



Report on

“Mini Compiler for If-Else and While Constructs in Python”

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INTRODUCTION

This Mini Compiler is Built for Python and handles the If-Elif-Else and the While Constructs.

Sample Input

#Start Program

```
import helloWorld
x=10
y=x+20
z=x*10
str = "No."
```

```
if(z==x):
    print(z)
else:
    print("No.")
```

#End Program

Sample Output

```
-----Token Sequence-----
1 T_NL
2 T_NL
3 T_IMPT T_helloWorld T_NL
4 T_x T_EQL T_10 T_NL
5 T_y T_EQL T_x T_PL T_20 T_NL
6 T_z T_EQL T_x T_ML T_10 T_NL
7 T_str T_EQL T_"No." T_NL
8 T_NL
9 T_If T_OP T_z T_EQ T_x T_CP T_Cln T_NL
10 T_ID T_Print T_OP T_z T_CP T_NL
11 T_Else T_Cln T_NL
12 T_ID T_Print T_OP T_"No." T_CP T_NL
13 T_NL
14 T_EOF
Valid Python Syntax
-----
```

-----Abstract Syntax Tree-----

```
NewLine(2)
import(1) NewLine(2)
hWorld =(2) NewLine(2)
  x 10 =(2) NewLine(2)
    y +(2) =(2) NewLine(2)
      x 20 z *(2) =(2) If(3)
        x 10 str "No." ==(2) BeginBlock(2) Else(1)
          z x Print(1) EndBlock BeginBlock(2)
            z Print(1) EndBlock
              "No."
```

-----Three Address Code-----

```
import helloWorld
T2 = 10
x = T2
T5 = x
T6 = 20
T7 = T5 + T6
y = T7
T10 = x
T11 = 10
T12 = T10 * T11
z = T12
T15 = z
T16 = - T15
k = T16
T19 = "No."
str = T19
T22 = z
T23 = x
T24 = T22 == T23
If False T24 goto L0
T25 = z
Print T25
goto L1
L0: T29 = "No."
Print T29
L1:
```

-----All Quads-----				
0	import	helloWorld	-	-
1	=	10	-	T2
2	=	T2	-	x
3	=	x	-	T5
4	=	20	-	T6
5	+	T5	T6	T7
6	=	T7	-	y
7	=	x	-	T10
8	=	10	-	T11
9	*	T10	T11	T12
10	=	T12	-	z
11	=	z	-	T15
12	-	T15	-	T16
13	=	T16	-	k
14	=	"No."	-	T19
15	=	T19	-	str
16	=	z	-	T22
17	=	x	-	T23
18	==	T22	T23	T24
19	If False	T24	-	L0
20	=	z	-	T25
21	Print	T26	-	-
22	goto	-	-	L1
23	Label	-	-	L0
24	=	"No."	-	T29
25	Print	T30	-	-
26	Label	-	-	L1

After Dead Code Elimination

-----All Quads-----				
0	import	hWorld	-	-
1	=	10	-	T2
2	=	T2	-	x
7	=	x	-	T10
8	=	10	-	T11
9	*	T10	T11	T12
10	=	T12	-	z
16	=	z	-	T22
17	=	x	-	T23
18	==	T22	T23	T24
19	If False	T24	-	L0
21	Print	T26	-	-

