Answer 1

| | I and the second | 1 |
|----------|--|---|
| 3(a)(i) | Any three from: A circuit is established at the start of the communication Between sender and receiver This lasts for the duration of the call/data transfer Then the links that make up the circuit are removed | 3 |
| 3(a)(ii) | Any two from: A dedicated channel // Not sharing channel Can use all bandwidth Two-way real time conversation No delay as no switching Data arrives in order it is sent | 2 |
| 3(b)(i) | Any three from: A circuit does not have to be established at the start of the communication The data to be sent is divided into packets That can travel along different routes From node to node Packets are reassembled in the correct order at the receiver's end Must wait until the last packet is received to put the data back together | 3 |
| 3(b)(ii) | Any two from: Communication is asynchronous Allows for error checking Real time transmission is not required Smaller amounts of data are sent (than voice calls) therefore dedicated line/higher bandwidth not required // can share the bandwidth Doesn't matter if data arrives out of order | 2 |

| 3(a)(i) | Any five from A layered model / stack | 5 |
|---------|--|---|
| | with 4 layers | |
| | Uses a set of protocols for transmission of data transport control protocol with internet protocol | |
| | Named layers of Application Layer, Transport Layer, Internet/Network Layer, | |
| | Data Link Layer two layers correct all four layers correct in the correct order | |

| 3(a)(ii) | One mark for protocol, four marks for description | |
|----------|--|--|
| | Protocol: BitTorrent | |
| | Description: Any four from | |
| | BitTorrent client software made available to friends and family's computers | |
| | a complete copy of the torrent/file to be shared is available on at least one computer | |
| | the torrent/file is split into small pieces | |
| | rare pieces are given priority for downloading | |
| | the torrent descriptor file is made available | |
| | a computer/user joins (the swarm) by using the BitTorrent software to load the torrent descriptor file | |
| | the computer/user can now download a piece of the torrent/file | |
| | once a computer has a piece/ the whole file, it can become a seed (and upload to other members of the swarm) | |
| | pieces of the torrent/file are both downloaded and uploaded (by each member of the of the swarm) | |
| | a server called a tracker keeps records of all the computers in the swarm | |
| | the tracker shares their IP addresses allowing connection to each other | |

| 3(b)(i) | One mark for method Method: Packet switching Three marks for description, any three from | 4 |
|----------|---|---|
| | Description: file is divided into packets and necessary data added to each packet, e.g. header which are sent independently of each other and do not need to take the same route packets are reassembled at the destination missing / corrupt packets can be resent | |
| 3(b)(ii) | One mark for benefit, one mark for drawback Benefit: packets can be rerouted if there are problems // packets can take the least congested route // transmission errors can be detected // missing / corrupt packets can be resent | 2 |
| | Drawback: packets can be dropped / delayed | |

| 3(a) | Set of rules For (successful) transmission (and receipt) of data | 2 |
|------|---|---|
| 3(b) | One mark for protocol, one mark for purpose must match protocol Any two pairs from Protocol: POP3 Purpose: downloading email Protocol: SMTP Purpose: sending/transferring email Protocol: IMAP Purpose: downloading email // storing/organising emails on an email server Protocol: HTTP/HTTPS Purpose: accessing email using a browser | 4 |

| 3(c)(i) | One mark for method, max two marks for description Method: Circuit switching Description: Set up for the duration of the conversation Set up before communication starts Maintained throughout the transmission All data travels down the same route Dropped at the end of the transmission Complete bandwidth used | 3 |
|----------|--|---|
| 3(c)(ii) | One mark for benefit, one mark for limitation Benefit: Manav and Miora can see other in real time // better synchronisation // full bandwidth available Drawback: Bandwidth / channel not available to other users // extra time required to set up circuit at start of conversation // alternative route not available without restarting the conversation // less secure as easier to intercept data if only one channel used // failure single route used means failure of transmission | 2 |

| 2(a) | Circuit switching | 1 |
|------|---|---|
| 2(b) | 1 mark Any real-time application e.g. video conferencing // live streaming of a concert | 3 |
| | Justification 1 mark per bullet to max 2 ∞ reduced latency ∞ there are little/no delays in sending/receiving data once the circuit is set up ∞ because (stringent) error checking (as used in packet switching) is not required ∞ circuit made available is dedicated to this communication stream | |

