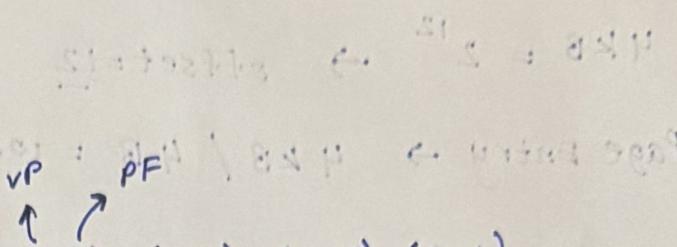


Abdul Munam Aziz

23L-3019

BCS-4F



Q1- Virtual Address = 6 bit  
page/frame = 16 byte

(a) Addresses =  $2^6 = 64$   $2^6 - 1 = 63$  addresses

displacement bits  $\rightarrow$  16 bytes  $\therefore 2^4 = 16$   $\therefore 4$  bits

• Binary of  $(17)_{10} = 010001$

• Page Number =  $17 \% 16 = 1$

•  $\text{PF} = 3 \rightarrow \text{PF base} = 3 \times 16 = 48$

Physical Address =  $48 + 1 = 49$

$(49)_{10} = \underline{\underline{(110001)}_2}$

(b)  $(35)_{10} = \underline{\underline{(100011)}_2}$

$$35 \% 16 = 3$$

• Page Number =  $35 \% 16 = 3 \Rightarrow \text{offset} = 3$  (Mapping (2,11) used)

•  $\text{PF} = 11 \rightarrow \text{PF base} = 11 \times 16 = 176$

• Physical Address =  $176 + 3 = 179$

-  $(179)_{10} = \underline{\underline{(10110011)}_2}$

Q2- L.P, P.P = 4 KB , Page Entry = 4 byte

(a)  $4 \text{ KB} = 2^{12} \rightarrow \underline{\text{offset}} = 12$

Page Entry  $\rightarrow 4 \text{ KB} / 4 \text{ B} = 1024 = 2^{\underline{10}}$  entries

max pages =  $64 - 12 = 42$   $\Rightarrow 2^{\underline{52}}$  pages

$\lceil 52 / 10 \rceil = \underline{\underline{6 \text{ Levels}}}$  |  $52 = 10 + 10 + 10 + 10 + 10 + 2$

(b)  $\underline{\underline{\text{Pld}}} = 2 \ 10 \ 10 \ 10 \ 10 \ 10 \ 12$

Q3- 4 Logical Pages

V.A = 16 =  $2^4$  ~~if~~

$(0,11), (1,35), (2,3), (3,1)$

Page size = 256 bytes =  $2^8$  ~~if 8~~ offset = 8 bits

VPN =  $16 - 8 = 8$  bits  $\underline{\underline{2^8 = 256 \text{ pages}}}$

Binary

(a)  $\underline{\underline{772}}_{10} = (0000001100000100)_2$

11 = 00001011

$\frac{772}{256} \quad 772 \% 256$   
· Page = 3 , offset = 4

35 = 00100011

· Frame = 1  $\rightarrow$  PF base =  $1 \times 256 = 256$

3 = 00000011

· Physical Address = 260

1 = 00000001

$\times \underline{\underline{\hspace{1cm}}}$

Pld

$(260)_{10} = (\underline{\underline{00000001}} \underline{\underline{000000100}})_2$

