# **Problem 1: Liquidations Problem**

### Part (a): Calculating the Health Factor

From given Information we can calculate these:

- Exchange rates:
- 1 ETH = 1500 USDC
- 1 ETH =  $100 \text{ AXS} \rightarrow 1 \text{ AXS} = 0.01 \text{ ETH}$

Since Bob holds 1000 USDC, we convert this to ETH:

Applying the collateral factor of 0.9:

$$0.6667 \times 0.9 = 0.6 ETH$$

- Debt in ETH: 0.2 ETH
- Debt in AXS (converted to ETH):  $20 \times 0.01 = 0.2$  ETH
- Total Debt: 0.2 + 0.2 = 0.4 ETH

Health Factor = BorrowCapacity / value(TotalDebt)

$$= 0.6 / 0.4 = 1.5$$

Since 1.5 > 1, Bob's position is healthy and does not need liquidation.

### Part (b): When the AXS Exchange Rate Changes

New AXS exchange rate: 1 AXS = 1/48 ETH ≈ 0.02083 ETH

New AXS debt in ETH:  $20 \times 0.02083 = 0.4167$  ETH

New total debt: 0.2 + 0.4167 = 0.6167 ETH

Step 2: Compute New Health Factor

 $0.6 / 0.6167 \approx 0.973$ 

Since 0.973 < 1, Bob's collateral must be liquidated.

#### Part (c): How Much USDC Needs to Be Liquidated?

A liquidator is willing to clear Bob's ETH debt in exchange for USDC collateral at a rate of 1520 USDC/ETH.

To restore a healthy position, let x be the amount of ETH debt to be cleared.

New debt after liquidation: 0.6167 - x

Solving for x:

$$0.6167 - x <= 0.6$$

$$=> x >= 0.0167 ETH$$

Choose x = 0.0167 ETH (the min value of x)

USDC to be liquidated =  $0.0167 \times 1520 \approx 25.38$  USDC

Rounding up, Bob will lose at least 26 USDC from his collateral to restore a healthy position.

## Problem 2: Slippage

We are tasked with showing that slippage (s) is always positive and approximates x/X, assuming X >> x. Below is the step-by-step derivation:

We are given the following formulas:

1. Constant Product Formula (Uniswap invariant):

$$x * y = k$$

2. Amount of Y tokens Alice needs to send to buy x amount of X tokens:

$$y = Y * x / (X - x)$$

3. Exchange rate Alice gets from Uniswap:

$$y/x = Y/(X-x)$$

4. Open market exchange rate:

$$p = Y / X$$

5. Slippage definition:

$$s = ((y / x) - p) / p$$

$$s = ((Y / (X - x)) - (Y / X)) / (Y / X)$$

$$s = (Y * (1 / (X - x) - 1 / X)) / (Y / X)$$

$$s = ((1/(X-x)-1/X)) * X$$

Using the identity (1 / (X - x) - 1 / X) = x / (X \* (X - x)):

$$s = X * (x / (X * (X - x))) = x / (X - x)$$

If X >> x, then  $X - x \approx X$ , so we can approximate:

$$s \approx x / X$$

Since x, X > 0, it follows that s > 0, meaning the exchange rate on Uniswap is always worse than on the open market.

However, as the liquidity pool size X increases, s decreases, meaning larger liquidity pools reduce slippage.

# Problem 3: Sandwich Attacks on Uniswap

### Part (a): Effect of Sam's Tx1 on Alice's Transaction

Initially, the pool has reserves x tokens of type X and y tokens of type Y, with k = xy. Alice's transaction Tx sends  $\alpha x$  tokens of type X, increasing the X reserves to  $x + \alpha x$ . Uniswap maintains the constant product xy = k, so Alice receives:

$$y' = (xy) / (x + \alpha x)$$

Now, when Sam front-runs Alice with Tx1, he sends  $\beta$ x tokens of type X to the pool first. This increases the X reserves to x +  $\beta$ x and reduces the Y reserves accordingly:

$$y1 = (xy) / (x + \beta x)$$

Then, Alice's Tx executes, sending αx more tokens of type X. The new X reserves become:

$$x + \beta x + \alpha x$$

and Uniswap ensures the constant product:

$$y' = (xy) / (x + \beta x + \alpha x)$$

The amount of Y Alice receives is now:

$$y' - y1 = (xy) / (x + \beta x + \alpha x) - (xy) / (x + \beta x)$$

Since  $x + \beta x + \alpha x > x + \alpha x$ , Alice now receives fewer Y tokens than she would have without Sam's front-running.

### Part (b): How Does Sam's Profit Depend on β?

Sam's profit comes from the arbitrage between his front-run and back-run transactions. In Tx1, Sam increases the X reserves by  $\beta x$ , which reduces the Y reserves and makes Y more expensive.

Alice then swaps her X, getting fewer Y tokens than before.

In Tx2, Sam swaps back the Y he bought earlier, benefiting from the temporary price impact

he created.

Increasing  $\beta$  makes Alice's swap worse (for her), which means Sam can extract more profit up to a point.

However, if  $\beta$  is too large, Sam's Tx1 creates excessive slippage, reducing his own profit.

Conclusion: Sam's profit initially increases with  $\beta$  but can decrease if  $\beta$  is too large due to excessive price impact.

### Part (c): Can Victor Mount the Attack Without Sam?

Yes! Since Victor, the validator, controls the transaction ordering, they can copy Sam's strategy.

However, Victor does not have X tokens to initiate the front-run.

If Victor has access to flash loans, they can temporarily borrow X tokens to execute Tx1 and repay in Tx2, keeping all profit for themselves.

Conclusion: If flash loans are available, Victor can replace Sam and execute the sandwich attack solo.

### Part (d): Can Alice Use MEV-Boost to Protect Herself?

MEV-Boost allows Alice to submit her transaction privately to a trusted block builder instead of the public mempool.

This prevents Sam (or any searcher) from seeing her transaction in advance, blocking sandwich attacks.

If Alice uses MEV-Boost, the only party that can still manipulate the order is the builder itself.

However, if Alice picks a trusted, honest builder, she can avoid the attack altogether.

Conclusion: MEV-Boost can help protect Alice, but it depends on the integrity of the builder she selects.