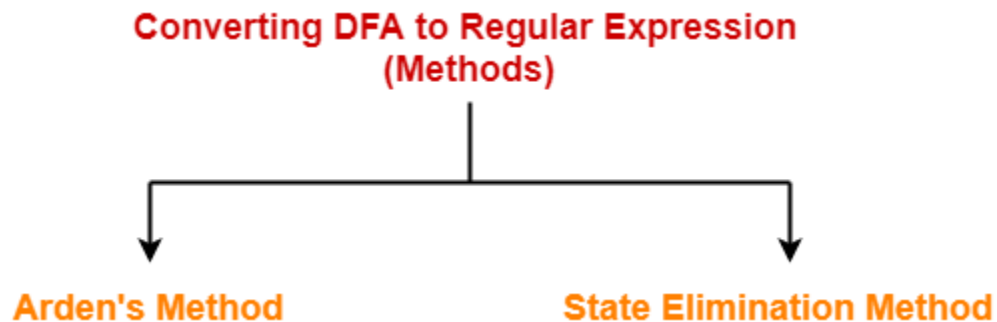


The two popular methods for converting a DFA to its regular expression are-



1. Arden's Method
2. State Elimination Method

In this article, we will discuss State Elimination Method.

State Elimination Method-

This method involves the following steps in finding the regular expression for any given DFA-

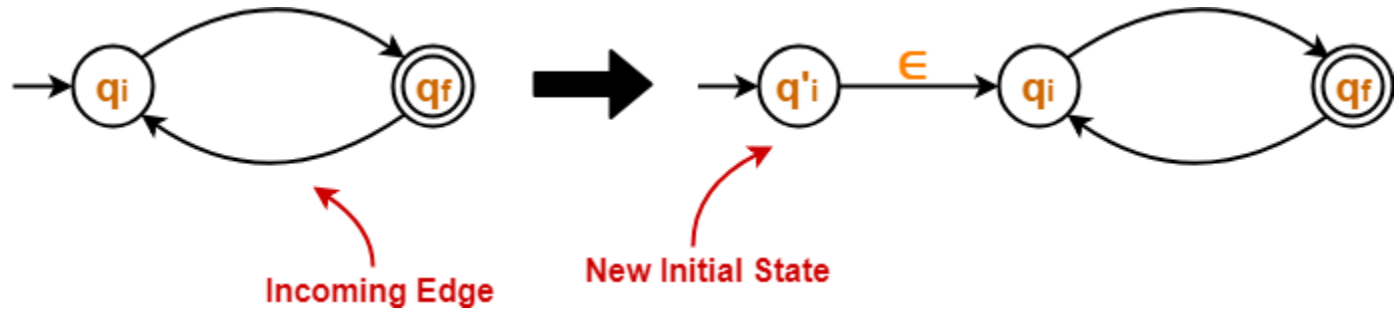
Step-01:

Thumb Rule

The initial state of the DFA must not have any incoming edge.

- If there exists any incoming edge to the initial state, then create a new initial state having no incoming edge to it.

Example-



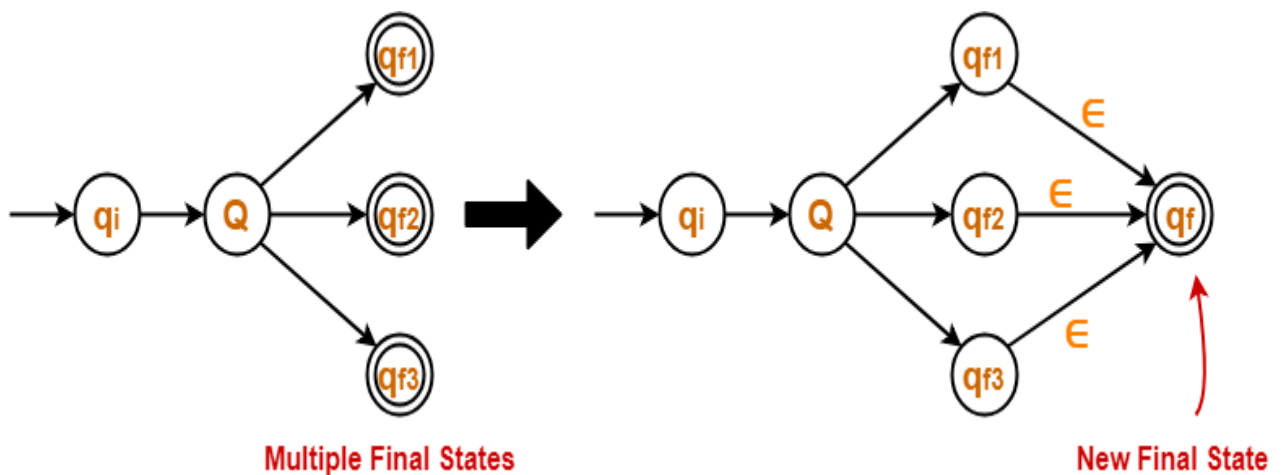
Step-02:

Thumb Rule

There must exist only one final state in the DFA.

If there exists multiple final states in the DFA, then convert all the final states into non-final states and create a new single final state.

