



DESIGN AND ANALYSIS OF ALGORITHMS

Presented to Sir Waheed Abro

Abstract

The complete report for the project.

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Abstract

This document is the proper documentation of the project developed for the course of Design and Analysis of Algorithms submitted to our instructor Dr. Waheed Abro. The project includes an introduction of the project and the group members, the project design of programming, the experimental setup, results and discussions, conclusions and references used.

Introduction

This project includes a visual representation of multiple sorting algorithms. These are executed on test files with input. The visualization of the sorts is a key feature. We receive the execution time, and the space complexity. The group members are Mustafa Zahid 20K-1045 and Munir Abbasi 20K-0244.

Project Programming Design

The programming language chosen was python. This language was used since python is easy to use and its libraries can be used to make an interactive GUI and manage the backend as well. We were also able to implement the algorithm for fast parallel sort. This was coded in C++ language as multiple processors were incorporated using threads and the OpenMP library. All relevant research paper logic are matched with our code in images.

Parallel fast sorting

```
int max_threads=dimensions*cbrt(dimensions);
num_buckets=cbrt(dimensions);
...
buckets = (struct bucket *) calloc(dimensions-cbrt(dimensions), sizeof(struct bucket));
```

Algorithm 4—parallel sort using $n^{4/3}$ processors

1. Partition the n input numbers into $n^{2/3}$ groups each having $n^{1/3}$ elements.

```
int buck_sum=0;
for (j=cur_id; j< num_buckets*num_threads; j=j+num_threads){
    buck_sum += buckets[j].element;
}
```

2. Within each group do

For each element, j , determine $\text{count}[j] = \#$ of i such that $c_i < c_j$
+ ($\#$ of $i \leq j$ such that $c_i = c_j$).

```
for (i=0; i< dimensions; i++){
    j = A[i]/w;
    if (j > num_buckets -1)
        j = num_buckets-1;
    k = j + cur_id*num_buckets;
    b_index = buckets[k].curr_ind++;
    B[b_index] = A[i];
}
```

```

        int j,k;
        int local_index;
        for(int i=0;i<=local_index;i++)
        binarySearch(A,j,k,i);

```

```

int binarySearch(int arr[], int l, int r, int x)

```

```

    if (r >= l) {
        int mid = l + (r - l) / 2;
        if (arr[mid] == x)
            return mid;
        if (arr[mid] > x)
            return binarySearch(arr, l, mid - 1, x);
        return binarySearch(arr, mid + 1, r, x);
    }
    return -1;

```

EX\MUNIRABBAS\university\semester 5\algo\ALGOCODE\parallelsort.cpp - Dev-C++ 5.11

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TM-GCC 4.9.2 64-bit Profiling

(globals)

parallelsort.cpp

```

68     buckets = (struct bucket *) calloc(dimensions-cbrt(dimensions), sizeof(struct bucket)); //define the bucket structure used
69     t1 = omp_get_wtime();
70     int j,k;
71     int local_index;
72     #pragma omp parallel
73     {
74         num_threads = omp_get_num_threads();
75         int real_bucket_index; // [0 : num_buckets * num_threads)
76         int cur_id = omp_get_thread_num();
77         workload = dimensions/num_threads; //as per paper divide total buckets numbers per processor
78         int previous_index;
79         #pragma omp for private(i,local_index)
80         for (i=0; i< dimensions;i++){
81             local_index = A[i]/w;
82             if (local_index > num_buckets-1)
83                 local_index = num_buckets-1;
84             real_bucket_index = local_index + cur_id*num_buckets;
85             buckets[real_bucket_index].element++; //implement on the local bucket registers
86         }
87         int buck_sum=0;
88         for (j=cur_id; j< num_buckets*num_threads; j=j+num_threads){
89             buck_sum += buckets[j].element;
90         }
91         global_n_elem[cur_id]=buck_sum;
92         #pragma omp barrier
93         #pragma omp master
94

```

EX\MUNIRABBAS\university\semester 5\algo\ALGOCODE\parallelsort.cpp - Dev-C++ 5.11

