

Luna's Magic Reference

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1 Environment

1.1 Vimrc

```

1 set ru nu ts=4 sts=4 sw=4 si sm hls is ar bs=2 mouse=a
2 syntax on
3 nm <F3> :vsplit %<.in <CR>
4 nm <F4> :!gedit % <CR>
5 au BufEnter *.cpp set cin
6 au BufEnter *.cpp nm <F5> :!time ./%< <CR>|nm <F7> :!
   gdb ./%< <CR>|nm <F8> :!time ./%< <CR>|nm <F9> :!g++ % -o % -g -std=gnu++14 -O2 -DLOCAL &&
   size %< <CR>
7 au BufEnter *.java nm <F5> :!time java %< <CR>|nm <F8>
   :!time java %< <CR>|nm <F9> :!javac % <CR>

```

2 Data Structure

2.1 KD tree

```

1 /* kd_tree : finds the k-th closest point in  $O(kn^{1-\frac{1}{k}})$ .
2 Usage : Stores the data in p[]. Call function init (n,
   k). Call min_kth (d, k). (or max_kth) (k is 1-
   based)
3 Note : Switch to the commented code for Manhattan
   distance.
4 Status : SPOJ-FAILURE Accepted.*/
5 template <int MAXN = 200000, int MAXK = 2>
6 struct kd_tree {
7     int k, size;
8     struct point { int data[MAXN], id; } p[MAXN];
9     struct kd_node {
10         int l, r; point p, dmin, dmax;
11         kd_node() {}
12         kd_node (const point &rhs) : l (-1), r (-1), p (rhs)
13             , dmin (rhs), dmax (rhs) {}
14     void merge (const kd_node &rhs, int k) {
15         for (register int i = 0; i < k; ++i) {
16             dmin.data[i] = std::min (dmin.data[i], rhs.dmin.
17                 data[i]);
18             dmax.data[i] = std::max (dmax.data[i], rhs.dmax.
19                 data[i]); } }
20     long long min_dist (const point &rhs, int k) const {
21         register long long ret = 0;
22         for (register int i = 0; i < k; ++i) {
23             if (dmin.data[i] <= rhs.data[i] && rhs.data[i] <=
24                 dmax.data[i]) continue;
25             ret += std::min (1ll * (dmin.data[i] - rhs.data[i]
26                 ) * (dmin.data[i] - rhs.data[i]),
27                 1ll * (dmax.data[i] - rhs.data[i]) * (dmax.
28                     data[i] - rhs.data[i]));
29             // ret += std::max (0, rhs.data[i] - dmax.data[i])
30             // + std::max (0, dmin.data[i] - rhs.data[i]);
31         } return ret; }
32     long long max_dist (const point &rhs, int k) {
33         long long ret = 0;
34         for (int i = 0; i < k; ++i) {
35             int tmp = std::max (std::abs (dmin.data[i] - rhs.
36                 data[i]), std::abs (dmax.data[i] - rhs.data[i]
37                 ));
38             ret += 1ll * tmp * tmp; }
39             // ret += std::max (std::abs (rhs.data[i] - dmax.
40             // data[i]) + std::abs (rhs.data[i] - dmin.data[i]));
41         } return ret; } } tree[MAXN * 4];
42 struct result {
43     long long dist; point d; result() {}
44     result (const long long &dist, const point &d) :
45         dist (dist), d (d) {}
46     bool operator > (const result &rhs) const { return
47         dist > rhs.dist || (dist == rhs.dist && d.id >
48             rhs.d.id); }
49     bool operator < (const result &rhs) const { return
50         dist < rhs.dist || (dist == rhs.dist && d.id <
51             rhs.d.id); } };
52 long long sqrdist (const point &a, const point &b) {
53     long long ret = 0;
54     for (int i = 0; i < k; ++i) ret += 1ll * (a.data[i]
55         - b.data[i]) * (a.data[i] - b.data[i]);
56     // for (int i = 0; i < k; ++i) ret += std::abs (a.
57     // data[i] - b.data[i]);
58     return ret; }
59 int alloc() { tree[size].l = tree[size].r = -1;
60     return size++; }
61 void build (const int &depth, int &rt, const int &l,
62     const int &r) {
63     if (l > r) return;
64     register int middle = (l + r) >> 1;
65     std::nth_element (p + l, p + middle, p + r + 1, [=]
66         (const point &a, const point &b) { return a.
67             data[depth] < b.data[depth]; });
68     tree[rt] = alloc(); = kd_node (p[middle]);
69     if (l == r) return;
70     build ((depth + 1) % k, tree[rt].l, l, middle - 1);
71     build ((depth + 1) % k, tree[rt].r, middle + 1, r);
72     if (!tree[rt].l) tree[rt].merge (tree[tree[rt].l], k
73         );
74     if (!tree[rt].r) tree[rt].merge (tree[tree[rt].r], k
75         );
76     std::priority_queue<result>, std::vector<result>, std
77     ::less<result>> heap_l;
78     std::priority_queue<result>, std::vector<result>, std
79     ::greater<result>> heap_r;
80 void min_kth (const int &depth, const int &rt, const
81     int &m, const point &d) {
82     result tmp = result (sqrdist (tree[rt].p, d), tree[
83         rt].p);
84     if ((int)heap_l.size() < m) heap_l.push (tmp);
85     else if (tmp < heap_l.top()) {
86         heap_l.pop();
87         heap_l.push (tmp); }

```

```

61 int x = tree[rt].l, y = tree[rt].r;
62 if (~x && ~y && sqrdist (d, tree[x].p) > sqrdist (d,
63     tree[y].p)) std::swap (x, y);
64 if (~x && ((int)heap_l.size() < m || tree[x].
65     min_dist (d, k) < heap_l.top().dist))
66     min_kth ((depth + 1) % k, x, m, d);
67 if (~y && ((int)heap_l.size() < m || tree[y].
68     min_dist (d, k) < heap_l.top().dist))
69     min_kth ((depth + 1) % k, y, m, d);
70 void max_kth (const int &depth, const int &rt, const
71     int &m, const point &d) {
72     result tmp = result (sqrdist (tree[rt].p, d), tree[
73         rt].p);
74     if ((int)heap_r.size() < m) heap_r.push (tmp);
75     else if (tmp > heap_r.top()) {
76         heap_r.pop();
77         heap_r.push (tmp); }
78 int x = tree[rt].l, y = tree[rt].r;
79 if (~x && ~y && sqrdist (d, tree[x].p) < sqrdist (d,
80     tree[y].p)) std::swap (x, y);
81 if (~x && ((int)heap_r.size() < m || tree[x].
82     max_dist (d, k) > heap_r.top().dist))
83     max_kth ((depth + 1) % k, x, m, d);
84 if (~y && ((int)heap_r.size() < m || tree[y].
85     max_dist (d, k) > heap_r.top().dist))
86     max_kth ((depth + 1) % k, y, m, d); }
87 void init (int n, int k) { this->k = k; size = 0;
88     int rt = 0; build (0, rt, 0, n - 1); }
89 result min_kth (const point &d, const int &m) {
90     heap_l = decltype (heap_l) (); min_kth (0, 0, m,
91         d); return heap_l.top (); }
92 result max_kth (const point &d, const int &m) {
93     heap_r = decltype (heap_r) (); max_kth (0, 0, m,
94         d); return heap_r.top (); } }

```

2.2 Splay

```

1 void push_down (int x) {
2     if (~n[x].c[0]) push (n[x].c[0], n[x].t);
3     if (~n[x].c[1]) push (n[x].c[1], n[x].t);
4     n[x].t = tag (); }
5 void update (int x) {
6     n[x].m = gen (x);
7     if (~n[x].c[0]) n[x].m = merge (n[n[x].c[0]].m, n[x].
8         m);
9     if (~n[x].c[1]) n[x].m = merge (n[x].m, n[n[x].c[1]].
10         m); }
11 void rotate (int x, int k) {
12     push_down (x); push_down (n[x].c[k]);
13     int y = n[x].c[k]; n[x].c[k] = n[y].c[k ^ 1]; n[y].c[
14         k ^ 1] = x;
15     if (n[x].f != -1) n[n[x].f].c[n[n[x].f].c[1] == x] =
16         y;
17     n[y].f = n[x].f; n[x].f = y; if (~n[x].c[k]) n[n[x].c
18         [k]].f = x;
19     update (x); update (y); }
20 void splay (int x, int s = -1) {
21     push_down (x);
22     while (n[x].f != s) {
23         if (n[n[x].f].f != s) rotate (n[n[x].f].f, n[n[x].
24             f].f.c[1] == n[x].f);
25         rotate (n[x].f, n[n[x].f].c[1] == x); }
26     update (x);
27     if (s == -1) root = x; }

```

2.3 Link-cut tree

```

1 void access (int x) {
2     int u = x, v = -1;
3     while (u != -1) {
4         splay (u); push_down (u);
5         if (~n[u].c[1]) n[n[u].c[1]].f = -1, n[n[u].c[1]].p
6             = u;
7         n[u].c[1] = v;
8         if (~v) n[v].f = u, n[v].p = -1;
9         update (u); u = n[v = u].p; }
10 splay (x); }

```

3 Formula

3.1 Zellers congruence

```

1 /* Zeller's congruence : converts between a calendar
2 date and its Gregorian calendar day. (y >= 1) (0 =
3 Monday, 1 = Tuesday, ..., 6 = Sunday) */
4 int get_id (int y, int m, int d) {
5     if (m < 3) { --y; m += 12; }
6     return 365 * y + y / 4 - y / 100 + y / 400 + (153 * (
7         m - 3) + 2) / 5 + d - 307; }
8 std::tuple <int, int, int> date (int id) {
9     int x = id + 1789995, n, i, j, y, m, d;
10    n = 4 * x / 146097; x -= (146097 * n + 3) / 4;
11    i = (4000 * (x + 1)) / 1461001; x -= 1461 * i / 4 -
12        31;
13    j = 80 * x / 2447; d = x - 2447 * j / 80;
14    x = j / 11;
15    m = j + 2 - 12 * x; y = 100 * (n - 49) + i + x;
16    return std::make_tuple (y, m, d); }

```

3.2 Lattice points below segment

```

1 /* Euclidean-like algorithm : computes the sum of
2  $\sum_{i=0}^{n-1} \lfloor \frac{a+bi}{m} \rfloor$ . */
3 long long solve (long long n, long long a, long long b,
4     long long m) {
5     if (b == 0) return n * (a / m);
6     if (a >= m) return n * (a / m) + solve (n, a % m, b,
7         m);

```

```

3  int ans = 1, mul = x; while (n) {
4      if (n & 1) ans = int (1ll * ans * mul % mod);
5      mul = int (1ll * mul * mul % mod); n >= 1; }
6  return ans; }
7  long long mul_mod (long long x, long long y, long long
   mod) {
8      long long t = (x * y - (long long) ((long double) x /
        mod * y + 1E-3) * mod) % mod;
9      return t < 0 ? t + mod : t; }
10 long long lfpm (long long x, long long n, long long
    mod) {
11     long long ans = 1, mul = x; while (n) {
12         if (n & 1) ans = mul_mod (ans, mul, mod);
13         mul = mul_mod (mul, mul, mod); n >= 1; }
14     return ans; }

```

4.2 Euclidean algorithm

```

1 /* Euclidean algorithm : solves for ax + by = gcd (a,
2    b). */
3 void euclid (const long long &a, const long long &b,
4             long long &x, long long &y) {
5     if (b == 0) x = 1, y = 0;
6     else euclid (b, a % b, y, x), y -= a / b * x; }
7
8 long long inverse (long long x, long long m) {
9     long long a, b; euclid (x, m, a, b); return (a % m +
10    m) % m; }

```

4.3 Discrete Fourier transform

```

1  /* Discrete Fourier transform : the nafarious you-know
   -what thing.
2  Usage : call init for the suggested array size, and
   solve for the transform. (use f!=0 for the inverse
   ) */
3  template <int MAXN = 1000000>
4  struct dft {
5      typedef std::complex <double> complex;
6      complex e[2][MAXN];
7      int init (int n) {
8          int len = 1;
9          for (; len <= 2 * n; len <= 1);
10         for (int i = 0; i < len; ++i) {
11             e[0][i] = complex (cos (2 * PI * i / len), sin (2
12                 * PI * i / len));
13             e[1][i] = complex (cos (2 * PI * i / len), -sin (2
14                 * PI * i / len)); }
15         return len; }
16     void solve (complex *a, int n, int f) {
17         for (int i = 0; i < n; ++i) {
18             if (i > j) std::swap (a[i], a[j]);
19             for (int t = n >> 1; (j ^= t) < t; t >>= 1); }
20         for (int i = 2; i <= n; i <= 1)
21             for (int j = 0; j <= n; j += i)
22                 for (int k = 0; k < (i >> 1); ++k) {
23                     complex A = a[j + k];
24                     complex B = e[f][n / i * k] * a[j + k + (i >> 1)
25                         ];
26                     a[j + k] = A + B;
27                     a[j + k + (i >> 1)] = A - B; }
28         if (f == 1) {
29             for (int i = 0; i < n; ++i) a[i] = complex (a[i].
30                 real () / n, a[i].imag ()); } }

```

```

1  /* Fast Walsh-Hadamard transform : binary operation
   transform. */
2  void fwt (int *a, int n, int w) {
3      for (int i = 1; i < n; i <= 1)
4          for (int j = 0; j < n; j += i < 1) {
5              for (int k = 0; k < i; ++k) {
6                  int x = a[j + k], y = a[i + j + k];
7                  if (w) {
8                      /* xor : a[j + k] = (x + y) / 2, a[i + j + k] = (x
                       - y) / 2, and : a[j + k] = x - y, or : a[i +
                       j + k] = y - x; */
9                  } else {
10                     /* xor : a[j + k] = x + y, a[i + j + k] = x - y,
                        and : a[j + k] = x + y, or : a[i + j + k] = x
                        + y; */
11                 }
12             }
13         }
14     }
15 }

```

```

1 /* Number theoretic transform : NTT for any module.
2 Usage : Perform NTT on 3 modules and call crt () to
3 merge the result. */
4 template <int MAXN = 1000000>
5 struct ntt {
6     int MOD[3] = {1045430273, 1051721729, 1053818881},
7     PRT[3] = {3, 6, 7};
8     void solve (int *a, int n, int f = 0, int mod =
9         998244353, int prt = 3) {
10         for (int i = 0, j = 0; i < n; ++i) {
11             if (i > j) std::swap (a[i], a[j]);
12             for (int t = n >> 1; (j ^= t) < t; t >>= 1);
13             for (int i = 2; i <= n; i <= 1) {
14                 static int exp[MAXN]; exp[0] = 1;
15                 exp[1] = fpm (prt, (mod - 1) / i, mod);
16                 if (f == 1) exp[1] = fpm (exp[1], mod - 2, mod);
17                 for (int k = 2; k < (i >> 1); ++k) {
18                     exp[k] = int (1ll * exp[k - 1] * exp[1] % mod);
19                 }
20                 for (int j = 0; j < n; j += i) {
21                     for (int k = 0; k < (i >> 1); ++k) {
22                         int &pa = a[j + k], &pb = a[j + k + (i >> 1)];
23                         int A = pa, B = int (1ll * pb * exp[k] % mod);
24                         pa = (A + B) % mod;
25                         pb = (A - B + mod) % mod; } } } }

```

4.1 Fast power module

```
1 /* Fast power module :  $x^n$  */
2 int fpm (int x, int n, int mod) {
```



```

22 if (f == 1) {
23     int rev = fpm(n, mod - 2, mod);
24     for (int i = 0; i < n; ++i) a[i] = int (1LL * a[i]
25         * rev % mod); } }
26 int crt (int *a, int mod) {
27     static int inv[3][3];
28     for (int i = 0; i < 3; ++i) for (int j = 0; j < 3;
29         ++j)
30         inv[i][j] = (int) inverse (MOD[i], MOD[j]);
31     static int x[3];
32     for (int i = 0; i < 3; ++i) { x[i] = a[i];
33         for (int j = 0; j < i; ++j) {
34             int t = (x[i] - x[j] + MOD[i]) % MOD[i];
35             if (t < 0) t += MOD[i];
36             x[i] = int (1LL * t * inv[j][i] % MOD[i]); } }
37     int sum = 1, ret = x[0] % mod;
38     for (int i = 1; i < 3; ++i) {
39         sum = int (1LL * sum * MOD[i - 1] % mod);
40         ret += int (1LL * x[i] * sum % mod);
41         if (ret >= mod) ret -= mod; }
42     return ret; } }

