. Set g[][].w to the weight of the edge. Run solve
  ().

3 The first result is the answer, the second one is the number of matching pairs. Obtain the matching with match[].

4 Note: 1-based. */

5 struct weighted_blossom {
6 static const int INF = INT_MAX, MAXN = 400;
7 struct edge{ int u, v, w; edge (int u = 0, int v = 0, int w = 0): u(u), v(v), w(w) {} };

2 int w = 0): u(u), v(v), w(w) {} };
         int n, n_x;
edge g[MAXN * 2 + 1][MAXN * 2 + 1];
int lab[MAXN * 2 + 1], match[MAXN * 2 + 1], slack[
MAXN * 2 + 1], st[MAXN * 2 + 1], pa[MAXN * 2 +
          int> q;
          int e_delta (const edge &e) { return lab[e.u] + lab[e
    .v] - g[e.u][e.v].w * 2; }
void update_slack (int u, int x) { if (!slack[x] ||
    e_delta (g[u][x]) < e_delta (g[slack[x]][x]))
    slack[x] = u; }</pre>
        28
29
         void augment (int u, int v) {
  for (; ; ) {
    int xnv = st[match[u]]; set_match (u, v);
    if (!xnv) return; set_match (xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv; }
int get_lca (int u, int v) {
    static int t = 0;
  for (++t; u || v; std::swap (u, v)) {
    if (u == 0) continue; if (vis[u] == t) return u;
    vis[u] = t; u = st[match[u]]; if (u) u = st[pa[u]];
  }
return 0: }
          void augment (int u, int v) {
         return 0; }
void add_blossom (int u, int lca, int v) {
  int b = n + 1; while (b <= n_x && st[b]) ++b;
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca]; flower[b].clear ();
  flower[b].push_back (lca);
  for (int x = u, y; x != lca; x = st[pa[y]]) {
    flower[b].push_back (x), flower[b].push_back (y =
        st[match[x]]), q_push (y); }
std::reverse (flower[b].begin () + 1, flower[b].end
    ()):
```

```
for (int i = 0; i < pr; i += 2) {
  int xs = flower[b][i], xns = flower[b][i + 1];
  pa[xs] = g[xns][xs].u; S[xs] = 1, S[xns] = 0;
  slack[xs] = 0, set_slack(xns); q_push(xns); }
S[xr] = 1, pa[xr] = pa[b];
for {size_t i = pr + 1; i < flower[b].size (); ++i)</pre>
          int xs = flower[b][i]; S[xs] = -1, set_slack(xs); }
st[b] = 0; }
         st[b] = 0; }
pool on_found_edge (const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1; int nu = st[match[v]];
    slack[v] = slack[nu] = 0; S[nu] = 0, q_push(nu);
  } else if(S[v] == 0) {
  int lca = get_lca(u, v);
  if (!lca) return augment(u, v), augment(v, u), true
     103
110
112
115
          return false; }
       seturn raise; ;
sd::pair <long long, int> solve () {
  memset (match + 1, 0, sizeof (int) * n); n_x = n;
  int n_matches = 0; long long tot_weight = 0;
  for (int u = 0; u <= n; ++u) st[u] = u, flower[u].</pre>
118
120
          clear();
int w_max = 0;
121
      123
124
```

### 6.5 Maximum flow

```
/* Sparse graph maximum flow : isap.*/
2 template <int MAXN = 1000, int MAXM = 100000>
3 struct isap {
4 struct flow_edge_list {
5 int size, begin[MAXN], dest[MAXM], next[MAXM], flow[
MAXM];
         maxmj;
void clear (int n) { size = 0; std::fill (begin,
    begin + n, -1); }
flow_edge_list (int n = MAXN) { clear (n); }
void add_edge (int u, int v, int f) {
  dest[size] = v; next[size] = begin[u]; flow[size] =
    f; begin[u] = size++;
  dest[size] = u: next[size] = begin[v]: flow[size] =
```

```
e (p != s) { p = pre[p]; e.flow[cur[p]] -
dflow; e.flow[cur[p] ^ 1] += dflow; } }
                    while
              } else {
  int mindist = n + 1
     int mindist = n + 1;
for (int i = e.begin[u]; ~i; i = e.next[i])
  if (e.flow[i] && mindist > d[e.dest[i]]) {
    mindist = d[e.dest[i]]; cur[u] = i; }
  if (!--gap[d[u]]) return maxflow;
  gap[d[u] = mindist + 1]++; u = pre[u]; } }
return maxflow; };
/* Dense graph maximum flow : dinic. */
template <int MAXN = 1000, int MAXM = 100000>
struct dinic. //
42
43
50
52
       if (d[e.dest[k]] == d[u] + 1 && e.flow[k] > 0) {
   int flow = dfs (e, e.dest[k], std::min (e.flow[k],
        ext));
   if (flow > 0) {
      e.flow[k] -= flow, e.flow[k ^ 1] += flow;
      ret += flow, ext -= flow; } }
   if (!~k) d[u] = -1; return ret; }
   int solve (flow_edge_list &e, int n_, int s_, int t_)
           int ans = 0; n = n_; s = s_; dinic::t = t_;
while (bfs (e)) {
  for (int i = 0; i < n; ++i) w[i] = e.begin[i];
    ans += dfs (e, s, INF); }
return ans; };</pre>
```

## 6.6 Minimum cost flow

24 26 27

```
int x = queue[head];
for (int i = e.begin[x]; ~i; i = e.next[i]) {
   int y = e.dest[i];
   if (e.flow[i] && dist[y] > dist[x] + e.cost[i]) {
      dist[y] = dist[x] + e.cost[i]; prev[y] = i;
      if (!occur[y]) {
        occur[x] = false; }
   return dist[t] < INF; }

std::pair <int, int> solve (cost_flow_edge_list &e, int n_, int s_, int t_) {
   n = n_; s = s_; t = t_; std::pair <int, int> ans = std::make_pair (0, 0);

while (augment (e)) {
   int num = INF;
   for (int i = t; i != s; i = e.dest[prev[i] ^ 1]) {
      num = std::min (num, e.flow[prev[i]); }
   ans.first += num;
   for (int i = t; i != s; i = e.dest[prev[i] ^ 1]) {
      e.flow[prev[i]] -= num; e.flow[prev[i] ^ 1] += num
      ans.second += num * e.cost[prev[i]]: }
}
                                                                          int x = queue[head];
ans.second += num * e.cost[prev[i]]; } }
securn ans; } ;

                                                                                  ans.second += num * e.cost[prev[i]]; } }
```

```
int modlable()
    int delta = INF;
for (int i = 0; i < n; i++) {
   if (!visit[i] && slack[i] < delta) delta = slack[i]</pre>
51
    slack[i] = INF; }
if (delta == INF) return 1;
for (int i = 0; i < n; i++) if (visit[i]) dis[i] +=
    delta;
return 0; }</pre>
  e.flow[i] -= delta; e.flow[i ^ 1] += delta; left
    -= delta;
        if (!left) { visit[x] = false; return flow; }
else
   std::fill (visit + 1, visit + t + 1, 0);
} while (dfs (e, s, INF)); } while (!modlable ());
return std::make_pair (tf, tc);
```

## 6.7 Stoer Wagner algorithm

## 6.8 DN maximum clique

```
(0) {} };

17 std::vector <StepCount> S;
17 std::vector <StepCount> S;
18 bool cut1 (const int pi, const ColorClass &A) {
19    for (int i = 0; i < (int) A.size (); ++i)
20    if (e[pi][A[i]]) return true; return false; }
21 void cut2 (const Vertices &A, Vertices & B) {
22    for (int i = 0; i < (int) A.size () - 1; ++i)
23    if (e[A.back().i][A[i].i]) B.push_back(A[i].i); }
24 void color_sort (Vertices & R) {
25    int j = 0, maxno = 1, min_k = std::max ((int) QMAX.
26    size () - (int) Q.size() + 1, 1);
27    for (int i = 0; i < (int) R.size (); ++i) {
28        int pi = R[i].i, k = 1; while (cut1(pi, C[k])) ++k;
29    if (k > maxno) maxno = k, C[maxno + 1].clear();
```

