

CS-2006 Operating system

Quiz: 5 Memory management

Total-Marks:10

Q:1- Combined Paging and Segmentation

[5]

A system uses segmentation with paging. Each segment can be up to 1 MB in size.  
Page size is 4 KB.

- The virtual address is **32 bits** long
  - Each Page Table Entry (PTE) is 4 bytes
  - There are 16 segments per process
1. How many bits are required to specify the **segment number**?

4 bits

2. How many pages are there in **each segment**?

256 pages

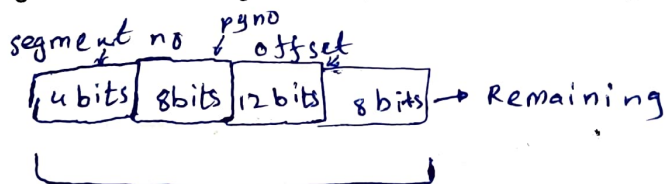
3. How many bits are needed to specify the **page number** within a segment?

8 bits

4. How many bits are needed for the **page offset**?

$d = 12$

5. Show the **final virtual address structure** by breaking it down into:  
Segment number, Page number, Offset and Remaining bits (if any)



**Q:2-** Consider a system with an 8-bit virtual address space, and a 32-byte page/frame size. The page table for a process maps virtual page numbers (VPN) to physical frame numbers (PFN) as follows: (0 → 5), (1 → 3), (2 → 12), (3 → 7), (4 → 15).

[5]

1. Translate the following virtual addresses (given in decimal) to physical addresses, indicating:
  - a. Virtual Page Number (VPN), b. Offset, c. Frame Number (PFN), d. Final Physical Address
2. Write the final physical address in both decimal and binary.
3. Indicate if any address is invalid (i.e., the VPN is not mapped in the page table).

**Given Virtual Addresses:**

(a) 50 (b) 100 (c) 145 (d) 200

<p>(a) 50</p> $\text{Pg no} = 50 / 32$ $= 1$ $\text{offset} = 50 \% 32$ $= 18$ $\text{frame} = 3$ <p>final PA = 114 ✓</p> <p>(b) 100 in binary <math>1100100</math> <sup>p=3 d=4</sup></p> $\text{pg no} = 3 \quad \text{frame} = 7$ $\text{offset} = 4$ $\text{PA} = (\text{frame} \times \text{pg size}) + \text{offset}$ $\text{PA} = 228$	<p>(c) 145</p> $\text{Pg} = 4$ $\text{offset} = 9$ $\text{frame} = 15$ $\text{PA} = 497$ ✓ <p>(d) 200</p> $\text{Pg} = 6$ $\text{offset} = 8$ <p>invalid as 6 is not mapped</p>
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