```

4.6 Polynomial operation

```

1 template <int MAXN = 1000000>
2 struct polynomial {
3     ntt <MAXN> tr;
4     /* inverse : finds a polynomial b so that
5         a(x)b(x) ≡ 1 mod x^n mod mod.
6     Note : n must be a power of 2. 2x max length. */
7     void inverse (int *a, int *b, int n, int mod, int prt) {
8         static int c[MAXN]; b[0] = ::inverse (a[0], mod); b
9         [1] = 0;
10        for (int m = 2, i, m <= n; m <= 1) {
11            std::copy (a, a + m, c);
12            std::fill (b + m, b + m + m, 0); std::fill (c + m,
13                c + m + m, 0);
14            tr.solve (c, m + m, 0, mod, prt); tr.solve (b, m +
15                m, 0, mod, prt);
16            for (int i = 0; i < m + m; ++i) b[i] = 1LL * b[i] *
17                (2 - 1LL * b[i] * c[i] % mod + mod) % mod;
18            tr.solve (b, m + m, 1, mod, prt); std::fill (b + m,
19                b + m + m, 0); } }
20    /* sqrt : finds a polynomial b so that
21        b^2(x) ≡ a(x) mod x^n mod mod.
22    Note : n ≥ 2 must be a power of 2. 2x max length. */
23    void sqrt (int *a, int *b, int n, int mod, int prt) {
24        static int d[MAXN], ib[MAXN]; b[0] = 1; b[1] = 0;
25        int i2 = ::inverse (2, mod), m, i;
26        for (int m = 2; m <= n; m <= 1) {
27            std::copy (a, a + m, d);
28            std::fill (d + m, d + m + m, 0); std::fill (b + m,
29                b + m + m, 0);
30            tr.solve (d, m + m, 0, mod, prt); inverse (b, ib, m
31                , mod, prt);
32            tr.solve (ib, m + m, 0, mod, prt); tr.solve (b, m +
33                m, 0, mod, prt);
34            for (int i = 0; i < m + m; ++i) b[i] = (1LL * b[i]
35                * i2 + 1LL * i2 * d[i] % mod * ib[i]) % mod;
36            tr.solve (b, m + m, 1, mod, prt); std::fill (b + m,
37                b + m + m, 0); } }
38    /* divide : given polynomial a(x) and b(x) with degree
39        n and m respectively, finds a(x) = d(x)b(x) + r(x)
40        with deg(d) ≤ n - m and deg(r) < m. 4x max length
41        required. */
42    void divide (int *a, int n, int *b, int m, int *d,
43        int *r, int mod, int prt) {
44        static int u[MAXN], v[MAXN]; while (!b[m - 1]) --m;
45        int p = 1, t = n - m + 1; while (p < t <= 1) p <=
46        1;
47        std::fill (u, u + p, 0); std::reverse_copy (b, b + m
48            , u); inverse (u, v, p, mod, prt);
49        std::fill (v + t, v + p, 0); tr.solve (v, p, 0, mod,
50            prt); std::reverse_copy (a, a + n, u);
51        std::fill (u + t, u + p, 0); tr.solve (u, p, 0, mod,
52            prt);
53        for (int i = 0; i < p; ++i) u[i] = 1LL * u[i] * v[i]
54            % mod;
55        tr.solve (u, p, 1, mod, prt); std::reverse (u, u + t
56            );
57        for (p = 1; p < n; p <= 1); std::fill (u + t, u + p
58            , 0);
59        tr.solve (u, p, 0, mod, prt); std::copy (b, b + m, v
60            );
61        std::fill (v + m, v + p, 0); tr.solve (v, p, 0, mod,
62            prt);
63        for (int i = 0; i < p; ++i) u[i] = 1LL * u[i] * v[i]
64            % mod;
65        tr.solve (u, p, 1, mod, prt);
66        for (int i = 0; i < m; ++i) r[i] = (a[i] - u[i] +
67            mod) % mod;
68        std::fill (r + m, r + p, 0); } }

```

4.7 Chinese remainder theorem

```

1 /* Chinese remainder theorem : finds positive integers
2    x = out.first + k * out.second that satisfies x %
3    in[i].second = in[i].first. */
4 struct crt {
5     long long fix (const long long &a, const long long &b
6         ) { return (a % b + b) % b; }
7     bool solve (const std::vector <std::pair <long long,
8         long long>> &in, std::pair <long long, long long>
9         &out) {
10         out = std::make_pair (1LL, 1LL);
11         for (int i = 0; i < (int) in.size (); ++i) {
12             long long n, u;
13             euclid (out.second, in[i].second, n, u);
14             long long divisor = std::__gcd (out.second, in[i].
15                 second);

```

```

10 if ((in[i].first - out.first) % divisor) return
11     false;
12 n *= (in[i].first - out.first) / divisor;
13 n = fix (n, in[i].second);
14 out.first += out.second * n;
15 out.second *= in[i].second / divisor;
16 out.first = fix (out.first, out.second); }
17 return true; } }

```

4.8 Linear Recurrence

```

1 /* Linear recurrence : finds the n-th element of a
2    linear recurrence.
3 Usage : vector <int> - first n terms, vector <int> -
4    transition function, calc (k) : the kth term mod
5    MOD.
6 Example : In : {2, 1}, {2, 1} :
7           a1 = 2, a2 = 1, an = 2a_{n-1} + a_{n-2}, Out : calc (3) = 5,
8           calc (10007) = 959155122 (MOD 1E9+7) */
9 struct linear_rec {
10     const int LOG = 30, MOD = 1E9 + 7; int n;
11     std::vector <int> first, trans;
12     std::vector <std::vector <int>> bin;
13     std::vector <int> add (std::vector <int> &a, std::
14         vector <int> &b) {
15         std::vector <int> result(n * 2 + 1, 0);
16         for (int i = 0; i <= n; ++i) for (int j = 0; j <= n;
17             ++j)
18             if ((result[i + j] += 1LL * a[i] * b[j] % MOD) >=
19                 MOD) result[i + j] -= MOD;
20         for (int i = 2 * n; i > n; --i) {
21             for (int j = 0; j < n; ++j)
22                 if ((result[i - 1 - j] += 1LL * result[i] * trans[
23                     j] % MOD) >= MOD) result[i - 1 - j] -= MOD;
24             result[i] = 0; }
25         result.erase(result.begin() + n + 1, result.end());
26         return result; }
27     linear_rec (const std::vector <int> &first, const std
28         ::vector <int> &trans) : first(first), trans(
29             trans) {
30         n = first.size(); std::vector <int> a(n + 1, 0); a
31         [1] = 1; bin.push_back(a);
32         for (int i = 1; i < LOG; ++i) bin.push_back(add(bin
33             [i - 1], bin[i - 1])); }
34     int solve (int k) {
35         std::vector <int> a(n + 1, 0); a[0] = 1;
36         for (int i = 0; i < LOG; ++i) if (k >> i & 1) a =
37             add(a, bin[i]);
38         int ret = 0;
39         for (int i = 0; i < n; ++i) if ((ret += (long long)
40             a[i + 1] * first[i] % MOD) >= MOD) ret -= MOD;
41         return ret; } }

```

4.9 Berlekamp Massey algorithm

```

1 /* Berlekamp Massey algorithm : Complexity: O(n^2)
2    Requirement: const MOD, inverse(int)
3 Input: the first elements of the sequence
4 Output: the recursive equation of the given sequence
5 Example In: {1, 1, 2, 3}
6 Example Out: {1, 1000000006, 1000000006} (MOD = 1e9+7)
7 */
8 struct berlekamp-massey {
9     struct Poly { std::vector <int> a; Poly() { a.clear()
10         ; }
11         Poly (std::vector <int> &a) : a (a) {}
12         int length () const { return a.size(); }
13         Poly move (int d) { std::vector <int> na (d, 0);
14             na.insert (na.end (), a.begin (), a.end ());
15             return Poly (na); }
16         int calc(std::vector <int> &d, int pos) { int ret =
17             0;
18             for (int i = 0; i < (int) a.size (); ++i) {
19                 if ((ret += 1LL * d[pos - i] * a[i] % MOD) >= MOD)
20                     ret -= MOD; } }
21         return ret; }
22     Poly operator - (const Poly &b) {
23         std::vector <int> na (std::max (this -> length (),
24             b.length ());
25         for (int i = 0; i < (int) na.size (); ++i) {
26             int aa = i < this -> length () ? this -> a[i] : 0;
27             int bb = i < b.length () ? b.a[i] : 0;
28             na[i] = (aa + MOD - bb) % MOD; }
29         return Poly (na); }
30     Poly operator * (const int &c, const Poly &p) {
31         std::vector <int> na (p.length ());
32         for (int i = 0; i < (int) na.size (); ++i) {
33             na[i] = 1LL * c * p.a[i] % MOD; }
34         return na; }
35     std::vector <int> solve(vector<int> a) {
36         int n = a.size (); Poly s, b;
37         s.a.push_back (1); b.a.push_back (1);
38         for (int i = 0, j = -1, ld = 1; i < n; ++i) {
39             int d = s.calc(a, i); if (d) {
40                 if ((s.length () - 1) * 2 <= i) {
41                     Poly ob = b; b = s;
42                     s = s - 1LL * d * inverse (ld) % MOD * ob.move (i
43                         - j);
44                     j = i; ld = d;
45                 } else {
46                     s = s - 1LL * d * inverse (ld) % MOD * b.move (i
47                         - j); } } }
48         return s.a; } }

```

4.10 Baby step giant step algorithm

```

1 /* Baby step giant step algorithm : Solves a^x = b mod c
2    in O(√c). */
3 struct bsgs {
4     int solve (int a, int b, int c) {

```

```

4| std::unordered_map <int, int> bs;
5| int m = (int) sqrt ((double) c) + 1, res = 1;
6| for (int i = 0; i < m; ++i) {
7|     if (bs.find (res) == bs.end ()) bs[res] = i;
8|     res = int (1LL * res * a % c); }
9| int mul = 1, inv = (int) inverse (a, c);
10| for (int i = 0; i < m; ++i) mul = int (1LL * mul *
11|     inv % c);
12| res = b % c;
13| for (int i = 0; i < m; ++i) {
14|     if (bs.find (res) != bs.end ()) return i * m + bs[
15|         res];
16|     res = int (1LL * res * mul % c); }
17| return -1; } };
```

4.11 Miller Rabin primality test

```

1| /* Miller Rabin : tests whether a certain integer is
2|    prime. */
3| struct miller_rabin {
4|     int BASE[12] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29,
5|         31, 37};
6|     bool check (const long long &p, const long long &b) {
7|         long long n = p - 1;
8|         for (; ~n & 1; n >= 1);
9|         long long res = llfpm (b, n, p);
10|         for (; n != p - 1 && res != 1 && res != p - 1; n <=
11|             1)
12|             res = mul_mod (res, res, p);
13|         return res == p - 1 || (n & 1) == 1; }
14|     bool solve (const long long &n) {
15|         if (n < 2) return false;
16|         if (n < 4) return true;
17|         if (n & 1) return false;
18|         for (int i = 0; i < 12 && BASE[i] < n; ++i) if (!
19|             check (n, BASE[i])) return false;
20|         return true; } };
```

4.12 Pollard's Rho algorithm

```

1| /* Pollard Rho : factorizes an integer. */
2| struct pollard_rho {
3|     miller_rabin is_prime;
4|     const long long thr = 13E9;
5|     long long facize (const long long &n, const long long
6|         &seed) {
7|         long long x = rand () % (n - 1) + 1, y = x;
8|         for (int head = 1, tail = 2; ; ) {
9|             x = mul_mod (x, x, n);
10|            x = (x + seed) % n;
11|            if (x == y) return n;
12|            long long ans = std::gcd (std::abs (x - y), n);
13|            if (ans > 1 && ans < n) return ans;
14|            if (++head == tail) { y = x; tail <= 1; } } }
15|     void search (const long long &n, std::vector <long
16|         long> &div) {
17|         if (n > 1) {
18|             if (is_prime.solve (n)) div.push_back (n);
19|             else {
20|                 long long fac = n;
21|                 for (; fac >= n; fac = facize (n, rand () % (n -
22|                     1) + 1));
23|                 search (n / fac, div); search (fac, div); } } }
24|     std::vector <long long> solve (const long long &n) {
25|         std::vector <long long> ans;
26|         if (n > thr) search (n, ans);
27|         else {
28|             long long rem = n;
29|             for (long long i = 2; i * i <= rem; ++i)
30|                 while (!(rem % i)) { ans.push_back (i); rem /= i; }
31|             if (rem > 1) ans.push_back (rem); }
32|         return ans; } };
```

5 Geometry

```

1| #define cd const double &
2| const double EPS = 1E-8, PI = acos (-1);
3| int sgn (cd x) { return x < -EPS ? -1 : x > EPS; }
4| int cmp (cd x, cd y) { return sgn (x - y); }
5| double sqr (cd x) { return x * x; }
```

5.1 Point

```

1| #define cp const point &
2| struct point {
3|     double x, y;
4|     explicit point (cd x = 0, cd y = 0) : x (x), y (y) {}
5|     int dim () const { return sgn (y) == 0 ? sgn (x) < 0
6|         : sgn (y) < 0; }
7|     point unit () const { double l = sqrt (x * x + y * y)
8|         ; return point (x / l, y / l); }
9|     //counter-clockwise
10|     point rot90 () const { return point (-y, x); }
11|     //clockwise
12|     point_rot90 () const { return point (y, -x); }
13|     point_rot (cd t) const {
14|         double c = cos (t), s = sin (t);
15|         return point (x * c - y * s, x * s + y * c); } }
16| bool operator == (cp a, cp b) { return cmp (a.x, b.x)
17|     == 0 && cmp (a.y, b.y) == 0; }
18| bool operator != (cp a, cp b) { return cmp (a.x, b.x)
19|     != 0 || cmp (a.y, b.y) != 0; }
20| bool operator < (cp a, cp b) { return (cmp (a.x, b.x)
21|     == 0) ? cmp (a.y, b.y) < 0 : cmp (a.x, b.x) < 0; }
22| point operator - (cp a, cp b) { return point (-a.x, -a.y); }
23| point operator + (cp a, cp b) { return point (a.x + b.
24|     x, a.y + b.y); }
25| point operator - (cp a, cp b) { return point (a.x - b.
26|     x, a.y - b.y); }
```

```

20| point operator * (cp a, cd b) { return point (a.x * b,
21|     a.y * b); }
22| point operator / (cp a, cd b) { return point (a.x / b,
23|     a.y / b); }
24| double dot (cp a, cp b) { return a.x * b.x + a.y * b.y
25|     ; }
26| double det (cp a, cp b) { return a.x * b.y - a.y * b.x
27|     ; }
28| double dis2 (cp a, cp b = point ()) { return sqr (a.x
29|     - b.x) + sqr (a.y - b.y); }
30| double dis (cp a, cp b = point ()) { return sqrt (dis2
31|     (a, b)); }
```

5.2 Line

```

1| #define cl const line &
2| struct line {
3|     point s, t;
4|     explicit line (cp s = point (), cp t = point ()) : s
5|         (s), t (t) {} }
6| bool point_on_segment (cp a, cl b) { return sgn (det (
7|     a - b.s, b.t - b.s)) == 0 && sgn (dot (b.s - a, b.
8|     t - a)) <= 0; }
9| bool two_side (cp a, cp b, cl c) { return sgn (det (a
10|     - c.s, c.t - c.s)) * sgn (det (b - c.s, c.t - c.s))
11|     < 0; }
12| bool intersect_judgment (cl a, cl b) {
13|     if (point_on_segment (b.s, a) || point_on_segment (b.
14|         t, a)) return true;
15|     if (point_on_segment (a.s, b) || point_on_segment (a.
16|         t, b)) return true;
17|     return two_side (a.s, a.t, b) && two_side (b.s, b.t,
18|         a); }
19| point line_intersect (cl a, cl b) {
20|     double s1 = det (a.t - a.s, b.s - a.s), s2 = det (a.t
21|         - a.s, b.t - a.s);
22|     return (b.s * s2 - b.t * s1) / (s2 - s1); }
23| double point_to_line (cp a, cl b) { return fabs (det (
24|     b.t - b.s, a - b.s)) / dis (b.s, b.t); }
25| point project_to_line (cp a, cl b) { return b.s + (b.t
26|     - b.s) * (dot (a - b.s, b.t - b.s) / dis2 (b.t, b.
27|         s)); }
28| double point_to_segment (cp a, cl b) {
29|     if (sgn (dot (b.s - a, b.t - b.s)) <= 0) return fabs (det (b.t - a, b.
30|         s - b.s)) / dis (b.s, b.t);
31|     return std::min (dis (a, b.s), dis (a, b.t)); }
32| bool in_polygon (cp p, const std::vector <point> &po)
33| {
34|     int n = (int) po.size (), counter = 0;
35|     for (int i = 0; i < n; ++i) {
36|         point a = po[i], b = po[(i + 1) % n];
37|         //Modify the next line if necessary.
38|         if (point_on_segment (p, line (a, b))) return true;
39|         int x = sgn (det (p - a, b - a)), y = sgn (a.y - p.y
40|             ), z = sgn (b.y - p.y);
41|         if (x > 0 && y <= 0 && z > 0) counter++;
42|         if (x < 0 && z <= 0 && y > 0) counter--; }
43|     return counter != 0; }
44| double polygon_area (const std::vector <point> &a) {
45|     double ans = 0.0;
46|     for (int i = 0; i < (int) a.size (); ++i) ans += det
47|         (a[i], a[(i + 1) % a.size ()]) / 2.0;
48|     return ans; }
```

5.3 Circle

```

1| #define cc const circle &
2| struct circle {
3|     point c; double r;
4|     explicit circle (point c = point (), double r = 0) :
5|         c (c), r (r) {} }
6| bool operator == (cc a, cc b) { return a.c == b.c &&
7|     cmp (a.r, b.r) == 0; }
8| bool operator != (cc a, cc b) { return !(a == b); }
9| bool in_circle (cp a, cc b) { return cmp (dis (a, b.c)
10|     , b.r) <= 0; }
11| circle make_circle (cp a, cp b) { return circle ((a +
12|     b) / 2, dis (a, b) / 2); }
13| circle make_circle (cp a, cp b, cp c) { point p =
14|     circumcenter (a, b, c); return circle (p, dis (p,
15|         a)); }
16| //In the order of the line vector.
17| std::vector <point> line_circle_intersect (cl a, cc b)
18| {
19|     if (cmp (point_to_line (b.c, a), b.r) > 0) return std
20|         ::vector <point> ();
21|     double x = sqrt (sqr (b.r) - sqr (point_to_line (b.c,
22|         a)));
23|     return std::vector <point> ({project_to_line (b.c, a)
24|         + (a.s - a.t).unit () * x, project_to_line (b.c,
25|         a) - (a.s - a.t).unit () * x}); }
26| double circle_intersect_area (cc a, cc b) {
27|     double d = dis (a.c, b.c);
28|     if (sgn (d - (a.r + b.r)) >= 0) return 0;
29|     if (sgn (d - abs (a.r - b.r)) <= 0) {
30|         double r = std::min (a.r, b.r); return r * r * PI; }
31|     double x = (d * d + a.r * a.r - b.r * b.r) / (2 * d),
32|         t1 = acos (min (1., max (-1., x / a.r))), t2 =
33|         acos (min (1., max (-1., (d - x) / b.r)));
34|     return a.r * a.r * t1 + b.r * b.r * t2 - d * a.r *
35|         sin (t1); }
36| //Counter-clockwise with respect of vector OaOb.
37| std::vector <point> circle_intersect (cc a, cc b) {
38|     if (a.c == b.c || cmp (dis (a.c, b.c), a.r + b.r) > 0
39|         || cmp (dis (a.c, b.c), std::abs (a.r - b.r)) <
40|         0) return std::vector <point> ();
41|     point r = (b.c - a.c).unit ();
42|     double d = dis (a.c, b.c);
43|     double x = ((sqr (a.r) - sqr (b.r)) / d + d) / 2;
44|     double h = sqrt (sqr (a.r) - sqr (x));
45|     if (sgn (h) == 0) return std::vector <point> ({a.c +
46|         r * x}); }
```

```

30 return std::vector<point> ({a.c + r * x - r.rot90 ()
   * h, a.c + r * x + r.rot90 () * h}); }
31 //Counter-clockwise with respect of point a.
32 std::vector<point> tangent (cp a, cc b) { circle p =
   make_circle (a, b.c); return circle_intersect (p,
   b); }
33 //Counter-clockwise with respect of point Oa.
34 std::vector<line> extangent (cc a, cc b) {
35 std::vector<line> ret;
36 if (cmp (dis (a.c, b.c), std::abs (a.r - b.r)) <= 0)
   return ret;
37 if (sgn (a.r - b.r) == 0) {
38 point dir = b.c - a.c; dir = (dir * a.r / dis (dir))
   .rot90 ();
39 ret.push_back (line (a.c - dir, b.c - dir));
40 ret.push_back (line (a.c + dir, b.c + dir));
41 } else {
42 point p = (b.c * a.r - a.c * b.r) / (a.r - b.r);
43 std::vector pp = tangent (p, a), qq = tangent (p, b)
   ;
44 if (pp.size () == 2 && qq.size () == 2) {
45 if (cmp (a.r, b.r) < 0) std::swap (pp[0], pp[1]),
   std::swap (qq[0], qq[1]);
46 ret.push_back (line (pp[0], qq[0]));
47 ret.push_back (line (pp[1], qq[1])); } }
48 return ret; }
49 //Counter-clockwise with respect of point Oa.
50 std::vector<line> intangent (cc c1, cc c2) {
51 point p = (b.c * a.r + a.c * b.r) / (a.r + b.r);
52 std::vector pp = tangent (p, a), qq = tangent (p, b);
53 if (pp.size () == 2 && qq.size () == 2) {
54 ret.push_back (line (pp[0], qq[0]));
55 ret.push_back (line (pp[1], qq[1])); }
56 return ret; }