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(globals)

parallelsort.cpp

```

93     #pragma omp barrier
94     #pragma omp master
95     {
96         for (j=1; j<num_buckets; j++){
97             global_starting_position[j] = global_starting_position[j-1] + global_n_elem[j-1];
98             buckets[j].start_ind = buckets[j-1].start_ind + global_n_elem[j-1];
99             buckets[j].curr_ind = buckets[j-1].curr_ind + global_n_elem[j-1]; //formula application for calculation of the offset
100         }
101     }
102     #pragma omp barrier
103     for (j=cur_id*num_buckets; j< num_buckets*num_threads; j=j+num_threads){
104         int previous_index = j-num_buckets; //barrier for each bucket
105         buckets[j].start_ind = buckets[previous_index].start_ind + buckets[previous_index].element;
106         buckets[j].curr_ind = buckets[previous_index].curr_ind + buckets[previous_index].element;
107         // binarySearch(A,j,k,i)
108     }
109     #pragma omp barrier
110     #pragma omp for private(i, b_index)
111     for (i=0; i< dimensions ;i++){
112         j = A[i]/w;
113         if (j > num_buckets -1)
114             j = num_buckets-1; //final bucket sort
115         k = j + cur_id*num_buckets;
116         b_index = buckets[k].curr_ind++;
117         B[b_index] = A[i];
118     }
119

```

Then evaluate $\text{count}[j] = \text{the sum (over } k) \text{ of count}[j, k]$.

8. Do a bucket sort of all n elements.

9. END of Algorithm 4.

```
    }  
    #pragma omp for private(i)  
    for(i=0; i<num_buckets; i++)  
        qsort(B+global_starting_position[i], global_n_elem[i], sizeof(int), cmpfunc);  
}  
  
total = omp_get_wtime() - t1;  
tmp = A;  
A = B;  
B = tmp;  
if (dimensions <= 40) {  
    printf("A \n");  
    for(i=0; i<dimensions; i++) {  
        printf("%d ", A[i]);  
    }  
    printf("\n");  
}  
printf("Sorting %d elements took %f seconds\n", dimensions, total);  
int sorted = 1;  
for(i=0; i<dimensions-1; i++) {
```