Example-



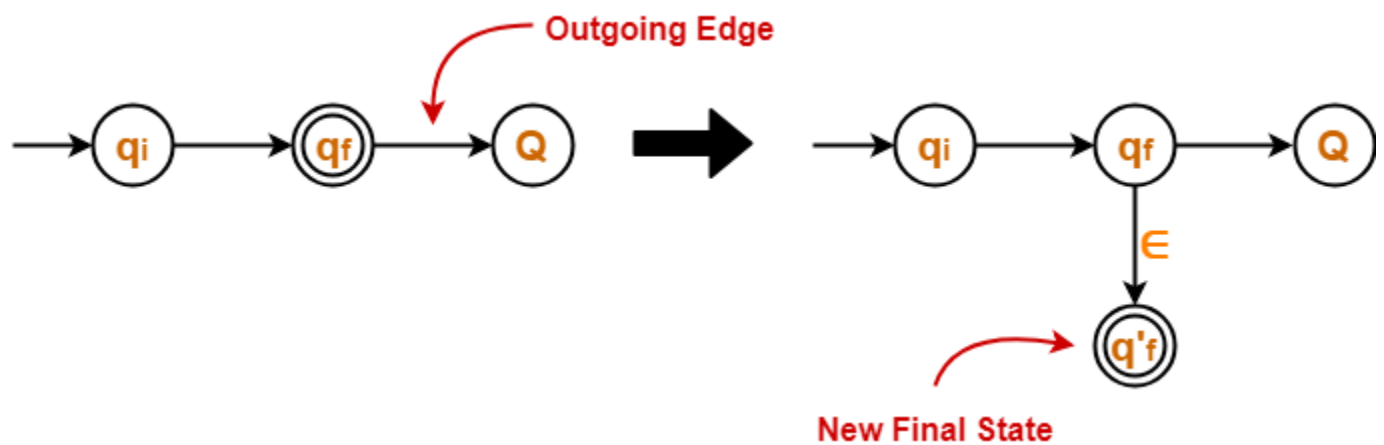
Step-03:

Thumb Rule

The final state of the DFA must not have any outgoing edge.

- If there exists any outgoing edge from the final state, then create a new final state having no outgoing edge from it.

Example-



Step-04:

- Eliminate all the intermediate states one by one.
- These states may be eliminated in any order.

In the end,

- Only an initial state going to the final state will be left.
- The cost of this transition is the required regular expression.

NOTE

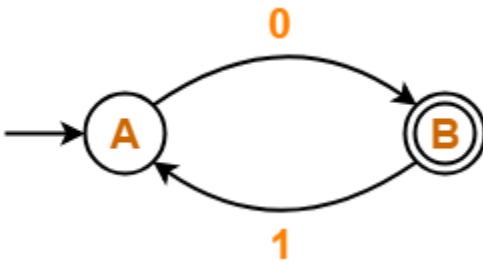
The state elimination method can be applied to any finite automata.

(NFA, ϵ -NFA, DFA etc)

PRACTICE PROBLEMS BASED ON CONVERTING DFA TO REGULAR EXPRESSION-

Problem-01:

Find regular expression for the following DFA-

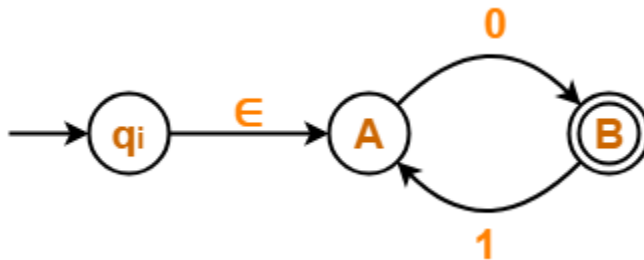


Solution-

Step-01:

- Initial state A has an incoming edge.
- So, we create a new initial state q_i .

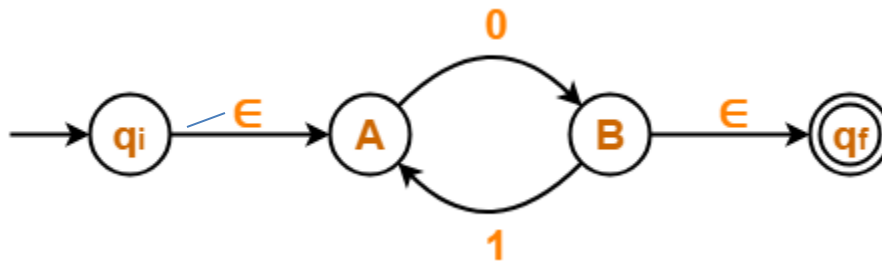
The resulting DFA is-



Step-02:

- Final state B has an outgoing edge.
- So, we create a new final state q_f

The resulting DFA is-



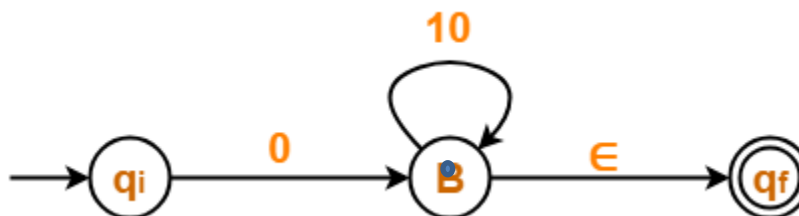
Step-03:

Now, we start eliminating the intermediate states.

First, let us eliminate state A.

- There is a path going from state q_i to state B via state A.
- So, after eliminating state A, we put a direct path from state q_i to state B having cost $\epsilon \cdot 0 = 0$
- There is a loop on state B using state A.
- So, after eliminating state A, we put a direct loop on state B having cost $1 \cdot 0 = 10$.

Eliminating state A, we get-

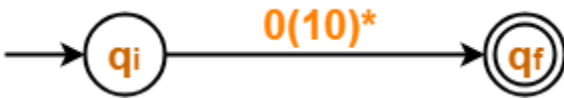


Step-04:

Now, let us eliminate state B.

- There is a path going from state q_i to state q_f via state B.
- So, after eliminating state B, we put a direct path from state q_i to state q_f having cost $0 \cdot (10)^* \cdot \epsilon = 0(10)^*$

Eliminating state B, we get-



From here,

Regular Expression = $0(10)^*$
--

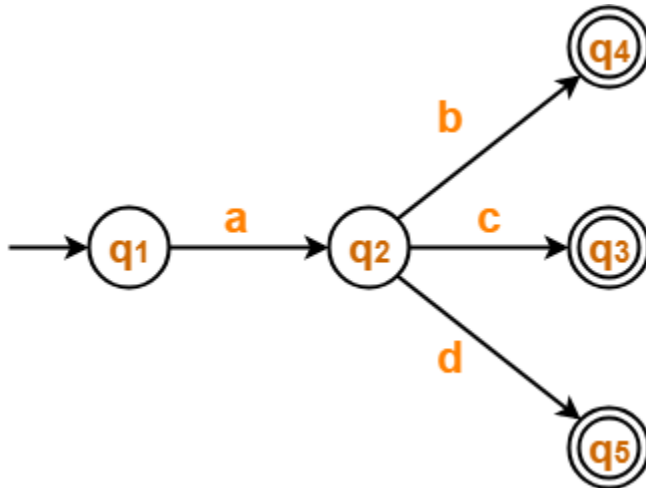
NOTE-

In the above question,

- If we first eliminate state B and then state A, then regular expression would be $= (01)^*0$.
- This is also the same and correct.

