Suzune Nisiyama

July 18, 2018

## MIT License

Copyright (c) 2018 Nisiyama-Suzune

Permission is hereby granted, free of charge, to any person obtaining a copy of this software and associated documentation files (the "Software"), to deal in the Software without restriction, including without limitation the rights to use, copy, modify, merge, publish, distribute, sublicense, and/or sell copies of the Software, and to permit persons to whom the Software is furnished to do so, subject to the following conditions:

The above copyright notice and this permission notice shall be included in all copies or substantial portions of the Software. THE SOFTWARE IS PROVIDED "AS IS", WITHOUT WARRANTY OF ANY KIND, EXPRESS OR IMPLIED, INCLUDING BUT NOT LIMITED TO THE WARRANTIES OF MERCHANTABILITY, FITNESS FOR A PARTICULAR PURPOSE AND NONINFRINGEMENT. IN NO EVENT SHALL THE AUTHORS OR COPYRIGHT HOLDERS BE LIABLE FOR ANY CLAIM, DAMAGES OR OTHER LIABILITY, WHETHER IN AN ACTION OF CONTRACT, TORT OR OTHERWISE, ARISING FROM, OUT OF OR IN CONNECTION WITH THE SOFTWARE OR THE USE OR OTHER DEALINGS IN THE SOFTWARE.

			me hack	
1	Environment 1.1 Vimrc	<b>2</b> 2	8.7 Bu 8.8 Pr	ultiplication hack
2	Data Structure           2.1 KD tree	2 2 2 2	8.10 Me 8.11 2-5	eanning tree counting
3	Formula 3.1 Zellers congruence	2 2 2 3 3		dix alculus table
4	Number theory 4.1 Fast power module	3 3 3 3 4 4 4 4 4 4		
5	Geometry         5.1 Point          5.2 Line          5.3 Circle          5.4 Centers of a triangle          5.5 Fermat point          5.6 Convex hull          5.7 Half plane intersection          5.8 Nearest pair of points          5.9 Minimum circle          5.10 Intersection of a polygon and a circle          5.11 Union of circles          5.12 3D point          5.13 3D line          5.14 3D convex hull	4 5 5 5 5 5 6 6 6 6 6 7 7		
6	Graph 6.1 Hopcoft-Karp algorithm 6.2 Kuhn-Munkres algorithm 6.3 Blossom algorithm 6.4 Weighted blossom algorithm 6.5 Maximum flow 6.6 Minimum cost flow 6.7 Stoer Wagner algorithm 6.8 DN maximum clique 6.9 Dominator tree	7 7 7 8 8 8 9 9		
7	String 7.1 Suffix Array	10 10 10 10 11		
8	Tips           8.1 Java	<b>11</b> 11 11		

8.3

# **Environment**

### 1.1 Vimrc

```
7 au BufEnter *.java nm <F5> :!time java %< <CR>|nm <F8>
:!time java %< < %<.in <CR>|nm <F9> :!javac % <CR
```

## Data Structure

## 2.1 KD tree

```
/* kd_tree : finds the k-th closest point in O(kn^{1-\frac{1}{k}}). Usage : Stores the data in p[]. Call function init (n, k). Call min_kth (d, k). (or max_kth) (k is 1-based)
     2 Usage
 | based)
| Note : Switch to the commented code for Manhattan distance.
| Status : SPOJ-FAILURE Accepted.*/
| template <int MAXN = 200000, int MAXK = 2>
| struct kd_tree {
| int k, size;
| struct point { int data[MAXK], id; } p[MAXN];
| struct kd_node {
| int l, r; point p, dmin, dmax;
| kd_node (const point &rhs) : l (-1), r (-1), p (rhs) |
| , dmin (rhs), dmax (rhs) {}
| void merge (const kd_node &rhs, int k) {
| for (register int i = 0; i < k; ++i) {
| dmin.data[i] = std::max (dmax.data[i], rhs.dmin. data[i]); }
| long long min_dist (const point &rhs, int k) const {
| register long long ret = 0; |
| for (register int i = 0; i < k; ++i) {
| if (dmin.data[i] = rhs.data[i] & rhs.data[i] <= dmax.data[i] > continue; |
| ret += std::min (111 * (dmin.data[i] - rhs.data[i]) * (dmax. data[i] - rhs.data[i]) * (dmax. data[i] - rhs.data[i]) * (ret += std::max (0, rhs.data[i] - rhs.data[i]); |
| ret += std::max (0, dmin.data[i] - rhs.data[i]); |
| ret trun ret; |
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long max_dist (const point &rhs, int k) {
| long long ret = 0; |
| for (int i = 0; i < k; ++i) {
| int tmp = std::max (std::abs (dmin.data[i] - rhs.data[i]) |
| ret += std::max (std::abs (dmin.data[i] - rhs.data[i]) |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] - rhs.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] |
| ret += std::max (std::abs (rhs.data[i] - dmax.data[i] |
| r
      3 Note
                                                           Switch to the commented code for Manhattan
                                            distance.
21
23 //
28
                                         | ]));
ret += 111 * tmp * tmp; }
  ret += std::max (std::abs (rhs.data[i] - dmax.
data[i]) + std::abs (rhs.data[i] - dmin.data[i]));
                   return ret; } tree[MAXN * 4];
struct result {
 long long dist; point d; result() {}
 result (const long long &dist, const point &d) :
    dist (dist), d (d) {}
 bool operator > (const result &rhs) const { return
 32
34
                35
 38
                             if ("tree[rt].r) tree[rt].merge (tree[tree[rt].r], k
                  if (~tree[rt].r) tree[rt].merge (tree[tree[rt].r], k
    ); }
std::priority_queue<result, std::vector<result>, std
    ::less <result>> heap_1;
std::priority_queue<result, std::vector<result>, std
    ::greater <result>> heap_r;
void _min_kth (const int &depth, const int &rt, const
    int &m, const point &d) {
    result tmp = result (sqrdist (tree[rt].p, d), tree[
        rt].p);
    if ((int)heap_1.size() < m) heap_1.push (tmp);
    else if (tmp < heap_1.top()) {
        heap_1.push (tmp); }</pre>
 55
```

```
62
74
75
80
```

### Splay 2.2

```
void push_down (int x) {
   if ( n[x].c[0]) push (n[x].c[0], n[x].t);
   if ( n[x].c[1]) push (n[x].c[1], n[x].t);
   if ( n[x].t = tag (); }
   void update (int x) {
      n[x].m = gen (x);
   if ( n[x].c[0]) n[x].m = merge (n[n[x].c[0]].m, n[x].
    m);
if (\tilde{n}[x].c[1]) n[x].m = merge (n[x].m, n[n[x].c[1]].
  11
12
```

## Link-cut tree

```
= u;

n[u].c[1] = v;

if (~v) n[v].f = u, n[v].p = -1;

update (u); u = n[v = u].p; }

splay (x); }
```

## Formula

### Zellers congruence 3.1

```
/* Zeller's congruence : converts between a calendar
```

## 3.2 Lattice points below segment

```
1 /* Euclidean-like algorithm : computes the sum of
     \sum_{i=0}^{n-1} \left[ \frac{a+bi}{m} \right] . \star /
```

```
if (b >= m) return (n - 1) * n / 2 * (b / m) + solve
    (n, a, b % m, m);
return solve ((a + b * n) / m, (a + b * n) % m, m, b)
```

### Adaptive Simpson's method 3.3

```
|\mathbf{r}| / \star Adaptive Simpson's method : integrates f in [1, r].
  struct simpson {
  double area (double (*f) (double), double 1, double r
   double m = 1 + (r - 1) / 2;
return (f (1) + 4 * f (m) + f (r)) * (r - 1) / 6; }
double solve (double (*f) (double), double 1, double
    r, double eps, double a) {
    double m = 1 + (r - 1) / 2;
    double left = area (f, 1, m), right = area (f, m, r)
    :
```

### 3.4Neural network

```
1 /* Neural network : machine learning. */
2 template <int ft = 6, int n = 6, int MAXDATA = 100000>
3 struct network {
      double wp[n][ft], w[n], avg[ft + 1], sig[ft + 1], val
        network () {
std::mt19937_64 mt (time (0));
std::uniform_real_distribution <double> urdp (0, 2 *
sqrt (ft));
      network
        double
      void desc (double *x, double t, double eta) {
  double o = compute (x), del = (o - t); // * o * (1 -
         o)
for (int i = 0; i < n; ++i) for (int j = 0; j < ft;
21
           wp[i][j] -= eta * del * val[i] * (1 - val[i]) * w[i
    ] * x[j];
22
      26
     dn; ++j)
sig[i] += (data[j][i] - avg[i]) * (data[j][i]
for (int i = 0; i < ft + 1; ++i) sig[i] = sqrt (sig[i] / dn);
for (int i = 0; i < ft + 1; ++i) for (int j = 0; j < dn; ++j)
data[j][i] = (data[j][i] - avg[i]) / sig[i];
for (int cnt = 0; cnt < epoch; ++cnt) for (int test = 0; test < dn; ++test)
desc (data[test], data[test][ft], eta);
double predict (double *x) {
  for (int i = 0; i < ft; ++i) x[i] = (x[i] - avg[i])
  / sig[i];
  return compute (x) * sig[ft] + avg[ft]; }
</pre>
33
      return compute (x) * sig[ft] + avg[ft]; }
std::string to_string () {
  std::ostringstream os; os.precision (20); os << std</pre>
        stat.ostringscrem = -,
    ::fixed;
for (int i = 0; i < n; ++i) for (int j = 0; j < ft;
    ++j) os << wp[i][j] << "_";
for (int i = 0; i < n; ++i) os << w[i] << "_";
for (int i = 0; i < ft + 1; ++i) os << avg[i] << "_"</pre>
42
         for (int i = 0; i < ft + 1; ++i) os << sig[i] << "_"
      return os.str (); }
void read (const std::string &str) {
    std::istringstream is (str);
    for (int i = 0; i < n; ++i) for (int j = 0; j < ft;
        ++j) is >> wp[i][j];
    for (int i = 0; i < n; ++i) is >> w[i];
    for (int i = 0; i < ft + 1; ++i) is >> avg[i];
    for (int i = 0; i < ft + 1; ++i) is >> sig[i]; } };
```

### Number theory 4

### 4.1 Fast power module

```
1 /* Fast power module : x^n */ 2 int fpm (int x, int n, int mod) {
```

```
int ans = 1, mul = x; while (n) {
  if (n & 1) ans = int (111 * ans * mul * mod);
  mul = int (111 * mul * mul * mod); n >>= 1; }
  return ans; }
```

## 4.2 Euclidean algorithm

```
\frac{1}{2} /* Euclidean algorithm : solves for ax + by = gcd (a,
  b). */
void euclid (const long long &a, const long long &b,
    long long &x, long long &y) {
    if (b == 0) x = 1, y = 0;
    else euclid (b, a % b, y, x), y -= a / b * x; }
  long long inverse (long long x, long long m) {
  long long a, b; euclid (x, m, a, b); return (a % m +
     m) % m; }
```