Experimental results

Parallel sort code execution

```
E:\MUNIRABBAS\university\semester 5\algo\ALGOCODE\parallelsort.exe  
Give length of array to sort  
99  
number of buckets is 4  
  
Unsorted data  
41 67 34 0 69 24 78 58 62 64 5 45 81 27 61  
91 95 42 27 36 91 4 2 53 92 82 21 16 18  
95 47 26 71 38 69 12 67 99 35 94 3 11 22  
33 73 64 41 11 53 68 47 44 62 57 37 59 23  
41 29 78 16 35 90 42 88 6 40 42 64 48 46  
5 90 29 70 50 6 1 93 48 29 23 84 54 56  
40 66 76 31 8 44 39 26 23 37 38 18 82 29  
  
Sorted final bucket is  
0 1 2 3 4 5 5 6 6 8 11 11 12 16 16  
18 18 21 22 23 23 23 24 26 26 27 27 29 29  
29 29 31 33 34 35 35 36 37 37 38 38 39 40  
40 41 41 41 42 42 42 44 44 45 46 47 47 48  
48 50 53 53 54 56 57 58 59 61 62 62 64 64  
64 66 67 67 68 69 69 70 71 73 76 78 78 81  
82 82 84 88 90 90 91 91 92 93 94 95 95 99  
  
Sorting 99 elements took 0.000000 seconds  
-----  
Process exited after 1.499 seconds with return value 0  
Press any key to continue . . .
```

```
E:\MUNIRABBASI\university\semester 5\algo\ALGOCODE\parallelsort.exe
Give length of array to sort
43
number of buckets is 3

Unsorted data
41 67 34 0 69 24 78 58 62 64 5 45 81 27 61
91 95 42 27 36 91 4 2 53 92 82 21 16 18
95 47 26 71 38 69 12 67 99 35 94 3 11 22

Sorted final bucket is
0 2 3 4 5 11 12 16 18 21 22 24 26 27 27
34 35 36 38 41 42 45 47 53 58 61 62 64 67
67 69 69 71 78 81 82 91 91 92 94 95 95 99

Sorting 43 elements took 0.000000 seconds

-----
Process exited after 0.5775 seconds with return value 0
Press any key to continue . . .
```

```
Give length of array to sort
399
number of buckets is 7

Unsorted data
41 67 34 0 69 24 78 58 62 64 5 45 81 27 61
91 95 42 27 36 91 4 2 53 92 82 21 16 18
95 47 26 71 38 69 12 67 99 35 94 3 11 22
33 73 64 41 11 53 68 47 44 62 57 37 59 23
41 29 78 16 35 90 42 88 6 40 42 64 48 46
5 90 29 70 50 6 1 93 48 29 23 84 54 56
40 66 76 31 8 44 39 26 23 37 38 18 82 29
41 33 15 39 58 4 30 77 6 73 86 21 45 24
72 70 29 77 73 97 12 86 90 61 36 55 67 55
74 31 52 50 50 41 24 66 30 7 91 7 37 57
87 53 83 45 9 9 58 21 88 22 46 6 30 13
68 0 91 62 55 10 59 24 37 48 83 95 41 2
50 91 36 74 20 96 21 48 99 68 84 81 34 53
99 18 38 0 88 27 67 28 93 48 83 7 21 10
17 13 14 9 16 35 51 0 49 19 56 98 3 24
8 44 9 89 2 95 85 93 43 23 87 14 3 48