Problem-02:

Find regular expression for the following DFA-

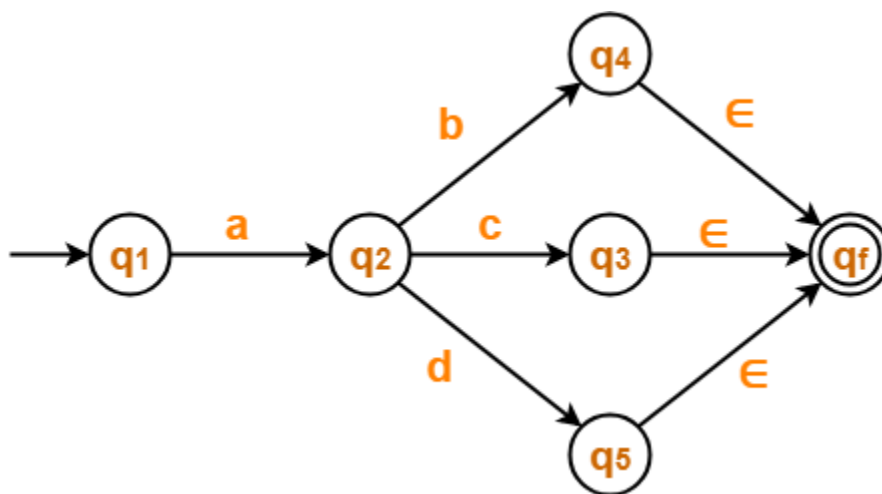


Solution-

Step-01:

- There exist multiple final states.
- So, we convert them into a single final state.

The resulting DFA is-

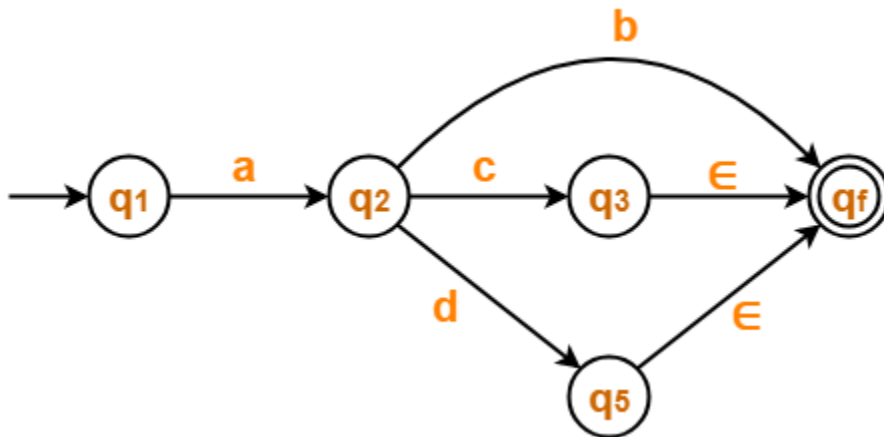


Step-02:

Now, we start eliminating the intermediate states.

First, let us eliminate state q_4 .

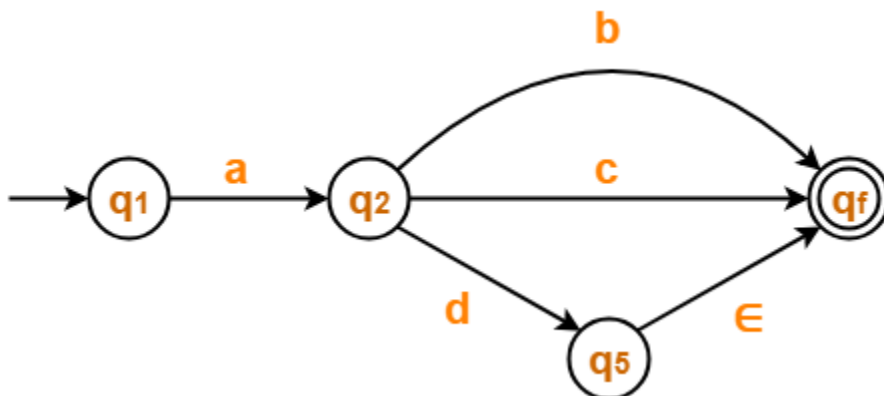
- There is a path going from state q_2 to state q_f via state q_4 .
- So, after eliminating state q_4 , we put a direct path from state q_2 to state q_f having cost $b \cdot \epsilon = b$.



Step-03:

Now, let us eliminate state q_3 .

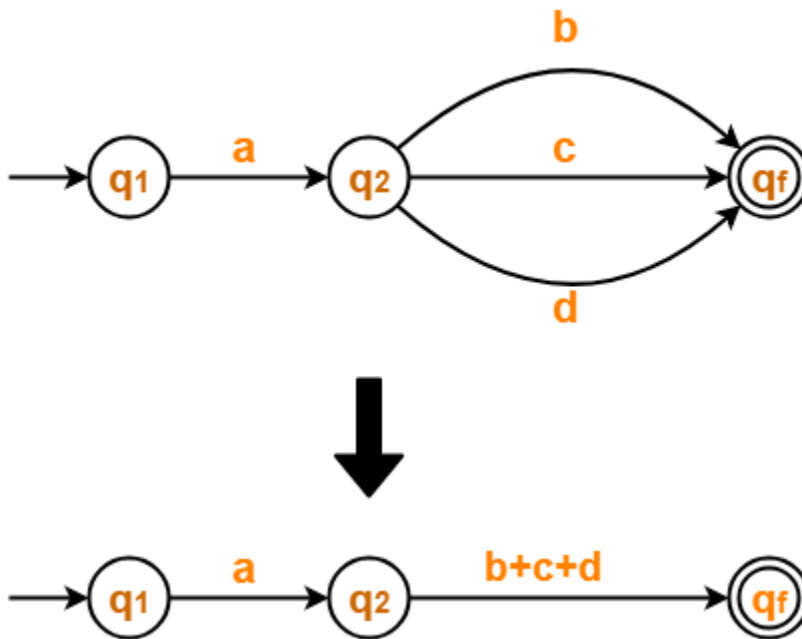
- There is a path going from state q_2 to state q_f via state q_3 .
- So, after eliminating state q_3 , we put a direct path from state q_2 to state q_f having cost $c \cdot \epsilon = c$.



