# **Algebra**

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# **Contents**

1 Polynomials				
	1.1	Division Algorithm	5	
	1.2	The Fundamental Theorem of Algebra	7	
	1.3	Quadratic Polynomials	10	
	1.4	Roots of Cubic Polynomials	12	
	1.5	Lagrange Interpolation	14	

# Polynomials

**Definition 1.0.1.** A Polynomial P(x) is an one variable expression or function of the form

$$P(x) = \sum_{i=0}^{n} a_i x^i = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$$

where  $a_0, a_1, \dots, a_n$  are constants and  $n \in \mathbb{N}$ . The constants  $a_i$  are called the *coefficients* of the polynomial. We will denote A[x] as the set of all polynomials with  $a_i \in A$ . If  $n \neq 0$  then n is called the *degree* of the polynomial P(x) and write  $\deg P(x) = n$ . If  $a_n = 1$  then we say that the polynomial is *monic*. r is called a *root* of the polynomial P(x) if and only if P(r) = 0.

## §1.1 Division Algorithm

**Theorem 1.1.1** (The Division Algorithm)

Given two polynomial A(x) and B(x) there exists unique polynomials Q(x) and R(x) with deg  $R(x) < \deg B(x)$  such that,

$$A(x) = Q(x)B(x) + R(x)$$

The polynomials Q(x) and R(x) are known as the *quotient* and the *remainder*, respectively. If the remainder R(x) = 0 then we say that B(x) divides A(x) and write  $B(x) \mid A(x)$ .

**Proof:** We will first prove the uniqueness of the polynomials Q(x) and R(x). Assume,

$$A(x) = Q_1(x)B(x) + R_1(x), \quad \deg R_1(x) < \deg B(x)$$
  
 $A(x) = Q_2(x)B(x) + R_2(x), \quad \deg R_2(x) < \deg B(x)$ 

Now,

$$(Q_1(x) - Q_2(x)) B(x) + (R_1(x) - R_2(x)) = 0$$

Let  $q(x) = Q_1(x) - Q_2(x)$  and  $r(x) = R_1(x) - R_2(x)$ . Now,

$$q(x)B(x) + r(x) = 0$$

$$\implies q(x)B(x) = -r(x)$$

If  $q(x) \neq 0$  then  $\deg r(x) = \deg q(x) + \deg B(x) \geq \deg B(x)$ . But that is impossible since  $\deg R_2(x) < \deg B(x) \implies \deg (R_2(x) - R_1(x)) < \deg B(x)$ . Thus q(x) must be zero. Consequently r(x) will also be zero. Therefore  $R_1(x) = R_2(x)$  and  $Q_1(x) = Q_2(x)$ .

Now we will prove the existence of the polynomials Q(x) and R(x). Notice the following algorithm,

- 1:  $A(x) \leftarrow a_n x^n + a_{n-1} x^{n-1} + \dots + a_0$
- 2:  $B(x) \leftarrow b_n x^n + b_{n-1} b^{n-1} + \dots + b_0$
- 3:  $R(x) \leftarrow A(x)$
- 4: while  $\deg R(x) \ge \deg B(x)$  do
- 5:  $a \leftarrow \text{leading coefficient of } R(x)$
- 6:  $b \leftarrow \text{leading coefficient of } B(x)$

- 7:  $d \leftarrow \deg R(x) \deg B(x)$
- 8:  $Q(x) = Q(x) + \left(\frac{a}{b}\right) x^d$
- 9:  $R(x) \leftarrow R(x) \left(\frac{a}{b}\right) x^d B(x)$ **output** Q(x) and R(x)

In each iteration of the while loop,  $\deg R(x)$  is decreasing (mono-variant) and the polynomial Q(x)B(x) + R(x) always stays equal to A(x) (invariant). At some point we will eventually get  $\deg R(x) \leq \deg B(x)$  which proves the existence of Q(x) and R(x).