| 3(a) | 1 mark per bullet point to max 2 | 2 |
|------|--|---|
| | Provide a set of standards for transmission of data that gives a known/accepted set of rules for transmitting and receiving data This enables communication/compatibility between devices from different manufacturers/platforms etc. | |

| 7(a) | 1 mark per bullet point | 4 |
|----------|---|---|
| | Application | |
| | Transport | |
| | Internet / Network | |
| | Data Link | |
| | - Data Link | |
| 7(b)(i) | 1 mark per bullet point to max 2 | 2 |
| | Packet switching makes best use of the available (channel) capacity | |
| | by using alternative routes | |
| | which is more secure / robust | |
| | as packets to / from different sources and destinations can share the | |
| | same route | |
| 7(b)(ii) | 1 mark per bullet point to max 2 | 2 |
| | | |
| | To store data about packet | |
| | and its routing // to ensure it reaches its destination | |
| | to ensure that message can be properly reconstructed | |
| 7(b)(ii) | 1 mark per item to max 3 | 3 |
| | For example: | |
| | IP address of sender | |
| | IP address of destination | |
| | IP version | |
| | Number of packets the message consists of | |
| | ID number of that packet | |
| | Protocol used | |
| | Packet length | |
| | Time to live // max number of hops | |
| | Synchronisation data | |
| | Source port | |
| | Destination Port | |
| | Checksum | |
| | | |

| 3(b) | 1 n | nark for each correct pair of letters in the right order max 3 | 3 |
|------|-----|--|---|
| | 1 | Computer X sends a connection request to Computer Y. | |
| | 2 | Computer Y sends ready or busy signal. | |
| | 3 | If busy, Computer X waits and then resends the connection request to Computer Y. | |
| | 4 | D | |
| | 5 | A | |
| | 6 | С | |
| | 7 | В | |

| 3(c)(i) | 1 mark for each la | ayer | | | |
|----------|--------------------|-----------------------|----------------------------|----------------------|--|
| | | Protocol | Layer | | |
| | | ТСР | Transport | | |
| | | IP | Internet/Network | | |
| | | SMTP | Application | | |
| 3(c)(ii) | Any three points | from: | | | |
| | ∞ BitTorrer | nt client software m | ade available | | |
| | | nputer must keep a | complete copy of the t | orrent/file to be | |
| | shared | ile is split into sma | Il niocos | | |
| | | |) by using the BitTorren | t software to load | |
| | | descriptor file | , by doing the Dictorion | i continui o to iouu | |
| | | • | vnload a piece of the file | е | |
| | | | ce it can become a see | ed and upload (to | |
| | | embers of the swar | , | | |
| | | | th downloaded and upl | oaded (by each | |
| | | of the of the swarr | eps records of all the c | omputers in the | |
| | swarm | danied a tracker ke | opo records of all the c | | |
| | | | addresses allowing the | m to connect to | |

| 4(a) | 1 mark per correct row | | | 5 |
|------|---|-----------------|---------|---|
| | Responsibility | ТСР | IP | |
| | Correct routing | | ✓ | |
| | Host to host communication | ✓ | | |
| | Communication between networks | | ✓ | |
| | Retransmitting missing packets | ✓ | | |
| | Reassembling packets into the correct order | ✓ | | |
| 4(b) | 1 mark for name, 1 mark for matching use, max 4 for 2 | 2 protocols | | 4 |
| | POP3/IMAP (1) receiving emails // download email SMTP (1) sending emails (1) | ils from a serv | er (1) | |
| | FTP (1) allows files to be transferred from one cor | • | ner (1) | |
| | HTTP/HTTPS (1) transfer of web pages/hypertext Bit Torrent (1) used for peer-to-peer file sharing (1 | ` ' | | |
| 4(c) | Internet / Network (layer) | | | 1 |

| 4(d)(i) | 1 mark per bullet to max 4 | 4 |
|-----------|--|---|
| | Message data / payload IP version number Internet header length Type of service Explicit Congestion Notification Total length/size of packet (in bytes) Identification / sequence / packet number Fragmentation flags Fragmentation offset Time to live // number of hops Protocol (Header) checksum Source (IP) address Destination (IR) address | |
| 4(d)(ii) | Destination (IP) address 1 mark per benefit, 1 mark per expansion, max 4 for 2 benefits | 4 |
| | For example : alternative route available in case of network problem If packet fails to arrive then only that packet has to be resent | |
| 4(d)(iii) | 1 mark per bullet to max 2 Network ID // (IP Address of) network destination (1) Subnetmask (1) Routing metric // data to decide best route (IP Addresses of possible) next hop / Gateway Interface | 2 |