Step-04:

Now, let us eliminate state q_5 .

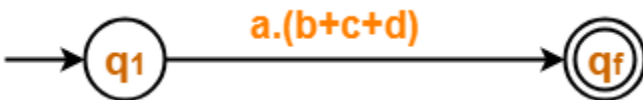
- There is a path going from state q_2 to state q_f via state q_5 .
- So, after eliminating state q_5 , we put a direct path from state q_2 to state q_f having cost $d \cdot \epsilon = d$.



Step-05:

Now, let us eliminate state q_2 .

- There is a path going from state q_1 to state q_f via state q_2 .
- So, after eliminating state q_2 , we put a direct path from state q_1 to state q_f having cost $a \cdot (b+c+d)$.

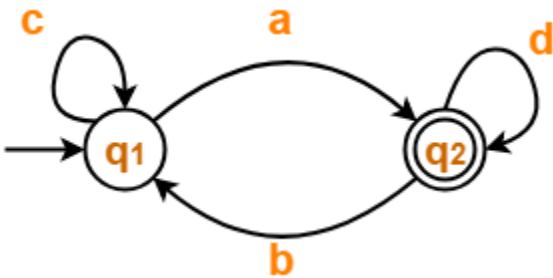


From here,

Regular Expression = $a(b+c+d)$

Problem-03:

Find regular expression for the following DFA-

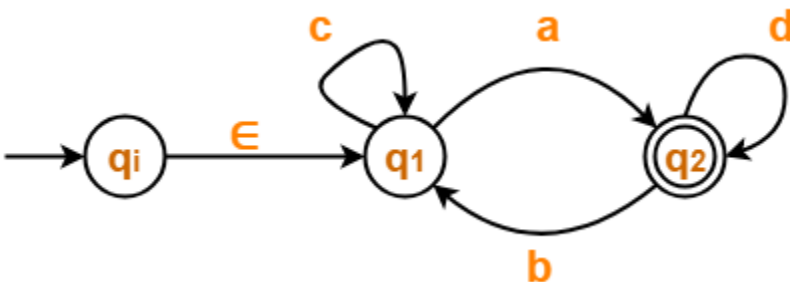


Solution-

Step-01:

- Initial state q_1 has an incoming edge.
- So, we create a new initial state q_i .

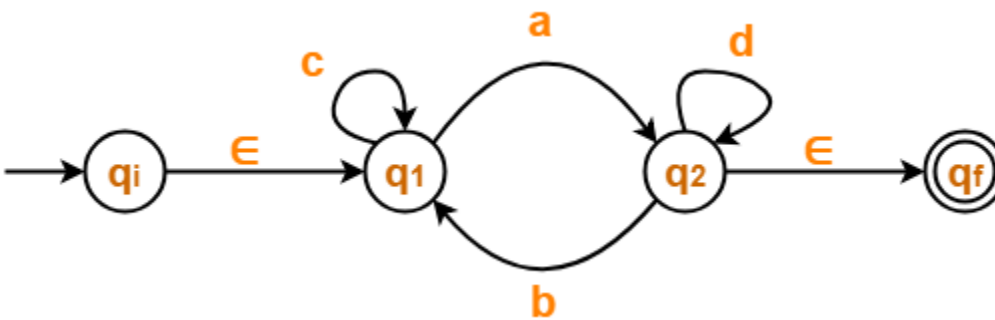
The resulting DFA is-



Step-02:

- Final state q_2 has an outgoing edge.
- So, we create a new final state q_f .

The resulting DFA is-



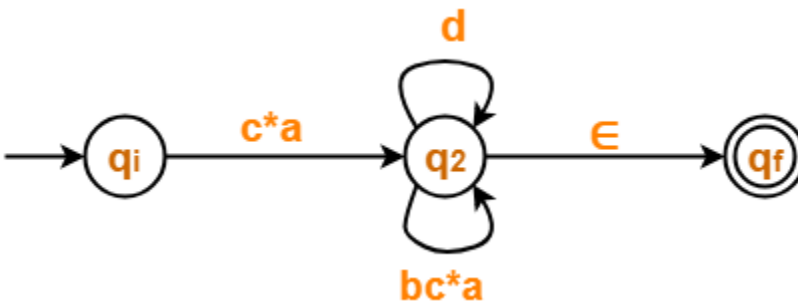
Step-03:

Now, we start eliminating the intermediate states.

First, let us eliminate state q_1 .

- There is a path going from state q_i to state q_2 via state q_1 .
- So, after eliminating state q_1 , we put a direct path from state q_i to state q_2 having cost $\epsilon.c*.a = c*a$
- There is a loop on state q_2 using state q_1 .
- So, after eliminating state q_1 , we put a direct loop on state q_2 having cost $b.c*.a = bc*a$

Eliminating state q_1 , we get-

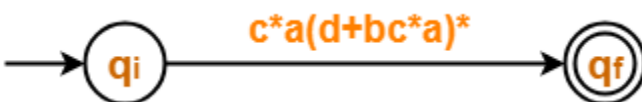


Step-04:

Now, let us eliminate state q_2 .

