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Environment

1.1 Vimrc

```
set ru nu ts=4 sts=4 sw=4 si sm hls is ar bs=2 mouse=a syntax on nm <F3> :vsplit %<.in <CR>
nm <F4> :!gedit % <CR>
s au BufEnter *.cpp set cin
au BufEnter *.cpp nm <F5> :!time ./%< <CR>|nm <F7> :!
    gdb ./%< <CR>|nm <F8> :!time ./%< < %<.in <CR>|nm <F9> :!g++ % -o %< -g -std=gnu++14 -O2 -DLOCAL && size %< <CR>
```

Data Structure

2.1 KD tree

```
/* kd_tree : finds the k-th closest point in O(kn^{1-\frac{1}{k}}). Usage : Stores the data in p[]. Call function init (n, k). Call min_kth (d, k). (or max_kth) (k is 1-based)
  2 Usage
3 Note
                      Switch to the commented code for Manhattan
                distance
23 //
               idata[i], std::abs (dmax.data[i] ins.data[i]
]));
ret += 111 * tmp * tmp; }
  ret += std::max (std::abs (rhs.data[i] - dmax.data[i]) + std::abs (rhs.data[i] - dmin.data[i]));
       return ret; } tree[MAXN * 4];
struct result {
long long dist; point d; result() {}
result (const long long &dist, const point &d) :
    dist (dist), d (d) {}
bool operator > (const result &rhs) const { return
      35
38
           if ("tree[rt].r) tree[rt].merge (tree[tree[rt].r], k
      if ("tree[rt].r) tree[rt].merge (tree[tree[tree]...,
); }
std::priority_queue<result, std::vector<result>, std
::less <result>> heap_1;
std::priority_queue<result, std::vector<result>, std
::greater <result>> heap_r;
void _min_kth (const int &depth, const int &rt, const
    int &m, const point &d) {
  result tmp = result (sqrdist (tree[rt].p, d), tree[
    rt].p);
if ((int)heap_l.size() < m) heap_l.push (tmp);
else if (tmp < heap_l.top()) {
  heap_l.pop();
  heap_l.push (tmp); }</pre>
55
```

```
62
74
75
80
```

2.2Splay

```
void push_down (int x) {
  if (~n[x].c[0]) push (n[x].c[0], n[x].t);
  if (~n[x].c[1]) push (n[x].c[1], n[x].t);
  if (~n[x].t = tag (); )
  void update (int x) {
        \dot{m} = gen (x);
\dot{n}[x].c[0]) n[x].m = merge (n[n[x].c[0]].m, n[x].
  if ("n[x].c[1]) n[x].m = merge (n[x].m, n[n[x].c[1]].
m); }
```

2.3Link-cut tree

```
= u;

n[u].c[1] = v;

if (~v) n[v].f = u, n[v].p = -1;

update (u); u = n[v = u].p; }

splay (x); }
```

3 Formula

3.1 Zellers congruence

```
/* Zeller's congruence : converts between a calendar date and its Gregorian calendar day. (y >= 1) (0 = Monday, 1 = Tuesday, ..., 6 = Sunday) */
int get_id (int y, int m, int d) {
  if (m < 3) { --y; m += 12; }
  return 365 * y + y / 4 - y / 100 + y / 400 + (153 * ( m - 3) + 2) / 5 + d - 307; }
  std::tuple <int, int, int> date (int id) {
    int x = id + 1789995, n, i, j, y, m, d;
    n = 4 * x / 146097; x -= (146097 * n + 3) / 4;
    i = (4000 * (x + 1)) / 1461001; x -= 1461 * i / 4 - 31;
    j = 80 * x / 2447; d = x - 2447 * j / 80;
    x = j / 11;
    m = j + 2 - 12 * x; y = 100 * (n - 49) + i + x;
    return std::make_tuple (y, m, d); }
```

3.2 Lattice points below segment

```
/* Euclidean-like algorithm : computes the sum of
         \sum_{i=0}^{n-1} \left[ \frac{a+bi}{m} \right] \cdot \star /
long long solve(long long n, long long a, long long b,
long long m) {
   if (b == 0) return n * (a / m);
```

3.3 Adaptive Simpson's method

```
/* Adaptive Simpson's method : integrates f in [1, r].

*/
struct simpson {
    double area (double (*f) (double), double 1, double r
    ) {
        double m = 1 + (r - 1) / 2;
        return (f (1) + 4 * f (m) + f (r)) * (r - 1) / 6; }

    double solve (double (*f) (double), double 1, double r, double solve (double a) {
        double m = 1 + (r - 1) / 2;
        double left = area (f, 1, m), right = area (f, m, r)

        if (fabs (left + right - a) <= 15 * eps) return left + right + (left + right - a) / 15.0;
        return solve (f, 1, m, eps / 2, left) + solve (f, m, r, eps / 2, right); }

        double solve (double (*f) (double), double 1, double r, double eps) {
            return solve (f, 1, r, eps, area (f, 1, r)); } };
```

3.4 Neural network

```
1 /* Neural network : ft features, n layers, m neurons
       12
       for (int i = 0; i < rt + i; ++i; avg[i] - Sig[i] -
0; }
double compute (double *x) {
  for (int j = 0; j < m; ++j) {
    val[0][j] = 0; for (int k = 0; k < ft; ++k) val[0][
    j] += wp[0][j][k] * x[k];
    val[0][j] = 1 / (1 + exp (-val[0][j])); }
  for (int i = 1; i < n; ++i) for (int j = 0; j < m;
    i+i )</pre>
      double
17
    18
       for (int j = 0; j < m; ++j) del[n - 1][j] = w[j] *
   delo * val[n - 1][j] * (1 - val[n - 1][j]);
for (int i = n - 2; i >= 0; --i) for (int j = 0; j <</pre>
25
    28
32 //
33
35
                 in; ++j)
i] += (data[j][i] - avg[i]) * (data[j][i] - avg
    / sig[l];
return compute (x) * sig[ft] + avg[ft]; }
std::string to_string () {
    std::ostringstream os; os << std::fixed << std::
        setprecision (16);
    for (int i = 0; i < n; ++i) for (int j = 0; j < m;
        ++j) for (int k = 0; k < (i ? m : ft); ++k)
    os << wp[i][j][k] << "'";
    for (int i = 0; i < m; ++i) os << w[i] << """;
    for (int i = 0; i < ft + 1; ++i) os << avg[i] << """;
}</pre>
```

4 Number theory

4.1 Fast power module

```
/* Fast power module : x<sup>n</sup> */
2 int fpm (int x, int n, int mod) {
3 int ans = 1, mul = x; while (n) {
4 if (n & 1) ans = int (111 * ans * mul * mod);
5 mul = int (111 * mul * mod); n >>= 1; }
7 long long mul_mod (long long x, long long y, long long mod) {
8 long long t = (x * y - (long long) ((long double) x / mod * y + 1E-3) * mod) * mod;
9 return t < 0 ? t + mod : t; }
10 long long ans = 1, mul = x; while (n) {
11 long long ans = 1, mul = x; while (n) {
12 if (n & 1) ans = mul_mod (ans, mul, mod);
13 mul = mul_mod (mul, mul, mod); n >>= 1; }
14 return ans; }
```

4.2 Euclidean algorithm

```
/* Euclidean algorithm : solves for ax + by = gcd (a, b). */
void euclid (const long long &a, const long long &b, long long &x, long long &y) {
   if (b == 0) x = 1, y = 0;
   else euclid (b, a % b, y, x), y -= a / b * x; }
   if long long inverse (long long x, long long m) {
    long long a, b; euclid (x, m, a, b); return (a % m + m) % m; }
```

4.3 Discrete Fourier transform

4.4 Fast Walsh-Hadamard transform

4.5 Number theoretic transform

```
/* Number theoretic transform : NTT for any module.
Usage : Perform NTT on 3 modules and call crt () to
merge the result. */
template <int MAXN = 1000000>
struct ntt {
```

4.6 Polynomial operation

```
| template <int MAXN = 1000000>
 Note: n must be a power of 2. 2x max length. */
void inverse (int *a, int *b, int n, int mod, int prt
   11
12
13
    22
                                     0, mod, prt); tr.solve (b, m +
25
26
          with deg(d) \leq n-m and deg(r) < m. 4x max length required. */
    when aeg(a) \ge n-m and aeg(r) < m. 4x max length required. */
void divide (int *a, int n, int *b, int m, int *d, int *r, int mod, int prt) {
static int u[MAXN], v[MAXN]; while (!b[m - 1]) --m; int p = 1, t = n - m + 1; while (p < t << 1) p <<= 1.1...
     31
      prt);
for (int i = 0; i < p; ++i) u[i] = 1LL * u[i] * v[i]
              % mod:
     tr.solve (u, p, 1, mod, prt); std::reverse (u, u + t
   ); std::copy (u, u + t, d);
for (p = 1; p < n; p <<= 1); std::fill (u + t, u + p
   , 0);
tr.solve (u, p, 0, mod, prt); std::copy (b, b + m, v
   );</pre>
      std::fill (v + m, v + p, 0); tr.solve (v, p, 0, mod,
      prt);
for (int i = 0; i < p; ++i) u[i] = 1LL * u[i] * v[i]</pre>
             % mod;
     % mod;
tr.solve (u, p, 1, mod, prt);
for (int i = 0; i < m; ++i) r[i] = (a[i] - u[i] +
    mod) % mod;
std::fill (r + m, r + p, 0); } };
```

4.7 Chinese remainder theorem

```
/* Chinese remainder theroem: finds positive integers
    x = out.first + k * out.second that satisfies x %
    in[i].second = in[i].first. */

struct crt {

long long fix (const long long &a, const long long &b) { return (a % b + b) % b; }

bool solve (const std::vector <std::pair <long long, long long> &in, std::pair <long long, long long> &out) {

out = std::make_pair (1LL, 1LL);

for (int i = 0; i < (int) in.size (); ++i) {

long long n, u;

euclid (out.second, in[i].second, n, u);

long long divisor = std::_gcd (out.second, in[i]. second);

if ((in[i].first - out.first) % divisor) return false;

n *= (in[i].first - out.first) / divisor;

n = fix (n, in[i].second);

out.first += out.second * n;

out.second *= in[i].second / divisor;

out.first = fix (out.first, out.second);

return true; } ;
```

