

CSE 604

Artificial Intelligence

Chapter 5: Adversarial Search

Adapted from slides available in Russell & Norvig's textbook webpage

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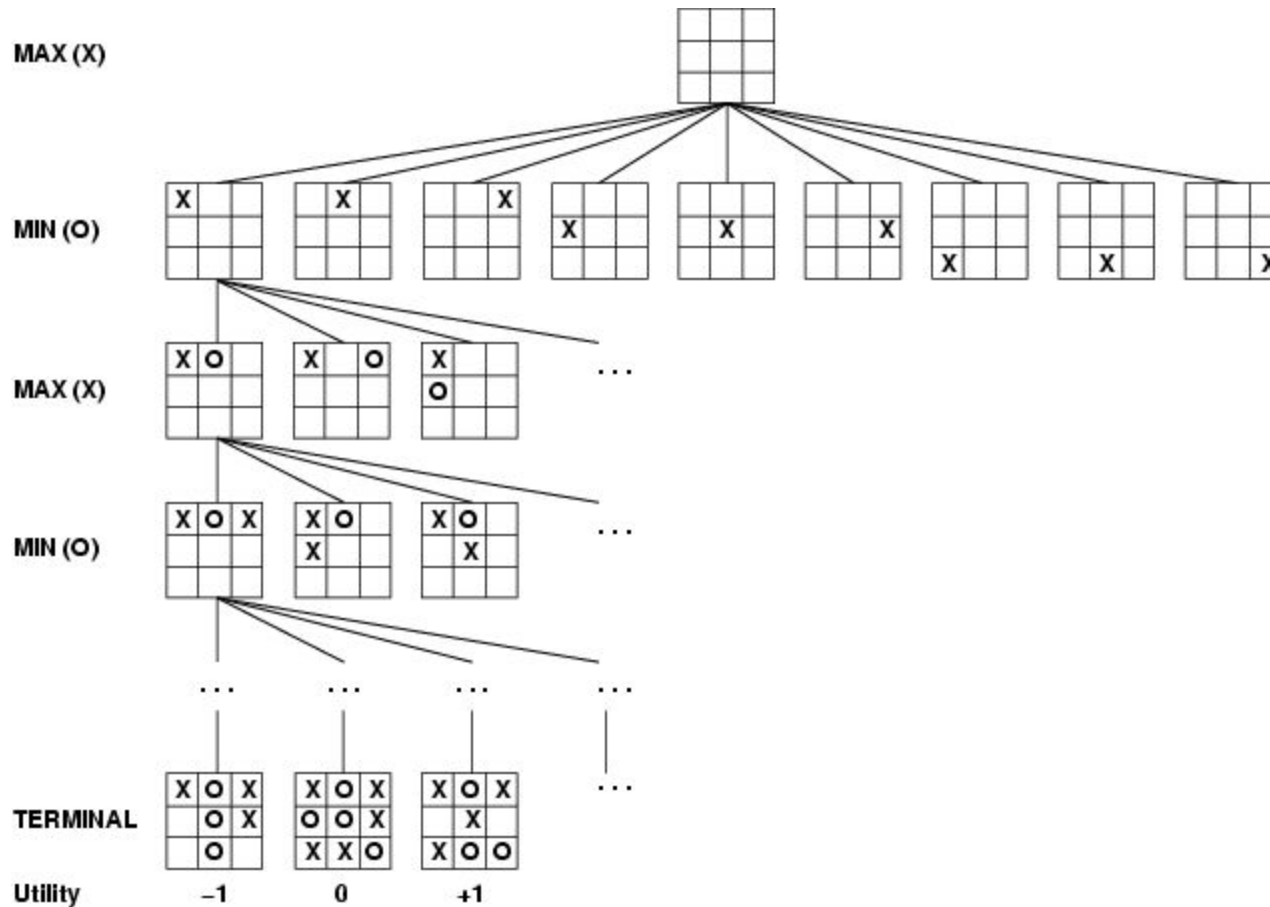
Outline

