

Given:

$$a_1 = 2, a_2 = 3, a_3 = 1$$

$$x \equiv 2 \cdot 20 \cdot 2 + 3 \cdot 15 \cdot 3 + 1 \cdot 12 \cdot 3 \pmod{60}$$

$$x \equiv 80 + 135 + 36 \equiv 251 \pmod{60}$$

$$x \equiv 251 \pmod{60} \Rightarrow x \equiv 11 \pmod{60}$$

Q. 4: Carmichael number:

$$a^{n-1} \equiv 1 \pmod{n}$$

Step-1:

$$561 = 3 \times 187 = 3 \times 11 \times 17$$

561 prime factors: 3, 11, 17

Step-2:

prime factorization: 3, 11, 17

1. n is square-free (no repeated prime factors)

2. For every time p dividing $n, p-1$ divides

$$n-1$$

$$p-1 \text{ divide } 561-1 = 560$$

$$(1) 60 \mid 560 \text{ and } 11 \mid 560 \text{ and } 17 \mid 560$$

Step-3:

$$a^{560} \equiv 1 \pmod{561}$$

• $a=2$

$$2^{560} \pmod{561} = 1$$

• $a=10$:

$$10^{560} \pmod{561} = 1$$

• $a=13$:

$$13^{560} \pmod{561} = 1$$

Q.5:

Step-1

$$\phi(p) = p-1$$

$$\phi(p-1) = \{1, 2, \dots, p-1\} \pmod{p}$$

$$\{g^1, g^2, \dots, g^{p-1}\} \pmod{p}$$

$$p-1 \pmod{p}$$

$$p=17, \phi(17)=16$$

$$g^k \not\equiv 1 \pmod{17} \text{ for all } k < 16, \text{ but } g^{16} \equiv 1 \pmod{17}$$

Step-2:

$$g^{16/9} \not\equiv 1 \pmod{17}, 4 \text{ of } 16$$

$$16 = 2^4 \Rightarrow 4=2 \Rightarrow g^8, g^4, g^2 \not\equiv 1 \pmod{17}$$

Step-3:

$$g = 3$$

$$\bullet 3^1 \bmod 17 = 9$$

$$\bullet 3^4 = (3^4)^{\vee} = 81 \bmod 17 = 13$$

$$\bullet 3^8 = (3^4)^{\vee} = 13^{\vee} = 169 \bmod 17 = 16$$

$$\bullet 3^{16} = (3^8)^{\vee} = 16^{\vee} = 256 \bmod 17 = 16$$

powers = 1, only $3^{16} \equiv 1 \bmod 17$: 01 = 00

Q-6:

$$3^x \equiv 13 \bmod 17$$

step-by-step solutions

$$\xrightarrow{x} \frac{3^x \bmod 17}{3}$$

1

2

3^4

4

$$3^4 \equiv 81 \bmod 17 = 13$$

$$x = 4 \text{ since, } 3^4 \equiv 13 \bmod 17$$

∴ 400 $\bmod 17$ = 13

$$400 \bmod 17 = 3^4 \bmod 17 = 81 \bmod 17 = 13$$

Q-7:

- p : a Large prime $\nmid 10$ and 16 .
- g : a primitive root modulo p .
- $g^a \bmod p$: party A's public key mod p .
- $g^b \bmod p$: party B's public key mod p .

steps of the protocol:

1. public parameters.

2. private secrets.

3. Exchange.

4. shared secret computation.

Discrete logarithm:

- $g^a \bmod p$ \leftarrow (any a mod 10 or 16)
- $g^a \bmod p^{10} \leftarrow$ (any a mod 10)
- $g^b \bmod p^{16} \leftarrow$ (any b mod 16)
- $g^a \equiv A \pmod{p^{16}}$

Security based on the DLP:

• $p > 2^{2048}$

• $p \equiv 1 \pmod{4}$, $16 \nmid p-1$

• DLP in \mathbb{Z}_p^* .