**Remark.** Notice that, if  $A(x), B(x) \in \mathbb{R}[x]$  then  $Q(x), R(x) \in \mathbb{R}[x]$ . This implies that, if  $A(x), B(x) \in \mathbb{R}[x]$  and  $B(x) \mid A(x)$  then  $\frac{A(x)}{B(x)} \in \mathbb{R}[x]$ 

For example, if  $B(x) = x^2 - x + 1$  and  $A(x) = x^5 + x^3 + 2x$  then,

$$x^5 + x^3 + 2x = (x^3 + x^2 + x)(x^2 - x + 1) + x$$

In this example, the remainder R(x) = x and the quotient  $Q(x) = x^3 + x^2 + x$ .

#### **Theorem 1.1.2** (Remainder Theorem)

If P(x) is a polynomial and a is a constant then the remainder upon dividing P(x) by the linear polynomial x - a is equal to P(a).

**Proof:** From the Division Algorithm we know that there exists polynomials Q(x) and R(x) such that,

$$P(x) = Q(x)(x - a) + R(x)$$

Since  $\deg R(x) < \deg(x-a) = 1$ , R(x) must be a constant polynomial. Let us assume, R(x) = r. Now letting x = a we get,

$$P(a) = Q(a) \times (a - a) + r \implies P(a) = r$$

Therefore P(a) is the remainder upon dividing P(x) by x - a. QED

#### **Theorem 1.1.3** (Factor Theorem)

The number z will be a root of the polynomial P(x) if and only if P(x) is divisible by x-z.

**Proof:** We will first prove that,  $P(z) = 0 \implies (x - z) \mid P(x)$ . Let us assume that r is the remainder upon dividing P(x) by x - z. Now we know from the Remainder Theorem that, P(z) = r. But since z is a root of P(x), P(z) = r = 0. Therefore since r = 0, we must have  $(x - z) \mid P(x)$ . Using similar arguments one can also prove the converse.

#### Corollary 1.1.3.1

The number  $-\frac{b}{a}$  where  $a, b \in \mathbb{R}$  will be a root of the polynomial P(x) if and only if the polynomial P(x) is divisible by ax + b.

If P(x) has the root z then the Factor Theorem guarantees that there exists a polynomial Q(x) such that,

$$P(x) = (x - z) Q(x)$$

Now if,

$$P(x) = (x - z)^m Q'(x), \quad Q'(z) \neq 0$$

then we say that z is root of P(x) of multiplicity m.

For example, in the polynomial  $P(x) = (x-2)^2 (x-3)$  the root 2 has multiplicity 2 and the root 3 has multiplicity 1.

### §1.2 The Fundamental Theorem of Algebra

#### **Theorem 1.2.1** (The Fundamental Theorem of Algebra)

The Fundamental Theorem of Algebra states that, every polynomial P(x) in  $\mathbb{C}[x]$  has at least one root in  $\mathbb{C}$ 

#### Corollary 1.2.1.1

If  $P(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$  is a polynomial of degree n then,

$$P(x) = k(x - z_1)(x - z_2) \cdots (x - z_n)$$

where,  $k = a_n$  and  $z_i \in \mathbb{C}$ . The numbers  $z_1, z_2 \cdots z_n$  are not necessarily distinct.

**Proof:** This is an immediate consequence of The Fundamental Theorem of Algebra and Factor Theorem.

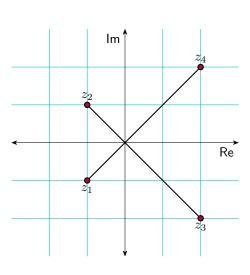


Figure 1.1: The 4 complex roots of the polynomial  $x^4 - 2x^3 + 2x^2 + 8x + 16$ 

#### **Theorem 1.2.2** (Complex Conjugate Root Theorem)

If  $P(x) \in \mathbb{R}[x]$  and z = a + bi where  $a, b \in \mathbb{R}$  is a complex root of the polynomial P(x) then  $\overline{z} = a - bi$  is also a root of the polynomial P(x).