22	goto	-	-	L1
23	Label	-	-	L0
25	Print	T30	-	-
26	Label	-	-	L1

-----All Symbol Tables-----

Scope Name	Type	Declaration	Last Used Line
(0, 1) helloWorld	PackageName	3	3
(0, 1) 10	Constant	4	6
(0, 1) x	Identifier	4	10
(0, 1) 20	Constant	5	5
(0, 1) y	Identifier	5	5
(0, 1) z	Identifier	6	11
(0, 1) k	Identifier	7	7
(0, 1) "No."	Constant	8	8
(0, 1) str	Identifier	8	8
(0, 1) T2	ICGTempVar	-1	-1
(0, 1) T5	ICGTempVar	-1	-1
(0, 1) T6	ICGTempVar	-1	-1
(0, 1) T7	ICGTempVar	-1	-1
(0, 1) T10	ICGTempVar	-1	-1
(0, 1) T11	ICGTempVar	-1	-1
(0, 1) T12	ICGTempVar	-1	-1
(0, 1) T15	ICGTempVar	-1	-1
(0, 1) T16	ICGTempVar	-1	-1
(0, 1) T19	ICGTempVar	-1	-1
(0, 1) T22	ICGTempVar	-1	-1
(0, 1) T23	ICGTempVar	-1	-1
(0, 1) T24	ICGTempVar	-1	-1
(0, 1) T25	ICGTempVar	-1	-1
(0, 1) T26	ICGTempVar	-1	-1
(0, 1) L0	ICGTempLabel	-1	-1
(0, 1) T29	ICGTempVar	-1	-1
(0, 1) T30	ICGTempVar	-1	-1
(0, 1) L1	ICGTempLabel	-1	-1
(1, 4) "No."	Constant	13	13

ARCHITECTURE OF THE LANGUAGE

Python has a very flexible syntax and we have tried to incorporate as much as possible from our experience of using python into the grammar. But we have not taken care of semicolons, so a semicolon will result in an error while parsing. All lines of code terminate upon seeing a newline character. We have taken care of the following:

- If-Elif-Else and While constructs
- Print Statements
- pass, break and void returns
- Function definitions and Calls
- Lists
- All arithmetic operators and all boolean operators except standalone '!' ('!=' is taken care of)
- Single Line Comments (#)

Semantically we have checked the following :

- Whether any variable used on the RHS is defined and in the current scope or any Enclosing Scope of the current scope.
- Whether a variable being indexed is List
- Whether all expressions in the If and While Clauses are Boolean expressions

LITERATURE SURVEY

1. Lex and Yacc Doc by Tom Niemann
2. Official Bison Documentation : <https://www.gnu.org/software/bison/manual/>
3. Stackoverflow : <https://stackoverflow.com/>

THE CONTEXT FREE GRAMMAR

Grammar to Parse While And If Statements in Python

Grammar	Expression	Legend
StartDebugger	StartParse T_EndOfFile	T_Import : "import"
constant	T_Number T_String	T_Print : "print"
T_ID	[_a-zA-Z][_a-zA-Z0-9]*	T_Pass : "pass"
term	T_ID constant list_index	T_If : "if"
list_index	T_ID T_OB constant T_CB	T_While : "while"
StartParse	T_NL StartParse finalStatements T_NL StartParse finalStatements T_NL	T_Break : "break"
basic_smt	pass_smt break_smt import_smt assign_smt bool_exp arith_exp print_smt return_smt	T_And : "and"
arith_exp	term arith_exp T_PL arith_exp arith_exp T_MN arith_exp arith_exp T_ML arith_exp arith_exp T_DV arith_exp T_OP arith_exp T_CP	T_Or : "or"
ROP	T_GT T_LT T_ELT T_EGT	T_Not : "not"
bool_exp	bool_term bool_term T_Or bool_term arith_exp T_LT arith_exp bool_term T_And bool_term arith_exp T_GT arith_exp arith_exp T_ELT arith_exp arith_exp T_EGT arith_exp arith_exp T_In T_ID	T_elif : "elif"
bool_term	bool_factor arith_exp T_EQ arith_exp T_True T_False	T_Else : "else"
bool_factor	T_Not bool_factor T_OP bool_exp T_CP	T_Def : "def"
import_smt	T_Import T_ID	T_In : "in"
assign_smt	T_ID T_EQL arith_exp T_ID T_EQL bool_exp T_ID T_EQL func_call T_ID T_EQL T_OB T_CB	T_Cln : ":"
pass_smt	T_Pass	T_GT : ">"
break_smt	T_Break	T_LT : "<"
return_smt	T_Return	T_EGT : ">="
print_smt	T_Print T_OP constant T_CP	T_ELT : "<="
finalStatements	basic_smt cmpd_smt func_def func_call	T_EQ : "=="
cmpd_smt	if_smt while_smt	T_NEQ : "!="
if_smt	T_If bool_exp T_Cln start_suite T_If bool_exp T_Cln start_suite elif_stmts	T_True : "True"
elif_stmts	T_elif bool_exp T_Cln start_suite elif_stmts else_smt	T_False : "False"
else_smt	T_Else T_Cln start_suite	T_PL : "+"
while_smt	T_While bool_exp T_Cln start_suite	T_MN : "-"
start_suite	basic_smt T_NL ID finalStatements suite	T_ML : "***"
suite	T_NL ND finalStatements suite T_NL end_suite	T_DV : "/"
end_suite	DD finalStatements DD	T_OP : "("
args	T_ID args_list	T_CP : ")"
args_list	T_Comma T_ID args_list	T_OB : "["
call_list	T_Comma term call_list	T_CB : "]"
call_args	term call_list	T_Comma : ","
func_def	T_Def T_ID T_OP args T_CP T_Cln start_suite	T_EQL : "=="
func_call	T_ID T_OP call_args T_CP	T_NL : "\n"
		T_String : String
		T_Number : Number
		ND : Nodent
		ID : Indent
		DD : Dedent
		T_EOF : End Of File