0 58 18 80 96 98 81 89 98 9 57 72 22 38
92 38 79 90 57 58 91 15 88 56 11 2 34 72
55 28 46 62 86 75 33 69 42 44 16 81 98 22
51 21 99 57 76 92 89 75 12 0 10 3 69 61
88 1 89 55 23 2 85 82 85 88 26 17 57 32
32 69 54 21 89 76 29 68 92 25 55 34 49 41
12 45 60 18 53 39 23 79 96 87 29 49 37 66
49 93 95 97 16 86 5 88 82 55 34 14 1 16
71 86 63 13 55 85 53 12 8 32 45 13 56 21
58 46 82 81 44 96 22 29 61 35 50 73 66 44
59 92 39 53 24 54 10 45 49 86 13 74 22 68
18 87 5 58 91 2 25 77 14 14 24 34 74 72
59 33 70 87 97 18

Sorted final bucket is
```

```
E:\MUNIRABBASI\university\semester 5\algo\ALGOCODE\parallelsort.exe
17 18 18 18 18 18 18 18 19 20 21 21 21 21
21 21 21 21 22 22 22 22 22 22 23 23 23 23
23 23 24 24 24 24 24 24 24 24 25 25 26 26
27 27 27 28 28 29 29 29 29 29 29 29 29 30
30 30 31 31 32 32 32 32 33 33 33 33 34 34
34 34 34 35 35 35 35 36 36 36 37 37 37 37
37 38 38 38 38 38 39 39 39 39 40 40 41 41
41 41 41 41 41 42 42 42 42 43 44 44 44 44
44 44 45 45 45 45 45 46 46 46 46 47 47 47
48 48 48 48 48 48 49 49 49 49 50 50 50 50
50 50 51 51 52 53 53 53 53 53 53 54 54 54
54 55 55 55 55 55 55 55 56 56 56 56 57 57
57 57 57 57 58 58 58 58 58 58 58 59 59 59
59 59 60 61 61 61 61 62 62 62 62 63 64 64
64 66 66 66 66 67 67 67 67 68 68 68 68 68
69 69 69 69 70 70 70 70 71 71 72 72 72 72
73 73 73 73 74 74 74 74 75 75 76 76 76 77
77 77 78 78 79 79 80 81 81 81 81 81 82 82
82 82 82 83 83 83 84 84 85 85 85 85 86 86
86 86 86 86 87 87 87 87 87 88 88 88 88 88
88 88 89 89 89 89 90 90 90 90 91 91 91 91
91 91 91 92 92 92 92 92 93 93 93 93 94 94
95 95 95 95 96 96 96 96 97 97 97 97 98 98
98 98 99 99 99 99

Sorting 399 elements took 0.000000 seconds

-----
Process exited after 0.1446 seconds with return value 0
Press any key to continue . . .

total = omp_get_wtime() - t1;
tmp = A;
A = B;
B = tmp;
if (dimensions <= 400) {
    //printf("\t");
    printf("\tSorted final bucket is  \n");
    for(i=0;i<dimensions;i++) {
        printf("%d\t",A[i]);
    }
    printf("\n\n");
}
printf("Sorting %d elements took %f seconds\n", dimensions,total);
```

Note how sorting elements takes such small time that the system clock can only show 0.0000000 since the time taken is so much less of the parallel sort. Now normal bucket sort takes how much time lets see.