**Proof 1:** We have to show that,  $P(z) = 0 \implies P(\overline{z}) = 0$ . Let  $\mathbb{C}' = \{ki : k \in \mathbb{R}\}$  and let  $\mathbb{R}$  be the set of real numbers. Now,

$$z^{k} + \overline{z}^{k} = (a+bi)^{k} + (a-bi)^{k}$$

$$= \sum_{j=0}^{k} {k \choose j} b^{j} a^{k-j} i^{j} + \sum_{j=0}^{k} {k \choose j} b^{j} a^{k-j} (-i)^{j}$$

$$= \sum_{j=0}^{k} {k \choose j} b^{j} a^{k-j} \left\{ i^{j} + (-i)^{j} \right\}$$

Notice,  $i^j + (-i)^j$  will be zero if j is odd. If j is even then  $i^j + (-i)^j = 2i^j = 2(-1)^{\frac{j}{2}}$ . Therefore,

$$z^{k} + \overline{z}^{k} = \sum_{j=0}^{k} {k \choose j} b^{j} a^{k-j} \left\{ i^{j} + (-i)^{j} \right\}$$
$$= \sum_{l=0}^{\left\lfloor \frac{k}{2} \right\rfloor} {k \choose 2l} b^{2l} a^{k-2l} \left\{ 2(-1)^{l} \right\}$$
$$= \sum_{l=0}^{\left\lfloor \frac{k}{2} \right\rfloor} {k \choose 2l} 2(-1)^{l} b^{2l} a^{k-2l}$$

**Remark.** The set,  $\{2l: 0 \le l \le \lfloor \frac{k}{2} \rfloor\}$ , contains all even integers (including zero) less than or equal to k.

Therefore,  $z^k + \overline{z}^k \in \mathbb{R}$  for all  $0 \le k \in \mathbb{Z}$ . This implies that,

$$P(z) + P(\overline{z}) = \sum_{i=0}^{n} a_i \left( z^i + \overline{z}^i \right) \in \mathbb{R}$$

But since P(z) = 0,  $P(z) + P(\overline{z}) \in \mathbb{R}$  implies  $P(\overline{z}) \in \mathbb{R}$ . Now,

$$\begin{split} z^k - \overline{z}^k &= (a+bi)^k - (a-bi)^k \\ &= \sum_{j=0}^k \binom{k}{j} b^j a^{k-j} i^j - \sum_{j=0}^k \binom{k}{j} b^j a^{k-j} (-i)^j \\ &= \sum_{j=0}^k \binom{k}{j} b^j a^{k-j} \left\{ i^j - (-i)^j \right\} \end{split}$$

If j is even then  $i^j - (-i)^j$  will be equal to zero. If j is odd that is j = 2l - 1 for some  $l \in \mathbb{N}$  then  $i^j - (-i)^j = i^{2l-1} \left(1 - (-1)^{2l-1}\right) = 2i^{2l-1} = 2(-1)^l i^{-1} = 2(-1)^l i^3 = 2(-1)^{l+1} i$ . Therefore,

$$z^{k} - \overline{z}^{k} = \sum_{j=0}^{k} {k \choose j} b^{j} a^{k-j} \left\{ i^{j} - (-i)^{j} \right\}$$

$$= \sum_{l=1}^{\left \lfloor \frac{k+1}{2} \right \rfloor} {k \choose 2l-1} b^{2l-1} a^{k-2l+1} 2(-1)^{l+1} i$$

$$= \left( \sum_{l=1}^{\left \lfloor \frac{k+1}{2} \right \rfloor} {k \choose 2l-1} b^{2l-1} a^{k-2l+1} 2(-1)^{l+1} \right) i$$

**Remark.** The set  $\left\{2l-1:1\leq l\leq \left\lfloor\frac{k+1}{2}\right\rfloor\right\}$  contains all odd positive integers less than or equal to k.