DESIGN STRATEGY

The Design Strategy we have implemented is that firstly, every links back to the Symbol Table. This means that the required nodes of the Abstract Syntax tree and a the required Quadruples in the Intermediate code have a link to the Symbol table. All the compiler generated Temporaries are also stored in the symbol table.

The Symbol Table stores 'Records' having 4 columns as you can see in the sample output, ie,

- ❖ Scope : Scope of each variable contained in record
- ❖ Name : Value/Name of each variable contained in record
- ❖ Type : Type of each variable contained in record
 - PackageName
 - Func_Name
 - Identifier
 - Constant
 - ListTypeID
 - ICGTempVar
 - ICGTempLabel
- ❖ Declaration : Line of Declaration of each variable contained in record
- ❖ Last Use Line

The scope is a function of Indentation depth and to make it unique we have a tuple of the scope of the parent and the current scope calculated using the Indentation depth.

For the Abstract Syntax Tree, We have 2 Types of Nodes, Leaf nodes and Internal nodes. The nodes can have variable number of children (0-3) depending upon the construct it represents. Take the example of the If-Else Statement,

```

                If
      Condition  CodeBlock  Else
```

To display the AST, We take the AST and store it as a matrix of levels. As we can see in the sample output, we have printed each level of the AST. All Internal nodes also have a number enclosed in brackets next to them, which represents the number of children they have in the next level. Leaf nodes in the AST representing identifiers, constants, Lists, packages point to a record in the symbol table.

The Intermediate code is generated by recursively stepping through the AST. Each line is stored as a quadruple (Operation, Arg1, Arg2, Result) so that it we may easily optimize code in the following steps.

We Eliminate dead code, specifically unused variables in the whole program. For example, if we have the following lines of code:

```
a = 10
b = 10
c = a + b
```

And these 3 variables are not used on any other RHS, then these 3 lines of code are Eliminated during optimization. All the dead variables in the code are removed. We iterate through the Quads to do this.

To take care of the Indentation based code structure and scoping we have implemented a stack and we use 3 tokens. The top of the stack always points to the current indentation value, if that value on scanning the next line, doesn't change, it implies that we're in the same scope and hence, we return the token 'ND', i.e, 'No-dent'. If the value increases, it means we are entering a sub-scope and we return the token 'ID' , i.e. 'Indent' and finally, if the value decreases, it means we're returning to one of the enclosing scope and we return 'DD', i.e. 'Dedent'.