## 4.3 Discrete Fourier transform

```
1 /* Discrete Fourier transform : the nafarious you-know
  -what thing.

2 Usage: call init for the suggested array size, and solve for the transform. (use f!=0 for the inverse
        template <int MAXN = 1000000>
 3 template <int MAXN = 1000000>
4 struct dft {
5    typedef std::complex <double> complex;
6    complex e[2][MAXN];
7    int init (int n) {
8        int len = 1;
9        for (; len <= 2 * n; len <<= 1);
10        for (int i = 0; i < len; ++i) {
11        e[0][i] = complex (cos (2 * PI * i / len), sin (2 * PI * i / len));
12        e[1][i] = complex (cos (2 * PI * i / len), -sin (2 * PI * i / len));
13        return len; }
14        void solve (complex *a, int n, int f) {</pre>
11
             return len; }
void solve (complex *a, int n, int f) {
  for (int i = 0, j = 0; i < n; ++i) {
    if (i > j) std::swap (a[i], a[j]);
    for (int t = n >> 1; (j ^= t) < t; t >>= 1); }
  for (int i = 2; i <= n; i <<= 1)
    for (int j = 0; j < n; j += i)
    for (int k = 0; k < (i >> 1); ++k) {
      complex A = a[j + k];
      complex B = e[f][n / i * k] * a[j + k + (i >> 1)
      ];
}
                 23
```

## Fast Walsh-Hadamard transform

```
1 /* Fast Walsh-Hadamard transform : binary operation
  /* Fast Walsh-Hadamara cransform. transform. */
void fwt (int *a, int n, int w) {
  for (int i = 1; i < n; i <<= 1)
    for(int j = 0; j < n; j += i << 1) {
      for(int k = 0; k < i; ++k) {
        int x = a[j + k], y = a[i + j + k];
        if (w) {
           }else{
/* xor : a[j + k] = x + y, a[i + j + k] = x - y,
    and : a[j + k] = x + y, or : a[i + j + k] = x
    + y; */
```

### Number theoretic transform 4.5

```
int mod, int prt) {
       inv[i][j] = (int) inverse (MOD[i], MOD[j]);
static int x[3];
for (int i = 0; i < 3; ++i) { x[i] = a[i];
    for (int j = 0; j < i; ++j) {
        int t = (x[i] - x[j] + MOD[i]) % MOD[i];</pre>
30
```

```
if (t < 0) t += MOD[i];
    x[i] = int (1LL * t * inv[j][i] % MOD[i]); } 
int sum = 1, ret = x[0] % mod;

for (int i = 1; i < 3; ++i) {
    sum = int (1LL * sum * MOD[i - 1] % mod);

ret += int (1LL * x[i] * sum % mod);

if (ret >= mod) ret -= mod; }

return ret; } ;
```

## 4.6 Chinese remainder theorem

## 4.7 Linear Recurrence

## 4.8 Berlekamp Massey algorithm

## 4.9 Baby step giant step algorithm

```
| /* Baby step giant step algorithm : Solves a^x = b \mod c | in O(\sqrt{c}). */
| struct bsgs {
| int solve (int a, int b, int c) {
| std::unordered_map <int, int> bs;
| int m = (int) sqrt ((double) c) + 1, res = 1;
| for (int i = 0; i < m; ++i) {
| if (bs.find (res) == bs.end ()) bs[res] = i;
| res = int (1LL * res * a % c); }
| int mul = 1, inv = (int) inverse (a, c);
| for (int i = 0; i < m; ++i) mul = int (1LL * mul * inv % c);
| res = b % c;
| for (int i = 0; i < m; ++i) {
| if (bs.find (res) != bs.end ()) return i * m + bs[ res];
| res = int (1LL * res * mul % c); }
| return -1; };
```

## 4.10 Miller Rabin primality test

## 4.11 Pollard's Rho algorithm

## 5 Geometry

```
#define cd const double &
const double EPS = 1E-8, PI = acos (-1);
int sgn (cd x) { return x < -EPS ? -1 : x > EPS; }
```

```
4 int cmp (cd x, cd y) { return sgn (x - y); } 5 double sqr (cd x) { return x * x; }
```

## 5.1 Point

## 5.2 Line

## 5.3 Circle

```
#define cc const circle &
struct circle {
    point c; double r;
    explicit circle (point c = point (), double r = 0) :
        c (c), r (r) {};
    bool operator == (cc a, cc b) { return a.c == b.c &&
        cmp (a.r, b.r) == 0; }
    bool operator != (cc a, cc b) { return ! (a == b); }
    bool in_circle (cp a, cc b) { return cmp (dis (a, b.c) , b.r) <= 0; }
    circle make_circle (cp a, cp b) { return circle ((a + b) / 2, dis (a, b) / 2); }</pre>
```

```
std::vector <point> line_circle_intersect (cl a, cc b)
     if (cmp (point_to_line (b.c, a), b.r) > 0) return std
    ::vector <point> ();
double x = sqrt (sqr (b.r) - sqr (point_to_line (b.c,
a)));
b); \overline{} 33 //Counter-clockwise with respect of point O_a.
if (pp.size () == 2 && qq.size () == 2) {
   if (cmp (a.r, b.r) < 0) std::swap (pp[0], pp[1]),
      std::swap (qq[0], qq[1]);
   ret.push_back (line (pp[0], qq[0]));
   ret.push_back (line (pp[1], qq[1])); }
return ret; }
//Counter-clockwise with respect of point Oa.</pre>
    //Counter-clockwise with respect of point O_a. std::vector <line> intangent (cc c1, cc c2) { point p = (b.c * a.r + a.c * b.r) / (a.r + b.r); std::vector pp = tangent (p, a), qq = tangent (pif (pp.size () == 2 && qq.size () == 2) { ret.push_back (line (pp[0], qq[0])); ret.push_back (line (pp[1], qq[1])); } return ret; }
                                                                                tangent (p, b);
```

## 5.4 Centers of a triangle

```
point incenter (cp a, cp b, cp c) {
   double p = dis (a, b) + dis (b, c) + dis (c, a);
   return (a * dis (b, c) + b * dis (c, a) + c * dis (a, b)) / p; }

   point circumcenter (cp a, cp b, cp c) {
    point p = b - a, q = c - a, s (dot (p, p) / 2, dot (q, q) / 2);
   return a + point (det (s, point (p.y, q.y)), det (point (p.x, q.x), s)) / det (p, q); }

   point orthocenter (cp a, cp b, cp c) { return a + b + c - circumcenter (a, b, c) * 2; }
```

## 5.5 Fermat point

```
/* Fermat point : finds a point P that minimizes |PA| + |PB| + |PC|. */
2 point fermat_point (cp a, cp b, cp c) {
3    if (a == b) return a; if (b == c) return b; if (c == a) return c;
4    double ab = dis (a, b), bc = dis (b, c), ca = dis (c, a);
5    double cosa = dot (b - a, c - a) / ab / ca;
6    double cosb = dot (a - b, c - b) / ab / bc;
7    double cosc = dot (b - c, a - c) / ca / bc;
8    double sq3 = PI / 3.0; point mid;
9    if (sgn (cosa + 0.5) < 0) mid = a;
10    else if (sgn (cosc + 0.5) < 0) mid = b;
11    else if (sgn (det (b - a, c - a)) < 0) mid = line_intersect (line (a, b + (c - b).rot (sq3)), line (b, c + (a - c).rot (sq3)));
13    else mid = line_intersect (line (a, c + (b - c).rot (sq3)));
14    return mid; }
```

## 5.6 Convex hull

//Counter-clockwise, with minimum number of points.
bool turn\_left (cp a, cp b, cp c) { return sgn (det (b - a, c - a)) >= 0; }

```
3 std::vector <point> convex_hull (std::vector <point> a
       int cnt = 0; std::sort (a.begin (), a.end ());
std::vector <point> ret (a.size () << 1, point ());
for (int i = 0; i < (int) a.size (); ++i) {
  while (cnt > 1 && turn_left (ret[cnt - 2], a[i], ret
        [cnt - 1])) --cnt;
  ret[cnt++] = a[i]; }
int fixed = cnt;
for (int i = (int) a.size () - 1; i >= 0; --i) {
  while (cnt > fixed && turn_left (ret[cnt - 2], a[i],
        ret[cnt - 1])) --cnt;
  ret[cnt++] = a[i]; }
return std::vector <point> (ret.begin (), ret.begin
  () + cnt - 1); }
```

## Half plane intersection

```
if (two_side (c[i], c[j], p)) ret.push_back (
    line_intersect (p, line (c[i], c[j]))); }
return ret; }
in /* Offline half plane intersection : complexity
             O(n\log n).
| D(n log n). */
| bool turn_left (cl 1, cp p) { return turn_left (l.s, 1 .t, p); }
| int cmp (cp a, cp b) { return a.dim () != b.dim () ? (
| a.dim () < b.dim () ? -1 : 1) : -sgn (det (a, b));
}
std::vector <point> half_plane_intersect (std::vector line> h) {
typedef std::pair <point, line> polar;
std::vector <point> g; g.resize (h.size ());
for (int i = 0; i < (int) h.size (); ++i) g[i] = std ::make_pair (h[i].t - h[i].s, h[i]);
sort (g.begin (), g.end (), [&] (const polar &a, const polar &b) {
   if (cmp (a.first, b.first) == 0) return sgn (det (a. second.t - a.second.s)) < 0;
     20
     line_intersect (ret[fore], ret[fore + 1]))) ++
fore;
if (rear - fore < 2) return std::vector <point> ();
std::vector <point> ans; ans.resize (rear + 1);
for (int i = 0; i < rear + 1; ++i) ans[i] =
    line_intersect (ret[i], ret[(i + 1) % (rear + 1)
]);
return ans; }</pre>
```

# 5.8 Nearest pair of points

```
1 /* Nearest pair of points : [1, r), need to sort p
print | first. */
2 double solve (std::vector <point> &p, int l, int r) {
3 if (l + 1 >= r) return INF;
4 int m = (l + r) / 2; double mx = p[m].x; std::vector
      );
for (int i = 1; i < r; ++i)
    if (sqr (p[i].x - mx) < ret) v.push_back (p[i]);
sort (v.begin (), v.end (), [&] (cp a, cp b) { return
        a.y < b.y; });
for (int i = 0; i < v.size (); ++i)
for (int j = i + 1; j < v.size (); ++j) {
    if (sqr (v[i].y - v[j].y) > ret) break;
    ret = min (ret, dis2 (v[i] - v[j])); }
return ret; }
```

### 5.9Minimum circle

```
circle minimum_circle (std::vector <point> p) {
  circle ret; std::random_shuffle (p.begin (), p.end ()
for (int i = 0; i < (int) p.size (); ++i) if (!
```