- There is a path going from state q_i to state q_f via state q_2 .
- So, after eliminating state q_2 , we put a direct path from state q_i to state q_f having cost $c^*a(d+bc^*a)^*\epsilon = c^*a(d+bc^*a)^*$

Eliminating state q_2 , we get-

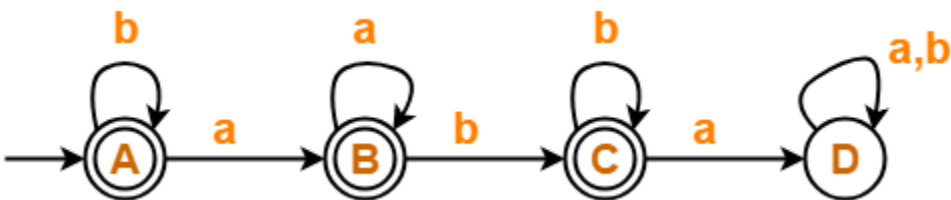


From here,

Regular Expression = $c^*a(d+bc^*a)^*$

Problem-04:

Find regular expression for the following DFA-

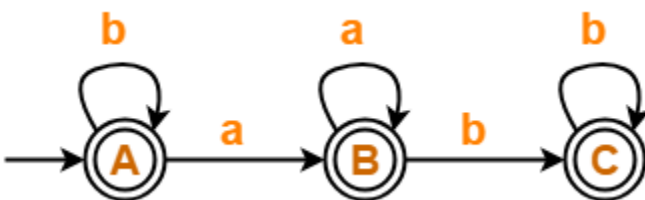


Solution-

Step-01:

- State D is a dead state as it does not reach to any final state.
- So, we eliminate state D and its associated edges.

The resulting DFA is-

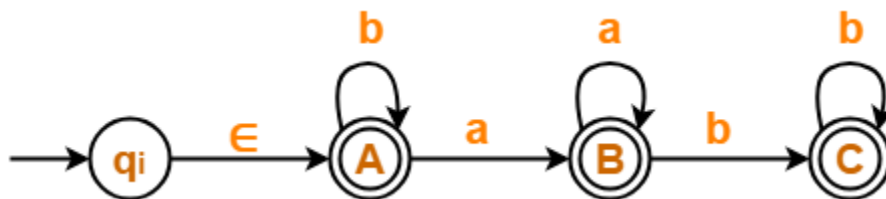


Step-02:

- Initial state A has an incoming edge (self loop).

- So, we create a new initial state q_i .

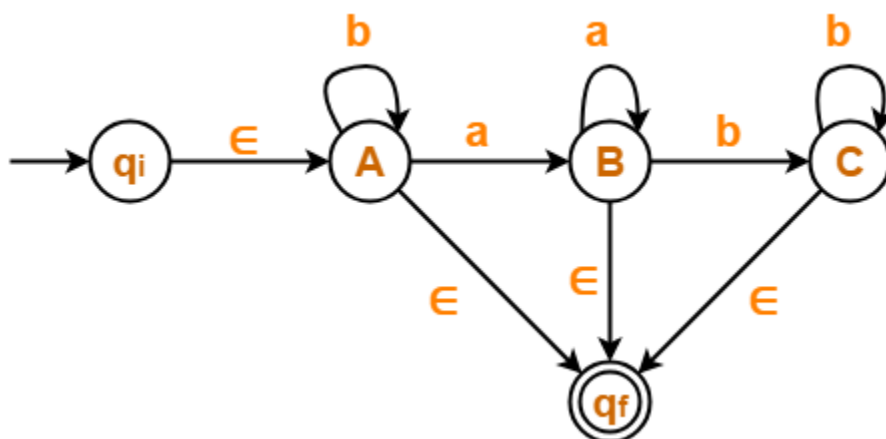
The resulting DFA is-



Step-03:

- There exist multiple final states.
- So, we convert them into a single final state.

The resulting DFA is-



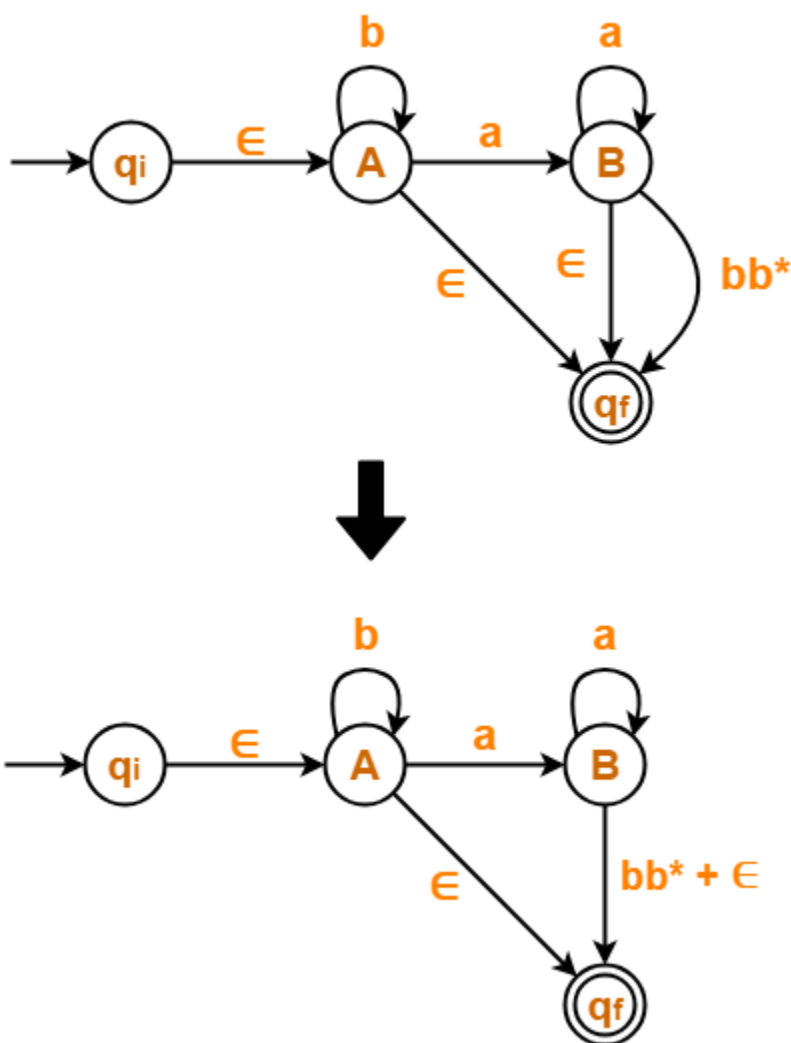
Step-04:

Now, we start eliminating the intermediate states.