| 3(b)(i) | 1 mark per bullet point to max 3: | 3 |
|----------|---|---|
| | allows applications to exchange data establishes and maintains a connection until exchange of data is complete determines how to break application data into packets adds sequence / packet number to (TCP) header sends packets to and accepts packets from the network / Internet layer manages flow control // manages congestion avoidance acknowledges all packets that arrive detects when a packet has not arrived at destination handles retransmission of dropped packets reassembles packets into the correct order | |
| 3(b)(ii) | 1 mark per bullet point to max 2 | 2 |
| | encapsulates data into datagram passes datagram to the network access layer (for transmission on the LAN)// passes datagram to the transport layer (on arrival at destination) Defines the addressing method e.g. subnetting, NAT | |

| | 1 | | |
|-----------|---|---|---|
| 3(b)(iii) | HTTP(S) // FTP // POP3 // SMTP // UDP // etc | | 1 |
| 3(c) | 1 mark for appropriate pr | rotocol in each layer | 3 |
| | Layer | Protocol | |
| | Application | HTTP(S) // FTP // POP3 // SMTP // UDP etc | |
| | Transport | TCP | |
| | Internet | IP | |
| 3(d) | ∑ Packets are forward ∑ Packets may take di ∑ Missing packets are | p packets ed size n a header on IP address, sequence number etc. ed from one LAN to the other LAN | 4 |

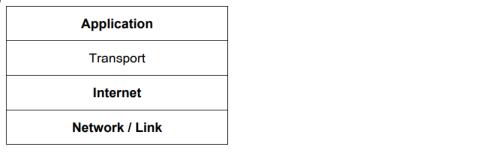
| 2(b)(i) | connection estwhich lasts for | oint to max 2: uit/channel/(physical) path tablished before/at the start of the duration of connection // circuit r smitted along the same route | | ommunication |
|----------|--|--|-------------------|------------------|
| 2(b)(ii) | 1 mark for each row | : | | |
| | | Statements | Circuit switching | Packet switching |
| | | Shares bandwidth | | ✓ |
| | | Data may arrive out of order | | ✓ |
| | | Data can be corrupted | ✓ | ✓ |
| | | Data are less likely to get lost | either√ | or√ |

| 5(a)(i) | Packet: Both web page and web page request are split into packets Each packet is sent individually from device to device | 1 | 2 |
|-----------|---|---|-------|
| 5(a)(ii) | Router: Transmit packets Contain connections to many other routers When packets arrive at router, router decides where next to send packet 1 mark for any valid point | | Max 2 |
| 5(a)(iii) | TCP/IP: Is the protocol Rules for communication between web server and browser | 1 | 2 |
| | | | |

| 5(b)(i) | Two from: Picture and sound not synchronised 1 Interruptions // video not continuous 1 Can be degraded by other competing traffic 1 | Max 2 |
|-----------|--|-------|
| 5(b)(ii) | Dedicated communications channel between the two communicating devices 1 Established prior to start of communication // removal of links at end of communication 1 | 2 |
| 5(b)(iii) | In packet switching, packets can take different routes and may not arrive in order Will arrive in order (only one route) As packets can take many different routes / share paths with others can be delayed Dedicated circuit has full bandwidth No loss of synch 1 mark for any valid point | Max 3 |

| 5(a) | | Option 1 | Option 2 | | 3 |
|-----------|---|--|--|----------------------------|-------|
| | | Application Layer | Application Layer | | |
| | | Transport | Transport (Layer) | 1 | |
| | | Internet | Network (Layer) | 1 | |
| | | Network Interface | (Data) Link (Layer) | 1 | |
| 5(b)(i) | Peer-to- | -peer | | | 1 |
| 5(b)(ii) | File sha | ring | | | 1 |
| 5(b)(iii) | • Tor • File • BitT Allo A p Pee • Ond dow Lee • Cer the | owing them to work as seeds of eer can act as a 'seed' – used er downloading file can get pie ce a peer has a piece of the find who aded eches download much more the | es available to other peers / users or leeches. It to upload pieces of a file exces from different seeds simulate it can become a seed for the man they upload expers records of all the peers ('s | itaneously e parts | Max 4 |
| 5(c) | HTTP/H Used for FTP Used for SMTP Used for POP3 | r sending email messages | server to client | 1 1 1 1 1 1 | Max 4 |