```
30| C[k].push_back (pi); if (k < min_k) R[j++].i = pi; }
31| if (j > 0) R[j - 1].d = 0;
32| for (int k = min_k; k <= maxno; ++k)
33| for (int i = 0; i < (int) C[k].size (); ++i)
34| R[j].i = C[k][i], R[j++].d = k; }
35| void expand_dyn (Vertices &R) {
36| S[level].i1 = S[level].i1 + S[level - 1].i1 - S[level | ].i2;
             Q.push_back (R.back ().i); Vertices Rp; cut2 (R, Rp
Q.pusn_back (R.back ().1); vertices Rp; Cut2 (R, Rp);

if ((int) Rp.size ()) {
   if ((float) S[level].i1 / ++pk < Tlimit)
        degree_sort (Rp);

   degree_sort (Rp);

color_sort (Rp); ++S[level].i1, ++level;
   expand_dyn (Rp); --level;

expand_dyn (Rp); --level;

belse if ((int) Q.size () > (int) QMAX.size ())
   QMAX = Q;

Q.pop_back (); } else return; R.pop_back (); } }

desc_degrees (V); std::sort(V.begin (), V.end (),
   desc_degree); init_colors (V);

sexpand_dyn (V); sz = (int) QMAX.size ();

for (int i = 0; i < (int) QMAX.size ();

sexpand_dyn (V); sz = (int) QMAX.size ();

for (int i = 0; i < (int) QMAX.size ();

maxclique[i] = QMAX[i]; }

void degree_sort (Vertices & R) {
   set_degrees (R); std::sort(R.begin (), R.end (),
   desc_degree); }

Maxilians (descript Rp); storn genet int sa genet float
 desc_degree); }

desc_degree); }

Maxclique (const BB *conn, const int sz, const float tt = .025) : pk (0), level (1), Tlimit (tt) for (int i = 0; i < sz; i++) V.push_back (Vertex (i)); e = conn, C.resize (sz + 1), S.resize (sz + 1); };
                                             int ans, sol[N]; for (...) e[x][y] = e[y][x]
 59| Maxclique mc (e, n); mc.mcqdyn (sol, ans); //0-based.
60| for (int i = 0; i < ans; ++i) std::cout << sol[i] <<
                          std::endl;
```

## 6.9 Dominator tree

```
/* Dominator tree : finds the immediate dominator (
   idom[]) of each node, idom[x] will be x if x does
   not have a dominator, and will be -1 if x is not
reachable from s. */
2 template <int MAXN = 100000, int MAXM = 100000>
 21
22
```

# String

#### Suffix Array 7.1

```
/* Suffix Array : sa[i] - the beginning position of
the ith smallest suffix, rk[i] - the rank of the
suffix beginning at position i. height[i] - the
longest common prefix of sa[i] and sa[i - 1]. */
template <int MAXN = 1000000, int MAXC = 26>
struct suffix_array {
int rk[MAXN], height[MAXN], sa[MAXN];
int cmp (int *x, int a, int b, int d) {
creturn x[a] == x[b] && x[a + d] == x[b + d]; }
void doubling (int *a, int n) {
static int sRank[MAXN], tmpA[MAXN], tmpB[MAXN];
int m = MAXC, *x = tmpA, *y = tmpB;
```

```
(int i = 0; i < m; ++i) sRank[i] = 0;
(int i = 0; i < n; ++i) ++sRank[x[i] = a[i]];
(int i = 1; i < m; ++i) sRank[i] += sRank[i -</pre>
11
12
     for (int i = n - 1; i >= 0; --i) sa[--sRank[x[i]]] =
     for (int i = 0; i < m; ++i) sRank[i] = 0;
for (int i = 0; i < n; ++i) ++sRank[x[i]];
for (int i = 1; i < m; ++i) sRank[i] += sRank[i -
1];
   height[rk[i]] = cur; } };
```

#### 7.2Suffix Automaton

```
/* Suffix automaton : head - the first state. tail -
the last state. Terminating states can be reached
via visiting the ancestors of tail. state::len -
the longest length of the string in the state.
state::right - 1 - the first location in the
string where the state can be reached. state::
parent - the parent link. state::dest - the
automaton link. */
template <int MAXN = 1000000, int MAXC = 26>
struct suffix_automaton {
struct state {
int len. right: state *parent. *dest[MAXC]:
          struct state {
  int len, right; state *parent, *dest[MAXC];
  state (int len = 0, int right = 0) : len (len),
      right (right), parent (NULL) {
    memset (dest, 0, sizeof (dest)); }
} node_pool[MAXN * 2], *tot_node, *null = new state()
         slse {
    state *q = p -> dest[token];
    if (p -> len + 1 == q -> len) {
        np -> parent = q;
    } else {
        state *nq = new (tot_node++) state (*q);
        nq -> len = p -> len + 1;
        np -> parent = q -> parent = nq;
    while (p && p -> dest[token] == q) {
        p -> dest[token] = nq, p = p -> parent;
    } } }
              tail = np == null ? np -> parent : np; }
          void init () {
  tot_node = node_pool;
  head = tail = new (to
          head = tail = new (tot_node++) state (); }
suffix_automaton () { init (); } };
```

## 7.3 Palindromic tree

```
1 /* Palindromic tree : extend () - returns whether the
tree has generated a new node. odd, even - the
root of two trees. last - the node representing
tree has generated a new node. odd, even - the
root of two trees. last - the node representing
the last char. node::len - the palindromic string
length of the node. */
template <int MAXN = 1000000, int MAXC = 26>
struct palindromic_tree {
struct node {
node *child[MAXC], *fail; int len;
node (int len): fail (NULL), len (len) {
memset (child, NULL, sizeof (child)); }
} node_pool[MAXN * 2], *tot_node;
int size, text(MAXN);
node *odd, *even, *last;
node *odd, *even, *last;
node *match (node *now) {
for (; text[size - now -> len - 1] != text[size];
now = now -> fail);
return now; }

bool extend (int token) {
text[++size] = token; node *now = match (last);
if (now -> child[token])
return last = now -> child[token], false;
last = now -> child[token] = new (tot_node++) node (
now -> len + 2);
if (now == odd) last -> fail = even;
else {
match (now -> fail):
12
                       else {
  now = match (now -> fail);
                       last -> fail =
return true; }
                                                                                                     = now -> child[token]; }
              return true; ;
void init() {
  text[size = 0] = -1; tot_node = node_pool;
  last = even = new (tot_node++) node (0); odd = new (
    tot_node++) node (-1);
  even -> fail = odd; }
palindromic_tree () { init (); } };
```

## 7.4 Regular expression

```
std::string str = ("The_the_there");
std::regex pattern ("(th|Th)[\\w]*",
s std::regex_match (str, match, pattern);
auto mbegin = std::sregex_iterator (str.begin (), str.
end (), pattern);
auto mend = std::sregex_iterator ();
std::cout << "Found_" << std::distance (mbegin, mend)
```

1 /\* Java reference : References on Java IO, structures,

### 8 Tips 8.1 Java