Q-8:

1. substitution ciphering

Alphabet: $A \rightarrow D, B \rightarrow E, \dots, X \rightarrow A, Y \rightarrow B, Z \rightarrow C$

plaintext: HELLOWORLD

ciphertext: KHOORZRUOGC

result: "KHOORZRUOGC"

plaintext: HELLOWORLD

digraphs: HE, LL, OW, OR, LD

• HE \rightarrow H (row 2, col 2), E (row 3, col 1) \rightarrow CF

• LL \rightarrow LX \rightarrow L (4, 1) \times (5, 4) \rightarrow VS
L \rightarrow with filler X \rightarrow VS

• OW \rightarrow O (1, 2) W (5, 3) \rightarrow NW

• OR \rightarrow O (1, 2), R (1, 5) \rightarrow NM

• LD \rightarrow L (4, 1), D (2, 5) \rightarrow TR

result: "CFUSVSNWMTTR"

Q-9: (a) Using RSA with p=13, q=17

• Encryption function: $E(x) = (ax + b) \bmod 26$

$$E(x) = (5x + 8) \bmod 26$$

$$\bullet a = 5$$

$$\bullet b = 8$$

Plaintext: "Dept of ICT, MBSTU"

• x is the plaintext letter ($A=0, B=1, \dots, Z=25$)

• $\gcd(a, 26) = 1$ so, $a=5$ is valid.

Part-a: Encrypt the plaintext in vTL.

Step-1

Dept of ICT MBSTU \Rightarrow DEPT OF ICT MBSTU

Step-2

$$A=0, B=1 \dots, Z=25$$

Part-B: Decryption in vTL

Step-1

$$D(y) = a^{-1} \cdot (y - b) \bmod 26$$

$$\bullet a=5$$

$$\bullet a^{-1} \bmod 26 = 21, 5 \cdot 21 = 105 \equiv 1 \bmod 26$$

Decrypted plaintext:

DEPT OF ICT MBSTU

Q-10: Encryption process:

1. key and block structure with function
2. key derivation: $(K + S_{0,1}) = C_1 \oplus$
3. substitution step: $S = S_{0,1}$
4. permutation step: $P = P_{0,1}$

Decryption process:

1. key derivation
2. inverse permutation: P^{-1}
3. inverse substitution: S^{-1}

Basic cryptanalysis of SP-shift cipher

- simple and lightweight for educational use
- PRNG-driven substitution and permutation observe direct letter-to-letter-mapping
- multi-round application increases confusion and diffusion.

improvements

- use a better PRNG or a key schedule algorithm.
- introduce inter-block chaining

Q Is 1729 a Carmichael number?

Ans, 1729 is a Carmichael number.

A Carmichael number is a composite number n that satisfies the congruence

$$a^{n-1} \equiv 1 \pmod{n}$$

for all integers a that are relatively prime to n .

- 1729 is not prime $1729 = 7 \times 13 \times 19$

- All the prime factors of 1729 are distinct.

- For Carmichael numbers, Korselt's criterion says:

A number n is a Carmichael number if and only if:

1. n is composite,

A number n is a Carmichael number if and only if:

- for every prime divisor p of n , $p-1$ divides $n-1$

Prime factors : 7, 13, 19

$$\bullet n-1 = 1728$$

$$\bullet 7-1 = 6, \text{ and } 6 \mid 1728$$

$$\bullet 13-1 = 12, \text{ and } 12 \mid 1728$$

$$\bullet 19-1 = 18, \text{ and } 18 \mid 1728$$

So, 1729 is a Carmichael number.

1729 is also famous as the Hardy-Ramanujan number, being the smallest number.

$$1729 = 1^3 + 12^3 = 9^3 + 10^3$$

2. Primitive Root or generator of \mathbb{Z}_{23}^*

- \mathbb{Z}_{23}^* is the multiplicative group modulo 23.
- since 23 is prime, \mathbb{Z}_{23}^* has order $23-1 = 22$.
- A primitive root modulo 23 is an integer g such that the smallest κ for which $g^\kappa \equiv 1 \pmod{23}$ is $\kappa = 22$.