First, let us eliminate state C.

- There is a path going from state B to state q_f via state C.
- So, after eliminating state C, we put a direct path from state B to state q_f having cost $b.b^*. \epsilon = bb^*$

Eliminating state C, we get-

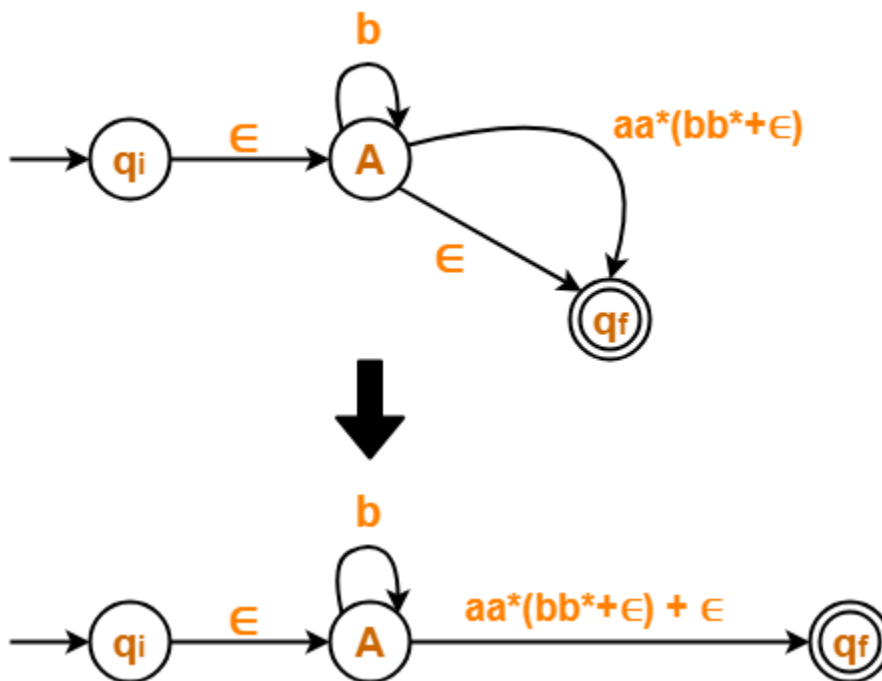


Step-05:

Now, let us eliminate state B.

- There is a path going from state A to state q_f via state B.
- So, after eliminating state B, we put a direct path from state A to state q_f having cost $a.a^*. (bb^* + \epsilon) = aa^*(bb^* + \epsilon)$

Eliminating state B, we get-

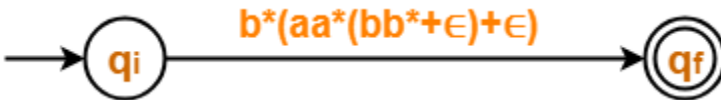


Step-06:

Now, let us eliminate state A.

- There is a path going from state q_i to state q_f via state A.
- So, after eliminating state A, we put a direct path from state q_i to state q_f having cost $\epsilon.b^*. (aa^*(bb^*+\epsilon)+\epsilon) = b^*(aa^*(bb^*+\epsilon)+\epsilon)$

Eliminating state A, we get-



From here,

Regular Expression = $b^*(aa^*(bb^*+\epsilon)+\epsilon)$

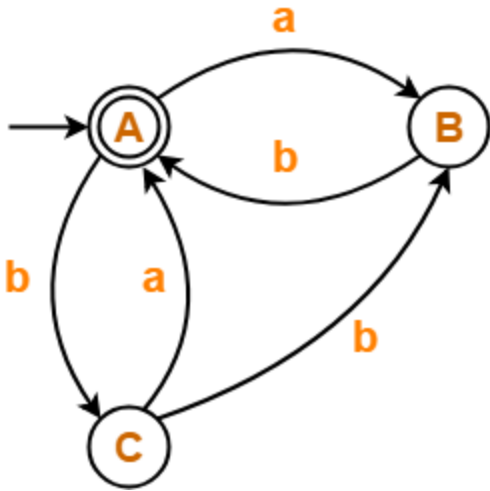
We know, $bb^* + \epsilon = b^*$

So, we can also write-

Regular Expression = $b^*(aa^*b^*+\epsilon)$

Problem-05:

Find regular expression for the following DFA-

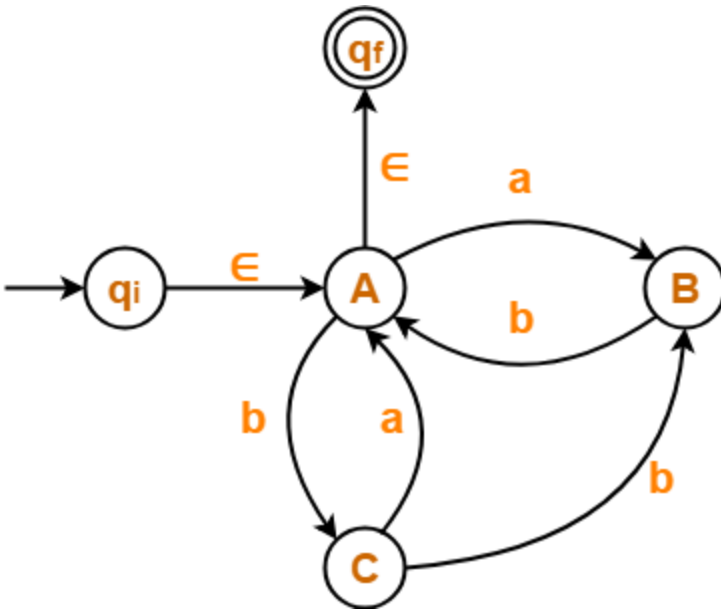


Solution-

Step-01:

- Since initial state A has an incoming edge, so we create a new initial state q_i .
- Since final state A has an outgoing edge, so we create a new final state q_f .