## 5.10 Intersection of a polygon and a circle

```
struct polygon_circle_intersect {
  double sector_area (cp a, cp b, const double &r) {
   double c = (2.0 * r * r - dis2 (a, b)) / (2.0 * r *
  r);
return r * r * acos (c) / 2.0; }
double area (cp a, cp b, const double &r) {
    double dA = dot (a, a), dB = dot (b, b), dC =
        point_to_segment (point (), line (a, b));
    if (sgn (dA - r * r) <= 0 && sgn (dB - r * r) <= 0)
        return det (a, b) / 2.0;
    point tA = a.unit () * r, tB = b.unit () * r;
    if (sgn (dC - r) > 0) return sector_area (tA, tB, r)
```

### Union of circles 5.11

```
| template <int MAXN = 500> struct union_circle {
  int C; circle c[MAXN]; double area[MAXN]
struct event {
   point p; double ang; int delta;
event (cp p = point (), double ang = 0, int delta =
    0) : p(p), ang(ang), delta(delta) {}
bool operator < (const event &a) { return ang < a.</pre>
   bool ope:
 void addevent(cc a, cc b, std::vector <event> &evt,
     intersect (c[i], c[j]))
addevent (c[i], c[j], evt, cnt);
if (evt.empty ()) area[cnt] += PI * c[i].r * c[i].r
```

## 5.12 3D point

## 5.13 3D line

## 5.14 3D convex hull

## 6 Graph

## 6.1 Hopcoft-Karp algorithm

```
/* Hopcoft-Karp algorithm : unweighted maximum matching for bipartition graphs with complexity O(m\sqrt{n}). */
2 template <int MAXN = 100000, int MAXM = 100000>
3 struct hopcoft_karp {
4 using edge_list = std::vector <int> [MAXN];
5 int mx[MAXN], my[MAXM], lv[MAXN];
6 bool dfs (edge_list <MAXN, MAXM> &e, int x) {
```

## 6.2 Kuhn-Munkres algorithm

```
/* Kuhn Munkres algorithm : weighted maximum ming algorithm for bipartition graphs with complexity O(N^3).

Note: The graph is 1-based. */

template <int MAXN = 500>

template <int MAXN = 500>

template <int MAXN] [MAXN], lx[MAXN], ly[MAXN], m[MAXN],

way[MAXN], sl[MAXN];

bool u[MAXN];

void hungary(int x) {

m[0] = x; int j0 = 0;

std::fill (sl, sl + n + 1, INF); std::fill (u, u + n + 1, false);

do {

u[j0] = true; int i0 = m[j0], d = INF, j1 = 0;

for (int j = 1; j <= n; ++j)

if (u[j] == false) {

int cur = -w[i0][j] - lx[i0] - ly[j];

if (sl[j] < d) { d = sl[j]; j1 = j; } }

for (int j = 0; j <= n; ++j) {

if (u[j]) { lx[m[j]] += d; ly[j] -= d; }

else sl[j] -= d; }

j0 = j1; } while (m[j0] != 0);

do {

int j1 = way[j0]; m[j0] = m[j1]; j0 = j1; }

y while (j0); }

int solve() {

for (int i = 1; i <= n; ++i) m[i] = lx[i] = ly[i] = way[i] = 0;

for (int i = 1; i <= n; ++i) hungary (i);

int sum = 0; for (int i = 1; i <= n; ++i) sum += w[m [i]][i];

return sum; } };
```

## 6.3 Blossom algorithm

```
if (d[dest] == -1)
if (match[dest] == -1) {
  solve (root, loc); match[loc] = dest;
  match[dest] = loc; return 1;
}
  match[dest] = loc; return 1;
} else {
  fa[dest] = loc; fa[match[dest]] = dest;
  d[dest] = 1; d[match[dest]] = 0;
  *qtail++ = match[dest];
} else if (d[ufs.find (dest)] == 0) {
  int b = lca (loc, dest, root);
  contract (loc, dest, b); contract (dest, loc, b)
  return 0; }
int solve (int n cost)
int solve (int n, const edge_list &e) {
  std::fill (fa, fa + n, 0); std::fill (c1, c1 + n, 0)
   std::fill (c2, c2 + n, 0); std::fill (match, match +
n, -1);
int re = 0; for (int i = 0; i < n; i++)
if (match[i] == -1) if (bfs (i, n, e)) ++re; else
    match[i] = -2;
return re; } };</pre>
```

## Weighted blossom algorithm

```
1 /* Weighted blossom algorithm (vfleaking ver.) : maximum matching for general weighted graphs with
  complexity O(n^3). 2 Usage : Set n to the size of the vertices. Run init () . Set g[][].w to the weight of the edge. Run solve
  The first result is the answer, the second one is the number of matching pairs. Obtain the matching with
  match[].
4 Note: 1-based.
  4 Note: 1-Dased. */
struct weighted_blossom {
static const int INF = INT_MAX, MAXN = 400;
struct edge{ int u, v, w; edge (int u = 0, int v = 0,
    int w = 0): u(u), v(v), w(w) {} };
        int n, n_x;
edge g[MAXN * 2 + 1][MAXN * 2 + 1];
int lab[MAXN * 2 + 1], match[MAXN * 2 + 1], slack[
MAXN * 2 + 1], st[MAXN * 2 + 1], pa[MAXN * 2 +
        std::vector <int> flower[MAXN * 2 + 1]; std::queue <
    int> q;
int e_delta (const edge &e) { return lab[e.u] + lab[e
    .v] - g[e.u][e.v].w * 2; }
void update_slack (int u, int x) { if (!slack[x] ||
    e_delta (g[u][x]) < e_delta (g[slack[x]][x]))
    slack[x] = u; }
void set_slack (int x) { slack[x] = 0; for (int u =
    1; u <= n; ++u) if (g[u][x].w > 0 && st[u] != x &&
    S[st[u]] == 0)
    update_slack(u, x); }
void a push (int x) {
      21
26
28
29

    void augment (int u, int v) {
    for (;;) {
        int xnv = st[match[u]]; set_match (u, v);
        if (!xnv) return; set_match (xnv, st[pa[xnv]]);
        u = st[pa[xnv]], v = xnv; }

int get_lca (int u, int v) {
    static int t = 0;
    for (++t; u || v; std::swap (u, v)) {
        if (u == 0) continue; if (vis[u] == t) return u;
        vis[u] = t; u = st[match[u]]; if (u) u = st[pa[u]];

    return 0;
}

return 0;
}

void add blossom (int u, int lca, int v) {

        return 0; }
void add_blossom (int u, int lca, int v) {
  int b = n + 1; while (b <= n_x && st[b]) ++b;
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca]; flower[b].clear ();
  flower[b].push_back (lca);
  for (int x = n_x x = lca; x = st[na[v]]) {
           (i);

for (int x = v, y; x != lca; x = st[pa[y]]) {

flower[b].push_back (x), flower[b].push_back (y = st[match[x]]), q_push(y); }
            set_st (b, b);

for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].w

= 0;
           for (int x = 1; x <= n; ++x) flower_from[b][x] = 0;
for (size_t i = 0; i < flower[b].size (); ++i){
  int xs = flower[b][i];
  for (int x = 1; x <= n, ++x) if (g[b][x].w == 0</pre>
                           (int x = 1; x <= n_x; ++x) if (g[b][x].w == 0
|| e_delta(g[xs][x]) < e_delta(g[b][x]))
```

```
60
63
         tut xr = riower_from[b][g[b][pa[b]].u], pr = ge
    b, xr);
for (int i = 0; i < pr; i += 2) {
  int xs = flower[b][i], xns = flower[b][i + 1];
  pa[xs] = g[xns][xs].u; S[xs] = 1, S[xns] = 0;
  slack[xs] = 0, set_slack(xns); q_push(xns); }
s[xr] = 1, pa[xr] = pa[b];</pre>
       S[xr] = 1, pa[xr] = pa[b];
for (size_t i = pr + 1; i < flower[b].size (); ++i)
       {
  int xs = flower[b][i]; S[xs] = -1, set_slack(xs); }
st[b] = 0; }
pool on_found_edge (const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1; int nu = st[match[v]];
    slack[v] = slack[nu] = 0; S[nu] = 0, q_push(nu);
} else if(S[v] == 0) {
  int lca = get_lca(u, v);
  if (!lca) return augment(u, v), augment(v, u), true
    6.5
              Maximum flow
```

87

99 100

115

120

```
/* Sparse graph maximum flow : isap.*/
template <int MAXN = 1000, int MAXM = 100000>
      struct isap {
        f; begin[u] = size++;
dest[size] = u; next[size] = begin[v]; flow[size] =
    0; begin[v] = size++; };
int pre[MAXN] d[MAXN], gap[MAXN], cur[MAXN];
int solve (flow_edge_list &e, int n, int s, int t) {
  for (int i = 0; i < n; ++i) { pre[i] = d[i] = gap[i] }
    = 0; cur[i] = e.begin[i]; }
  gap[0] = n; int u = pre[s] = s, v, maxflow = 0;
  while (d[s] < n) {
    v = n; for (int i = cur[u]; ~i; i = e.next[i])</pre>
13
```

```
if (e.flow[i] && d[u] == d[e.dest[i]] + 1) {
    v = e.dest[i]; cur[u] = i; break; }
if (v < n) {
    pre[v] = u; u = v;
    if (v == t) {
        int dflow = INF, p = t; u = s;
        while (p != s) { p = pre[p]; dflow = std::min (
            dflow, e.flow[cur[p]]); }
    maxflow += dflow; p = t;
    while (p != s) { p = pre[p]; e.flow[cur[p]] -=
            dflow; e.flow[cur[p] ^ 1] += dflow; }
} else {</pre>
    19
   22
   25
                         else {
int mindist = n + 1;
  int mindist = n + 1;
for (int i = e.begin[u]; ~i; i = e.next[i])
if (e.flow[i] && mindist > d[e.dest[i]]) {
    mindist = d[e.dest[i]]; cur[u] = i; }
if (!--gap[d[u]]) return maxflow;
gap[d[u] = mindist + 1]++; u = pre[u]; } }
return maxflow; };
template <int MAXN = 1000, int MAXM = 100000>
el struct dinic {
n = n_; s = s_; dinic::t = t_;
                 while (bfs (e)) {
  for (int i = 0; i < n; ++i) w[i] = e.begin[i];
  ans += dfs (e, s, INF); }
return ans; };</pre>
```