4.8 Linear Recurrence

4.10 Baby step giant step algorithm

```
| /* Baby step giant step algorithm : Solves a^x = b \mod c in O(\sqrt{c}). */
| struct bsgs {
| int solve (int a, int b, int c) {
| std::unordered_map <int, int> bs;
| int m = (int) sqrt ((double) c) + 1, res = 1;
| for (int i = 0; i < m; ++i) {
| if (bs.find (res) == bs.end ()) bs[res] = i;
| res = int (1LL * res * a % c); }
| int mul = 1, inv = (int) inverse (a, c);
| for (int i = 0; i < m; ++i) mul = int (1LL * mul * inv % c);
| res = b % c;
| for (int i = 0; i < m; ++i) {
| if (bs.find (res) != bs.end ()) return i * m + bs[ res];
| res = int (1LL * res * mul % c); }
| return -1; };
```

4.11 Pell equation

4.12 Quadric residue

```
/* Quadric residue: finds solution for x^2 = n \mod p (0 \le a < p) with prime p in O(\log p) complexity. */

struct quadric {

void multiply(long long &c, long long &d, long long a , long long b, long long w, long long p) {

int cc = (a * c + b * d % p * w) % p;

int dd = (a * d + b * c) % p; c = cc, d = dd; }

bool solve(int n, int p, int &x) {

if (n == 0) return x = 0, true; if (p == 2) return x

= 1, true;

if (power (n, p / 2, p) == p - 1) return false;

long long c = 1, d = 0, b = 1, a, w;

do { a = rand() % p; w = (a * a - n + p) % p;

if (w == 0) return x = a, true;

} while (power (w, p / 2, p) != p - 1);

for (int times = (p + 1) / 2; times; times >>= 1) {

if (times & 1) multiply (c, d, a, b, w, p);

multiply (a, b, a, b, w, p);

return x = c, true; };
```

4.13 Miller Rabin primality test

4.14 Pollard's Rho algorithm

```
1 /* Pollard's Rho : factorizes an integer. */
2 struct pollard_rho {
3   miller_rabin is_prime;
4   const long long thr = 13E9;
5   long long facize (const long long &n, const long long &seed) {
```

${f 5}$ Geometry

```
#define cd const double &
const double EPS = 1E-8, PI = acos (-1);
int sgn (cd x) { return x < -EPS ? -1 : x > EPS; }
int cmp (cd x, cd y) { return sgn (x - y); }
double sqr (cd x) { return x * x; }
double msqrt (cd x) { return sgn (x) <= 0 ? 0 : sqrt (x); }</pre>
```

5.1 Point

5.2 Line

```
18| return std::min (dis (a, b.s), dis (a, b.t)); }
19| bool in_polygon (cp p, const std::vector <point> & po)
          int n = (int) po.size (), counter = 0;
for (int i = 0; i < n; ++i) {
  point a = po[i], b = po[(i + 1) % n];
  //Modify the next line if necessary.
  if (point_on_segment (p, line (a, b))) return true;
  int x = sgn (det (p - a, b - a)), y = sgn (a.y - p.
      ), z = sgn (b.y - p.y);
  if (x > 0 && y <= 0 && z > 0) counter++;
  if (x < 0 && z <= 0 && y > 0 counter--; }
  return counter != 0; }
double polygon area (const std: vector point > &a) {
```

5.3 Circle

```
if (cmp (point_to_line (b.c, a), b.r) > 0) return std
    ::vector <point> ();
double x = msqrt (sqr (b.r) - sqr (point_to_line (b.c
    , a)));
point s = project_to line (b.c.)
              point s = project_to_line (b.c, a), u = (a.t - a.s).
    unit ();
if (sgn (x) == 0) return std::vector <point> ({s});
return std::vector <point> ({s - u * x, s + u * x});
28
 29
r * x});

return std::vector <point> ({a.c + r * x - r.rot90 () * h, a.c + r * x + r.rot90 () * h}); }

//Counter-clockwise with respect of point a.

std::vector <point> tangent (cp a, cc b) { circle p = make_circle (a, b.c); return circle_intersect (p, b); }
make_circle (a, b.c); return circle_intersect (p, b); }

33 //Counter-clockwise with respect of point Oa.

34 std::vector <line> extangent (cc a, cc b) {

35 std::vector <line> ret;

36 if (cmp (dis (a.c, b.c), std::abs (a.r - b.r)) <= 0) return ret;

37 if (sgn (a.r - b.r) == 0) {

38 point dir = b.c - a.c; dir = (dir * a.r / dis (dir)) ...

39 ret.push_back (line (a.c - dir, b.c - dir));

40 ret.push_back (line (a.c + dir, b.c + dir));

41 } else {

42 point p = (b.c * a.r - a.c * b.r) / (a.r - b.r);

53 std::vector <point> pp = tangent (p, a), qq = tangent (p, b);

44 if (pp.size () == 2 && qq.size () == 2) {

45 if (cmp (a.r, b.r) < 0) std::swap (pp[0], pp[1]),

46 std::swap (qq[0], qq[1]);

77 ret.push_back (line (pp[0], qq[0]));

88 return ret; }

48 return ret; }
ret.push_back (line (pp[1], qq[1])); } }

return ret; }

// (Counter-clockwise with respect of point Oa.

// (Std::vector <line> intangent (cc c1, cc c2) {

// (Std::vector <line> ret;

// (point p = (b.c * a.r + a.c * b.r) / (a.r + b.r);

// (for the counter point) pp = tangent (p, a), qq = tangent

// (p, b);

// (pp.size () == 2 && qq.size () == 2) {

// (std) fine (pp[0], qq[0]));

// (ret.push_back (line (pp[0], qq[0]));

// (ret.push_back (line (pp[1], qq[1]));

// (ret.push_back (line (pp[1], qq[1]));

// (ret.push_back (line (pp[1], qq[1]));

// (ret.push_back (line (pp[1], qq[1]));
```

5.4 Centers of a triangle

```
point incenter (cp a, cp b, cp c) {
  double p = dis (a, b) + dis (b, c) + dis (c, a);
  return (a * dis (b, c) + b * dis (c, a) + c * dis (a, b)) / p; }
```

5.5 Fermat point

```
/* Fermat point : finds a point P that minimizes |PA| + |PB| + |PC|. */
2 point fermat point (cp a, cp b, cp c) {
3 if (a == b) return a; if (b == c) return b; if (c == a) return c;
4 double ab = dis (a, b), bc = dis (b, c), ca = dis (c, a);
                double ab = dis (a, b), bc = dis (b, c), ca = dis (c, a);
double cosa = dot (b - a, c - a) / ab / ca;
double cosb = dot (a - b, c - b) / ab / bc;
double cosc = dot (b - c, a - c) / ca / bc;
double cosc = dot (b - c, a - c) / ca / bc;
double sq3 = PI / 3.0; point mid;
if (sqn (cosa + 0.5) < 0) mid = a;
else if (sqn (cosb + 0.5) < 0) mid = b;
else if (sqn (cosc + 0.5) < 0) mid = c;
else if (sqn (det (b - a, c - a)) < 0) mid =
    line_intersect (line (a, b + (c - b).rot (sq3)),
    line (b, c + (a - c).rot (sq3)));
else mid = line_intersect (line (a, c + (b - c).rot (sq3)));
return mid; }</pre>
```

5.6 Convex hull

5.7 Half plane intersection

```
1/* Online half plane intersection : complexity O(n)
    std::vector <point> ret;
if (c.empty ()) return ret;
for (int i = 0; i < (int) c.size (); ++i) {
  int j = (i + 1) % (int) c.size ();
  if (turn_left (p.s, p.t, c[i])) ret.push_back (c[i])</pre>
   if (two_side (c[i], c[j], p)) ret.push_back (
    line_intersect (p, line (c[i], c[j]))); }
return ret; }
/* Offline half plane intersection : complexity
\begin{array}{c} O(n\log n). \ */ \\ \text{11} \ bool \ turn \ left} \ (\text{cl 1, cp p}) \ \{ \ \text{return turn left} \ (\text{l.s, l. in cmp (cp a, cp b)} \ \{ \ \text{return a.dim () != b.dim () ? (a.dim () < b.dim () ? -1 : 1) : -sgn (det (a, b));} \end{array}
   else return cmp (a.first, b.first) < 0; });
h.resize (std::unique (g.begin (), g.end (), [] (
    const polar &a, const polar &b) { return cmp (a.
    first, b.first) == 0 }) - g.begin ());
for (int i = 0; i < (int) h.size (); ++i) h[i] = g[i</pre>
    22
      fore;
    if (rear - fore < 2) return std::vector <point> ();
std::vector <point> ans; ans.resize (rear + 1);
```

```
for (int i = 0; i < rear + 1; ++i) ans[i] =
    line_intersect (ret[i], ret[(i + 1) % (rear + 1)</pre>
 return ans:
```

5.8 Nearest pair of points

```
1 /* Nearest pair of points : [1, r), need to sort p
first. */
2 double solve (std::vector <point> &p, int 1, int r) {
3 if (1 + 1 >= r) return INF;
4 int m = (1 + r) / 2; double mx = p[m].x; std::vector
     for (int i = 1; i < r; ++i)
if (sqr (p[i].x - mx) < ret) v.push_back (p[i]);
sort (v.begin (), v.end (), [&] (cp a, cp b) { ret
a.y < b.y; } );
for (int i = 0; i < v.size (); ++i)
for (int j = i + 1; j < v.size (); ++j) {
   if (sqr (v[i].y - v[j].y) > ret) break;
   ret = min (ret, dis2 (v[i] - v[j])); }
return ret; }
```