```

5.4 Centers of a triangle

```

1 point incenter (cp a, cp b, cp c) {
2 double p = dis (a, b) + dis (b, c) + dis (c, a);
3 return (a * dis (b, c) + b * dis (c, a) + c * dis (a,
   b)) / p; }
4 point circumcenter (cp a, cp b, cp c) {
5 point p = b - a, q = c - a, s (dot (p, p) / 2, dot (q,
   q) / 2);
6 return a + point (det (s, point (p.y, q.y)), det (
   point (p.x, q.x), s)) / det (p, q); }
7 point orthocenter (cp a, cp b, cp c) { return a + b +
   c - circumcenter (a, b, c) * 2; }

```

5.5 Fermat point

```

1 /* Fermat point : finds a point P that minimizes
   |PA|+|PB|+|PC|. */
2 point fermat_point (cp a, cp b, cp c) {
3 if (a == b) return a; if (b == c) return b; if (c ==
   a) return c;
4 double ab = dis (a, b), bc = dis (b, c), ca = dis (c,
   a);
5 double cosa = dot (b - a, c - a) / ab / ca;
6 double cosb = dot (a - b, c - b) / ab / bc;
7 double cosc = dot (b - c, a - c) / ca / bc;
8 double sq3 = PI / 3.0; point mid;
9 if (sgn (cosa + 0.5) < 0) mid = a;
10 else if (sgn (cosb + 0.5) < 0) mid = b;
11 else if (sgn (cosc + 0.5) < 0) mid = c;
12 else if (sgn (det (b - a, c - a)) < 0) mid =
   line_intersect (line (a, b + (c - b).rot (sq3)),
   line (b, c + (a - c).rot (sq3)));
13 else mid = line_intersect (line (a, c + (b - c).rot (
   sq3)), line (c, b + (a - b).rot (sq3)));
14 return mid; }

```

5.6 Convex hull

```

1 //Counter-clockwise, with minimum number of points.
2 bool turn_left (cp a, cp b, cp c) { return sgn (det (b
   - a, c - a)) >= 0; }
3 std::vector<point> convex_hull (std::vector<point> a
   ) {
4 int cnt = 0; std::sort (a.begin (), a.end ());
5 static std::vector<point> ret; ret.resize (a.size ()
   << 1);
6 for (int i = 0; i < (int) a.size (); ++i) {
7 while (cnt > 1 && turn_left (ret[cnt - 2], a[i], ret
   [cnt - 1])) --cnt;
8 ret[cnt++] = a[i]; }
9 int fixed = cnt;
10 for (int i = (int) a.size () - 1; i >= 0; --i) {
11 while (cnt > fixed && turn_left (ret[cnt - 2], a[i],
   ret[cnt - 1])) --cnt;
12 ret[cnt++] = a[i]; }
13 return std::vector<point> (ret.begin (), ret.begin
   () + cnt - 1); }

```

5.7 Half plane intersection

```

1 /* Online half plane intersection : complexity O(n)
   each operation. */
2 std::vector<point> cut (const std::vector<point> &c,
   line p) {
3 std::vector<point> ret;
4 if (c.empty ()) return ret;
5 for (int i = 0; i < (int) c.size (); ++i) {
6 int j = (i + 1) % (int) c.size ();
7 if (turn_left (p.s, p.t, c[i])) ret.push_back (c[i])
   ;
8 if (two_side (c[i], c[j], p)) ret.push_back (
   line_intersect (p, line (c[i], c[j]))); }
9 return ret; }

```

```

10 /* Offline half plane intersection : complexity
   O(n log n). */
11 bool turn_left (cl l, cp p) { return turn_left (l.s, l
   .t, p); }
12 int cmp (cp a, cp b) { return a.dim () != b.dim () ? (
   a.dim () < b.dim () ? -1 : 1) : -sgn (det (a, b)); }
13 std::vector<point> half_plane_intersect (std::vector<vector
   <line> h) {
14 typedef std::pair<point, line> polar;
15 std::vector<polar> g; g.resize (h.size ());
16 for (int i = 0; i < (int) h.size (); ++i) g[i] = std
   ::make_pair (h[i].t - h[i].s, h[i]);
17 sort (g.begin (), g.end (), [&] (const polar &a,
   const polar &b) {
18 if (cmp (a.first, b.first) == 0) return sgn (det (a.
   second.t - a.second.s, b.second.t - a.second.s))
   < 0;
19 else return cmp (a.first, b.first) < 0; });
20 h.resize (std::unique (g.begin (), g.end ()), [] (
   const polar &a, const polar &b) { return cmp (a.
   first, b.first) == 0 } - g.begin ());
21 for (int i = 0; i < (int) h.size (); ++i) h[i] = g[i
   ].second;
22 int fore = 0, rear = -1; std::vector<line> ret (h.
   size (), line ());
23 for (int i = 0; i < (int) h.size (); ++i) {
24 while (fore < rear && !turn_left (h[i],
   line_intersect (ret[rear - 1], ret[rear]))) --
   rear;
25 while (fore < rear && !turn_left (h[i],
   line_intersect (ret[fore], ret[fore + 1]))) ++
   fore;
26 ret[++rear] = h[i]; }
27 while (rear - fore > 1 && !turn_left (ret[fore],
   line_intersect (ret[rear - 1], ret[rear]))) --
   rear;
28 while (rear - fore > 1 && !turn_left (ret[rear],
   line_intersect (ret[fore], ret[fore + 1]))) ++
   fore;
29 if (rear - fore < 2) return std::vector<point> ();
30 std::vector<point> ans; ans.resize (rear + 1);
31 for (int i = 0; i < rear + 1; ++i) ans[i] =
   line_intersect (ret[i], ret[(i + 1) % (rear + 1)
   ]);
32 return ans; }

```

5.8 Nearest pair of points

```

1 /* Nearest pair of points : [l, r), need to sort p
   first. */
2 double solve (std::vector<point> &p, int l, int r) {
3 if (l + 1 >= r) return INF;
4 int m = (l + r) / 2; double mx = p[m].x; std::vector
   <point> v;
5 double ret = std::min (solve(p, l, m), solve(p, m, r)
   );
6 for (int i = l; i < r; ++i)
7 if (sqr (p[i].x - mx) < ret) v.push_back (p[i]);
8 sort (v.begin (), v.end (), [&] (cp a, cp b) { return
   a.y < b.y; });
9 for (int i = 0; i < v.size (); ++i)
10 for (int j = i + 1; j < v.size (); ++j) {
11 if (sqr (v[i].y - v[j].y) > ret) break;
12 ret = min (ret, dis2 (v[i] - v[j])); }
13 return ret; }

```

5.9 Minimum circle

```

1 circle minimum_circle (std::vector<point> p) {
2 circle ret; std::random_shuffle (p.begin (), p.end ()
   );
3 for (int i = 0; i < (int) p.size (); ++i) if (!
   in_circle (p[i], ret)) {
4 ret = circle (p[i], 0); for (int j = 0; j < i; ++j)
   if (!in_circle (p[j], ret)) {
5 ret = make_circle (p[j], p[i]); for (int k = 0; k <
   j; ++k)
6 if (!in_circle (p[k], ret)) ret = make_circle (p[i]
   ], p[j], p[k]); } }
7 return ret; }

```

5.10 Intersection of a polygon and a circle

```

1 struct polygon_circle_intersect {
2 double sector_area (cp a, cp b, const double &r) {
3 double c = (2.0 * r * r - dis2 (a, b)) / (2.0 * r *
   r);
4 return r * r * acos (c) / 2.0; }
5 double area (cp a, cp b, const double &r) {
6 double dA = dot (a, a), dB = dot (b, b), dC =
   point_to_segment (point (), line (a, b));
7 if (sgn (dA - r * r) <= 0 && sgn (dB - r * r) <= 0)
   return det (a, b) / 2.0;
8 point tA = a.unit () * r, tB = b.unit () * r;
9 if (sgn (dC - r) > 0) return sector_area (tA, tB, r)
   ;
10 std::vector<point> ret = line_circle_intersect (
   line (a, b), circle (point (), r));
11 if (sgn (dA - r * r) > 0 && sgn (dB - r * r) > 0)
   return sector_area (tA, ret[0], r) + det (ret[0],
   ret[1]) / 2.0 + sector_area (ret[1], tB, r);
12 if (sgn (dA - r * r) > 0) return det (ret[0], b) /
   2.0 + sector_area (tA, ret[0], r);
13 else return det (a, ret[1]) / 2.0 + sector_area (ret
   [1], tB, r); }
14 double solve (const std::vector<point> &p, cc c) {
15 double ret = 0.0;
16 for (int i = 0; i < (int) p.size (); ++i) {
17 int s = sgn (det (p[i] - c.c, p[(i + 1) % p.size ()
   ] - c.c));
18 }

```



```

19 if (s > 0) ret += area (p[i] - c.c, p[(i + 1) % p.
    size ()] - c.c, c.r);
20 else ret -= area (p[(i + 1) % p.size ()] - c.c, p[i]
    - c.c, c.r);
21 return std::abs (ret); } };

```

5.11 Union of circles

```

1 template <int MAXN = 500> struct union_circle {
2 int C; circle c[MAXN]; double area[MAXN];
3 struct event {
4 point p; double ang; int delta;
5 event (cp p = point (), double ang = 0, int delta =
6 0) : p(p), ang(ang), delta(delta) {}
7 bool operator < (const event &a) { return ang < a.
    ang; }
8 void addevent(cc a, cc b, std::vector <event> &evt,
    int &cnt) {
9 double d2 = dis2 (a.c, b.c), d_ratio = ((a.r - b.r)
    * (a.r + b.r) / d2 + 1) / 2;
10 p_ratio = sqrt (std::max (0., -(d2 - sqr(a.r - b.r)
    ) * (d2 - sqr(a.r + b.r)) / (d2 * d2 * 4)));
11 point d = b.c - a.c, p = d.rot(Pi / 2), q0 = a.c + d
    * d_ratio + p * p_ratio, q1 = a.c + d * d_ratio
    - p * p_ratio;
12 double ang0 = atan2 ((q0 - a.c).y, (q0 - a.c).x),
    ang1 = atan2 ((q1 - a.c).x, (q1 - a.c).y);
13 evt.emplace_back(q1, ang1, 1); evt.emplace_back(q0,
    ang0, -1); cnt += ang1 > ang0;
14 bool same(cc a, cc b) { return sgn (dis (a.c, b.c))
    == 0 && sgn (a.r - b.r) == 0; }
15 bool overlap(cc a, cc b) { return sgn (a.r - b.r -
    dis (a.c, b.c)) >= 0; }
16 bool intersect(cc a, cc b) { return sgn (dis (a.c, b.
    c) - a.r - b.r) < 0; }
17 void solve() {
18 std::fill (area, area + C + 2, 0);
19 for (int i = 0; i < C; ++i) {
20 int cnt = 1; std::vector <event> evt;
21 for (int j = 0; j < i; ++j) if (same (c[i], c[j]))
    ++cnt;
22 for (int j = 0; j < C; ++j) if (j != i && !same (c[
    i], c[j]) && overlap (c[j], c[i])) ++cnt;
23 for (int j = 0; j < C; ++j) if (j != i && !overlap
    (c[j], c[i]) && !intersect (c[i], c[j]) &&
    intersect (c[i], c[j])) {
24 addevent (c[i], c[j], evt, cnt);
25 if (evt.empty ()) area[cnt] += Pi * c[i].r * c[i].r
    ;
26 else {
27 std::sort (evt.begin (), evt.end ());
28 evt.push_back (evt.front ());
29 for (int j = 0; j + 1 < (int) evt.size (); ++j) {
30 cnt += evt[j].delta; area[cnt] += det(evt[j].p,
    evt[j + 1].p) / 2;
31 double ang = evt[j + 1].ang - evt[j].ang; if (ang
    < 0) ang += Pi * 2;
32 area[cnt] += ang * c[i].r * c[i].r / 2 - sin(ang)
    * c[i].r * c[i].r / 2; } } } };

```

5.12 3D point

```

1 #define cp3 const point3 &
2 struct point3 {
3 double x, y, z;
4 explicit point3 (cd x = 0, cd y = 0, cd z = 0) : x (x
    ), y (y), z (z) {} };
5 point3 operator + (cp3 a, cp3 b) { return point3 (a.x
    + b.x, a.y + b.y, a.z + b.z); }
6 point3 operator - (cp3 a, cp3 b) { return point3 (a.x
    - b.x, a.y - b.y, a.z - b.z); }
7 point3 operator * (cp3 a, cd b) { return point3 (a.x *
    b, a.y * b, a.z * b); }
8 point3 operator / (cp3 a, cd b) { return point3 (a.x /
    b, a.y / b, a.z / b); }
9 double dot (cp3 a, cp3 b) { return a.x * b.x + a.y * b
    .y + a.z * b.z; }
10 point3 det (cp3 a, cp3 b) { return point3 (a.y * b.z -
    a.z * b.y, -a.x * b.z + a.z * b.x, a.x * b.y - a.
    y * b.x); }
11 double dis2 (cp3 a, cp3 b = point3 ()) { return sqr (a
    .x - b.x) + sqr (a.y - b.y) + sqr (a.z - b.z); }
12 double dis (cp3 a, cp3 b = point3 ()) { return sqrt (
    dis2 (a, b)); }
13 //right-handed, if x+ -> y+ is right-handed
14 point3 rotate(cp3 p, cp3 axis, double w) {
15 double x = axis.x, y = axis.y, z = axis.z;
16 double s = x * x + y * y + z * z, ss = sqrt(s), cosw
    = cos(w), sinw = sin(w);
17 double a[4][4]; memset(a, 0, sizeof (a));
18 a[3][3] = 1;
19 a[0][0] = ((y * y + z * z) * cosw + x * x) / s;
20 a[0][1] = x * y * (1 - cosw) / s + z * sinw / ss;
21 a[0][2] = x * z * (1 - cosw) / s - y * sinw / ss;
22 a[1][0] = x * y * (1 - cosw) / s - z * sinw / ss;
23 a[1][1] = ((x * x + z * z) * cosw + y * y) / s;
24 a[1][2] = y * z * (1 - cosw) / s + x * sinw / ss;
25 a[2][0] = x * z * (1 - cosw) / s + y * sinw / ss;
26 a[2][1] = y * z * (1 - cosw) / s - x * sinw / ss;
27 a[2][2] = ((x * x + y * y) * cosw + z * z) / s;
28 double ans[4] = {0, 0, 0, 0}, c[4] = {p.x, p.y, p.z,
    1};
29 for (int i = 0; i < 4; ++i) for (int j = 0; j < 4; ++
    j)
30 ans[i] += a[j][i] * c[j];
31 return point3 (ans[0], ans[1], ans[2]);
32 }

```

5.13 3D line

```

1 #define cl3 const line3 &
2 struct line3 {
3 point3 s, t;
4 explicit line3 (cp3 s = point3 (), cp3 t = point3 ())
    : s (s), t (t) {} };
5 point3 line_plane_intersection (cl3 a, cl3 b) { return
    a.s + (a.t - a.s) * dot (b.s - a.s, b.t - b.s) /
    dot (a.t - a.s, b.t - b.s); }
6 line3 plane_intersection (cl3 a, cl3 b) {
7 point3 p = det (a.t - a.s, b.t - b.s), q = det (a.t -
    a.s, p), s = line_plane_intersection (line3 (a.s
    , a.s + q), b);
8 return line3 (s, s + p); }
9 point3 project_to_plane (cp3 a, cl3 b) { return a + (b
    .t - b.s) * dot (b.t - b.s, b.s - a) / dis2 (b.t -
    b.s); }