Step-2

$$g^{\frac{22}{p}} \not\equiv 1 \pmod{23}$$

- Prime divisors of 22: 2, 11
- so, we must check:
 - $g^1 \not\equiv 1 \pmod{23}$
 - $g^2 \not\equiv 1 \pmod{23}$

Step-3

$$g = 5$$

$$\Rightarrow 5^2 = 25 \equiv 2 \pmod{23} \Rightarrow \not\equiv 1$$

$$\bullet 5^4 \pmod{23}:$$

$$\bullet 5^2 = 25 \equiv 2$$

$$\bullet 5^4 = (5^2)^2 = 2^2 = 4$$

$$\bullet 5^8 = (5^4)^2 = 4^2 = 16$$

$$\bullet 5^{11} = 5^8 \cdot 5^2 \cdot 5 = 16 \cdot 2 \cdot 5 = 160 \pmod{23}$$

$$\bullet 160 \pmod{23} = 160 - 6 \cdot 23 = 160 - 138 = 22 \not\equiv 1$$

5 is a primitive root modulo 23

roots mod 23 include: 5, 7, 10, 11, 14, 15, 17, 19, 20, 21.

3. Is $\mathbb{Z}_{11}, +, \cdot >$ a Ring?

Yes, $\mathbb{Z}_{11} = (\mathbb{Z}_{11}, +, \cdot)$ is a ring.

Under addition (+):

- 1. Closure: $a + b \in \mathbb{Z}_{11}$
- 2. Associativity: $(a+b)+c = a+(b+c)$
- 3. Identity: $0 \in \mathbb{Z}_{11}$ such that $a+0=a$
- 4. Inverses: Every $a \in \mathbb{Z}_{11}$ has an additive inverse $-a \in \mathbb{Z}_{11}$
- 5. Commutativity: $a+b = b+a$

Under Multiplication (\circ)

- 1. Closure: $a \circ b \in \mathbb{Z}_{11}$
- 2. Associativity: $(ab)c = a(bc)$

3. Distributive laws:

- Left: $a \circ (b+c) = ab + ac$
- Right: $(a+b)c = ac + bc$

Yes, $(\mathbb{Z}_{11}; +; \circ)$ is a ring.

4. Is $\langle \mathbb{Z}_{37}, + \rangle$, $\langle \mathbb{Z}_{35}, \times \rangle$ are abelian groups?

1. $\langle \mathbb{Z}_{37}, + \rangle$

This is the set of integers modulo 37 under addition.

- $a+b \bmod 37 \in \mathbb{Z}_{37}$
- Addition mod 37 is associative.
- 0 is the additive identity.
- Every element $a \in \mathbb{Z}_{37}$ has an inverse $-a \bmod 37$.
- $a+b = b+a \bmod 37$

Yes, $\langle \mathbb{Z}_{37}, + \rangle$ is an abelian group.

2. $\langle \mathbb{Z}_{35}, \times \rangle$

This is the set $\mathbb{Z}_{35} = \{0, 1, \dots, 34\}$ under multiplication mod 35

- 1 is the multiplicative identity.
- $\gcd(5, 35) = 5$, so 5 has no inverse in \mathbb{Z}_{35} .

• closure, associativity, commutativity all hold but not all elements have inverse.

NO, $\langle \mathbb{Z}_{35}, \times \rangle$ is not a group, let alone an abelian one.

• $\langle \mathbb{Z}_{37}, + \rangle$: yes, abelian group.

• $\langle \mathbb{Z}_{35}, \times \rangle$: NO, not a group.

5. Let's take $p=2$ and $n=3$ that makes the GF $(p^n)^n = \text{GF}(2^3)$ then solve this with polynomial arithmetic approach.

Step-1

we want to construct GF(2³), a finite field with 2³=8 elements.

• Elements in GF(2³) are polynomials over GF(2) of degree less than 3.

$$\cdot f_2(0, 1, x, x+1, x^2, x^2+1, x^2+x, x^2+x+1)$$

Step-3:

$$f(x) = x^3 + x + 1$$

Step-3:

$$\cdot (x^2 + x + 1) + (x^2 + 1) = x$$

multiplication Example:

Let,

$$a(x) = x + 1$$

$$b(x) = x^2 + x$$

$$a(x) \cdot b(x) = (x + 1) \cdot (x^2 + x) = x^3 + x^2 + x^2 + x = x^3 + 2x^2 + x$$

Step-2:

$$f(x) = x^3 + x + 1$$

Note:

$$x^3 \equiv x + 1 \pmod{f(x)} \Rightarrow x^3 + x \equiv (x + 1) + x = 1$$

Result:

$$(x+1)(x^2+x) \equiv 1 \pmod{x^3+x+1}$$