5.9 Minimum circle

```
circle minimum_circle (std::vector <point> p) {
circle ret; std::random_shuffle (p.begin (), p.end ()
      for (int i = 0; i < (int) p.size (); ++i) if (!
    in_circle (p[i], ret)) {
    ret = circle (p[i], 0); for (int j = 0; j < i; ++j)
        if (!in_circle (p[j], ret)) {
        ret = make_circle (p[j], p[i]); for (int k = 0; k < i; ++k)</pre>
    for (int i = 0; i
```

5.10Intersection of a polygon and a circle

```
struct polygon_circle_intersect {
  double sector_area (cp a, cp b, const double &r) {
   double c = (2.0 * r * r - dis2 (a, b)) / (2.0 * r *
        r);
   roughly const double on the const double on 
                                            r);
return r * r * acos (c) / 2.0; }
double area (cp a, cp b, const double &r) {
    double dA = dot (a, a), dB = dot (b, b), dC =
        point_to_segment (point (), line (a, b));
    if (sgn (dA - r * r) <= 0 && sgn (dB - r * r) <= 0)
        return det (a, b) / 2.0;
point tA = a.unit () * r, tB = b.unit () * r;
    if (sgn (dC - r) > 0) return sector_area (tA, tB, r)
```

Union of circles 5.11

```
| template <int MAXN = 500> struct union_
                                                                                 circle {
   template <int MAXN = 500> struct union_circle {
  int C; circle c[MAXN]; double area[MAXN];
  struct event {
  point p; double ang; int delta;
  event (cp p = point (), double ang = 0, int delta =
      0) : p(p), ang(ang), delta(delta) {}
  bool operator < (const event &a) { return ang < a.
      ang; }
}</pre>
    };
void addevent(cc a, cc b, std::vector <event> &evt,
```

```
for (int j = 0; j < C; ++j) if (j != i && !same (c[
    i], c[j]) && overlap (c[j], c[i])) ++cnt;
for (int j = 0; j < C; ++j) if (j != i && !overlap
    (c[j], c[i]) && !overlap (c[i], c[j]) &&
    intersect (c[i], c[j]))
addevent (c[i], c[j], evt, cnt);
if (evt.empty ()) area[cnt] += PI * c[i].r * c[i].r</pre>
23
                31
```

5.123D point

```
1);
for (int i = 0; i < 4; ++i) for (int j = 0; j < 4; ++
29
   ans[i] += a[j][i] * c[j];
return point3 (ans[0], ans[1], ans[2]);
```

5.133D line

5.14 3D convex hull

```
1 /* 3D convex hull : initializes n and p / outputs face
template <int MAXN = 500>
struct convex hull3 {
   double mix (cp3 a, cp3 b, cp3 c) { return dot (det (a b), c); }
   double volume (cp3 a, cp3 b, cp3 c, cp3 d) { return mix (b - a, c - a, d - a); }
   struct tri {
}
```

6 Graph

```
template <int MAXN = 100000, int MAXM = 100000>
struct edge_list {
  int size, begin[MAXN], dest[MAXM], next[MAXM];
  void clear (int n) { size = 0; std::fill (begin, begin + n, -1); }
  edge_list (int n = MAXN) { clear (n); }
  void add_edge (int u, int v) { dest[size] = v; next[ size] = begin[u]; begin[u] = size++; } };
  template <int MAXN = 100000, int MAXM = 100000>
  struct cost_edge_list {
   int size, begin[MAXN], dest[MAXM], next[MAXM], cost[ MAXM];
  void clear (int n) { size = 0; std::fill (begin, begin + n, -1); }
  cost_edge_list (int n = MAXN) { clear (n); }
  void add_edge (int u, int v, int c) { dest[size] = v; next[size] = begin[u]; cost[size] = c; begin[u] = size++; };
}
```

6.1 Hopcoft-Karp algorithm

```
| /* Hopcoft-Karp algorithm : unweighted maximum matching for bipartition graphs with complexity O(m\sqrt{n}). */

| template < int MAXN = 100000, int MAXM = 100000>
| struct hopcoft_karp {
| using edge_list = std::vector < int> [MAXN];
| int mx[MAXN], my[MAXM], lv[MAXN];
| bool dfs (edge_list < MAXN, MAXM> &e, int x) {
| for (int i = e.begin[x]; ~i; i = e.next[i]) {
| int y = e.dest[i], w = my[y];
| if (!~w || (lv[x] + 1 == lv[w] && dfs (e, w))) {
| mx[x] = y; my[y] = x; return true; } {
| lv[x] = -1; return false; } {
| int solve (edge_list < MAXN, MAXM> &e, int n, int m) {
| std::fill (mx, mx + n, -1); std::fill (my, my + m, -1); {
| for (int ans = 0; ;) {
| std::vector < int> q; {
| for (int i = 0; i < n; ++i) |
| if (mx[i] = -1); {
| for (int head = 0; head < (int) q.size(); ++head) {
| int x = q[head]; for (int i = e.begin[x]; ~i; i = e.next[i]) {
| int y = e.dest[i], w = my[y]; |
| if (~w && lv[w] < 0) { lv[w] = lv[x] + 1; q. push_back (w); } {
| using y = e.dest[i], w = my[y]; |
| if (~w && lv[w] < 0) { lv[w] = lv[x] + 1; q. push_back (w); } {
| int d = 0; for (int i = 0; i < n; ++i) if (!~mx[i] && dfs (e, i)) ++d; |
| if (d == 0) return ans; else ans += d; } };
```

6.2 Kuhn-Munkres algorithm

```
/* Kuhn Munkres algorithm : weighted maximum matching
    on bipartition graphs.
lote : the graph is 1-based. */
stemplate <int MAXN = 500>
struct kuhn_munkres {
    int n, w[MAXN][MAXN], lx[MAXN], ly[MAXN], m[MAXN],
    wav[MAXN];

bool u[MAXN];

void hungary(int x) {
    m[0] = x; int j0 = 0;
    std::fill (sl, sl + n + 1, INF); std::fill (u, u + n + 1, false);

do {
    u[j0] = true; int i0 = m[j0], d = INF, j1 = 0;
    for (int j = 1; j <= n; ++j)
    if (u[j] == false) {
        int cur = -w[i0][j] - lx[i0] - ly[j];
        if (cur < sl[j]) { sl[j] = cur; way[j] = j0; }
    if (sl[j] < d) { d = sl[j]; j1 = j; } }
    for (int j = 0; j <= n; ++j) {
        if (u[j]) { lx[m[j]] += d; ly[j] -= d; }
        if (u[j]) { lx[m[j]] += d; ly[j] -= d; }
        if (u[j]) { lx[m[j]] != 0);

do {
    int j1 = way[j0]; m[j0] = m[j1]; j0 = j1;
}</pre>
```

6.3 Blossom algorithm

```
1 /* Blossom algorithm : maximum match for general graph
 template <int MAXN = 500, int MAXM = 250000>
3 struct blossom {
     12
    33
          int dest = e.dest[1];
if (match[dest] == -2 || ufs.find (loc) == ufs.
    find (dest)) continue;
if (d[dest] == -1)
if (match[dest] == -1) {
    solve (root, loc); match[loc] = dest;
    match[dest] = loc; return 1;
}
    } else {
    fa[dest] = loc; fa[match[dest]] = dest;
    d[dest] = 1; d[match[dest]] = 0;
    *qtail++ = match[dest];
} else if (d[ufs.find (dest)] == 0) {
    int b = lca (loc, dest, root);
    contract (loc, dest, b); contract (dest, loc, b)
    return 0;
}
return 0;
}
int solve (int n, const edge_list <MAXN, MAXM> &e) {
    std::fill (fa, fa + n, 0); std::fill (c1, c1 + n, 0)
                else
47
51
       std: fill (c2, c2 + n, 0); std::fill (match, match +
       std:::::: (B2 + n, 0); std:::::: (match, match
    n, -1);
int re = 0; for (int i = 0; i < n; i++)
if (match[i] == -1) if (bfs (i, n, e)) ++re; else
    match[i] = -2;
return re; };</pre>
```

6.4 Weighted blossom algorithm

```
/* Weighted blossom algorithm (vfleaking ver.):
    maximum matching for general weighted graphs with complexity O(n³).

2 Usage: Set n to the size of the vertices. Run init ()
    . Set g[][].w to the weight of the edge. Run solve ().

3 The first result is the answer, the second one is the number of matching pairs. Obtain the matching with match[].

4 Note: 1-based. */
    struct weighted blossom {
    static const int INF = INT_MAX, MAXN = 400;
    struct edge{ int u, v, w; edge (int u = 0, int v = 0, int w = 0): u(u), v(v), w(w) { } ;
    int n, n_x;
    edge g[MAXN * 2 + 1][MAXN * 2 + 1];
    int lab[MAXN * 2 + 1], match[MAXN * 2 + 1], slack[MAXN * 2 + 1], int flower_from[MAXN * 2 + 1][MAXN * 1], S[MAXN * 2 + 1];
    int flower_from[MAXN * 2 + 1];
    std::vector <int> flower[MAXN * 2 + 1]; std::queue < int> q;
    int e_delta (const edge &e) { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
```

```
b, xr);
for (int i = 0; i
      for (int i = 0; i < pr; i += 2) {
  int xs = flower[b][i], xns = flower[b][i + 1];
  pa[xs] = g[xns][xs].u; S[xs] = 1, S[xns] = 0;
  slack[xs] = 0, set_slack(xns); q_push(xns); }
  S[xr] = 1, pa[xr] = pa[b];
  for (size_t i = pr + 1; i < flower[b].size (); ++i)</pre>
     int xs = flower[b][i]; S[xs] = -1, set_slack(xs); }
st[b] = 0; }
bool on_found_edge (const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1; int nu = st[match[v]];
    slack[v] = slack[nu] = 0; S[nu] = 0, q_push(nu);
} else if(S[v] == 0) {
  int lca = get_lca(u, v);
  if (!lca) return augment(u, v), augment(v, u), true
```

```
return false; }
std::pair <long long, int> solve () {
   memset (match + 1, 0, sizeof (int) * n); n_x = n;
   int n_matches = 0; long long tot_weight = 0;
   for (int u = 0; u <= n; ++u) st[u] = u, flower[u].
      clear();
   int w_max = 0;
   for (int u = 1; u <= n; ++u) for (int v = 1; v <= n;
      ++v) for (int v = 1; v <= n;</pre>
```