```

5.14 3D convex hull

```

1 /* 3D convex hull : initializes n and p / outputs face
    . */
2 template <int MAXN = 500>
3 struct convex_hull3 {
4 double mix (cp3 a, cp3 b, cp3 c) { return dot (det (a
    , b), c); }
5 double volume (cp3 a, cp3 b, cp3 c, cp3 d) { return
    mix (b - a, c - a, d - a); }
6 struct tri {
7 int a, b, c;
8 tri() {}
9 tri(int _a, int _b, int _c) : a(_a), b(_b), c(_c) {}
10 double area() const { return dis (det (p[b] - p[a],
    p[c] - p[a])) / 2; }
11 point3 normal() const { return det (p[b] - p[a], p[c]
    - p[a]).unit (); }
12 double dis (cp3 p0) const { return dot (normal (),
    p0 - p[a]); } };
13 int n; std::vector <point3> p;
14 std::vector <tri> face, tmp;
15 int mark[MAXN][MAXN], time;
16 void add (int v) {
17 ++time; tmp.clear ();
18 for (int i = 0; i < (int) face.size (); ++i) {
19 int a = face[i].a, b = face[i].b, c = face[i].c;
20 if (sgn (volume (p[v], p[a], p[b], p[c])) > 0)
21 mark[a][b] = mark[b][a] = mark[a][c] = mark[c][a]
    = mark[b][c] = mark[c][b] = time;
22 else tmp.push_back (face[i]); }
23 face.clear (); face = tmp;
24 for (int i = 0; i < (int) tmp.size (); ++i) {
25 int a = face[i].a, b = face[i].b, c = face[i].c;
26 if (mark[a][b] == time) face.emplace_back (v, b, a)
    ;
27 if (mark[b][c] == time) face.emplace_back (v, c, b)
    ;
28 if (mark[c][a] == time) face.emplace_back (v, a, c)
    ; } }
29 void reorder () {
30 for (int i = 2; i < n; ++i) {
31 point3 tmp = det (p[i] - p[0], p[i] - p[1]);
32 if (sgn (dis (tmp))) {
33 std::swap (p[i], p[2]);
34 for (int j = 3; j < n; ++j)
35 if (sgn (volume (p[0], p[1], p[2], p[j]))) {
36 std::swap (p[j], p[3]); return; } } } }
37 void build_convex () {
38 reorder (); face.clear ();
39 face.emplace_back (0, 1, 2);
40 face.emplace_back (0, 2, 1);
41 for (int i = 3; i < n; ++i) add(i); } };

```

6 Graph

```

1 template <int MAXN = 100000, int MAXM = 100000>
2 struct edge_list {
3 int size, begin[MAXN], dest[MAXM], next[MAXM];
4 void clear (int n) { size = 0; std::fill (begin,
    begin + n, -1); }
5 edge_list (int n = MAXN) { clear (n); }
6 void add_edge (int u, int v) { dest[size] = v; next[
    size] = begin[u]; begin[u] = size++; } };
7 template <int MAXN = 100000, int MAXM = 100000>
8 struct cost_edge_list {
9 int size, begin[MAXN], dest[MAXM], next[MAXM], cost[
    MAXM];
10 void clear (int n) { size = 0; std::fill (begin,
    begin + n, -1); }
11 cost_edge_list (int n = MAXN) { clear (n); }
12 void add_edge (int u, int v, int c) { dest[size] = v;
    next[size] = begin[u]; cost[size] = c; begin[u]
    = size++; } };

```

6.1 Hopcroft-Karp algorithm

```

1 /* Hopcroft-Karp algorithm : unweighted maximum
    matching for bipartition graphs with complexity
    O(m*sqrt(n)). */
2 template <int MAXN = 100000, int MAXM = 100000>
3 struct hopcroft_karp {
4 using edge_list = std::vector <int> [MAXN];
5 int mx[MAXN], my[MAXM], lv[MAXN];
6 bool dfs (edge_list <MAXN, MAXM> &e, int x) {
7 for (int i = e.begin[x]; ~i; i = e.next[i]) {
8 int y = e.dest[i], w = my[y];
9 if (!w || (lv[x] + 1 == lv[w] && dfs (e, w))) {
10 mx[x] = y; my[y] = x; return true; } }
11 lv[x] = -1; return false; }
12 int solve (edge_list <MAXN, MAXM> &e, int n, int m) {
13 std::fill (mx, mx + n, -1); std::fill (my, my + m,
    -1);

```



```

46     *qtail++ = match[dest];
47     } else if (d[u].find(dest) == 0) {
48         int b = lca(loc, dest, root);
49         contract(loc, dest, b); contract(dest, loc, b)
50         ; } }
51     return 0; }
52 int solve (int n, const edge_list <MAXN, MAXN> &e) {
53     std::fill (fa, fa + n, 0); std::fill (c1, c1 + n, 0)
54     ;
55     std::fill (c2, c2 + n, 0); std::fill (match, match +
56         n, -1);
57     int re = 0; for (int i = 0; i < n; i++)
58         if (match[i] == -1) if (bfs (i, n, e)) ++re; else
59             match[i] = -2;
60     return re; } };

```

6.4 Weighted blossom algorithm

```

1  /* Weighted blossom algorithm (vfleaking ver.) :
2     maximum matching for general weighted graphs with
3     complexity  $O(n^3)$ .
4  Usage : Set n to the size of the vertices. Run init ()
5     . Set g[u][v].w to the weight of the edge. Run solve
6     ().
7  The first result is the answer, the second one is the
8     number of matching pairs. Obtain the matching with
9     match[].
10 Note : 1-based. */
11 struct weighted_blossom {
12     static const int INF = INT_MAX, MAXN = 400;
13     struct edge { int u, v, w; edge (int u = 0, int v = 0,
14         int w = 0) : u(u), v(v), w(w) {} };
15     int n, n_x;
16     edge g[MAXN * 2 + 1][MAXN * 2 + 1];
17     int lab[MAXN * 2 + 1], match[MAXN * 2 + 1], slack[
18         MAXN * 2 + 1], st[MAXN * 2 + 1], pa[MAXN * 2 +
19         1];
20     int flower_from[MAXN * 2 + 1][MAXN + 1], S[MAXN * 2 +
21         1], vis[MAXN * 2 + 1];
22     std::vector<int> flower[MAXN * 2 + 1];
23     std::queue<int> q;
24     int e_delta (const edge &e) { return lab[e.u] + lab[e
25         .v] - g[e.u][e.v].w * 2; };
26     void update_slack (int u, int x) { if (!slack[x] ||
27         e_delta (g[u][x]) < e_delta (g[slack[x]][x]))
28         slack[x] = u; };
29     void set_slack (int x) { slack[x] = 0; for (int u =
30         1; u <= n; ++u) if (g[u][x].w > 0 && st[u] != x &&
31         S[st[u]] == 0)
32         update_slack(u, x); }
33     void q_push (int x) {
34         if (x <= n) q.push (x);
35         else for (size_t i = 0; i < flower[x].size (); ++i)
36             q.push (flower[x][i]); }
37     void set_st (int x, int b) {
38         st[x] = b; if (x > n) for (size_t i = 0; i < flower[
39             x].size (); ++i) set_st (flower[x][i], b); }
40     int get_pr (int b, int xr) {
41         int pr = std::find (flower[b].begin (), flower[b].
42             end (), xr) - flower[b].begin ();
43         if (pr % 2 == 1) { std::reverse (flower[b].begin ()
44             + 1, flower[b].end ()); return (int) flower[b].
45             size () - pr; }
46         else return pr; }
47     void set_match (int u, int v) {
48         match[u] = g[u][v].v; if (u > n) {
49             edge e = g[u][v]; int xr = flower_from[u][e.u], pr
50             = get_pr (u, xr);
51             for (int i = 0; i < pr; ++i) set_match (flower[u][i
52                 ], flower[u][i ^ 1]);
53             set_match (xr, v); std::rotate (flower[u].begin (),
54                 flower[u].begin () + pr, flower[u].end ()); }
55         }
56     void augment (int u, int v) {
57         for (; ; ) {
58             int xnv = st[match[u]]; set_match (u, v);
59             if (!xnv) return; set_match (xnv, st[pa[xnv]]);
60             u = st[pa[xnv]], v = xnv; } }
61     int get_lca (int u, int v) {
62         static int t = 0;
63         for (++t; u || v; std::swap (u, v)) {
64             if (u == 0) continue; if (vis[u] == t) return u;
65             vis[u] = t; u = st[match[u]]; if (u) u = st[pa[u]];
66         }
67         return 0; }
68     void add_blossom (int u, int lca, int v) {
69         int b = n + 1; while (b <= n_x && st[b]) ++b;
70         if (b > n_x) ++n_x;
71         lab[b] = 0, S[b] = 0;
72         match[b] = match[lca]; flower[b].clear ();
73         flower[b].push_back (lca);
74         for (int x = u, y; x != lca; x = st[pa[y]]) {
75             flower[b].push_back (x), flower[b].push_back (y =
76                 st[match[x]]), q.push (y); }
77         std::reverse (flower[b].begin () + 1, flower[b].end
78             ());
79         for (int x = v, y; x != lca; x = st[pa[y]]) {
80             flower[b].push_back (x), flower[b].push_back (y =
81                 st[match[x]]), q.push (y); }
82         set_st (b, b);
83         for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].w
84             = 0;
85         for (int x = 1; x <= n; ++x) flower_from[b][x] = 0;
86         for (size_t i = 0; i < flower[b].size (); ++i) {
87             int xs = flower[b][i];
88             for (int x = 1; x <= n_x; ++x) if (g[b][x].w == 0
89                 || e_delta (g[xs][x]) < e_delta (g[b][x]))
90                 g[b][x] = g[xs][x], g[x][b] = g[x][xs];
91             for (int x = 1; x <= n; ++x) if (flower_from[xs][x])
92                 flower_from[b][x] = xs; }
93         set_slack (b); }
94     void expand_blossom (int b) {

```

```

63 for (size_t i = 0; i < flower[b].size (); ++i)
64   set_st (flower[b][i], flower[b][i]);
65 int xr = flower_from[b][g[b][pa[b]].u], pr = get_pr(
66   b, xr);
67 for (int i = 0; i < pr; i += 2) {
68   int xs = flower[b][i], xns = flower[b][i + 1];
69   pa[xs] = g[xns][xs].u; S[xs] = 1, S[xns] = 0;
70   slack[xs] = 0, set_slack(xns); q_push(xns);
71   S[xr] = 1, pa[xr] = pa[b];
72   for (size_t i = pr + 1; i < flower[b].size (); ++i)
73     {
74       int xs = flower[b][i]; S[xs] = -1, set_slack(xs);
75       st[b] = 0;
76     }
77 bool on_found_edge (const edge &e) {
78   int u = st[e.u], v = st[e.v];
79   if (S[v] == -1) {
80     pa[v] = e.u, S[v] = 1; int nu = st[match[v]];
81     slack[v] = slack[nu] = 0; S[nu] = 0, q_push(nu);
82   } else if (S[v] == 0) {
83     int lca = get_lca(u, v);
84     if (!lca) return augment(u, v), augment(v, u), true;
85     else add_blossom(u, lca, v);
86   }
87   return false;
88 }
89 bool matching () {
90   memset (S + 1, -1, sizeof (int) * n_x);
91   memset (slack + 1, 0, sizeof (int) * n_x);
92   q = std::queue<int> ();
93   for (int x = 1; x <= n_x; ++x) if (st[x] == x && !
94     match[x]) pa[x] = 0, S[x] = 0, q_push (x);
95   if (q.empty ()) return false;
96   for (; ; ) {
97     while (q.size ()) {
98       int u = q.front (); q.pop ();
99       if (S[st[u]] == 1) continue;
100       for (int v = 1; v <= n_x; ++v) if (g[u][v].w > 0 &&
101         st[u] != st[v]) {
102         if (e_delta (g[u][v]) == 0) {
103           if (on_found_edge (g[u][v])) return true;
104           else update_slack (u, st[v]);
105         }
106         int d = INF;
107         for (int b = n + 1; b <= n_x; ++b) if (st[b] == b &&
108           S[b] == 1) d = std::min (d, lab[b] / 2);
109         for (int x = 1; x <= n_x; ++x) if (st[x] == x &&
110           slack[x]) {
111           if (S[x] == -1) d = std::min (d, e_delta (g[slack[
112             x]][x]));
113           else if (S[x] == 0) d = std::min (d, e_delta (g[
114             slack[x]][x]) / 2);
115         }
116         for (int u = 1; u <= n; ++u) {
117           if (S[st[u]] == 0) {
118             if (lab[u] <= d) return 0;
119             lab[u] -= d;
120           } else if (S[st[u]] == 1) lab[u] += d;
121         }
122         for (int b = n + 1; b <= n_x; ++b)
123           if (st[b] == b) {
124             if (S[st[b]] == 0) lab[b] += d * 2;
125             else if (S[st[b]] == 1) lab[b] -= d * 2;
126           }
127         q = std::queue<int> ();
128         for (int x = 1; x <= n_x; ++x)
129           if (st[x] == x && slack[x] && st[slack[x]] != x &&
130             e_delta(g[slack[x]][x]) == 0)
131             if (on_found_edge (g[slack[x]][x])) return true;
132         for (int b = n + 1; b <= n_x; ++b) if (st[b] == b
133           && S[b] == 1 && lab[b] == 0) expand_blossom(b);
134       }
135       return false;
136     }
137     std::pair<long long, int> solve () {
138       memset (match + 1, 0, sizeof (int) * n); n_x = n;
139       int n_matches = 0; long long tot_weight = 0;
140       for (int u = 0; u <= n; ++u) st[u] = u, flower[u].
141         clear();
142       int w_max = 0;
143       for (int u = 1; u <= n; ++u) for (int v = 1; v <= n;
144         ++v) {
145         flower_from[u][v] = (u == v ? u : 0); w_max = std::
146           max (w_max, g[u][v].w);
147       }
148       for (int u = 1; u <= n; ++u) lab[u] = w_max;
149       while (matching ()) ++n_matches;
150       for (int u = 1; u <= n; ++u) if (match[u] && match[
151         u] < u) tot_weight += g[u][match[u]].w;
152       return std::make_pair (tot_weight, n_matches);
153     }
154     void init () { for (int u = 1; u <= n; ++u) for (int
155       v = 1; v <= n; ++v) g[u][v] = edge (u, v, 0); }
156   };

```

6.5 Maximum flow

```

1 /* Sparse graph maximum flow : isap.*/
2 template<int MAXN = 1000, int MAXM = 100000>
3 struct isap {
4   struct flow_edge_list {
5     int size, begin[MAXN], dest[MAXM], next[MAXM], flow[
6       MAXM];
7     void clear (int n) { size = 0; std::fill (begin,
8       begin + n, -1); }
9     flow_edge_list (int n = MAXN) { clear (n); }
10    void add_edge (int u, int v, int f) {
11      dest[size] = v; next[size] = begin[u]; flow[size] =
12      f; begin[u] = size++;
13      dest[size] = u; next[size] = begin[v]; flow[size] =
14      0; begin[v] = size++; }
15    int pre[MAXN], d[MAXN], gap[MAXN], cur[MAXN];
16    int solve (flow_edge_list &e, int n, int s, int t) {
17      for (int i = 0; i < n; ++i) { pre[i] = d[i] = gap[i]
18        = 0; cur[i] = e.begin[i]; }
19      gap[0] = n; int u = pre[s] = s, v, maxflow = 0;
20      while (d[s] < n) {
21        v = n; for (int i = cur[u]; ~i; i = e.next[i])
22          if (e.flow[i] && d[u] == d[e.dest[i]] + 1) {
23            v = e.dest[i]; cur[u] = i; break; }
24        if (v < n) {
25          pre[v] = u; u = v;
26          if (v == t) {

```

```

22   int dflow = INF, p = t; u = s;
23   while (p != s) { p = pre[p]; dflow = std::min (
24     dflow, e.flow[cur[p]]); }
25   maxflow += dflow; p = t;
26   while (p != s) { p = pre[p]; e.flow[cur[p]] -=
27     dflow; e.flow[cur[p] ^ 1] += dflow; }
28   } else {
29     int mindist = n + 1;
30     for (int i = e.begin[u]; ~i; i = e.next[i])
31       if (e.flow[i] && mindist > d[e.dest[i]]) {
32         mindist = d[e.dest[i]]; cur[u] = i;
33       }
34     if (!--gap[d[u]]) return maxflow;
35     gap[d[u] = mindist + 1]++; u = pre[u]; } }
36   return maxflow; } };
37 /* Dense graph maximum flow : dinic. */
38 template<int MAXN = 1000, int MAXM = 100000>
39 struct dinic {
40   struct flow_edge_list {
41     int size, begin[MAXN], dest[MAXM], next[MAXM], flow[
42       MAXM];
43     void clear (int n) { size = 0; std::fill (begin,
44       begin + n, -1); }
45     flow_edge_list (int n = MAXN) { clear (n); }
46     void add_edge (int u, int v, int f) {
47       dest[size] = v; next[size] = begin[u]; flow[size] =
48       f; begin[u] = size++;
49       dest[size] = u; next[size] = begin[v]; flow[size] =
50       0; begin[v] = size++; }
51     int n, s, t, d[MAXN], w[MAXN], q[MAXN];
52     int bfs (flow_edge_list &e) {
53       std::fill (d, d + n, -1);
54       int l, r; q[l = r = 0] = s, d[s] = 0;
55       for (; l <= r; l++)
56         for (int k = e.begin[q[l]]; ~k; k = e.next[k])
57           if (!d[e.dest[k]] && e.flow[k] > 0) d[e.dest[k]]
58             = d[q[l]] + 1, q[++r] = e.dest[k];
59       return ~d[t] ? 1 : 0;
60     }
61     int dfs (flow_edge_list &e, int u, int ext) {
62       if (u == t) return ext; int k = w[u], ret = 0;
63       for (; ~k; k = e.next[k], w[u] = k) {
64         if (ext == 0) break;
65         if (d[e.dest[k]] == d[u] + 1 && e.flow[k] > 0) {
66           int flow = dfs (e, e.dest[k], std::min (e.flow[k],
67             ext));
68           if (flow > 0) {
69             e.flow[k] -= flow, e.flow[k ^ 1] += flow;
70             ret += flow, ext -= flow; }
71         }
72         if (!k) d[u] = -1; return ret; }
73     int solve (flow_edge_list &e, int n_, int s_, int t_) {
74       int ans = 0; n = n_; s = s_; dinic::t = t_;
75       while (bfs (e)) {
76         for (int i = 0; i < n; ++i) w[i] = e.begin[i];
77         ans += dfs (e, s, INF); }
78       return ans; } };