Normal bucket Sort for 399 elements.

```
E:\MUNIRABBASI\university\semester 5\algo\Untitled1.exe
Bucket[3]: 21 21 21 21 21 21 21 21 22 22 22 22 22 22 22 23 23 23 23 23 23 24 24 24 24 24 24 25 25 26 26 26 27 27 27
Bucket[4]: 28 28 29 29 29 29 29 29 29 29 30 30 30 31 31 32 32 32 33 33 33 33 34 34 34 34 34 34
Bucket[5]: 35 35 35 35 36 36 36 37 37 37 37 37 38 38 38 38 39 39 39 39 40 40 41 41 41 41 41 41 41
Bucket[6]: 42 42 42 42 43 44 44 44 44 44 44 45 45 45 45 45 46 46 46 46 47 47 48 48 48 48 48
Bucket[7]: 49 49 49 49 50 50 50 50 50 51 51 52 53 53 53 53 53 53 53 54 54 55 55 55 55 55 55 55
Bucket[8]: 56 56 56 56 57 57 57 57 57 57 58 58 58 58 58 58 59 59 59 59 60 61 61 61 62 62 62 62
Bucket[9]: 63 64 64 64 66 66 66 66 67 67 67 67 68 68 68 68 68 69 69 69 69 69
Bucket[10]: 70 70 70 71 71 72 72 72 72 73 73 73 74 74 74 74 75 75 76 76 76
Bucket[11]: 77 77 77 77 78 78 79 79 80 81 81 81 81 82 82 82 82 83 83
Bucket[12]: 84 84 85 85 85 85 86 86 86 86 86 87 87 87 87 87 88 88 88 88 88 88 89 89 89 89 90 90 90 90
Bucket[13]: 91 91 91 91 91 91 91 92 92 92 92 93 93 93 93 94 95 95 95 95 95 96 96 96 96 97 97
Bucket[14]: 98 98 98 98 99 99 99 99

0.812311 is the total seconds needed for sorting 399 elements

Sorted array: 0 0 0 0 0 0 1 1 1 2 2 2 2 2 2 3 3 3 3 4 4 5 5 5 5 6 6 6 6 7 7 7 8 8 8 9 9 9 9 10 10 10 10 11 11 11 12 12
12 12 12 13 13 13 13 13 14 14 14 14 14 15 15 16 16 16 16 16 16 17 17 18 18 18 18 18 18 19 20 21 21 21 21 21 21 21
22 22 22 22 22 23 23 23 23 23 23 24 24 24 24 24 24 25 25 26 26 26 27 27 27 28 28 29 29 29 29 29 29 29 30 30 30
31 31 32 32 32 33 33 33 33 34 34 34 34 34 35 35 35 35 36 36 36 37 37 37 37 38 38 38 38 39 39 39 39 40 40 41 41
41 41 41 41 41 42 42 42 42 43 44 44 44 44 44 45 45 45 45 45 46 46 46 46 47 47 48 48 48 48 48 49 49 49 49 50
50 50 50 51 51 51 52 53 53 53 53 53 53 54 54 54 55 55 55 55 55 56 56 56 57 57 57 57 57 58 58 58 58
58 58 59 59 59 59 60 61 61 61 61 62 62 62 62 63 64 64 66 66 66 66 67 67 67 67 68 68 68 68 69 69 69 69 70 70 70
71 71 72 72 72 72 73 73 73 73 74 74 74 74 75 75 76 76 76 77 77 77 77 78 78 79 79 80 81 81 81 81 82 82 82 82 83 83
83 84 84 85 85 85 85 86 86 86 86 86 87 87 87 87 87 88 88 88 88 88 88 89 89 89 89 90 90 90 91 91 91 91 91
91 92 92 92 92 92 93 93 93 93 94 95 95 95 95 95 96 96 96 97 97 97 98 98 98 98 99 99 99 99

-----
Process exited after 1.179 seconds with return value 0
Press any key to continue . . .
```

```
int t1 = omp_get_wtime();
BucketSort(array);
int total = omp_get_wtime() - t1;

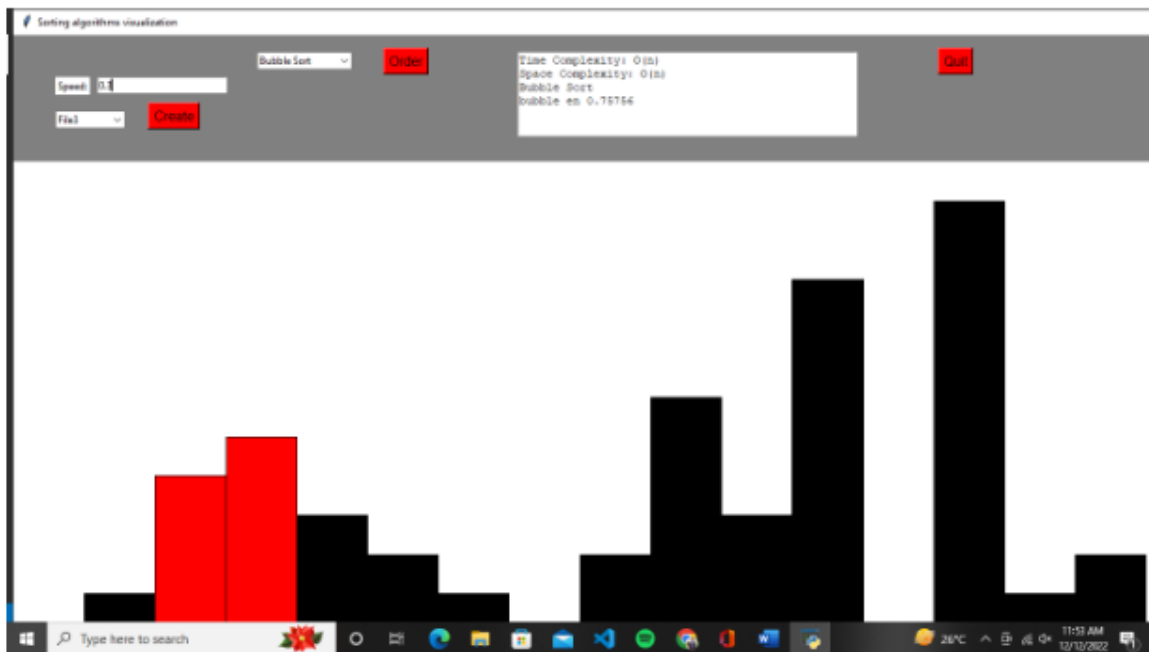
printf("\n\n%f is the total seconds needed for sorting 399 elements\n\n",total);
```