The resulting DFA is-



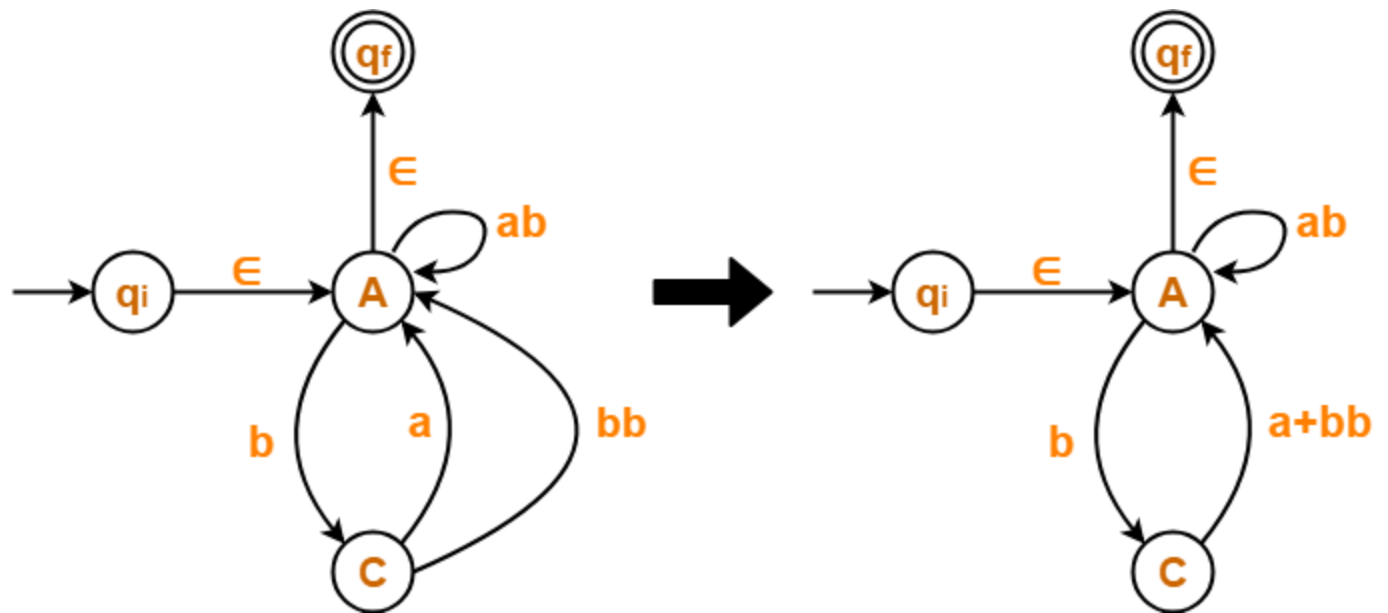
Step-02:

Now, we start eliminating the intermediate states.

First, let us eliminate state B.

- There is a path going from state C to state A via state B.
- So, after eliminating state B, we put a direct path from state C to state A having cost $b.b = bb$.
- There is a loop on state A using state B.
- So, after eliminating state B, we put a direct loop on state A having cost $a.b = ab$.

Eliminating state B, we get-

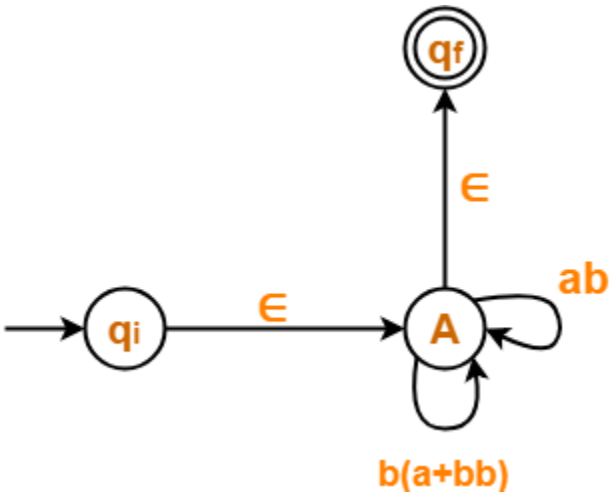


Step-03:

Now, let us eliminate state C.

- There is a loop on state A using state C.
- So, after eliminating state C, we put a direct loop on state A having cost $b.(a+bb) = b(a+bb)$

Eliminating state C, we get-

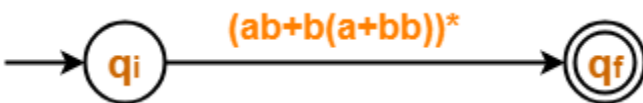


Step-04:

Now, let us eliminate state A.

- There is a path going from state q_i to state q_f via state A.
- So, after eliminating state A, we put a direct path from state q_i to state q_f having cost $\epsilon.(ab + b(a+bb))^*\epsilon = (ab + b(a+bb))^*$

Eliminating state A, we get-

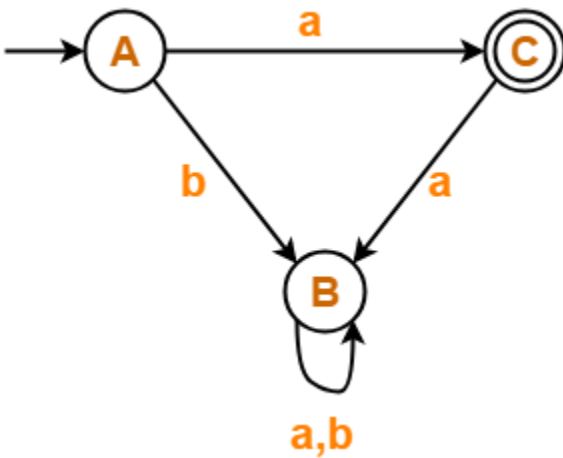


From here,

Regular Expression = $(ab + b(a+bb))^*$

Problem-06:

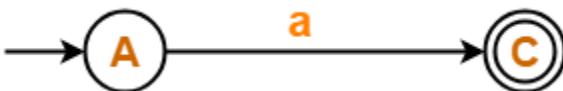
Find regular expression for the following DFA-



Solution-

- State B is a dead state as it does not reach to the final state.
- So, we eliminate state B and its associated edges.