```

6.6 Minimum cost flow

```

1 /* Sparse graph minimum cost flow : EK.*/
2 template<int MAXN = 1000, int MAXM = 100000>
3 struct minimum_cost_flow {
4   struct cost_flow_edge_list {
5     int size, begin[MAXN], dest[MAXM], next[MAXM], cost[
6       MAXM], flow[MAXM];
7     void clear (int n) { size = 0; std::fill (begin,
8       begin + n, -1); }
9     cost_flow_edge_list (int n = MAXN) { clear (n); }
10    void add_edge (int u, int v, int c, int f) {
11      dest[size] = v; next[size] = begin[u]; cost[size] =
12      c; flow[size] = f; begin[u] = size++;
13      dest[size] = u; next[size] = begin[v]; cost[size] =
14      -c; flow[size] = 0; begin[v] = size++; }
15    int n, s, t, prev[MAXN], dist[MAXN], occur[MAXN];
16    bool augment (cost_flow_edge_list &e) {
17      std::vector<int> queue;
18      std::fill (dist, dist + n, INF); std::fill (occur,
19        occur + n, 0);
20      dist[s] = 0; occur[s] = true; queue.push_back (s);
21      for (int head = 0; head < (int)queue.size(); ++head)
22        {
23          int x = queue[head];
24          for (int i = e.begin[x]; ~i; i = e.next[i]) {
25            int y = e.dest[i];
26            if (e.flow[i] && dist[y] > dist[x] + e.cost[i]) {
27              dist[y] = dist[x] + e.cost[i]; prev[y] = i;
28              if (!occur[y]) {
29                occur[y] = true; queue.push_back (y); }
30            }
31            occur[x] = false;
32          }
33          return dist[t] < INF; }
34      std::pair<int, int> solve (cost_flow_edge_list &e,
35        int n_, int s_, int t_) {
36        n = n_; s = s_; t = t_; std::pair<int, int> ans =
37        std::make_pair (0, 0);
38        while (augment (e)) {
39          int num = INF;
40          for (int i = t; i != s; i = e.dest[prev[i] ^ 1]) {
41            num = std::min (num, e.flow[prev[i]]); }
42          ans.first += num;
43          for (int i = t; i != s; i = e.dest[prev[i] ^ 1]) {
44            e.flow[prev[i]] -= num; e.flow[prev[i] ^ 1] += num;
45            ans.second += num * e.cost[prev[i]]; }
46          return ans; } };
47    /* Dense graph minimum cost flow : zkw. */
48    template<int MAXN = 1000, int MAXM = 100000>
49    struct zkw_flow {
50      struct cost_flow_edge_list {
51        int size, begin[MAXN], dest[MAXM], next[MAXM], cost[
52          MAXM], flow[MAXM];
53        void clear (int n) { size = 0; std::fill (begin,
54          begin + n, -1); }
55        cost_flow_edge_list (int n = MAXN) { clear (n); }
56        void add_edge (int u, int v, int c, int f) {

```



```

45 dest[size] = v; next[size] = begin[u]; cost[size] =
   c; flow[size] = f; begin[u] = size++;
46 dest[size] = u; next[size] = begin[v]; cost[size] =
   -c; flow[size] = 0; begin[v] = size++; } }
47 int n, s, t, tf, tc, dis[MAXN], slack[MAXN], visit[
   MAXN];
48 int modlable() {
49     int delta = INF;
50     for (int i = 0; i < n; i++) {
51         if (!visit[i] && slack[i] < delta) delta = slack[i];
52     }
53     if (delta == INF) return 1;
54     for (int i = 0; i < n; i++) if (visit[i]) dis[i] +=
       delta;
55     return 0;
56 int dfs (cost_flow_edge_list &e, int x, int flow) {
57     if (x == t) { tf += flow; tc += flow * (dis[s] - dis
       [t]); return flow; }
58     visit[x] = 1; int left = flow;
59     for (int i = e.begin[x]; ~i; i = e.next[i])
60         if (e.flow[i] > 0 && !visit[e.dest[i]]) {
61             int y = e.dest[i];
62             if (dis[y] + e.cost[i] == dis[x]) {
63                 int delta = dfs (e, y, std::min (left, e.flow[i])
                   );
64                 e.flow[i] -= delta; e.flow[i ^ 1] += delta; left
                   -= delta;
65                 if (!left) { visit[x] = false; return flow; }
66             } else
67                 slack[y] = std::min (slack[y], dis[x] + e.cost[i]
                   - dis[x]);
68             return flow - left;
69     std::pair <int, int> solve (cost_flow_edge_list &e,
       int n, int s, int t) {
70         n = n; s = s; t = t; tf = tc = 0;
71         std::fill (dis + 1, dis + t + 1, 0);
72         do { do {
73             std::fill (visit + 1, visit + t + 1, 0);
74         } while (dfs (e, s, INF)); } while (!modlable ());
75         return std::make_pair (tf, tc);
76     } };
```

6.7 Stoer Wagner algorithm

```

1 /* Stoer Wagner algorithm : Finds the minimum cut of
   an undirected graph. (1-based) */
2 template <int MAXN = 500>
3 struct stoer_wagner {
4     int n, edge[MAXN][MAXN];
5     int dist[MAXN];
6     bool vis[MAXN], bin[MAXN];
7     stoer_wagner () {
8         memset (edge, 0, sizeof (edge));
9         memset (bin, false, sizeof (bin)); }
10    int contract (int &s, int &t) {
11        memset (dist, 0, sizeof (dist));
12        memset (vis, false, sizeof (vis));
13        int i, j, k, mincut, maxc;
14        for (i = 1; i <= n; i++) {
15            k = -1; maxc = -1;
16            for (j = 1; j <= n; j++)
17                if (!bin[j] && !vis[j] && dist[j] > maxc) {
18                    k = j; maxc = dist[j]; }
19            if (k == -1) return mincut;
20            s = t; t = k; mincut = maxc; vis[k] = true;
21            for (j = 1; j <= n; j++) if (!bin[j] && !vis[j])
22                dist[j] += edge[k][j];
23            return mincut; }
24    int solve () {
25        int mincut, i, j, s, t, ans;
26        for (mincut = INF, i = 1; i < n; i++) {
27            ans = contract (s, t); bin[t] = true;
28            if (mincut > ans) mincut = ans;
29            if (mincut == 0) return 0;
30            for (j = 1; j <= n; j++) if (!bin[j])
31                edge[s][j] = (edge[j][s] += edge[j][t]); }
32        return mincut; } };
```

6.8 DN maximum clique

```

1 /* DN maximum clique : n <= 150 */
2 typedef bool BB[N]; struct Maxclique {
3     const BB *e; int pk, level; const float Tlimit;
4     struct Vertex { int i, d; Vertex (int i) : i(i), d(0)
       {} };
5     typedef std::vector <Vertex> Vertices; Vertices V;
6     typedef std::vector <int> ColorClass; ColorClass QMAX,
       Q;
7     std::vector <ColorClass> C;
8     static bool desc_degree (const Vertex &vi, const Vertex
       &vj) { return vi.d > vj.d; }
9     void init_colors (Vertices &v) {
10         const int max_degree = v[0].d;
11         for (int i = 0; i < (int) v.size(); ++i) v[i].d = std
           ::min (i, max_degree + 1); }
12     void set_degrees (Vertices &v) {
13         for (int i = 0, j; i < (int) v.size(); ++i)
14             for (v[i].d = j = 0; j < (int) v.size(); ++j)
15                 v[i].d += e[v[i].i][v[j].i]; }
16     struct StepCount { int i1, i2; StepCount () : i1 (0), i2
       (0) {} };
17     std::vector <StepCount> S;
18     bool cut1 (const int pi, const ColorClass &A) {
19         for (int i = 0; i < (int) A.size(); ++i)
20             if (e[pi][A[i]]) return true; return false; }
21     void cut2 (const Vertices &A, Vertices &B) {
22         for (int i = 0; i < (int) A.size() - 1; ++i)
23             if (e[A.back().i][A[i].i]) B.push_back (A[i].i); }
24     void color_sort (Vertices &R) {
25         int j = 0, maxno = 1, min_k = std::max ((int) QMAX.
           size () - (int) Q.size () + 1, 1);
```

```

26 C[1].clear (); C[2].clear ();
27 for (int i = 0; i < (int) R.size (); ++i) {
28     int pi = R[i].i, k = 1; while (cut1(pi, C[k])) ++k;
29     if (k > maxno) maxno = k, C[maxno + 1].clear();
30     C[k].push_back (pi); if (k < min_k) R[j++].i = pi; }
31 if (j > 0) R[j - 1].d = 0;
32 for (int k = min_k; k <= maxno; ++k)
33     for (int i = 0; i < (int) C[k].size (); ++i)
34         R[j].i = C[k][i], R[j++].d = k; }
35 void expand_dyn (Vertices &R) {
36     S[level].i1 = S[level].i1 + S[level - 1].i1 - S[level
       ].i2;
37     S[level].i2 = S[level - 1].i1;
38     while ((int) R.size ()) {
39         if ((int) Q.size () + R.back ().d > (int) QMAX.size
           ()) {
40             Q.push_back (R.back ().i); Vertices Rp; cut2 (R, Rp
           );
41             if ((int) Rp.size ()) {
42                 if ((float) S[level].i1 / ++pk < Tlimit)
43                     degree_sort (Rp);
44                 color_sort (Rp); ++S[level].i1, ++level;
45                 expand_dyn (Rp); --level;
46             } else if ((int) Q.size () > (int) QMAX.size ())
47                 QMAX = Q;
48             Q.pop_back (); } else return; R.pop_back (); } }
49 void mcqdyn (int *maxclique, int &sz) {
50     set_degrees (V); std::sort (V.begin (), V.end (),
51         desc_degree); init_colors (V);
52     for (int i = 0; i < (int) V.size () + 1; ++i) S[i].i1
       = S[i].i2 = 0;
53     expand_dyn (V); sz = (int) QMAX.size ();
54     for (int i = 0; i < (int) QMAX.size (); ++i)
55         maxclique[i] = QMAX[i]; }
56 void degree_sort (Vertices &R) {
57     set_degrees (R); std::sort (R.begin (), R.end (),
58         desc_degree); }
59 Maxclique (const BB *conn, const int sz, const float
       tt = .025) : pk (0), level (1), Tlimit (tt) {
60     for (int i = 0; i < sz; ++i) V.push_back (Vertex (i));
61     e = conn, C.resize (sz + 1), S.resize (sz + 1); }
62 BB e[N]; int ans, sol[N]; for (...) e[x][y] = e[y][x]
       = true;
63 Maxclique mc (e, n); mc.mcqdyn (sol, ans); //0-based.
64 for (int i = 0; i < ans; ++i) std::cout << sol[i] <<
       std::endl;
```

6.9 Dominator tree

```

1 /* Dominator tree : finds the immediate dominator (
   idom[]) of each node, idom[x] will be x if x does
   not have a dominator, and will be -1 if x is not
   reachable from s. */
2 template <int MAXN = 100000, int MAXM = 100000>
3 struct dominator_tree {
4     using edge_list = std::vector <int> [MAXN];
5     int dfn[MAXN], sdom[MAXN], idom[MAXN], fa[MAXN], f[
       MAXN], fa[MAXN], smin[MAXN], stamp;
6     void predfs (int x, const edge_list <MAXN, MAXM> &
       succ) {
7         idf[n] = stamp++; x;
8         for (int i = succ.begin[x]; ~i; i = succ.next[i]) {
9             int y = succ.dest[i];
10            if (dfn[y] < 0) { f[y] = x; predfs (y, succ); } } }
11    int getfa (int x) {
12        if (fa[x] == x) return x;
13        int ret = getfa (fa[x]);
14        if (dfn[sdom[smin[fa[x]]]] < dfn[sdom[smin[x]]])
15            smin[x] = smin[fa[x]];
16        return fa[x] = ret; }
17    void solve (int s, int n, const edge_list <MAXN, MAXM>
       & succ) {
18        std::fill (dfn, dfn + n, -1); std::fill (idom, idom
          + n, -1);
19        static edge_list <MAXN, MAXM> pred, tmp; pred.clear
          (n);
20        for (int i = 0; i < n; ++i) for (int j = succ.begin[
          i]; ~j; j = succ.next[j])
21            pred.add_edge (succ.dest[j], i);
22        stamp = 0; tmp.clear (n); predfs (s, succ);
23        for (int i = 0; i < stamp; ++i) fa[id[i]] = smin[id[
          i]];
24        for (int o = stamp - 1; o >= 0; --o) {
25            int x = id[o];
26            if (o) {
27                sdom[x] = f[x];
28                for (int i = pred.begin[x]; ~i; i = pred.next[i])
29                    {
30                        int p = pred.dest[i];
31                        if (dfn[p] < 0) continue;
32                        if (dfn[p] > dfn[x]) { getfa (p); p = sdom[smin[p]
                          ]; }
33                        if (dfn[sdom[x]] > dfn[p]) sdom[x] = p; }
34                tmp.add_edge (sdom[x], x); }
35            while (~tmp.begin[x]) {
36                int y = tmp.dest[tmp.begin[x]];
37                tmp.begin[x] = tmp.next[tmp.begin[x]]; getfa (y);
38                if (x != sdom[smin[y]]) idom[y] = smin[y];
39                else idom[y] = x; }
40            for (int v : succ[x]) if (f[v] == x) fa[v] = x; }
41        idom[s] = s; for (int i = 1; i < stamp; ++i) {
42            int x = id[i]; if (idom[x] != sdom[x]) idom[x] =
              idom[idom[x]]; } }
```

6.10 Tarjan

```

1 /* Tarjan : strongly-connected components. */
2 template <int MAXN = 1000000>
3 struct tarjan {
4     int comp[MAXN], size;
5     int dfn[MAXN], ind, low[MAXN], ins[MAXN], stk[MAXN],
       stks;
```



```

6 void dfs (const edge_list <MAXN, MAXM> &e, int i) {
7   dfn[i] = low[i] = ind++;
8   ins[i] = 1; stk[stks++] = i;
9   for (int x = e.begin[i]; ~x; x = e.next[x]) {
10    int j = e.dest[x]; if (!dfn[j]) {
11     dfs (j);
12     if (low[i] > low[j]) low[i] = low[j];
13     if (low[j] >= dfn[i]); //vertex-biconnected
14     if (low[j] > dfn[i]); //edge-biconnected
15    } else if (ins[j] && low[i] > dfn[j])
16     low[i] = dfn[j];
17    if (dfn[i] == low[i]) { //strongly-connected
18     for (int j = -1; j != i;
19          j = stk[--stks], ins[j] = 0, comp[j] = size);
20     ++size; }
21 void solve (const edge_list <MAXN, MAXM> &e, int n) {
22   size = ind = stks = 0;
23   std::fill (dfn, dfn + n, -1);
24   for (int i = 0; i < n; ++i) if (!dfn[i])
25     dfs (e, i); }

```

7 String

7.1 Manacher

```

1 /* Manacher : Odd parliindromes only. */
2 for (int i = 1, j = 0; i != (n << 1) - 1; ++i) {
3   int p = i >> 1, q = i - p, r = ((j + 1) >> 1) + 1[j]
4     - 1;
5   l[i] = r < q ? 0 : std::min (r - q + 1, l[(j << 1) -
6     i]);
7   while (p - l[i] != -1 && q + l[i] != n
8     && s[p - l[i]] == s[q + l[i]]) l[i]++;
9   if (q + l[i] - 1 > r) j = i;
10  a += l[i]; }

```

7.2 Suffix Array

```

1 /* Suffix Array : sa[i] - the beginning position of
2   the ith smallest suffix, rk[i] - the rank of the
3   suffix beginning at position i. height[i] - the
4   longest common prefix of sa[i] and sa[i - 1]. */
5 template <int MAXN = 1000000, int MAXC = 26>
6 struct suffix_array {
7   int rk[MAXN], height[MAXN], sa[MAXN];
8   int cmp (int *x, int a, int b, int d) {
9     return x[a] == x[b] && x[a + d] == x[b + d]; }
10  void doubling (int *a, int n) {
11    static int sRank[MAXN], tmpA[MAXN], tmpB[MAXN];
12    int m = MAXC, *x = tmpA, *y = tmpB;
13    for (int i = 0; i < m; ++i) sRank[i] = 0;
14    for (int i = 0; i < n; ++i) ++sRank[x[i] = a[i]];
15    for (int i = 1; i < m; ++i) sRank[i] += sRank[i -
16      1];
17    for (int i = n - 1; i >= 0; --i) sa[--sRank[x[i]]] =
18      i;
19    for (int d = 1, p = 0; p < n; m = p, d <= 1) {
20      p = 0; for (int i = n - d; i < n; ++i) y[p++] = i;
21      for (int i = 0; i < n; ++i) if (sa[i] >= d) y[p++]
22        = sa[i] - d;
23      for (int i = 0; i < m; ++i) sRank[i] = 0;
24      for (int i = 0; i < n; ++i) ++sRank[x[i]];
25      for (int i = 1; i < m; ++i) sRank[i] += sRank[i -
26        1];
27      for (int i = n - 1; i >= 0; --i) sa[--sRank[x[y[i]
28        ]]] = y[i];
29      std::swap (x, y); x[sa[0]] = 0; p = 1; y[n] = -1;
30      for (int i = 1; i < n; ++i)
31        x[sa[i]] = cmp (y, sa[i], sa[i - 1], d) ? p - 1 :
32          p++; }
33 void solve (int *a, int n) {
34   a[n] = -1; doubling (a, n);
35   for (int i = 0; i < n; ++i) rk[sa[i]] = i;
36   int cur = 0;
37   for (int i = 0; i < n; ++i)
38     if (rk[i]) {
39       if (cur) cur--;
40       for (; a[i + cur] == a[sa[rk[i] - 1] + cur]; ++cur)
41         ;
42       height[rk[i]] = cur; } }

```

7.3 Suffix Automaton

```

1 /* Suffix automaton : head - the first state. tail -
2   the last state. Terminating states can be reached
3   via visiting the ancestors of tail. state::len -
4   the longest length of the string in the state.
5   state::right - 1 - the first location in the
6   string where the state can be reached. state::
7   parent - the parent link. state::dest - the
8   automaton link. */
9 template <int MAXN = 1000000, int MAXC = 26>
10 struct suffix_automaton {
11   struct state {
12     int len, right; state *parent, *dest[MAXC];
13     state (int len = 0, int right = 0) : len (len),
14       right (right), parent (NULL) {
15       memset (dest, 0, sizeof (dest)); }
16   } node_pool[MAXN * 2], *tot_node, *null = new state();
17   state *head, *tail;
18 void extend (int token) {
19   state *p = tail;
20   state *np = tail -> dest[token] ? null : new (
21     tot_node++) state (tail -> len + 1, tail -> len
22     + 1);
23   while (p && !p -> dest[token]) p -> dest[token] = np;
24   if (!p) np -> parent = head;
25   else {