6.5Maximum flow

103

123

```
/* Sparse graph maximum flow : isap.*/
template <int MAXN = 1000, int MAXM = 100000>
13
50
```

```
if (flow > 0) {
    e.flow[k] -= flow, e.flow[k ^ 1] += flow;
    ret += flow, ext -= flow; } }
if (!~k) d[u] = -1; return ret; }
int solve (flow_edge_list &e, int n_, int s_, int t_)
    {
    int ans = 0; n = n_; s = s_; dinic::t = t_;
    while (bfs (e)) {
    for (int i = 0; i < n; ++i) w[i] = e.begin[i];
    ans += dfs (e, s, INF); }
    return ans; };
}</pre>
```

6.6 Minimum cost flow

```
-c; flow[size] - v, Degint,
int n, t, tf, tc, dis[MAXN], slack[MAXN], visit[
    MAXN];
int modlable() {
    int delta = INF;
    for (int i = 0; i < n; i++) {
        if (!visit[i] && slack[i] < delta) delta = slack[i]
   e.flow[i] -= delta; e.flow[i ^ 1] += delta; left
    -= delta;
if (!left) { visit[x] = false; return flow; }
  slack[y] = std::min (slack[y], dis[y] + e.cost[i]
```

6.7 Stoer Wagner algorithm

```
6.8
           DN maximum clique
 /* DN maximum clique : n <= 150 */
typedef bool BB[N]; struct Maxclique {
3 const BB *e; int pk, level; const float Tlimit;
4 struct Vertex { int i, d; Vertex (int i) : i(i), d(0)
{ } };
typedef std::vector <Vertex> Vertices; Vertices V;
typedef std::vector <int> ColorClass; ColorClass QMAX, O:
BB e[N]; int ans, sol[N]; for (...) e[x][y] = e[y][x]
= true:
```

```
59| Maxclique mc (e, n); mc.mcqdyn (sol, ans); //0-based.
60| for (int i = 0; i < ans; ++i) std::cout << sol[i] << std::endl;</pre>
```

6.9 Dominator tree

6.10 Tarjan

7 String

7.1 Manacher

7.2 Suffix Array

7.3 Suffix Automaton

7.4 Palindromic tree

```
/* Palindromic tree : extend () - returns whether the
tree has generated a new node. odd, even - the
root of two trees. last - the node representing
the last char. node::len - the palindromic string
length of the node. */
template <int MAXN = 1000000, int MAXC = 26>
struct palindromic_tree {
struct node {
node *child[MAXC], *fail; int len;
node (int len) : fail (NULL), len (len) {
memset (child, NULL, sizeof (child)); }
} node pool[MAXN * 2], *tot_node;
int size, text[MAXN];
node *odd, *even, *last;
node *match (node *now) {
for (; text[size - now -> len - 1] != text[size];
now = now -> fail);
return now; }
bool extend (int token) {
text[++size] = token; node *now = match (last);
if (now -> child[token])
return last = now -> child[token], false;
```

```
18
     if (now == odd) last -> fail = even;
else {
  now = match (now -> fail);
  last -> fail = now -> child[token]; }
  return true; }

void init() {
  text[size = 0] = -1; tot_node = node_pool;
  last = even = new (tot_node++) node (0); odd = new (
     tot_node++) node (-1);
  even -> fail = odd; }

palindromic_tree () { init (); } };
```

Regular expression

```
std::regex_match (str, match, pattern);
 | The word is match[0], backreferences are match[i]
| wp to match.size ().
| match.prefix () and match.suffix () give the prefix
| and the suffix.
| match.length () gives length and match.position ()
| gives position of the match. */ }
| std::regex_replace (str, pattern, "sh$1");
| //$n is the backreference, $& is the entire match, $\frac{1}{2}$
| is the prefix, $' is the suffix, $\frac{1}{2}$ is the $\frac{1}{2}$ sign.
```

8 Tips 8.1 Java

```
doubleValue () / toPlainString () : converts to other
  types.

20 Arrays: Arrays.sort (T [] a); Arrays.sort (T [] a, int fromIndex, int toIndex); Arrays.sort (T [] a, int fromIndex, int toIndex, Comperator <? super T>
 | Int fromindex, int toindex, Comperator <? super T>
| comperator);
| LinkedList <E> : addFirst / addLast (E) / getFirst /
| getLast / removeFirst / removeLast () / clear () /
| add (int, E) / remove (int) / size () / contains
| / removeFirstOccurrence / removeLastOccurrence (E)
| ListIterator <E> listIterator (int index) : returns an iterator :
                          iterator
 iterator :
23    E next / previous () : accesses and iterates.
24    hasNext / hasPrevious () : checks availablity.
25    nextIndex / previousIndex () : returns the index of a subsequent call.
26    add / set (E) / remove () : changes element.
27    PriorityQueue <E> (int initcap, Comparator <? super E> comparator) : add (E) / clear () / iterator () / peek () / poll () / size ()
28    TreeMap <K, V> (Comparator <? super K> comparator) :
        Map.Entry <K, V> ceilingEntry / floorEntry / higherEntry / lowerEntry (K): getKey / getValue () / setValue (V) : entries.
29    clear () / put (K, V) / get (K) / remove (K) / size ()
  StringBuilder: StringBuilder (string) / append (int, string, ...) / insert (int offset, ...) charAt (int) / setCharAt (int, char) / delete (int, int) / reverse () / length () / toString ()
```

```
public Point (int xx, int yy) {
    x = xx;
    y = yy; };
public static class Cmp implements Comparator <Point>
        return;
53 */
54 /* or
        public static class Point implements Comparable <</pre>
            Point> {
public int x; public int y;
public Point () {
x = 0;
y = 0;
          public lint x, pre-
public Point () {
  x = 0;
  y = 0; }
public Point (int xx, int yy) {
  x = xx;
  y = yy; }
public int compareTo (Point p) {
  if (x < p.x) return -1;
  if (x == p.x) {
   if (y < p.y) return 0; }
  return 1; }
public boolean equalTo (Point p) {
  return (x == p.x && y == p.y); }
public int hashCode () {
  return x + y; };
//</pre>
73 */
74 //Faster IO
hasMoreTokens()) {
  try {
    String line = reader.readLine();
    tokenizer = new StringTokenizer (line);
    } catch (IOException e) {
      throw new RuntimeException (e); }
    return tokenizer.nextToken(); }
    public BigInteger nextBigInteger() {
      return new BigInteger (next (), 10); /* radix */ }
    public int nextInt() {
      return Integer.parseInt (next()); }
    public double nextDouble() {
      return Double.parseDouble (next()); }
    public static void main (String[] args) {
      InputReader in = new InputReader (System.in);
    }
```

8.2Random numbers

```
std::mt19937_64 mt (time (0));
std::uniform_int_distribution <int> uid (1, 100);
std::uniform_real_distribution <double> urd (1, 100);
std::cout << uid (mt) << "_" << urd (mt) << "\n";
```

8.3 Formatting

```
std::cout << std::fixed; // std::cout << std::
    scientific;</pre>
```

Read hack 8.4

Stack hack 8.5

```
//C++
#pragma comment (linker, "/STACK:36777216")
//G++
int __size__ = 256 << 20;
char *_p_ = (char*) malloc(__size__) + __size__;
__asm__ ("movl_%0,_%*esp\n" :: "r"(_p__));</pre>
```

8.6 Time hack

```
clock_t t = clock ();
std::cout << 1. * t / CLOCKS_PER_SEC << "\n";</pre>
```

8.7 **Builtin functions**

- _builtin_clz: Returns the number of leading 0-bits in x, starting at the most significant bit position. If x is 0, the result is
- undefined. __builtin_ctz: Returns the number of trailing 0-bits in x, starting at the least significant bit position. If x is 0, the result is
- undefined. __builtin_clrsb: Returns the number of leading redundant sign bits in x, i.e. the number of bits following the most significant bit that are identical to it. There are no special cases for 0 or
- other values.
 _builtin.popcount: Returns the number of 1-bits in x.
 _builtin.parity: Returns the parity of x, i.e. the number of
- 1-bits in x modulo 2. _builtin.bswap16, _builtin.bswap32, _builtin.bswap64: Returns x with the order of the bytes (8 bits as a group) reversed.
- bitset::Find_first(), bitset::Find_next(idx): set built in functions.

Prufer sequence

In combinatorial mathematics, the Prufer sequence of a labeled tree is a unique sequence associated with the tree. The sequence for a tree on n vertices has length n-2.

One can generate a labeled tree's Prufer sequence by iteratively One can generate a labeled tree's Fruier sequence by iteratively removing vertices from the tree until only two vertices remain. Specifically, consider a labeled tree T with vertices 1, 2, ..., n. At step i, remove the leaf with the smallest label and set the ith element of the Prufer sequence to be the label of this leaf's neighbour.

One can generate a labeled tree from a sequence in three steps. The tree will have n+2 nodes, numbered from 1 to n+2. For each node set its degree to the number of times it appears in the sequence plus

1. Next, for each number in the sequence a[i], find the first (lowestnumbered) node, j, with degree equal to 1, add the edge (j, a[i]) to the tree, and decrement the degrees of j and a[i]. At the end of this loop two nodes with degree 1 will remain (call them u, v). Lastly, add the edge (u, v) to the tree.