For Error Handling, Whenever the parser encounters an error, It prints the Line no and column number of the error. We also display the following errors:

- Identifier <var> Not declared in scope
- Identifier <var> Not Indexable

All Comments are removed from the code before parsing

IMPLEMENTATION DETAILS

The Symbol table uses two Structures,

```
typedef struct record
{
    char *type;
    char *name;
    int declLineNo;
    int lastUseLine;
} record;

typedef struct STable
{
    int no;
    int noOfElems;
    int scope;
    record *Elements;
    int Parent;

} STable;
```

The “record” structure represents each record in the symbol table. Each symbol table can contain a maximum of “MAXRECS” records, MAXRECS is a macro.

The “STable” structure represents one symbol table. A new Symbol table is made for every scope. A Maximum of “MAXST” symbol tables and hence scopes can exist, MAXST is a macro.

The Abstract Syntax Tree uses one structure,

```
typedef struct ASTNode
{
    int nodeNo;
    /*Operator*/
    char *NType;
    int noOps;
    struct ASTNode** NextLevel;

    /*Identifier or Const*/
    record *id;
} node;
```

This ASTNode structure takes care of both leaf nodes as well as Internal

“Operator” Nodes. The respective values are set depending upon the type of node. Each node can have 0-3 children. We print the AST by first storing it in a Matrix of Order “MAXLEVELS” x “MAXCHILDREN” and printing the matrix Levelwise. This Matrix, is a matrix of pointers to the AST. The “noOps” element of the Node gives the number of children of that node.

The Three-Address Code is represented and stored as Quads that are given by the Structure,

```
typedef struct Quad
{
    char *R;
    char *A1;
    char *A2;
    char *Op;
    int I;
} Quad;
```

The last element, the integer ‘I’, is used during code optimization. All the Three-Address codes are stored as Quads in an array of Quads. There can be a maximum of “MAXQUADS” quadruples.

For the dead code elimination we continuously loop through the code till no more code can be eliminated. To check if a quadruple represents dead code we see if the result parameter/element of that quad appears as any arguments to any other subsequent quads which have not been eliminated, If not, we consider that quad to represent dead code and mark the element ‘I’ with the value “-1”.

The scope checking is handled by recursively stepping through enclosing scopes and finding the most recent definition of the variable. If no definition is found, we print the error.

Lastly, To compile the code and execute the code we have provided a makefile.

If you wish to only see the AST,

```
lex grammar.l
yacc -dv grammarAST.y
gcc lex.yy.c y.tab.c -g -ll -o TestAST.out
./TestAST.out < InputFile.txt
```

If you wish to only see the Intermediate code and Optimization

```
lex grammar.l
```

```
yacc -dv grammarICG.y
gcc lex.yy.c y.tab.c -g -ll -o TestICG.out
./TestICG.out < InputFile.txt
```

The makefile provided compiles the “grammar.y” file and prints everything.
./Test.out < Input.txt

RESULTS AND SHORTCOMINGS

The result achieved is that we have a mini compiler which parses grammar corresponding to basic python syntax and generates finally, an optimized intermediate representation.

The areas where our mini compiler falls short are,

- Doesn't handle semi colon.
- Only one very basic optimization is implemented that does not reduce code density by a huge margin.
- The Program has a few memory leaks, although most of them have been taken care of.