```

```

16   state *q = p -> dest[token];
17   if (p -> len + 1 == q -> len) {
18     np -> parent = q;
19   } else {
20     state *nq = new (tot_node++) state (*q);
21     nq -> len = p -> len + 1;
22     np -> parent = q -> parent = nq;
23     while (p && p -> dest[token] == q) {
24       p -> dest[token] = nq, p = p -> parent;
25     }
26   }
27   tail = np == null ? np -> parent : np; }
28 void init () {
29   tot_node = node_pool;
30   head = tail = new (tot_node++) state (); }
31 suffix_automaton () { init (); }

```

7.4 Palindromic tree

```

1 /* Palindromic tree : extend () - returns whether the
2   tree has generated a new node. odd, even - the
3   root of two trees. last - the node representing
4   the last char. node::len - the palindromic string
5   length of the node. */
6 template <int MAXN = 1000000, int MAXC = 26>
7 struct palindromic_tree {
8   struct node {
9     node *child[MAXC], *fail; int len;
10    node (int len) : fail (NULL), len (len) {
11      memset (child, NULL, sizeof (child)); }
12    } node_pool[MAXN * 2], *tot_node;
13    int size, text[MAXN];
14    node *odd, *even, *last;
15    node *match (node *now) {
16      for (; text[size] - now -> len - 1 != text[size];
17           now = now -> fail; )
18        ;
19      return now; }
20 bool extend (int token) {
21   text[size] = token; node *now = match (last);
22   if (now -> child[token])
23     return last = now -> child[token], false;
24   last = now -> child[token] = new (tot_node++) node (
25     now -> len + 2);
26   if (now == odd) last -> fail = even;
27   else {
28     now = match (now -> fail);
29     last -> fail = now -> child[token]; }
30   return true; }
31 void init () {
32   text[size] = 0; tot_node = node_pool;
33   last = even = new (tot_node++) node (0); odd = new (
34     tot_node++) node (-1);
35   even -> fail = odd; }
36 palindromic_tree () { init (); }

```

7.5 Regular expression

```

1 std::string str = ("The_the_there");
2 std::regex pattern ("(th|Th)[\\w]*", std::
3   regex_constants::optimize | std::regex_constants::
4   ECMAScript);
5 std::smatch match; //std::cmatch for char *
6
7 std::regex_match (str, match, pattern);
8
9 auto mbegin = std::sregex_iterator (str.begin (), str.
10   end (), pattern);
11 auto mend = std::sregex_iterator ();
12 std::cout << "Found_" << std::distance (mbegin, mend)
13   << "\nwords:\n";
14 for (std::sregex_iterator i = mbegin; i != mend; ++i)
15   {
16     match = *i;
17     /* The word is match[0], backreferences are match[i]
18       up to match.size ().
19     match.prefix () and match.suffix () give the prefix
20       and the suffix.
21     match.length () gives length and match.position ()
22       gives position of the match. */
23     std::regex_replace (str, pattern, "sh$1");
24     // $n is the backreference, $& is the entire match, $'
25     is the prefix, $' is the suffix, $$ is the $ sign.

```

8 Tips

8.1 Java

```

1 /* Java reference : References on Java IO, structures,
2   etc. */
3 import java.io.*;
4 import java.lang.*;
5 import java.math.*;
6 import java.util.*;
7 /* Common usage:
8 Scanner in = new Scanner (System.in);
9 Scanner in = new Scanner (new BufferedInputStream (
10   System.in));
11 in.nextInt () / in.nextBigInteger () / in.
12   nextBigDecimal () / in.nextDouble ()
13 in.nextLine () / in.hasNext ()
14 System.out.print (...);
15 System.out.println (...);
16 System.out.printf (...);
17 BigInteger : BigInteger.valueOf (int) / abs / negate
18   () / max / min / add / subtract / multiply /
19   divide / remainder (BigInteger) / gcd (BigInteger)
20   / modInverse (BigInteger mod) / modPow (
21   BigInteger ex, BigInteger mod) / pow (int ex) /
22   not () / and / or / xor (BigInteger) / shiftLeft /
23   shiftRight (int) / compareTo (BigInteger) /
24   intValue () / longValue () / toString (int radix)
25   / isProbablePrime (int certainty) /
26   nextProbablePrime ()

```

```

15 BigDecimal : consists of a BigInteger value and a
    scale. The scale is the number of digits to the
    right of the decimal point.
16 divide (BigDecimal) : exact divide.
17 divide (BigDecimal, int scale, RoundingMode
    roundingMode) : divide with roundingMode, which
    may be: CEILING / DOWN / FLOOR / HALF_DOWN /
    HALF_EVEN / HALF_UP / UNNECESSARY / UP.
18 BigDecimal setScale (int newScale, RoundingMode
    roundingMode) : returns a BigDecimal with newScale
19 doubleValue () / toString () : converts to other
    types.
20 Arrays : Arrays.sort (T [] a); Arrays.sort (T [] a,
    int fromIndex, int toIndex); Arrays.sort (T [] a,
    int fromIndex, int toIndex, Comparator <? super T>
    comparator);
21 LinkedList <E> : addFirst / addLast (E) / getFirst /
    getLast / removeFirst / removeLast () / clear () /
    add (int, E) / remove (int) / size () / contains
    / removeFirstOccurrence / removeLastOccurrence (E)
22 ListIterator <E> listIterator (int index) : returns an
    iterator :
23 E next / previous () : accesses and iterates.
24 hasNext / hasPrevious () : checks availability.
25 nextIndex / previousIndex () : returns the index of a
    subsequent call.
26 add / set (E) / remove () : changes element.
27 PriorityQueue <E> (int initCap, Comparator <? super E>
    comparator) : add (E) / clear () / iterator () /
    peek () / poll () / size ()
28 TreeMap <K, V> (Comparator <? super K> comparator) :
    Map.Entry <K, V> ceilingEntry / floorEntry /
    higherEntry / lowerEntry (K): getKey / getValue ()
    / setValue (V) : entries.
29 clear () / put (K, V) / get (K) / remove (K) / size
    ()
30 StringBuilder : StringBuilder (string) / append (int,
    string, ...) / insert (int offset, ...) charAt (
    int) / setCharAt (int, char) / delete (int, int) /
    reverse () / length () / toString ()
31 String : String.format (String, ...) / toLowerCase /
    toUpperCase () */
32 /* Examples on Comparator :
33 public class Main {
34     public static class Point {
35         public int x; public int y;
36         public Point () {
37             x = 0;
38             y = 0;
39         }
40         public Point (int xx, int yy) {
41             x = xx;
42             y = yy;
43         };
44     }
45     public static class Cmp implements Comparator <Point>
46     {
47         public int compare (Point a, Point b) {
48             if (a.x < b.x) return -1;
49             if (a.x == b.x) {
50                 if (a.y < b.y) return -1;
51                 if (a.y == b.y) return 0;
52             }
53             return 1;
54         };
55     }
56     public static void main (String [] args) {
57         Cmp c = new Cmp ();
58         TreeMap <Point, Point> t = new TreeMap <Point, Point>
59         (c);
60         return;
61     };
62 }
63 */
64 /* or :
65 public static class Point implements Comparable <
66     Point> {
67     public int x; public int y;
68     public Point () {
69         x = 0;
70         y = 0;
71     }
72     public Point (int xx, int yy) {
73         x = xx;
74         y = yy;
75     }
76     public int compareTo (Point p) {
77         if (x < p.x) return -1;
78         if (x == p.x) {
79             if (y < p.y) return -1;
80             if (y == p.y) return 0;
81         }
82         return 1;
83     }
84     public boolean equalTo (Point p) {
85         return (x == p.x && y == p.y);
86     }
87     public int hashCode () {
88         return x + y;
89     };
90 }
91 */
92 //Faster IO :
93 public class Main {
94     static class InputReader {
95         public BufferedReader reader;
96         public StringTokenizer tokenizer;
97         public InputReader (InputStream stream) {
98             reader = new BufferedReader (new InputStreamReader
99             (stream), 32768);
100             tokenizer = null;
101         }
102         public String next () {
103             while (tokenizer == null || !tokenizer.
104                 hasMoreTokens ()) {
105                 try {
106                     String line = reader.readLine();
107                     tokenizer = new StringTokenizer (line);
108                 } catch (IOException e) {
109                     throw new RuntimeException (e);
110                 }
111             }
112             return tokenizer.nextToken();
113         }
114         public BigInteger nextBigInteger () {
115             return new BigInteger (next (), 10); /* radix */
116         }
117         public int nextInt () {
118             return Integer.parseInt (next ());
119         }
120         public double nextDouble () {
121             return Double.parseDouble (next ());
122         }
123     }
124     public static void main (String [] args) {
125         InputReader in = new InputReader (System.in);
126     }
127 }

```

8.2 Random numbers

```

1 std::mt19937_64 mt (time (0));
2 std::uniform_int_distribution <int> uid (1, 100);
3 std::uniform_real_distribution <double> urd (1, 100);
4 std::cout << uid (mt) << " " << urd (mt) << "\n";

```

8.3 Read hack

```

1 #define __attribute__ ((optimize ("-O3")))
2 #define __inline__ __attribute__ ((__gnu_inline__,
3     __always_inline__, __artificial__))
4 int next_int () {
5     const int SIZE = 110000; static char buf[SIZE + 1];
6     static int p = SIZE;
7     register int ans = 0, f = 1, sgn = 1;
8     while ((p < SIZE || (p = 0, buf[fread (buf, 1, SIZE,
9         stdin)] = 0, buf[0])) && (isdigit (buf[p]) && (
10         ans = ans * 10 + buf[p] - '0', f = 0, 1) || f &&
11         (buf[p] == '-' && (sgn = 0, 1))) ++p;
12     return sgn ? ans : -ans;
13 }

```

8.4 Stack hack

```

1 //C++
2 #pragma comment (linker, "/STACK:36777216")
3 //G++
4 int _size_ = 256 << 20;
5 char * _p_ = (char*) malloc (_size_ + _size_);
6 __asm__ ("movl %0, %%esp\n" :: "r" (_p_));

```

8.5 Time hack

```

1 clock_t t = clock ();
2 std::cout << 1. * t / CLOCKS_PER_SEC << "\n";

```

8.6 Builtin functions

1. `__builtin_clz`: Returns the number of leading 0-bits in `x`, starting at the most significant bit position. If `x` is 0, the result is undefined.
2. `__builtin_ctz`: Returns the number of trailing 0-bits in `x`, starting at the least significant bit position. If `x` is 0, the result is undefined.
3. `__builtin_clrsb`: Returns the number of leading redundant sign bits in `x`, i.e. the number of bits following the most significant bit that are identical to it. There are no special cases for 0 or other values.
4. `__builtin_popcount`: Returns the number of 1-bits in `x`.
5. `__builtin_parity`: Returns the parity of `x`, i.e. the number of 1-bits in `x` modulo 2.
6. `__builtin_bswap16`, `__builtin_bswap32`, `__builtin_bswap64`: Returns `x` with the order of the bytes (8 bits as a group) reversed.
7. `bitset::Find_first()`, `bitset::Find_next(idx)`: `bitset` builtin functions.

8.7 Prufer sequence

In combinatorial mathematics, the Prufer sequence of a labeled tree is a unique sequence associated with the tree. The sequence for a tree on n vertices has length $n - 2$.

One can generate a labeled tree's Prufer sequence by iteratively removing vertices from the tree until only two vertices remain. Specifically, consider a labeled tree T with vertices $1, 2, \dots, n$. At step i , remove the leaf with the smallest label and set the i th element of the Prufer sequence to be the label of this leaf's neighbour.

One can generate a labeled tree from a sequence in three steps. The tree will have $n + 2$ nodes, numbered from 1 to $n + 2$. For each node set its degree to the number of times it appears in the sequence plus 1. Next, for each number in the sequence $a[i]$, find the first (lowest-numbered) node, j , with degree equal to 1, add the edge $(j, a[i])$ to the tree, and decrement the degrees of j and $a[i]$. At the end of this loop two nodes with degree 1 will remain (call them u, v). Lastly, add the edge (u, v) to the tree.

The Prufer sequence of a labeled tree on n vertices is a unique sequence of length $n - 2$ on the labels 1 to n - this much is clear. Somewhat less obvious is the fact that for a given sequence S of length $n - 2$ on the labels 1 to n , there is a unique labeled tree whose Prufer sequence is S .

8.8 Spanning tree counting

Kirchhoff's Theorem: the number of spanning trees in a graph G is equal to *any* cofactor of the Laplacian matrix of G , which is equal to the difference between the graph's degree matrix (a diagonal matrix with vertex degrees on the diagonals) and its adjacency matrix (a $(0,1)$ -matrix with 1's at places corresponding to entries where the vertices are adjacent and 0's otherwise).

The number of edges with a certain weight in a minimum spanning tree is fixed given a graph. Moreover, the number of its arrangements can be obtained by finding a minimum spanning tree, compressing connected components of other edges in that tree into a point, and then applying Kirchhoff's theorem with only edges of the certain weight in the graph. Therefore, the number of minimum spanning trees in a graph can be solved by multiplying all numbers of arrangements of edges of different weights together.

8.9 Mobius inversion

8.9.1 Mobius inversion formula

$$[x = 1] = \sum_{d|x} \mu(d)$$

$$x = \sum_{d|x} \mu(d)$$

8.9.2 Gcd inversion

$$\begin{aligned}
\sum_{a=1}^n \sum_{b=1}^n \gcd^2(a, b) &= \sum_{d=1}^n d^2 \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} [\gcd(i, j) = 1] \\
&= \sum_{d=1}^n d^2 \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \\
&= \sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \sum_{i=1}^{\lfloor \frac{n}{dt} \rfloor} [t|i] \sum_{j=1}^{\lfloor \frac{n}{dt} \rfloor} [t|j] \\
&= \sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \left\lfloor \frac{n}{dt} \right\rfloor^2
\end{aligned}$$

The formula can be computed in $O(n \log n)$ complexity. Moreover, let $l = dt$, then

$$\sum_{d=1}^n d^2 \sum_{t=1}^{\lfloor \frac{n}{d} \rfloor} \mu(t) \left\lfloor \frac{n}{dt} \right\rfloor^2 = \sum_{l=1}^n \left\lfloor \frac{n}{l} \right\rfloor^2 \sum_{d|l} d^2 \mu\left(\frac{l}{d}\right)$$

Let $f(l) = \sum_{d|l} d^2 \mu\left(\frac{l}{d}\right)$. It can be proven that $f(l)$ is multiplicative. Besides, $f(p^k) = p^{2k} - p^{2k-2}$.

Therefore, with linear sieve the formula can be computed in $O(n)$ complexity.

8.10 2-SAT

In terms of the implication graph, two literals belong to the same strongly connected component whenever there exist chains of implications from one literal to the other and vice versa. Therefore, the two literals must have the same value in any satisfying assignment to the given 2-satisfiability instance. In particular, if a variable and its negation both belong to the same strongly connected component, the instance cannot be satisfied, because it is impossible to assign both of these literals the same value. As Aspvall et al. showed, this is a necessary and sufficient condition: a 2-CNF formula is satisfiable if and only if there is no variable that belongs to the same strongly connected component as its negation.

This immediately leads to a linear time algorithm for testing satisfiability of 2-CNF formulae: simply perform a strong connectivity analysis on the implication graph and check that each variable and its negation belong to different components. However, as Aspvall et al. also showed, it also leads to a linear time algorithm for finding a satisfying assignment, when one exists. Their algorithm performs the following steps:

Construct the implication graph of the instance, and find its strongly connected components using any of the known linear-time algorithms for strong connectivity analysis.

Check whether any strongly connected component contains both a variable and its negation. If so, report that the instance is not satisfiable and halt.

Construct the condensation of the implication graph, a smaller graph that has one vertex for each strongly connected component, and an edge from component i to component j whenever the implication graph contains an edge uv such that u belongs to component i and v belongs to component j . The condensation is automatically a directed acyclic graph and, like the implication graph from which it was formed, it is skew-symmetric.

Topologically order the vertices of the condensation. In practice this may be efficiently achieved as a side effect of the previous step, as components are generated by Kosaraju's algorithm in topological order and by Tarjan's algorithm in reverse topological order.

For each component in the reverse topological order, if its variables do not already have truth assignments, set all the literals in the component to be true. This also causes all of the literals in the complementary component to be set to false.