Thus,  $z^k - \overline{z}^k \in \mathbb{C}'$  for all  $k \in \mathbb{N}$ . Now,

$$P(z) - P(\overline{z}) = \sum_{i=0}^{n} a_i \left( z^i - \overline{z}^i \right)$$

$$\implies P(z) - P(\overline{z}) = \sum_{i=1}^{n} a_i \left( z^i - \overline{z}^i \right) + a_0 \left( z^0 - \overline{z}^0 \right)$$

$$\implies P(z) - P(\overline{z}) = \sum_{i=1}^{n} a_i \left( z^i - \overline{z}^i \right) \in \mathbb{C}'$$

But since P(z) = 0,  $P(z) - P(\overline{z}) \in \mathbb{C}'$  implies  $P(\overline{z}) \in \mathbb{C}'$ . And so,  $P(\overline{z}) \in \mathbb{R} \cup \mathbb{C}' \implies P(\overline{z}) \in \{0\} \implies P(\overline{z}) = 0$ . QED

**Proof 2(wiki):** Since P(z) = 0,

$$P(z) = \sum_{k=0}^{n} a_k z^k = 0$$

Now using the properties of complex conjugates,

$$P(\overline{z}) = \sum_{k=0}^{n} a_k \overline{z}^k = \sum_{k=0}^{n} a_k \overline{z^k} = \sum_{k=0}^{n} \overline{a_k z^k} = \overline{\sum_{k=0}^{n} a_k z^k} = \overline{P(z)} = \overline{0} = 0$$

Therefore,  $P(\overline{z}) = 0$ .

#### Corollary 1.2.2.1

If z is a complex root of the polynomial P(x) of multiplicity m then  $\bar{z}$  is also a complex root of the polynomial P(x) of multiplicity m. That is, complex conjugate roots have the same multiplicity.

**Proof:** If  $z \in \mathbb{R}$  then obviously z and  $\bar{z}$  will have the same multiplicity as  $z = \bar{z}$ . Let us assume  $z \notin \mathbb{R}$  and let m and n be the multiplicity of z and  $\bar{z}$  respectively. Without loss of generality, we can assume n < m. Now, let

$$P(x) = (x-z)^m (x-\bar{z})^n Q(x)$$

Now,

$$P(x) = (x-z)^n (x-\bar{z})^n (x-z)^{m-n} Q(x)$$

$$\Longrightarrow \frac{P(x)}{(x-z)^n (x-\bar{z})^n} = (x-z)^{m-n} Q(x)$$

Let,  $R(x) = \frac{P(x)}{(x-z)^n(x-\bar{z})^n}$ . Since  $P(x) \in \mathbb{R}[x]$  and  $(x-z)^n(x-\bar{z})^n \in \mathbb{R}[x]$ ,  $R(x) \in \mathbb{R}[x]$ . Therefore,  $R(x) = (x-z)^{m-n}Q(x) \in \mathbb{R}[x]$ . As z is a root of R(x) and  $R(x) \in \mathbb{R}[x]$ ,  $\bar{z}$  must also be a root of R(x) which implies the multiplicity of  $\bar{z} > n$ . But that contradicts our assumption that  $\bar{z}$  has multiplicity n. Therefore, m and n must be equal.

#### Corollary 1.2.2.2

Every polynomial P(x) in  $\mathbb{R}[x]$  can be expressed in the form,

$$P(x) = f_1^{e_1}(x) f_2^{e_2}(x) \cdots f_n^{e_n}(x)$$

where the polynomials  $f_i(x)$  are either linear or quadratic polynomials in  $\mathbb{R}[x]$  and  $e_i \in \mathbb{N}$ 

#### Corollary 1.2.2.3

If  $P(x) \in \mathbb{R}[x]$  and deg P(x) is odd then P(x) has at least on real root.

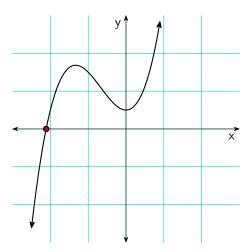


Figure 1.2: The real root of the cubic polynomial  $f(x) = x^3 + 2x^2 + 0.5$ 

# §1.3 Quadratic Polynomials

**Definition 1.3.1.** A quadratic polynomial is a polynomial of the form,

$$P(x) = ax^2 + bx + c$$

where a, b, c are constants and  $a \neq 0$ .