- Optimal decisions
- α - β pruning
- Imperfect, real-time decisions

Games vs. search problems

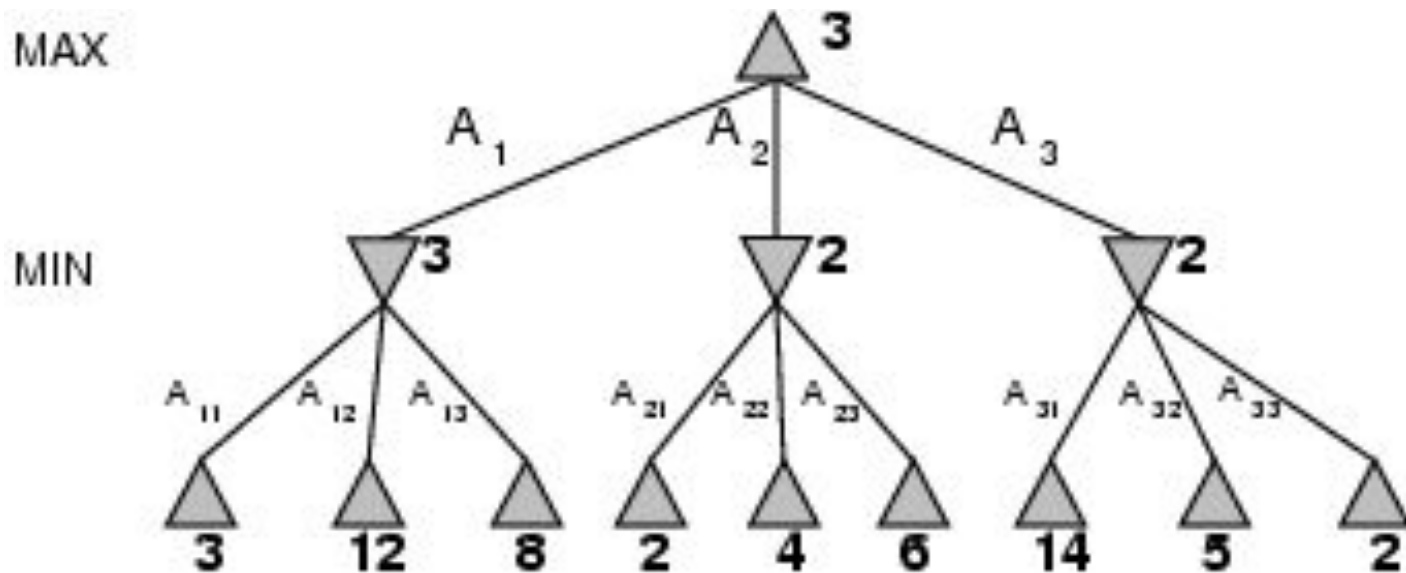
- "Unpredictable" opponent specifying a move for every possible opponent reply
- Time limits unlikely to find goal, must approximate

Game tree (2-player, deterministic, turns)



Minimax

- Perfect play for deterministic games
- Idea: choose move to position with highest **minimax value**
= best achievable payoff against best play
- E.g., 2-ply game:



Minimax algorithm

function MINIMAX-DECISION(*state*) *returns an action*

$v \leftarrow \text{MAX-VALUE}(\textit{state})$

return the *action* in SUCCESSORS(*state*) with value *v*

function MAX-VALUE(*state*) *returns a utility value*

if TERMINAL-TEST(*state*) **then return** UTILITY(*state*)

$v \leftarrow -\infty$

for *a, s* in SUCCESSORS(*state*) **do**

$v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(s))$

return *v*

function MIN-VALUE(*state*) *returns a utility value*

if TERMINAL-TEST(*state*) **then return** UTILITY(*state*)