The Prufer sequence of a labeled tree on n vertices is a unique sequence of length n-2 on the labels 1 to n - this much is clear. Somewhat less obvious is the fact that for a given sequence S of length n-2 on the labels 1 to n, there is a unique labeled tree whose Prufer sequence

$\begin{array}{c} \text{is } S. \\ \textbf{8.9} \end{array}$ Spanning tree counting

Kirchhoff's Theorem: the number of spanning trees in a graph G is equal to any cofactor of the Laplacian matrix of G, which is equal to the difference between the graph's degree matrix (a diagonal matrix with vertex degrees on the diagonals) and its adjacency matrix (a (0,1)matrix with 1's at places corresponding to entries where the vertices are adjacent and 0's otherwise).

The number of edges with a certain weight in a minimum spanning tree is fixed given a graph. Moreover, the number of its arrangements can be obtained by finding a minimum spanning tree, compressing connected components of other edges in that tree into a point, and then applying Kirrchoff's theorem with only edges of the certain weight in the graph. Therefore, the number of minimum spanning trees in a graph can be solved by multiplying all numbers of arrangements of edges of different weights together.

Mobius inversion 8.10

Mobius inversion formula 8.10.1

$$[x = 1] = \sum_{d|x} \mu(d)$$
$$x = \sum_{d|x} \mu(d)$$

8.10.2Gcd inversion

$$\begin{split} \sum_{a=1}^n \sum_{b=1}^n gcd^2(a,b) &= \sum_{d=1}^n d^2 \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [gcd(i,j) = 1] \\ &= \sum_{d=1}^n d^2 \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{t \mid gcd(i,j)} \mu(t) \\ &= \sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} [t|i] \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [t|j] \\ &= \sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \lfloor \frac{n}{dt} \rfloor^2 \end{split}$$

The formula can be computed in O(nlogn) complexity. Moreover, let l = dt, then

$$\sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \lfloor \frac{n}{dt} \rfloor^2 = \sum_{l=1}^n \lfloor \frac{n}{l} \rfloor^2 \sum_{d|l} d^2 \mu(\frac{l}{d})$$

Let $f(l)=\sum_{d|l}d^2\mu(\frac{l}{d})$. It can be proven that f(l) is multiplicative. Besides, $f(p^k)=p^{2k}-p^{2k-2}$.

Therefore, with linear sieve the formula can be computed in O(n)complexity.

8.11 2-SAT

In terms of the implication graph, two literals belong to the same strongly connected component whenever there exist chains of implications from one literal to the other and vice versa. Therefore, the two literals must have the same value in any satisfying assignment to the given 2-satisfiability instance. In particular, if a variable and its negation both belong to the same strongly connected component, the instance cannot be satisfied, because it is impossible to assign both of these literals the same value. As Aspvall et al. showed, this is a necessary and sufficient condition: a 2-CNF formula is satisfiable if and only if there is no variable that belongs to the same strongly connected component as its negation.

This immediately leads to a linear time algorithm for testing satisfiability of 2-CNF formulae: simply perform a strong connectivity analysis on the implication graph and check that each variable and its negation belong to different components. However, as Aspvall et al. also showed, it also leads to a linear time algorithm for finding a satisfying assignment, when one exists. Their algorithm performs the following

Construct the implication graph of the instance, and find its strongly connected components using any of the known linear-time algorithms for strong connectivity analysis.

Check whether any strongly connected component contains both a variable and its negation. If so, report that the instance is not satisfiable

Construct the condensation of the implication graph, a smaller graph that has one vertex for each strongly connected component, and an edge from component i to component j whenever the implication graph contains an edge uv such that u belongs to component i and v belongs to component j. The condensation is automatically a directed acyclic graph and, like the implication graph from which it was formed, it is skew-symmetric.

Topologically order the vertices of the condensation. In practice this may be efficiently achieved as a side effect of the previous step, as components are generated by Kosaraju's algorithm in topological order and by Tarjan's algorithm in reverse topological order.

For each component in the reverse topological order, if its variables do not already have truth assignments, set all the literals in the component to be true. This also causes all of the literals in the complementary component to be set to false.

8.12Interesting numbers

Binomial Coefficients

$$\binom{n}{k} = (-1)^k \binom{k-n-1}{k}, \quad \sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

$$\sum_{k=0}^n \binom{k}{m} = \binom{n+1}{m+1}$$

$$\sqrt{1+z} = 1 + \sum_{k=1}^\infty \frac{(-1)^{k-1}}{k \times 2^{2k-1}} \binom{2k-2}{k-1} z^k$$

$$\sum_{k=0}^r \binom{r-k}{m} \binom{s+k}{n} = \binom{r+s+1}{m+n+1}$$

$$C_{n,m} = \binom{n+m}{m} - \binom{n+m}{m-1}, n \ge m$$

$$\binom{n}{k} \equiv [n\&k = k] \pmod{2}$$

$$\binom{n_1 + \dots + n_p}{m} = \sum_{k_1 + \dots + k_p = m} \binom{n_1}{k_1} \dots \binom{n_p}{k_p}$$

8.12.2 Fibonacci Numbers

$$F(z) = \frac{z}{1-z-z^2}$$

$$f_n = \frac{\phi^n - \hat{\phi}^n}{\sqrt{5}}, \phi = \frac{1+\sqrt{5}}{2}, \hat{\phi} = \frac{1-\sqrt{5}}{2}$$

$$\sum_{k=1}^n f_k = f_{n+2} - 1, \quad \sum_{k=1}^n f_k^2 = f_n f_{n+1}$$

$$\sum_{k=0}^n f_k f_{n-k} = \frac{1}{5}(n-1)f_n + \frac{2}{5}nf_{n-1}$$

$$\frac{f_{2n}}{f_n} = f_{n-1} + f_{n+1}$$

$$f_1 + 2f_2 + 3f_3 + \dots + nf_n = nf_{n+2} - f_{n+3} + 2]$$

$$\gcd(f_m, f_n) = f_{\gcd(m, n)}$$

$$f_n^2 + (-1)^n = f_{n+1}f_{n-1}$$

$$f_{n+k} = f_n f_{k+1} + f_{n-1}f_k$$

$$f_{2n+1} = f_n^2 + f_{n+1}^2$$

$$(-1)^k f_{n-k} = f_n f_{k-1} - f_{n-1}f_k$$

$$f_n = f_n f_n f_{k-1} - f_{n-1}f_k$$

$$f_n = f_n f_n f_{k-1} - f_{n-1}f_k$$

$$f_$$

8.12.3 Lucas Numbers

$$L_0 = 2, L_1 = 1, L_n = L_{n-1} + L_{n-2} = \left(\frac{1+\sqrt{5}}{2}\right)^n + \frac{1-\sqrt{5}}{2}^n$$
$$L(x) = \frac{2-x}{1-x-x^2}$$

8.12.4 Catlan Numbers

$$c_1 = 1, c_n = \sum_{i=0}^{n-1} c_i c_{n-1-i} = c_{n-1} \frac{4n-2}{n+1} = \frac{\binom{2n}{n}}{n+1}$$
$$= \binom{2n}{n} - \binom{2n}{n-1}, c(x) = \frac{1-\sqrt{1-4x}}{2x}$$

8.12.5 Stirling Cycle Numbers

Divide n elements into k non-empty cycles.

$$\begin{split} s(n,0) &= 0, s(n,n) = 1, s(n+1,k) = s(n,k-1) - ns(n,k) \\ & s(n,k) = (-1)^{n-k} {n \brack k} \\ & {n+1 \brack k} = n {n \brack k} + {n \brack k-1}, {n+1 \brack 2} = n! H_n \\ & x^{\underline{n}} = x(x-1)...(x-n+1) = \sum_{k=0}^n {n \brack k} (-1)^{n-k} x^k \\ & x^{\overline{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^n {n \brack k} x^k \end{split}$$

8.12.6 Stirling Subset Numbers

Divide n elements into k non-empty subsets.

$${n+1 \choose k} = k {n \choose k} + {n \choose k-1}$$

$$x^n = \sum_{k=0}^n {n \choose k} x^{\underline{k}} = \sum_{k=0}^n {n \choose k} (-1)^{n-k} x^{\overline{k}}$$

$$m! {n \choose m} = \sum_{k=0}^m {m \choose k} k^n (-1)^{m-k}$$

$$\sum_{k=1}^n k^p = \sum_{k=0}^p {p \choose k} (n+1)^{\underline{k}}$$

For a fixed k, generating functions:

$$\sum_{n=0}^{\infty} {n \brace k} x^{n-k} = \prod_{r=1}^{k} \frac{1}{1 - rx}$$

8.12.7 Motzkin Numbers

Draw non-intersecting chords between n points on a circle. Pick n numbers $k_1,k_2,...,k_n \in \{-1,0,1\}$ so that $\sum_i^a k_i (1 \le a \le n)$ is non-negative and the sum of all numbers is 0.

$$M_{n+1} = M_n + \sum_{i=0}^{n-1} M_i M_{n-1-i} = \frac{(2n+3)M_n + 3nM_{n-1}}{n+3}$$
$$M_n = \sum_{i=0}^{\lfloor \frac{n}{2} \rfloor} {n \choose 2k} Catlan(k)$$

$$M(X) = \frac{1 - x - \sqrt{1 - 2x - 3x^2}}{2x^2}$$

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8.12.8 Eulerian Numbers

Permutations of the numbers 1 to n in which exactly k elements are greater than the previous element.

