Some Filesystem Design Notes

[An Example of the components Required for the implementation]

1.1 Directory Structure (Define the type of Directory Structure you expect to use)

For example:

File Allocation (FA) Structure: logical or user level directory of files and folders

Describe what the FA structure is, what it contains, and how it fits into the overall design of the filesystem.

Show the format and detail of entries in the structure/table with enough detail to completely specify the data structures required for implementation. Describe storage sizes required for entries and types of identifiers. Include an example of the content of the directory structure for a set of files/directories. Differentiate the differences and descriptive details of directory data structures vs. file data structures in the logical directory.

Physical Allocation (directory of storage units on the disk and how they are related/allocated to files) You will need to decide how the disk is to be organized in terms of blocks of byte that become the 'allocation units' (size of a storage unit). As a file is created and grows it is by allocation units.

Describe storage unit size, how each storage unit on the disk is allocated and associated with the file, directory or system space to which it belongs (let's call them storage containers). Show an example of allocation of the units to each of the functional storage containers Describe how you address storage units on the disk.

Describe the integration of the logical and physical directories that are defined above.

Definitions

Include any definitions that describe naming conventions or other terms you use for referencing storage units, files, directories, etc.

2.1 Details of Implementation

Overview

Describe storage unit sizes, overhead for storage, location of various filesystem entities on the virtual disk, etc.

3.1 File Interface Code Definitions and Pseudo-code

Define how your filesystem functions will reference files (for example file handles when open) and the data structures you will design to support the method of referencing the files. [do you need tables of data structures to manage open files, positions in files, etc.?]

Define how folders are referenced.

Define how data are read/written and or buffered between disk and user application.

Define and describe each of the functions that you have implemented to manipulate the filesystem, files and folders. Include the function calls and parameter definitions. Be sure to include documentation of any special features or functionality you are including.

Include a description of the function (and its functionality) with details in English. Also include pseudocode of the function implementation.

4.1 Notes on programming language issues and related use of File I/O functions of the run-time library (for example C run-time library) or Memory Mapped Files

Whether you choose to implement your system with direct file I/O or through memory mapped files, you will have to define and construct the functions that perform the actual manipulation of the data on your virtual disk (this includes directory data and file data). Your will need to create the equivalent of the functions that the OS system library provides for device I/O or the equivalents of the functions that the C run-time library provides.

That is you need to implement the real (physical) file operations for your file system. You are to explain and detail the system or run-time library functions you will be using and how you will use them. Include enough detail to show you understand the functions, their features and limitations. Show prototypes and argument definitions for each function.

1. For example, describe the use of the C Stream I/O Functions that will be called within your 'wrapper functions'.

Sample discussion:

The **fread** function is used to read data from a file (identified by a FILE structure), without formatting.

The **fread** function reads up to count items of size bytes from the input stream and stores them in buffer. The file pointer associated with stream (if there is one) is increased by the number of bytes actually read. If the given stream is opened in text mode, carriage return—linefeed pairs are replaced with single linefeed characters. The replacement has no effect on the file pointer or the return value. The file-pointer position is indeterminate if an error occurs. The value of a partially read item cannot be determined.

Prototype:

size_t fread (void * Buffer, size_t Size, size_t Count, FILE * Stream);

Arguments:

Stream: Pointer to an open FILE structure.

Buffer: The address where the data read is to be stored.

Count: Number of data elements to be read.

Size: Size of each data element to be read (in bytes.)

Returns:

size_t Number of data elements read.

The **fwrite** function is used to write data to a file (identified by a FILE structure), without formatting , ...

The fopen function is used to open a file (identified by FILE structure), ...

The **fclose** function is used to close an open FILE stream. If the stream was being used for output, any buffered data is written first, ...

The **fseek** function is used to change the position of a file pointer (identified by FILE stream) , ...

The sscanf function reads data from buffer into the location given by each argument, ...

The sprintf function formats and stores a series of characters and values in buffer, ...

2. Example – using the memory mapped files facility

As an alternative to standard file I/O, the kernel provides an interface that allows an application to map a file into memory, meaning that there is a one-to-one correspondence between a memory address and a word in the file. A programmer can then access the file directly through memory, identically to any other chunk of memory-resident data—it is even possible to allow writes to the memory region to transparently map back to the file on disk.

POSIX.1 standardizes the mmap() system call for mapping objects into memory. mmap(): void * mmap (void *addr,

```
size_t len,
int prot,
int flags,
int fd,
off_t offset);
```

A call to mmap() asks the kernel to map len bytes of the object represented by the file descriptor fd, starting at offset bytes into the file, into memory. If addr is included, it indicates a preference to use that starting address in memory.

```
A typical use is (explore the meaning of the parameters) p = mmap (0, len, PROT_READ, MAP_SHARED, fd, 0);
```

You will need to construct the file manipulation functions to act on the mapped memory region as though blocks of a file are read or written.