SNAPSHOTS

Input File with Comments

```
import hWorld
x=10
y=10
#Comment1
x+y
listX = []
def F1(A, B, C):
    while(listX[2]==y):
        c=0
        z=10
        b=z
        if(z==b):
            c=10+b
        else:
            c=10+c
        w=21
#Comment2
m = F1(10, 10, 10)
if(x==y):
    x=10
else:
    x=10
if(x==y):
    x=10
else:
    y=10
```

Token Sequence

```

-----Token Sequence-----
1 T IMPT T hWorld T NL
2 T x T EQL T 10 T NL
3 T y T EQL T 10 T NL
4 T NL
5 T NL
6 T x T PL T y T NL
7 T listX T EQL T OB T CB T NL
8 T NL
9 T Def T F1 T OP T A T Comma T B T Comma T C T CP T Cln T NL
10 T ID T While T OP T listX T OB T 2 T CB T EQ T y T CP T Cln T NL
11 T ID T c T EQL T 0 T NL
12 T ND T z T EQL T 10 T NL
13 T ND T b T EQL T z T NL
14 T ND T If T OP T z T EQ T b T CP T Cln T NL
15 T ID T c T EQL T 10 T PL T b T NL
16 T DD T Else T Cln T NL
17 T ID T c T EQL T 10 T PL T c T NL
18 T DD T w T EQL T 21 T NL
19 T DD T NL
20 T NL
21 T m T EQL T F1 T OP T 10 T Comma T 10 T Comma T 10 T CP T NL
22 T If T OP T x T EQ T y T CP T Cln T NL
23 T ID T x T EQL T 10 T NL
24 T Else T Cln T NL
25 T ID T x T EQL T 10 T NL
26 T NL
27 T If T OP T x T EQ T y T CP T Cln T NL
28 T ID T x T EQL T 10 T NL
29 T Else T Cln T NL
30 T ID T y T EQL T 10 T NL
31 T NL
32 T EOF
Valid Python Syntax

```

Abstract Syntax Tree

```

-----Abstract Syntax Tree-----
NewLine(2)BD
import(1) NewLine(2)
hWorld = (2) NewLine(2)
  x 10 = (2) NewLine(2)
    y 10 + (2) NewLine(2)
      x y listX NewLine(2)
        Func Name(3) NewLine(2)
          F1 A, B, C BeginBlock(2) = (2) NewLine(2)
            While(2) EndBlock m Func Call(2) If(3) If(3)
              == (2) BeginBlock(2) F1 10, 10, 10 == (2) BeginBlock(2) Else(1) == (2) BeginBlock(2) Else(1)
                ListIndex(2) y = (2) Next(2) x y = (2) EndBlock BeginBlock(2) x y = (2) EndBlock BeginBlock(2)
                  listX 2 c 0 = (2) Next(2) x 10 = (2) EndBlock x 10 = (2) EndBlock
                    z 10 = (2) Next(2) x 10 y 10
                      b z If(3) EndBlock
                        == (2) BeginBlock(2) Else(1)
                          z b = (2) EndBlock BeginBlock(2)
                            c + (2) = (2) EndBlock(1)
                              10 b c + (2) = (2)
                                10 c w 21

```

Intermediate Code

```

import hWorld
T2 = 10
x = T2
T5 = 10
y = T5
T8 = x
T9 = y
T10 = T8 + T9
Begin Function F1
T15 = listX[2]
T16 = y
T17 = T15 == T16
L0: If False T17 goto L1
T18 = 0
c = T18
T21 = 10
z = T21
T24 = z
b = T24
T27 = z
T28 = b
T29 = T27 == T28
If False T29 goto L2
T30 = 10
T31 = b
T32 = T30 + T31
c = T32
goto L3
L2: T37 = 10
T38 = c
T39 = T37 + T38
c = T39
T42 = 21
w = T42
L3: goto L0
L1: End Function F1
Push Param 10
Push Param 10
Push Param 10
(T63)Call Function F1, 3
Pop Params for Function F1, 3
m = T63
T66 = x
T67 = y
T68 = T66 == T67
If False T68 goto L6
T69 = 10
x = T69
goto L7
L6: T74 = 10
x = T74
L7: T81 = x
T82 = y
T83 = T81 == T82
If False T83 goto L8
T84 = 10
x = T84
goto L9
L8: T89 = 10
y = T89