```
etc. */
import java.io.*;
import java.lang.*;
import java.math.*;
import java.util.*;
/* Common usage:
    | Stanfort Java.utll.*;
| A common usage:
| Scanner in = new Scanner (System.in);
| Scanner in = new Scanner (new BufferedInputStream (
| System.in));
| in.nextInt () / in.nextBigInteger () / in.
| nextBigDecimal () / in.nextDouble ()
| 9 in.nextInt () / in.nextBigInteger () / in.
| nextBigDecimal () / in.nextDouble ()
| 10 in.nextLine () / in.hasNext ()
| 11 System.out.print (...);
| 12 System.out.println (...);
| 13 System.out.printf (...);
| 14 BigInteger : BigInteger.valueOf (int) / abs / negate
| () / max / min / add / subtract / multiply /
| divide / remainder (BigInteger) / gcd (BigInteger)
| / modInverse (BigInteger mod) / modPow (
| BigInteger ex, BigInteger mod) / pow (int ex) /
| not () / and / or / xor (BigInteger) / shiftLeft /
| shiftRight (int) / compareTo (BigInteger) /
| intValue () / longValue () / toString (int radix) /
| / isProbablePrime (int certainty) /
| nextProbablePrime ()
| 15 BigDecimal : consists of a BigInteger value and a
| scale. The scale is the number of digits to the
| right of the decimal point.
| 16 divide (BigDecimal) : exact divide.
| 17 divide (BigDecimal) : int scale, RoundingMode |
| roundingMode) : divide with roundingMode, which
| may be: CEILING / DOWN / FLOOR / HALF_DOWN /
| HALF_EVEN / HALF_UP / UNNECESSARY / UP.
| 18 BigDecimal setScale (int newScale, RoundingMode roundingMode) : returns a BigDecimal with newScale
  19 doubleValue () / toPlainString () : converts to other
  types.

20 Arrays: Arrays.sort (T [] a); Arrays.sort (T [] a, int fromIndex, int toIndex); Arrays.sort (T [] a, int fromIndex, int toIndex, Comperator <? super T>
 comperator);
linkedList <E> : addFirst / addLast (E) / getFirst /
getLast / removeFirst / removeLast () / clear () /
add (int, E) / remove (int) / size () / contains
/ removeFirstOccurrence / removeLastOccurrence (E)
listIterator <E> listIterator (int index) : returns an
                                iterator
             E next / previous () : accesses and iterates.
hasNext / hasPrevious () : checks availablity.
nextIndex / previousIndex () : returns the index of a
x = 0;
y = 0;
y = 0; }
public Point (int xx, int yy) {
x = xx;
y = yy; } };
public static class Cmp implements Comparator <Point>
```

```
44
45
             Point> {
public static class Point im
Point> {
public int x; public int y;
public Point () {
    x = 0;
    y = 0; 1
           public static class Point implements Comparable <</pre>
              y = 0; }
public Point (int xx, int yy) {
x = xx;
   59
60
61
              x = xx;
y = yy; }
public int compareTo (Point p) {
if (x < p.x) return -1;
if (x == p.x) {
  if (y < p.y) return -1;
  if (y == p.y) return 0; }
return 1; }</pre>
              public boolean equalTo (Point p) {
  return (x == p.x && y == p.y); }
public int hashCode () {
  return x + y; } };
hasMoreTokens()) {
  try {
    String line = reader.readLine();
    tokenizer = new StringTokenizer (line);
    } catch (IOException e) {
      throw new RuntimeException (e); }
    return tokenizer.nextToken(); }
    public BigInteger nextBigInteger() {
      return new BigInteger (next (), 10); /* radix */ }
    public int nextInt() {
      return Integer.parseInt (next()); }
    public double nextDouble() {
      return Double.parseDouble (next()); }
    public static void main (String[] args) {
      InputReader in = new InputReader (System.in);
    }
}
```

#### 8.2 Random numbers

```
std::mt19937_64 mt (time (0));
std::uniform_int_distribution <int> uid (1, 100);
std::uniform_real_distribution <double> urd (1, 100);
std::cout << uid (mt) << "_" << urd (mt) << "\n";
```

#### 8.3 Read hack

#### 8.4 Stack hack

```
//C++
#pragma comment (linker, "/STACK:36777216")
//G++
int __size__ = 256 << 20;
char * _ p_ = (char*) malloc(__size__) + ___
              _size__ = 256 << 20;
*__p_ = (char*) malloc(__size__) + __size__;
__ ("movl_%0,_%%esp\n" :: "r"(__p__));
     asm
```

### Time hack

```
clock_t t = clock ();
std::cout << 1. * t / CLOCKS_PER_SEC << "\n";</pre>
```

#### 8.6 Multiplication hack