8.12.9 Harmonic Numbers

Sum of the reciprocals of the first n natural numbers.

$$\sum_{k=1}^{n} H_k = (n+1)H_n - n$$

$$\sum_{k=1}^{n} kH_k = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}$$

$$\sum_{k=1}^{n} {k \choose m} H_k = {n+1 \choose m+1} (H_{n+1} - \frac{1}{m+1})$$

8.12.10 Pentagonal Number Theorem

$$\prod_{n=1}^{\infty} (1-x^n) = \sum_{n=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2}$$

$$p(n) = p(n-1) + p(n-2) - p(n-5) - p(n-7) + \cdots$$

$$f(n,k) = p(n) - p(n-k) - p(n-2k) + p(n-5k) + p(n-7k) - \cdots$$

8.12.11 Bell Numbers
Divide a set that has exactly n elements.

$$B_n = \sum_{k=1}^n {n \brace k}, \quad B_{n+1} = \sum_{k=0}^n {n \brack k} B_k$$

$$B_{p^m+n} \equiv mB_n + B_{n+1} \pmod{p}$$

$$B(x) = \sum_{n=0}^\infty \frac{B_n}{n!} x^n = e^{e^x - 1}$$

8.12.12 Bernoulli Numbers

$$B_n = 1 - \sum_{k=0}^{n-1} {n \choose k} \frac{B_k}{n-k+1}$$

$$G(x) = \sum_{k=0}^{\infty} \frac{B_k}{k!} x^k = \frac{1}{\sum_{k=0}^{\infty} \frac{x^k}{(k+1)!}}$$

$$\sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k n^{m-k+1}$$

8.12.13 Sum of Powers

$$\begin{split} \sum_{i=1}^{n} i^2 &= \frac{n(n+1)(2n+1)}{6}, \quad \sum_{i=1}^{n} i^3 = (\frac{n(n+1)}{2})^2 \\ &\sum_{i=1}^{n} i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} \\ &\sum_{i=1}^{n} i^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12} \end{split}$$

8.12.14 Sum of Squares

Denote $r_k(n)$ the ways to form n with k squares. If:

$$n = 2^{a_0} p_1^{2a_1} \cdots p_r^{2a_r} q_1 b_1 \cdots q_s b_s$$

where $p_i \equiv 3 \mod 4$, $q_i \equiv 1 \mod 4$, then

$$r_2(n) = \begin{cases} 0 & \text{if any } a_i \text{ is a half-integer} \\ 4\prod_{i=1}^{r} (b_i + 1) & \text{if all } a_i \text{ are integers} \end{cases}$$

 $r_3(n) > 0$ when and only when n is not $4^a(8b+7)$.

8.12.15 Derangement

$$D_1 = 0, D_2 = 1, D_n = n! \left(\frac{1}{0!} - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + \frac{(-1)^n}{n!}\right)$$
$$D_n = (n-1)(D_{n-1} + D_{n-2})$$

8.12.16 Tetrahedron Volume

If U, V, W, u, v, w are lengths of edges of the tetrahedron (first three form a triangle; u opposite to U and so on)

$$V = \frac{\sqrt{4u^2v^2w^2 - \sum_{cyc} u^2(v^2 + w^2 - U^2)^2 + \prod_{cyc} (v^2 + w^2 - U^2)}}{12}$$

Appendix

9.1 Calculus table

9.1.1 $ax + b \ (a \neq 0)$

1.
$$\int \frac{x}{ax+b} dx = \frac{1}{a^2} (ax+b-b \ln |ax+b|) + C$$

2.
$$\int \frac{x^2}{ax+b} dx = \frac{1}{a^3} \left(\frac{1}{2} (ax+b)^2 - 2b(ax+b) + b^2 \ln|ax+b| \right) + C$$

3.
$$\int \frac{\mathrm{d}x}{x(ax+b)} = -\frac{1}{b} \ln \left| \frac{ax+b}{x} \right| + C$$

4.
$$\int \frac{\mathrm{d}x}{x^2(ax+b)} = -\frac{1}{bx} + \frac{a}{b^2} \ln \left| \frac{ax+b}{x} \right| + C$$

5.
$$\int \frac{x}{(ax+b)^2} dx = \frac{1}{a^2} \left(\ln|ax+b| + \frac{b}{ax+b} \right) + C$$

7.
$$\int \frac{\mathrm{d}x}{x(ax+b)^2} = \frac{1}{b(ax+b)} - \frac{1}{b^2} \ln \left| \frac{ax+b}{x} \right| + C$$

9.1.2 $\sqrt{ax+b}$

1.
$$\int \sqrt{ax+b} dx = \frac{2}{3a} \sqrt{(ax+b)^3} + C$$

2.
$$\int x\sqrt{ax+b}dx = \frac{2}{15a^2}(3ax-2b)\sqrt{(ax+b)^3} + 6$$

3.
$$\int x^2 \sqrt{ax+b} dx = \frac{2}{105a^3} (15a^2x^2 - 12abx + 8b^2) \sqrt{(ax+b)^3} + 6a^2 + 6a$$

4.
$$\int \frac{x}{\sqrt{ax+b}} dx = \frac{2}{3a^2} (ax-2b)\sqrt{ax+b} + C$$

7.
$$\int \frac{\mathrm{d}x}{x^2 \sqrt{ax+b}} = -\frac{\sqrt{ax+b}}{bx} - \frac{a}{2b} \int \frac{\mathrm{d}x}{x\sqrt{ax+b}}$$

8.
$$\int \frac{\sqrt{ax+b}}{x} dx = 2\sqrt{ax+b} + b \int \frac{dx}{x\sqrt{ax+b}}$$

8.
$$\int \frac{\sqrt{ax+b}}{x} dx = 2\sqrt{ax+b} + b \int \frac{dx}{x\sqrt{ax+b}}$$
9.
$$\int \frac{\sqrt{ax+b}}{x^2} dx = -\frac{\sqrt{ax+b}}{x} + \frac{a}{2} \int \frac{dx}{\sqrt{ax+b}}$$

9.1.3 $x^2 \pm a^2$

1.
$$\int \frac{\mathrm{d}x}{x^2 + a^2} = \frac{1}{a} \arctan \frac{x}{a} + C$$

2.
$$\int \frac{\mathrm{d}x}{(x^2+a^2)^n} = \frac{x}{2(n-1)a^2(x^2+a^2)^{n-1}} + \frac{2n-3}{2(n-1)a^2} \int \frac{\mathrm{d}x}{(x^2+a^2)^{n-1}}$$

3.
$$\int \frac{\mathrm{d}x}{x^2 - a^2} = \frac{1}{2a} \ln \left| \frac{x - a}{x + a} \right| + C$$

3.
$$\int \frac{1}{x^2 - a^2} = \frac{1}{2a} \ln \left| \frac{1}{x + a} \right| + C$$
9.1.4 $ax^2 + b$ $(a > 0)$

1.
$$\int \frac{dx}{ax^2 + b} = \begin{cases} \frac{1}{\sqrt{ab}} \arctan \sqrt{\frac{a}{b}}x + C & (b > 0) \\ \frac{1}{2\sqrt{-ab}} \ln \left| \frac{\sqrt{ax} - \sqrt{-b}}{\sqrt{ax} + \sqrt{-b}} \right| + C & (b < 0) \end{cases}$$
2.
$$\int \frac{x}{ax^2 + b} dx = \frac{1}{2a} \ln \left| ax^2 + b \right| + C$$
3.
$$\int \frac{x^2}{ax^2 + b} dx = \frac{x}{a} - \frac{b}{a} \int \frac{dx}{ax^2 + b}$$
4.
$$\int \frac{dx}{x(ax^2 + b)} = \frac{1}{2b} \ln \frac{x^2}{|ax^2 + b|} + C$$
5.
$$\int \frac{dx}{x^2(ax^2 + b)} = -\frac{1}{bx} - \frac{b}{a} \int \frac{dx}{ax^2 + b}$$

2.
$$\int \frac{x}{1-x^2+b} dx = \frac{1}{2a} \ln |ax^2+b| + C$$

3.
$$\int \frac{x^2}{ax^2+b} dx = \frac{x}{a} - \frac{b}{a} \int \frac{dx}{ax^2+b}$$

5.
$$\int \frac{\mathrm{d}x}{2(-2+1)} = -\frac{1}{hx} - \frac{a}{h} \int \frac{\mathrm{d}x}{2+1}$$

$$\int \frac{1}{x^2(ax^2+b)} = -bx - b \int \frac{1}{ax^2+b}$$

$$\begin{array}{l}
0. \quad \int \frac{x}{x^{3}(ax^{2}+b)} = \frac{2b^{2} \ln \frac{x}{x^{2}} - \frac{2bx^{2}}{2b(ax^{2}+b)} + \frac{1}{2b} \int \frac{dx}{ax^{2}+b} \\
7. \quad \int \frac{dx}{(ax^{2}+b)^{2}} = \frac{x}{2b(ax^{2}+b)} + \frac{1}{2b} \int \frac{dx}{ax^{2}+b} \\
9.1.5 \quad ax^{2} + bx + c \quad (a > 0) \\
1. \quad \frac{dx}{ax^{2}+bx+c} = \begin{cases} \frac{2}{\sqrt{4ac-b^{2}}} \arctan \frac{2ax+b}{\sqrt{4ac-b^{2}}} + C \quad (b^{2} < 4ac) \\ \frac{1}{\sqrt{b^{2}-4ac}} \ln \left| \frac{2ax+b-\sqrt{b^{2}-4ac}}{2ax+b+\sqrt{b^{2}-4ac}} \right| + C \quad (b^{2} > 4ac) \\
2. \quad \int \frac{x}{ax^{2}+bx+c} dx = \frac{1}{2a} \ln |ax^{2} + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^{2}+bx+c} \\
0.1.6 \quad \sqrt{x^{2}+a^{2}} \quad \sqrt{a^{2}+a^{2}} \quad (a > 0)
\end{cases}$$

2.
$$\int \frac{x}{ax^2 + bx + c} dx = \frac{1}{2a} \ln|ax^2 + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^2 + bx + c}$$

9.1.6 $\sqrt{x^2 + a^2}$ (a > 0)

1.
$$\int \frac{\mathrm{d}x}{\sqrt{x^2 + a^2}} = \operatorname{arsh} \frac{x}{a} + C_1 = \ln(x + \sqrt{x^2 + a^2}) + C$$

2.
$$\int \frac{dx}{\sqrt{(x^2+a^2)^3}} = \frac{x}{a^2\sqrt{x^2+a^2}} + C$$

3.
$$\int \frac{x}{\sqrt{x^2 + a^2}} dx = \sqrt{x^2 + a^2} + C$$

4.
$$\int \frac{x}{\sqrt{2x^2+3x^2}} dx = -\frac{1}{\sqrt{2x^2+3x^2}} + C$$

9.
$$\int \sqrt{x^2 + a^2} dx = \frac{x}{2} \sqrt{x^2 + a^2} + \frac{a^2}{2} \ln(x + \sqrt{x^2 + a^2}) + C$$