### 6.6Minimum cost flow

```
/* Sparse graph minimum cost flow : EK. */
template <int MAXN = 1000, int MAXM = 100000>
struct minimum_cost_flow {
    struct cost_flow_edge_list {
    int size, begin[MAXN], dest[MAXM], next[MAXM], cost[
        MAXM], flow[MAXXM];
    void clear (int n) { size = 0; std::fill (begin, begin + n, -1); }
    cost_flow_edge_list (int n = MAXN) { clear (n); }
    void add_edge (int u, int v, int c, int f) {
    dest[size] = v; next[size] = begin[u]; cost[size] = c; flow[size] = f; begin[u] = size++;
    dest[size] = u; next[size] = begin[v]; cost[size] = -c; flow[size] = 0; begin[v] = size++; };
    int n, s, t, prev[MAXN], dist[MAXN], occur[MAXN];
    bool augment (cost_flow_edge_list &e) {
        std::vector <int> queue;
        std::fill (dist, dist + n, INF); std::fill (occur, occur + n, 0);
    dist[s] = 0; occur[s] = true; queue.push_back (s);
    for (int head = 0; head < (int) queue.size(); ++head)
    int x = queue[head]:</pre>
16
            ans.second += num * e.cost[prev[i]]; } }
as return ans; } ;;
37 /* Dense graph minimum cost flow: zkw. */
38 template <int MAXN = 1000, int MAXM = 100000>
39 struct zkw_flow {
            struct cost_flow_edge_list {
  int size, begin[MAXN], dest[MAXM], next[MAXM], cost[
     MAXM], flow[MAXM];
```

```
int modlable() {
  int delta = INF;
  for (int i = 0; i < n; i++) {
    if (!visit[i] && slack[i] < delta) delta = slack[i]</pre>
51
    slack[i] = INF; }
if (delta == INF) return 1;
for (int i = 0; i < n; i++) if (visit[i]) dis[i] +=
    delta;
return 0; }</pre>
  e.flow[i] -= delta; e.flow[i ^ 1] += delta; left
__= delta;
  if (!left) { visit[x] = false; return flow; }
   } };
```

## Stoer Wagner algorithm

## 6.8 DN maximum clique

```
_{1} /* DN maximum clique : n <= 150 *.
     1/* DN Maximum Clique : n - 150 %
2 typedef bool BB[N]; struct Maxclique {
3 const BB *e; int pk, level; const float Tlimit;
4 struct Vertex { int i, d; Vertex (int i) : i(i), d(0)
{ } };
5 typedef std::vector <Vertex> Vertices; Vertices V;
6 typedef std::vector <int> ColorClass; ColorClass QMAX,
          Q;
std::vector <ColorClass> C;
    | 15| v[i].d += e[v[i].i][v[j].i]; }
| 16| struct StepCount { int i1, i2; StepCount(): i1 (0), (0) {} }; |
| 17| std::vector <StepCount> S; |
| 18| bool cut1 (const int pi, const ColorClass &A) { |
| 19| for (int i = 0; i < (int) A.size (); ++i) |
| 20| if (e[pi][A[i]]) return true; return false; } |
| 21| void cut2 (const Vertices &A, Vertices & B) { |
| 22| for (int i = 0; i < (int) A.size () - 1; ++i) |
| 23| if (e[A.back().i][A[i].i]) B.push_back(A[i].i); }
```

## 6.9 Dominator tree

# 7 String

# 7.1 Suffix Array

```
/* Suffix Array : sa[i] - the beginning position of the ith smallest suffix, rk[i] - the rank of the suffix beginning at position i. height[i] - the longest common prefix of sa[i] and sa[i - 1]. */
```

Page 10

## 7.2 Suffix Automaton

## 7.3 Palindromic tree

```
/* Palindromic tree : extend () - returns whether the tree has generated a new node. odd, even - the root of two trees. last - the node representing the last char. node::len - the palindromic string length of the node. */

template <int MAXN = 1000000, int MAXC = 26>

struct palindromic_tree {

struct node {

node *child[MAXC], *fail; int len;

node (int len) : fail (NULL), len (len) {

memset (child, NULL, sizeof (child)); }

node_pool[MAXN * 2], *tot_node;

int size, text[MAXN];

node *match (node *now) {

for (; text[size - now -> len - 1] != text[size];

now = now -> fail);

return now; }

bool extend (int token) {

text[++size] = token; node *now = match (last);

if (now -> child[token])

return last = now -> child[token], false;

last = now -> child[token] = new (tot_node++) node (

now -> len + 2);

if (now == odd) last -> fail = even;

else {

now = match (now -> fail);
```

## 7.4 Regular expression

```
std::string str = ("The_the_there");
std::regex pattern ("(th|Th)[\\w]*", std::
    regex_constants::optimize | std::regex_constants::
    ECMAScript);
std::smatch match; //std::cmatch for char *

std::regex_match (str, match, pattern);
auto mbegin = std::sregex_iterator (str.begin (), str.end (), pattern);
sauto mend = std::sregex_iterator ();
std::cout << "Found_" << std::distance (mbegin, mend) << "_words:\n";
for (std::sregex_iterator i = mbegin; i != mend; ++i)
{
    match = *i;
    /* The word is match[0], backreferences are match[i]
    up to match.size ().
match.prefix () and match.suffix () give the prefix and the suffix.
match.length () gives length and match.position ()
    gives position of the match. */ }
std::regex_replace (str, pattern, "sh$1");
if //$n is the backreference, $& is the entire match, $`
    is the prefix, $' is the suffix, $$is the $ sign.</pre>
```

# 8 Tips8.1 Java

```
1 /* Java reference : References on Java IO, structures,
    etc. */
2 import java.io.*;
3 import java.lang.*;
    | import java.math.*;
| import java.util.*;
    doubleValue () / toPlainString () : converts to other types.

Nrrays: Arrays.sort (T [] a); Arrays.sort (T [] a, int fromIndex, int toIndex); Arrays.sort (T [] a, int fromIndex, int toIndex, Comperator <? super T>
 int fromIndex, int toIndex, Comperator <? super T>
    comperator);
21 LinkedList <E> : addFirst / addLast (E) / getFirst /
    getLast / removeFirst / removeLast () / clear () /
    add (int, E) / remove (int) / size () / contains
    / removeFirstOccurrence / removeLastOccurrence (E)
22 ListIterator <E> listIterator (int index) : returns an
    iterator :
            E next / previous () : accesses and iterates.
hasNext / hasPrevious () : checks availablity.
nextIndex / previousIndex () : returns the index of a
            hasNext
 25  nextIndex / previousIndex () : returns the index of a
    subsequent call.
26  add / set (E) / remove () : changes element.
27  PriorityQueue <E> (int initcap, Comparator <? super E>
        comparator) : add (E) / clear () / iterator () /
        peek () / poll () / size ()
28  TreeMap <K, V> (Comparator <? super K> comparator) :
        Map.Entry <K, V> ceilingEntry / floorEntry /
        higherEntry / lowerEntry (K): getKey / getValue ()
        / setValue (V) : entries.
29  clear () / put (K, V) / get (K) / remove (K) / size
        ()
 ()
30 StringBuilder: StringBuilder (string) / append (int, string, ...) / insert (int offset, ...) charAt ( int) / setCharAt (int, char) / delete (int, int) / reverse () / length () / toString ()
31 String: String.format (String, ...) / toLowerCase / toUpperCase () */
32 /* Examples on Comparator:
33 public class Main {
```

```
public int x; public int y;
public Point () {
         x = 0;
y = 0; }
public Point (int xx, int yy) {
x = xx;
       y = yy; } };
public static class Cmp implements Comparator <Point>
      46
         Point> {
    public int x; public int y;
    public Point () {
        x = 0;
        y = 0; }
    public Point ()
      public static class Point implements Comparable <
55
         public Point (int xx, int yy) {
  x = xx;
         x = xx;
y = yy; }
public int compareTo (Point p) {
if (x < p.x) return -1;
if (x == p.x) {
  if (y < p.y) return -1;
  if (y == p.y) return 0; }
return 1; }
public boolean equalTo (Point p)</pre>
         public boolean equalTo (Point p) {
  return (x == p.x && y == p.y); }
public int hashCode () {
  return x + y; } };
73| */
74| //Faster IO :
hasMoreTokens()) {

try {
   String line = reader.readLine();
   tokenizer = new StringTokenizer (line);
   } catch (IOException e) {
    throw new RuntimeException (e); }

return tokenizer.nextToken(); }

public BigInteger nextBigInteger() {
   return new BigInteger (next (), 10); /* radix */ }

public int nextInt() {
   return Integer.parseInt (next()); }

public double nextDouble() {
   return Double.parseDouble (next()); }

public static void main (String[] args) {
   InputReader in = new InputReader (System.in);
}
                     Random numbers
```

public static class Point {

```
std::mt19937_64 mt (time (0));
std::uniform_int_distribution <int> uid (1, 100);
std::uniform_real_distribution <double> urd (1, 100);
std::cout << uid (mt) << "_" << urd (mt) << "\n";
```

## 8.3 Read hack

## 8.4 Stack hack

```
1 //C++
2 #pragma comment (linker, "/STACK:36777216")
3 //G++
4 int __size__ = 256 << 20;
5 char * _p_ = (char*) malloc(__size__) + __size__;
6 __asm__ ("movl_%0,_%%esp\n" :: "r"(__p__));</pre>
```

## 8.5 Time hack

```
1 clock_t t = clock ();
2 std::cout << 1. * t / CLOCKS_PER_SEC << "\n";</pre>
```

## Multiplication hack

long long mul\_mod (long long x, long long y, long long
mod) {
long long t = (x \* y - (long long) ((long double) x /
mod \* y + 1E-3) \* mod) % mod;
return t < 0 ? t + mod : t; }</pre>

### 8.7 Builtin functions

\_builtin\_clz: Returns the number of leading 0-bits in x, starting at the most significant bit position. If x is 0, the result is

undefined.\_\_builtin\_ctz: Returns the number of trailing 0-bits in x, starting at the least significant bit position. If x is 0, the result is

undefined.
\_builtin\_clrsb: Returns the number of leading redundant sign
\_builtin\_clrsb: Returns the number of leading redundant sign
for following the most significant bits in x, i.e. the number of bits following the most significant bit that are identical to it. There are no special cases for 0 or

other values.
\_builtin\_popcount: Returns the number of 1-bits in x.
\_builtin\_parity: Returns the parity of x, i.e. the number of

1-bits in x modulo 2. \_builtin.bswap16, \_builtin.bswap32, \_builtin.bswap64: Returns x with the order of the bytes (8 bits as a group) reversed.

 $\begin{array}{ccc} \text{bitset::Find\_first(), bitset::Find\_next(idx): bitset built-in functions.} \\ Prufer sequence \end{array}$ 

## 8.8

In combinatorial mathematics, the Prufer sequence of a labeled tree is a unique sequence associated with the tree. The sequence for a tree on n vertices has length n-2.

One can generate a labeled tree's Prufer sequence by iteratively removing vertices from the tree until only two vertices remain. Specifically, consider a labeled tree T with vertices 1, 2, ..., n. At step i, remove the leaf with the smallest label and set the *i*th element of the Prufer sequence to be the label of this leaf's neighbour.