```
long long mul_mod (long long x, long long y, long long
mod) {
long long t = (x * y - (long long) ((long double) x /
mod * y + 1E-3) * mod) % mod;
return t < 0 ? t + mod : t; }</pre>
```

#### 8.7 Builtin functions

- \_builtin\_clz: Returns the number of leading 0-bits in x, starting at the most significant bit position. If x is 0, the result is
- undefined. \_\_builtin\_ctz: Returns the number of trailing 0-bits in x, starting at the least significant bit position. If x is 0, the result is
- undefined.
  \_builtin\_clrsb: Returns the number of leading redundant sign bits in x, i.e. the number of bits following the most significant bit that are identical to it. There are no special cases for 0 or
- other values.

  \_builtin\_popcount: Returns the number of 1-bits in x.
  \_builtin\_parity: Returns the parity of x, i.e. the number of
- 1-bits in x modulo 2. \_builtin\_bswap16, \_builtin\_bswap32, \_builtin\_bswap64: Returns x with the order of the bytes (8 bits as a group) reversed.
- bitset::Find\_first(), bitset::Find\_next(idx): bitset built-in functions.

## Prufer sequence

In combinatorial mathematics, the Prufer sequence of a labeled tree is a unique sequence associated with the tree. The sequence for a tree on n vertices has length n-2.

One can generate a labeled tree's Prufer sequence by iteratively removing vertices from the tree until only two vertices remain. Specifically, consider a labeled tree T with vertices 1, 2, ..., n. At step i, remove the leaf with the smallest label and set the *i*th element of the Prufer sequence to be the label of this leaf's neighbour.

One can generate a labeled tree from a sequence in three steps. The tree will have n+2 nodes, numbered from 1 to n+2. For each node set its degree to the number of times it appears in the sequence plus 1. Next, for each number in the sequence a[i], find the first (lowestnumbered) node, j, with degree equal to 1, add the edge (j, a[i]) to the tree, and decrement the degrees of j and a[i]. At the end of this loop two nodes with degree 1 will remain (call them u, v). Lastly, add the edge (u, v) to the tree.

The Prufer sequence of a labeled tree on n vertices is a unique sequence of length n-2 on the labels 1 to n - this much is clear. Somewhat less obvious is the fact that for a given sequence S of length n-2 on the labels 1 to n, there is a unique labeled tree whose Prufer sequence

## is S. **8.9** Spanning tree counting

Kirchhoff's Theorem: the number of spanning trees in a graph G is equal to *any* cofactor of the Laplacian matrix of G, which is equal to the difference between the graph's degree matrix (a diagonal matrix with vertex degrees on the diagonals) and its adjacency matrix (a (0,1)matrix with 1's at places corresponding to entries where the vertices are adjacent and 0's otherwise).

The number of edges with a certain weight in a minimum spanning tree is fixed given a graph. Moreover, the number of its arrangements can be obtained by finding a minimum spanning tree, compressing connected components of other edges in that tree into a point, and then applying Kirrchoff's theorem with only edges of the certain weight in the graph. Therefore, the number of minimum spanning trees in a graph can be solved by multiplying all numbers of arrangements of edges of different weights together.

#### 8.10Mobius inversion

#### 8.10.1 Mobius inversion formula

$$[x = 1] = \sum_{d|x} \mu(d)$$
$$x = \sum_{d|x} \mu(d)$$

#### 8.10.2 Gcd inversion

$$\begin{split} \sum_{a=1}^{n} \sum_{b=1}^{n} gcd^{2}(a,b) &= \sum_{d=1}^{n} d^{2} \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [gcd(i,j) = 1] \\ &= \sum_{d=1}^{n} d^{2} \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{t \mid gcd(i,j)} \mu(t) \\ &= \sum_{d=1}^{n} d^{2} \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} [t|i] \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [t|j] \\ &= \sum_{d=1}^{n} d^{2} \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \lfloor \frac{n}{dt} \rfloor^{2} \end{split}$$

The formula can be computed in O(nlogn) complexity. Moreover, let l = dt, then

$$\sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \lfloor \frac{n}{dt} \rfloor^2 = \sum_{l=1}^n \lfloor \frac{n}{l} \rfloor^2 \sum_{d|l} d^2 \mu(\frac{l}{d})$$

Let  $f(l) = \sum_{d|l} d^2 \mu(\frac{l}{d})$ . It can be proven that f(l) is multiplicative. Besides,  $f(p^k) = p^{2k} - p^{2k-2}$ .

Therefore, with linear sieve the formula can be computed in O(n)complexity.

## 8.11 Numbers

### 8.11.1 Binomial Coefficients

$$\binom{n}{k} = (-1)^k \binom{k-n-1}{k}, \quad \sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

$$\sum_{k=0}^n \binom{k}{m} = \binom{n+1}{m+1}$$

$$\sqrt{1+z} = 1 + \sum_{k=1}^\infty \frac{(-1)^{k-1}}{k \times 2^{2k-1}} \binom{2k-2}{k-1} z^k$$

$$\sum_{k=0}^r \binom{r-k}{m} \binom{s+k}{n} = \binom{r+s+1}{m+n+1}$$

$$C_{n,m} = \binom{n+m}{m} - \binom{n+m}{m-1}, n \ge m$$

$$\binom{n}{k} \equiv [n\&k = k] \pmod{2}$$

$$\binom{n_1 + \dots + n_p}{m} = \sum_{k_1 + \dots + k_p = m} \binom{n_1}{k_1} \dots \binom{n_p}{k_p}$$

## 8.11.2 Fibonacci Numbers

$$f(z) = \frac{1 - z - z^2}{1 - z - z^2}$$

$$f_n = \frac{\phi^n - \hat{\phi}^n}{\sqrt{5}}, \phi = \frac{1 + \sqrt{5}}{2}, \hat{\phi} = \frac{1 - \sqrt{5}}{2}$$

$$\sum_{k=1}^n f_k = f_{n+2} - 1, \quad \sum_{k=1}^n f_k^2 = f_n f_{n+1}$$

$$\sum_{k=0}^n f_k f_{n-k} = \frac{1}{5}(n-1)f_n + \frac{2}{5}nf_{n-1}$$

$$\frac{f_{2n}}{f_n} = f_{n-1} + f_{n+1}$$

$$f_1 + 2f_2 + 3f_3 + \dots + nf_n = nf_{n+2} - f_{n+3} + 2]$$