We cannot even deduce mathematically how fast the parallel bucket sort is as compared to a normal bucket sort as the clock time of parallel one is 0.000000, for sake of execution let's assume it is 0.1446 as mentioned in program end. For normal bucket it is less, but extremely large when compared to parallel bucket one as time is 1.179 seconds and 0.812311 for bucket.

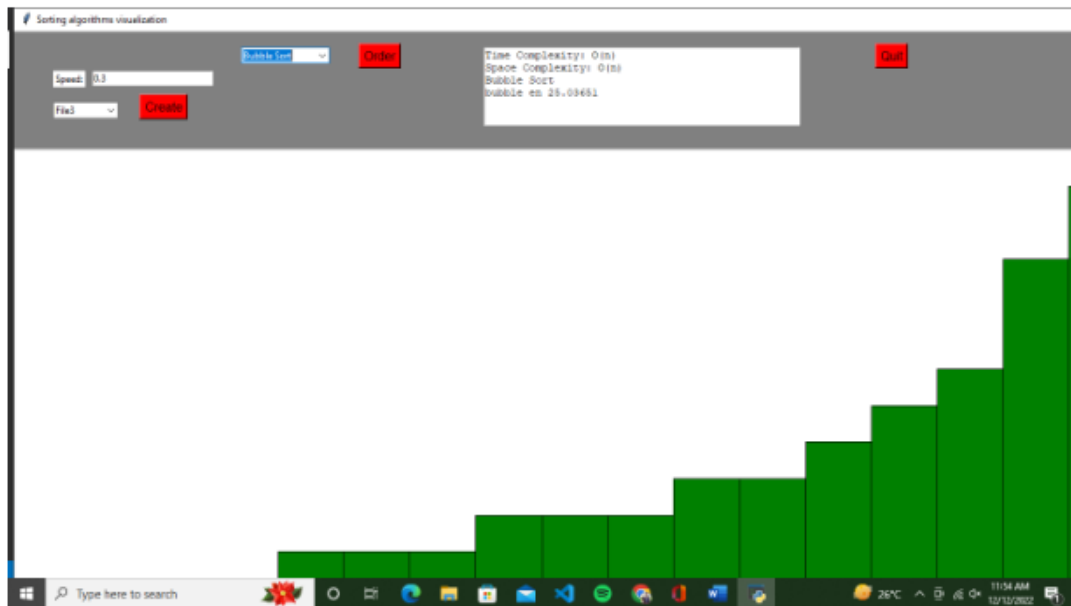
Experimental Setup

The setup for the programming was done by using Visual Studio. The code was written here. These are some screenshots for all the 10 implemented sorts.

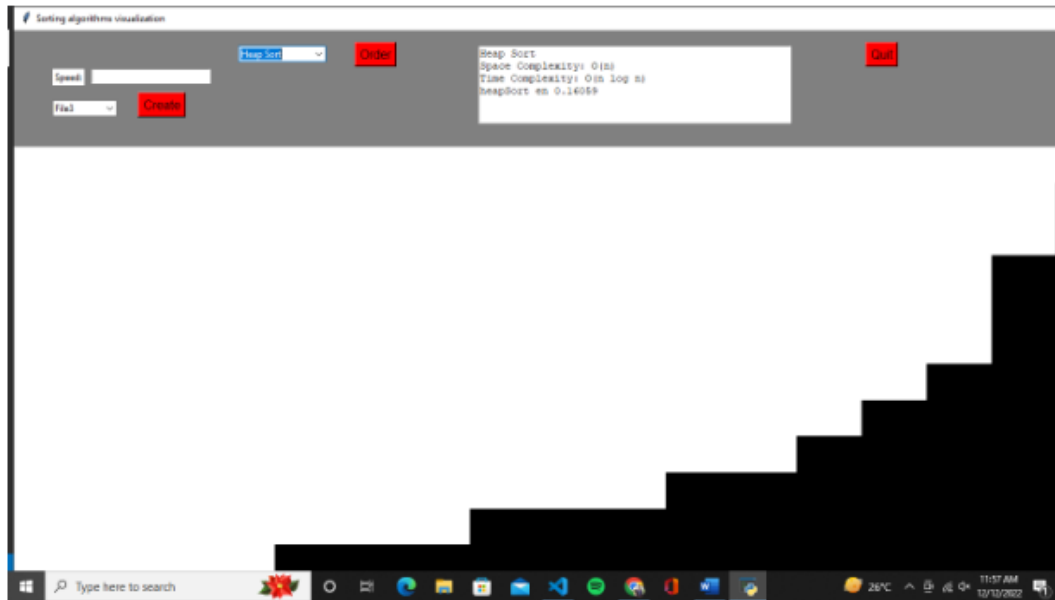
Bubble Sort in progress



Bubble Sort completed



Heap sort



Results and discussions

Let us compile some results for each sort.

Algorithm	Time Complexity			Space Complexity	Stable?
	Best	Average	Worst		
<u>Bubble Sort</u>	$\Omega(n)$	$\theta(n^2)$	$O(n^2)$	$O(1)$	Yes
<u>Insertion Sort</u>	$\Omega(n)$	$\theta(n^2)$	$O(n^2)$	$O(1)$	Yes
<u>Heap Sort</u>	$\Omega(n \log(n))$	$\theta(n \log(n))$	$O(n \log(n))$	$O(1)$	No