$v \leftarrow \infty$

for *a, s* in SUCCESSORS(*state*) **do**

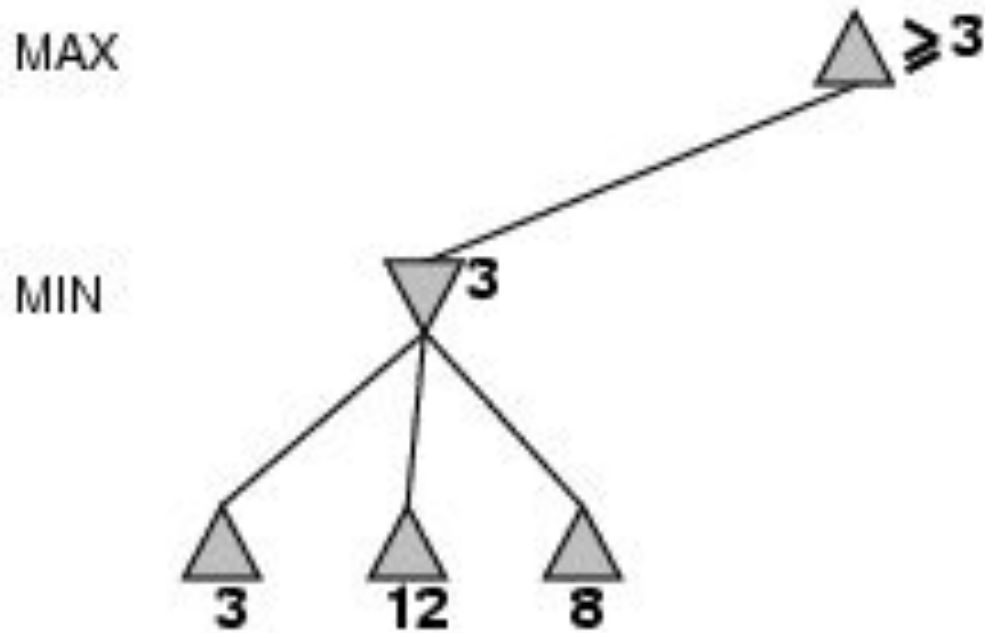
$v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(s))$

return *v*

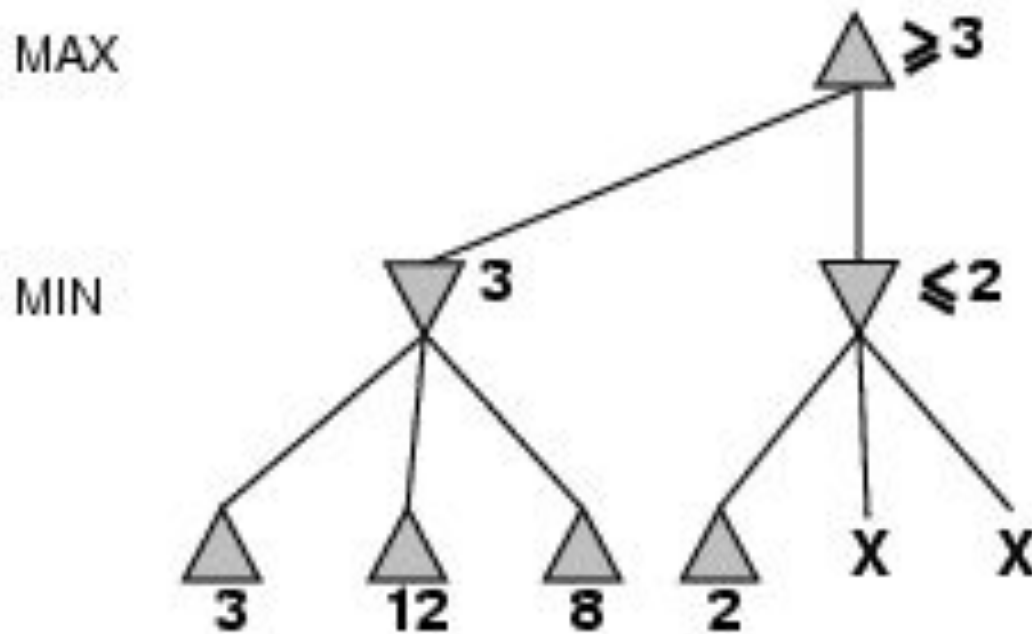
Properties of minimax

- Complete? Yes (if tree is finite)
- Optimal? Yes (against an optimal opponent)
- Time complexity? $O(b^m)$
- Space complexity? $O(bm)$ (depth-first exploration)
- For chess, $b \approx 35$, $m \approx 100$ for "reasonable" games
exact solution completely infeasible