$$\gcd(f_m, f_n) = f_{\gcd(m, n)}$$

$$f_n^2 + (-1)^n = f_{n+1}f_{n-1}$$

$$f_{n+k} = f_n f_{k+1} + f_{n-1}f_k$$

$$f_{2n+1} = f_n^2 + f_{n+1}^2$$

$$(-1)^k f_{n-k} = f_n f_{k-1} - f_{n-1}f_k$$
Modulo  $f_n, f_{mn+r} \equiv \begin{cases} f_r, & m \bmod 4 = 0; \\ (-1)^{r+1} f_{n-r}, & m \bmod 4 = 2; \\ (-1)^{r+1+n} f_{n-r}, & m \bmod 4 = 3. \end{cases}$ 

### 8.11.3 Lucas Numbers

$$L_0 = 2, L_1 = 1, L_n = L_{n-1} + L_{n-2} = \left(\frac{1+\sqrt{5}}{2}\right)^n + \frac{1-\sqrt{5}}{2}^n$$

$$L(x) = \frac{2-x}{1-x-x^2}$$

### 8.11.4 Catlan Numbers

$$c_1 = 1, c_n = \sum_{i=0}^{n-1} c_i c_{n-1-i} = c_{n-1} \frac{4n-2}{n+1} = \frac{\binom{2n}{n}}{n+1}$$
$$= \binom{2n}{n} - \binom{2n}{n-1}, c(x) = \frac{1-\sqrt{1-4x}}{2x}$$

## 8.11.5 Stirling Cycle Numbers

Divide n elements into k non-empty cycles.

$$\begin{split} s(n,0) &= 0, s(n,n) = 1, s(n+1,k) = s(n,k-1) - ns(n,k) \\ & s(n,k) = (-1)^{n-k} {n \brack k} \\ & {n+1 \brack k} = n {n \brack k} + {n \brack k-1}, {n+1 \brack 2} = n! H_n \\ & x^{\underline{n}} = x(x-1)...(x-n+1) = \sum_k {n \brack k} (-1)^{n-k} x^k \\ & x^{\overline{n}} = x(x+1)...(x+n-1) = \sum_k {n \brack k} x^k \end{split}$$

### 8.11.6 Stirling Subset Numbers

Divide n elements into k non-empty subsets.

$${n+1 \choose k} = k {n \choose k} + {n \choose k-1}$$
$$x^n = \sum_k {n \choose k} x^{\underline{k}} = \sum_k {n \choose k} (-1)^{n-k} x^{\overline{k}}$$

$$m! {n \brace m} = \sum_{k} {m \choose k} k^n (-1)^{m-k}$$

For a fixed k, generating functions

$$\sum_{n=0}^{\infty} {n \brace k} x^{n-k} = \prod_{r=1}^k \frac{1}{1-rx}$$

### 8.11.7 Motzkin Numbers

Draw non-intersecting chords between n points on a circle.

Pick n numbers  $k_1,k_2,...,k_n\in\{-1,0,1\}$  so that  $\sum_i^a k_i (1\leq a\leq n)$  is non-negative and the sum of all numbers is 0.

$$M_{n+1} = M_n + \sum_{i=0}^{n-1} M_i M_{n-1-i} = \frac{(2n+3)M_n + 3nM_{n-1}}{n+3}$$

$$M_n = \sum_{i=0}^{\lfloor \frac{n}{2} \rfloor} {n \choose 2k} Catlan(k)$$

$$M(X) = \frac{1 - x - \sqrt{1 - 2x - 3x^2}}{2x^2}$$

## 8.11.8 Eulerian Numbers

Permutations of the numbers 1 to n in which exactly k elements are greater than the previous element.

## 8.11.9 Harmonic Numbers

Sum of the reciprocals of the first n natural numbers.

$$\sum_{k=1}^{n} H_k = (n+1)H_n - n$$

$$\sum_{k=1}^{n} kH_k = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}$$

$$\sum_{k=1}^{n} {n \choose k} H_k = {n+1 \choose m+1} (H_{n+1} - \frac{1}{m+1})$$

### 8.11.10 Pentagonal Number Theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = \sum_{n=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2}$$

$$p(n) = p(n-1) + p(n-2) - p(n-5) - p(n-7) + \cdots$$

$$f(n,k) = p(n) - p(n-k) - p(n-2k) + p(n-5k) + p(n-7k) - \cdots$$

## 8.11.11 Bell Numbers

Divide a set that has exactly n elements.

$$B_n = \sum_{k=1}^n {n \brace k}, \quad B_{n+1} = \sum_{k=0}^n {n \brack k} B_k$$

$$B_{p^m+n} \equiv mB_n + B_{n+1} \pmod{p}$$

$$B(x) = \sum_{n=0}^\infty \frac{B_n}{n!} x^n = e^{e^x - 1}$$

## 8.11.12 Bernoulli Numbers

$$B_n = 1 - \sum_{k=0}^{n-1} {n \choose k} \frac{B_k}{n-k+1}$$

$$G(x) = \sum_{k=0}^{\infty} \frac{B_k}{k!} x^k = \frac{1}{\sum_{k=0}^{\infty} \frac{x^k}{(k+1)!}}$$

$$\sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k n^{m-k+1}$$

## 8.11.13 Sum of Powers

$$\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}, \quad \sum_{i=1}^{n} i^3 = (\frac{n(n+1)}{2})^2$$

$$\sum_{i=1}^{n} i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

$$\sum_{i=1}^{n} i^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12}$$

## 8.11.14 Sum of Squares

Denote  $r_k(n)$  the ways to form n with k squares. If :

$$n = 2^{a_0} p_1^{2a_1} \cdots p_r^{2a_r} q_1 b_1 \cdots q_s b_s$$

where  $p_i \equiv 3 \mod 4$ ,  $q_i \equiv 1 \mod 4$ , then

$$r_2(n) = \begin{cases} 0 & \text{if any } a_i \text{ is a half-integer} \\ 4 \prod_{1}^{r} (b_i + 1) & \text{if all } a_i \text{ are integers} \end{cases}$$

 $r_3(n) > 0$  when and only when n is not  $4^a(8b+7)$ .

## 9 Appendix

## 9.1 Calculus table

**9.1.1** 
$$ax + b \ (a \neq 0)$$
  
1.  $\int \frac{x}{ax+b} dx = \frac{1}{a^2} (ax+b-b \ln |ax+b|) + C$ 

2. 
$$\int \frac{x^2}{ax+b} dx = \frac{1}{a^3} \left( \frac{1}{2} (ax+b)^2 - 2b(ax+b) + b^2 \ln|ax+b| \right) + C$$

3. 
$$\int \frac{\mathrm{d}x}{x(ax+b)} = -\frac{1}{b} \ln \left| \frac{ax+b}{x} \right| + C$$

4. 
$$\int \frac{\mathrm{d}x}{x^2(ax+b)} = -\frac{1}{bx} + \frac{a}{b^2} \ln \left| \frac{ax+b}{x} \right| + C$$

5. 
$$\int \frac{x}{(ax+b)^2} dx = \frac{1}{a^2} \left( \ln|ax+b| + \frac{b}{ax+b} \right) + C$$

6. 
$$\int \frac{x^2}{(ax+b)^2} dx = \frac{1}{a^3} \left( ax+b-2b \ln|ax+b| - \frac{b^2}{ax+b} \right) + C$$

## **9.1.2** $\sqrt{ax+b}$

1. 
$$\int \sqrt{ax+b} dx = \frac{2}{3a} \sqrt{(ax+b)^3} + C$$

2. 
$$\int x\sqrt{ax+b} dx = \frac{3a^{-4}}{15a^2} (3ax-2b)\sqrt{(ax+b)^3} + C$$

3. 
$$\int x^2 \sqrt{ax+b} dx = \frac{2}{105a^3} (15a^2x^2 - 12abx + 8b^2) \sqrt{(ax+b)^3} + C$$

4. 
$$\int \frac{x}{\sqrt{ax+b}} dx = \frac{2}{3a^2} (ax-2b) \sqrt{ax+b} + C$$

5. 
$$\int \frac{x^2}{\sqrt{ax+b}} dx = \frac{2}{15a^3} (3a^2x^2 - 4abx + 8b^2) \sqrt{ax+b} + C$$

2. 
$$\int x\sqrt{ax+b} dx = \frac{2}{15a^2} (3ax-2b)\sqrt{(ax+b)^3} + C$$
3. 
$$\int x^2\sqrt{ax+b} dx = \frac{2}{105a^3} (15a^2x^2 - 12abx + 8b^2)\sqrt{(ax+b)^3} + C$$
4. 
$$\int \frac{x}{\sqrt{ax+b}} dx = \frac{2}{3a^2} (ax-2b)\sqrt{ax+b} + C$$
5. 
$$\int \frac{2}{\sqrt{ax+b}} dx = \frac{2}{15a^3} (3a^2x^2 - 4abx + 8b^2)\sqrt{ax+b} + C$$
6. 