5 (a) (i)



- (ii) software / module / program / code
- [1]
- (b) (i) For example:

check packet port ... [1]

to identify the application type [1]

check packet destination socket ... [1]

so that packet sent to correct application [1]

check incoming packet sequence number ... [1]

to ensure data is reassembled in correct order [1]

recalculate checksum of packet ... [1]

to ensure integrity of packet [1]

if packet checksum invalid ... [1]

send message to have packet retransmitted [1]

[Max 2 tasks]

[Max 4]

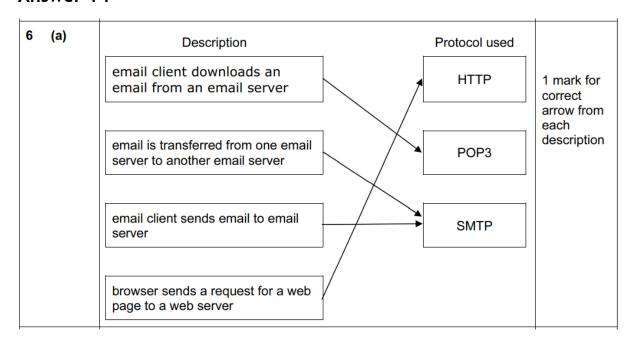
[1]

[1]

[1]

(ii) HTTP/HTTPS [1]

(iii) POP3 [1]



| (b) | peer-to-peer | 1 |
|---------|---|-----------|
| (c) (i) | Tracker: central server that: stores details of other computers that have all / part of file to be downloaded | 1 |
| | // has data on those peers downloading and uploading file // shares IP addresses with other clients in swarm allowing them to connect | 1 |
| (ii) | Seed: peer computer that has 100% of file // is uploading downloaded content | 1 |
| (iii) | Swarm: all the connected peer computers that have all or part of the file to be downloaded / uploaded // share a torrent | 1 |
| | n shale a torrent | Total: 11 |

| (b) (i) | Server: central computer stores files that are to be downloaded | 1 |
|---------|--|---|
| (ii) | Command: user can send action/instruction | 1 |
| | (or by example, e.g. change directory) that are carried out on server | 1 |
| (iii) | Anonymous: allows user to access files user does not need to identify themselves to server | 1 |

Answer 16

3 (a) Any point 1 mark

sender's IP address receiver's IP address packet sequence number checksum

[Max 2]

(b) Any point 1 mark

email has been split up into packets packet has destination address packets pass through many different routers in journey packets don't take same route routers use IP addresses packets reassembled at destination to rebuild email

[Max 3]

| (c) | Ar | ny point 1 mark | |
|-----|-----------------|---|--|
| | tin so pa | nail message is only read when all of it is received ne delays due to lost/delayed packets not significant sending different packets by different routes is not issue/is efficient ackets arriving out of order not an issue requirement for a continuous circuit (circuit switching) | |
| | ПО | requirement for a continuous circuit (circuit switching) | [Max 2] |
| (d) | Ci | rcuit switching | [1] |
| (e) | e. | g. real-time video/video conferencing | [1] |
| | Any | point 1 mark | |
| | full | uit made available is dedicated to this communication stream bandwidth available / no sharing lost packets | |
| | | ranteed quality of service | [Max 2] |
| | | dedicated circuit/channel/physical path which lasts for duration of connection | [1] [1] |
| | (b) | e.g. cs: gives dedicated circuit ps: split into packets/chunks ps: sends packets on individual routes cs: whole bandwidth available // ps: shares bandwidth cs: faster data transfer cs: packets arrive in order they are sent cs: packets cannot get lost cs: better for a real-time application ps: packets may arrive out of order so delay until packet order restored ps: packets may get lost so retransmission causes delays | [1] [1] [1] [1] [1] [1] [1] [1] [max. 6] |
| | (c) | web page divided into packets/chunks each packet has destination address router looks at IP address and decides where to send packet next for most efficient path packets can take different routes home computer reassembles packets to rebuild web page | [1] [1] [1] [1] [1] [max. 3] |