The resulting DFA is-

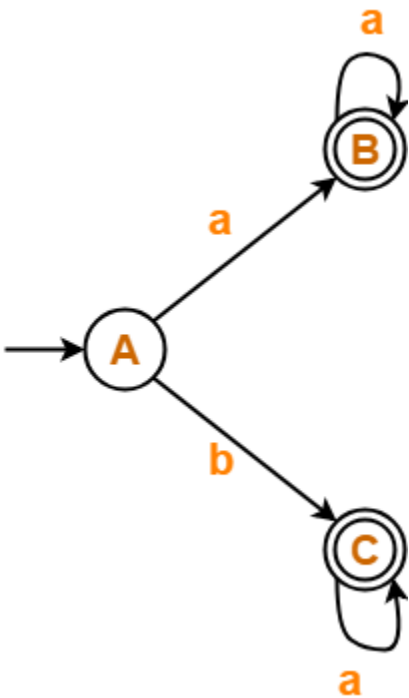


From here,

Regular Expression = a

Problem-07:

Find regular expression for the following DFA-

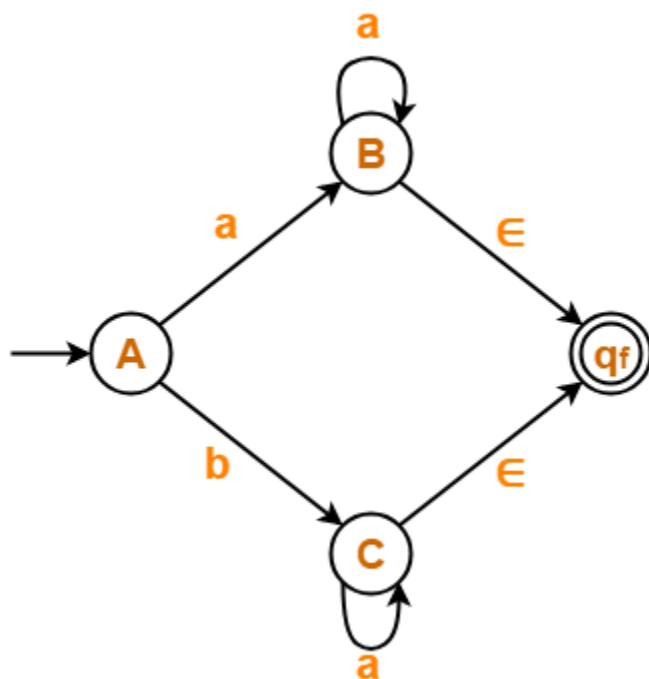


Solution-

Step-01:

- There exist multiple final states.
- So, we create a new single final state.

The resulting DFA is-



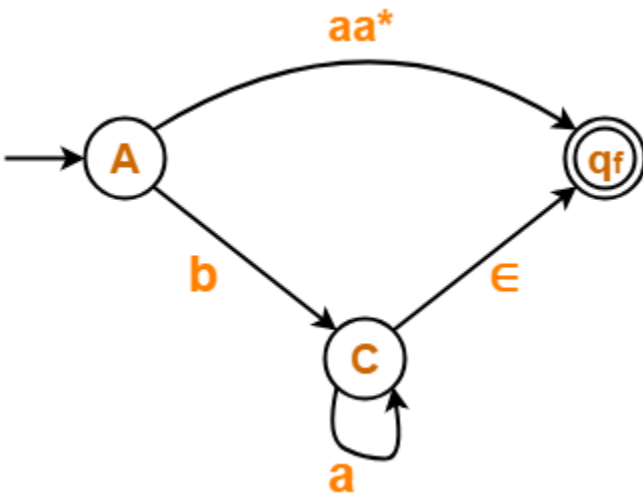
Step-02:

Now, we start eliminating the intermediate states.

First, let us eliminate state B.

- There is a path going from state A to state q_f via state B.
- So, after eliminating state B, we put a direct path from state A to state q_f having cost $a.a^*.ε = aa^*$.

Eliminating state B, we get-

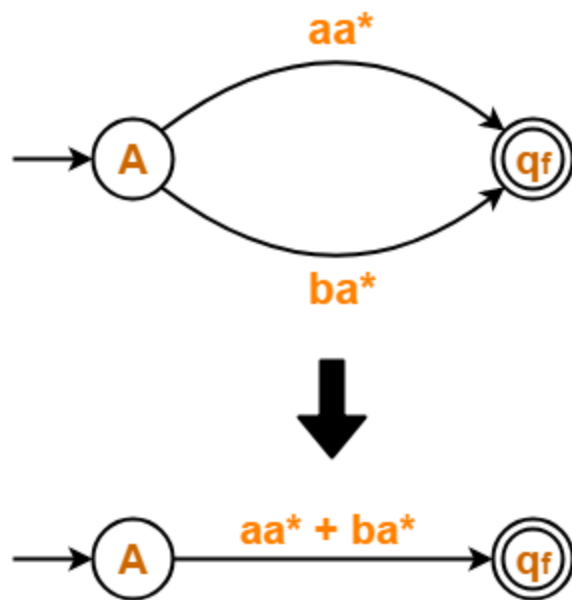


Step-03:

Now, let us eliminate state C.

- There is a path going from state A to state q_f via state C.
- So, after eliminating state C, we put a direct path from state A to state q_f having cost $b.a^*.ε = ba^*$.

Eliminating state C, we get-



From here,

Regular Expression = $aa^* + ba^*$
--