0000

1011  
0000  
0000

(b)  $(456)_{10} = 00000001111011000$

$456/256 = 1$        $456 \% 256 = 200$

Page = 1      offset = 200

Frame = 35  $\Rightarrow$  PF base =  $35 \times 256 = 8960$

Physical Address =  $8960 + 200 = 9160$

$(9160)_{10} = (0010\ 0011\ 1100\ 1000)_2$

\* Assumption of  
256 k bits, not  
256 k bytes

Q4- Logical Address = 20 bits

Memory size = 256 k bits =  $32\text{ kB} = 2^{15}$

page bits = 8

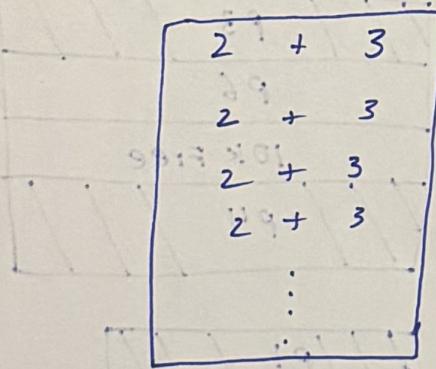
displace bits / offset = 12

(a) Page size =  $2^{12} = 4096$  bytes = 32768 bits

(b) Max Pages =  $2^8 = 256$  pages

(c)  $(256 \times 1024) \div (32768) = 8$  frames

(d) 5 bits in total



11 2 for attribute  
3 for frame/page

e) 5 bits

f) 1024 k bits = 12.8 kB =  $2^{17}$

frame bits = 5

$2^5 = 32$  frames

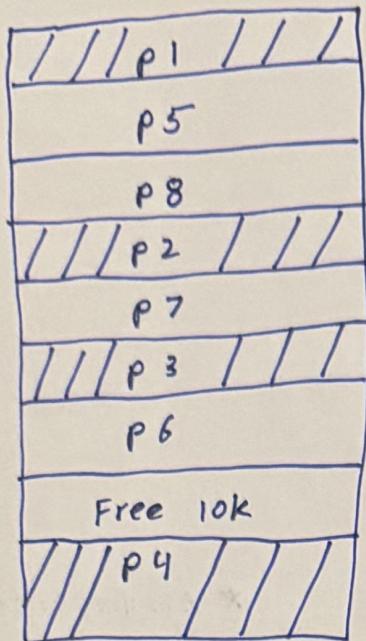
$5 + 2 = 7$  bits

5  $\rightarrow$  frame  
2  $\rightarrow$  attribute

Q5-

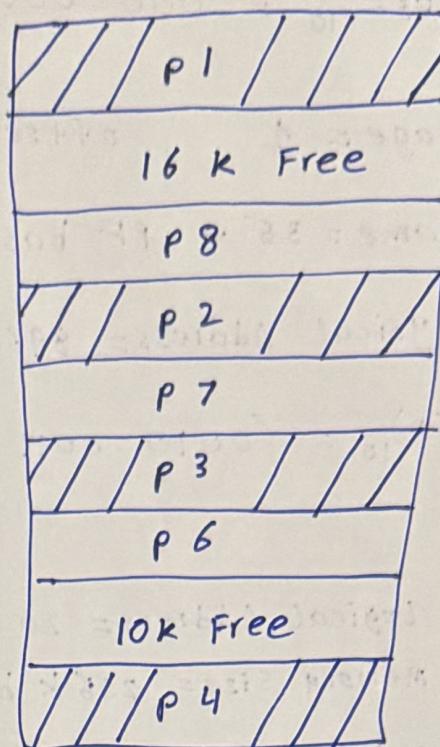
(a)

(iv)



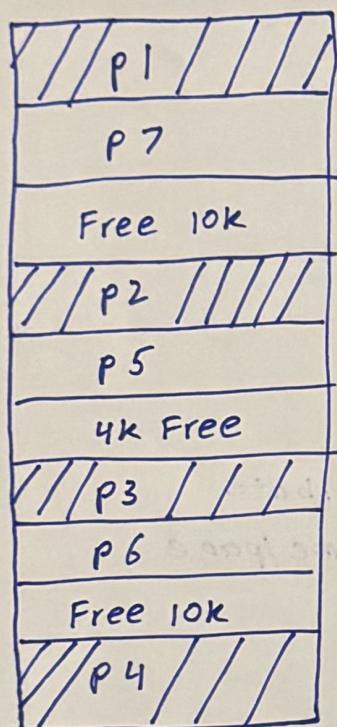
FIRST FIT

(vi)