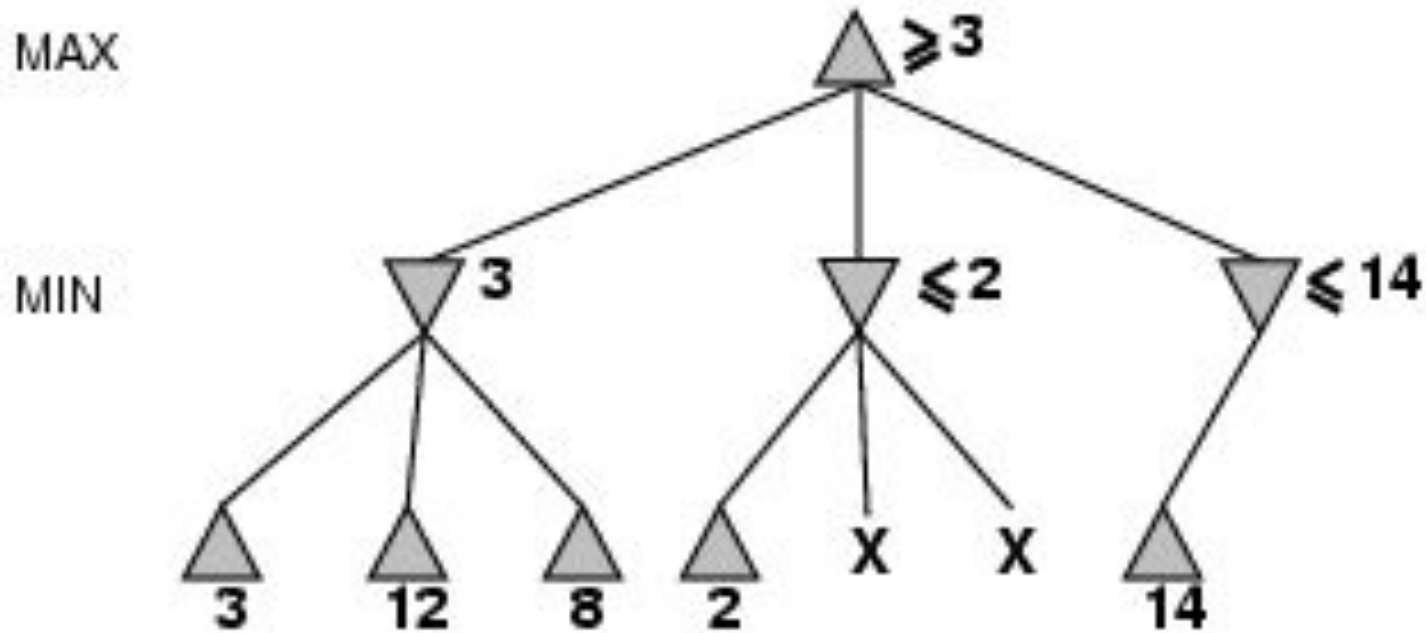
α - β pruning example



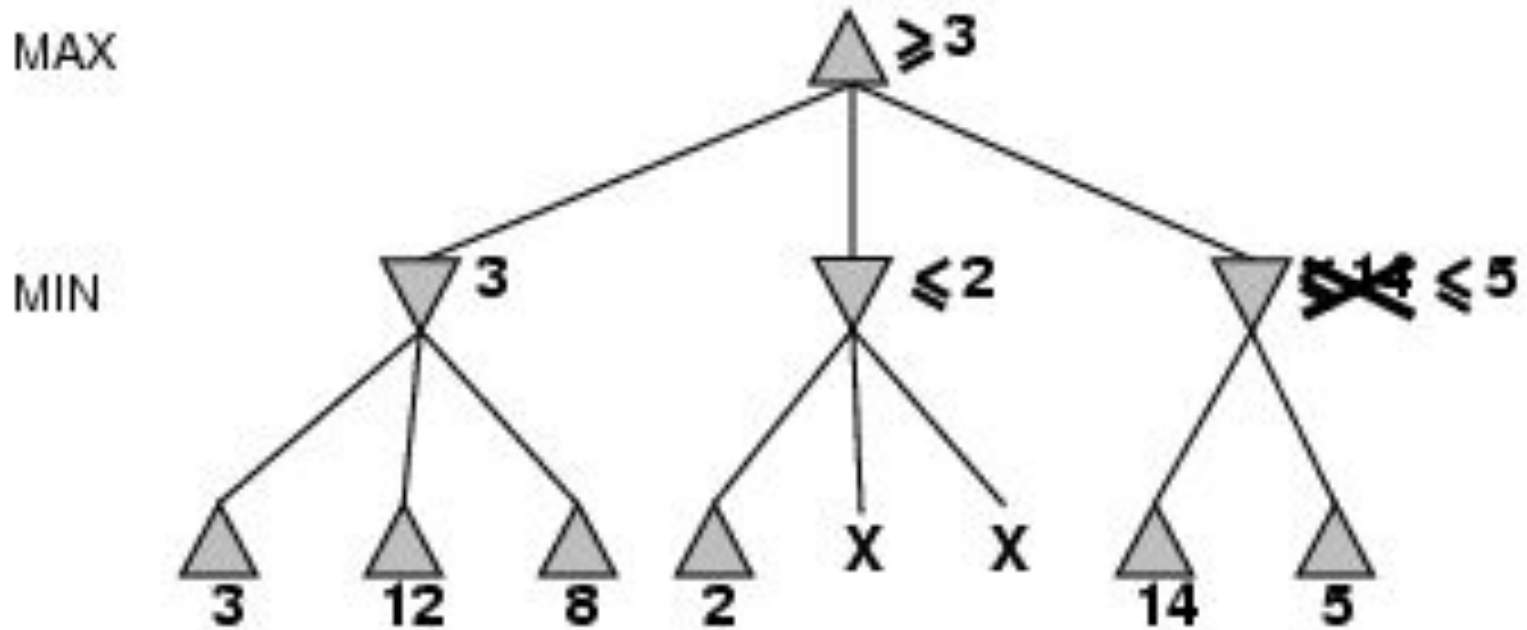
α - β pruning example



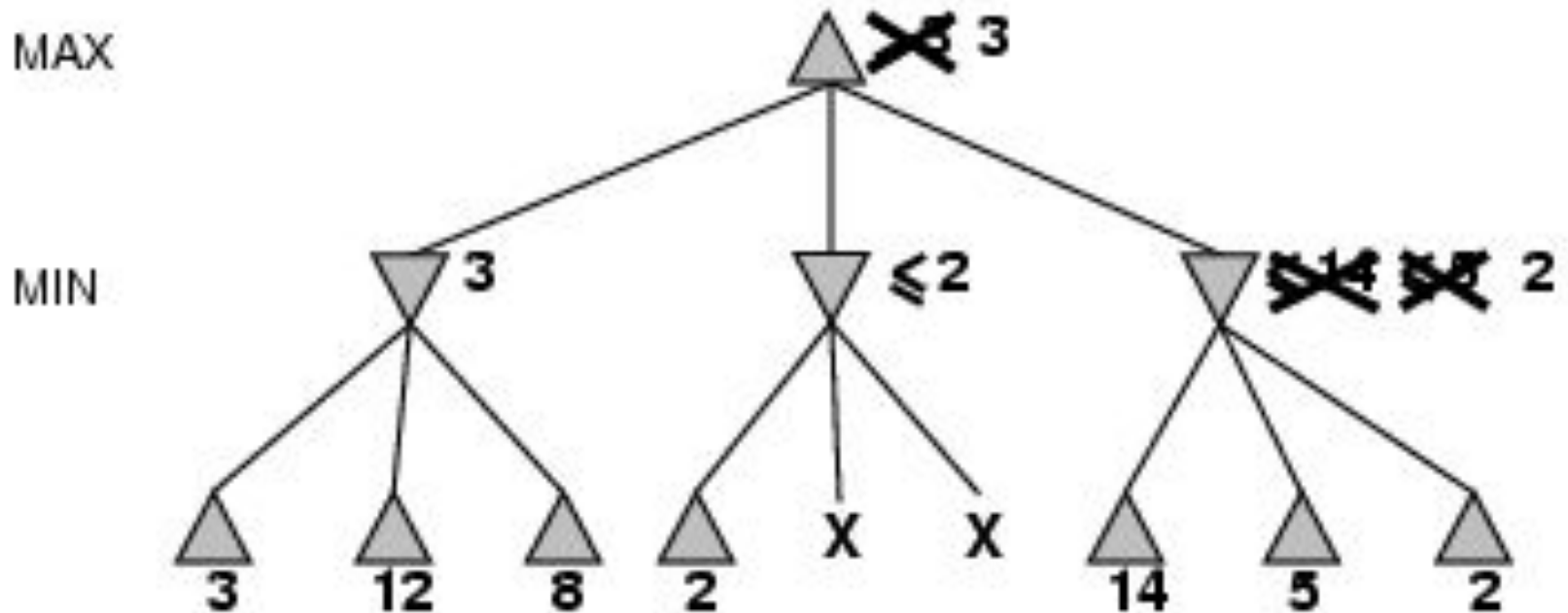
α - β pruning example



α - β pruning example



α - β pruning example