One can find the roots of a quadratic polynomial using the well known quadratic formula,

$$x_{1,2} = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

The value  $\Delta = b^2 - 4ac$  is called the *discriminant* of the quadratic polynomial. The discriminant gives us the following informations about the roots of the quadratic polynomial,

- $\Delta > 0 \iff x_1, x_2 \in \mathbb{R} \text{ and } x_1 \neq x_2$
- $\Delta = 0 \iff x_1, x_2 \in \mathbb{R} \text{ and } x_1 = x_2$
- $\Delta < 0 \iff x_1, x_2 \in \mathbb{C} \text{ and } x_1 \neq x_2$

#### Theorem 1.3.1

The value  $P\left(-\frac{b}{2a}\right)$  is either the maximum (if a>0) or the minimum value (if a<0) of the quadratic polynomial,  $P(x)=ax^2+bx+c$ 

**Proof:** 

$$P(x) = ax^{2} + bx + c$$

$$\implies P(x) = a\left(x^{2} + 2\frac{b}{2a}x + \frac{b^{2}}{4a^{2}}\right) + c - \frac{b^{2}}{4a}$$

$$\implies P(x) = a\left(x + \frac{b}{2a}\right)^{2} + \left(c - \frac{b^{2}}{4a}\right)$$

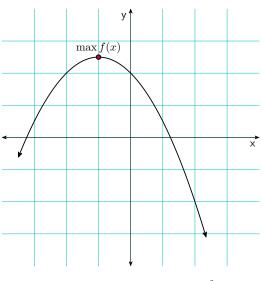
If a < 0 then,

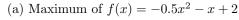
$$P(x) = \left(c - \frac{b^2}{4a}\right) - |a| \left(x + \frac{b}{2a}\right)^2$$

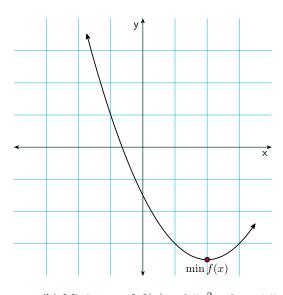
P(x) will reach its maximum when  $x + \frac{b}{2a} = 0 \implies x = -\frac{b}{2a}$ . If a > 0 then,

$$P(x) = \left(c - \frac{b^2}{4a}\right) + a\left(x + \frac{b}{2a}\right)^2$$

P(x) will reach its minimum when  $x + \frac{b}{2a} = 0 \implies x = -\frac{b}{2a}$ .







(b) Minimum of  $f(x) = 0.5x^2 - 2x - 1.5$ 

# §1.4 Roots of Cubic Polynomials

Finding the roots of a cubic polynomial is quite hard. So, we are going to first try to solve the cubic polynomial,

$$f(x) = x^3 + px + q$$

where  $p, q \in \mathbb{R}$ . Setting p = -3ab and  $q = a^3 + b^3$  we get,

$$f(x) = x^3 + a^3 + b^3 - 3abx$$

Now using the formula,

$$a^{3} + b^{3} + c^{3} - 3abc = (a+b+c)(a^{2} + b^{2} + c^{2} - ab - bc - ca)$$

we get,

$$f(x) = (x + a + b)(x^{2} - (a + b)x + a^{2} + b^{2} - ab)$$

Therefore the 3 roots of f are,

$$x_{1} = -a - b$$

$$x_{2}, x_{3} = \frac{a + b \pm \sqrt{a^{2} + b^{2} + 2ab - 4a^{2} - 4b^{2} + 4ab}}{2}$$

$$= \frac{a + b \pm \sqrt{-3a^{2} - 3b^{2} + 6ab}}{2}$$

$$= \frac{a + b \pm \sqrt{-3(a - b)^{2}}}{2}$$

$$= \frac{a + b \pm \sqrt{3}i(a - b)}{2}$$

$$= \frac{(1 + \sqrt{3}i)a \pm (1 - \sqrt{3}i)b}{2}$$

$$= \frac{1 + \sqrt{3}i}{2}a \pm \frac{1 - \sqrt{3}i}{2}b$$

Now we have express the root in terms of p, q that is, we have to express a, b in terms of p, q. Now,