10.
$$\int \sqrt{(x^2 + a^2)^3} dx = \frac{x}{8} (2x^2 + 5a^2) \sqrt{x^2 + a^2} + \frac{3}{8} a^4 \ln(x + \sqrt{x^2 + a^2}) + C$$

11.
$$\int x\sqrt{x^2+a^2} dx = \frac{1}{3}\sqrt{(x^2+a^2)^3} + C$$

12.
$$\int x^2 \sqrt{x^2 + a^2} dx = \frac{x}{8} (2x^2 + a^2) \sqrt{x^2 + a^2} - \frac{a^4}{8} \ln(x + \sqrt{x^2 + a^2}) + C$$

13.
$$\int \frac{\sqrt{x^2 + a^2}}{x} dx = \sqrt{x^2 + a^2} + a \ln \frac{\sqrt{x^2 + a^2} - a}{|x|} + C$$

14.
$$\int \frac{\sqrt{x^2 + a^2}}{x^2} dx = -\frac{\sqrt{x^2 + a^2}}{x} + \ln(x + \sqrt{x^2 + a^2}) + C$$

9.1.7 $\sqrt{x^2-a^2}$ (a>0)

1.
$$\int \frac{\mathrm{d}x}{\sqrt{x^2 - a^2}} = \frac{x}{|x|} \operatorname{arch} \frac{|x|}{a} + C_1 = \ln \left| x + \sqrt{x^2 - a^2} \right| + C$$

3.
$$\int \frac{x}{\sqrt{x^2-x^2}} dx = \sqrt{x^2-a^2} + C$$

4.
$$\int \frac{x}{\sqrt{(-2-2)^3}} dx = -\frac{1}{\sqrt{(2-2-2)^3}} + C$$

6.
$$\int \frac{x^2}{\sqrt{(x^2 - a^2)^3}} dx = -\frac{x}{\sqrt{x^2 - a^2}} + \ln|x + \sqrt{x^2 - a^2}| + C$$
7.
$$\int \frac{dx}{x\sqrt{x^2 - a^2}} = \frac{1}{a} \arccos \frac{a}{|x|} + C$$

7.
$$\int \frac{\mathrm{d}x}{x\sqrt{x^2-a^2}} = \frac{1}{a} \arccos \frac{a}{|x|} + C$$

$$\int \sqrt{x^2 - a^2} \, dx = \frac{x}{2} \sqrt{x^2 - a^2} - \frac{a^2}{2} \ln|x + \sqrt{x^2 - a^2}| + C$$

$$x^{2}\sqrt{x^{2}-a^{2}} = \frac{a^{2}}{2} \left(x^{2} - a^{2} - \frac{a^{2}}{2} \ln|x + \sqrt{x^{2}-a^{2}}| + C \right)$$

$$10. \int \sqrt{(x^{2} - a^{2})^{3}} dx = \frac{x}{8} (2x^{2} - 5a^{2}) \sqrt{x^{2} - a^{2}} + \frac{3}{8} a^{4} \ln|x + \sqrt{x^{2} - a^{2}}| + C$$

$$11. \int x\sqrt{x^{2} - a^{2}} dx = \frac{1}{3} \sqrt{(x^{2} - a^{2})^{3}} + C$$

$$12. \int x\sqrt{x^{2} - a^{2}} dx = \frac{1}{3} \sqrt{(x^{2} - a^{2})^{3}} + C$$

$$13. \int x\sqrt{x^{2} - a^{2}} dx = \frac{1}{3} \sqrt{(x^{2} - a^{2})^{3}} + C$$

11.
$$\int x \sqrt{x^2 - a^2} dx = \frac{1}{2} \sqrt{(x^2 - a^2)^3} + C$$

12.
$$\int x^2 \sqrt{x^2 - a^2} \, dx = \frac{x}{8} (2x^2 - a^2) \sqrt{x^2 - a^2} - \frac{a^4}{8} \ln|x + \sqrt{x^2 - a^2}| + C$$

13.
$$\int \frac{\sqrt{x^2 - a^2}}{x} dx = \sqrt{x^2 - a^2} - a \arccos \frac{a}{|x|} + C$$

14.
$$\int \frac{\sqrt{x^2 - a^2}}{x^2} dx = -\frac{\sqrt{x^2 - a^2}}{x} + \ln|x + \sqrt{x^2 - a^2}| + C$$

9.1.8 $\sqrt{a^2 - x^2}$ (a > 0)1. $\int \frac{dx}{\sqrt{a^2 - x^2}} = \arcsin \frac{x}{a} + C$

1.
$$\int \frac{\mathrm{d}x}{\sqrt{a^2 - x^2}} = \arcsin \frac{x}{a} + C$$

2.
$$\frac{dx}{\sqrt{(a^2 - x^2)^3}} = \frac{x}{a^2 \sqrt{a^2 - x^2}} + C$$
3.
$$\int \frac{x}{\sqrt{a^2 - x^2}} dx = -\sqrt{a^2 - x^2} + C$$

3.
$$\int \frac{x}{\sqrt{a^2 - x^2}} dx = -\sqrt{a^2 - x^2} + C$$

4.
$$\int \frac{x}{\sqrt{(a^2-x^2)^3}} dx = \frac{1}{\sqrt{a^2-x^2}} + C$$

9.
$$\int \sqrt{a^2 - x^2} dx = \frac{x}{2} \sqrt{a^2 - x^2} + \frac{a^2}{2} \arcsin \frac{x}{a} + C$$

10.
$$\int \sqrt{(a^2 - x^2)^3} dx = \frac{x}{8} (5a^2 - 2x^2) \sqrt{a^2 - x^2} + \frac{3}{8} a^4 \arcsin \frac{x}{a} + C$$

11.
$$\int x\sqrt{a^2-x^2} dx = -\frac{1}{2}\sqrt{(a^2-x^2)^3} + C$$

11.
$$\int \sqrt{x^2 - x^2} \, dx = -\frac{1}{3} \sqrt{(a^2 - x^2)^3 + C}$$
12.
$$\int x^2 \sqrt{a^2 - x^2} \, dx = \frac{x}{8} (2x^2 - a^2) \sqrt{a^2 - x^2} + \frac{a^4}{8} \arcsin \frac{x}{a} + C$$
13.
$$\int \frac{\sqrt{a^2 - x^2}}{x} \, dx = \sqrt{a^2 - x^2} + a \ln \frac{a - \sqrt{a^2 - x^2}}{|x|} + C$$

13.
$$\int \frac{\sqrt{a^2 - x^2}}{x} dx = \sqrt{a^2 - x^2} + a \ln \frac{a - \sqrt{a^2 - x^2}}{|x|} + C$$

14.
$$\int \frac{x}{\sqrt{a^2 - x^2}} dx = -\frac{\sqrt{a^2 - x^2}}{x} - \arcsin \frac{x}{a} + C$$

9.1.9 $\sqrt{\pm ax^2 + bx + c}$ (a > 0)

1.
$$\int \frac{\mathrm{d}x}{\sqrt{ax^2 + bx + c}} = \frac{1}{\sqrt{a}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

2.
$$\int \sqrt{ax^2 + bx + c} dx = \frac{2ax + b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac - b^2}{8\sqrt{a^3}} \ln|2ax| + b + \frac{a^2}{4a} + \frac{a^2}{4a}$$

$$2\sqrt{a\sqrt{ax^2+bx+c}}+C$$

$$2\sqrt{a}\sqrt{ax^{2} + bx + c}| + C$$
3.
$$\int \frac{x}{\sqrt{ax^{2} + bx + c}} dx = \frac{1}{a}\sqrt{ax^{2} + bx + c} - \frac{b}{2\sqrt{a^{3}}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^{2} + bx + c}| + C$$

4.
$$\int \frac{\mathrm{d}x}{\sqrt{c+bx-ax^2}} = -\frac{1}{\sqrt{a}}\arcsin\frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

4.
$$\int \frac{dx}{\sqrt{c+bx-ax^2}} = -\frac{1}{\sqrt{a}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

5. $\int \sqrt{c+bx-ax^2} dx = \frac{2ax-b}{4a} \sqrt{c+bx-ax^2} + \frac{b^2+4ac}{axcin} = \frac{2ax-b}{axcin} + C$

$$\frac{b^2 + 4ac}{8\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + c$$

6.
$$\int \frac{b^2 + 4ac}{8\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$
6.
$$\int \frac{x}{\sqrt{c + bx - ax^2}} dx = -\frac{1}{a}\sqrt{c + bx - ax^2} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$

9.1.10 $\sqrt{\pm \frac{x-a}{x-b}}$ & $\sqrt{(x-a)(x-b)}$

1.
$$\int \sqrt{\frac{x-a}{x-b}} dx = (x-b)\sqrt{\frac{x-a}{x-b}} + (b-a)\ln(\sqrt{|x-a|} + \sqrt{|x-b|}) + C$$

2.
$$\int \sqrt{\frac{x-a}{b-x}} dx = (x-b)\sqrt{\frac{x-a}{b-x}} + (b-a)\arcsin\sqrt{\frac{x-a}{b-x}} + C$$

3.
$$\int \frac{\mathrm{d}x}{\sqrt{(x-a)(b-x)}} = 2\arcsin\sqrt{\frac{x-a}{b-x}} + C \ (a < b)$$

4.
$$\int \sqrt{(x-a)(b-x)} \, \mathrm{d}x \, = \, \frac{2x-a-b}{4} \sqrt{(x-a)(b-x)} \, + \, \frac{(b-a)^2}{4} \arcsin \sqrt{\frac{x-a}{b-x}} \, + \, C$$