```

```

-----All Quads-----
0      import hWorld - -
1      = 10 - T2
2      = T2 - x
3      = 10 - T5
4      = T5 - y
5      = x - T8
6      = y - T9
7      + T8 T9 T10
8      BeginF F1 - -
9      =[illegal]listX 2 T15
10     = y - T16
11     == T15 T16 T17
12     Label - - L0
13     If False T17 - L1
14     = 0 - T18
15     = T18 - c
16     = 10 - T21
17     = T21 - z
18     = z - T24
19     = T24 - b
20     = z - T27
21     = b - T28
22     == T27 T28 T29
23     If False T29 - L2
24     = 10 - T30
25     = b - T31
26     + T30 T31 T32
27     = T32 - c
28     goto - - L3
29     Label - - L2
30     = 10 - T37
31     = c - T38
32     + T37 T38 T39
33     = T39 - c
34     = 21 - T42
35     = T42 - w
36     Label - - L3
37     goto - - L0
38     Label - - L1
39     EndF F1 - -
40     Param 10 - -
41     Param 10 - -
42     Param 10 - -
43     Call F1 3 T63
44     = T63 - m
45     = x - T66
46     = y - T67
47     == T66 T67 T68
48     If False T68 - L6
49     = 10 - T69
50     = T69 - x
51     goto - - L7
52     Label - - L6
53     = 10 - T74
54     = T74 - x
55     Label - - L7
56     = x - T81
57     = y - T82
58     == T81 T82 T83
59     If False T83 - L8
60     = 10 - T84
61     = T84 - x
62     goto - - L9
63     Label - - L8
64     = 10 - T89
65     = T89 - y
66     Label - - L9

```

Optimized Code (Quads Removed)

```

-----Code-----All Quads-----
0      import  hWorld  -      -
1      =      10      -      T2
2      =      T2      -      x
3      =      10      -      T5
4      =      T5      -      y
8      BeginFg F1      -      -
9      =[] listX      2      T15
10     =      y      -      T16
11     ==     T15      T16      T17
12     Label -      -      L0
13     If False      T17      -      L1
14     =      0      -      T18
15     = CORI T18      -      c
16     =      10      -      T21
17     =      T21      -      z
18     =      z      -      T24
19     =      T24      -      b
20     =      z      -      T27
21     =      b      -      T28
22     == CCBD T27      T28      T29
23     If False      T29      -      L2
24     =      10      -      T30
25     =      b      -      T31
26     +      T30      T31      T32
27     =      T32      -      c
28     goto -      -      L3
29     Label -      -      L2
30     = Move 10      -      T37
31     =      c      -      T38
32     +      T37      T38      T39
33     =      T39      -      c
36     Label -      -      L3
37     goto -      -      L0
38     Label -      -      L1
39     EndFg F1      -      -
40     Param  10      -      -
41     Param  10      -      -
42     Param  10      -      -
43     Call  F1      3      T63
45     =      x      -      T66
46     =      y      -      T67
47     == Kaldi Patch T66      T67      T68
48     If False      T68      -      L6
49     =      10      -      T69
50     =      T69      -      x
51     goto -      -      L7
52     Label -      -      L6
53     =      10      -      T74
54     =      T74      -      x
55     Label -      -      L7
56     =      x      -      T81
57     =      y      -      T82
58     == png T81      T82      T83
59     If False      T83      -      L8
60     =      10      -      T84
61     =      T84      -      x
62     goto -      -      L9
63     Label -      -      L8
64     =      10      -      T89
65     =      T89      -      y
66     Label -      -      L9
-----