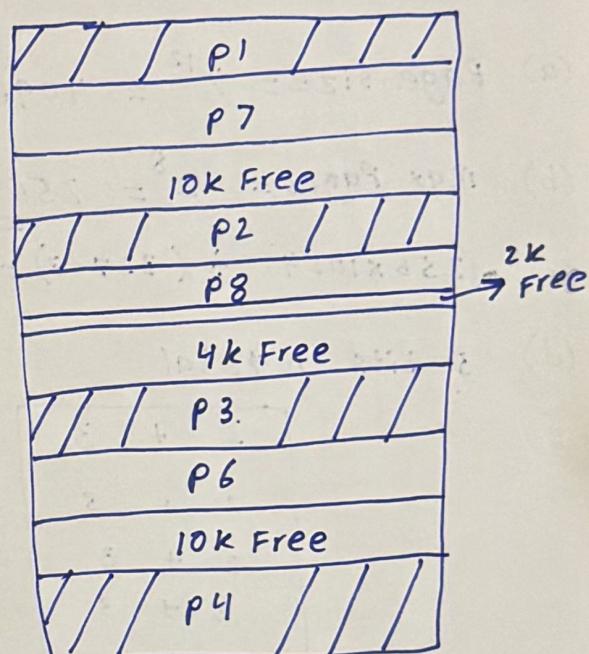


BEST FIT

(iv)

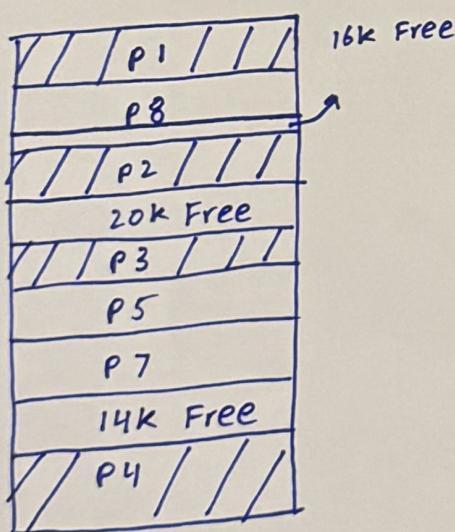


(vi)

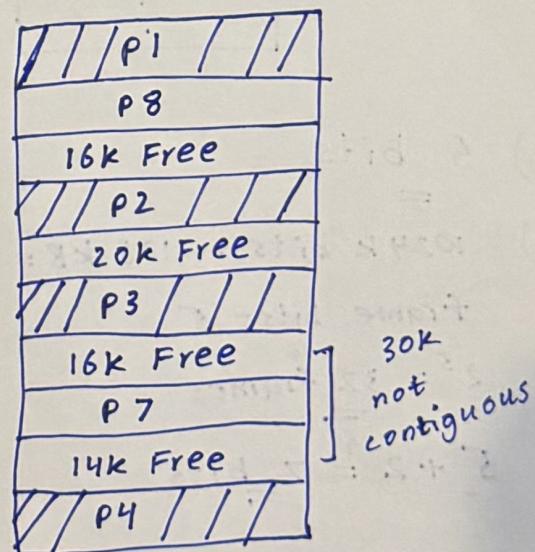


WORST FIT

(iv)

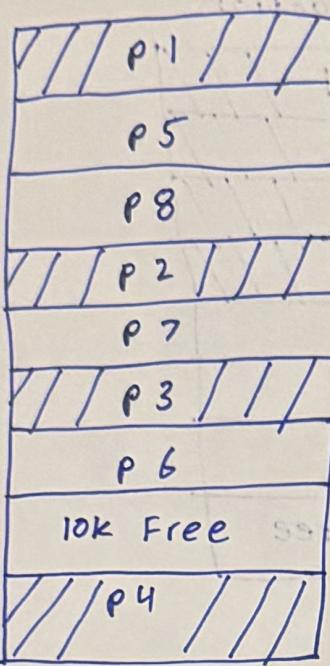


(vi)