$$q = a^{3} + b^{3}$$
$$p = -3ab$$
$$\implies a^{3}b^{3} = -\frac{p^{3}}{27}$$

Let  $u = a^3$  and  $v = b^3$ . Notice that u and v are the roots of the quadratic polynomial,

$$P(x) = x^{2} - (u+v)x + uv = x^{2} - q^{3}x - \frac{p^{3}}{27}$$

Using the quadratic equation we get,

$$u, v = \frac{q^3 \pm \sqrt{q^6 + \frac{4}{27}p^3}}{2}$$
$$a, b = \sqrt[3]{\frac{q^3}{2} \pm \frac{\sqrt{q^6 + \frac{4}{27}p^3}}{2}}$$
$$= \sqrt[3]{\frac{q^3}{2} \pm \sqrt{\frac{q^6}{4} + \frac{p^3}{27}}}$$

So let's now try to solve the cubic equation using the results we've got so far,

$$f(x) = x^3 - x^2 - 2x + 1$$

First we have to use substitution to transform the polynomial into another polynomial of the form,

$$x^3 + px + q$$

Since every polynomial can be uniquely defined by its coefficients, we can associate or express a polynomial of degree n by a unique point in n+1 dimensional space. For example we can express the polynomial,

$$(1) \cdot x^3 + (-1) \cdot x^2 + (-2) \cdot x + (1)$$

as,

$$(1) \cdot x^3 + (-1) \cdot x^2 + (-2) \cdot x + (1) \to (1, -1, -2, 1)$$

Likewise, the point (5, -1, 0, 1) can be used to represent the polynomial,

$$(5,-1,0,1) \rightarrow 5x^3 - x^2 + 1$$

Notice that,

$$(x_3, x_2, \dots, x_0) + (y_3, y_2, \dots, y_0) = (x_3 + y_3, x_2 + y_2, \dots, x_0 + y_0)$$

Say that there exists a polynomial u(x) = (x+n) where  $n \in \mathbb{C}$  such that,

$$x^{3} - x^{2} - 2x + 1 = u(x)^{3} + pu(x) + q$$

This equation can also be represented as,

$$(1,-1,-2,1) = (1,3n,3n^2,n^3) + (0,0,p,pn) + (0,0,0,q) \implies (1,-1,-2,1) = (1,3n,3n^2+p,n^3+pn+q)$$

This gives us the system of equation,

$$3n = -1$$
$$3n^2 + p = -2$$
$$n^3 + pn + q = 1$$

Therefore,

$$n = -\frac{1}{3}$$
$$p = -\frac{7}{3}$$
$$q = \frac{7}{27}$$

Thus,

$$x^{3} - x^{2} - 2x + 1 = \left(x - \frac{1}{3}\right)^{3} - \frac{7}{3}\left(x - \frac{1}{3}\right) + \frac{7}{27} \implies f(x) = g\left(x - \frac{1}{3}\right)$$

where

$$g(x) = x^3 - \frac{7}{3}x + \frac{7}{27}$$

Now we can use the results we've proved earlier to solve the polynomial g(x). After that we add  $\frac{1}{3}$  to the 3 roots of g(x). The 3 numbers we will get by adding  $\frac{1}{3}$  are the 3 roots of f(x).

# §1.5 Lagrange Interpolation

#### **Theorem 1.5.1** (Lagrange Interpolation)

Let  $\alpha_0, \alpha_1, \dots, \alpha_n$  be distinct real numbers and  $\beta_0, \beta_1, \dots, \beta_n$  be another set of n+1 real numbers. Then there exists a unique polynomial,

$$P(x) = \sum_{i=0}^{n} \left( \prod_{\substack{j=0\\j\neq i}}^{n} \frac{x - \alpha_j}{\alpha_i - \alpha_j} \right) \beta_i$$

with deg  $P(x) \le n$  such that  $P(\alpha_k) = \beta_k$  for all  $0 \le k \le n$ .