Algorithm	Time Complexity			Space Complexity	Stable?
	Best	Average	Worst		
<u>Quick Sort</u>	$\Omega(n \log(n))$	$\theta(n \log(n))$	$O(n^2)$	$O(n)$	No
<u>Merge Sort</u>	$\Omega(n \log(n))$	$\theta(n \log(n))$	$O(n \log(n))$	$O(n)$	Yes
<u>Bucket Sort</u>	$\Omega(n + k)$	$\theta(n + k)$	$O(n^2)$	$O(n)$	Depends
<u>Radix Sort</u>	$\Omega(n.k)$	$\theta(n.k)$	$O(n.k)$	$O(n + k)$	Yes
<u>Count Sort</u>	$\Omega(n + k)$	$\theta(n + k)$	$O(n + k)$	$O(k)$	Yes

Space complexity restrictions

When we have a priority of using less space, then it is best to use bubble sort, insertion sort and heap sort. However in the least space for an unsorted array, it is best to use heap sort and on average and worst cases it will perform better than the other two. Heap is best if the stability of the sort is not a problem.

Time complexity restrictions

When we have a time restriction, it is best to use merge sort. It gives a fixed complexity of running time. It is stable as well. If stability is not an issue, we can use heap sort.

File 1=> Big Range, unique data

File 2=>Small range, repeated numbers

File 3=>Already sorted data

File 4=>Big range, repeated numbers

File 5=>Nearly sorted

Algorithm	Time Taken (seconds)				
	File 1	File 2	File 3	File 4	File 5
<u>Bubble Sort</u>	0.32318	0.11126	0.00908	1.42703	0.01392
<u>Insertion Sort</u>	0.32601	0.12561	0.01283	1.37864.	0.01659
<u>Heap Sort</u>	0.09544	0.04476	0.07625	0.25297	0.03766
<u>Quick Sort</u>	0.57128	0.23902	0.92755	2.17132	0.23413
<u>Merge Sort</u>	0.19647	0.09125	0.13542	0.47491	0.08237
<u>Bucket Sort</u>	0.13652	0.05219	0.07081	0.14463	0.04503
<u>Radix Sort</u>	0.62434	0.31078	0.62312	0.6209	0.30793
<u>Count Sort</u>	0.13651	0.13381	0.11689	0.28774	0.28388

- For the modified quicksort or 7.4.5, we make the following deductions for the value of K.
 - $T(N)$ = running time of quicksort + running time of insertion sort (for the remaining N/K arrays of size K (or less than K))
 - $N (\log N - \log K \text{ and } n/k \text{ for in bracket due to division}) + N/K (K^2 \text{ as } n \text{ square}) = > N \log(N/K) + NK$.
 - In theory, K has to be chosen such that : $NK + N \log (N/K) = N \log N$ for same complexity==>
 - $K = \log N$ (K can't be more than $\log N$ otherwise the total running time will be more than $N \log N$ which deems the change pointless).
 - In practice, K should be the largest list length on which insertion sort is faster than quicksort which can be observed as well through experimental analysis.

Finding => Quicksort algorithm is efficient if the size of the input is very large. But, insertion sort is more efficient than quick sort in case of small arrays as the number of comparisons and swaps are less compared to quicksort. So we combine the two algorithms to sort efficiently using both approaches.

- For the 8.2.4 or the modified counting sort/range. We make use of counting sort. We know that in counting sort we maintain the sum of the occurrences of all the elements in the array. So $\text{array}[b] - \text{array}[a-1]$ will return the answer in $O(1)$ time.

3. The parallel bucket sort is almost 800% faster (with the assumed values which are still very big 0.8 normal vs 0.1 parallel) than the normal bucket sort. We can be sure that this factor depends upon the number of threads used which can be the buckets as per the paper.
4. Heap sort was best for sorting files that had big range and unique data along with those with a small range and repeated numbers.
5. On already sorted data bubbles sort was fastest, had a close call with insertion sort.
6. Bucket sort was best on big range with repeated numbers.
7. Bubble sort was fastest on nearly sorted data with an again close call with insertion sort.

Conclusion

We can conclude that different sorts are feasible in different scenarios as we see from the time complexities. Another important finding is that we can design modified algorithms as per our needs to get the accurate results while keeping in mind our constraints and like time and space.

References

1. www.youtube.com
2. Geeksforgeeks.com
3. Wikipedia.com