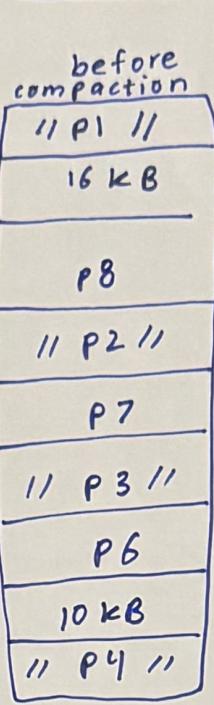


(b) Assuming that best-fit and worst-fit not required b/c of compaction, so only 2 maps assumed for answer (1 for (iv), 1 for (vi))

(iv)

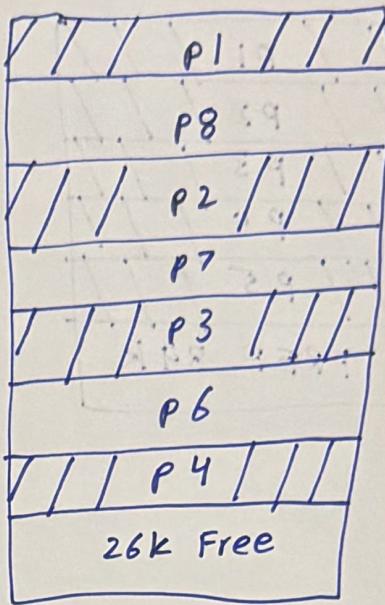


FIRST FIT



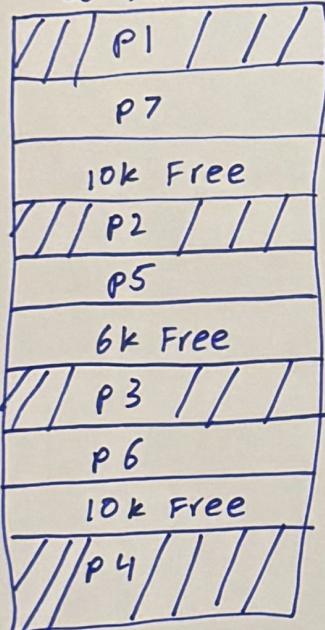
(vi)

After compaction



→ can't insert P9, no space

before compaction



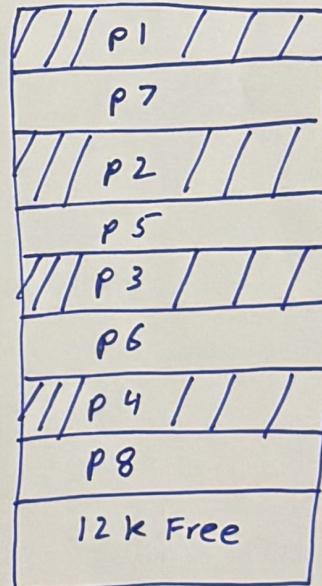
BEST FIT

P8 EXISTS, 2nd compaction

⋮ same until P8

P8  
16+12 = 28 K Free

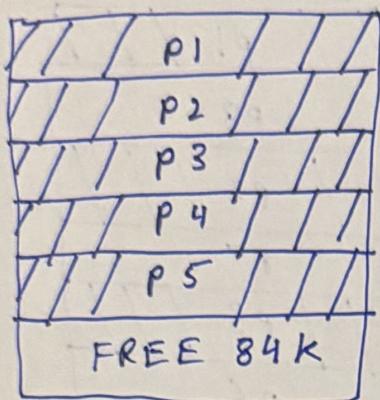
After compaction + insertion of P8 in released area



WORST

FIT

After compaction + inserting  
P5



leaves  
P5 levels +  
compaction

