STEPS OF NORMALIZATION

Step Description

- 1 Begin with a table (UNF) with all data combined
- 2 Apply 1NF: Eliminate repeating groups / multivalued columns
- 3 Apply 2NF: Eliminate partial dependencies (if composite PK)
- 4 Apply 3NF: Eliminate transitive dependencies
- 5 Optionally, apply BCNF or higher (if determinant is not a candidate key)
- NORMAL FORMS (NF)
- 1NF (First Normal Form)
 - Atomic values only (no arrays, comma-separated values)
 - Remove repeating groups
 - Ensure each row is uniquely identified (usually with a primary key)
- 2NF (Second Normal Form)
 - Table should be in 1NF
 - No partial dependency on composite primary key
 - All non-key columns depend on the entire primary key
- ✓ 3NF (Third Normal Form)
 - Table should be in 2NF
 - No transitive dependency
 - Every non-key attribute must depend only on the primary key
- BCNF (Boyce-Codd Normal Form)
 - Stricter version of 3NF
 - Every determinant must be a candidate key

▼ FULL WORKING EXAMPLE – COVERS ALL NFs UPTO 3NF

© Scenario: College Course Registration

A university maintains a single table like this:

StudentID StudentName Course Instructor InstructorPhone

S101	Alice	Java	Smith	9876543210
S101	Alice	DBMS	John	9123456789
S102	Bob	DBMS	John	9123456789

UNF (Unnormalized Form)

- Multivalued: One student has multiple courses in same row OR multiple rows repeating full details
- Repetition of student name, instructor details

● 1NF (First Normal Form)

Make sure each field is atomic and single valued.

StudentID StudentName Course Instructor InstructorPhone

S101	Alice	Java	Smith	9876543210
S101	Alice	DBMS	John	9123456789
S102	Bob	DBMS	John	9123456789

Now: ✓ Atomic values. But: X Still data repetition → Instructor details are repeated.

2NF (Second Normal Form)

Condition: Remove partial dependency

Let's assume composite key = (StudentID, Course)

- StudentName depends only on StudentID → violates 2NF
- InstructorPhone depends only on Instructor → violates both PKs

Split into 3 tables:

Table: Student

StudentID StudentName

S101 Alice

S102 Bob

💄 Table: CourseRegistration

StudentID Course

S101 Java

S101 DBMS

S102 DBMS

📳 Table: Instructor

Course Instructor InstructorPhone

Java Smith 9876543210

DBMS John 9123456789

✓ All non-key attributes now fully depend on **entire PK** of their respective tables.

3NF (Third Normal Form)

Check for transitive dependencies

In Instructor table:

- Course → Instructor
- Instructor → InstructorPhone → X Transitive Dependency

Fix it by separating Instructor details

Table: Course

Course Instructor

Java Smith

Course Instructor

DBMS John

L Table: Instructor

Instructor InstructorPhone

Smith 9876543210

John 9123456789

✓ Now everything is in 3NF

TINAL STRUCTURE

- 1. Student(StudentID, StudentName)
- 2. CourseRegistration(StudentID, Course)
- 3. Course(Course, Instructor)
- 4. Instructor(Instructor, InstructorPhone)

BCNF (Boyce-Codd Normal Form) Example

First, What's the Problem?

Even if a table is in **3NF**, there can still be **anomalies** if:

A **non-primary attribute (determinant)** determines another attribute AND

That determinant is not a candidate key

This is where **BCNF** is needed.

EXAMPLE: College Room Allocation

Table: RoomAllocation

StudentID RoomNo Warden

S1 R1 W1

S2 R1 W1

S3 R2 W2

S4 R3 W3

Functional Dependencies

Let's identify the Functional Dependencies (FDs):

- 1. StudentID → RoomNo (Each student is assigned one room)
- 2. RoomNo → Warden <a> (Each room has one warden)

Primary Key = StudentID

So all **non-key attributes** should depend **only** on StudentID

- → Let's check:
 - RoomNo depends on StudentID
 - But Warden depends on RoomNo, not StudentID X

✓ Is it in 3NF?

Check Transitive Dependency:

StudentID → RoomNo → Warden

So StudentID → Warden is **transitive** → We can move RoomNo → Warden to a new table.

Split Tables for 3NF:

◆ Table 1: StudentRoom

StudentID RoomNo

S1 R1

S2 R1

S3 R2

S4 R3

Table 2: RoomWarden

RoomNo Warden

R1 W1

R2 W2

R3 W3

☑ Now all non-key attributes depend **only** on the PK.

So this is in **3NF**.

X But... Not Yet in BCNF

Look at RoomNo → Warden

- RoomNo is a determinant
- But is RoomNo a candidate key in original table?

NO – because in the original RoomAllocation, the **PK was StudentID**, and RoomNo was not unique

So 3NF allowed this, but BCNF does not

✓ To Fix for BCNF

Split the original table into:

Table 1: StudentRoom (StudentID → RoomNo)

StudentID RoomNo

- S1 R1
- S2 R1
- S3 R2
- S4 R3

Table 2: RoomWarden (RoomNo → Warden)

RoomNo Warden

- R1 W1
- R2 W2
- R3 W3

Now, both tables are in **BCNF**:

- Every determinant is a candidate key
- No anomalies left

★ Key Takeaways for BCNF:

Concept	Rule
Determinant	Attribute that determines others (e.g., A → B)
Candidate Key	A column or combination that can uniquely identify each row
When to Apply BCNF	If a non-candidate key is a determinant
3NF ≠ BCNF if	There is a non-key determinant
Goal of BCNF	Eliminate all anomalies even from non-key dependencies

▼ TRICKY CONCEPTS TO WATCH FOR:

Tricky Concept	Why It's Confusing
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Functional vs Transitive FD is direct (A \rightarrow B), transitive is indirect (A \rightarrow B \rightarrow C,

Dependency so $A \rightarrow C$)

2NF applies only to composite keys If no composite key → 2NF automatically satisfied

BCNF stricter than 3NF

All determinants must be candidate keys, not just

PΚ

Denormalization vs 1NF removes

multi-valued fields

Multivalued dependency vs

Repeating Group

Multivalued is more advanced → applies in 4NF

© TRICKY MCQs

Q1. Which of the following tables is automatically in 2NF?

- A. A table with only atomic columns
- B. A table with a single-column primary key
- C. A table with transitive dependency
- D. A table with composite primary key and partial dependency
- Answer: B
- 🖈 Because 2NF issues arise only in composite primary keys.

\triangleleft Q2. A \rightarrow B and B \rightarrow C. What does this imply?

- A. Transitive dependency exists
- B. Functional dependency is violated
- C. 2NF is maintained
- D. Candidate key is B
- Answer: A
- 📌 This is a classic **transitive dependency** → must be removed in 3NF.

Q3. Which of the following violates 3NF? A. Every attribute is fully dependent on PK B. There is a column dependent on a non-key attribute

- C. There is no repeating group
- D. All columns are atomic
- Answer: B
- \star If a column depends on a non-key column, it breaks 3NF \to transitive dependency.
- Q4. Which NF ensures there are no partial and no transitive dependencies?
- A. 2NF
- B. 3NF
- C. 1NF
- D. BCNF
- Answer: B
- ★ 2NF removes partial, 3NF removes both partial and transitive.
- Q5. A determinant must be a candidate key is a rule of:
- A. 3NF
- B. 2NF
- C. BCNF
- D. 4NF
- Answer: C
- This is the strict rule of BCNF.