Proof: Let,

$$D_k(x) = \prod_{\substack{j=0\\j\neq k}}^n \frac{x - \alpha_j}{\alpha_k - \alpha_j} = \frac{(x - \alpha_0)(x - \alpha_1)\cdots(x - \alpha_{k-1})(x - \alpha_{k+1})\cdots(x - \alpha_n)}{(\alpha_k - \alpha_0)(\alpha_k - \alpha_1)\cdots(\alpha_k - \alpha_{k-1})(\alpha_k - \alpha_{k+1})\cdots(\alpha_k - \alpha_n)}$$

If  $x = \alpha_k$  then  $D_k(x) = 1$  else if  $x = \alpha_i$  where  $i \neq k$  then  $D_k(x) = 0$ . Thus the polynomial,

$$P(x) = \sum_{k=0}^{n} D_k(x)\beta_k$$

will be equal to  $\beta_k$  for all  $x = \alpha_k$ . It is also clear that the polynomial P(x) has degree at most n since deg  $D_k(x) = n$  for all  $0 \le k \le n$ .

Now suppose that there exists two polynomials  $P_1(x)$  and  $P_2(x)$ , with degree at most n, such that,

$$P_1(\alpha_k) = P_2(\alpha_k) = \beta_k, \ 0 \le k \le n$$

Therefore the polynomial  $Q(x) = P_1(x) - P_2(x)$  has n+1 distinct roots. But that is impossible since we know that  $\deg Q(x) \leq n$  and a polynomial of degree n has at most n distinct roots. This proves that the polynomial P(x) must be unique, that is, P(x) is the only polynomial, with degree at most n, such that,  $P(\alpha_k) = \beta_k$  for all  $0 \leq k \leq n$ 

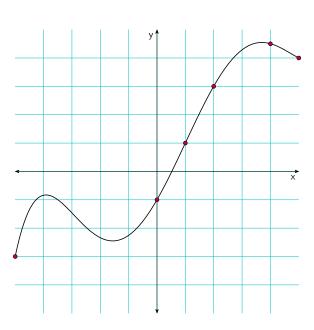


Figure 1.4: Plot of a Lagrange Polynomial

Figure 1.4 shows the Lagrange polynomial going through the points,

$$\{(1,1),(2,3),(0,-1),(5,4),(-5,-3),(4,4.5)\}$$

We can easily compute Lagrange polynomials in python using sympy.

#### Problem 1.5.1

Let P(x) be a polynomial of degree n such that,  $P(k) = 2^k$  for all  $0 \le k \le n$ . Find P(n+1).

**Solution:** From Theorem 1.5.1 we have,

$$P(x) = \sum_{k=0}^{n} 2^k D_k(x)$$

where,

$$D_k(x) = \frac{x(x-1)\cdots(x-k+1)(x-k-1)(x-k-2)\cdots(x-n+1)(x-n)}{(k)(k-1)\cdots(1)(-1)(-2)\cdots(k-n+1)(k-n)}$$
$$= (-1)^{n-k} \frac{x(x-1)\cdots(x-k+1)(x-k-1)(x-k-2)\cdots(x-n+1)(x-n)}{k!(n-k)!}$$

Therefore,

$$P(n+1) = \sum_{k=0}^{n} (-1)^{n-k} 2^k \frac{(n+1)n(n-1)\cdots(n-k+2)(n-k)(n-k-1)\cdots1}{k!(n-k)!}$$

$$= \sum_{k=0}^{n} (-1)^{n-k} 2^k \frac{(n+1)!}{k!(n-k)!(n-k+1)}$$

$$= \sum_{k=0}^{n} (-1)^{n-k} 2^k \binom{n+1}{k}$$

$$= (-1) \left(\sum_{k=0}^{n+1} \binom{n+1}{k} 2^k (-1)^{n-k+1}\right) + 2^{n+1}$$

$$= (-1) (2-1)^{n+1} + 2^{n+1}$$

$$= 2^{n+1} - 1$$