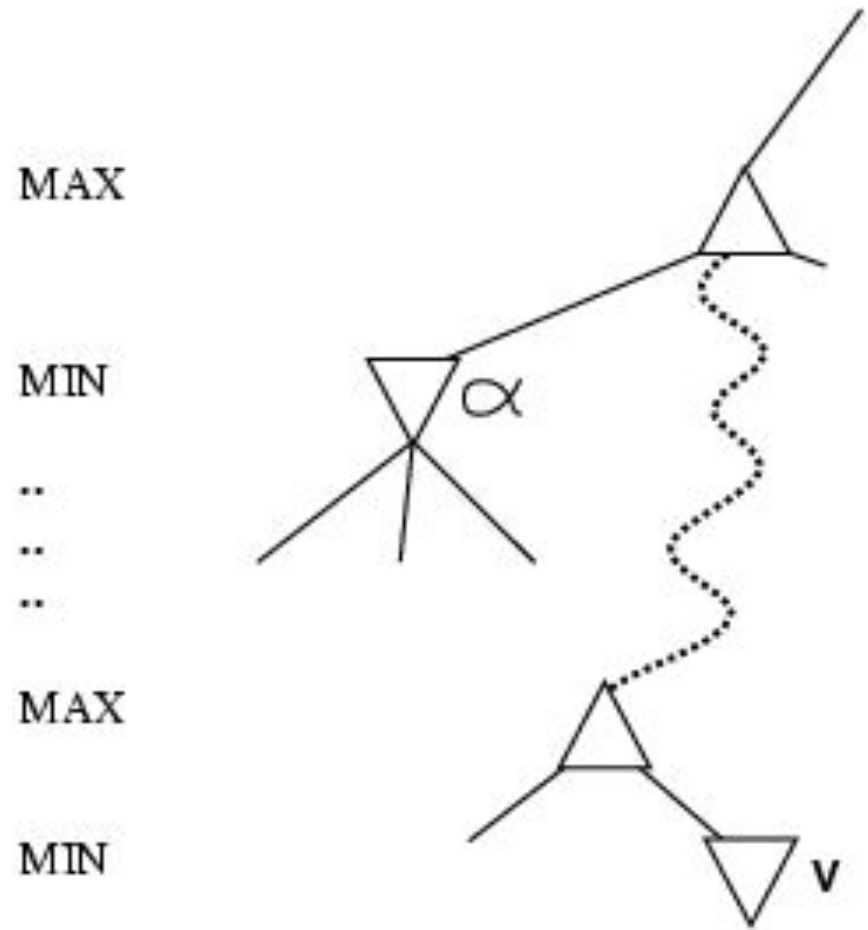


Properties of α - β

- Pruning **does not** affect final result
- Good **move ordering** improves effectiveness of pruning
 - Try the moves that are “likely to be best” first
 - E.g., in chess, try captures, threats, forward moves, backward moves in that order
- With "perfect ordering," time complexity $\approx O(b^{m/2})$
doubles depth of search
- A simple example of the value of reasoning about which computations are relevant (a form of **metareasoning**)

Why is it called α - β ?