One can generate a labeled tree from a sequence in three steps. The tree will have n+2 nodes, numbered from 1 to n+2. For each node set its degree to the number of times it appears in the sequence plus 1. Next, for each number in the sequence a[i], find the first (lowestnumbered) node, j, with degree equal to 1, add the edge (j, a[i]) to the tree, and decrement the degrees of j and a[i]. At the end of this loop two nodes with degree 1 will remain (call them u, v). Lastly, add the edge (u, v) to the tree.

The Prufer sequence of a labeled tree on n vertices is a unique sequence of length n-2 on the labels 1 to n - this much is clear. Somewhat less obvious is the fact that for a given sequence S of length n-2 on the labels 1 to n, there is a unique labeled tree whose Prufer sequence

Spanning tree counting

Kirchhoff's Theorem: the number of spanning trees in a graph G is equal to any cofactor of the Laplacian matrix of G, which is equal to the difference between the graph's degree matrix (a diagonal matrix with vertex degrees on the diagonals) and its adjacency matrix (a (0,1)matrix with 1's at places corresponding to entries where the vertices are adjacent and 0's otherwise).

The number of edges with a certain weight in a minimum spanning tree is fixed given a graph. Moreover, the number of its arrangements can be obtained by finding a minimum spanning tree, compressing connected components of other edges in that tree into a point, and then applying Kirrchoff's theorem with only edges of the certain weight in the graph. Therefore, the number of minimum spanning trees in a graph can be solved by multiplying all numbers of arrangements of edges of different weights together.

### 8.10 Mobius inversion

### 8.10.1 Mobius inversion formula

$$[x = 1] = \sum_{d|x} \mu(d)$$
$$x = \sum_{d|x} \mu(d)$$

### 8.10.2Gcd inversion

$$\sum_{a=1}^{n} \sum_{b=1}^{n} gcd^{2}(a,b) = \sum_{d=1}^{n} d^{2} \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [gcd(i,j) = 1]$$

$$= \sum_{d=1}^{n} d^{2} \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{t \mid gcd(i,j)} \mu(t)$$

$$= \sum_{d=1}^{n} d^{2} \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} [t \mid i] \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [t \mid j]$$

$$= \sum_{d=1}^{n} d^{2} \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \lfloor \frac{n}{dt} \rfloor^{2}$$

The formula can be computed in O(nlogn) complexity.

$$\sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \big\lfloor \frac{n}{dt} \big\rfloor^2 = \sum_{l=1}^n \big\lfloor \frac{n}{l} \big\rfloor^2 \sum_{d|l} d^2 \mu(\frac{l}{d})$$

Let  $f(l) = \sum_{d|l} d^2 \mu(\frac{l}{d})$ . It can be proven that f(l) is multiplicative. Besides,  $f(p^k) = p^{2k} - p^{2k-2}$ .

Therefore, with linear sieve the formula can be computed in O(n)complexity.

### 2-SAT 8.11

In terms of the implication graph, two literals belong to the same strongly connected component whenever there exist chains of implications from one literal to the other and vice versa. Therefore, the two literals must have the same value in any satisfying assignment to the given 2-satisfiability instance. In particular, if a variable and its negation both belong to the same strongly connected component, the instance cannot be satisfied, because it is impossible to assign both of these literals the same value. As Aspvall et al. showed, this is a necessary and sufficient condition: a 2-CNF formula is satisfiable if and only if there is no variable that belongs to the same strongly connected component as its negation.

This immediately leads to a linear time algorithm for testing satisfiability of 2-CNF formulae: simply perform a strong connectivity analysis on the implication graph and check that each variable and its negation belong to different components. However, as Aspvall et al. also showed, it also leads to a linear time algorithm for finding a satisfying assignment, when one exists. Their algorithm performs the following steps

Construct the implication graph of the instance, and find its strongly connected components using any of the known linear-time algorithms for strong connectivity analysis.

Check whether any strongly connected component contains both a variable and its negation. If so, report that the instance is not satisfiable

Construct the condensation of the implication graph, a smaller graph that has one vertex for each strongly connected component, and an edge from component i to component j whenever the implication graph contains an edge uv such that u belongs to component i and v belongs to component j. The condensation is automatically a directed acyclic graph and, like the implication graph from which it was formed, it is skew-symmetric.

Topologically order the vertices of the condensation. In practice this may be efficiently achieved as a side effect of the previous step, as components are generated by Kosaraju's algorithm in topological order and by Tarjan's algorithm in reverse topological order.

For each component in the reverse topological order, if its variables do not already have truth assignments, set all the literals in the component to be true. This also causes all of the literals in the complementary component to be set to false.

### 8.12Numbers

## 8.12.1 Binomial Coefficients

$$\binom{n}{k} = (-1)^k \binom{k-n-1}{k}, \quad \sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$$

$$\sum_{k=0}^n \binom{k}{m} = \binom{n+1}{m+1}$$

$$\sqrt{1+z} = 1 + \sum_{k=1}^\infty \frac{(-1)^{k-1}}{k \times 2^{2k-1}} \binom{2k-2}{k-1} z^k$$

$$\sum_{k=0}^r \binom{r-k}{m} \binom{s+k}{n} = \binom{r+s+1}{m+n+1}$$

$$C_{n,m} = \binom{n+m}{m} - \binom{n+m}{m-1}, n \ge m$$

$$\binom{n}{k} \equiv [n\&k = k] \pmod{2}$$

$$\binom{n_1 + \dots + n_p}{m} = \sum_{k_1 + \dots + k_n = m} \binom{n_1}{k_1} \dots \binom{n_p}{k_p}$$

## Fibonacci Numbers

$$F(z) = \frac{z}{1 - z - z^2}$$

$$f_n = \frac{\phi^n - \hat{\phi}^n}{\sqrt{5}}, \phi = \frac{1 + \sqrt{5}}{2}, \hat{\phi} = \frac{1 - \sqrt{5}}{2}$$

$$\sum_{k=1}^n f_k = f_{n+2} - 1, \quad \sum_{k=1}^n f_k^2 = f_n f_{n+1}$$

$$\sum_{k=0}^n f_k f_{n-k} = \frac{1}{5}(n-1)f_n + \frac{2}{5}nf_{n-1}$$

$$\frac{f_{2n}}{f_n} = f_{n-1} + f_{n+1}$$

$$f_1 + 2f_2 + 3f_3 + \dots + nf_n = nf_{n+2} - f_{n+3} + 2]$$

$$\gcd(f_m, f_n) = f_{\gcd(m, n)}$$

$$f_n^2 + (-1)^n = f_{n+1}f_{n-1}$$

$$f_{n+k} = f_n f_{k+1} + f_{n-1}f_k$$

$$f_{2n+1} = f_n^2 + f_{n+1}^2$$

$$(-1)^k f_{n-k} = f_n f_{k-1} - f_{n-1}f_k$$

Modulo 
$$f_n, f_{mn+r} \equiv \begin{cases} f_r, & m \bmod 4 = 0; \\ (-1)^{r+1} f_{n-r}, & m \bmod 4 = 1; \\ (-1)^n f_r, & m \bmod 4 = 2; \\ (-1)^{r+1+n} f_{n-r}, & m \bmod 4 = 3. \end{cases}$$

## 8.12.3 Lucas Numbers

$$L_0 = 2, L_1 = 1, L_n = L_{n-1} + L_{n-2} = \left(\frac{1+\sqrt{5}}{2}\right)^n + \frac{1-\sqrt{5}}{2}\right)^n$$
  
$$L(x) = \frac{2-x}{1-x-x^2}$$

### 8.12.4Catlan Numbers

$$c_1 = 1, c_n = \sum_{i=0}^{n-1} c_i c_{n-1-i} = c_{n-1} \frac{4n-2}{n+1} = \frac{\binom{2n}{n}}{n+1}$$
$$= \binom{2n}{n} - \binom{2n}{n-1}, c(x) = \frac{1-\sqrt{1-4x}}{2x}$$

### Stirling Cycle Numbers 8.12.5

Divide n elements into k non-empty cycles.

$$\begin{split} s(n,0) &= 0, s(n,n) = 1, s(n+1,k) = s(n,k-1) - ns(n,k) \\ & s(n,k) = (-1)^{n-k} {n \brack k} \\ & {n \brack k} = n {n \brack k} + {n \brack k-1}, {n+1 \brack 2} = n! H_n \\ & x^{\underline{n}} = x(x-1)...(x-n+1) = \sum_k {n \brack k} (-1)^{n-k} x^k \\ & x^{\overline{n}} = x(x+1)...(x+n-1) = \sum_k {n \brack k} x^k \end{split}$$

## Stirling Subset Numbers

Divide n elements into k non-empty subsets.

For a fixed k, generating functions

$$\sum_{n=0}^{\infty} {n \brace k} x^{n-k} = \prod_{r=1}^{k} \frac{1}{1-rx}$$

### Motzkin Numbers

Draw non-intersecting chords between n points on a circle.

Pick n numbers  $k_1, k_2, ..., k_n \in \{-1, 0, 1\}$  so that  $\sum_{i=1}^{n} k_i (1 \le a \le n)$ is non-negative and the sum of all numbers is 0.

$$M_{n+1} = M_n + \sum_{i=0}^{n-1} M_i M_{n-1-i} = \frac{(2n+3)M_n + 3nM_{n-1}}{n+3}$$

$$M_n = \sum_{i=0}^{\lfloor \frac{n}{2} \rfloor} {n \choose 2k} Catlan(k)$$

$$M(X) = \frac{1 - x - \sqrt{1 - 2x - 3x^2}}{2x^2}$$

### 8.12.8 **Eulerian Numbers**

Permutations of the numbers 1 to n in which exactly k elements are

### Harmonic Numbers 8.12.9

Sum of the reciprocals of the first n natural numbers.

$$\sum_{k=1}^{n} H_k = (n+1)H_n - n$$

$$\sum_{k=1}^{n} kH_k = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}$$

$$\sum_{k=1}^{n} {k \choose m} H_k = {n+1 \choose m+1} (H_{n+1} - \frac{1}{m+1})$$

### 8.12.10Pentagonal Number Theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = \sum_{n=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2}$$
$$p(n) = p(n-1) + p(n-2) - p(n-5) - p(n-7) + \cdots$$

 $f(n,k) = p(n) - p(n-k) - p(n-2k) + p(n-5k) + p(n-7k) - \cdots$ 

8.12.11 Bell Numbers

Divide a set that has exactly n elements. 
$$B_n=\sum_{k=1}^n {n\brace k},\ B_{n+1}=\sum_{k=0}^n {n\choose k}B_k$$

$$B_{p^m+n} \equiv mB_n + B_{n+1} \pmod{p}$$

$$B(x) = \sum_{n=0}^{\infty} \frac{B_n}{n!} x^n = e^{e^x - 1}$$

### Bernoulli Numbers 8.12.12

$$B_n = 1 - \sum_{k=0}^{n-1} {n \choose k} \frac{B_k}{n-k+1}$$

$$G(x) = \sum_{k=0}^{\infty} \frac{B_k}{k!} x^k = \frac{1}{\sum_{k=0}^{\infty} \frac{x^k}{(k+1)!}}$$

$$\sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k n^{m-k+1}$$

### 8.12.13

$$\begin{split} \sum_{i=1}^n i^2 &= \frac{n(n+1)(2n+1)}{6}, \quad \sum_{i=1}^n i^3 = (\frac{n(n+1)}{2})^2 \\ &\sum_{i=1}^n i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} \\ &\sum_{i=1}^n i^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12} \end{split}$$

## 8.12.14 Sum of Squares

Denote  $r_k(n)$  the ways to form n with k squares. If:

$$n = 2^{a_0} p_1^{2a_1} \cdots p_r^{2a_r} q_1 b_1 \cdots q_s b_s$$

where  $p_i \equiv 3 \mod 4$ ,  $q_i \equiv 1 \mod 4$ , then

$$r_2(n) = \begin{cases} 0 & \text{if any } a_i \text{ is a half-integer} \\ 4\prod_{i=1}^{r}(b_i+1) & \text{if all } a_i \text{ are integers} \end{cases}$$

 $r_3(n) > 0$  when and only when n is not  $4^a(8b+7)$ .