8.11 Interesting numbers

8.11.1 Binomial Coefficients

$$\begin{aligned}
\binom{n}{k} &= (-1)^k \binom{k-n-1}{k}, \quad \sum_{k \leq n} \binom{r+k}{k} = \binom{r+n+1}{n} \\
\sum_{k=0}^n \binom{k}{m} &= \binom{n+1}{m+1} \\
\sqrt{1+z} &= 1 + \sum_{k=1}^{\infty} \frac{(-1)^{k-1}}{k \times 2^{2k-1}} \binom{2k-2}{k-1} z^k \\
\sum_{k=0}^r \binom{r-k}{m} \binom{s+k}{n} &= \binom{r+s+1}{m+n+1} \\
C_{n,m} &= \binom{n+m}{m} - \binom{n+m}{m-1}, n \geq m \\
\binom{n}{k} &\equiv [n \& k = k] \pmod{2} \\
\binom{n_1 + \dots + n_p}{m} &= \sum_{k_1 + \dots + k_p = m} \binom{n_1}{k_1} \dots \binom{n_p}{k_p}
\end{aligned}$$

8.11.2 Fibonacci Numbers

$$\begin{aligned}
F(z) &= \frac{z}{1-z-z^2} \\
f_n &= \frac{\phi^n - \hat{\phi}^n}{\sqrt{5}}, \phi = \frac{1+\sqrt{5}}{2}, \hat{\phi} = \frac{1-\sqrt{5}}{2} \\
\sum_{k=1}^n f_k &= f_{n+2} - 1, \quad \sum_{k=1}^n f_k^2 = f_n f_{n+1} \\
\sum_{k=0}^n f_k f_{n-k} &= \frac{1}{5}(n-1)f_n + \frac{2}{5}n f_{n-1} \\
\frac{f_{2n}}{f_n} &= f_{n-1} + f_{n+1} \\
f_1 + 2f_2 + 3f_3 + \dots + n f_n &= n f_{n+2} - f_{n+3} + 2 \\
\gcd(f_m, f_n) &= f_{\gcd(m, n)} \\
f_n^2 + (-1)^n &= f_{n+1} f_{n-1} \\
f_{n+k} &= f_n f_{k+1} + f_{n-1} f_k \\
f_{2n+1} &= f_n^2 + f_{n+1}^2 \\
(-1)^k f_{n-k} &= f_n f_{k-1} - f_{n-1} f_k \\
\text{Modulo } f_n, f_{mn+r} &\equiv \begin{cases} f_r, & m \bmod 4 = 0; \\ (-1)^{r+1} f_{n-r}, & m \bmod 4 = 1; \\ (-1)^n f_r, & m \bmod 4 = 2; \\ (-1)^{r+1+n} f_{n-r}, & m \bmod 4 = 3. \end{cases}
\end{aligned}$$

8.11.3 Lucas Numbers

$$\begin{aligned}
L_0 &= 2, L_1 = 1, L_n = L_{n-1} + L_{n-2} = \left(\frac{1+\sqrt{5}}{2}\right)^n + \left(\frac{1-\sqrt{5}}{2}\right)^n \\
L(x) &= \frac{2-x}{1-x-x^2}
\end{aligned}$$

8.11.4 Catalan Numbers

$$\begin{aligned}
c_1 &= 1, c_n = \sum_{i=0}^{n-1} c_i c_{n-1-i} = c_{n-1} \frac{4n-2}{n+1} = \frac{\binom{2n}{n}}{n+1} \\
&= \binom{2n}{n} - \binom{2n}{n-1}, c(x) = \frac{1-\sqrt{1-4x}}{2x}
\end{aligned}$$

8.11.5 Stirling Cycle Numbers

Divide n elements into k non-empty cycles.

$$s(n, 0) = 0, s(n, n) = 1, s(n+1, k) = s(n, k-1) - ns(n, k)$$

$$s(n, k) = (-1)^{n-k} \begin{bmatrix} n \\ k \end{bmatrix}$$

$$\begin{bmatrix} n+1 \\ k \end{bmatrix} = n \begin{bmatrix} n \\ k \end{bmatrix} + \begin{bmatrix} n \\ k-1 \end{bmatrix}, \begin{bmatrix} n+1 \\ 2 \end{bmatrix} = n! H_n$$

$$x^{\underline{n}} = x(x-1)\dots(x-n+1) = \sum_k \begin{bmatrix} n \\ k \end{bmatrix} (-1)^{n-k} x^k$$

$$x^{\overline{n}} = x(x+1)\dots(x+n-1) = \sum_k \begin{bmatrix} n \\ k \end{bmatrix} x^k$$

8.11.6 Stirling Subset Numbers

Divide n elements into k non-empty subsets.

$$\begin{aligned}
\left\{ \begin{matrix} n+1 \\ k \end{matrix} \right\} &= k \left\{ \begin{matrix} n \\ k \end{matrix} \right\} + \left\{ \begin{matrix} n \\ k-1 \end{matrix} \right\} \\
x^n &= \sum_k \left\{ \begin{matrix} n \\ k \end{matrix} \right\} x^{\underline{k}} = \sum_k \left\{ \begin{matrix} n \\ k \end{matrix} \right\} (-1)^{n-k} x^{\overline{k}} \\
m! \left\{ \begin{matrix} n \\ m \end{matrix} \right\} &= \sum_k \binom{m}{k} k^n (-1)^{m-k}
\end{aligned}$$

For a fixed k , generating functions :

$$\sum_{n=0}^{\infty} \left\{ \begin{matrix} n \\ k \end{matrix} \right\} x^{n-k} = \prod_{r=1}^k \frac{1}{1-rx}$$

8.11.7 Motzkin Numbers

Draw non-intersecting chords between n points on a circle.

Pick n numbers $k_1, k_2, \dots, k_n \in \{-1, 0, 1\}$ so that $\sum_i^a k_i (1 \leq a \leq n)$ is non-negative and the sum of all numbers is 0.

$$M_{n+1} = M_n + \sum_{i=1}^{n-1} M_i M_{n-1-i} = \frac{(2n+3)M_n + 3nM_{n-1}}{n+3}$$

$$M_n = \sum_{i=0}^{\lfloor \frac{n}{2} \rfloor} \binom{n}{2i} \text{Catalan}(i)$$

$$M(X) = \frac{1-x-\sqrt{1-2x-3x^2}}{2x^2}$$

8.11.8 Eulerian Numbers

Permutations of the numbers 1 to n in which exactly k elements are greater than the previous element.

$$\begin{aligned}\langle n \rangle_k &= (k+1) \langle n-1 \rangle_k + (n-k) \langle n-1 \rangle_{k-1} \\ x^n &= \sum_k \langle n \rangle_k \binom{x+k}{n} \\ \langle n \rangle_m &= \sum_{k=0}^m \binom{n+1}{k} (m+1-k)^n (-1)^k\end{aligned}$$

8.11.9 Harmonic Numbers

Sum of the reciprocals of the first n natural numbers.

$$\begin{aligned}\sum_{k=1}^n H_k &= (n+1)H_n - n \\ \sum_{k=1}^n kH_k &= \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4} \\ \sum_{k=1}^n \binom{k}{m} H_k &= \binom{n+1}{m+1} \left(H_{n+1} - \frac{1}{m+1} \right)\end{aligned}$$

8.11.10 Pentagonal Number Theorem

$$\begin{aligned}\prod_{n=1}^{\infty} (1-x^n) &= \sum_{n=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2} \\ p(n) &= p(n-1) + p(n-2) - p(n-5) - p(n-7) + \dots \\ f(n, k) &= p(n) - p(n-k) - p(n-2k) + p(n-5k) + p(n-7k) - \dots\end{aligned}$$

8.11.11 Bell Numbers

Divide a set that has exactly n elements.

$$\begin{aligned}B_n &= \sum_{k=1}^n \{n \atop k\}, \quad B_{n+1} = \sum_{k=0}^n \binom{n}{k} B_k \\ B_{p^m+n} &\equiv mB_n + B_{n+1} \pmod{p} \\ B(x) &= \sum_{n=0}^{\infty} \frac{B_n}{n!} x^n = e^{e^x-1}\end{aligned}$$

8.11.12 Bernoulli Numbers

$$\begin{aligned}B_n &= 1 - \sum_{k=0}^{n-1} \binom{n}{k} \frac{B_k}{n-k+1} \\ G(x) &= \sum_{k=0}^{\infty} \frac{B_k}{k!} x^k = \frac{1}{\sum_{k=0}^{\infty} \frac{x^k}{(k+1)!}} \\ \sum_{k=1}^n k^m &= \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k n^{m-k+1}\end{aligned}$$

8.11.13 Sum of Powers

$$\begin{aligned}\sum_{i=1}^n i^2 &= \frac{n(n+1)(2n+1)}{6}, \quad \sum_{i=1}^n i^3 = \left(\frac{n(n+1)}{2}\right)^2 \\ \sum_{i=1}^n i^4 &= \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} \\ \sum_{i=1}^n i^5 &= \frac{n^2(n+1)^2(2n^2+2n-1)}{12}\end{aligned}$$

8.11.14 Sum of Squares

Denote $r_k(n)$ the ways to form n with k squares. If :

$$n = 2^{a_0} p_1^{2a_1} \dots p_r^{2a_r} q_1 b_1 \dots q_s b_s$$

where $p_i \equiv 3 \pmod{4}$, $q_i \equiv 1 \pmod{4}$, then

$$r_2(n) = \begin{cases} 0 & \text{if any } a_i \text{ is a half-integer} \\ 4 \prod_{i=1}^r (b_i + 1) & \text{if all } a_i \text{ are integers} \end{cases}$$

$r_3(n) > 0$ when and only when n is not $4^a(8b+7)$.

8.11.15 Derangement

$$\begin{aligned}D_1 &= 0, D_2 = 1, D_n = n! \left(\frac{1}{0!} - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + \frac{(-1)^n}{n!} \right) \\ D_n &= (n-1)(D_{n-1} + D_{n-2})\end{aligned}$$

8.11.16 Tetrahedron Volume

If U, V, W, u, v, w are lengths of edges of the tetrahedron (first three form a triangle; u opposite to U and so on)

$$V = \frac{\sqrt{4u^2v^2w^2 - \sum_{cyc} u^2(v^2+w^2-U^2)^2 + \prod_{cyc} (v^2+w^2-U^2)}}{12}$$

9 Appendix

9.1 Calculus table

$$\begin{aligned}\left(\frac{u}{v}\right)' &= \frac{u'v - uv'}{v^2} & (\operatorname{arcsec} x)' &= \frac{1}{x\sqrt{1-x^2}} \\ (a^x)' &= (\ln a)a^x & (\tanh x)' &= \operatorname{sech}^2 x \\ (\tan x)' &= \sec^2 x & (\coth x)' &= -\operatorname{csch}^2 x \\ (\cot x)' &= -\csc^2 x & (\operatorname{sech} x)' &= -\operatorname{sech} x \tanh x \\ (\sec x)' &= \tan x \sec x & (\operatorname{csch} x)' &= -\operatorname{csch} x \coth x \\ (\csc x)' &= -\cot x \csc x & (\operatorname{arcsinh} x)' &= \frac{1}{\sqrt{1+x^2}} \\ (\arcsin x)' &= \frac{1}{\sqrt{1-x^2}} & (\operatorname{arccosh} x)' &= \frac{1}{\sqrt{x^2-1}} \\ (\arccos x)' &= -\frac{1}{\sqrt{1-x^2}} & (\operatorname{arctanh} x)' &= \frac{1}{1-x^2} \\ (\arctan x)' &= \frac{1}{1+x^2} & (\operatorname{arccoth} x)' &= \frac{1}{x^2-1} \\ (\operatorname{arccot} x)' &= -\frac{1}{1+x^2} & (\operatorname{arccsch} x)' &= -\frac{1}{|x|\sqrt{1+x^2}} \\ (\operatorname{arccsc} x)' &= -\frac{1}{x\sqrt{1-x^2}} & (\operatorname{arcsech} x)' &= -\frac{1}{x\sqrt{1-x^2}}\end{aligned}$$

9.1.1 $ax+b$ ($a \neq 0$)

$$\begin{aligned}1. \int \frac{x}{ax+b} dx &= \frac{1}{a^2} (ax+b - b \ln |ax+b|) + C \\ 2. \int \frac{x^2}{ax+b} dx &= \frac{1}{a^3} \left(\frac{1}{2} (ax+b)^2 - 2b(ax+b) + b^2 \ln |ax+b| \right) + C \\ 3. \int \frac{dx}{x(ax+b)} &= -\frac{1}{b} \ln \left| \frac{ax+b}{x} \right| + C \\ 4. \int \frac{x^2}{x^2(ax+b)} dx &= -\frac{1}{bx} + \frac{a}{b^2} \ln \left| \frac{ax+b}{x} \right| + C \\ 5. \int \frac{x}{(ax+b)^2} dx &= \frac{1}{a^2} \left(\ln |ax+b| + \frac{b}{ax+b} \right) + C \\ 6. \int \frac{x^2}{(ax+b)^2} dx &= \frac{1}{a^3} \left(ax+b - 2b \ln |ax+b| - \frac{b^2}{ax+b} \right) + C \\ 7. \int \frac{dx}{x(ax+b)^2} &= \frac{1}{b(ax+b)} - \frac{1}{b^2} \ln \left| \frac{ax+b}{x} \right| + C\end{aligned}$$

9.1.2 $\sqrt{ax+b}$

$$\begin{aligned}1. \int \sqrt{ax+b} dx &= \frac{2}{3a} \sqrt{(ax+b)^3} + C \\ 2. \int x \sqrt{ax+b} dx &= \frac{2}{15a^2} (3ax-2b) \sqrt{(ax+b)^3} + C \\ 3. \int x^2 \sqrt{ax+b} dx &= \frac{2}{105a^3} (15a^2x^2 - 12abx + 8b^2) \sqrt{(ax+b)^3} + C \\ 4. \int \frac{x}{\sqrt{ax+b}} dx &= \frac{2}{3a^2} (ax-2b) \sqrt{ax+b} + C \\ 5. \int \frac{x^2}{\sqrt{ax+b}} dx &= \frac{2}{15a^3} (3a^2x^2 - 4abx + 8b^2) \sqrt{ax+b} + C \\ 6. \int \frac{dx}{x \sqrt{ax+b}} &= \begin{cases} \frac{1}{\sqrt{b}} \ln \left| \frac{\sqrt{ax+b}-\sqrt{b}}{\sqrt{ax+b}+\sqrt{b}} \right| + C & (b > 0) \\ \frac{2}{\sqrt{-b}} \arctan \sqrt{\frac{ax+b}{-b}} + C & (b < 0) \end{cases} \\ 7. \int \frac{dx}{x^2 \sqrt{ax+b}} &= -\frac{\sqrt{ax+b}}{bx} - \frac{a}{2b} \int \frac{dx}{x \sqrt{ax+b}} \\ 8. \int \frac{\sqrt{ax+b}}{x} dx &= 2\sqrt{ax+b} + b \int \frac{dx}{x \sqrt{ax+b}} \\ 9. \int \frac{\sqrt{ax+b}}{x^2} dx &= -\frac{\sqrt{ax+b}}{x} + \frac{a}{2} \int \frac{dx}{x \sqrt{ax+b}}\end{aligned}$$

9.1.3 $x^2 \pm a^2$

$$\begin{aligned}1. \int \frac{dx}{x^2+a^2} &= \frac{1}{a} \arctan \frac{x}{a} + C \\ 2. \int \frac{dx}{(x^2+a^2)^n} &= \frac{1}{2(n-1)a^2} \frac{1}{(x^2+a^2)^{n-1}} + \frac{2n-3}{2(n-1)a^2} \int \frac{dx}{(x^2+a^2)^{n-1}} \\ 3. \int \frac{dx}{x^2-a^2} &= \frac{1}{2a} \ln \left| \frac{x-a}{x+a} \right| + C\end{aligned}$$

9.1.4 ax^2+b ($a > 0$)

$$\begin{aligned}1. \int \frac{dx}{ax^2+b} &= \begin{cases} \frac{1}{\sqrt{ab}} \arctan \sqrt{\frac{a}{b}} x + C & (b > 0) \\ \frac{1}{2\sqrt{-ab}} \ln \left| \frac{\sqrt{ax+b}-\sqrt{-b}}{\sqrt{ax+b}+\sqrt{-b}} \right| + C & (b < 0) \end{cases} \\ 2. \int \frac{x}{ax^2+b} dx &= \frac{1}{2a} \ln |ax^2+b| + C \\ 3. \int \frac{x^2}{ax^2+b} dx &= \frac{x}{a} - \frac{b}{a} \int \frac{dx}{ax^2+b} \\ 4. \int \frac{dx}{x(ax^2+b)} &= \frac{1}{2b} \ln \left| \frac{x^2}{ax^2+b} \right| + C \\ 5. \int \frac{dx}{x^2(ax^2+b)} &= -\frac{1}{bx} - \frac{a}{b} \int \frac{dx}{ax^2+b} \\ 6. \int \frac{dx}{x^3(ax^2+b)} &= \frac{a}{2b^2} \ln \frac{|ax^2+b|}{x^2} - \frac{1}{2bx^2} + C \\ 7. \int \frac{dx}{(ax^2+b)^2} &= \frac{x}{2b(ax^2+b)} + \frac{1}{2b} \int \frac{dx}{ax^2+b}\end{aligned}$$

9.1.5 ax^2+bx+c ($a > 0$)

$$\begin{aligned}1. \int \frac{dx}{ax^2+bx+c} &= \begin{cases} \frac{2}{\sqrt{4ac-b^2}} \arctan \frac{2ax+b}{\sqrt{4ac-b^2}} + C & (b^2 < 4ac) \\ \frac{1}{\sqrt{b^2-4ac}} \ln \left| \frac{2ax+b-\sqrt{b^2-4ac}}{2ax+b+\sqrt{b^2-4ac}} \right| + C & (b^2 > 4ac) \end{cases} \\ 2. \int \frac{x}{ax^2+bx+c} dx &= \frac{1}{2a} \ln |ax^2+bx+c| - \frac{b}{2a} \int \frac{dx}{ax^2+bx+c}\end{aligned}$$

9.1.6 $\sqrt{x^2+a^2}$ ($a > 0$)

$$\begin{aligned}1. \int \frac{dx}{\sqrt{x^2+a^2}} &= \operatorname{arsh} \frac{x}{a} + C_1 = \ln(x + \sqrt{x^2+a^2}) + C \\ 2. \int \frac{dx}{\sqrt{(x^2+a^2)^3}} &= \frac{1}{a^2} \frac{x}{\sqrt{x^2+a^2}} + C \\ 3. \int \frac{x}{\sqrt{x^2+a^2}} dx &= \sqrt{x^2+a^2} + C \\ 4. \int \frac{x}{\sqrt{(x^2+a^2)^3}} dx &= -\frac{1}{\sqrt{x^2+a^2}} + C \\ 5. \int \frac{x^2}{\sqrt{x^2+a^2}} dx &= \frac{x}{2} \sqrt{x^2+a^2} - \frac{a^2}{2} \ln(x + \sqrt{x^2+a^2}) + C \\ 6. \int \frac{x^2}{\sqrt{(x^2+a^2)^3}} dx &= -\frac{x}{\sqrt{x^2+a^2}} + \ln(x + \sqrt{x^2+a^2}) + C \\ 7. \int \frac{dx}{x \sqrt{x^2+a^2}} &= \frac{1}{a} \ln \left| \frac{\sqrt{x^2+a^2}-a}{|x|} \right| + C \\ 8. \int \frac{dx}{x^2 \sqrt{x^2+a^2}} &= -\frac{\sqrt{x^2+a^2}}{a^2 x} + C \\ 9. \int \sqrt{x^2+a^2} dx &= \frac{x}{2} \sqrt{x^2+a^2} + \frac{a^2}{2} \ln(x + \sqrt{x^2+a^2}) + C \\ 10. \int \sqrt{(x^2+a^2)^3} dx &= \frac{x}{8} (x^2+a^2) \sqrt{x^2+a^2} + \frac{3}{8} a^4 \ln(x + \sqrt{x^2+a^2}) + C\end{aligned}$$