- α is the value of the best (i.e., highest-value) choice found so far at any choice point along the path for *max*
- If v is worse than α , *max* will avoid it
prune that branch
- Define β similarly for *min*



The α - β algorithm

function ALPHA-BETA-SEARCH(*state*) *returns an action*

inputs: *state*, current state in game

$v \leftarrow \text{MAX-VALUE}(\text{state}, -\infty, +\infty)$

return the *action* in SUCCESSORS(*state*) with value v

function MAX-VALUE(*state*, α , β) *returns a utility value*

inputs: *state*, current state in game

α , the value of the best alternative for MAX along the path to *state*

β , the value of the best alternative for MIN along the path to *state*

if TERMINAL-TEST(*state*) **then return** UTILITY(*state*)

$v \leftarrow -\infty$

for a, s in SUCCESSORS(*state*) **do**

$v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(s, \alpha, \beta))$

if $v \geq \beta$ **then return** v

$\alpha \leftarrow \text{MAX}(\alpha, v)$

return v

The α - β algorithm

```
function MIN-VALUE(state,  $\alpha$ ,  $\beta$ ) returns a utility value
  inputs: state, current state in game
            $\alpha$ , the value of the best alternative for MAX along the path to state
            $\beta$ , the value of the best alternative for MIN along the path to state

  if TERMINAL-TEST(state) then return UTILITY(state)
   $v \leftarrow +\infty$ 
  for  $a, s$  in SUCCESSORS(state) do
     $v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(s, \alpha, \beta))$ 
    if  $v \leq \alpha$  then return  $v$ 
     $\beta \leftarrow \text{MIN}(\beta, v)$ 
  return  $v$ 
```


Imperfect Real Time Decisions

Even with alpha-beta pruning, it is infeasible to grow the whole game tree!

Standard approach:

- evaluation function
= estimated desirability of position
- cut off search
e.g., depth limit or iterative deepening
- forward pruning
e.g., Beam search

Evaluation functions

- For chess, typically linear weighted sum of **features**

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$$

- e.g., $w_1 = 9$ with

$f_1(s) = (\text{number of white queens}) - (\text{number of black queens}),$

$w_2 = 5$ with

$f_2(s) = (\text{number of white rooks}) - (\text{number of black rooks}),$

etc.

Cutting off search

- We can use a modified algorithm *MinimaxCutoff*
- *MinimaxCutoff* is identical to *MinimaxValue* except
 - *Terminal?* is replaced by *Cutoff?*
 - *Utility* is replaced by *Eval*
- Does it work in practice?
 - Suppose we have 100 secs, explore 10^4 nodes/sec
 10^6 nodes per move
 - $b^m = 10^6, b=35 \quad m=4$
- 4-ply lookahead is a hopeless chess player!
 - 4-ply \approx human novice
 - 8-ply \approx typical PC, human master
 - 12-ply \approx Deep Blue, Kasparov

Deterministic games in practice

- **Checkers:** Chinook ended 40-year-reign of human world champion Marion Tinsley in 1994.
- **Chess:** Deep Blue defeated human world champion Garry Kasparov in a six-game match in 1997. Deep Blue searches 200 million positions per second, uses very sophisticated evaluation, and undisclosed methods for extending some lines of search up to 40 ply.
- **Othello:** human champions refuse to compete against computers, who are too good.
- **Go:** Until recently, human champions refused to compete against computers, who were too bad (in Go, $b > 300$).
 - But in 2016, Google's AlphaGo defeated human world champion Lee Sedol.
 - In 2017, AlphaGo Zero defeated the previous version of AlphaGo 100-0