$$\int \frac{dx}{x\sqrt{ax+b}} = \begin{cases} \frac{1}{\sqrt{b}} \ln \left| \frac{\sqrt{ax+b} - \sqrt{b}}{\sqrt{ax+b} + \sqrt{b}} \right| + C \quad (b > 0) \\ \frac{2}{\sqrt{-b}} \arctan \sqrt{\frac{ax+b}{-b}} + C \quad (b < 0) \end{cases}$$
7. 
$$\int \frac{dx}{x^2\sqrt{ax+b}} = -\frac{\sqrt{ax+b}}{bx} - \frac{a}{2b} \int \frac{dx}{x\sqrt{ax+b}}$$
8. 
$$\int \frac{\sqrt{ax+b}}{\sqrt{ax+b}} dx = 2\sqrt{ax+b} + b \int \frac{dx}{x\sqrt{ax+b}}$$
9. 
$$\int \frac{\sqrt{ax+b}}{\sqrt{ax+b}} dx = -\frac{\sqrt{ax+b}}{\sqrt{ax+b}} + \frac{a}{2b} \int \frac{dx}{x\sqrt{ax+b}}$$

7. 
$$\int \frac{\mathrm{d}x}{-2\sqrt{x-1+b}} = -\frac{\sqrt{ax+b}}{bx} - \frac{a}{2b} \int \frac{\mathrm{d}x}{x\sqrt{ax+b}}$$

8. 
$$\int \frac{\sqrt{ax+b}}{x} dx = 2\sqrt{ax+b} + b \int \frac{dx}{x\sqrt{ax+b}}$$

9. 
$$\int \frac{\sqrt{ax+b}}{x^2} dx = -\frac{\sqrt{ax+b}}{x} + \frac{a}{2} \int \frac{dx}{x\sqrt{ax+b}}$$

## **9.1.3** $x^2 \pm a^2$

1. 
$$\int \frac{\mathrm{d}x}{x^2 + a^2} = \frac{1}{a} \arctan \frac{x}{a} + C$$

2. 
$$\int \frac{\mathrm{d}x}{(x^2+a^2)^n} = \frac{x}{2(n-1)a^2(x^2+a^2)^{n-1}} + \frac{2n-3}{2(n-1)a^2} \int \frac{\mathrm{d}x}{(x^2+a^2)^{n-1}}$$

3. 
$$\int \frac{\mathrm{d}x}{x^2 - a^2} = \frac{1}{2a} \ln \left| \frac{x - a}{x + a} \right| + C$$

9.1.4 
$$ax^2 + b (a > 0)$$
  
1.  $\int \frac{dx}{ax^2 + b} = \begin{cases} \frac{1}{\sqrt{ab}} \arctan \sqrt{\frac{a}{b}} x + C & (b > 0) \\ \frac{1}{2\sqrt{-ab}} \ln \left| \frac{\sqrt{ax} - \sqrt{-b}}{\sqrt{ax} + \sqrt{-b}} \right| + C & (b < 0) \end{cases}$   
2.  $\int \frac{x}{ax^2 + b} dx = \frac{1}{2a} \ln \left| ax^2 + b \right| + C$   
3.  $\int \frac{x^2}{ax^2 + b} dx = \frac{x}{a} - \frac{b}{a} \int \frac{dx}{ax^2 + b}$ 

2. 
$$\int \frac{x}{ax^2 + b} dx = \frac{1}{2a} \ln |ax^2 + b| + C$$

3. 
$$\int \frac{x^2}{ax^2 + b} dx = \frac{x}{a} - \frac{b}{a} \int \frac{dx}{ax^2 + b}$$

4. 
$$\int \frac{dx}{x(ax^2+b)} = \frac{1}{2b} \ln \frac{x^2}{|ax^2+b|} + C$$

$$\begin{array}{l} 4. \quad \int \frac{\mathrm{d}x}{x(ax^2+b)} = \frac{1}{2b} \ln \frac{x^2}{|ax^2+b|} + C \\ 5. \quad \int \frac{\mathrm{d}x}{x^2(ax^2+b)} = -\frac{1}{bx} - \frac{a}{b} \int \frac{\mathrm{d}x}{ax^2+b} \\ \end{array}$$

7. 
$$\int \frac{dx}{(ax^{2}+b)} = \frac{2b^{2}}{2b(ax^{2}+b)} + \frac{1}{2b} \int \frac{dx}{ax^{2}+b}$$
9.1.5 
$$ax^{2} + bx + c \quad (a > 0)$$
1. 
$$\frac{dx}{ax^{2}+bx+c} = \begin{cases} \frac{2}{\sqrt{4ac-b^{2}}} \arctan \frac{2ax+b}{\sqrt{4ac-b^{2}}} + C \quad (b^{2} < 4ac) \\ \frac{1}{\sqrt{b^{2}-4ac}} \ln \left| \frac{2ax+b-\sqrt{b^{2}-4ac}}{2ax+b+\sqrt{b^{2}-4ac}} \right| + C \quad (b^{2} > 4ac) \end{cases}$$
2. 
$$\int \frac{x}{ax^{2}+bx+c} dx = \frac{1}{2a} \ln |ax^{2} + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^{2}+bx+c}$$
0.1.6 
$$\int \frac{dx}{ax^{2}+bx+c} dx = \frac{1}{2a} \ln |ax^{2} + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^{2}+bx+c}$$

## **9.1.6** $\sqrt[a]{x^2 + a^2}$ (a > 0)

1. 
$$\int \frac{\mathrm{d}x}{\sqrt{x^2 + a^2}} = \operatorname{arsh} \frac{x}{a} + C_1 = \ln(x + \sqrt{x^2 + a^2}) + C$$

2. 
$$\int \frac{dx}{\sqrt{(x^2 + a^2)^3}} = \frac{x}{a^2 \sqrt{x^2 + a^2}} + C$$
3. 
$$\int \frac{x}{\sqrt{x^2 + a^2}} dx = \sqrt{x^2 + a^2} + C$$

3. 
$$\int \frac{x}{\sqrt{2a^2-a^2}} dx = \sqrt{x^2 + a^2} + C$$

4. 
$$\int \frac{x}{\sqrt{(-2+-2)^3}} dx = -\frac{1}{\sqrt{-2+-2}} + C$$

$$\int \frac{\sqrt{x^2 + a^2}}{\sqrt{(x^2 + a^2)^3}} dx = -\frac{x}{\sqrt{x^2 + a^2}} + \ln(x + \sqrt{x^2 + a^2}) + C$$

7. 
$$\int \frac{\mathrm{d}x}{x\sqrt{x^2+a^2}} = \frac{1}{a} \ln \frac{\sqrt{x^2+a^2}-a}{|x|} + C$$

9. 
$$\int \sqrt{x^2 + a^2} dx = \frac{x}{2} \sqrt{x^2 + a^2} + \frac{a^2}{2} \ln(x + \sqrt{x^2 + a^2}) + C$$

10. 
$$\int \sqrt{(x^2 + a^2)^3} \, dx = \frac{x}{8} (2x^2 + 5a^2) \sqrt{x^2 + a^2} + \frac{3}{8} a^4 \ln(x + \sqrt{x^2 + a^2}) + C$$

11. 
$$\int x\sqrt{x^2+a^2}dx = \frac{1}{3}\sqrt{(x^2+a^2)^3} + C$$

12. 
$$\int x^2 \sqrt{x^2 + a^2} \, dx = \frac{x}{8} (2x^2 + a^2) \sqrt{x^2 + a^2} - \frac{a^4}{8} \ln(x + \sqrt{x^2 + a^2}) + C$$

13. 
$$\int \frac{\sqrt{x^2 + a^2}}{x} dx = \sqrt{x^2 + a^2} + a \ln \frac{\sqrt{x^2 + a^2} - a}{|x|} + C$$

13. 
$$\int \frac{\sqrt{x^2 + a^2}}{x} dx = \sqrt{x^2 + a^2} + a \ln \frac{\sqrt{x^2 + a^2} - a}{|x|} + C$$
14. 
$$\int \frac{\sqrt{x^2 + a^2}}{x^2} dx = -\frac{\sqrt{x^2 + a^2}}{x} + \ln(x + \sqrt{x^2 + a^2}) + C$$

## **9.1.7** $\sqrt{x^2-a^2}$ (a>0)

$$1. \ \int \frac{\mathrm{d}x}{\sqrt{x^2 - a^2}} \, = \, \frac{x}{|x|} \mathrm{arch} \, \frac{|x|}{a} \, + \, C_1 \, = \, \ln \left| x + \sqrt{x^2 - a^2} \right| + C$$

3. 
$$\int \frac{x}{\sqrt{2-a^2}} dx = \sqrt{x^2 - a^2} + C$$

4. 
$$\int \frac{x}{\sqrt{(x^2-a^2)^3}} dx = -\frac{1}{\sqrt{x^2-a^2}} + C$$

4. 
$$\int \frac{x}{\sqrt{(x^2 - a^2)^3}} dx = -\frac{1}{\sqrt{x^2 - a^2}} + C$$
5. 
$$\int \frac{x^2}{\sqrt{x^2 - a^2}} dx = \frac{x}{2} \sqrt{x^2 - a^2} + \frac{a^2}{2} \ln|x + \sqrt{x^2 - a^2}| + C$$

6. 
$$\int \frac{x^2}{\sqrt{(x^2 - a^2)^3}} dx = -\frac{x}{\sqrt{x^2 - a^2}} + \ln|x + \sqrt{x^2 - a^2}| + C$$
7. 
$$\int \frac{dx}{x\sqrt{x^2 - a^2}} = \frac{1}{a} \arccos \frac{a}{|x|} + C$$

7. 
$$\int \frac{\mathrm{d}x}{\sqrt{-2}-2} = \frac{1}{a} \arccos \frac{a}{|x|} + C$$

9. 
$$\int \sqrt{x^2 - a^2} dx = \frac{x}{2} \sqrt{x^2 - a^2} - \frac{a^2}{2} \ln|x + \sqrt{x^2 - a^2}| + C$$

9. 
$$\int \sqrt{x^2 - a^2} dx = \frac{x}{2} \sqrt{x^2 - a^2} - \frac{a^2}{2} \ln|x + \sqrt{x^2 - a^2}| + C$$
10. 
$$\int \sqrt{(x^2 - a^2)^3} dx = \frac{x}{8} (2x^2 - 5a^2) \sqrt{x^2 - a^2} + \frac{3}{8} a^4 \ln|x + \sqrt{x^2 - a^2}| + C$$
11. 
$$\int x \sqrt{x^2 - a^2} dx = \frac{1}{3} \sqrt{(x^2 - a^2)^3} + C$$

11. 
$$\int x\sqrt{x^2-a^2} dx = \frac{1}{3}\sqrt{(x^2-a^2)^3} + C$$