## 8.12.15 Derangement

$$D_1 = 0, D_2 = 1, D_n = n! \left( \frac{1}{0!} - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + \frac{(-1)^n}{n!} \right)$$
$$D_n = (n-1)(D_{n-1} + D_{n-2})$$

## 8.12.16 Tetrahedron Volume

If U, V, W, u, v, w are lengths of edges of the tetrahedron (first three form a triangle; u opposite to U and so on)

$$V = \frac{\sqrt{4u^2v^2w^2 - \sum_{cyc} u^2(v^2 + w^2 - U^2)^2 + \prod_{cyc} (v^2 + w^2 - U^2)}}{12}$$

### 9 **Appendix** 9.1 Calculus table

1. 
$$\int \frac{x}{ax+b} dx = \frac{1}{2}(ax+b-b\ln|ax+b|) + C$$

$$\begin{array}{ll} \textbf{9.1.1} & ax+b \ \left(a\neq 0\right) \\ \text{1.} & \int \frac{x}{ax+b} \, \mathrm{d}x = \frac{1}{a^2} (ax+b-b \ln |ax+b|) + C \\ \text{2.} & \int \frac{x^2}{ax+b} \, \mathrm{d}x = \frac{1}{a^3} \left(\frac{1}{2} (ax+b)^2 - 2b(ax+b) + b^2 \ln |ax+b|\right) + C \\ \text{3.} & \int \frac{\mathrm{d}x}{x(ax+b)} = -\frac{1}{b} \ln \left|\frac{ax+b}{x}\right| + C \\ \text{4.} & \int \frac{\mathrm{d}x}{x^2(ax+b)} = -\frac{1}{bx} + \frac{a}{b^2} \ln \left|\frac{ax+b}{x}\right| + C \end{array}$$

3. 
$$\int \frac{\mathrm{d}x}{\pi(ax+b)} = -\frac{1}{b} \ln \left| \frac{ax+b}{x} \right| + C$$

4. 
$$\int \frac{\mathrm{d}x}{x^2(ax+b)} = -\frac{1}{bx} + \frac{a}{b^2} \ln \left| \frac{ax+b}{x} \right| + C$$

Luna's Magic Reference

5. 
$$\int \frac{x}{(ax+b)^2} dx = \frac{1}{a^2} \left( \ln|ax+b| + \frac{b}{ax+b} \right) + C$$
6. 
$$\int \frac{x}{(ax+b)^2} dx = \frac{1}{a^3} \left( ax+b-2b \ln|ax+b| - \frac{b^2}{ax+b} \right) + C$$
7. 
$$\int \frac{dx}{(ax+b)^2} = \frac{1}{b(ax+b)} - \frac{1}{b^2} \ln\left|\frac{ax+b}{x}\right| + C$$
9.1.2 
$$\sqrt{ax+b}$$
1. 
$$\int \sqrt{ax+b} dx = \frac{2}{3a} \sqrt{(ax+b)^3} + C$$
2. 
$$\int x\sqrt{ax+b} dx = \frac{2}{15a^2} (3ax-2b) \sqrt{(ax+b)^3} + C$$
3. 
$$\int x^2 \sqrt{ax+b} dx = \frac{2}{15a^3} (15a^2x^2 - 12abx + 8b^2) \sqrt{(ax+b)^3} + C$$
4. 
$$\int \frac{x}{\sqrt{ax+b}} dx = \frac{2}{3a^2} (ax-2b) \sqrt{ax+b} + C$$
5. 
$$\int \frac{x^2}{\sqrt{ax+b}} dx = \frac{2}{15a^3} (3a^2x^2 - 4abx + 8b^2) \sqrt{ax+b} + C$$
6. 
$$\int \frac{dx}{x\sqrt{ax+b}} = \begin{cases} \frac{1}{\sqrt{b}} \ln\left|\frac{\sqrt{ax+b-\sqrt{b}}}{\sqrt{ax+b+\sqrt{b}}}\right| + C \quad (b > 0) \\ \frac{2}{\sqrt{-b}} \arctan \sqrt{\frac{ax+b}{-b}} + C \quad (b < 0) \end{cases}$$
7. 
$$\int \frac{dx}{x\sqrt{ax+b}} dx = \frac{2}{x\sqrt{ax+b}} dx = \frac{2}{x\sqrt{ax+b}} + \frac{1}{b} \int \frac{dx}{\sqrt{ax+b}}$$
8. 
$$\int \frac{\sqrt{ax+b}}{\sqrt{ax+b}} dx = 2\sqrt{ax+b} + b \int \frac{dx}{\sqrt{ax+b}}$$

5. 
$$\int \frac{\sqrt{x^2}}{\sqrt{ax+b}} dx = \frac{2}{15a^3} (3a^2x^2 - 4abx + 8b^2) \sqrt{ax+b} + C$$

6. 
$$\int \frac{\mathrm{d}x}{x\sqrt{ax+b}} = \begin{cases} \frac{1}{\sqrt{b}} \ln \left| \frac{\sqrt{ax+b} - \sqrt{b}}{\sqrt{ax+b} + \sqrt{b}} \right| + C & (b > 0) \\ \frac{2}{\sqrt{-b}} \arctan \sqrt{\frac{ax+b}{-b}} + C & (b < 0) \end{cases}$$

8. 
$$\int \frac{\sqrt{ax+b}}{x} dx = 2\sqrt{ax+b} + b \int \frac{dx}{x\sqrt{ax+b}}$$

9. 
$$\int \frac{\sqrt{ax+b}}{x^2} dx = -\frac{\sqrt{ax+b}}{x} + \frac{a}{2} \int \frac{dx}{x\sqrt{ax+b}}$$

**9.1.3** 
$$x^2 \pm a^2$$

1. 
$$\int \frac{\mathrm{d}x}{x^2 + a^2} = \frac{1}{a} \arctan \frac{x}{a} + C$$

2. 
$$\int \frac{\mathrm{d}x}{(x^2+a^2)^n} = \frac{x}{2(n-1)a^2(x^2+a^2)^{n-1}} + \frac{2n-3}{2(n-1)a^2} \int \frac{\mathrm{d}x}{(x^2+a^2)^{n-1}}$$

3. 
$$\int \frac{\mathrm{d}x}{x^2 - a^2} = \frac{1}{2a} \ln \left| \frac{x - a}{x + a} \right| + C$$

**9.1.4** 
$$ax^2 + b \ (a > 0)$$

1. 
$$\int \frac{\mathrm{d}x}{ax^2 + b} = \begin{cases} \frac{1}{\sqrt{ab}} \arctan \sqrt{\frac{a}{b}}x + C & (b > 0) \\ \frac{1}{2\sqrt{-ab}} \ln \left| \frac{\sqrt{ax} - \sqrt{-b}}{\sqrt{ax} + \sqrt{-b}} \right| + C & (b < 0) \end{cases}$$
2. 
$$\int \frac{x}{ax^2 + b} \, \mathrm{d}x = \frac{1}{2a} \ln \left| ax^2 + b \right| + C$$

2. 
$$\int \frac{x}{ax^2+b} dx = \frac{1}{2a} \ln |ax^2+b| + C$$

3. 
$$\int \frac{x^2}{ax^2 + b} dx = \frac{x}{a} - \frac{b}{a} \int \frac{dx}{ax^2 + b}$$

4. 
$$\int \frac{\mathrm{d}x}{x(ax^2+b)} = \frac{1}{2b} \ln \frac{x^2}{|ax^2+b|} + C$$

5. 
$$\int \frac{dx}{x^2(ax^2+b)} = -\frac{1}{bx} - \frac{a}{b} \int \frac{dx}{ax^2+b}$$

$$6. \int \frac{dx}{x^3(ax^2+b)} = \frac{a}{2b^2} \ln \frac{|ax^2+b|}{x^2} - \frac{1}{2bx^2} + C$$

$$7. \int \frac{dx}{(ax^2+b)^2} = \frac{x}{2b(ax^2+b)} + \frac{1}{2b} \int \frac{dx}{ax^2+b}$$

7. 
$$\int \frac{dx}{(ax^2+b)^2} = \frac{x}{2b(ax^2+b)} + \frac{1}{2b} \int \frac{dx}{ax^2+b}$$

9.1.5 
$$ax^{2} + bx + c \quad (a > 0)$$

$$1. \quad \frac{dx}{ax^{2} + bx + c} = \begin{cases} \frac{2}{\sqrt{4ac - b^{2}}} \arctan \frac{2ax + b}{\sqrt{4ac - b^{2}}} + C \quad (b^{2} < 4ac) \\ \frac{1}{\sqrt{b^{2} - 4ac}} \ln \left| \frac{2ax + b - \sqrt{b^{2} - 4ac}}{2ax + b + \sqrt{b^{2} - 4ac}} \right| + C \quad (b^{2} > 4ac) \end{cases}$$

$$2. \quad \int \frac{x}{ax^{2} + bx + c} dx = \frac{1}{2a} \ln |ax^{2} + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^{2} + bx + c}$$

2. 
$$\int \frac{x}{ax^2 + bx + c} dx = \frac{1}{2a} \ln|ax^2 + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^2 + bx + c}$$

# **9.1.6** $\sqrt{x^2 + a^2}$ (a > 0)

1. 
$$\int \frac{\mathrm{d}x}{\sqrt{x^2 + a^2}} = \operatorname{arsh} \frac{x}{a} + C_1 = \ln(x + \sqrt{x^2 + a^2}) + C$$