```


All Symbol Tables

-----All Symbol Tables-----				
Scope	Name	Type	Declaration	Last Used Line
(0, 1)	hWorld	PackageName	1	1
(0, 1)	10	Constant	2	3
(0, 1)	x	Identifier	2	27
(0, 1)	y	Identifier	3	27
(0, 1)	listX	ListTypeID	7	10
(0, 1)	F1	Func Name	9	9
(0, 1)	T2	ICGTempVar	-1	-1
(0, 1)	T5	ICGTempVar	-1	-1
(0, 1)	T8	ICGTempVar	-1	-1
(0, 1)	T9	ICGTempVar	-1	-1
(0, 1)	T10	ICGTempVar	-1	-1
(0, 1)	T15	ICGTempVar	-1	-1
(0, 1)	T16	ICGTempVar	-1	-1
(0, 1)	T17	ICGTempVar	-1	-1
(0, 1)	L0	ICGTempLabel	-1	-1
(0, 1)	L1	ICGTempLabel	-1	-1
(0, 1)	T18	ICGTempVar	-1	-1
(0, 1)	T21	ICGTempVar	-1	-1
(0, 1)	T24	ICGTempVar	-1	-1
(0, 1)	T27	ICGTempVar	-1	-1
(0, 1)	T28	ICGTempVar	-1	-1
(0, 1)	T29	ICGTempVar	-1	-1
(0, 1)	L2	ICGTempLabel	-1	-1
(0, 1)	T30	ICGTempVar	-1	-1
(0, 1)	T31	ICGTempVar	-1	-1
(0, 1)	T32	ICGTempVar	-1	-1
(0, 1)	L3	ICGTempLabel	-1	-1
(0, 1)	T37	ICGTempVar	-1	-1
(0, 1)	T38	ICGTempVar	-1	-1
(0, 1)	T39	ICGTempVar	-1	-1
(0, 1)	T42	ICGTempVar	-1	-1
(0, 1)	T63	ICGTempVar	-1	-1
(0, 1)	T66	ICGTempVar	-1	-1
(0, 1)	T67	ICGTempVar	-1	-1
(0, 1)	T68	ICGTempVar	-1	-1
(0, 1)	L6	ICGTempLabel	-1	-1
(0, 1)	T69	ICGTempVar	-1	-1
(0, 1)	L7	ICGTempLabel	-1	-1
(0, 1)	T74	ICGTempVar	-1	-1
(0, 1)	T81	ICGTempVar	-1	-1
(0, 1)	T82	ICGTempVar	-1	-1
(0, 1)	T83	ICGTempVar	-1	-1
(0, 1)	L8	ICGTempLabel	-1	-1
(0, 1)	T84	ICGTempVar	-1	-1
(0, 1)	L9	ICGTempLabel	-1	-1
(0, 1)	T89	ICGTempVar	-1	-1
(0, 2)	2	Constant	10	10
(0, 2)	10	Constant	21	21
(0, 2)	m	Identifier	21	21
(1, 3)	0	Constant	11	11
(1, 3)	c	Identifier	11	11
(1, 3)	10	Constant	12	12
(1, 3)	z	Identifier	12	14
(1, 3)	b	Identifier	13	15
(1, 3)	21	Constant	18	18
(1, 3)	w	Identifier	18	18
(2, 4)	10	Constant	15	15
(2, 4)	c	Identifier	15	17
(2, 16)	10	Constant	17	17
(2, 16)	c	Identifier	17	17
(1, 4)	10	Constant	23	23
(1, 4)	x	Identifier	23	23
(1, 8)	10	Constant	25	25
(1, 8)	x	Identifier	25	25
(1, 16)	10	Constant	28	28
(1, 16)	x	Identifier	28	28
(1, 32)	10	Constant	30	30
(1, 32)	y	Identifier	30	30

CONCLUSION

In Conclusion, A Compiler for Python was implemented. In addition to the constructs specified, the basic python constructs were implemented and function definitions and calls were supported.

The compiler also reports the basic errors and gives the line number and column number.

The Intermediate code was represented by quads which was later optimized to remove dead code.

FURTHER ENHANCEMENTS

The Compiler can be further enhanced by adding,

- Support for 'For' Loops
- Better Memory Management
- More efficient optimization techniques
- Semantic analysis for Parameter Matching
- Error Recovery