12. 
$$\int x^2 \sqrt{x^2 - a^2} \, dx = \frac{x}{8} (2x^2 - a^2) \sqrt{x^2 - a^2} - \frac{a^4}{8} \ln|x + \sqrt{x^2 - a^2}| + C$$

13. 
$$\int \frac{\sqrt{x^2 - a^2}}{x} dx = \sqrt{x^2 - a^2} - a \arccos \frac{a}{|x|} + C$$

14. 
$$\int \frac{\sqrt{x^2 - a^2}}{x^2} dx = -\frac{\sqrt{x^2 - a^2}}{x} + \ln|x + \sqrt{x^2 - a^2}| + C$$

## **9.1.8** $\sqrt{a^2 - x^2} \ (a > 0)$

1. 
$$\int \frac{\mathrm{d}x}{\sqrt{a^2 - x^2}} = \arcsin \frac{x}{a} + C$$

2. 
$$\frac{dx}{\sqrt{(a^2-x^2)^3}} = \frac{x}{a^2\sqrt{a^2-x^2}} + C$$

2. 
$$\frac{dx}{\sqrt{(a^2 - x^2)^3}} = \frac{x}{a^2 \sqrt{a^2 - x^2}} + C$$
3. 
$$\int \frac{x}{\sqrt{a^2 - x^2}} dx = -\sqrt{a^2 - x^2} + C$$

4. 
$$\int \frac{x}{\sqrt{(-2^2-2)^3}} dx = \frac{1}{\sqrt{2^2-2}} + C$$

7. 
$$\int \frac{\mathrm{d}x}{x\sqrt{a^2 - x^2}} = \frac{1}{a} \ln \frac{a - \sqrt{a^2 - x^2}}{|x|} + C$$

8. 
$$\int \frac{\mathrm{d}x}{x^2 \sqrt{a^2 - x^2}} = -\frac{\sqrt{a^2 - x^2}}{a^2 x} + C$$

9. 
$$\int \sqrt{a^2 - x^2 dx} = \frac{x}{2} \sqrt{a^2 - x^2 + \frac{a^2}{2}} \arcsin \frac{x}{a} + C$$

9. 
$$\int \sqrt{a^2 - x^2} dx = \frac{x}{2} \sqrt{a^2 - x^2} + \frac{a^2}{2} \arcsin \frac{x}{a} + C$$
10. 
$$\int \sqrt{(a^2 - x^2)^3} dx = \frac{x}{8} (5a^2 - 2x^2) \sqrt{a^2 - x^2} + \frac{3}{8} a^4 \arcsin \frac{x}{a} + C$$
11. 
$$\int x \sqrt{a^2 - x^2} dx = -\frac{1}{3} \sqrt{(a^2 - x^2)^3} + C$$

11. 
$$\int x\sqrt{a^2 - x^2} dx = -\frac{1}{3}\sqrt{(a^2 - x^2)^3 + C}$$

13. 
$$\int \frac{\sqrt{a^2 - x^2}}{x} dx = \sqrt{a^2 - x^2} + a \ln \frac{a - \sqrt{a^2 - x^2}}{|x|} + C$$

13. 
$$\int \frac{1}{x} dx = \sqrt{a^2 - x^2 + a \ln \frac{1}{|x|}} + \frac{1}{|x|}$$

14. 
$$\int \frac{\sqrt{a^2 - x^2}}{x^2} dx = -\frac{\sqrt{a^2 - x^2}}{x} - \arcsin \frac{x}{a} + C$$

# **9.1.9** $\sqrt{\pm ax^2 + bx + c}$ (a > 0)

1. 
$$\int \frac{\mathrm{d}x}{\sqrt{ax^2 + bx + c}} = \frac{1}{\sqrt{a}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

2. 
$$\int \sqrt{ax^2 + bx + c} dx = \frac{2ax + b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac - b^2}{8\sqrt{a^3}} \ln|2ax| + b + \frac{a^2}{4a} + \frac{a^2}{4a}$$

$$2\sqrt{a}\sqrt{ax^2 + bx + c} + C$$

2. 
$$\int \sqrt{ax^2 + bx + c} \, dx = \frac{2ax + b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac - b^2}{8\sqrt{a^3}} \ln|2ax + b| + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$
3. 
$$\int \frac{x}{\sqrt{ax^2 + bx + c}} \, dx = \frac{1}{a} \sqrt{ax^2 + bx + c} - \frac{b}{2\sqrt{a^3}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

4. 
$$\int \frac{dx}{\sqrt{c + hx - ax^2}} = -\frac{1}{\sqrt{a}} \arcsin \frac{2ax - b}{\sqrt{h^2 + 4ac}} + C$$

5. 
$$\int \sqrt{c + bx - ax^2} dx = \frac{2ax - b}{4a} \sqrt{c + bx - ax^2} + \frac{b^2 + 4ac}{a} + \frac{2ax - b}{a} + \frac{c}{a}$$

$$\frac{5+7ac}{8\sqrt{a^3}}$$
 arcsin  $\frac{2ab}{\sqrt{b^2+4ac}}+C$ 

$$\sqrt{ax^{2}+bx+c} = 2\sqrt{a^{3}}$$

$$C$$
4. 
$$\int \frac{dx}{\sqrt{c+bx-ax^{2}}} = -\frac{1}{\sqrt{a}}\arcsin\frac{2ax-b}{\sqrt{b^{2}+4ac}} + C$$
5. 
$$\int \sqrt{c+bx-ax^{2}} dx = \frac{2ax-b}{4a} \sqrt{c+bx-ax^{2}} + \frac{b^{2}+4ac}{8\sqrt{a^{3}}}\arcsin\frac{2ax-b}{\sqrt{b^{2}+4ac}} + C$$
6. 
$$\int \frac{x}{\sqrt{c+bx-ax^{2}}} dx = -\frac{1}{a}\sqrt{c+bx-ax^{2}} + \frac{b}{2\sqrt{a^{3}}}\arcsin\frac{2ax-b}{\sqrt{b^{2}+4ac}} + C$$

# **9.1.10** $\sqrt{\pm \frac{x-a}{x-b}}$ & $\sqrt{(x-a)(x-b)}$

1. 
$$\int \sqrt{\frac{x-a}{x-b}} dx = (x-b)\sqrt{\frac{x-a}{x-b}} + (b-a)\ln(\sqrt{|x-a|} + \sqrt{|x-b|}) + C$$

2. 
$$\int \sqrt{\frac{x-a}{b-x}} \, \mathrm{d}x = (x-b)\sqrt{\frac{x-a}{b-x}} + (b-a)\arcsin\sqrt{\frac{x-a}{b-x}} + C$$

3. 
$$\int \frac{\mathrm{d}x}{\sqrt{(x-a)(b-x)}} = 2\arcsin\sqrt{\frac{x-a}{b-x}} + C \ (a < b)$$

4. 
$$\int \sqrt{(x-a)(b-x)} dx = \frac{2x-a-b}{4} \sqrt{(x-a)(b-x)} + \frac{(b-a)^2}{4} \arcsin \sqrt{\frac{x-a}{b-x}} + C$$

#### 9.1.11Triangular function

- $\int \tan x dx = -\ln|\cos x| + C$  $\int \cot x dx = \ln|\sin x| + C$
- 3.  $\int \sec x dx = \ln \left| \tan \left( \frac{\pi}{4} + \frac{x}{2} \right) \right| + C = \ln \left| \sec x + \tan x \right| + C$
- 4.  $\int \csc x \, dx = \ln \left| \tan \frac{x}{2} \right| + C = \ln \left| \csc x \cot x \right| + C$

- 5.  $\int \sec^2 x dx = \tan x + C$ 6.  $\int \csc^2 x dx = -\cot x + C$ 7.  $\int \sec x \tan x dx = \sec x + C$ 8.  $\int \csc x \cot x dx = -\csc x + C$
- 9.  $\int \sin^2 x \, dx = \frac{x}{2} \frac{1}{4} \sin 2x + C$
- 10.  $\int \cos^2 x \, dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$
- 10.  $\int \cos x dx = \frac{1}{2} + \frac{1}{4} \sin x + C$ 11.  $\int \sin^n x dx = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x dx$ 12.  $\int \cos^n x dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x dx$ 13.  $\frac{dx}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\sin^{n-2} x}$ 14.  $\frac{dx}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x}$ 15.

$$\int \cos^m x \sin^n x dx$$

$$= \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x dx$$

$$= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x dx$$

- 16.  $\int \sin ax \cos bx dx = -\frac{1}{2(a+b)} \cos(a+b)x \frac{1}{2(a-b)} \cos(a-b)x + C$ 17.  $\int \sin ax \sin bx dx = -\frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$
- 18.  $\int \cos ax \cos bx dx = \frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$

18. 
$$\int \cos ax \cos bx dx = \frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$$
19. 
$$\int \frac{dx}{a+b \sin x} = \begin{cases} \frac{2}{\sqrt{a^2-b^2}} \arctan \frac{a \tan \frac{x}{2}+b}{\sqrt{a^2-b^2}} + C & (a^2 > b^2) \\ \frac{1}{\sqrt{b^2-a^2}} \ln \left| \frac{a \tan \frac{x}{2}+b-\sqrt{b^2-a^2}}{a \tan \frac{x}{2}+b+\sqrt{b^2-a^2}} \right| + C & (a^2 < b^2) \end{cases}$$