9.1.11 Triangular function

- 1. $\int \tan x dx = -\ln|\cos x| + C$ 2. $\int \cot x dx = \ln|\sin x| + C$
- 3. $\int \sec x dx = \ln \left| \tan \left(\frac{\pi}{4} + \frac{x}{2} \right) \right| + C = \ln \left| \sec x + \tan x \right| + C$
- 4. $\int \csc x \, dx = \ln \left| \tan \frac{x}{2} \right| + C = \ln \left| \csc x \cot x \right| + C$
- $\int \sec^2 x \, \mathrm{d}x = \tan x + C$
- 5. $\int \sec^{2} x dx = \tan x + C$ 6. $\int \csc^{2} x dx = -\cot x + C$ 7. $\int \sec x \tan x dx = \sec x + C$ 8. $\int \csc x \cot x dx = -\csc x + C$ 9. $\int \sin^{2} x dx = \frac{x}{2} \frac{1}{4} \sin 2x + C$ 10. $\int \cos^{2} x dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$

$$\begin{aligned} &11. & \int \sin^n x \, \mathrm{d}x = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x \, \mathrm{d}x \\ &12. & \int \cos^n x \, \mathrm{d}x = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x \, \mathrm{d}x \\ &13. & \frac{dx}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\sin^{n-2} x} \\ &14. & \frac{dx}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x} \\ &15. & \int \cos^m x \sin^n x \, \mathrm{d}x \\ &= \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x \, \mathrm{d}x \\ &= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, \mathrm{d}x \\ &= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, \mathrm{d}x \\ &16. & \int \sin ax \cos bx \, \mathrm{d}x = -\frac{1}{2(a+b)} \cos(a+b)x - \frac{1}{2(a-b)} \cos(a-b)x + C \\ &17. & \int \sin ax \sin bx \, \mathrm{d}x = -\frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C \\ &18. & \int \cos ax \cos bx \, \mathrm{d}x = \frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C \\ &19. & \int \frac{\mathrm{d}x}{a+b\sin x} = \begin{cases} \frac{2}{\sqrt{a^2-b^2}} \arctan \frac{a \tan \frac{x}{2} + b}{\sqrt{a^2-b^2}} + C & (a^2 > b^2) \\ \frac{1}{\sqrt{b^2-a^2}} \ln \left| \frac{a \tan \frac{x}{2} + b - \sqrt{b^2-a^2}}{a \tan \frac{x}{2} + b + \sqrt{b^2-a^2}} \right| + C & (a^2 < b^2) \end{cases} \\ &20. & \int \frac{\mathrm{d}x}{a+b\cos x} = \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan \left(\sqrt{\frac{a-b}{a+b}} \tan \frac{x}{2}\right) + C & (a^2 < b^2) \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln \left| \frac{\tan \frac{x}{2} + \sqrt{\frac{a+b}{b-a}}}{\tan \frac{x}{2} - \sqrt{\frac{a+b}{b-a}}} + C & (a^2 < b^2) \end{cases} \\ &21. & \int \frac{\mathrm{d}x}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{ab} \arctan \left(\frac{b}{b} \tan x + a\right) + C \\ &22. & \int \frac{\mathrm{d}x}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{2ab} \ln \left| \frac{b \tan x + a}{b \tan x} \right| + C \\ &23. & \int x \sin ax \, \mathrm{d}x = \frac{1}{a^2} \sin ax - \frac{1}{a} x \cos ax + C \\ &24. & \int x^2 \sin ax \, \mathrm{d}x = -\frac{1}{a^2} \cos ax + \frac{1}{a^2} \sin ax + C \end{cases} \\ &26. & \int x^2 \cos ax \, \mathrm{d}x = \frac{1}{a} x^2 \sin ax + \frac{2}{a^2} x \cos ax - \frac{2}{a^3} \sin ax + C \end{cases}$$

9.1.12 Inverse triangular function (a > 0)

$$\begin{array}{l} 1. \ \, \int \arcsin \frac{x}{a} \mathrm{d}x = x \arcsin \frac{x}{a} + \sqrt{a^2 - x^2} + C \\ 2. \ \, \int x \arcsin \frac{x}{a} \mathrm{d}x = (\frac{x^2}{2} - \frac{a^2}{4}) \arcsin \frac{x}{a} + \frac{x}{4} \sqrt{x^2 - x^2} + C \\ 3. \ \, \int x^2 \arcsin \frac{x}{a} \mathrm{d}x = \frac{x^3}{3} \arcsin \frac{x}{a} + \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 - x^2} + C \\ 4. \ \, \int \arccos \frac{x}{a} \mathrm{d}x = x \arccos \frac{x}{a} - \sqrt{a^2 - x^2} + C \\ 5. \ \, \int x \arccos \frac{x}{a} \mathrm{d}x = (\frac{x^2}{2} - \frac{a^2}{4}) \arccos \frac{x}{a} - \frac{x}{4} \sqrt{a^2 - x^2} + C \\ 6. \ \, \int x^2 \arccos \frac{x}{a} \mathrm{d}x = \frac{x}{3} \arccos \frac{x}{a} - \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 - x^2} + C \\ 7. \ \, \int \arctan \frac{x}{a} \mathrm{d}x = x \arctan \frac{x}{a} - \frac{a}{2} \ln(a^2 + x^2) + C \\ 8. \ \, \int x \arctan \frac{x}{a} \mathrm{d}x = \frac{1}{2} (a^2 + x^2) \arctan \frac{x}{a} - \frac{a}{2} x + C \\ 9. \ \, \int x^2 \arctan \frac{x}{a} \mathrm{d}x = \frac{x^3}{3} \arctan \frac{x}{a} - \frac{a}{6} x^2 + \frac{a^3}{6} \ln(a^2 + x^2) + C \\ \end{array}$$

9.1.13 Exponential function

```
1. \int a^x dx = \frac{1}{\ln a} a^x + C
2. \int e^{ax} dx = \frac{1}{a} a^{ax} + C
\begin{aligned} 2. & \int e^{ax} \, dx = \frac{1}{a} a^{ax} + C \\ 3. & \int x e^{ax} \, dx = \frac{1}{a^2} (ax - 1) a^{ax} + C \\ 4. & \int x^n e^{ax} \, dx = \frac{1}{a} x^n e^{ax} - \frac{n}{a} \int x^{n-1} e^{ax} \, dx \\ 5. & \int x a^x \, dx = \frac{x}{\ln a} a^x - \frac{1}{(\ln a)^2} a^x + C \\ 6. & \int x^n a^x \, dx = \frac{1}{\ln a} x^n a^x - \frac{n}{\ln a} \int x^{n-1} a^x \, dx \\ 7. & \int e^{ax} \sin bx \, dx = \frac{1}{a^2 + b^2} e^{ax} (a \sin bx - b \cos bx) + C \\ 6. & \int x^n \cos bx \, dx = \frac{1}{a^2 + b^2} e^{ax} (b \sin bx + a \cos bx) + C \end{aligned}
  \frac{n(n-1)b^2}{a^2+b^2n^2} \int e^{ax} \sin^{n-2} bx dx
10. \int e^{ax} \cos^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) +
             \frac{n(n-1)b^2}{a^2+b^2n^2} \int e^{ax} \cos^{n-2} bx dx
```

9.1.14 Logarithmic function

	$\int \ln x \mathrm{d}x = x \ln x - x + C$
	$\int \frac{\mathrm{d}x}{x \ln x} = \ln \left \ln x \right + C$
3.	$\int x^n \ln x dx = \frac{1}{n+1} x^{n+1} (\ln x - \frac{1}{n+1}) + C$
	$\int (\ln x)^n dx = x(\ln x)^n - n \int (\ln x)^{n-1} dx$
5.	$\int x^m (\ln x)^n dx = \frac{1}{m+1} x^{m+1} (\ln x)^n - \frac{n}{m+1} \int x^m (\ln x)^{n-1} dx$

9.2Regular expression

9.2.1Special pattern characters

Characters	Description
	Not newline
\t	Tab (HT)
\n	Newline (LF)
\v	Vertical tab (VT)
\f	Form feed (FF)
\r	Carriage return (CR)
\cletter	Control code
\xhh	ASCII character
\uhhhh	Unicode character
\0	Null
\int	Backreference
\d	Digit
\D	Not digit
\s	Whitespace
\S	Not whitespace
\ W	Word (letters, numbers and the underscore)
\W	Not word
\character	Character
[class]	Character class
[^class]	Negated character class

9.2.2Quantifiers

Characters	Times
*	0 or more
+	1 or more
?	0 or 1
{int}	int
{int,}	int or more
{min,max}	Between min and max

By default, all these quantifiers are greedy (i.e., they take as many characters that meet the condition as possible). This behavior can be overridden to ungreedy (i.e., take as few characters that meet the condition as possible) by adding a question mark (?) after the quantifier.

9.2.3Groups

Characters	Description
(subpattern)	Group with backreference
(?:subpattern)	Group without backreference

9.2.4Assertions

Characters	Description
^	Beginning of line
\$	End of line
\b	Word boundary
\B	Not a word boundary
(?=subpattern)	Positive lookahead
(?!subpattern)	Negative lookahead

9.2.5Alternative

A regular expression can contain multiple alternative patterns simply by separating them with the separator operator (|): The regular expression will match if any of the alternatives match, and as soon as 9.2.6 Character classes

Class	Description
[:alnum:]	Alpha-numerical character
[:alpha:]	Alphabetic character
[:blank:]	Blank character
[:cntrl:]	Control character
[:digit:]	Decimal digit character
[:graph:]	Character with graphical representation
[:lower:]	Lowercase letter
[:print:]	Printable character
[:punct:]	Punctuation mark character
[:space:]	Whitespace character
[:upper:]	Uppercase letter
[:xdigit:]	Hexadecimal digit character
[:d:]	Decimal digit character
[:w:]	Word character
[:s:]	Whitespace character

Please note that the brackets in the class names are additional to those opening and closing the class definition. For example:

[[:alpha:]] is a character class that matches any alphabetic character. [abc[:digit:]] is a character class that matches a, b, c, or a

[^[:space:]] is a character class that matches any character ex-