- $\int x\sqrt{x^2+a^2}dx = \frac{1}{3}\sqrt{(x^2+a^2)^3} + C$
- $\int x^2\sqrt{x^2+a^2}dx = \frac{x^3}{8}(2x^2+a^2)\sqrt{x^2+a^2} - \frac{a^4}{8}\ln(x+\sqrt{x^2+a^2}) + C$
- $\int \frac{\sqrt{x^2+a^2}}{x}dx = \sqrt{x^2+a^2} + a\ln\frac{\sqrt{x^2+a^2}-a}{|x|} + C$
- $\int \frac{\sqrt{x^2+a^2}}{x^2}dx = -\frac{\sqrt{x^2+a^2}}{x} + \ln(x+\sqrt{x^2+a^2}) + C$

9.1.7 $\sqrt{x^2-a^2}$ ($a > 0$)

- $\int \frac{dx}{\sqrt{x^2-a^2}} = \frac{1}{|x|}\operatorname{arch}\frac{|x|}{a} + C_1 = \ln|x+\sqrt{x^2-a^2}| + C$
- $\int \frac{dx}{\sqrt{(x^2-a^2)^3}} = -\frac{x}{a^2\sqrt{x^2-a^2}} + C$
- $\int \frac{x}{\sqrt{x^2-a^2}}dx = \sqrt{x^2-a^2} + C$
- $\int \frac{x}{\sqrt{(x^2-a^2)^3}}dx = -\frac{1}{\sqrt{x^2-a^2}} + C$
- $\int \frac{x^2}{\sqrt{x^2-a^2}}dx = \frac{x}{2}\sqrt{x^2-a^2} + \frac{a^2}{2}\ln|x+\sqrt{x^2-a^2}| + C$
- $\int \frac{x^2}{\sqrt{(x^2-a^2)^3}}dx = -\frac{x}{\sqrt{x^2-a^2}} + \ln|x+\sqrt{x^2-a^2}| + C$
- $\int \frac{dx}{x\sqrt{x^2-a^2}} = \frac{1}{a}\arccos\frac{a}{|x|} + C$
- $\int \frac{dx}{x^2\sqrt{x^2-a^2}} = \frac{\sqrt{x^2-a^2}}{a^2x} + C$
- $\int \sqrt{x^2-a^2}dx = \frac{x}{2}\sqrt{x^2-a^2} - \frac{a^2}{2}\ln|x+\sqrt{x^2-a^2}| + C$
- $\int \sqrt{(x^2-a^2)^3}dx = \frac{x^3}{8}(2x^2-5a^2)\sqrt{x^2-a^2} + \frac{3}{8}a^4\ln|x+\sqrt{x^2-a^2}| + C$
- $\int x\sqrt{x^2-a^2}dx = \frac{1}{3}\sqrt{(x^2-a^2)^3} + C$
- $\int x^2\sqrt{x^2-a^2}dx = \frac{x^3}{8}(2x^2-a^2)\sqrt{x^2-a^2} - \frac{a^4}{8}\ln|x+\sqrt{x^2-a^2}| + C$
- $\int \frac{\sqrt{x^2-a^2}}{x}dx = \sqrt{x^2-a^2} - a\arccos\frac{a}{|x|} + C$
- $\int \frac{\sqrt{x^2-a^2}}{x^2}dx = -\frac{\sqrt{x^2-a^2}}{x} + \ln|x+\sqrt{x^2-a^2}| + C$

9.1.8 $\sqrt{a^2-x^2}$ ($a > 0$)

- $\int \frac{dx}{\sqrt{a^2-x^2}} = \arcsin\frac{x}{a} + C$
- $\frac{dx}{\sqrt{(a^2-x^2)^3}} = \frac{x}{a^2\sqrt{a^2-x^2}} + C$
- $\int \frac{x}{\sqrt{a^2-x^2}}dx = -\sqrt{a^2-x^2} + C$
- $\int \frac{x}{\sqrt{(a^2-x^2)^3}}dx = \frac{1}{\sqrt{a^2-x^2}} + C$
- $\int \frac{x^2}{\sqrt{a^2-x^2}}dx = -\frac{x}{2}\sqrt{a^2-x^2} + \frac{a^2}{2}\arcsin\frac{x}{a} + C$
- $\int \frac{x^2}{\sqrt{(a^2-x^2)^3}}dx = \frac{x}{\sqrt{a^2-x^2}} - \arcsin\frac{x}{a} + C$
- $\int \frac{dx}{x\sqrt{a^2-x^2}} = \frac{1}{a}\ln\frac{a-\sqrt{a^2-x^2}}{|x|} + C$
- $\int \frac{dx}{x^2\sqrt{a^2-x^2}} = -\frac{\sqrt{a^2-x^2}}{a^2x} + C$
- $\int \sqrt{a^2-x^2}dx = \frac{x}{2}\sqrt{a^2-x^2} + \frac{a^2}{2}\arcsin\frac{x}{a} + C$
- $\int \sqrt{(a^2-x^2)^3}dx = \frac{x^3}{8}(5a^2-2x^2)\sqrt{a^2-x^2} + \frac{3}{8}a^4\arcsin\frac{x}{a} + C$
- $\int x\sqrt{a^2-x^2}dx = -\frac{1}{3}\sqrt{(a^2-x^2)^3} + C$
- $\int x^2\sqrt{a^2-x^2}dx = \frac{x^3}{8}(2x^2-a^2)\sqrt{a^2-x^2} + \frac{a^4}{8}\arcsin\frac{x}{a} + C$
- $\int \frac{\sqrt{a^2-x^2}}{x}dx = \sqrt{a^2-x^2} + a\ln\frac{a-\sqrt{a^2-x^2}}{|x|} + C$
- $\int \frac{\sqrt{a^2-x^2}}{x^2}dx = -\frac{\sqrt{a^2-x^2}}{x} - \arcsin\frac{x}{a} + C$

9.1.9 $\sqrt{\pm ax^2+bx+c}$ ($a > 0$)

- $\int \frac{dx}{\sqrt{ax^2+bx+c}} = \frac{1}{\sqrt{a}}\ln|2ax+b+2\sqrt{a}\sqrt{ax^2+bx+c}| + C$
- $\int \sqrt{ax^2+bx+c}dx = \frac{2ax+b}{4a}\sqrt{ax^2+bx+c} + \frac{4ac-b^2}{8\sqrt{a^3}}\ln|2ax+b+2\sqrt{a}\sqrt{ax^2+bx+c}| + C$
- $\int \frac{x}{\sqrt{ax^2+bx+c}}dx = \frac{1}{a}\sqrt{ax^2+bx+c} - \frac{b}{2\sqrt{a^3}}\ln|2ax+b+2\sqrt{a}\sqrt{ax^2+bx+c}| + C$
- $\int \frac{dx}{\sqrt{c+bx-ax^2}} = -\frac{1}{\sqrt{a}}\arcsin\frac{2ax-b}{\sqrt{b^2+4ac}} + C$
- $\int \sqrt{c+bx-ax^2}dx = \frac{2ax-b}{4a}\sqrt{c+bx-ax^2} + \frac{b^2+4ac}{8\sqrt{a^3}}\arcsin\frac{2ax-b}{\sqrt{b^2+4ac}} + C$
- $\int \frac{x}{\sqrt{c+bx-ax^2}}dx = -\frac{1}{a}\sqrt{c+bx-ax^2} + \frac{b}{2\sqrt{a^3}}\arcsin\frac{2ax-b}{\sqrt{b^2+4ac}} + C$

9.1.10 $\sqrt{\pm\frac{x-a}{x-b}}$ & $\sqrt{(x-a)(x-b)}$

- $\int \sqrt{\frac{x-a}{x-b}}dx = (x-b)\sqrt{\frac{x-a}{x-b}} + (b-a)\ln(|\sqrt{x-a}|+|\sqrt{x-b}|) + C$
- $\int \sqrt{\frac{x-a}{b-x}}dx = (x-b)\sqrt{\frac{x-a}{b-x}} + (b-a)\arcsin\sqrt{\frac{x-a}{b-x}} + C$
- $\int \frac{dx}{\sqrt{(x-a)(b-x)}} = 2\arcsin\sqrt{\frac{x-a}{b-x}} + C$ ($a < b$)
- $\int \sqrt{(x-a)(b-x)}dx = \frac{2x-a-b}{4}\sqrt{(x-a)(b-x)} + \frac{(b-a)^2}{4}\arcsin\sqrt{\frac{x-a}{b-x}} + C$ ($a < b$)

9.1.11 Triangular function

- $\int \tan x dx = -\ln|\cos x| + C$
- $\int \cot x dx = \ln|\sin x| + C$
- $\int \sec x dx = \ln\left|\tan\left(\frac{x}{2} + \frac{\pi}{2}\right)\right| + C = \ln|\sec x + \tan x| + C$
- $\int \csc x dx = \ln\left|\tan\frac{x}{2}\right| + C = \ln|\csc x - \cot x| + C$
- $\int \sec^2 x dx = \tan x + C$
- $\int \csc^2 x dx = -\cot x + C$
- $\int \sec x \tan x dx = \sec x + C$
- $\int \csc x \cot x dx = -\csc x + C$
- $\int \sin^2 x dx = \frac{x}{2} - \frac{1}{4}\sin 2x + C$
- $\int \cos^2 x dx = \frac{x}{2} + \frac{1}{4}\sin 2x + C$

- $\int \sin^n x dx = -\frac{1}{n}\sin^{n-1}x \cos x + \frac{n-1}{n}\int \sin^{n-2}x dx$
- $\int \cos^n x dx = \frac{1}{n}\cos^{n-1}x \sin x + \frac{n-1}{n}\int \cos^{n-2}x dx$
- $\frac{dx}{\sin^n x} = -\frac{1}{n-1}\frac{\cos x}{\sin^{n-1}x} + \frac{n-2}{n-1}\int \frac{dx}{\sin^{n-2}x}$
- $\frac{dx}{\cos^n x} = \frac{1}{n-1}\frac{\sin x}{\cos^{n-1}x} + \frac{n-2}{n-1}\int \frac{dx}{\cos^{n-2}x}$

$$\begin{aligned} & \int \cos^m x \sin^n x dx \\ &= \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x dx \\ &= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x dx \end{aligned}$$

- $\int \sin ax \cos bx dx = -\frac{1}{2(a+b)} \cos(a+b)x - \frac{1}{2(a-b)} \cos(a-b)x + C$
- $\int \sin ax \sin bx dx = -\frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$
- $\int \cos ax \cos bx dx = \frac{1}{2(a+b)} \sin(a+b)x + \frac{1}{2(a-b)} \sin(a-b)x + C$

$$\begin{aligned} 19. \int \frac{dx}{a+b\sin x} &= \begin{cases} \frac{2}{\sqrt{a^2-b^2}} \arctan \frac{a \tan \frac{x}{2} + b}{\sqrt{a^2-b^2}} + C & (a^2 > b^2) \\ \frac{1}{\sqrt{b^2-a^2}} \ln \left| \frac{a \tan \frac{x}{2} + b - \sqrt{b^2-a^2}}{a \tan \frac{x}{2} + b + \sqrt{b^2-a^2}} \right| + C & (a^2 < b^2) \end{cases} \\ 20. \int \frac{dx}{a+b\cos x} &= \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan \left(\sqrt{\frac{a-b}{a+b}} \tan \frac{x}{2} \right) + C & (a^2 > b^2) \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln \left| \frac{\tan \frac{x}{2} + \sqrt{\frac{a+b}{a-b}}}{\tan \frac{x}{2} - \sqrt{\frac{a+b}{a-b}}} \right| + C & (a^2 < b^2) \end{cases} \end{aligned}$$

- $\int \frac{dx}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{ab} \arctan \left(\frac{b}{a} \tan x \right) + C$
- $\int \frac{dx}{a^2 \cos^2 x - b^2 \sin^2 x} = \frac{1}{2ab} \ln \left| \frac{b \tan x + a}{b \tan x - a} \right| + C$
- $\int x \sin ax dx = -\frac{1}{a^2} \sin ax - \frac{1}{a} x \cos ax + C$
- $\int x^2 \sin ax dx = -\frac{1}{a} x^2 \cos ax + \frac{2}{a^2} x \sin ax + \frac{2}{a^3} \cos ax + C$
- $\int x \cos ax dx = \frac{1}{a^2} \cos ax + \frac{1}{a} x \sin ax + C$
- $\int x^2 \cos ax dx = \frac{1}{a} x^2 \sin ax + \frac{2}{a^2} x \cos ax - \frac{2}{a^3} \sin ax + C$

9.1.12 Inverse triangular function ($a > 0$)

- $\int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2-x^2} + C$
- $\int x \arcsin \frac{x}{a} dx = \left(\frac{x^2}{2} - \frac{a^2}{4}\right) \arcsin \frac{x}{a} + \frac{x}{4}\sqrt{x^2-x^2} + C$
- $\int x^2 \arcsin \frac{x}{a} dx = \frac{x^3}{3} \arcsin \frac{x}{a} + \frac{1}{9}(x^2+2a^2)\sqrt{a^2-x^2} + C$
- $\int \arccos \frac{x}{a} dx = x \arccos \frac{x}{a} - \sqrt{a^2-x^2} + C$
- $\int x \arccos \frac{x}{a} dx = \left(\frac{x^2}{2} - \frac{a^2}{4}\right) \arccos \frac{x}{a} - \frac{x}{4}\sqrt{a^2-x^2} + C$
- $\int x^2 \arccos \frac{x}{a} dx = \frac{x^3}{3} \arccos \frac{x}{a} - \frac{1}{9}(x^2+2a^2)\sqrt{a^2-x^2} + C$
- $\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} - \frac{a}{2} \ln(a^2+x^2) + C$
- $\int x \arctan \frac{x}{a} dx = \frac{1}{2}(a^2+x^2) \arctan \frac{x}{a} - \frac{a}{2} x + C$
- $\int x^2 \arctan \frac{x}{a} dx = \frac{x^3}{3} \arctan \frac{x}{a} - \frac{a}{6} x^2 + \frac{a^3}{6} \ln(a^2+x^2) + C$

9.1.13 Exponential function

- $\int a^x dx = \frac{1}{\ln a} a^x + C$
- $\int e^{ax} dx = \frac{1}{a} a^x + C$
- $\int x e^{ax} dx = \frac{1}{a^2} (ax-1) a^x + C$
- $\int x^n e^{ax} dx = \frac{1}{a} x^n e^{ax} - \frac{n}{a} \int x^{n-1} e^{ax} dx$
- $\int x a^x dx = \frac{x}{\ln a} a^x - \frac{1}{(\ln a)^2} a^x + C$
- $\int x^n a^x dx = \frac{1}{\ln a} x^n a^x - \frac{n}{\ln a} \int x^{n-1} a^x dx$
- $\int e^{ax} \sin bx dx = \frac{1}{a^2+b^2} e^{ax} (a \sin bx - b \cos bx) + C$
- $\int e^{ax} \cos bx dx = \frac{1}{a^2+b^2} e^{ax} (b \sin bx + a \cos bx) + C$
- $\int e^{ax} \sin^n bx dx = \frac{1}{a^2+b^2n^2} e^{ax} \sin^{n-1} bx (a \sin bx - nb \cos bx) + \frac{n(n-1)b^2}{a^2+b^2n^2} \int e^{ax} \sin^{n-2} bx dx$
- $\int e^{ax} \cos^n bx dx = \frac{1}{a^2+b^2n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) + \frac{n(n-1)b^2}{a^2+b^2n^2} \int e^{ax} \cos^{n-2} bx dx$

9.1.14 Logarithmic function

- $\int \ln x dx = x \ln x - x + C$
- $\int x \ln x dx = \ln|\ln x| + C$
- $\int x^n \ln x dx = \frac{1}{n+1} x^{n+1} (\ln x - \frac{1}{n+1}) + C$
- $\int (\ln x)^n dx = x (\ln x)^n - n \int (\ln x)^{n-1} dx$
- $\int x^m (\ln x)^n dx = \frac{1}{m+1} x^{m+1} (\ln x)^n - \frac{n}{m+1} \int x^m (\ln x)^{n-1} dx$

9.2 Regular expression

9.2.1 Special pattern characters

Characters	Description
.	Not newline
\t	Tab (HT)
\n	Newline (LF)
\v	Vertical tab (VT)
\f	Form feed (FF)
\r	Carriage return (CR)
\cletter	Control code
\xhh	ASCII character
\uhhhh	Unicode character
\0	Null
\int	Backreference
\d	Digit
\D	Not digit
\s	Whitespace
\S	Not whitespace
\w	Word (letters, numbers and the underscore)
\W	Not word
\character	Character
[class]	Character class
[^class]	Negated character class

9.2.2 Quantifiers

Characters	Times
*	0 or more
+	1 or more
?	0 or 1
{int}	int
{int,}	int or more
{min,max}	Between min and max

By default, all these quantifiers are greedy (i.e., they take as many characters that meet the condition as possible). This behavior can be overridden to ungreedy (i.e., take as few characters that meet the condition as possible) by adding a question mark (?) after the quantifier.

9.2.3 Groups

Characters	Description
(subpattern)	Group with backreference
(?:subpattern)	Group without backreference

9.2.4 Assertions

Characters	Description
^	Beginning of line
\$	End of line
\b	Word boundary
\B	Not a word boundary
(?=subpattern)	Positive lookahead
(?!subpattern)	Negative lookahead

9.2.5 Alternative

A regular expression can contain multiple alternative patterns simply by separating them with the separator operator (|): The regular

expression will match if any of the alternatives match, and as soon as one does.

9.2.6 Character classes

Class	Description
[:alnum:]	Alpha-numerical character
[:alpha:]	Alphabetic character
[:blank:]	Blank character
[:cntrl:]	Control character
[:digit:]	Decimal digit character
[:graph:]	Character with graphical representation
[:lower:]	Lowercase letter
[:print:]	Printable character
[:punct:]	Punctuation mark character
[:space:]	Whitespace character
[:upper:]	Uppercase letter
[:xdigit:]	Hexadecimal digit character
[:d:]	Decimal digit character
[:w:]	Word character
[:s:]	Whitespace character

Please note that the brackets in the class names are additional to those opening and closing the class definition. For example:

[[:alpha:]] is a character class that matches any alphabetic character.

[abc[:digit:]] is a character class that matches a, b, c, or a digit.

[^[:space:]] is a character class that matches any character except a whitespace.