20. 
$$\int \frac{dx}{a+b \cos x} = \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan \left(\sqrt{\frac{a-b}{a+b}} \tan \frac{x}{2}\right) + C & (a^2 > b^2) \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln \left| \frac{\tan \frac{x}{2}+\sqrt{\frac{a+b}{b-a}}}{\tan \frac{x}{2}-\sqrt{\frac{b+b}{b-a}}} \right| + C & (a^2 < b^2) \end{cases}$$
21. 
$$\int \frac{dx}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{ab} \arctan \left(\frac{b}{a} \tan x\right) + C$$
22. 
$$\int \frac{dx}{a^2 \cos^2 x - b^2 \sin^2 x} = \frac{1}{2ab} \ln \left| \frac{b \tan x + a}{b \tan x - a} \right| + C$$

- 22.  $\int \frac{\mathrm{d}x}{a^2\cos^2x b^2\sin^2x} = \frac{1}{2ab} \ln \left| \frac{b\tan x + a}{b\tan x a} \right| + C$

### 9.1.12 Inverse triangular function (a > 0)

- 1.  $\int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2 x^2} + C$

- 1.  $\int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2 x^2 + C}$ 2.  $\int x \arcsin \frac{x}{a} dx = (\frac{x^2}{2} \frac{a^2}{4}) \arcsin \frac{x}{a} + \frac{x}{4} \sqrt{x^2 x^2} + C$ 3.  $\int x^2 \arcsin \frac{x}{a} dx = \frac{x^3}{3} \arcsin \frac{x}{a} + \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 x^2} + C$ 4.  $\int \arccos \frac{x}{a} dx = x \arccos \frac{x}{a} \sqrt{a^2 x^2} + C$ 5.  $\int x \arccos \frac{x}{a} dx = (\frac{x^2}{2} \frac{a^2}{4}) \arccos \frac{x}{a} \frac{x}{4} \sqrt{a^2 x^2} + C$ 6.  $\int x^2 \arccos \frac{x}{a} dx = \frac{x^3}{3} \arccos \frac{x}{a} \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 x^2} + C$ 7.  $\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} \frac{a}{2} \ln(a^2 + x^2) + C$ 8.  $\int x \arctan \frac{x}{a} dx = \frac{1}{2} (a^2 + x^2) \arctan \frac{x}{a} \frac{a}{2} x + C$

- 9.  $\int x^2 \arctan \frac{x}{a} dx = \frac{x^3}{3} \arctan \frac{x}{a} \frac{a}{6}x^2 + \frac{a^3}{6} \ln(a^2 + x^2) + C$

## 9.1.13 Exponential function

- 1.13 Exponential

  1.  $\int a^x dx = \frac{1}{\ln a} a^x + C$ 2.  $\int e^{ax} dx = \frac{1}{a} a^{ax} + C$ 3.  $\int x e^{ax} dx = \frac{1}{a^2} (ax 1) a^{ax} + C$ 4.  $\int x^n e^{ax} dx = \frac{1}{a^2} x^n e^{ax} \frac{n}{a} \int x^{n-1} e^{ax} dx$ 5.  $\int x a^x dx = \frac{x}{\ln a} a^x \frac{1}{(\ln a)^2} a^x + C$
- 8.  $\int e^{ax} \cos bx dx = \frac{a + b}{a^2 + b^2} e^{ax} (b \sin bx + a \cos bx) + C$
- 9.  $\int e^{ax} \sin^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \sin^{n-1} bx (a \sin bx nb \cos bx) +$  $\frac{-a^{(n-1)b^2}}{a^2+b^2n^2} \int e^{ax} \sin^{n-2} bx dx$
- 10.  $\int e^{ax} \cos^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) +$  $\frac{a^2}{a^2+b^2n^2} \int e^{ax} \cos^{n-2} bx dx$

### 9.1.14 Logarithmic function

- 1.  $\int \ln x dx = x \ln x x + C$ 2.  $\int \frac{dx}{x \ln x} = \ln |\ln x| + C$

- $\begin{array}{l} x \ln x & x + \frac{1}{n+1}x^{n+1}(\ln x \frac{1}{n+1}) + C \\ 3. \int x^n \ln x dx = \frac{1}{n+1}x^{n+1}(\ln x)^n 1 dx \\ 4. \int (\ln x)^n dx = x(\ln x)^n n \int (\ln x)^{n-1} dx \\ 5. \int x^m (\ln x)^n dx = \frac{1}{m+1}x^{m+1}(\ln x)^n \frac{n}{m+1} \int x^m (\ln x)^{n-1} dx \end{array}$

### 9.2Regular expression

#### 9.2.1Special pattern characters

I	Parter characters
Characters	Description
	Not newline
\t	Tab (HT)
\n	Newline (LF)
\v	Vertical tab (VT)
\f	Form feed (FF)
\r	Carriage return (CR)
\cletter	Control code
\xhh	ASCII character
\uhhhh	Unicode character
\0	Null
\int	Backreference
\d	Digit
\D	Not digit
\s	Whitespace
\S	Not whitespace
\ W	Word (letters, numbers and the underscore)
\W	Not word
\character	Character
[class]	Character class
[^class]	Negated character class
•	

### Quantifiers

Characters	Times
*	0 or more
+	1 or more
?	0 or 1
{int}	int
{int,}	int or more
{min, max}	Between min and max

By default, all these quantifiers are greedy (i.e., they take as many characters that meet the condition as possible). This behavior can be overridden to ungreedy (i.e., take as few characters that meet the condition as possible) by adding a question mark (?) after the quantifier.

### 9.2.3 Groups

Characters	Description
(subpattern)	Group with backreference
(?:subpattern)	Group without backreference

#### 9.2.4Assertions

Characters	Description
^	Beginning of line
\$	End of line
\b	Word boundary
\B	Not a word boundary
(?=subpattern)	Positive lookahead
(?!subpattern)	Negative lookahead

#### Alternative 9.2.5

A regular expression can contain multiple alternative patterns simply by separating them with the separator operator (|): The regular expression will match if any of the alternatives match, and as soon as one does. 9.2.6 Character classes

Class	Description
[:alnum:]	Alpha-numerical character
[:alpha:]	Alphabetic character
[:blank:]	Blank character
[:cntrl:]	Control character
[:digit:]	Decimal digit character
[:graph:]	Character with graphical representation
[:lower:]	Lowercase letter
[:print:]	Printable character
[:punct:]	Punctuation mark character
[:space:]	Whitespace character
[:upper:]	Uppercase letter
[:xdigit:]	Hexadecimal digit character
[:d:]	Decimal digit character
[:W:]	Word character
[:s:]	Whitespace character

Please note that the brackets in the class names are additional to those opening and closing the class definition. For example:

[[:alpha:]] is a character class that matches any alphabetic char-

acter.
[abc[:digit:]] is a character class that matches a, b, c, or a digit.

[^[:space:]] is a character class that matches any character except a whitespace.