2. 
$$\int \frac{\mathrm{d}x}{\sqrt{(x^2+a^2)^3}} = \frac{x}{a^2\sqrt{x^2+a^2}} + \frac{1}{a^2\sqrt{x^2+a^2}}$$

9. 
$$\int \sqrt{x^2 + a^2} dx = \frac{x}{2} \sqrt{x^2 + a^2} + \frac{a^2}{2} \ln(x + \sqrt{x^2 + a^2}) + C$$

10. 
$$\int \sqrt{(x^2 + a^2)^3} dx = \frac{x}{8} (2x^2 + 5a^2) \sqrt{x^2 + a^2} + \frac{3}{8} a^4 \ln(x + \sqrt{x^2 + a^2}) + C$$

11. 
$$\int x\sqrt{x^2 + a^2} dx = \frac{1}{3}\sqrt{(x^2 + a^2)^3} + C$$

12. 
$$\int x^2 \sqrt{x^2 + a^2} \, dx = \frac{x}{8} (2x^2 + a^2) \sqrt{x^2 + a^2} - \frac{a^2}{8} \ln(x + \sqrt{x^2 + a^2}) + C$$

14. 
$$\int \frac{\sqrt{x^2 + a^2}}{x^2} dx = -\frac{\sqrt{x^2 + a^2}}{x} + \ln(x + \sqrt{x^2 + a^2}) + C$$

# **9.1.7** $\sqrt[a]{x^2 - a^2}$ (a > 0)

1. 
$$\int \frac{dx}{\sqrt{x^2 - a^2}} = \frac{x}{|x|} \operatorname{arch} \frac{|x|}{a} + C_1 = \ln \left| x + \sqrt{x^2 - a^2} \right| + C$$

3. 
$$\int \frac{x}{\sqrt{2-a^2}} dx = \sqrt{x^2 - a^2} + C$$

4. 
$$\int \frac{x}{\sqrt{1-x^2}} dx = -\frac{1}{\sqrt{1-x^2}} + C$$

$$6. \int \frac{x^2}{\sqrt{(x^2 - a^2)^3}} dx = -\frac{x}{\sqrt{x^2 - a^2}} + \ln|x + \sqrt{x^2 - a^2}| + C$$

$$7. \int \frac{dx}{x\sqrt{x^2 - a^2}} = \frac{1}{a} \arccos \frac{a}{|x|} + C$$

7. 
$$\int \frac{\mathrm{d}x}{x\sqrt{x^2 - a^2}} = \frac{1}{a} \arccos \frac{a}{|x|} + C$$

9. 
$$\int \sqrt{x^2 - a^2} dx = \frac{x}{2} \sqrt{x^2 - a^2} - \frac{a^2}{2} \ln|x + \sqrt{x^2 - a^2}| + C$$

10. 
$$\int \sqrt{(x^2 - a^2)^3} dx = \frac{x}{8} (2x^2 - 5a^2) \sqrt{x^2 - a^2} + \frac{3}{8} a^4 \ln|x + \sqrt{x^2 - a^2}| + C$$

11. 
$$\int x\sqrt{x^2 - a^2} dx = \frac{1}{3}\sqrt{(x^2 - a^2)^3} + C$$

12. 
$$\int x^2 \sqrt{x^2 - a^2} \, dx = \frac{x}{8} (2x^2 - a^2) \sqrt{x^2 - a^2} - \frac{a^4}{8} \ln|x + \sqrt{x^2 - a^2}| + C$$

13. 
$$\int \frac{\sqrt{x^2 - a^2}}{x} dx = \sqrt{x^2 - a^2} - a \arccos \frac{a}{|x|} + C$$
14. 
$$\int \frac{\sqrt{x^2 - a^2}}{x^2} dx = -\frac{\sqrt{x^2 - a^2}}{x} + \ln|x + \sqrt{x^2 - a^2}| + C$$

# 9.1.8 $\sqrt{a^2 - x^2}$ (a > 0)1. $\int \frac{dx}{\sqrt{a^2 - x^2}} = \arcsin \frac{x}{a} + C$

1. 
$$\int \frac{\mathrm{d}x}{\sqrt{a^2-x^2}} = \arcsin \frac{x}{a} +$$

2. 
$$\frac{dx}{\sqrt{(a^2 - x^2)^3}} = \frac{x}{a^2 \sqrt{a^2 - x^2}} + C$$
3. 
$$\int \frac{x}{\sqrt{a^2 - x^2}} dx = -\sqrt{a^2 - x^2} + C$$

3. 
$$\int \frac{x}{\sqrt{a^2-x^2}} dx = -\sqrt{a^2-x^2} + C$$

4. 
$$\int \frac{x}{\sqrt{x^2 - x^2}} dx = \frac{1}{\sqrt{x^2 - x^2}} + \frac{1}{\sqrt{x^2 - x^2}} +$$

6. 
$$\int \frac{x^2}{\sqrt{(a^2 - x^2)^3}} dx = \frac{x}{\sqrt{a^2 - x^2}} - \arcsin \frac{x}{a} + C$$

9. 
$$\int \sqrt{a^2 - x^2} dx = \frac{x}{2} \sqrt{a^2 - x^2} + \frac{a^2}{2} \arcsin \frac{x}{a} + C$$

9. 
$$\int \sqrt{a^2 - x^2} dx = \frac{x}{2} \sqrt{a^2 - x^2} + \frac{a^2}{2} \arcsin \frac{x}{a} + C$$
10. 
$$\int \sqrt{(a^2 - x^2)^3} dx = \frac{x}{8} (5a^2 - 2x^2) \sqrt{a^2 - x^2} + \frac{3}{8} a^4 \arcsin \frac{x}{a} + C$$

11. 
$$\int x\sqrt{a^2-x^2} dx = -\frac{1}{3}\sqrt{(a^2-x^2)^3} + C$$

12. 
$$\int x^2 \sqrt{a^2 - x^2} dx = \frac{x}{8} (2x^2 - a^2) \sqrt{a^2 - x^2} + \frac{a^4}{8} \arcsin \frac{x}{a} + C$$

11. 
$$\int \sqrt{a^2 - x^2} dx = -\frac{1}{3}\sqrt{(a^2 - x^2)^3 + C}$$
12. 
$$\int x^2 \sqrt{a^2 - x^2} dx = \frac{x}{8}(2x^2 - a^2)\sqrt{a^2 - x^2} + \frac{a^4}{8}\arcsin\frac{x}{a} + C$$
13. 
$$\int \frac{\sqrt{a^2 - x^2}}{x} dx = \sqrt{a^2 - x^2} + a\ln\frac{a - \sqrt{a^2 - x^2}}{|x|} + C$$

14. 
$$\int \frac{\sqrt{a^2 - x^2}}{x^2} dx = -\frac{\sqrt{a^2 - x^2}}{x} - \arcsin \frac{x}{a} + C$$

# **9.1.9** $\sqrt[x^2]{\pm ax^2 + bx + c}$ (a > 0)

$$\sqrt{ax^{2} + bx + c}$$
2. 
$$\int \sqrt{ax^{2} + bx + c} dx = \frac{2ax + b}{4a} \sqrt{ax^{2} + bx + c} + \frac{4ac - b^{2}}{8\sqrt{a^{3}}} \ln|2ax| + b + b$$

$$2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

2. 
$$\int \sqrt{ax^2 + bx + c} \, dx = \frac{2ax + b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac - b^2}{8\sqrt{a}^3} \ln|2ax + b| + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$
3. 
$$\int \frac{x}{\sqrt{ax^2 + bx + c}} \, dx = \frac{1}{a} \sqrt{ax^2 + bx + c} - \frac{b}{2\sqrt{a}^3} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

$$\begin{pmatrix}
\sqrt{ax^{2}+bx+c} & 2\sqrt{a} \\
C & \int \frac{dx}{\sqrt{c+bx-ax^{2}}} = -\frac{1}{\sqrt{a}} \arcsin \frac{2ax-b}{\sqrt{b^{2}+4ac}} + C \\
5. & \int \sqrt{c+bx-ax^{2}} dx = \frac{2ax-b}{4a} \sqrt{c+bx-ax^{2}} + \frac{b^{2}+4ac}{8\sqrt{a^{3}}} \arcsin \frac{2ax-b}{\sqrt{b^{2}+4ac}} + C \\
6. & \int \frac{x}{\sqrt{a^{3}+bx-ax^{2}}} dx = -\frac{1}{2} \sqrt{c+bx-ax^{2}} + \frac{b}{\sqrt{b^{2}+ac}}$$

5. 
$$\int \sqrt{c + bx - ax^2} dx = \frac{2ax - b}{4a} \sqrt{c + bx - ax^2} + \frac{2ax - b}{$$

$$\frac{b^2+4ac}{8\sqrt{a^3}}$$
 arcsin  $\frac{2ax-b}{\sqrt{b^2+4ac}}+C$ 

6. 
$$\int \frac{x}{\sqrt{c+bx-ax^2}} \, dx = -\frac{1}{a} \sqrt{c+bx-ax^2} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

# **9.1.10** $\sqrt{\pm \frac{x-a}{x-b}}$ & $\sqrt{(x-a)(x-b)}$

1. 
$$\int \sqrt{\frac{x-a}{x-b}} dx = (x-b)\sqrt{\frac{x-a}{x-b}} + (b-a)\ln(\sqrt{|x-a|} + \sqrt{|x-b|}) + C$$

2. 
$$\int \sqrt{\frac{x-a}{b-x}} \, \mathrm{d}x = (x-b) \sqrt{\frac{x-a}{b-x}} + (b-a) \arcsin \sqrt{\frac{x-a}{b-x}} + C$$

3. 
$$\int \frac{\mathrm{d}x}{\sqrt{(x-a)(b-x)}} = 2\arcsin\sqrt{\frac{x-a}{b-x}} + C \ (a < b)$$

$$\sqrt{(x-a)(0-x)} = \sqrt{\frac{x-a}{4}(x-a)(b-x)} + \frac{(b-a)^2}{4} \arcsin \sqrt{\frac{x-a}{b-x}} + C$$
4.  $\int \sqrt{(x-a)(b-x)} dx = \frac{2x-a-b}{4} \sqrt{(x-a)(b-x)} + \frac{(b-a)^2}{4} \arcsin \sqrt{\frac{x-a}{b-x}} + C$ 

# 9.1.11 Triangular function

1. 
$$\int \tan x \, dx = -\ln|\cos x| + C$$
2. 
$$\int \cot x \, dx = \ln|\sin x| + C$$

$$2. \int \cot x \, \mathrm{d}x = \ln|\sin x| + C$$

3. 
$$\int \sec x dx = \ln \left| \tan \left( \frac{\pi}{4} + \frac{x}{2} \right) \right| + C = \ln \left| \sec x + \tan x \right| + C$$

4. 
$$\int \csc x \, dx = \ln \left| \tan \frac{x}{2} \right| + C = \ln \left| \csc x - \cot x \right| + C$$

$$5. \int \sec^2 x \, \mathrm{d}x = \tan x + C$$

6. 
$$\int \csc^2 x \, \mathrm{d}x = -\cot x + C$$

7. 
$$\int \sec x \tan x dx = \sec x + C$$

5. 
$$\int \sec^2 x dx = \tan x + C$$
  
6.  $\int \csc^2 x dx = -\cot x + C$   
7.  $\int \sec x \tan x dx = \sec x + C$   
8.  $\int \csc x \cot x dx = -\csc x + C$   
9.  $\int \sin^2 x dx = \frac{x}{2} - \frac{1}{4} \sin 2x + C$ 

9. 
$$\int \sin^2 x dx = \frac{x}{2} - \frac{x}{4} \sin 2x + C$$
  
10.  $\int \cos^2 x dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$ 

10. 
$$\int \cos^2 x dx = \frac{1}{2} + \frac{1}{4} \sin^2 2x + C$$

10. 
$$\int \cos^{2} x dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$$
11. 
$$\int \sin^{n} x dx = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x dx$$
12. 
$$\int \cos^{n} x dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x dx$$
13. 
$$\frac{dx}{\sin^{n} x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\sin^{n-2} x}$$
14. 
$$\frac{dx}{\cos^{n} x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x}$$
15.

12. 
$$\int \cos^n x \, dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x \, dx$$

13. 
$$\frac{dx}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin n - 1} + \frac{n-2}{n-1} \int \frac{dx}{\sin n - 2} dx$$

14. 
$$\frac{dx}{\cos^n x} = \frac{1}{n-1} - \frac{\sin x}{n-1} + \frac{n-2}{n-1} \int \frac{dx}{n-2}$$

$$\int \cos^m x \sin^n x dx$$

$$= \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x dx$$

$$= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x dx$$

16. 
$$\int \sin ax \cos bx dx = -\frac{1}{2(a+b)} \cos(a+b)x - \frac{1}{2(a-b)} \cos(a-b)x + C$$
17. 
$$\int \sin ax \sin bx dx = -\frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$$

17. 
$$\int \sin ax \sin bx dx = -\frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a+b)} \sin(a-b)x + C$$

18. 
$$\int \cos ax \cos bx dx = \frac{2(a+b)}{2(a+b)} \sin(a+b)x + \frac{2(a+b)}{2(a+b)} \sin(a-b)x + C$$

17. 
$$\int \sin ax \sin bx dx = -\frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$$
18. 
$$\int \cos ax \cos bx dx = \frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$$
19. 
$$\int \frac{dx}{a+b \sin x} = \begin{cases} \frac{2}{\sqrt{a^2 - b^2}} \arctan \frac{a \tan \frac{x}{2} + b}{\sqrt{a^2 - b^2}} + C & (a^2 > b^2) \\ \frac{1}{\sqrt{b^2 - a^2}} \ln \left| \frac{a \tan \frac{x}{2} + b - \sqrt{b^2 - a^2}}{a \tan \frac{x}{2} + b + \sqrt{b^2 - a^2}} \right| + C & (a^2 < b^2) \end{cases}$$
20. 
$$\int \frac{dx}{a+b \cos x} = \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan \left(\sqrt{\frac{a-b}{a+b}} \tan \frac{x}{2}\right) + C & (a^2 < b^2) \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln \left| \frac{\tan \frac{x}{2} + \sqrt{\frac{a+b}{b-a}}}{\tan \frac{x}{2} - \sqrt{\frac{a+b}{b-a}}} \right| + C & (a^2 < b^2) \end{cases}$$

$$\left( \begin{array}{c} \left| \tan \frac{\pi}{2} - \sqrt{\frac{b-a}{b-a}} \right| \end{array} \right)$$

21. 
$$\int \frac{dx}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{ab} \arctan\left(\frac{b}{a} \tan x\right) + C$$
22. 
$$\int \frac{dx}{a^2 \cos^2 x - b^2 \sin^2 x} = \frac{1}{2ab} \ln \left| \frac{b \tan x + a}{b \tan x - a} \right| + C$$

23. 
$$\int x \sin ax dx = \frac{1}{a^2} \sin ax - \frac{1}{a} x \cos ax + C$$

- 24.  $\int x^2 \sin ax dx = -\frac{1}{a}x^2 \cos ax + \frac{2}{a^2}x \sin ax + \frac{2}{a^3}\cos ax + C$
- 25.  $\int x \cos ax dx = \frac{1}{a^2} \cos ax + \frac{1}{a} x \sin ax + C$ 26.  $\int x^2 \cos ax dx = \frac{1}{a} x^2 \sin ax + \frac{2}{a^2} x \cos ax \frac{2}{a^3} \sin ax + C$

## 9.1.12 Inverse triangular function (a > 0)

- 1.12 Inverse triangular function (a > 0)1.  $\int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2 x^2} + C$ 2.  $\int x \arcsin \frac{x}{a} dx = (\frac{x^2}{2} \frac{a^2}{4}) \arcsin \frac{x}{a} + \frac{x}{4} \sqrt{x^2 x^2} + C$ 3.  $\int x^2 \arcsin \frac{x}{a} dx = \frac{x^3}{3} \arcsin \frac{x}{a} + \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 x^2} + C$ 4.  $\int \arccos \frac{x}{a} dx = x \arccos \frac{x}{a} \sqrt{a^2 x^2} + C$ 5.  $\int x \arccos \frac{x}{a} dx = (\frac{x^2}{2} \frac{a^2}{4}) \arccos \frac{x}{a} \frac{x}{4} \sqrt{a^2 x^2} + C$ 6.  $\int x^2 \arccos \frac{x}{a} dx = \frac{x^3}{3} \arccos \frac{x}{a} \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 x^2} + C$ 7.  $\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} \frac{a}{2} \ln(a^2 + x^2) + C$ 8.  $\int x \arctan \frac{x}{a} dx = \frac{1}{2} (a^2 + x^2) \arctan \frac{x}{a} \frac{a}{2} x + C$

- 9.  $\int x^2 \arctan \frac{x}{a} dx = \frac{x^3}{3} \arctan \frac{x}{a} \frac{a}{6}x^2 + \frac{a^3}{6} \ln(a^2 + x^2) + C$ 9.1.13 Exponential function

- 1.  $\int a^{x} dx = \frac{1}{\ln a} a^{x} + C$ 2.  $\int e^{ax} dx = \frac{1}{a} a^{ax} + C$ 3.  $\int x e^{ax} dx = \frac{1}{a^{2}} (ax 1) a^{ax} + C$ 4.  $\int x^{n} e^{ax} dx = \frac{1}{a} x^{n} e^{ax} \frac{n}{a} \int x^{n-1} e^{ax} dx$

- 4.  $\int x e^{-\frac{1}{4}} dx = \frac{1}{\ln a} a^x \frac{1}{(\ln a)^2} a^x + C$ 5.  $\int x a^x dx = \frac{x}{\ln a} a^x \frac{1}{(\ln a)^2} a^x + C$ 6.  $\int x^n a^x dx = \frac{1}{\ln a} x^n a^x \frac{n}{\ln a} \int x^{n-1} a^x dx$ 7.  $\int e^{ax} \sin bx dx = \frac{1}{a^2 + b^2} e^{ax} (a \sin bx b \cos bx) + C$
- 8.  $\int e^{ax} \cos bx dx = \frac{a^2 + b^2}{a^2 + b^2} e^{ax} (b \sin bx + a \cos bx) + C$
- 9.  $\int e^{ax} \sin^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \sin^{n-1} bx (a \sin bx nb \cos bx) +$  $\frac{n(n-1)b^2}{a^2+b^2n^2} \int e^{ax} \sin^{n-2} bx dx$
- 10.  $\int e^{ax} \cos^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) +$  $\frac{-a^{(n-1)b^2}}{a^2+b^2n^2} \int e^{ax} \cos^{n-2} bx dx$

## 9.1.14 Logarithmic function

- 1.  $\int \ln x dx = x \ln x x + C$ 2.  $\int \frac{dx}{x \ln x} = \ln |\ln x| + C$ 3.  $\int x^n \ln x dx = \frac{1}{n+1} x^{n+1} (\ln x \frac{1}{n+1}) + C$ 4.  $\int (\ln x)^n dx = x (\ln x)^n n \int (\ln x)^{n-1} dx$ 5.  $\int x^m (\ln x)^n dx = \frac{1}{m+1} x^{m+1} (\ln x)^n \frac{n}{m+1} \int x^m (\ln x)^{n-1} dx$

### 9.2Regular expression

## Special pattern characters

Characters	Description
•	Not newline
\t	Tab (HT)
\n	Newline (LF)
\v	Vertical tab (VT)
\f	Form feed (FF)
\r	Carriage return (CR)
\cletter	Control code
\xhh	ASCII character
\uhhhh	Unicode character
\0	Null
\int	Backreference
\d	Digit
\D	Not digit
\s	Whitespace
\S	Not whitespace
\W	Word (letters, numbers and the underscore)
\W	Not word
\character	Character
[class]	Character class
[^class]	Negated character class

### 9.2.2Quantifiers

Characters	Times
*	0 or more
+	1 or more
?	0 or 1
{int}	int
{int,}	int or more
{min,max}	Between min and max

By default, all these quantifiers are greedy (i.e., they take as many characters that meet the condition as possible). This behavior can be overridden to ungreedy (i.e., take as few characters that meet the condition as possible) by adding a question mark (?) after the quantifier.

## Groups

	Characters	Description
(	subpattern)	Group with backreference
(?	:subpattern)	Group without backreference

### 9.2.4Assertions

Characters	Description
^	Beginning of line
\$	End of line
\b	Word boundary
\B	Not a word boundary
(?=subpattern)	Positive lookahead
(?!subpattern)	Negative lookahead

### 9.2.5Alternative

A regular expression can contain multiple alternative patterns simply by separating them with the separator operator ( $\mid$ ): The regular expression will match if any of the alternatives match, and as soon as one does

## Character classes

Class	Description
[:alnum:]	Alpha-numerical character
[:alpha:]	Alphabetic character
[:blank:]	Blank character
[:cntrl:]	Control character
[:digit:]	Decimal digit character
[:graph:]	Character with graphical representation
[:lower:]	Lowercase letter
[:print:]	Printable character
[:punct:]	Punctuation mark character
[:space:]	Whitespace character
[:upper:]	Uppercase letter
[:xdigit:]	Hexadecimal digit character
[:d:]	Decimal digit character
[:w:]	Word character
[:s:]	Whitespace character

Please note that the brackets in the class names are additional to those opening and closing the class definition. For example:

 $\[\ [\ [: \ \ ]\ ]\ ]\ is\ a\ character\ class\ that\ matches\ any\ alphabetic\ character\ class\ cla$ 

[abc[:digit:]] is a character class that matches a, b, c, or a digit.

[^[:space:]] is a character class that matches any character except a whitespace.